

Module 9

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Team 4: Pause&Ponder*

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1 Linear combinations of functions: examples

1.1 In-Class Exercise 1

Download `LinCombExample.java` and verify that $g(x)$ is being computed as above. Compile and execute to draw the function g . You will also need `Function.java` and `SimplePlotPanel.java`

Yes it is correct.

```
public static void main (String [] argv)
{
    Function G = new Function ("g(x)_vs_x");
    for (double x=0; x<=6.283; x+=0.1) {
        G.add (x, g(x));
    }

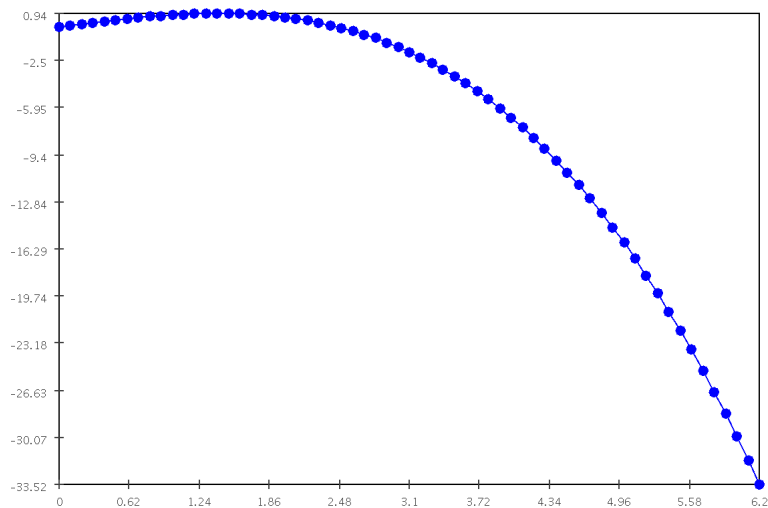
    // Draw the curve.
    G.show ();
}

static double g (double x)
{
    // Array of coefficients (scalars in the linear combination).
    double [] alpha = {0, 1.0, 0, -0.16666666666666666};

    // Here,  $f_1(x) = x$ ,  $f_2(x) = 0$ ,  $f_3(x) = x^3$ .
    double g = alpha[1] * x + alpha[3] * Math.pow(x,3) ;

    return g;
}
```

