



Theory of Answer Set Programming

Negation as Failure: Informal Introduction

Objectives



Objective

Explain the intuitive
meaning of negation as
failure



Stable Models of Programs with Negation

Recall: Syntax of Propositional Rules

- | A **(propositional) rule** is either
 - a propositional formula F that does not contain any implication symbol, or
 - a formula of the form $F \leftarrow G$ where F and G are implication-free.
- | A **propositional program** is a set of propositional rules.

Prolog vs. ASP

?p

p :- not q

q :- not p

P ← ¬q
q ← ¬P

Prolog does not terminate on
query p or q

?- p.

ERROR: Out of local stack

Exception: (729,178)

clingo returns

Answer: 1

p

Answer: 2

q

Finite ASP programs are
guaranteed to terminate

Negation as Failure

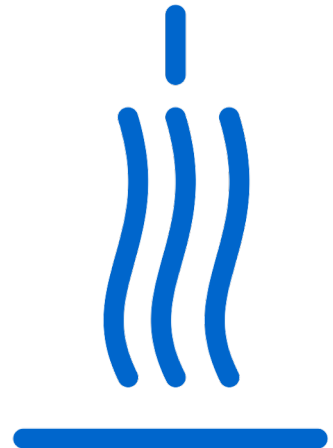
Q: How do we extend the definition of a stable model in the presence of negation?

$$\begin{array}{l} p, \\ q, \\ r \leftarrow p, \\ s \leftarrow q \end{array}$$
$$\{p, q, r, s\}$$
$$\begin{array}{l} p, \\ q, \\ r \leftarrow p \wedge \neg s, \\ s \leftarrow q. \end{array}$$
$$\{p, q, s\}$$
$$\begin{array}{l} p, \\ q, \\ r \leftarrow p \wedge \neg s, \end{array}$$
$$\{p, q, r\}$$
$$\begin{array}{l} p, \\ r \leftarrow p \wedge \neg s, \\ s \leftarrow q. \end{array}$$
$$\{p, r\}$$

Add r to the model if p is included under the condition that s is not included in the model and will not be included in the future.

Informal Reading: Rationality Principle

- | Informally, program Π can be viewed as a specification for stable models--sets of beliefs that could be held by a rational reasoner associated with Π .
- | Stable models will be represented by collections of atoms.
- | In forming such sets the reasoner must be guided by the following informal principles:
 - Satisfy the rules of Π .
 - If one believes in the body of a rule, one must also believe in its head.
 - Adhere to the “the rationality principle.”
 - “Believe nothing you are not forced to believe.”



Negation as Failure, cont'd

| $p \leftarrow q$ SM
 \emptyset

| $p \leftarrow \neg q$ $\exists p \zeta$

| $p \leftarrow \neg q$ $\exists p \zeta$

$q \leftarrow \neg p$ $\exists q \zeta$

| $p \leftarrow q$
 $q \leftarrow \neg r$ $\exists p, q \zeta$

| $p \leftarrow \neg q$ $\exists q \zeta$
 $q \leftarrow \neg r$

Lecture Wrap-Up

