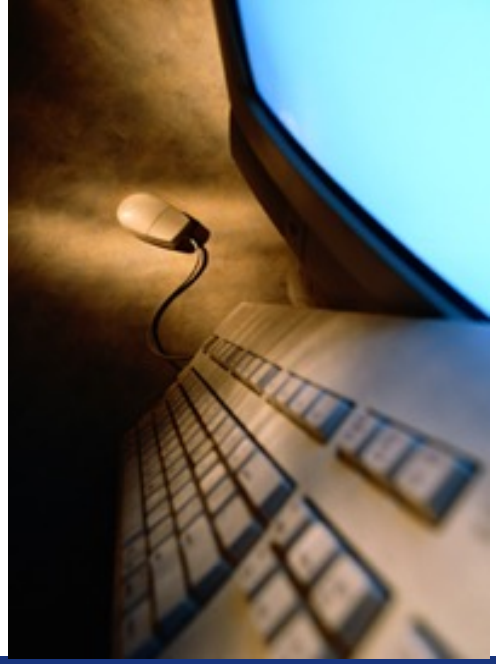
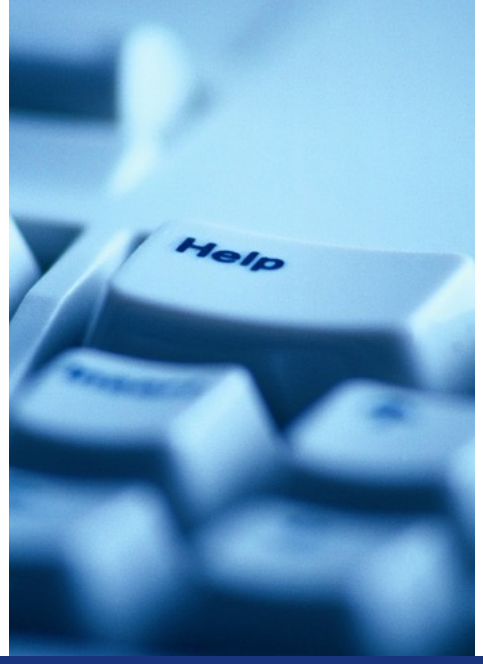
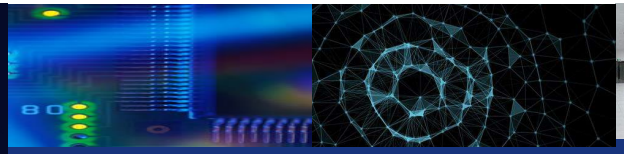



并行与分布式计算

Parallel & Distributed Computing

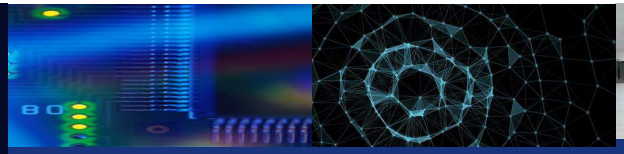
陈鹏飞
数据科学与计算机学院



Lecture 5 — OpenMP Programming

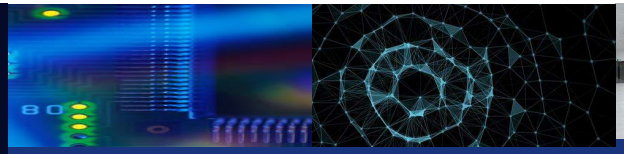
Pengfei Chen

School of Data and Computer Science



Outline:

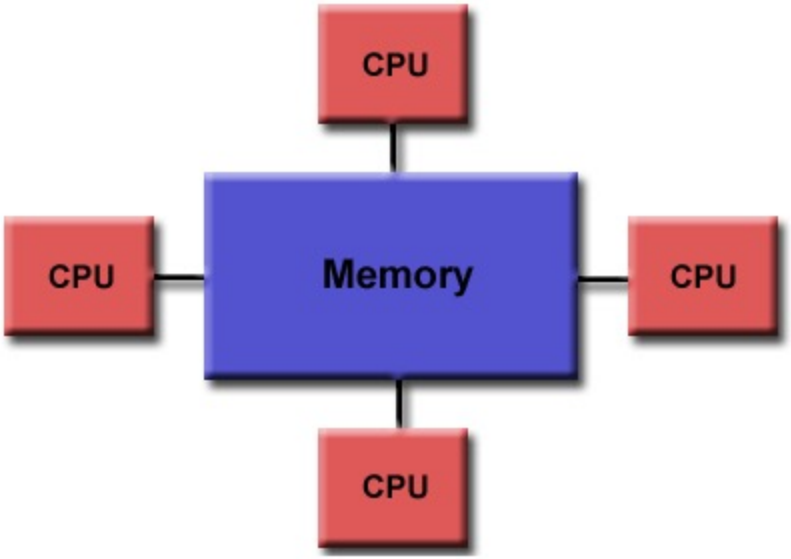
1. OpenMP Overview
2. Correctness
3. Task
4. Performance