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$g(x)$ vs x

1.2 In-Class Exercise 2

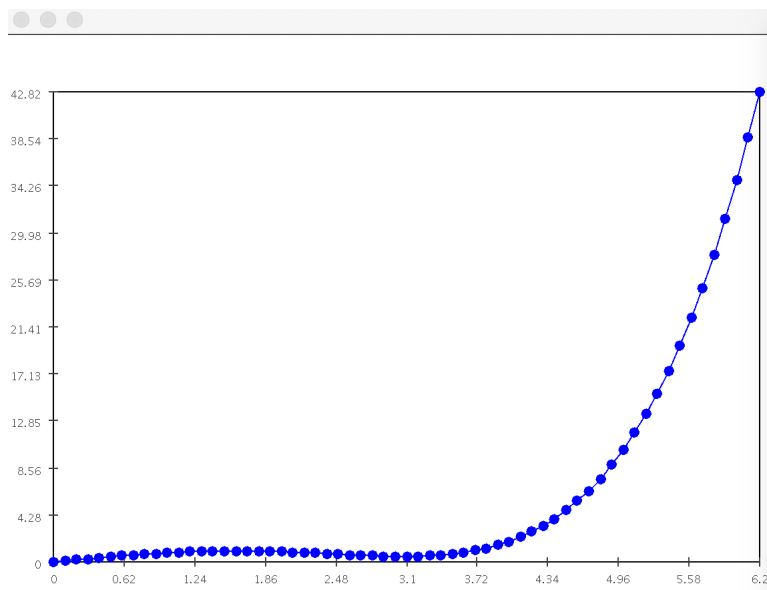
Download `LinCombExample2.java` and verify that $g(x)$ is being computed as above. Use different values of n in the program to include additional terms step by step. What is the function $g(x)$ remind you of?

```
static double g (double x)
{
    // Array of coefficients (scalars in the linear combination).
    double[] alpha = {0, 1.0, 0, -0.16666666666666666, 0, 0.008333333333333333, 0, -1.9841269841269839E-4,

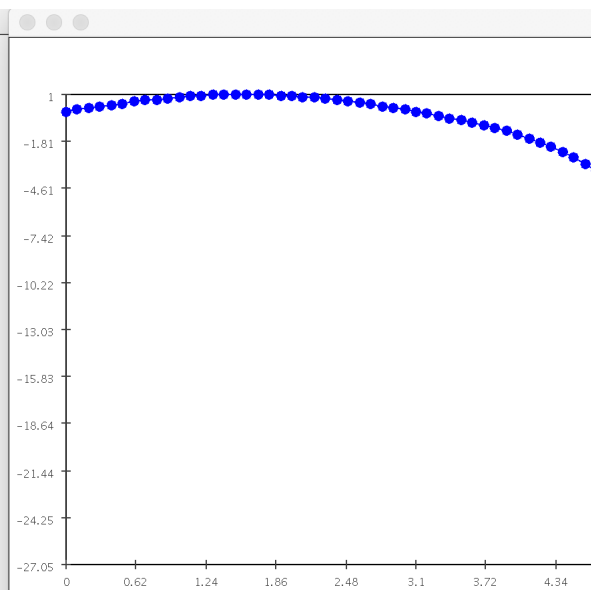
    // How many terms to include: Start with n=3, then change to n=5,
    // n=7, then n=13.
    int n = 3;

    double g = 0;
    for (int k=0; k<=n; k++) {
        g = g + alpha[k] * Math.pow (x, k);
    }

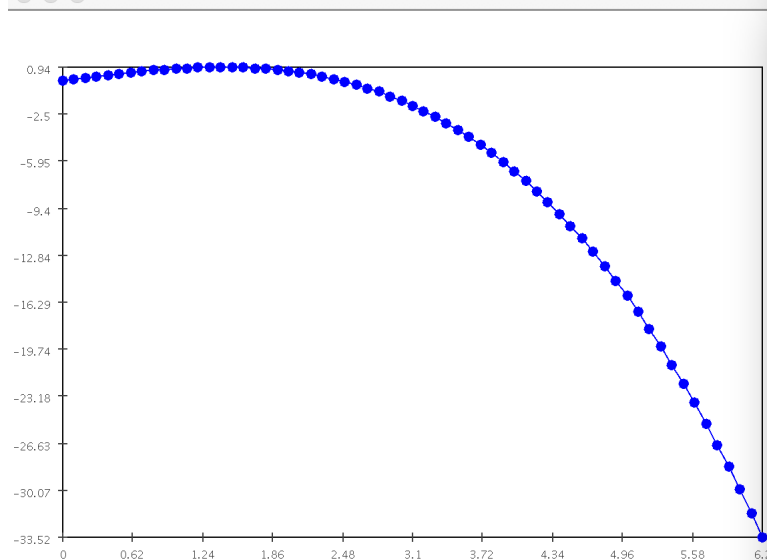
    return g;
}
```



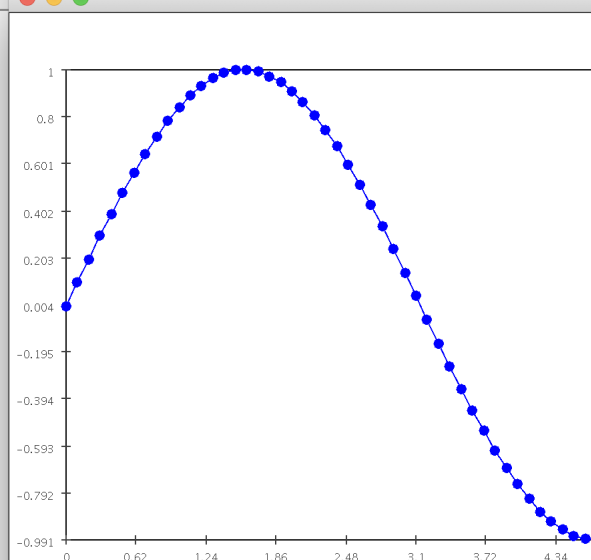
$g(x)$ vs x



$g(x)$ vs x



$g(x)$ vs x



$g(x)$ vs x

It is a sinusoidal (sin) function.

