Intro to Proofs

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# Contents

### Real Numbers

**Definition 0.1** (Properties of real numbers). Properties of  $\mathbb R$  are

- (d).  $\exists$  an order on  $\mathbb{R}$  which means  $\forall x, y \in \mathbb{R}, x < y$  or x > y, or x = y Ordering follows the following properties,
  - (1).  $x < y, y < z \Rightarrow x < z$  (transitivity)
  - (2).  $x < y \Rightarrow x + z < y + z, \forall z \in \mathbb{R}$
  - (3).  $x < y, z > 0 \Rightarrow xz < yz$

**Theorem 0.2.**  $xy = 0 \Leftrightarrow x = 0 \text{ or } y = 0$ 

**Proof.**  $\Leftarrow$  Without loss of generality take, x = 0 Then we get,

0y.

We can write this as,

$$(0+0)y = 0y + 0y.$$

So,

$$0y = 0y + 0y.$$

Or, m

 $\Rightarrow$ 

Assume the contrary that,  $x \neq 0$  and  $y \neq 0$  We have, xy = 0. Without loss of generality we take the multiplicative inverse of x so,

$$\frac{xy}{x} = \frac{0}{x}.$$

We showed that 0(k) = 0 so y = 0

Which contradicts our assumption, hence our assumptoin must be wrong and x=0 or y=0

**Theorem 0.3.** (-)x = -x

**Proof.** We start with (-1)x and add x to both sides so,

$$(-1)x + x = x(1-1) = 0x = 0.$$

So we showed that (-1)x is the additive identity of x. We know that the additive identity is unique for any x. Therefore, (-1)x = -x

#### Theorem 0.4. $\forall x < y, z < 0$

xz > yz.

**Proof.** If z < 0 then that means z = -k for some k > 0. We can write x < y as x - y < 0Now if we multiply both sides be k we get,

$$k(x - y) < 0$$

Now if k(x-y)=z' we can say that  $z'<0 \Rightarrow -z'>0$ Or that

$$(-1)k(x - y) > 0$$
$$z(x - y) > 0$$
$$xz > yz$$

**Theorem 0.5.**  $\forall x \in \mathbb{R} \text{ if } x \neq 0 \text{ then } x^2 > 0$ 

Theorem 0.6.  $x^2 = -(-x^2)$ 

Case 1, x > 0:

$$x \times x > x$$

$$x\times x>0x$$

$$x^2 > 0$$

Case 2, x < 0:

Then the additive inverse (-x) > 0

$$(-x)(-x) > (-x)0$$

$$(-)(-1)x^2 > 0$$

-(-1) = 1 as 1 is the additive inverse of -1

$$x^2 > 0$$

**Example.**  $\forall a, b > 0$ 

$$\frac{a+b}{2} \ge \sqrt{ab}$$

 $\Diamond$ 

Proof.

$$0 \le (\sqrt{a} - \sqrt{b})^2 = a - 2\sqrt{ab} + b.$$
$$2\sqrt{ab} \le a + b$$
$$\sqrt{ab} \le \frac{a+b}{2}$$

Example.  $x^2 - x + 1$ 

 $\Diamond$ 

**Theorem 0.7.**  $\forall x, y \in \mathbb{R}$  we have,

$$|x| \ge x$$
 and  $|x + y| \le |x| + |y|$ .

**Proof.** We use proof by cases.

#### Proof related to Sets

Theorem 0.8.

$$A \cup B \backslash (A \cap B) = (A \backslash B) \cup (B \backslash A).$$

**Proof.** We need to show that,

$$A \cup B \setminus (A \cap B) \subseteq (A \setminus B) \cup (B \setminus A)$$
.

and,

$$(A \backslash B) \cup (B \backslash A) \subseteq A \cup B \backslash (A \cap B).$$

Theorem 0.9.  $A \subseteq B \Leftrightarrow A \cup B = B$ 

**Proof.**  $\Rightarrow$  Take  $\forall x \in A \cup B$ , so either

Case 1,  $x \in A$ :

We know that by deifinition if,  $A \subseteq B$  then for  $x \in A, x \in B$  so  $x \in B$  Case 2,  $x \in B$ : If  $x \in B$  then we don't need to go further.

So we get  $\forall x \in A \cup B, x \in B$ 

 $\Leftarrow$ 

 $\forall x \in A \Rightarrow x \in A \cup B = B$ 

So,  $x \in B$  which means that,  $A \subseteq B$ 

### Disproofs

If we need to show existence,  $\exists x.P(x)$ . We can show using,

- 1. Direct constructions
- 2. Indirectly (contradiction). For instance we can show that,  $\forall x, P(x)$  is false

**Example.**  $\exists a, b, c \in R - Q \text{ s.t. } a^{bc} \in Q$ 

**Example.** Pigeonhole principle

Suppose there are m balls in n boxes,  $m > n \ge 1$  then,  $\exists$  a box where there are at least,  $\frac{m}{n} + 1$  balls

**Proof.** Assume pigeonhole is false.

Then, there are at most  $\frac{m}{n}$  balls in each box. In case 1 where  $\frac{m}{n} \notin N \Rightarrow$  total balls  $\leq n[\frac{m}{n}] = \frac{nm}{n} = m$  which is a contradiction.