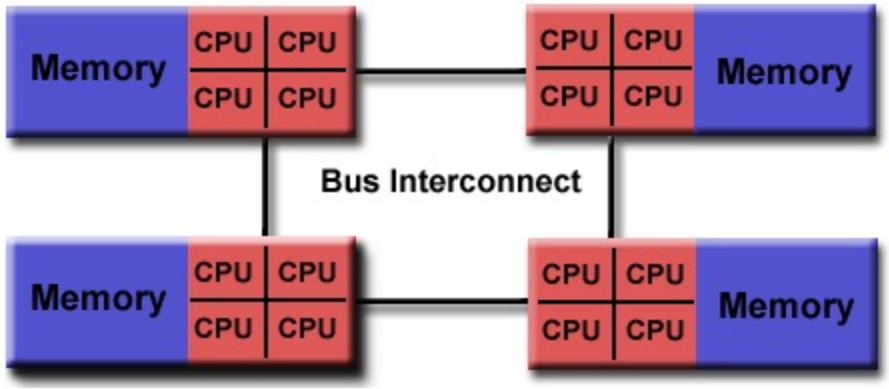
OpenMP programming

Overview

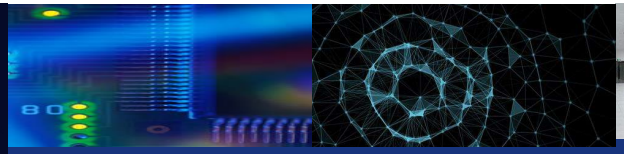
Architecture for Shared Memory Model



Uniform Memory Access



Non-uniform Memory Access



Thread Base Parallelism

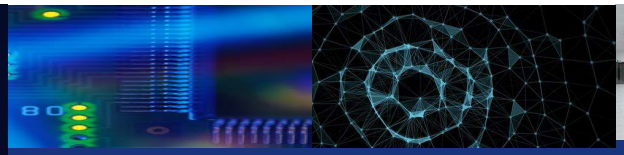
- OpenMP programs accomplish parallelism exclusively（仅仅）through the use of threads
- A thread of execution is the smallest unit of processing that can be scheduled by an operating system
  - The idea of a subroutine that can be scheduled to run autonomously might help explain what a thread is
- Threads exist within the resources of a single process
  - Without the process, they cease（停止）to exist
- Typically, the number of threads match the number of machine processors/cores
  - However, the actual use of threads is up to the application

Explicit Parallelism

- OpenMP is an explicit (not automatic) programming model, offering the programmer full control over parallelization
- Parallelization can be as simple as taking a serial program and inserting compiler directives....
- Or as complex as inserting subroutines to set multiple levels of parallelism, locks and even nested locks

What is OpenMP?

- An abbreviation for
  - Short version
    - Open Multi-Processing （开放多处理过程）
  - Long version
    - Open specifications for Multi-Processing via collaborative work between interested parties from the hardware and software industry, government and academia



OpenMP Overview

C$OMP FLUSH

#pragma omp critical

C$OMP THREADPRIVATE(/ABC/)

CALL OMP_SET_NUM_THREADS(10)

C$OM

C$OM

C$

C

#p

ED

OpenMP: An API for Writing Multithreaded Applications

- A set of compiler directives and library routines for parallel application programmers
- Greatly simplifies writing multi-threaded (MT) programs in Fortran, C and C++
- Standardizes last 20 years of SMP practice

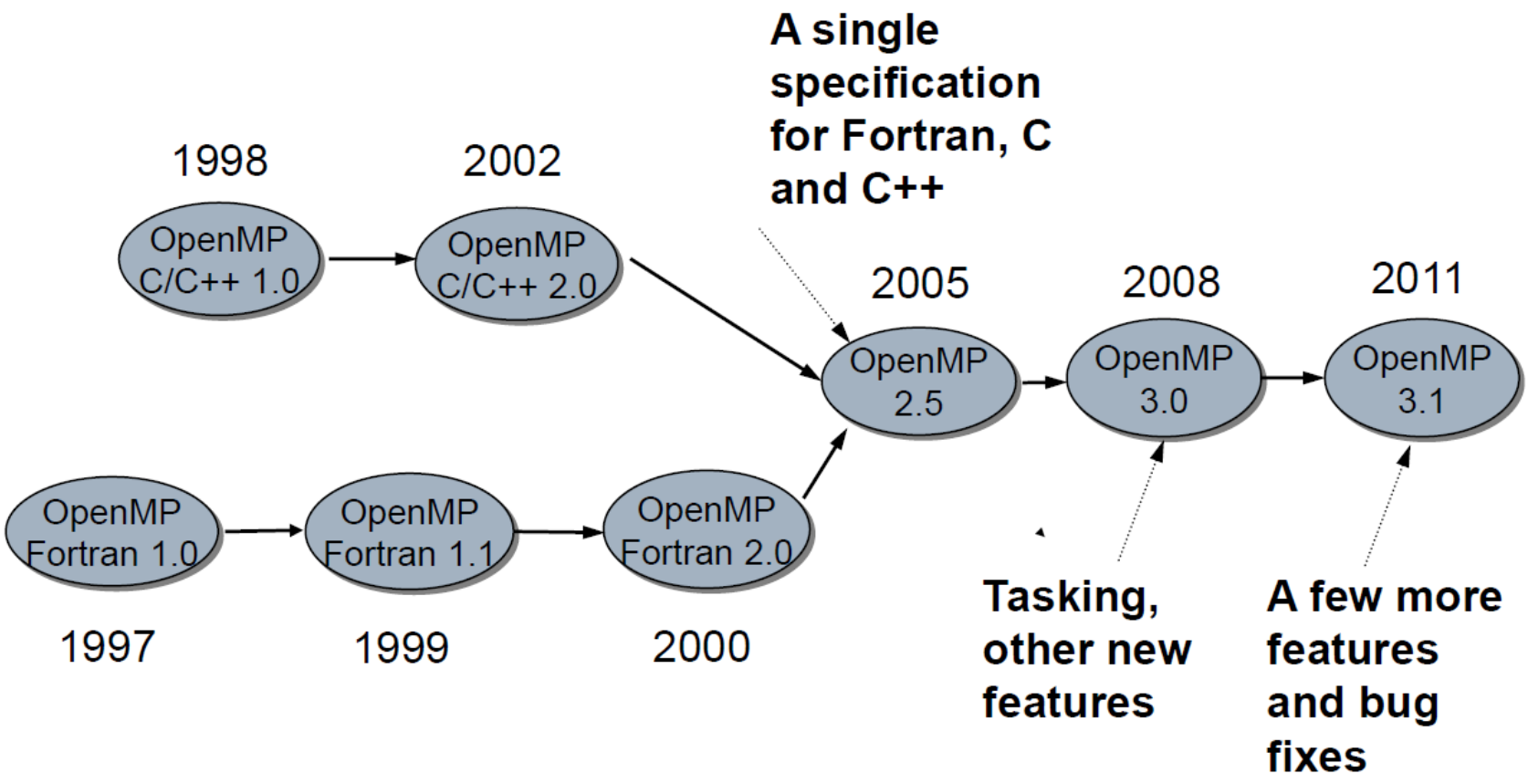
C$OMP PARALLEL COPYIN(/blk/)