1.3 In-Class Exercise 3

Download `LinCombExample3.java` use different values of n to print the error between the approximation function $g(x)$ and the function that $g(x)$ seeks to approximate.

- 1) What is the difference in error between $n=3$ and $n=13$? What is the ratio?
- 2) How does the ratio differ when the interval is $[0, \pi]$ versus $[0, 2\pi]$? $[0, \pi/2]$?
- 3) Why is the error at each point x multiplied by Δx , and what does this say about the totalError?

1) For $n = 3$, Total error: 44.039450550769644
and $n = 13$, Total error: 0.22858564290037903

Ratio : approx(192:1)

2) For $[0, \pi]$, Total error: 19.530699449230283

For $[0, 2\pi]$, Total error: 44.039450550769644

Ratio :

and

For $[0, \pi]$, Total error: 19.530699449230283

For $[0, \frac{\pi}{2}]$, Total error: 19.530699449230283

Ratio : 1

3) Why is the error at each point x multiplied by deltaX, and what does this say about the totalError?

1.4 In-Class Exercise 4

Look up the Taylor series for $\sin(x)$. Then, download SinTaylor.java and implement the factorial function. Compare the coefficients printed to those in the alpha[] array used in prior exercises.

```
public class SinTaylor {
    public static double lookup[];
    public static void main (String [] argv)
    {
        int range = 13;
        lookup = new double[range+1];           //Initialized with 0;

        for (int k=1; k<=range; k+=2) {
            double alpha = 1.0 / factorial (k);
            System.out.println ("alpha_(without_the_sign):_" + alpha);
        }
        static double factorial (int k)
        {
            if(k <= 1) {
                return 1;
            }
            if(lookup[k] != 0.0) {
                return lookup[k];
            }
            lookup[k] = k * factorial(k-1);
            return lookup[k];
        }
    }
}
```

Values of alpha for the odd powers of sin(x) is :

alpha (without the sign): 1.0
 alpha (without the sign): 0.16666666666666666
 alpha (without the sign): 0.008333333333333333
 alpha (without the sign): 1.984126984126984E-4
 alpha (without the sign): 2.7557319223985893E-6
 alpha (without the sign): 2.505210838544172E-8
 alpha (without the sign): 1.6059043836821613E-10

Compared to the previous alphas :

alpha: 0.0
 alpha: 1.0
 alpha: 0.0
 alpha: -0.16666666666666666
 alpha: 0.0
 alpha: 0.008333333333333333
 alpha: 0.0
 alpha: -1.9841269841269839E-4
 alpha: 0.0
 alpha: 2.7557319223985884E-6
 alpha: 0.0
 alpha: -2.5052108385441714E-8
 alpha: 0.0
 alpha: 1.605904383682161E-10

Since sin(x) is a odd function the taylor series just takes the odd powers.

The taylor series for sin(x) looks like :

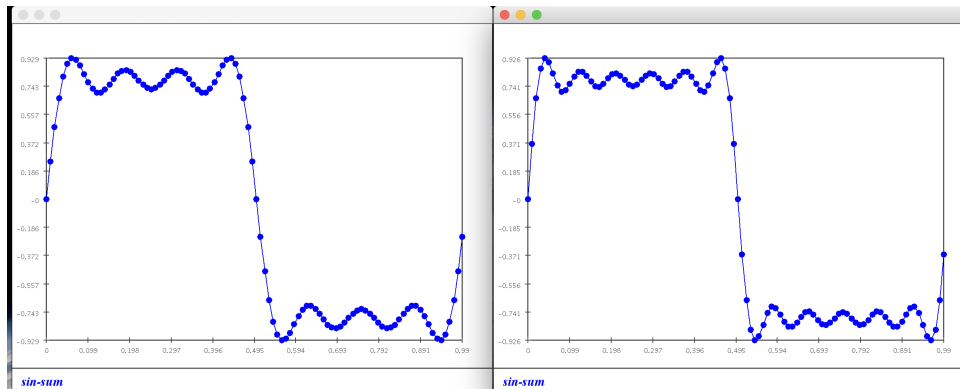
$$\begin{aligned}\sin x &= x - \frac{x^3}{3!} + \frac{x^5}{5!} - \frac{x^7}{7!} + \frac{x^9}{9!} - \dots \\ &= \sum_{n=1}^{\infty} (-1)^{(n-1)} \frac{x^{2n-1}}{(2n-1)!} \stackrel{\text{or}}{=} \sum_{n=0}^{\infty} (-1)^n \frac{x^{2n+1}}{(2n+1)!}\end{aligned}$$

