# Supervision work Supervision 11

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# 1 Ticks 1 and 1\*

#### 1.1 Tick 1

```
package uk.ac.cam.tma33.alg.tick1;
import uk.ac.cam.rkh23.Algorithms.Tick1.MaxCharHeapInterface;
import uk.ac.cam.rkh23.Algorithms.Tick1.EmptyHeapException|;
public class MaxCharHeap implements MaxCharHeapInterface{
char[] mHeap;
int mLength;
public MaxCharHeap(String s) {
        mLength = 0;
        mHeap = new char[1];
        char[] input = s.toCharArray();
        for (int i = 0; i < input.length; i++) {
                insert(input[i]);
        }
@Override
public char getMax() throws EmptyHeapException {
        if (mLength < 1) {
                // the heap is empty
                throw new EmptyHeapException();
        else {
                char result = mHeap[1];
```

```
swap (0, mLength);
                  mLength--;
                 heapDOWN(0);
                  // we return the head of the heap
                  return result;
         }
private void heapDOWN(int node){
         int sonLeft = node*2;
         int sonRight = node*2+1;
         int nextNode;
         if (sonLeft <= mLength) {</pre>
                  nextNode = sonLeft;
                  if (sonRight <= mLength && mHeap[sonRight]| >
                                            mHeap[sonLeft]) {
                          nextNode = sonRight;
                  if (mHeap[nextNode] > mHeap[node]){
                          swap(nextNode, node);
                          heapDOWN(nextNode);
                  }
         }
private void swap(int x, int y) {
         char aux = mHeap[x];
         \mathrm{mHeap}\,[\,\mathrm{x}\,] \ = \ \mathrm{mHeap}\,[\,\mathrm{y}\,]\,;
         mHeap[y] = aux;
private void resize() {
         int newSize = mHeap.length * 2;
         char[] newHeap = new char[newSize];
         for (int i = 1; i \le mLength; i++) {
                 newHeap[i] = mHeap[i];
         mHeap = newHeap;
private void heapUP(int node) {
         int father = node/2;
         if (father > 0 \&\& mHeap[node] > mHeap[father]) {
                  swap(node, father);
                 heapUP(father);
         }
```

```
@Override
public void insert(char i) {
    if (mLength + 1 >= mHeap.length) {
        // we double the size of the array
        resize();
    }
    char newChar = Character.toLowerCase(i);
    mLength++;
    mHeap[mLength] = newChar;
    heapUP(mLength);
}

@Override
public int getLength() {
    return mLength;
}
```

## 1.2 Tick 1\*

(a) Implementation of MaxHeap:

```
heapify();
}
private void heapify() {
        int N = mHeap. size();
        N /= 2;
        for (int i = N; i >= 0; i--) {
                 heapDOWN(i);
         }
}
@Override
public T getMax() throws EmptyHeapException {
         if (mHeap.size() = 0) {
                 // empty heap \Rightarrow we throw exception
                 throw new EmptyHeapException();
        }
        else {
                 T \text{ result} = mHeap.get(0);
                 swap(0, mHeap.size() - 1);
                 mHeap.remove(mHeap.size() - 1);
                 heapDOWN(0);
                 return result;
        }
}
private void heapDOWN(int node){
        int sonLeft = node*2+1;
        int sonRight = node*2+2;
        int nextNode;
         int length = mHeap.size();
         if (sonLeft < length) {</pre>
                 nextNode = sonLeft;
         if (sonRight < length && mHeap.get(sonLeft).
          compareTo(mHeap.get(sonRight)) < 0) {
                 nextNode = sonRight;
         if (mHeap.get(nextNode).compareTo(mHeap.get(node)) | > 0){
                 swap(nextNode, node);
                 heapDOWN(nextNode);
         }
}
}
```

```
protected void swap(int x, int y) {
        T \text{ aux} = m \text{Heap.get}(x);
        mHeap.set(x, mHeap.get(y));
        mHeap.set(y, aux);
private void heapUP(int node) {
        int father = (\text{node}-1)/2;
         if (father >= 0 \&\& mHeap.get(node).
                 compareTo(mHeap.get(father)) > 0) {
                 swap(node, father);
                 heapUP(father);
         }
}
@Override
public void insert(T i) {
        mHeap.add(i);
         // getting the number we just added to the right position
        heap UP(mHeap.size() - 1);
}
@Override
public int getLength() {
        return mHeap. size();
```