C$OMP DO lastprivate(XX)

Nthrds = OMP_GET_NUM_PROCS()

omp_set_lock(lck)

OpenMP Release History

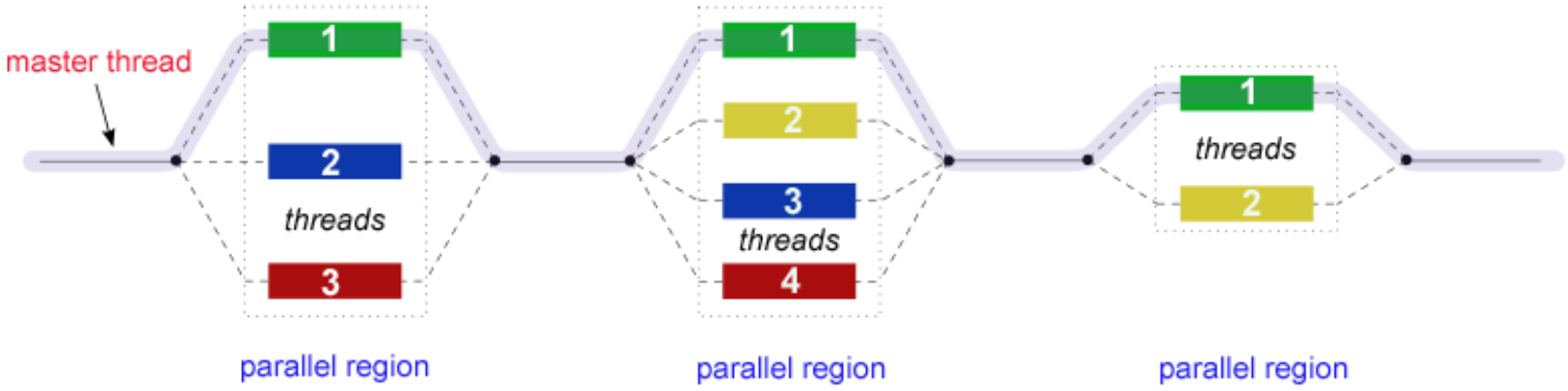


OpenMP Overview: How do Threads Interact?

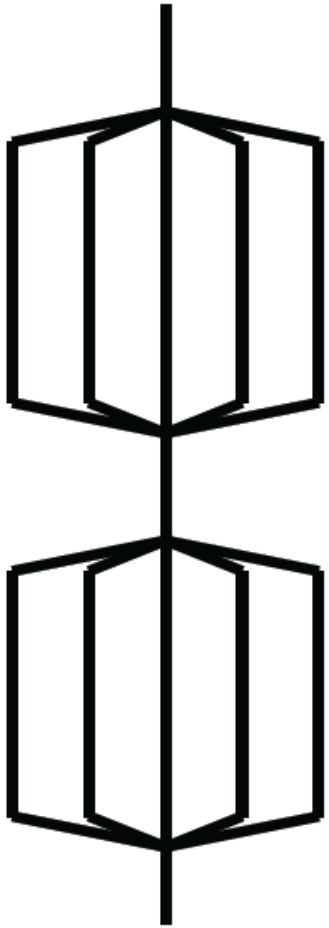
- OpenMP is a multi-threading, shared address model
  - Threads communicate by sharing variables
- Unintended sharing of data causes race conditions
  - Race condition: when the program's outcome changes as the threads are scheduled differently
- To control race conditions
  - Use synchronization to protect data conflicts
- Synchronization is expensive
  - Change how data is accessed to minimize the need for synchronization

Fork-Join Model

- **FORK**
  - The master thread then creates (or awakens) a team of parallel threads.
- **JOIN**
  - When the team threads complete the statements in the parallel region construct, they synchronize and terminate, leaving only the master thread.