1.5 In-Class Exercise 5

Download TrigPolyLinComb.java, which ends with k=3. Observe the result. Try adding additional terms until k=11.

```
public class TrigPolyLinComb {  
  
    public static void main (String[] argv)  
    {  
        int k = 1;  
        int range = 3;  
        double y = 0.0;  
        Function Fsum = new Function ("sin-sum");  
        for (double x=0; x<=1; x+=0.01) {  
            y = 0.0;  
            for(k = 1;k <= range; k += 2) {  
                y += 1.0/k * Math.sin (2*Math.PI*k*x);  
            }  
            // double y = Math.sin (2*Math.PI*x)  
            //           + 1.0/3.0 * Math.sin (2*Math.PI*3*x);  
            Fsum.add (x, y);  
        }  
        Fsum.show ();  
    }  
}
```

The graph for k = 7 and k = 11 looks like :



1.6 In-Class Exercise 6

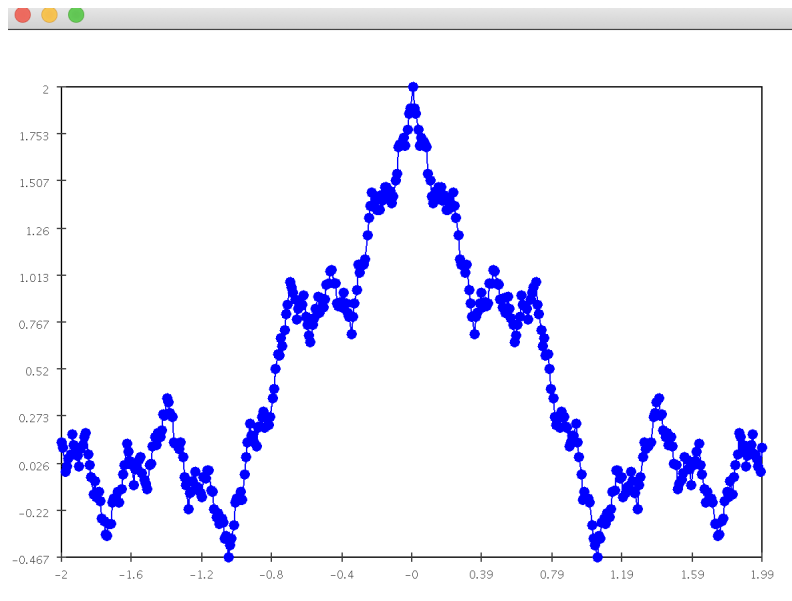
What is an example of a function $h(x)$ that no linear combination of the different $\sin(2\pi kx)$ functions could possibly approximate?

$$h(x) = \sin^2(2\pi kx) + \cos^2(2\pi kx)$$

1.7 In-Class Exercise 7

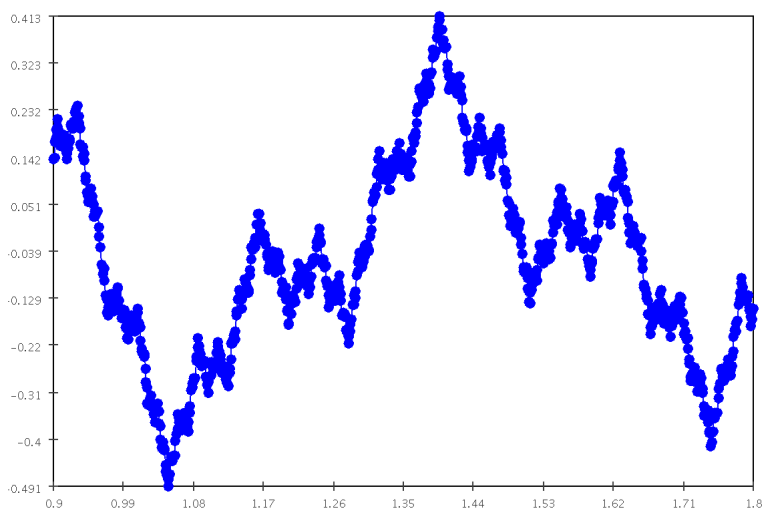
Download `Weierstrass.java` and execute. Then change the interval as directed.

