

1) Let $X \sim \text{Hypergeometric}(b, r, k)$.

$$\therefore \text{PMF of } X = P_X(x) = \begin{cases} \frac{\binom{b}{x} \binom{r}{k-x}}{\binom{b+r}{k}} & \text{for } x \in R_X \\ 0 & \text{o.w.} \end{cases}$$

where $R_X = \text{range of } X = \left\{ \max(0, k-r), \dots, \min(k, b) \right\}^{(\max(0, k-r)+1)}$

To prove that $P_X(x)$ is valid, we need to show that (i) $P_X(x) \geq 0, \forall x$.

$$(ii) \sum_{x \in R_X} P_X(x) = 1$$

Proof for (i)

$\binom{b}{x}, \binom{r}{k-x}, \binom{b+r}{k} \in$ set of non-negative integers.

I \because they represent no. of ways some objects can be combined.

For any value of x , $\frac{\binom{b}{x} \binom{r}{k-x}}{\binom{b+r}{k}} \geq 0 \Rightarrow P_X(x) \geq 0$. (Proved)

Proof for (ii)

$$\sum_{x \in R_X} p_X(x) = \sum_{x \in R_X} \frac{\binom{b}{x} \binom{r}{k-x}}{\binom{b+r}{k}}$$

$$= \binom{b+r}{k}^{-1} \sum_{x \in R_X} \binom{b}{x} \binom{r}{k-x} \dots (e1)$$

By binomial theorem,

$$\begin{aligned} \sum_{k=0}^{b+r} \binom{b+r}{k} y^k &= (1+y)^{b+r} \\ &= (1+y)^b (1+y)^r \\ &= \left(\sum_{i=0}^b \binom{b}{i} y^i \right) \left(\sum_{j=0}^r \binom{r}{j} y^j \right) \end{aligned}$$

] Again by
binomial
theorem

$$= \sum_{k=0}^{b+r} \left(\sum_{x=0}^k \binom{b}{x} y^x \cdot \binom{r}{k-x} y^{k-x} \right)$$

$\therefore x$ replaces i ; $(k-x)$ replaces j ; All the bounds
(D + n)

<contd. >

$$= \sum_{k=0}^{b+r} \left(\sum_{x=0}^k \binom{b}{x} \binom{r}{k-x} y^k \right)$$

$$= \sum_{k=0}^{b+r} \left(\sum_{x=0}^k \binom{b}{x} \binom{r}{k-x} y^k \right)$$

$$\Rightarrow \sum_{k=0}^{b+r} \binom{b+r}{k} y^k = \sum_{k=0}^{b+r} \left(\sum_{x=0}^k \binom{b}{x} \binom{r}{k-x} \right) y^k$$

$$\Rightarrow \binom{b+r}{k} = \sum_{\substack{k=0 \\ x=0}}^k \binom{b}{x} \binom{r}{k-x} \dots \text{(e2)}$$

Note: Equation (e2) is known as Vandermonde's

identity when $m, n, r \in \mathbb{N}_0 = \{0, 1, 2, 3, \dots\}$.
 You may check wikipedia for more interesting ways to prove this identity.

For $0 \leq x \leq k$,

$$P_{\substack{x \\ (0 \leq x \leq k)}}(x) = \binom{b+r}{k}^{-1} \binom{b+r}{k} \quad \left[\text{by (e1), (e2)} \right]$$

$$\Rightarrow P_{\substack{(x) \\ (0 \leq x \leq k)}} = 1 \dots (e3)$$

For $x < 0$,

$$P_{(x < 0)}(x) = 0. \quad \left[\begin{array}{l} \because \binom{b}{x} \text{ in (e1)} \\ \text{will always be zero.} \end{array} \right] \dots (e4)$$

For $x > k$,

$$P_{(x > k)}(x) = 0. \quad \left[\begin{array}{l} \because \binom{r}{k-x} \text{ in (e1)} \\ \text{will be } \binom{r}{\text{some -ve value}} \\ = 0. \end{array} \right] \dots (e5)$$

$$\begin{aligned} \therefore \sum_{x \in R_X} P_X(x) &= \sum_{x < 0} P_X(x) + \sum_{0 \leq x \leq k} P_X(x) + \sum_{x > k} P_X(x) \\ &= 0 + 1 + 0 \\ &\quad \text{I By (e3), (e4), (e5)} \\ &= 1. \text{ (Proved.)} \end{aligned}$$

