

A.

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ALEXANDRU

$f([], -1).$
 $f([H|T], s) :-$
 $f(T, s_1),$
 $s_1 > 0, !,$
 $s \text{ is } s_1 + H.$

$f([-1|T], s) :-$
 $f(T, s_1),$
 $s \text{ is } s_1.$

~~$f_{Aux}(T, H)$~~

$f_{Aux}([H|_], s, s_1) :-$
 $H > 0,$
 $s_1 < H, !,$
 $s \text{ is } H.$

$f_{Aux}(_, s, s_1) :-$
 $s \text{ is } s_1.$

now f will become.

$f([], -1).$
 $f([H|T], s) :-$
 $f(T, s_1),$
 $f_{Aux}([H|T], s, s_1).$

we have an auxiliary function replacing the recursive call, therefore the 2 cases are treated inside the newly added function. Moreover, in the old one we'll just have one branch from the previous 2.

2B

insert(E, L, [E|L]).

insert(E, [H|T], [H|R]) :-
insert(E, T, R).

Math. model.

$$\text{insert}(E, L_1 \dots L_m) = E \cup L_1 L_2 \dots L_m$$

$$= L_1 \cup \text{insert}(E, L_2 \dots L_m)$$

- insert(E - element, L - list, R - result list)
insert(i, i, o)

Math. model

$$\text{an}(l_1 l_2 \dots l_m, k) = l_1 \text{ if } k=1$$

$$\text{an}(l_2 \dots l_m, k), \text{ if } k \geq 1$$

$$\text{insert}(l_1, \text{an}(l_2 \dots l_m, k-1)) \text{ if } k > 1$$

an([E|_], 1, [E]).

an([_|T], k, R) :-
an(T, k, R).

an([H|T], k, R) :-
k > 1,
k1 is k-1,
an(T, k1, R),
insert(H, R, R).

- an(L - list, k - m. of elements, R - result list)
an(i, i, o)

Math. model:

$$\text{mylength}(l_1 \dots l_m) = 0, \text{ if } m=0$$

$$1 + \text{mylength}(l_2 \dots l_m), \text{ otherwise}$$

mylength([], 0).

mylength([_|T], LENGTH) :-
mylength(T, LENGTH2),
LENGTH is LENGTH2 + 1.

- mylength(L - list, N - number)
mylength(i, o)

Math model:

$\text{insertList}(l_1, l_m, \text{list}) = \text{list}$ if $m=0$
 $l_1 \cup \text{list}(l_2 \dots l_m, \text{list})$, otherwise

$\text{insertList}([], L, L)$.

$\text{insertList}([H|T], L, [H|R]) :-$

$\text{insertList}(T, L, R)$.

- $\text{insertList}(L: \text{list}, L: \text{list}, R: \text{result list})$
 $\text{insertList}(i, i, \circ)$

Math model:

$\text{strictlyAnc}(l_1, l_m) = \text{True}$, $m=0 \vee m=1$
 False , $l_1 > l_2$
 $\text{strictlyAnc}(l_2 \dots l_m)$, $l_2 > l_1$

$\text{strictlyAnc}([])$.

$\text{strictlyAnc}[_]$.

$\text{strictlyAnc}([H_1, H_2 | T]) :-$

$H_1 < H_2$,

$\text{strictlyAnc}([H_2 | T])$.

- $\text{strictlyAnc}(L: \text{list})$
- $\text{strictlyAnc}(i)$

$\text{oneSol}(L, k, \text{ARR}) :-$

$\text{an}(L, k, \text{ARR})$,

$\text{strictlyAnc}(\text{ARR})$.

- $\text{oneSol}(L: \text{list}, k: \text{number}, \text{ARR} - \text{result list})$
 $\text{oneSol}(i, i, \circ)$

$\text{allSol}[_ , 1, _]$.

$\text{allSol}(L, k, \text{RR}) :-$

$\text{findall}(\text{RPartial}, \text{oneSol}(L, k, \text{RPartial}), R_1)$,

$k_1 \text{ in } k-1$,

$\text{allSol}(L, k_1, k_2)$, $\text{insertList}(R_2, R_1, \text{RR})$.

wrapper(L, R): -
 myLength(L, LENGTH),
 allSol(L, LENGTH, R).

- call wrapper([1, 3, 6, 4], R)
- we use arrangements as in the given ex we can have [4, 8] which wouldn't have happened if we chose combinations
- we make allSol recursive so that we can call it for every k from the length of the list to 2 (1 is the stop case)
- for each arrangement found we check if it is strictly increasing

Math model:

$$\text{replaceNodeonOdd}(l_1 \dots l_m, \text{LEVEL}, E) = \begin{cases} E \cup \text{replaceNodeonOdd}(l_2 \dots l_m, \text{LEVEL}, E), & \text{if } l_1 = E \text{ and } l_1 \text{ atom and } \text{LEVEL mod } 2 = 1 \\ l_1 \cup \text{replaceNodeonOdd}(l_2 \dots l_m, \text{LEVEL}, E), & \text{if } l_1 \text{ atom} \end{cases}$$

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(defun replaceNodeonOdd (L LEVEL E)
  (cond
    ((and (atom L) (eq (mod LEVEL 2) 1)) (list E))
    ((atom L) (list L))
    (t (list (mapcon #'(lambda (s) (replaceNodeonOdd s
      (+ 1 LEVEL) X)) L)))))
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(replaceNodeonOdd '(a (b (g)) (c(d(e)) (f))) -1 'h))

we call the function with -1 because when it will first enter, L will be a list and therefore go on the last ~~term~~ case, therefore increasing the level such that when mapcon is applied and "a" comes to the function, the level will be 0 (as required).

if we find ^{on} the atom that we want on the odd level, we return the replace value, if we find an atom that is not on the level that we want (=not odd) and otherwise we apply mapcom for the entire list which will apply our function (as a lambda function) for every element in the list and creates ~~an~~ a list with the results returned from each node.

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