`xLeft=-2, xRight=2, deltaX=0.01`



Weierstrass

`xLeft=0.9, xRight=1.8, deltaX=0.001`



Weierstrass

2 Bernstein polynomials

2.1 In-Class Exercise 8

What does $\binom{n}{k}$ evaluate to in these four cases: $k=0, k=1, k=n-1, k=n$?

$\binom{n}{k}$, when $k = 0$; $\binom{n}{0} = 1$
 $\binom{n}{k}$, when $k = 1$; $\binom{n}{1} = n$
 $\binom{n}{k}$, when $k = n-1$; $\binom{n}{n-1} = \binom{n}{1} = n$
 $\binom{n}{k}$, when $k = n$; $\binom{n}{n} = \binom{n}{0} = 1$

2.2 In-Class Exercise 9

Look up the formula for $\binom{n}{k}$ and explain how it's derived. Write code in `Combinations.java` to compute it.

A k -combination of a set S is a subset of k distinct elements of S . If the set has n elements, the number of k -combinations is equal to the binomial coefficient

$$\binom{n}{k} = \frac{n(n-1)\cdots(n-k+1)}{k(k-1)\cdots 1}$$

```
public class Combinations {

    static int lookup[];
    public static void main (String[] argv)
    {
        // Try k=0,1,2,3,4,5.
        int k = 1;

        int r = numCombinations (5,k);
        System.out.println (r);
    }

    static int numCombinations (int n, int k)
    {
        lookup = new int [n+1];
        int n1 = factorial(n);
        lookup = new int [k+1];
        int k1 = factorial(k);
        lookup = new int [(n-k)+1];
        int nk = factorial(n-k);
        return (n1/(k1 * nk));
    }

    static int factorial (int k)
    {
        if(k <= 1) {
            return 1;
        }
        if(lookup[k] != 0) {
            return lookup[k];
        }
    }
}
```



```

        lookup[k] = k * factorial(k-1);
    }
    return lookup[k];
}
}

```

2.3 In-Class Exercise 10

Prove that $\binom{n}{k} = \binom{n-1}{k} + \binom{n-1}{k-1}$. Write recursive code in `CombinationsRecursive.java` to compute it and plot the shape vs. k . Why is it symmetric? Look up Pascal's triangle and draw a few rows. Explain what the above result has to do with Pascal's triangle.

The reason $\binom{n}{k}$ is embedded in Pascal's triangle is that

$$\binom{n}{k} = \binom{n-1}{k} + \binom{n-1}{k-1}$$

and $\binom{n}{0} - \binom{n}{n} = 1$ for all n . Obviously because there is only one possible combination to choose nothing from something and everything from everything. So the numbers $\binom{n}{k}$ satisfy the defining rule of Pascal's triangle. The way to think about it : How many ways can you choose k items from a set of n , where one of the items is marked? Who cares if one is marked, the answer is $\binom{n}{k}$, but we can also say that $\binom{n-1}{k-1}$ choices which include the marked item and $\binom{n-1}{k}$ which exclude the marked item so the answer will be to add both subsets to finish the set.

A more refined way to think of it is that the values in Pascal triangle can be expressed as :

[illegible]
$$\begin{array}{ccccccccccc}
n = 0 & & & & & & \binom{0}{0} & & & & \\
n = 1 & & & & & \binom{1}{0} & & \binom{1}{1} & & & \\
n = 2 & & & & \binom{2}{0} & & \binom{2}{1} & & \binom{2}{2} & & \\
n = 3 & & & \binom{3}{0} & & \binom{3}{1} & & \binom{3}{2} & & \binom{3}{3} & \\
n = 4 & & \binom{4}{0} & & \binom{4}{1} & & \binom{4}{2} & & \binom{4}{3} & & \binom{4}{4} \\
n = 5 & & \binom{5}{0} & & \binom{5}{1} & & \binom{5}{2} & & \binom{5}{3} & & \binom{5}{4} & & \binom{5}{5} \\
n = 6 & \binom{6}{0} & & \binom{6}{1} & & \binom{6}{2} & & \binom{6}{3} & & \binom{6}{4} & & \binom{6}{5} & & \binom{6}{6}
\end{array}$$

Since we know that a number is the sum of two numbers on the top, for example searching for $\binom{5}{2}$ we have to find the two numbers above it that is $\binom{4}{1}$ and $\binom{4}{2}$.

