

Practice of Answer Set Programming

Graph Problems in ASP (II)

Objectives



Objective

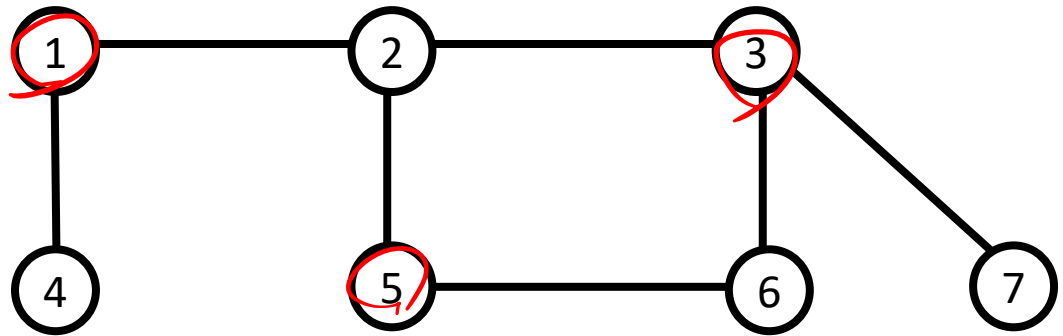
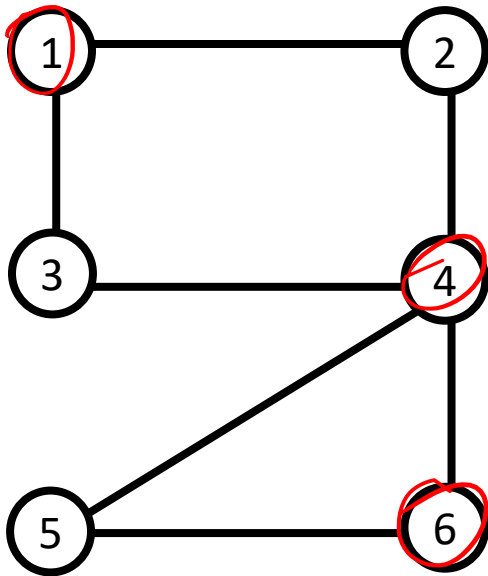
Use ASP to solve
graph-related
problems



Vertex Cover

Definition

- A **vertex cover** of a graph is a set of vertices such that every edge of the graph has at least one endpoint in that set
- We would like to find a vertex cover of size n



Example

```
% in/1 is a set consisting of n vertices of G  
{in(X) : vertex(X)} = n.
```

```
% covered/2 is the set of edges of G that have  
% an endpoint in in/1
```

```
covered(X,Y) :- edge(X,Y), in(X).
```

```
covered(X,Y) :- edge(X,Y), in(Y).
```

```
% every edge of G is in covered/2  
:- edge(X,Y), not covered(X,Y).
```