Relating Fork/Join to Code



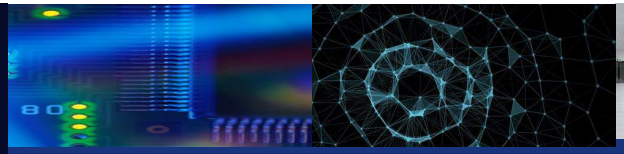
Sequential code

Parallel code

Sequential code

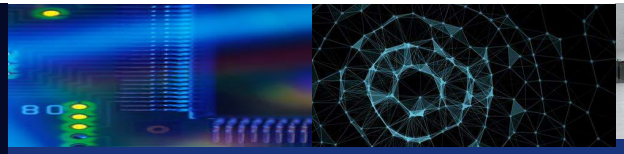
Parallel code

Sequential code



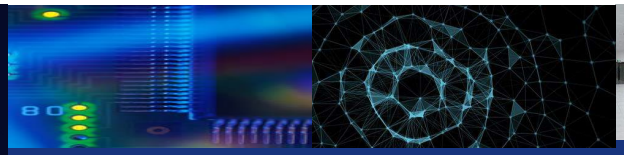
OpenMP Components

- **Three API components**
  - **Compiler directives**
  - **Runtime library routines**
  - **Environment variables**



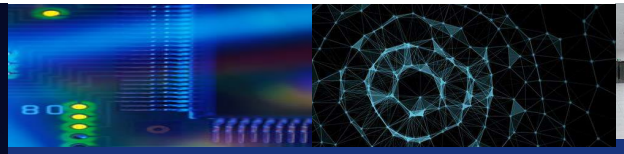
Syntax of Compiler Directives（指令）

- *pragma*: **a C/C++ compiler directive（编译开关）**
  - **(other compiler directives:** #include, #define, **…)**
  - **Stands for "pragmatic information（附注信息）"**
  - **A way for the programmer to communicate with the compiler**
  - **Pragmas are handled by the preprocessor**
  - **Compilers are free to ignore pragmas**
- **All OpenMP pragmas have the syntax:**
  - #pragma omp <directive-name> [clause, ...]
- **Pragmas appear immediately *before* relevant construct**



Hello World

```
#include <omp.h>
#include <stdio.h>
#include <stdlib.h>
int main (int argc, char *argv[]) {
  int nthreads, tid;
  /* Fork a team of threads giving them their own copies of variables */
#pragma omp parallel private(nthreads, tid)
  {
    /* Obtain thread number */
    tid = omp_get_thread_num();
    /* Only master thread does this */
    if (tid == 0) {
      nthreads = omp_get_num_threads();
      printf("Number of threads = %d\n", nthreads);
    }
    printf("Hello World from thread = %d\n", tid);
  }  /* All threads join master thread and disband */
}
```



Output – Non-deterministic!

```
[cong@lnxsrvg1 ~/cs133_examples]$ ./a.out
Hello World from thread = 5
Hello World from thread = 7
Number of threads = 16
Hello World from thread = 0
Hello World from thread = 8
Hello World from thread = 3
Hello World from thread = 1
Hello World from thread = 6
Hello World from thread = 15
Hello World from thread = 14
Hello World from thread = 10
Hello World from thread = 2
Hello World from thread = 9
Hello World from thread = 11
Hello World from thread = 13
Hello World from thread = 12
Hello World from thread = 4
[cong@lnxsrvg1 ~/cs133_examples]$ ./a.out
Hello World from thread = 6
Hello World from thread = 10
Hello World from thread = 8
Number of threads = 16
Hello World from thread = 0
Hello World from thread = 5
Hello World from thread = 14
Hello World from thread = 15
Hello World from thread = 13
Hello World from thread = 12
Hello World from thread = 4
Hello World from thread = 3
Hello World from thread = 9
Hello World from thread = 2
Hello World from thread = 7
Hello World from thread = 1
```

OpenMP Compilers Perform the Translations to Threads (e.g. pthreads)

```
int a, b;
main() {
  // serial segment
  #pragma omp parallel num_threads (8) private (a) shared (b)
  {
    // parallel segment
  }
  // rest of serial segment
}
```

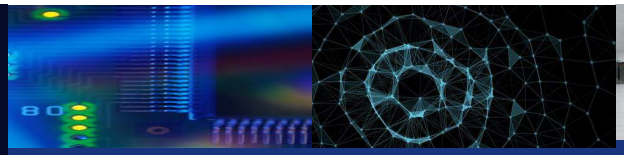
Sample OpenMP program

```
int a, b;
main() {
  // serial segment

  for (i = 0; i < 8; i++)
    pthread_create (........, internal_thread_fn_name, ...);

  for (i = 0; i < 8; i++)
    pthread_join (.......);

  // rest of serial segment

}

void *internal_thread_fn_name (void *packaged_argument) {
  int a;

  // parallel segment

}
```

Code inserted by the OpenMP compiler

Corresponding Pthreads translation

Grama et al., "Introduction to Parallel Computing," Addison Wesley, 2003



Matching Threads with CPUs

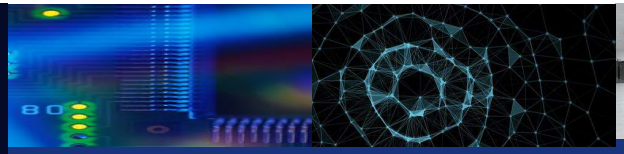
- **Function *omp_get_num_procs* returens the number of physical processors available to the parallel program**

```
int omp_get_num_procs(void);
```

- **Function *omp_set_num_threads* allow you to set the number of threads that should be active in parallel sections of code**

```
void omp_set_num_threads(int t);
```

  - **The function can be called with different arguments at different points in the program**



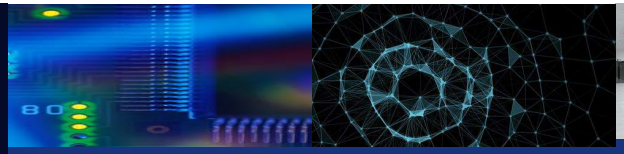
Pragma: *parallel for*

- **The compiler directive**

  #pragma omp parallel for

  tells the compiler that the *for* loop which immediately follows can be executed in parallel

