First Order Logic and KR Herbrand Models



Objectives



Objective

Explain how Herbrand models are defined and reduces the complexity of FOL reasoning

Herbrand Models

A Hebrand interpretation is a special case of first-order interpretation

- A Herbrand interpretation of signature σ (containing at least one object constant) is an interpretation of σ such that
 - its universe (Herbrand Universe) is the set of all ground (i.e., variable-free) terms of σ ,
 - every ground term is interpreted as itself ($t^{I} = t$)

$$\alpha^{x} = \alpha$$

Herbrand Models, cont'd

Herbrand interpretations of the signature {P, a}

$$|I| = 3a\zeta$$

$$Q^{I} = \alpha$$

$$P^{I}(\alpha) = \pm$$

$$|I| = 3a4$$

$$a^{T} = a$$

$$p^{T}(a) = f$$

$$2$$

Herbrand models of the signature {P, a} that satisfies formula P(a)



Herbrand Models, cont'd

Herbrand interpretations of the signature $\{P, a, b\}$ $|I| = \beta a, b \}$ $\alpha^{x} = \alpha$ $\beta^{x} = b$ $\rho^{x}(a) = f$ $\rho^{x}(b) = f$ $\rho^{x}(b) = f$ $\rho^{x}(b) = f$

$$|I| = 34, b$$

$$\alpha^{x} = a$$

$$b^{x} = b$$

$$P^{T}(a) = f$$

$$P^{T}(b) = f$$

$$p^{r}(a) = f$$

$$p^{r}(b) = x$$

$$p^{\mathrm{T}}(a) = \pm p^{\mathrm{T}}(b) = \pm p^{\mathrm{T}}(b)$$

Herbrand models of the signature {P, a, b} that satisfies formula P(a)

An Herbrand interpretation can be identified with the set of ground atoms to which it assigns the value true.

Exercise

1. Let F_1 be $P(a) \wedge \exists x \neg p(x)$.

Find the Herbrand models of F₁ whose signature is {a,P}.

$$|I| = \lambda a \lambda$$
 $\mathcal{P}^{r}(a) = t$

Find the Herbrand models of F₁ whose signature {a,b,P}

$$|T| = 3a, b$$
 $p^{T}(a) = f$

$$p^{T}(b) = f$$

$$|T| = \lambda a, b; \quad \text{or } p^{x}(a) = f \quad \text{or } p^{x}(b) = f \quad \text{or$$

2. Let F_2 be $P(a) \land \neg P(b) \land \exists x \neg P(x)$.

Find the Herbrand models of F_2 whose signature is $\{a,b,P\}$.

Entailment and Herbrand Models

Without functions, entailment restricted to Herbrand models is decidable

- Herbrand models are finitely enumerable

With functions, this is not the case

When the Herbrand universe is finite, quantified formulas can be identified with propositional formulas

- $\forall x P(x)$ vs. $P(a) \land P(b) \land P(c)$
- $-\exists x P(x) \text{ vs. } P(a) \vee P(b) \vee P(c)$

Wrap-Up