Graph

% Nodes

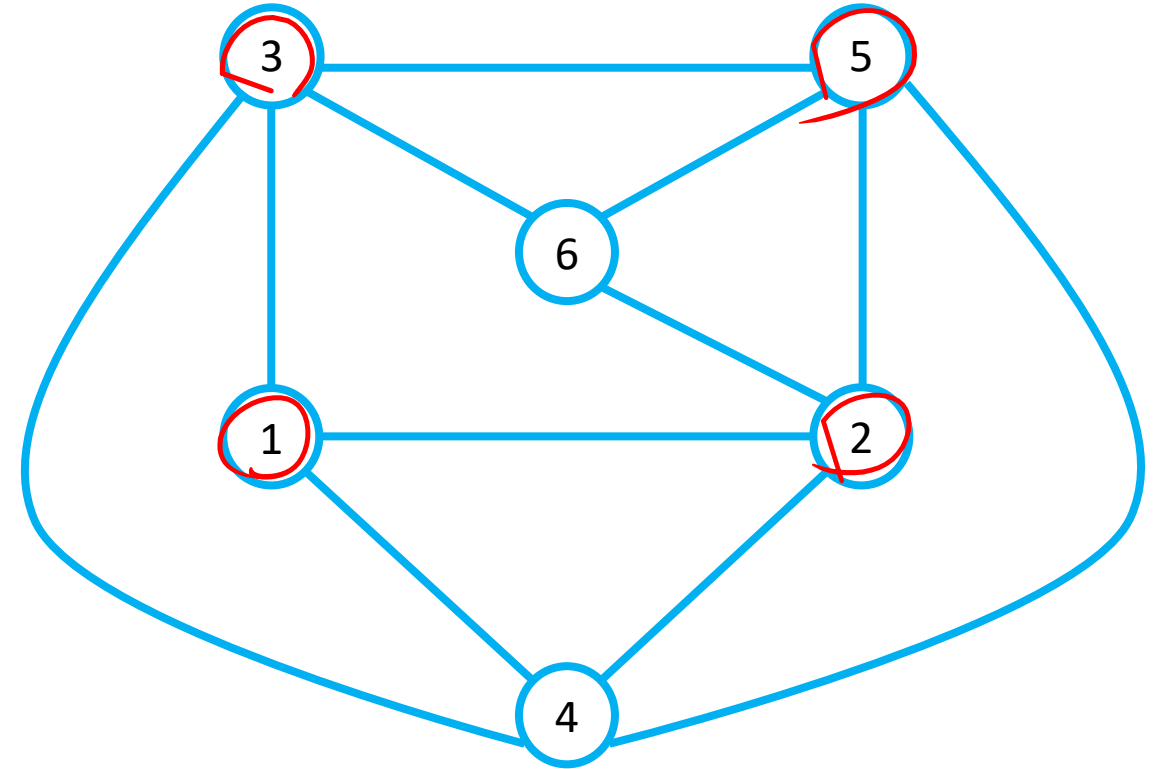
```
vertex(1..6).
```

% Edges

```
edge(1,(2;3;4)). edge(4,(1;2)).
```

```
edge(2,(4;5;6)). edge(5,(3;4;6)).
```

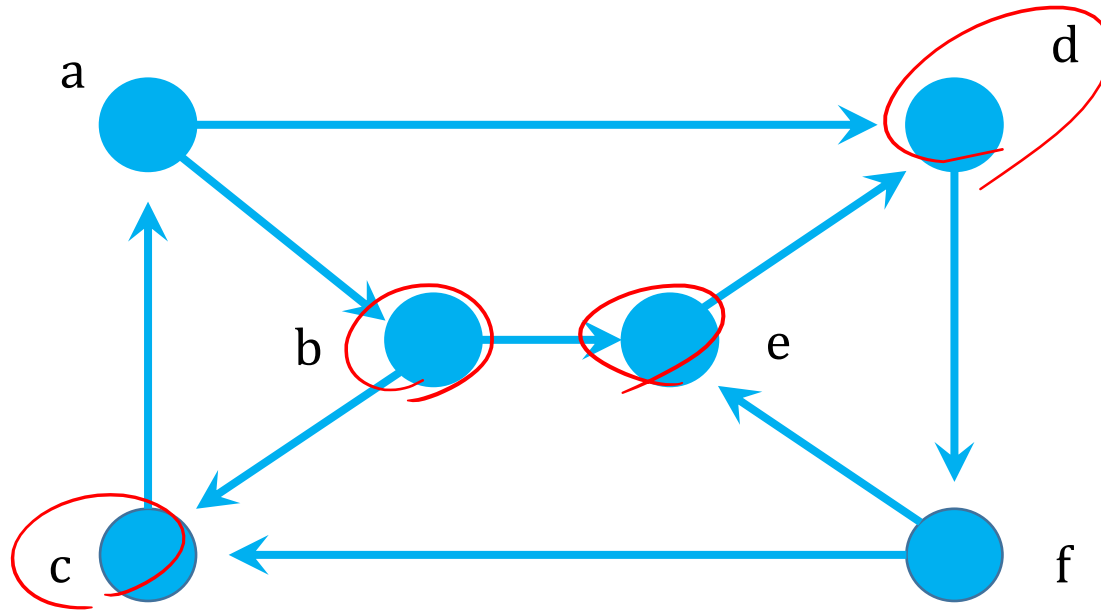
```
edge(3,(1;4;5)). edge(6,(2;3;5)).
```



