

# CS 2112 | Assignment 4 Written Problem

## Parsing and Fault Injection

Charles Qian {cq38}

Kelly Yu {kly24}

### Proving Asymptotic Complexity

For each of the following functions, show by giving a witness that it is  $O(n^2)$ , or else show that it isn't  $O(n^2)$  by arguing that no such witness can exist.

1.  $3n^2 + 10n + 1$

We choose a witness  $k = 5$  and  $n_0 = 9$ .  $g(n) = n^2$ .

For all  $n \geq n_0$ ,  $f(n) \leq k * g(n)$ :

$f(n) = 3n^2 + 10n + 1 \leq 5 * n^2$  for all  $n \geq 9$ .

2.  $2^n$

$2^n$  is NOT  $O(n)$  because the function  $2^n$  is raised to an exponent  $n$  that increases linearly, whereas function  $n^2$  is raised to a constant exponent 2. No matter what  $k$  and  $n_0$  we choose,  $2^n$  will always approach  $\infty$  faster than  $n^2$ .

By contradiction: Say that there is a  $k$  and  $n_0$  such that  $n \geq n_0$ ,  $f(n) \leq k * g(n)$ .

Then:

$$2^n \leq k * n^2$$

$$n \leq \log_2(k * n^2)$$

$k$  varies according to a non-constant, so  $k$  can never be constant.

3.  $n \lg n$

We choose a witness  $k = 1$  and  $n_0 = 0$ .  $g(n) = n^2$ .

For all  $n \geq n_0$ ,  $f(n) \leq k * g(n)$ :

$f(n) = n \lg n \leq 1 * n^2$  for all  $n \geq 0$ .

4.  $n^3/\lg n$

$n^3/\lg n$  is NOT  $O(n)$  because  $\lg n$  will always be less than  $n$ , so  $n^3/\lg n$  will always be greater than  $n^3/n$ , which is equal to  $n^2$ .

By contradiction: Say that there is a  $k$  and  $n_0$  such that  $n \geq n_0$ ,  $f(n) \leq k * g(n)$ .

Then:

$$n^3/\lg n \leq k * n^2$$

$$n/\lg n \leq k$$

The left hand side of the equation will ultimately approach  $\infty$  as  $n$  increases, so  $k$  is forced to vary according to this non-constant function, and  $k$  can therefore not be a constant.

5.  $f(n) + h(n)$ , where each of  $f(n)$  and  $h(n)$  are  $O(n^2)$ .

By definition of  $O(n^2)$ , we know that there is a  $k_1$  such that:

$$f(n) \leq k_1 * g(n)$$

and a  $k_2$  such that:

$$h(n) \leq k_2 * g(n)$$

Adding these two equations will produce:

$$f(n) + h(n) \leq k_1 * g(n) + k_2 * g(n)$$

$$f(n) + h(n) \leq (k_1 + k_2)g(n)$$

By the definition of  $O(n^2)$ , this means that  $f(n) + h(n)$  must also be  $O(n^2)$ .