6.
Sol.

The Matching Problem

We previously found that

Please check "The Matching Problem" discussed in the HOS Book on page no 100 [sec2.1.5: Prob 7]

$$P(X_N=0) = \frac{1}{2!} - \frac{1}{3!} + \frac{1}{4!} - \cdots (-1)^N \frac{1}{N!} \text{ for } N=1, 2, \dots$$

Using this we will find $P(X_N=k)$ for all $k \in \{0, 1, \dots, N\}$.
Let us first calculate $P(X_N=1)$.

This is the probability that exactly one person receives his/her hat. We can fix this person. Therefore, the rest of the people ($N-1$ people) won't receive their own hats.

So, there are N ways to choose the person who gets his/her own hat.

The probability that the chosen person gets his/her hat is equal to $\frac{1}{N}$.

The probability that none of the other $N-1$ people receive their own hats is $P(X_{N-1}=0)$.

So, we have

$$P(X_N=1) = N \cdot \frac{1}{N} \cdot P(X_{N-1}=0)$$

$$= P(X_{N-1}=0) = a_{N-1} \text{ (say)} \quad \text{--- (i)}$$

Similarly, for calculating $P(X_N=2)$, there are $\binom{N}{2}$ ways to choose two people who get their own hats.

$\frac{1}{N} \cdot \frac{1}{N-1}$ is the probability that the two chosen people receive their own hats and $P(X_{N-2}=0)$ is the probability that none of the other $N-2$ people receive their own hats.

$$P(X_N=2) = \binom{N}{2} \cdot \frac{1}{N} \cdot \frac{1}{N-1} \cdot P(X_{N-2}=0)$$

$$= \frac{1}{2} P(X_{N-2}=0) = \frac{1}{2} a_{N-2} \quad \text{--- (ii)}$$

In general we have,

$$\begin{aligned} P(X_N = k) &= \binom{N}{k} \cdot \frac{1}{N} \cdot \frac{1}{N-1} \cdots \frac{1}{N-k+1} \cdot P(X_{N-k} = 0) \\ &= \frac{1}{k!} P(X_{N-k} = 0) = \frac{1}{k!} p^{N-k} \text{ for } k=0, 1, 2, \dots, N \end{aligned}$$

Since $X \sim \text{Geometric}(p)$, we have:

$$P_X(k) = (1-p)^{k-1} p \text{ for } k=1, 2, \dots$$

$$\begin{aligned} \text{Thus, } P(X > m) &= \sum_{k=m+1}^{\infty} (1-p)^{k-1} p \\ &= \sum_{k=0}^{\infty} (1-p)^{k+m} p \\ &= p(1-p)^m \sum_{k=0}^{\infty} (1-p)^k \\ &= p(1-p)^m \frac{1}{1-(1-p)} \\ &= (1-p)^m \end{aligned}$$

$$\text{Similarly, } P(X > m+l) = (1-p)^{m+l}$$

$$\begin{aligned} \text{Therefore, } P(X > m+l \mid X > m) &= \frac{P(X > m+l) \text{ and } P(X > m)}{P(X > m)} \\ &= \frac{P(X > m+l)}{P(X > m)} \\ &= \frac{(1-p)^{m+l}}{(1-p)^m} \\ &= (1-p)^l \\ &= P(X > l) \end{aligned}$$