Graph

`vertex(a; b; c; d; e; f).`

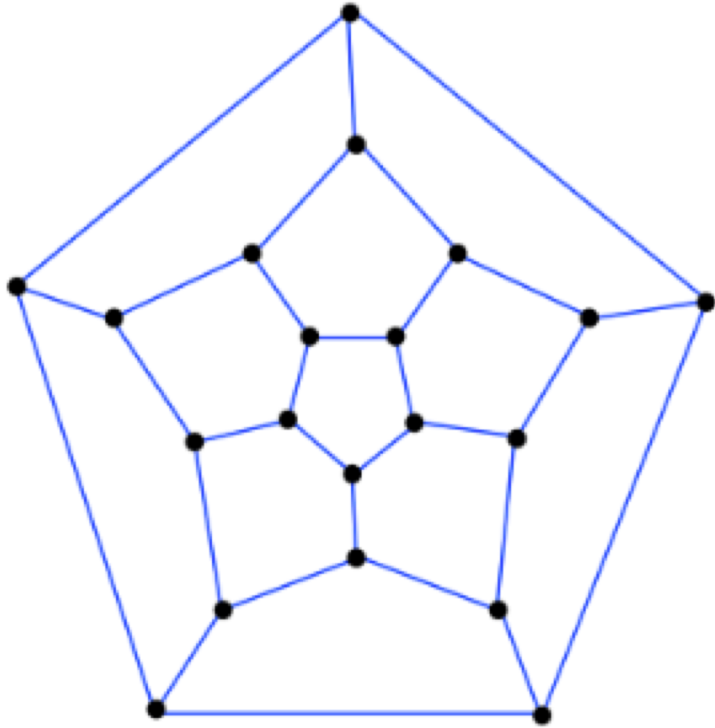
`edge(a,b; b,c; c,a; d,f; f,e; e,d; a,d; f,c; b,e).`



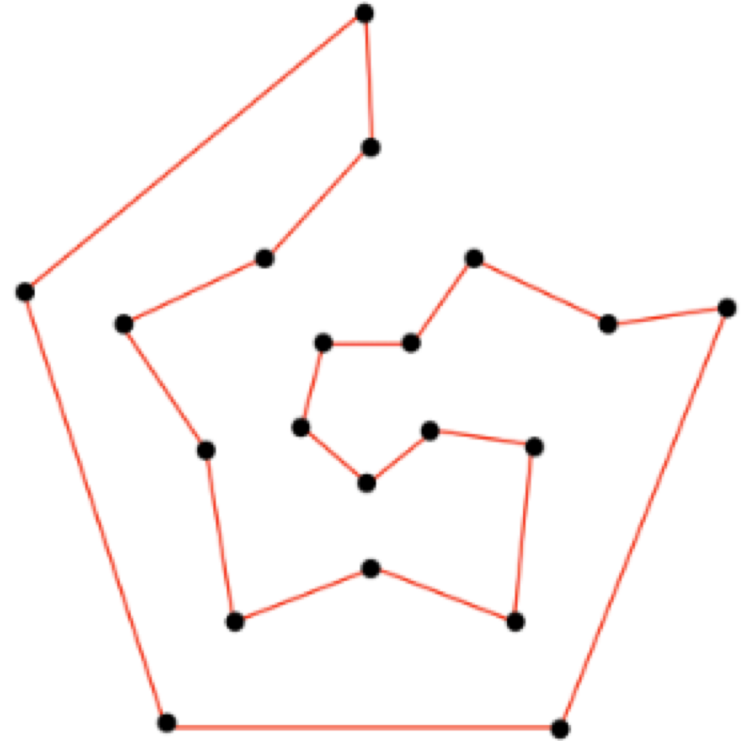
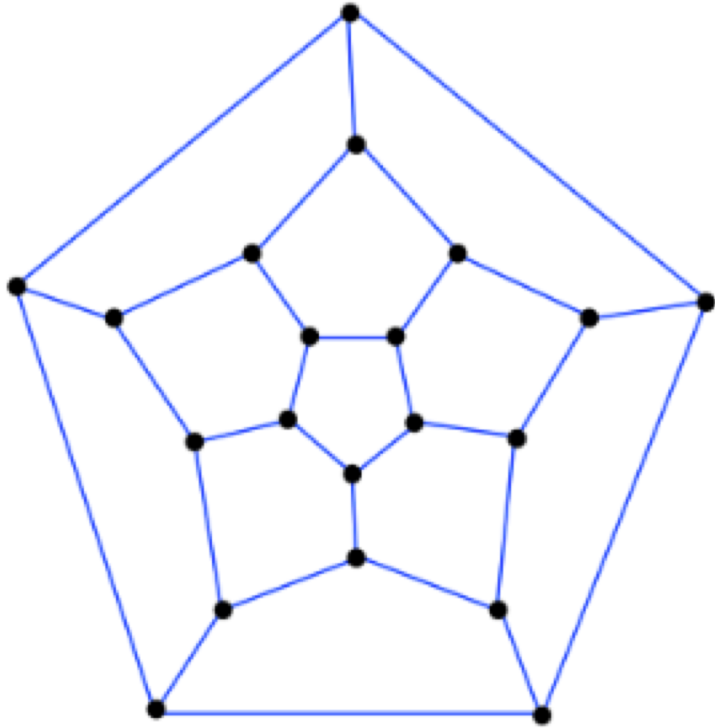