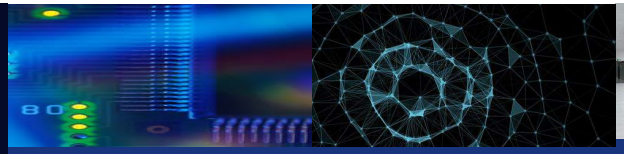
  - **The number of loop iterations must be computable at run time before loop executes**
  - **Loop must not contain a *break*, *return*, or *exit***
  - **Loop must not contain a *goto* to a label outside loop**



Example: *parallel for*

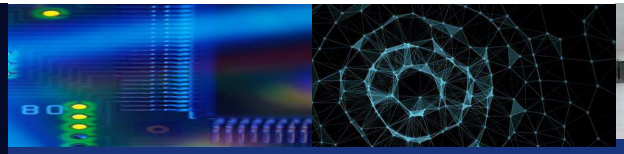
```
int a[1000], b[1000], s[1000];
…
#pragma omp parallel for
for (i = 0; i < 1000; i ++)
    s[i] = a[i] + b[i];
```

- **Threads are assigned an independent set of iterations**
- **Threads must wait at the end of construct**



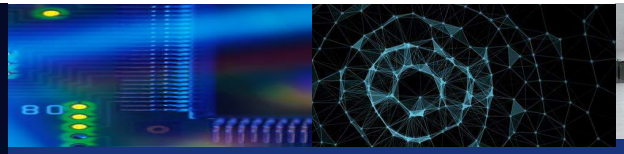
Which Loop to Make Parallel?

```
int main() {
    int i, j, k;
    float **a, **b;
    … // initialize a[][], b[][] as the 1-hop distance matrix
    for (k = 0; k < N; k++) {
        for (i = 0; i < N; i++)
            for (j = 0; j < N; j++)
                a[i][j] = min(a[i][j], b[i][k] + b[k][j]);
        … // copy a[][] to b[][]
    }
```



Which Loop to Make Parallel?

```
int main() {
    int i, j, k;
    float **a, **b;
    … // initialize b[][] as the 1-hop distance matrix
    for (k = 0; k < N; k++)          // Loop-carried dependences
        for (i = 0; i < N; i++)      // Can execute in parallel
            for (j = 0; j < N; j++)  // Can execute in parallel
                a[i][j] = min(a[i][j], b[i][k] + b[k][j]);
    … // copy a[][] to b[][]
```

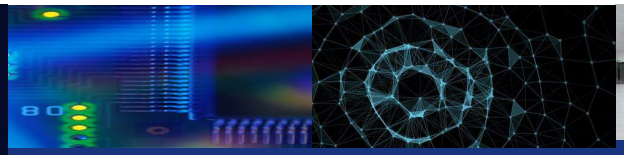


Minimizing Threading Overhead

- **There is a fork/join for every instance of**

```
#pragma omp parallel for
for (…) {
…
}
```

- **Since fork/join is a source of overhead, we want to maximize the amount of work done for each fork/join; i.e., the *grain size***
- **Hence we choose to make the middle loop parallel**



Almost Right, but Not Quite

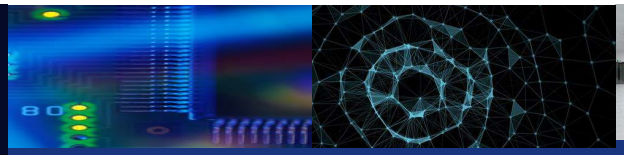
```
int main() {
    int i, j, k;
    float **a , **b;
    … // initialize b[][] as the 1-hop distance matrix
    for (k = 0; k < N; k++) {
        #pragma omp parallel for
        for (i = 0; i < N; i++)
            for (j = 0; j < N; j++)
                a[i][j] = min(a[i][j], b[i][k] + b[k][j]);
        … // copy a[][] to b[][]
    }
```

Problem: j is a shared variable

Clause: private

- **Clause**
  - **An optional, additional component to a pragma**
- **Private clause: directs compiler to make one or more variables private**

#pragma omp … private (<variable list>)



Problem Solved with private Clause

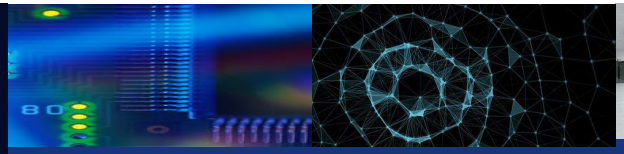
```
int main() {
    int i, j, k;
    float **a , **b;
    … // initialize b[][] as the 1-hop distance matrix
    for (k = 0; k < N; k++) {
        #pragma omp parallel for private (j)
        for (i = 0; i < N; i++)
            for (j = 0; j < N; j++)
                a[i][j] = min(a[i][j], b[i][k] + b[k][j]);
        … // copy a[][] to b[][]
    }
```

Tell compiler to make listed variables private

Another Example

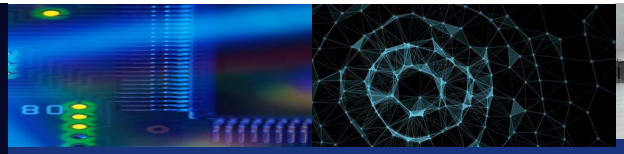
```
int i;
float *a, *b, *c, tmp;
…
for (i = 0; i < N; i++) {
tmp = a[i] / b[i];
c[i] = tmp * tmp;
}
```

Loop is perfectly parallelizable except for shared variable ***tmp***



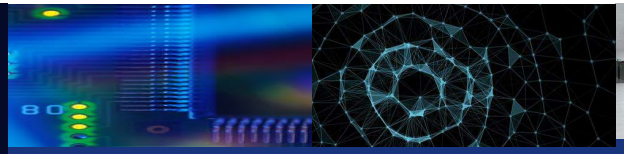
Solution

```
int i;
float *a, *b, *c, tmp;
…
#pragma omp parallel for private (tmp)
for (i = 0; i < N; i++) {
tmp = a[i] / b[i];
c[i] = tmp * tmp;
}
```