Proved

- Sol. 11
- (a). Time interval = 4×60 mins = 240 mins
 $\therefore \lambda = 240 \times \frac{1}{30} = 8$ (for a weekend)
- Thus, $X \sim \text{Poisson } (\lambda = 8)$
 since no emails received in the interval
- $$\therefore P(X=0) = \frac{e^{-\lambda} \lambda^0}{0!} = \frac{e^{-8} 8^0}{0!}$$
- $$= e^{-8}$$
- (b). Let D be the event that a weekday is chosen and let E be the event that a saturday or sunday is chosen.
- Then: $P(D) = \frac{5}{7}$ and $P(E) = \frac{2}{7}$
- Let, A be the event that I receive no emails during the chosen interval then:
- $$P(A|D) = e^{-\frac{1}{6} \cdot 60} = e^{-10}$$
- $$P(A|E) = e^{-\frac{1}{30} \cdot 60} = e^{-2}$$
- Therefore, $P(D|A) = \frac{P(A|D) \cdot P(D)}{P(A)} = \frac{e^{-10} \cdot \frac{5}{7}}{P(A|D)P(D) + P(A|E)P(E)}$
- $$= \frac{\frac{5}{7} \cdot e^{-10}}{\frac{5}{7} \cdot e^{-10} + \frac{2}{7} e^{-2}}$$

Jms

Sol. 21
(a)

Coupon collector's problem :-

Let X denote the (random) number of coupons that we need to purchase in order to complete our collection. We can write $X = X_0 + X_1 + X_2 + \dots + X_{N-1}$, where for any $i = 0, 1, 2, \dots, N-1$, X_i denotes the additional number of coupons that we need to purchase to pass from ~~i~~ to $i+1$ different types of coupons in our collection.

Thus we have $X_0 = 1$ since the first coupon is always a new one.

Also, when i distinct types of coupons have been collected, a new coupon purchased will be of a distinct type with probability equal to $\frac{N-i}{N}$, i.e., X_i will be a geometric random variable with success probability of $\frac{N-i}{N}$.

Thus we can write $X = X_0 + X_1 + \dots + X_{N-1}$ where

$$X_i \sim \text{Geometric} \left(\frac{N-i}{N} \right)$$

(b). By linearity of expectation, we have,

$$\begin{aligned} E[X] &= EX_0 + EX_1 + \dots + EX_{N-1} \\ &= 1 + \frac{N}{N-1} + \frac{N}{N-2} + \dots + \frac{N}{1} \\ &= N \left(1 + \frac{1}{2} + \dots + \frac{1}{N-1} + \frac{1}{N} \right) \end{aligned}$$

Ans.

22) (a)

$$R_X = \{1, 2, 4, 8, \dots\}.$$

$$P_X(1) = \frac{1}{2}$$

$$P_X(2) = \frac{1}{2} \cdot \frac{1}{2} = \frac{1}{4} = \frac{1}{2^2}$$

$$P_X(4) = \frac{1}{2} \cdot \frac{1}{2} \cdot \frac{1}{2} = \frac{1}{8} = \frac{1}{2^3}$$

$$P_X(2^{k-1}) = \frac{1}{2^k} \quad \text{for } k \in \{1, 2, 3, \dots\} = \mathbb{N}. \dots (e1)$$

$$\therefore EX = \sum_{x \in R_X} x \cdot P_X(x)$$

$$= 1 \cdot \frac{1}{2} + \sum_{k \in \mathbb{N}} \left(2^{k-1} \cdot \frac{1}{2^k} \right) \quad \begin{array}{l} \text{by} \\ (e1) \end{array}$$

$$= \sum_{k \in \mathbb{N}} \frac{1}{2}$$

$$= \frac{\infty}{2} \quad \left[\because \mathbb{N} \text{ is an infinite set.} \right]$$

$$= \infty \quad (\underline{\text{Ans.}})$$

22. (b)

$$P(X > 65) = P_X(128) + P_X(256) + \dots$$

$$= P_X(2^{8-1}) + P_X(2^{9-1}) + \dots$$

$$= \frac{1}{2^8} + \frac{1}{2^9} + \dots \quad] \text{ by (e1)}$$

$$= \frac{1}{2^8} \left(1 + \frac{1}{2} + \frac{1}{2^2} + \dots \right)$$

$$= \frac{1}{2^8} \cdot \frac{1}{\left(1 - \frac{1}{2}\right)} = \frac{1}{2^8} \cdot 2 = \frac{1}{2^7}$$

$$= \frac{1}{128} \quad (\underline{\text{Ans.}})$$

22. (c)

$$P_Y(2^{k-1}) = \begin{cases} \frac{1}{2^k}, & \text{for } k=1, 2, \dots, 30 \\ \frac{1}{2^{30}}, & \text{for } k=31 \\ 0, & \text{o.w.} \end{cases} \quad] \text{ like } X \text{ in (e1)}$$

$$\therefore EY = \sum_{\substack{y \in R_Y \\ y \in R_Y}} y \cdot P_Y(y) = \left(1 \cdot \frac{1}{2} + 2 \cdot \frac{1}{2^2} + \dots + 30 \cdot \frac{1}{2^30} \right) + 2^{30} \cdot \frac{1}{2^{30}} \quad \langle \text{P.T.O.} \rangle$$