This is musical!

The following code for CombinationsRecursive using Pascal's triangle. There are 3 versions to decrease comparisons.

```
import java.util.ArrayList;
```

```

import java.util.HashMap;

public class CombinationsRecursive {

    static int lookup[];
    static HashMap<ArrayList<Integer>,Integer> lookupTable;
    public static void main (String[] argv){
        int n = 5;
        Function C = new Function ("n-choose-k_vs_k");
        for (int k=0; k<=n; k++) {
            lookupTable = new HashMap<>();
            int r = numCombinationsRecursive (n,k);
            System.out.println("-----");
            C.add (k, r);
        }
        C.show ();
    }

    static int numCombinationsRecursive (int n, int k){
        if ((n==k) || (n==1) || (k==0)) {
            System.out.println("n:␣" + n + "␣k:␣" + k + "␣r:␣" + 1);
            return 1;
        }

        if(nkInMap(n,k)) {
            return getValue(n,k);
        }

        ArrayList<Integer> list = new ArrayList<Integer>(2);
        list.add(n);
        list.add(k);
        lookupTable.put(list, numCombinationsRecursive(n-1,k)+numCombinationsRecursive(n-1,k-1));
        System.out.println("n:␣" + n + "␣k:␣" + k + "␣r:␣" + lookupTable.get(list));
        return lookupTable.get(list);
    }

    private static int getValue(int n, int k) {
        ArrayList<Integer> list = new ArrayList<Integer>(2);
        list.add(n);
        list.add(k);
        return lookupTable.get(list);
    }

    private static boolean nkInMap(int n, int k) {
        ArrayList<Integer> list = new ArrayList<Integer>(2);
        list.add(n);
        list.add(k);
        if(lookupTable.containsKey(list)) {
            return true;
        }
        return false;
    }
}

```

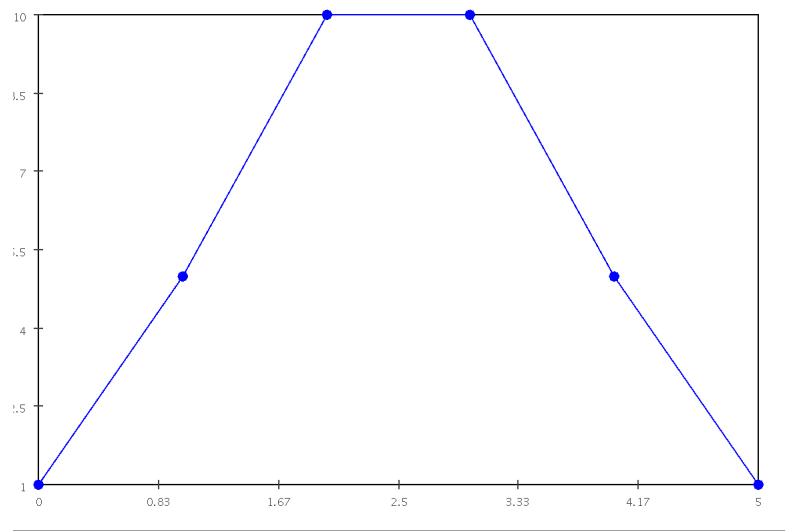
Output for $\binom{5}{2}$

```

n: 2 k: 2 r: 1
n: 1 k: 1 r: 1
n: 1 k: 0 r: 1
n: 2 k: 1 r: 2

```

n: 3 k: 2 r: 3
 n: 2 k: 0 r: 1
 n: 3 k: 1 r: 3
 n: 4 k: 2 r: 6
 n: 3 k: 0 r: 1
 n: 4 k: 1 r: 4
 n: 5 k: 2 r: 10



n-choose-k vs k

3 Proofs

3.1 In-Class Exercise 12

Prove that $\binom{n}{k} = \frac{n}{k} \binom{n-1}{k-1}$ and implement this approach as a second recursive method in `CombinationsComparison2.java`. Try $n = 10$ and $n = 20$. Why is the return type double? Implement an iterative version of the above tail recursion.