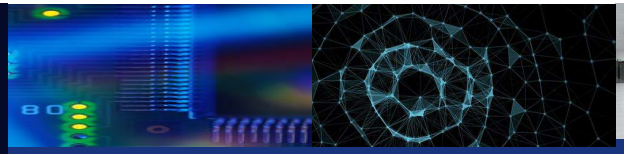
More About Private Variables (私有变量)

- Each thread has its own copy of the private variables
- If *j* is declared private, then inside the *for* loop no thread can access the "other" *j* (the *j* in shared memory)
- No thread can use a previously defined value of *j*
- No thread can assign a new value to the shared *j*
- Private variables are **undefined** at loop entry and loop exit, reducing execution time



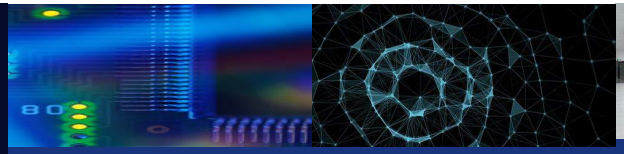
Clause: *firstprivate*

- **The *firstprivate* clause tells the compiler that the private variable should inherit the value of the shared variable upon loop entry**
- **The value is assigned once per thread, not once per loop iteration**



Example: *firstprivate*

```
a[0] = 0.0;
for (i = 1; i < N; i++)
    a[i] = alpha(i, a[i-1]);

#pragma omp parallel for firstprivate (a)
for (i = 0; i < N; i++)
    b[i] = beta(i, a[i]);
    a[i] = gamma(i);
    c[i] = delta(a[i], b[i]);
}
```



Clause: *firstprivate*

- pragma omp … firstprivate(x)
  - ***x* is a fundamental data type**
    - **Private *x* is directly copied from the shared *x***
  - ***x* is an array**
    - **Copy the data with *sizeof(x)* to the private memory**
  - ***x* is a pointer**
    - **Private *x* points to the same location as the shared *x***
  - ***x* is a class instance**
    - **Copy constructor is called to create the private *x***

Clause: *firstprivate*

- pragma omp … firstprivate(x)
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    - **Private *x* points to the same location as the shared *x***
  - ***x* is a class instance**
    - **Copy constructor is called to create the private *x***

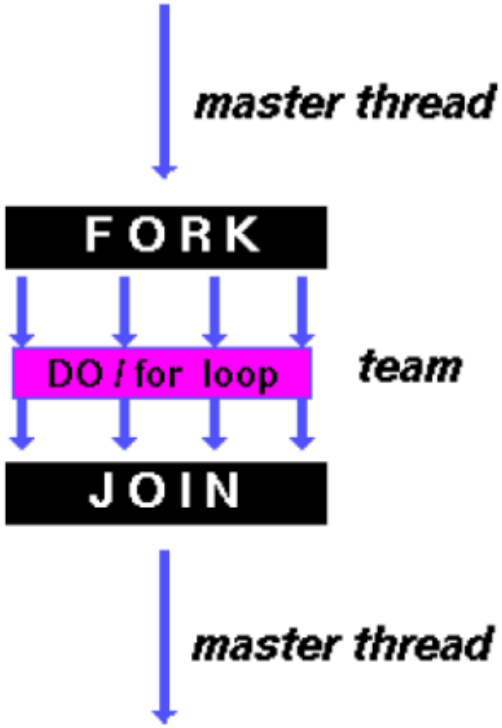
Clause: *lastprivate*

- **The *lastprivate* clause tells the compiler that the value of the private variable after the *sequentially last* loop iteration should be assigned to the shared variable upon loop exit**
  - **In other words, when the thread responsible for the sequentially last loop iteration exits the loop, its copy of the private variable is copied back to the shared variable**

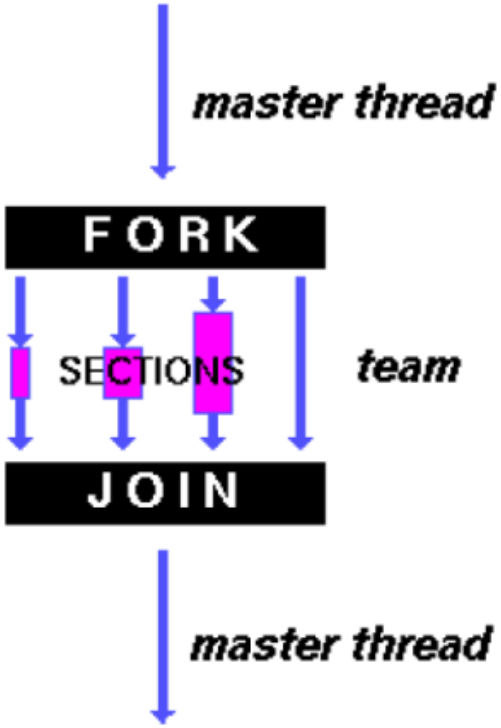
Example: *lastprivate*

```
#pragma omp parallel for lastprivate (x)
for (i = 0; i < N; i++)
x = foo(i);
y[i] = bar(i, x);
}
last_x = x; // == foo(N-1)
```