#### (b) Implementation of BottomUpMaxHeap:

```
}
public BottomUpMaxHeap(List<T> input){
        super(input);
}
@Override
public T getMax() throws EmptyHeapException {
        if (mHeap.size() == 0)  {
                 // empty heap \Rightarrow we throw exception
                 throw new EmptyHeapException();
        else {
                 mPath = new ArrayList<Integer >();
                 T result = mHeap.get(0);
                 swap(0, mHeap. size() - 1);
                 mHeap.remove(mHeap.size() - 1);
                 goDown(0);
                 T \text{ root} = mHeap.get(0);
                 int N = mPath.size() - 1;
                 for (int i = N; i > 0; i--) {
                         int node = mPath.get(i);
                          if (root.compareTo(mHeap.get(node))
                                           < 0) {}
                                  shiftUP(i);
                                  mHeap.set (node, root);
                                  break;
                          }
                 return result;
        }
}
private void shiftUP(int N) {
        for (int i = 0; i < N; i++) {
                 mHeap. set (mPath. get (i), mHeap. get (mPath. get (i+1)));
        }
private void goDown(int node) {
        int sonLeft = node*2+1;
        int sonRight = node*2+2;
        int nextNode;
        mPath.add(node);
```

\*note\* : i submitted the BottomUpMaxHeap, but it says that it throws an unexpected exception... I don't understand why. it works on my tests without any problem

# 2 Algorithms sheet

#### 2.1 Exercise 2

- (a) **Brute-force :** A brute-force strategy consists in just implementing the most inefficient algorithm for a specific problem. The solution has to be correct, though.
- (b) **Divide-and-conquer :** This strategy is based on splitting the problem into more sub-problems, that are easily to solve individually. The divide-and-conquer strategy usually is associated with a logarithmic execution time.
- (c) **Dynamic programming** is also based on splitting the problem in multiple sub-problems. Unlike the divide-and-conquer, this strategy aims to find a general formula for the answer, by applying it to small examples.
- (c) **Backtracking** is a programming strategy that tries multiple solutions for a problem, until it reaches a valid one. As the name indicates, this strategy involves "walking back" to previous solutions, to make changes onto them.

### 2.2 Exercise 3

The sudoku implementation:

```
package uk.ac.cam.tma33.s11;

public class Sudoku {

public static boolean ok(long a, long b, long x) {
    if ((a&x) != 0) return false;
    if ((b&x) != 0) return false;
    return true;
```

```
}
private static boolean full(long[] a, long[] b) {
        \mathbf{long} \ \mathbf{MAX} = 0;
        for (int i = 1; i \le 9; i++) {
                MAX += (1 << i) *1L;
                // generating the value MAX
        for (int i = 0; i < 8; i++) {
                 if (a[i] != MAX || b[i]!=MAX) {
                         // then the sudoku is not solved
                         return false;
                 }
        return true;
public static boolean Solve(long[][] A, long[] cols,
                 long[] rows, int x, int y) {
        if (x > 8) {
                 // We are at the end of the board
                 return full(cols, rows))
        for (long number = 1; number <= 9; number++) {
                 long noCode = 1 << number;
                 if (ok(cols[y], rows[x], noCode))  {
                         // number is not on the current row/
                         //column so we can proceed
                A[x][y] = number;
                int newX;
                 int newY;
                 if (y >= 8) {
                         // we are at the end of the row
                         newY = 0;
                         newX = x + 1; // get to the next row
                 else {
                         newY = y + 1; // next column, same | row
                         newX = x;
                 cols[y] += noCode; // we add the code to the row/
                 rows[x] += noCode; // column, to know
                                  //that we used this no
                 // we try to find a number to fit the next cell
                 boolean solved = Solve(A, cols, rows, newX, newY);
```

```
if (solved) return true;
                cols[y] -= noCode; //we substract the code this time,
                rows[x] -= noCode; // because we want
                                 //to try a different solution
                }
return false;
public static void main(String[] args) {
        long[][] gameBoard = new long[9][9]; // the gameboard
        long[] cols = new long[9];
        // cols[i] encodes the numbers on collumn i
        long[] rows = new long[9];
        // rows[i] encodes the numbers on row i
        for (int i = 0; i \le 8; i++) {
                cols[i] = 0; // we initialise the cols
                rows[i] = 0; // and rows arrays with 0
        boolean solved = Solve (gameBoard, cols, rows, 0, 0);
        // printing the solution
        for (int i = 0; i \le 8; i++) {
                for (int j = 0; j \le 8; j++){
                        System.out.print(Long.toString(
                                 gameBoard[i][j]) + "¬");
        System.out.println();
        }
}
```

#### 2.3 Exercise 5

(i) In the Foundations of Computer Science course, the strategy used to make change was based on a rather optimised backtracking algorithm, looking for any number that could be put in the result and trying it until the amount we had reached 0.