22. (c) <comtd>

$$\Rightarrow EY = \left(\frac{1}{2} + \frac{1}{2} + \dots \text{ 30 times} \right) + 1$$

$$= \frac{30}{2} + 1$$

$$= 15 + 1$$

$$= 16 . \quad (\underline{\text{Ans}})$$

25) (a)

Case $m = 1$

$$\begin{aligned} P(X \geq m) &= P(X \geq 1) = 0.4 + 0.3 + 0 \\ &= 0.6 \geq \frac{1}{2}. \end{aligned}$$

$$\begin{aligned} P(X \leq m) &= P(X \leq 1) = 0.4 + 0 \\ &= 0.4 < \frac{1}{2}. \end{aligned}$$

Reached a contradiction.

Case $m = 2$

$$\begin{aligned} P(X \geq m) &= P(X \geq 2) \\ &= 0.3 + 0.3 + 0 \\ &= 0.6 \geq \frac{1}{2}. \end{aligned}$$

$$\begin{aligned} P(X \leq m) &= P(X \leq 2) \\ &= 0.4 + 0.3 + 0 \\ &= 0.7 \geq \frac{1}{2}. \end{aligned}$$

$\therefore m = 2$ is valid.

Case $m = 3$ $P(X \geq m) = P(X \geq 3) = 0.3 + 0 = 0.3 < \frac{1}{2}$.

Reached a contradiction.

\therefore Median of $X = m = 2$. (Ans.)

25. (b)

PMF of X

$$= P_X(k) = \begin{cases} \frac{1}{6}, & \text{for } k=1, 2, \dots, 6 \\ 0, & \text{o.w.} \end{cases}$$

Case $m=1$

$$P(X \leq m) = P(X \leq 1) = \frac{1}{6} + 0 = \frac{1}{6} < \frac{1}{2}.$$

contradiction.

Case $m=2$

$$P(X \leq m) = P(X \leq 2) = \frac{1}{6} + \frac{1}{6} + 0 = \frac{1}{3} < \frac{1}{2}.$$

contradiction.

Case $m=3$

$$\begin{aligned} P(X \leq m) &= P(X \leq 3) = \frac{1}{6} + \frac{1}{6} + \frac{1}{6} + 0 \\ &= \frac{1}{2} \geq \frac{1}{2}. \end{aligned}$$

$$P(X \geq m) = P(X \geq 3) = \dots = \frac{1}{2} \geq \frac{1}{2}.$$

$\therefore m=3$ is valid.

Case $m=4$

$$P(X \leq 4) = \frac{4}{6} = \frac{2}{3} \geq \frac{1}{2}.$$

$$P(X \geq 4) = \frac{3}{6} = \frac{1}{2} \geq \frac{1}{2}.$$

$\therefore m=4$ is valid.

$\langle \text{R.T.O.} \rangle$

25. (b) (contd.)

Case m = 5

$$P(X \leq 5) = \frac{5}{6} > \frac{1}{2}.$$

$$P(X \geq 5) = \frac{1}{6} = \frac{1}{3} < \frac{1}{2}. \text{ contradiction.}$$

Case m = 6

$$P(X \leq 6) = \frac{6}{6} = 1 \geq \frac{1}{2}.$$

$$P(X \geq 6) = \frac{1}{6} < \frac{1}{2}. \text{ contradiction.}$$

\therefore Medians of X are $3, 4$. (Ans.)