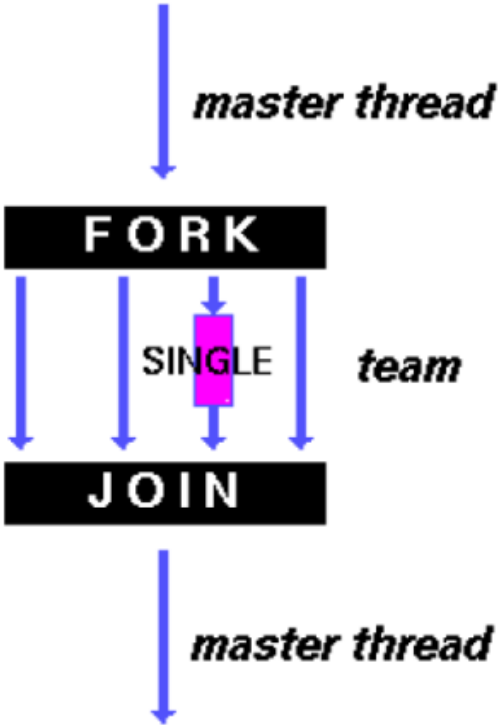
Work-Sharing Constructs



a type of "data parallelism".



a type of "functional parallelism".



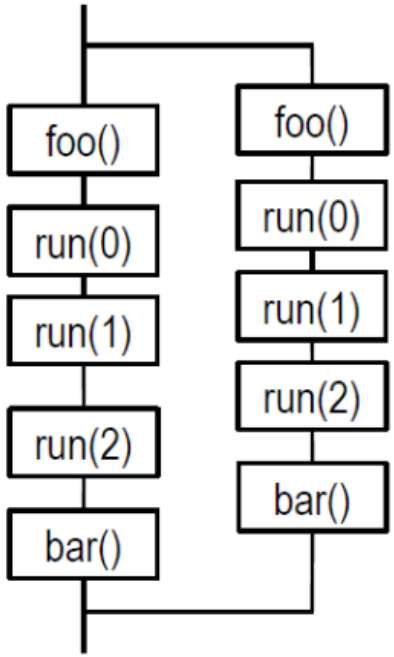
serializes a section of code

Pragma: *parallel*

- **In the effort to increase grain size, sometimes the code that should be executed in parallel goes beyond a single *for* loop**
  - **The *parallel* pragma is used when a block of code should be executed in parallel**
  - **SPMD-style programming**
  - **Single program, multiple data**

Pragma parallel ***and Pragma*** for

- #pragma omp for
  - **Used inside a block of code already marked by *parallel***
  - **Distribute the iterations to the active threads**

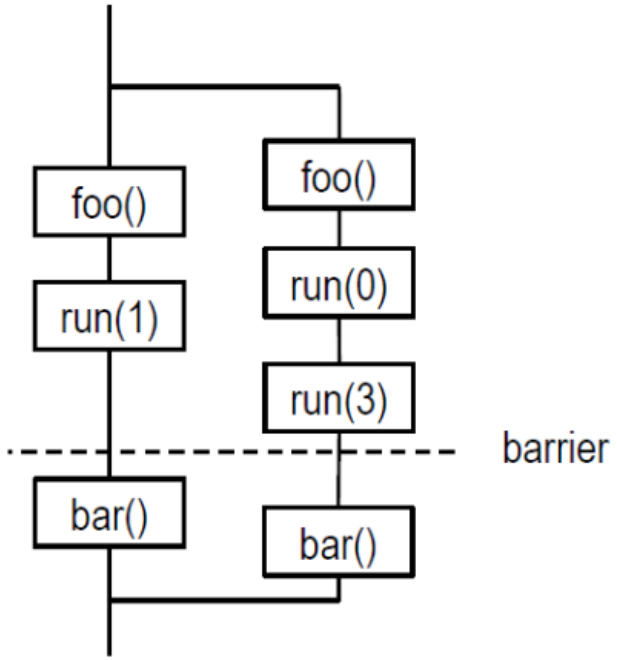


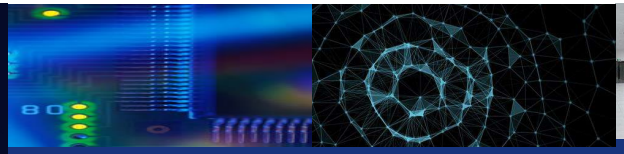

```
#pragma omp parallel \
 num_threads(2)
{
    foo();
    for (int i = 0; i < 3; i++)
        run(i);
    bar();
}
```

Pragma parallel ***and Pragma*** for

- #pragma omp for
  - **Used inside a block of code already marked by *parallel***
  - **Distribute the iterations to the active threads**

```
#pragma omp parallel \
 num_threads(2)
{
     foo();
     #pragma omp for
     for (int i = 0; i < 3; i++)
          run(i);
     bar();
}
```





Pragma: single

- **The *single* pragma is used inside a parallel block of code**
  - **It tells the compiler that only a single thread should execute the statement or block of code immediately following**
  - **May be useful when dealing with sections of code that are not thread safe (such as I/O)**
  - **Threads in the team that do not execute the** single **directive, wait at the end of the enclosed code block, unless a** nowait **clause is specified.**

Example: *master* ***and*** *single nowait*

```
tid = omp_get_thread_num();
if (tid == 0) {
    nthreads = omp_get_num_threads();
    printf("Number of threads = %d\n", nthreads);
}
```

=

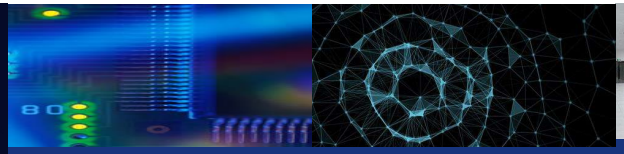
```
#pragma omp master
{
    nthreads = omp_get_num_threads();
    printf("Number of threads = %d\n", nthreads);
}
```

