

Max-flow and Circulation Comparison

Max-flow

- $G = (V, E)$ = directed graph,
- Two distinguished nodes:
 - s = source, t = sink.
- $c(e)$ = capacity of edge e .
- $0 \leq f(e) \leq c(e)$
- $\sum_{e \text{ in to } v} f(e) = \sum_{e \text{ out of } v} f(e)$
- max flow = min cut
- Algorithms:
- Generic augmenting path:
 - $O(m \cdot \text{val}(f^*))$.
- Capacity scaling:
 - $O(m^2 \log C)$
- *Shortest augmenting path:
 - $O(m^2 n)$.
- * Preflow-Push:
 - $O(m n^2)$ or $O(n^3)$.

Circulation with demands

- Node supply and demands $d(v)$, $v \in V$.

- G
- demand if $d(v) > 0$;
 - supply if $d(v) < 0$;
 - transshipment if $d(v) = 0$

- Conservation

$$\sum_{e \text{ in to } v} f(e) - \sum_{e \text{ out of } v} f(e) = d(v)$$

- Necessary condition to have a circulation

$$\sum_{v: d(v) > 0} d(v) = \sum_{v: d(v) < 0} -d(v) =: D$$

- Convert to network flow:

- G'
- Add new source s and sink t .
 - For each v with $d(v) < 0$, add edge (s, v) with capacity $-d(v)$.
 - For each v with $d(v) > 0$, add edge (v, t) with capacity $d(v)$.
 - Claim: G has circulation iff G' has max flow of value D (saturates all edges leaving s and entering t)

- with Demands and Lower Bound:

$$l(e) \leq f(e) \leq c(e)$$

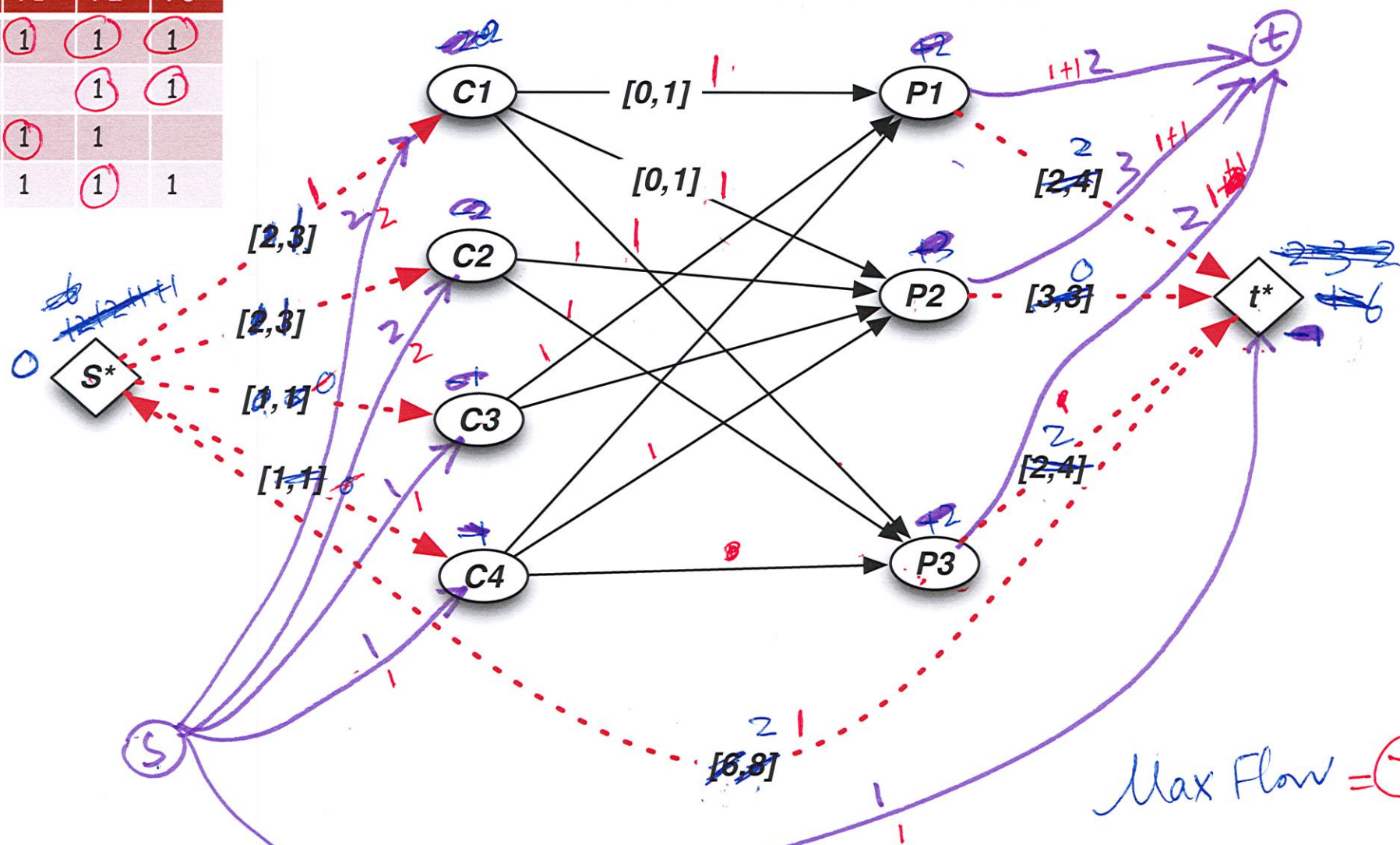
- Transfer each edge $e: (v, w)$:

- $d(v) = d(v) + l(e)$; $d(w) = d(w) - l(e)$; $c(e) = c(e) - l(e)$



Convert to Network Flow Problem

	3	3	2
	P1	P2	P3
C1	1	1	1
C2		1	1
C3	1	1	
C4	1	1	1



$$D = \sum_{d \in \mathcal{D}} z + 3 + 2 = 7 = \sum_{s \in \mathcal{S}} (-2, -2, -1, -1, -1) = 7$$

demand supply.

Survey Design Example

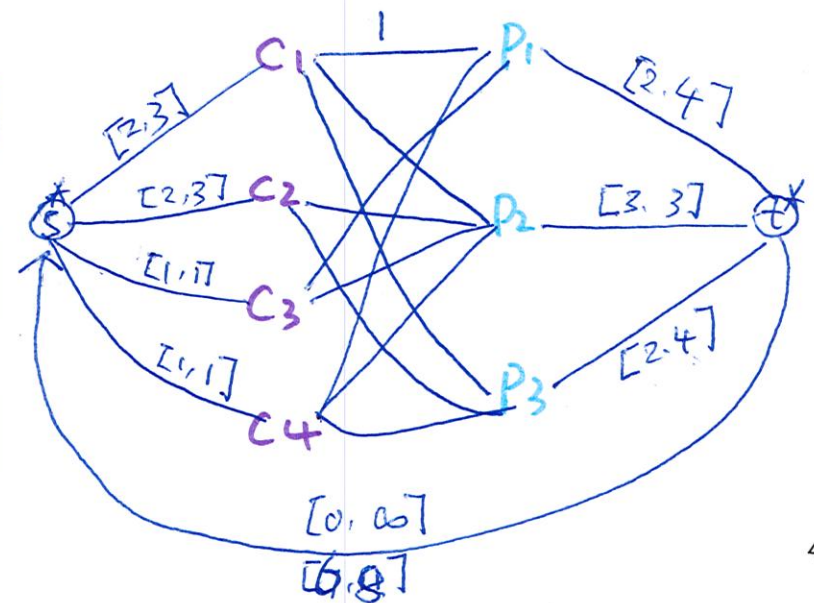
Survey design.

- Design survey asking 4 consumers about 3 products.
- Can only survey consumer i about product j if they own it, see table below.
- ✓ Ask consumer 1, consumer 2 each between 2 and 3 questions.
- ✓ Ask consumer 3, consumer 4 each 1 question only.
- ✓ Ask between 2 and 4 consumers about product 1.
- ✓ Ask 3 consumers about product 2.
- ✓ Ask between 2 and 4 consumers about product 3.

link: possible asking
flow: actual question

Goal. Design a survey that meets these specs, if possible.

	P1	P2	P3
C1	1	1	1
C2		1	1
C3	1	1	
C4	1	1	1



Project Selection

Projects with prerequisites.

- Set P of possible projects. Project v has associated revenue p_v .
 - some projects **generate** money: create interactive e-commerce interface, redesign web page
 - others **cost** money: upgrade computers, get site license



- Set of **prerequisites** E . If $(v, w) \in E$, can't do project v and unless also do project w .
- A subset of projects $A \subseteq P$ is **feasible** if the prerequisite of every project in A also belongs to A .

X A. BW X
 -10 B. UC ✓
 90 C. UC. BW ✓
 0 D. ~~UC~~

Project selection. Choose a **feasible** subset of projects to maximize **revenue**.

Project	Prerequisites	Profit
P1		-10 ✓
P2	P1	15 -
P3	P1, P2	-5
P4	P2	10 -
P5	P3	20

⊖C

$P_v > 0$ 45
 $P_v < 0$ -15