Hamiltonian Cycle

Introduction



Introduction



Graph

% Nodes

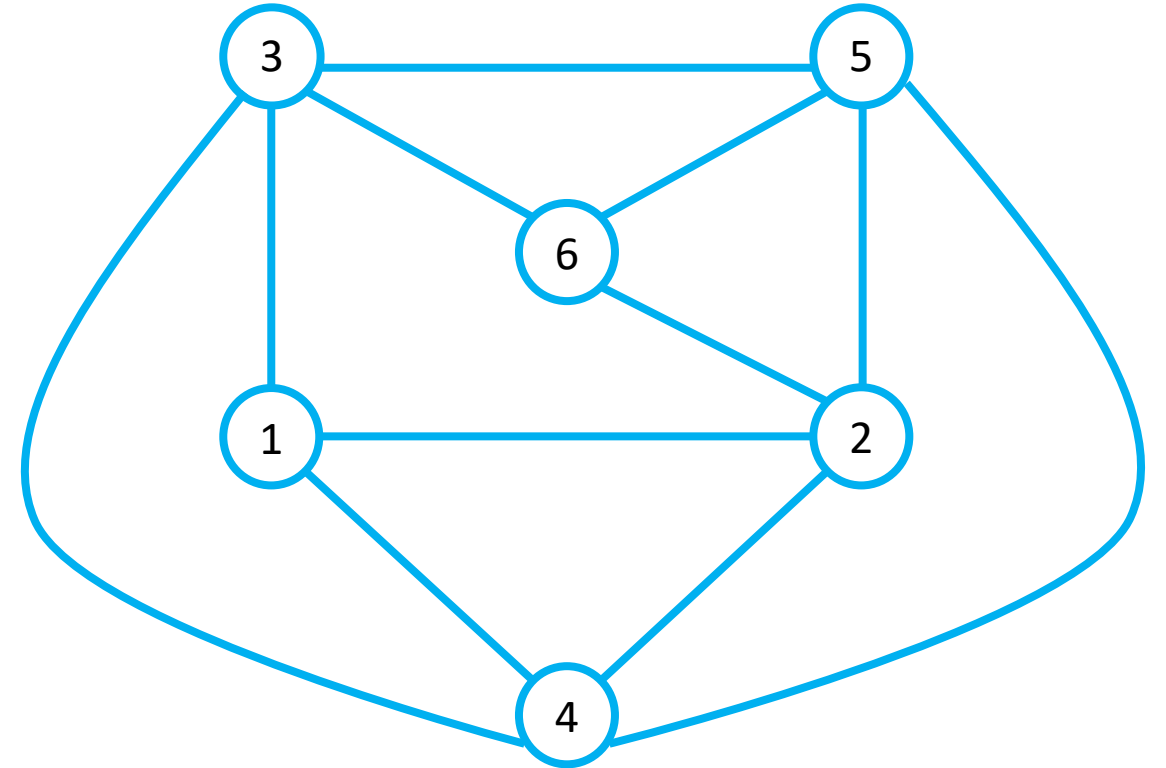
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% Edges

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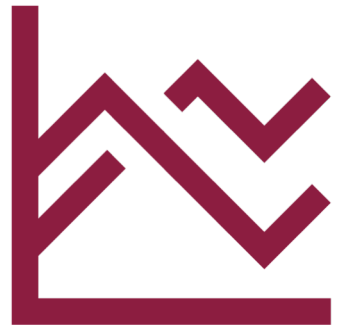
```
edge(3,(1;4;5)). edge(6,(2;3;5)).
```



Hamiltonian Cycle

A Hamiltonian cycle in a graph G is a subgraph C of G that has the same set V of vertices as G and satisfies two conditions:

1. Every vertex has only one incoming edge and only one outgoing edge in the cycle.
2. Every vertex is reachable in C from some fixed vertex, say 1.



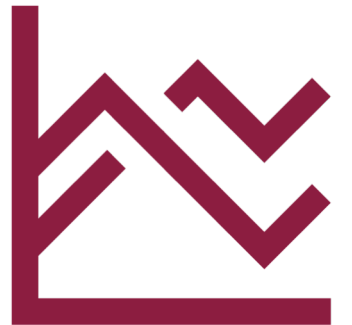
Hamiltonian Cycle in ASP

```
% in(X,Y) is the set of edges included in the cycle and
% satisfies Condition 1.
{in(X,Y) : edge(X,Y)} = 1 :- vertex(X).
{in(X,Y) : edge(X,Y)} = 1 :- vertex(Y).

% Define reachability recursively
reachable(X) :- in(1,X).
reachable(Y) :- in(X,Y), reachable(X).

% Satisfies Condition 2
:- not reachable(X), vertex(X).

% Display
#show in/2.
```



Wrap-Up