#### (ii) Backtracking approach:

```
package uk.ac.cam.tma33.s11;
public class MakeChange {
public int[] makeChange(int amt, int[] coins, int[] Sol) {
        int N = coins.length;
        if (amt == 0) {
                // job done, we finished
                return Sol;
        else if (amt < 0) {
                // WHOOPS, too far ahead
                return null;
        }
        else {
                // we try to add one more coin to the solu|tion
                int lng = Sol.length;
                // resizing the Sol array
                int[] aux = new int[lng + 1];
                for (int i = 0; i < lng; i++) {
                        aux[i] = Sol[i];
                Sol = aux;
                // trying different coins to see if they work
                for (int i = 0; i < N; i++) {
                         if (amt - coins[i] >= 0)
                                 // we find a coin that
                                 //could go in
                                 Sol[lng] = coins[i];
                                 int [] result = makeChange
                                         amt-coins [i], coins, Sol);
                                 if (result != null) return result;
                         }
                }
return null;
}
```

This implementation has a complexity of  $O(2^n)$ , where n is the number of coins available. We have this complexity due to the fact that, for every coin we have to options: to use it or not to use it.

#### Dynamic-programming approach:

```
package uk.ac.cam.tma33.s11;
public class DynamicMakingChange {
public int makeChange(int amt, int[] coins) {
        int N = coins.length;
         int[][] Sol = new int[amt+2][N + 2];
         for (int i = 0; i \le N; i++) {
                 Sol[i][0] = 1;
         for (int i = 1; i \le amt; i++) {
                  Sol[0][i] = 0;
         for (int i = 1; i \le N; i++) {
                  for (int j = 1; j \le amt; j++) {
                           if (j - coins[i-1] >= 0) {
                                   // we can make change with // this coin
                                    Sol[i][j] = Sol[i-1][j] + Sol[i][j] - coins[i-1]];
                           }
else {
                                    // we can't use coin \#(i-1)
                                    Sol[i][j] = Sol[i-1][j];
                  }
return Sol[amt][N];
```

This approach has a complexity of O(n\*amt), as it has to go through every element in the matrix exactly one time.

#### 2.4 Exercise 6

```
package uk.ac.cam.tma33.s11;
public class MinMax {
private int getSmaller(int x, int y) {
        if (x < y) return x;
        else return y;
public int getMin(int[] a, int left, int right) {
        int mid = (left + right)/2;
        if (right < left) return Integer.MAX_VALUE;</pre>
        if (right - left < 1) 
                 // return the minimum of the 2
                return getSmaller(a[right],a[left]);
        }
        else {
                 // we continue to divide the interval
                 int x = getMin(a, left, mid);
                 int y = getMin(a, mid+1, right);
                 return getSmaller(a[mid],getSmaller(x,y));
        }
private int getLarger(int x, int y) {
        if (x > y) return x;
        else return y;
public int getMax(int[] a, int left, int right) {
        int mid = (left + right)/2;
        if (right < left) return Integer.MIN_VALUE;</pre>
        if (right - left < 1) {
                // return the maximum of the 2
                return getLarger(a[right],a[left]);
        }
        else {
                 // we continue to divide the interval
                int x = getMin(a, left, mid);
                 int y = getMin(a, mid+1, right);
                 return getLarger(a[mid],getLarger(x,y));
        }
}
```

}

The number of operations that each of the 2 methods (getMin and getMax) executes is defined by the following recursive rule : T(n) = T(n/2) + T(n/2) + 2,  $\forall n \in \mathbb{N}, n \geq 2$ .

So, 
$$T(n) = 2^{k-1} + 2^k - 2$$
, where  $k = log(n)$ .

$$T(n) = \frac{n}{2} + n - 2 = \frac{3n}{2} - 2 = O(n).$$