n=5	k=0	p=1	q=1	r=1.0	s=1.0	
numCalls=12		numCallsRecursive=1		numCallsRecursive2=1		numIterations=0
n=5	k=1	p=5	q=5	r=5.0	s=5.0	
numCalls=23		numCallsRecursive=10		numCallsRecursive2=3		numIterations=1
n=5	k=2	p=10	q=10	r=10.0	s=10.0	
numCalls=34		numCallsRecursive=23		numCallsRecursive2=6		numIterations=3
n=5	k=3	p=10	q=10	r=10.0	s=10.0	
numCalls=45		numCallsRecursive=36		numCallsRecursive2=10		numIterations=6
n=5	k=4	p=5	q=5	r=5.0	s=5.0	
numCalls=56		numCallsRecursive=45		numCallsRecursive2=15		numIterations=10
n=5	k=5	p=1	q=1	r=1.0	s=1.0	
numCalls=68		numCallsRecursive=46		numCallsRecursive2=16		numIterations=15
TotalnumCalls=68						
TotalnumCallsRecursive=46						
TotalnumCallsRecursive2=16						
TotalnumIterations=15						

```

static double numCombinationsIterative (int n, int k)
{
    double result = 1.0;
    for(int i= 0;i<=n;i++) {
        if((k-i) > 0) {
            result *= ((n-i)*(1.0))/((k-i)*(1.0));
            numIterations++;
        }
    }
    return result;
}

```

```

static double numCombinationsRecursive2 (int n, int k)
{
    numCallsRecursive2 ++;
    if(n==k || n==1 || k==0) {
        return 1.0;
    }

    if(nkInDoubleMap(n,k)) {
        return getDoubleValue(n,k);
    }
}

```

```

        ArrayList<Integer> list = new ArrayList<Integer>(2);
        list.add(n);
        list.add(k);
        double ans = (n * 1.0)/(k * 1.0);
        lookupNewTable.put(list, ans * numCombinationsRecursive2(n-1,k-1));
        return lookupNewTable.get(list);
    }

```

The return type is double so that the division can give the best approximate value of the division involved.

3.2 In-Class Exercise 16

For any two numbers a,b such that $a < b$ and $t \in [0, 1]$ prove that $(1 - t)a + tb \in [a, b]$. Then, use this fact and Proposition 9.1 to prove 9.2.

To prove : $b_{n,k}(t) \geq 0$ for all n,k and $t \in [0, 1]$.

$b_{n,k}(t) = (1 - t)b_{n-1,k}(t) + tb_{n-1,k-1}(t)$ we can use this recursive proposition to prove the above property.

Base case : It is easily seen that the functions $B_{0,1}(t) = 1 - t$ and $B_{1,1}(t) = t$ are both non-negative for $0 \leq t \leq 1$.

Initial Hypothesis(IH) : If we assume that all Bernstein polynomials of degree less than k are non-negative, then by using the recursive definition of the Bernstein polynomial,

we can write $B_{i,k}(t) = (1 - t)B_{i,k-1}(t) + tB_{i-1,k-1}(t)$ and argue that

$B_{i,k}(t)$ is also non-negative for $0 \leq t \leq 1$, since **all components on the right-hand side of the equation are non-negative** components for $0 \leq t \leq 1$.

By induction, all Bernstein polynomials are non-negative for $0 \leq t \leq 1$.

In this process, we have also shown that each of the Bernstein polynomials is positive when $0 < t < 1$.

3.3 In-Class Exercise 19

Prove the below assertion.

$$t^k \cdot (1 - t)^{n-k} = t^k \cdot (1 - t)^{n+1-k} + t^{k+1} \cdot (1 - t)^{n+1-(k+1)}$$

In terms of Pascals triangle all we need to prove is :

$$\mathbf{k\text{-th term at level n} = k\text{-th term at level n+1} + (k+1)\text{-st term at level n+1}$$

This can be done through an induction proof.