≈

```
#pragma omp single nowait
{
    nthreads = omp_get_num_threads();
    printf("Number of threads = %d\n", nthreads);
}
```

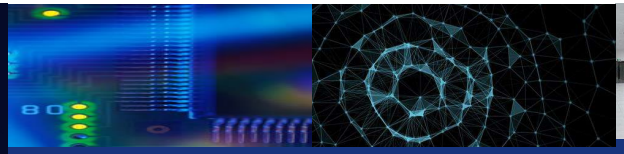
Pragma: *sections* ***and*** *section*

- **Directive** sections **specifies that the enclosed section(s) of code are to be divided among the threads in the team.**
- **Independent section directives are nested within a** sections **directive.**
  - **Each** section **is executed once by a thread in the team.**
  - **Different sections may be executed by different threads.**
  - **It is possible for a thread to execute more than one section if it is quick enough and the implementation permits such.**



Example: *sections* ***and*** *section*

```
#pragma omp parallel shared(a,b,c,d) private(i)
{
#pragma omp sections
   {
#pragma omp section
     {
      for (i=0; i<N; i++)
        c[i] = a[i] + b[i];
     }
#pragma omp section
     {
      for (i=0; i<N; i++)
        d[i] = a[i] * b[i];
     }
   }  /* end of sections */
 }  /* end of parallel section */
```

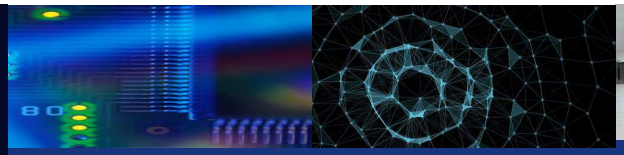


Clause: *reduction*（归并）

- **Reductions are so common that OpenMP provides a reduction clause for the *parallel*, *for*, and *sections***

  #pragma omp … reduction (op : list)

  - **A private copy of each list variable is created and initialized depending on the *op***
    - **The identity value *op* (e.g., 0 for addition)**
  - **These copies are updated locally by threads**
  - **At end of construct, local copies are combined through *op* into a single value and combine the value in the original shared variable**



C/C++ Reduction Operation

结合率

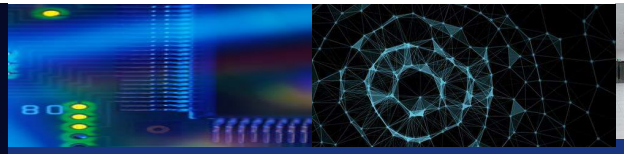
- Reduction with an associative binary operator $\oplus$

$$a_1 \oplus a_2 \oplus a_3 \oplus \ldots \oplus a_n$$

- A range of associative and commutative operators can be used with reduction
- Initial values are the ones that make sense

Operator	Initial Value
+	0
*	1
-	0
^	0

Operator	Initial Value
&	~0
|	0
&&	1
||	0

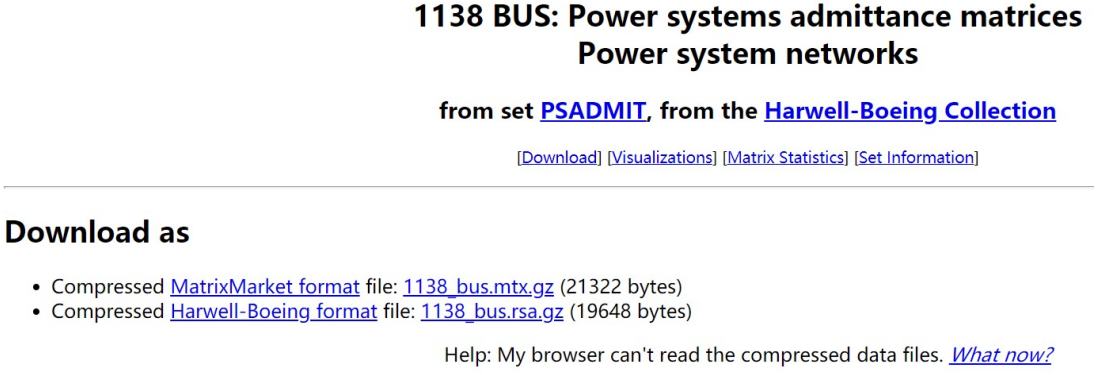


Strengths and Weaknesses of OpenMP

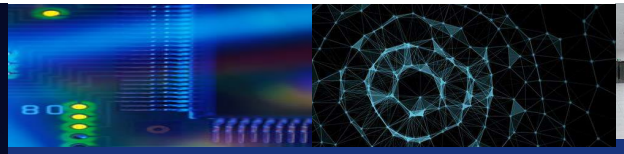
- **Strengths**
  - **Incremental parallelization & sequential equivalence**
  - **Well-suited for domain decompositions**
  - **Available on *nix and Windows**
- **Weaknesses**
  - **Not well-tailored for functional decompositions**
  - **Compilers do not have to check for such errors as deadlocks and race conditions**

Homework

1. Consider a sparse matrix stored in the compressed row format (you may find a description of this format on the web or any suitable text on sparse linear algebra). Write an OpenMP program for computing the product of this matrix with a vector. Download sample matrices from the Matrix Market (http://math.nist.gov/MatrixMarket/) and test the performance of your implementation as a function of matrix size and number of threads.



2. Implement a producer-consumer framework in OpenMP using sections to create a single producer task and a single consumer task. Ensure appropriate synchronization using locks. Test your program for a varying number of producers and consumers.



Thank You !