25. (c) $X \sim \text{Geometric}(p)$, where $0 < p < 1$.

$$\begin{aligned} \Rightarrow \text{PMF of } X = P_X(k) &= (1-p)^{k-1} p \\ &= q^{k-1} \cdot p \quad \left[\begin{array}{l} \text{Let } q = (1-p) \\ \dots (e0) \end{array} \right] \quad \dots (e3) \end{aligned}$$

Let m is a median of X .

$$\therefore P(X \geq m) \geq \frac{1}{2} \quad \dots (e1)$$

$$\text{and } P(X \leq m) \geq \frac{1}{2}. \quad \dots (e2)$$

< P.T.O. >

25. (c) < contd >

$$P(X \leq m) = \sum_{k=1}^{\lfloor m \rfloor} q^{k-1} p$$

[By (e0).
 $\lfloor m \rfloor$ is taken in case
 to tighten the upper
 bound.]

$$= p \cdot \sum_{k=1}^{\lfloor m \rfloor} q^{k-1}$$

$$= p \cdot \frac{(1 - q^{\lfloor m \rfloor - 1 + 1})}{(1 - q)}$$

$$= p \cdot \frac{(1 - q^{\lfloor m \rfloor})}{(1 - q)}$$

$$\Rightarrow P(X \leq m) = (1 - q^{\lfloor m \rfloor}).$$

[By (e3) and
 $\because p > 0$.
 ... (e4)]

$$(1 - q^{\lfloor m \rfloor}) \geq \frac{1}{2}$$

[By (e2) and (e4)]

$$\Rightarrow q^{\lfloor m \rfloor} \leq \frac{1}{2}$$

$$\Rightarrow \lg q^{\lfloor m \rfloor} \leq \lg \frac{1}{2} = -1$$

$$\Rightarrow \lfloor m \rfloor \lg q \leq -1$$

$$\Rightarrow -\lfloor m \rfloor \lg q \geq 1$$

< P.T.O. >

25.(c) <contd.>

$$\Rightarrow \lfloor m \rfloor \lg q^{-1} \geq 1$$

$$\Rightarrow \lfloor m \rfloor \geq \frac{1}{\lg(\frac{1}{q})} \quad \dots (e5)$$

Now, $P(X \geq m) = \sum_{\substack{m \\ K=\lceil m \rceil}}^{\infty} q^{K-1} p$

By (e0).
 $\lceil m \rceil$ is used
 to tighten the
 lower bound.

$$= q^{\lceil m \rceil - 1} p + q^{\lceil m \rceil + 1 - 1} p + q^{\lceil m \rceil + 2 - 1} p + \dots$$

$$= q^{\lceil m \rceil - 1} p (1 + q + q^2 + \dots \infty)$$

$$= q^{\lceil m \rceil - 1} \cdot p \cdot \frac{1}{1-q} \quad \because q < 1.$$

$$\Rightarrow P(X \geq m) = q^{\lceil m \rceil - 1}$$

By (e3) and
 ... (e6) $\because p > 0$.

$$q^{\lceil m \rceil - 1} \geq \frac{1}{2} \quad \text{By (e1) and (e6)}$$

$$\Rightarrow (\lceil m \rceil - 1) \lg q \geq -1$$

$$\Rightarrow (\lceil m \rceil - 1) \lg \frac{1}{q} \leq 1$$

< P.T.O. >

25. (c) <contd.>

$$\Rightarrow \lceil m \rceil - 1 \leq \frac{1}{\lg \frac{1}{\alpha}}$$

$$\Rightarrow \lceil m \rceil \leq \frac{1}{\lg \frac{1}{\alpha}} + 1. \quad \dots (e7)$$

\therefore Any $m \in R_X$ that satisfies

$$\bullet \quad \lfloor m \rfloor \geq \frac{1}{\lg(\frac{1}{\alpha})} \quad] \text{ By (e5)}$$

$$\text{and } \lceil m \rceil \leq \frac{1}{\lg(\frac{1}{\alpha})} + 1 \quad] \text{ By (e7)}$$

is a median of X . (Ans.)

3>

Case 3.1 : Poisson distribution.

Ref: Wikipedia article on 'Infinite divisibility (probability)

Case 3.2 Binomial distribution Binomial(n, p)

where

n = total number of nodes,

p = probability of a node being idle.

Case 3.3 : Hypergeometric distribution.