Base case : It is easily seen that the equation is true

$$\text{L.H.S.} \Rightarrow t^0 \cdot (1 - t)^0 = 1$$

$$\text{R.H.S.} \Rightarrow t^0 \cdot (1 - t)^{0+1-0} + t^{0+1} \cdot (1 - t)^{0+1-(0+1)}$$

$$= 1 - t + t = 1$$

for $n = 0, k = 0$ and similarly for $n = 1, k = 0, k = 1$.

Initial Hypothesis(IH) : If we assume that all further equations are similarly equal for $n=n+1$ then by using the recursive definition of the Bernstein polynomial,

we can write $t^k \cdot (1-t)^{i-k} = t^k \cdot (1-t)^{i+1-k} + t^{k+1} \cdot (1-t)^{i+1-(k+1)}$ and argue that

...

By induction, all

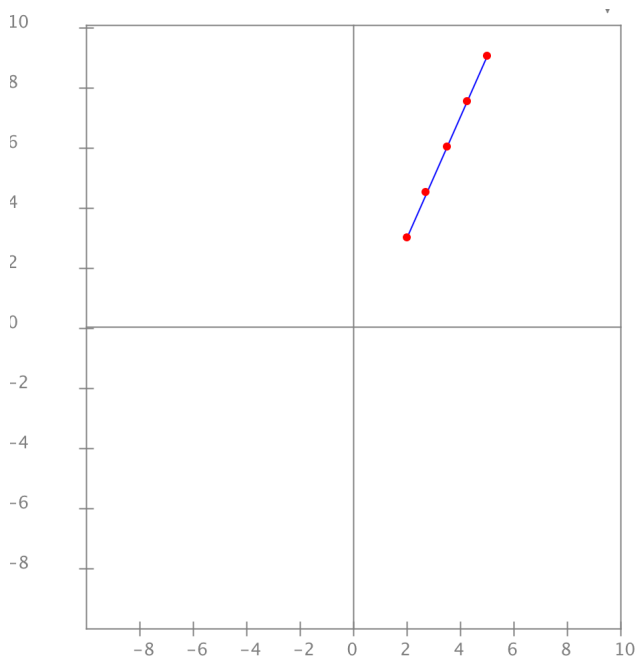
3.4 In-Class Exercise 25

Compile and execute `DrawLineParametric.java`. Notice that the computed points are always between the end points.

* Prove that if $x_0 < x_1$ then $x_0 \leq x(t) \leq x_1$ and if $x_1 < x_0$, then $x_1 \leq x(t) \leq x_0$. Obviously, the same will be true for $y(t)$.

* Prove that each point $(x(t), y(t))$ is on the line segment between the end points.

* What happens when $t < 0$ or $t > 1$? Experiment with your program, and then explain.



1)

The parametric equation says that if there are two points $(x_0, y_0), (x_1, y_1)$:

$$x = x_1 + (\text{change in } x) * t$$

$$y = y_1 + (\text{change in } y) * t$$

where t ranges from $[0 \text{ to } 1]$ where change in y is $y_1 - y_0$ and change in x is $x_1 - x_0$

When the points are being dealt in R^3 i.e. 3 dimensions on the real vector space, the only way to get an equation of line is the parametric equation.

Proof :

$$(y - y_1) = \frac{y_2 - y_1}{x_2 - x_1}(x - x_1)$$

can be written as :

$$\frac{y - y_1}{y_2 - y_1} = \frac{x - x_1}{x_2 - x_1}$$

Since they are equal

$$\begin{aligned} x - x_1 &= (x_2 - x_1)t \\ \Rightarrow x &= x_1 + (x_2 - x_1)t \end{aligned}$$

Similarly for $y(t)$.

2) We can prove this by the following example :

Parametrize the line segment that connects the points (2, 3) and (7, 9).

It's asking for a line segment. Taking the parameterization:

$$\begin{aligned} x(t) &= 2 + 5t \\ y(t) &= 3 + 6t \end{aligned}$$

We know that when $t = 0$ we're at the point (2, 3) and when $t = 1$ we're at the point (7, 9). If we restrict the parameter so that $0 \leq t \leq 1$ then we find only the line segment that lies between those two points.

3) When the value of t goes less than 0 or more than 1 it extends the point in the direction of the line. Shown below.

