

**Birla Institute of Technology & Science, Pilani**  
**Data Structures & Algo (CS F211)**  
**Online Lab test 2() OPEN book      Duration: 90 min      Date: 25<sup>th</sup> April, 2015**

A completely connected component  $K_n$  of a Graph  $G(V,E)$  is a connected component of the graph containing  $n$  vertices such that every pair of vertices in  $K_n$  have an edge connecting them.

**Problem statement:**

Given an undirected graph, find the number of completely connected components of the graph of size 1,2,3,4, i.e, count of  $K_1$   $K_2$   $K_3$   $K_4$ .

**File formats:**

**Input : input.txt**

line 1: number of vertices( $n$ ) [numbered from 0 to  $n$ ]

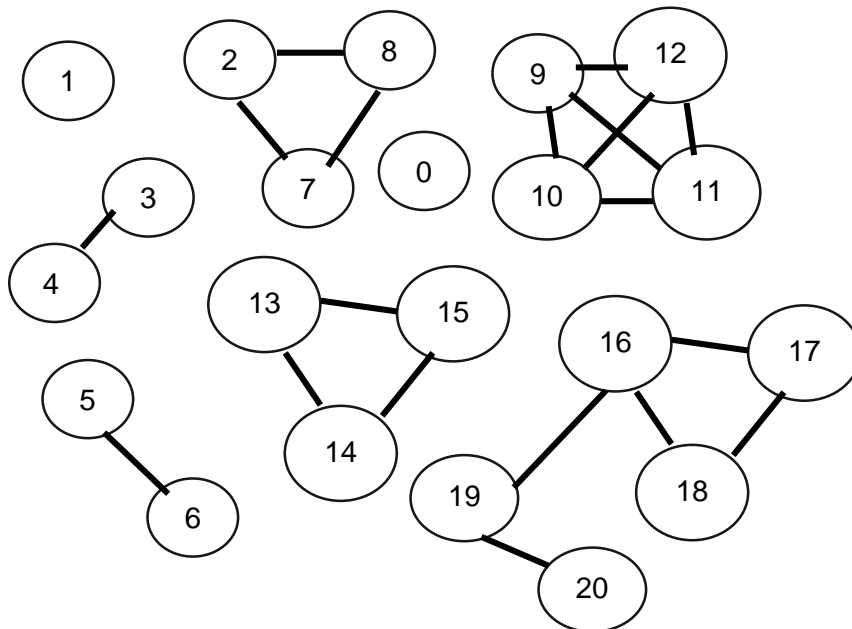
line 2: number of edges( $m$ )

next  $m$  lines, an edge between nodes  $u$  and  $v$

**Output : output.txt**

Four lines, containing the count of  $K_1$ ,  $K_2$ ,  $K_3$ ,  $K_4$  in the four lines.

**A sample scenario**



**Output for the above graph:**

2  
2  
2  
1

$K_1=\{1\}, \{0\}$   
 $K_2=\{5,6\}, \{3,4\}$   
 $K_3=\{2,7,8\}, \{13,14,15\}$   
 $K_4=\{9,10,11,12\}$

**Sample Input:**

21  
19  
3 4  
2 8  
7 8  
5 6  
2 7  
9 11  
10 11  
12 11  
13 15  
14 15  
9 10  
9 12  
10 12  
13 14  
16 17  
17 18  
16 19  
19 20  
16 18

## Overview of the solution:

The graph is represented as adjacency list. The relevant code for creating the graph is already provided to you.

The solution for this problem is achieved via depth first traversal.

In the depth first traversal (DFT), the depth first search (DFS) starting at node  $v$  is used to identify connected component in which the vertex is part of. The code for depth first traversal is already provided to you.

In the depth first search, you must keep track of the count of vertices in the current connected component. If the count is less than or equal to 4, this DFS call should return the list of vertices (node numbers) in the connected component. If count is larger than 4, then it should return -1 in all cells.

The DFT code verifies whether the connected component is  $K_1$ ,  $K_2$ ,  $K_3$ ,  $K_4$  and appropriately modifies the count.

### Codes to write:

- 1) **int isNeighbourOf(Graph G, int u, int v)**: This function determines whether the vertices identified by  $u$  and  $v$  are neighbours or not. This function returns 1 (for true) or 0 (for false).
- 2) **int DFS(Graph G, int v, int result[4])**: This function performs DFS search starting at node with node number as  $v$ . The list of the vertices visited are maintained in the array **result**. If only one node is in this connected component, only **result[0]** will be having valid node number. If two nodes are in this connected component, **result[0]** and **result[1]** will be having valid node numbers. If three nodes are in this connected component, **result[0]**, **result[1]** and **result[2]** will be having valid node numbers.
- 3) **int checkK3(Graph G, int result[4])**: This function verifies whether the three vertices in the array **result** form a  $K_3$  or not. This function returns 1 (for true) or 0 (for false).
- 4) **int checkK4(Graph G, int result[4])**: This function verifies whether the four vertices in the array **result** form a  $K_4$  or not. This function returns 1 (for true) or 0 (for false).

All code is provided in a file named **driver.c**

**Deliverables:** **driver.c** , **input.txt**, **output.txt**