In case 2 where  $\frac{m}{n} \in N$  there are at most  $\frac{m}{n} - 1$  balls in each box. So total number of balls are  $\frac{nm}{n} - n = m - n$  which is contradictory.

To disprove  $\forall x P(x)$  we can show that,  $\exists P(x)$ 

**Example.** 100 can't be written as the sum of two even integers and an odd integer.

**Proof.** Suppose it's false  $\Rightarrow \exists a, b, c \in Z \text{ s.t. } 2|a, 2|b, 2 \not/c \text{ and } 100 = a+b+c$ But,  $2|a, 2|b \Rightarrow 2|a+b$  but  $2 / c \Rightarrow 2 / (a+b) + c = 100$ So we get, 2 / 100 which is a contradiction.

Which means that the original statement is true.

**Example.** ∄ the smallest positive real number

The smallest positive real number is defined as  $x \in R$  s.t. x > 0 and  $\forall y > 0, x \le 0$ y

**Proof.** Let's assume it is true which mean that  $\exists x \in R \text{ s.t. } x > 0$  and  $\forall y > 0, x \leq y$ 

We know that  $x > 0 \Rightarrow \frac{x}{2} > 0$ So if we set  $y = \frac{x}{2}$  then we get

$$x \leq \frac{x}{2}$$
.

Which is a contradiction.

Hence it cannot be the case that there exists the smalest positive number.

**Example.**  $\not\exists f(x)$ : a polynomial with integer coefficients s.t.  $\forall n, f(n)$  is prime  $\diamond$ 

**Proof.** Consider the general form of a polynomial,

$$f(x) = a_1 x^n + \dots + a_n$$

Case 1:  $a_n = 0$ 

If  $a_n = 0$  then for any x > 1 we can take x common and get

$$f(x) = x(a_1x^{n-1} + \dots + a_{n-1})$$

So we get a factor  $x \neq 1$ 

Case 2:  $a_n = 1$ 

In this case we can just plug x=0 and we get f(x) is neither prime or composite

Case 3:  $a_n > 1$  ????

**Example.** Let  $f(x) = x^3 + 2x - 5$  then  $\exists$  unique  $x_0 \in [1,2]$  s.t.  $f(x_0) = 0$ 

**Proof.** Using intermediate value theorem.

$$f(1) = -2$$

$$f(2) = 7$$

So because -2 < 0 < 7 we know that there must exists an  $x_0 \in [1, 2]$  s.t. this is the case.

To show unique we need to show its strictly increasing. Or in other words, we need to show for every  $x_1, x_2 \in [1, 2], x_1 \leq x_2 \Rightarrow f(x_1) \leq f(x_2)$ So we need to show that,

$$x_1^3 + 2x_1 - 5 \le x_2^3 + 2x_2 - 5$$

$$x_1^3 + 2x_1 \le x_2^3 + 2x_2$$

$$(x_1^3 - x_2^3) + 2(x_1 - x_2) \le 0$$

It is enough to show that both  $x_1^3 - x_2^3$  and  $x_1 - x_2$  are smaller than or equal to 0.

$$x_1^3 - x_2^3 = (x_1 - x_2)(x_1^2 + x_1x_2 + x_2^2) = k(x_1 - x_2) \le 0$$

Similarly,

$$2(x_1 - x_2) \le 0$$
 a  $x_1 - x_2 \le 0$ 

So we have,  $x_1^3 - x^3 + 2(x_1 - x_2) \le 0$ 

Which tells us that our function is strictly increasing which implies that we only have a unique  $x_0 \in [1, 2]$ 

## **Mathematical Induction**

**Theorem 8.10** (Properties of Natural Numbers). (a).  $1 \in N$ 

- (b).  $\forall k \in \mathbb{N}, \exists k+1 \in \mathbb{N}$
- (c).  $\forall k \in N \{1\}, \exists | n \in N, \text{ s.t. } k = n + 1 \in N$
- (d). Needs to be well-ordered.

An ordered set S is well-ordered if,

$$\forall A \in S \text{ s.t. } A \neq \phi, \exists x = \min A$$

Or, 
$$\exists x \in A \text{ s.t. } y \in A, x \leq y$$

**Example.**  $\mathbb{Q}$  is not well-ordered as it does not have a minimum

**Example.**  $\mathbb{Z}$  is not well-ordered as it does not have a minimum

**Axiom 8.11.** N is well-ordered

**Theorem 8.12.** If  $A \subseteq S$  and S is a well-ordered set then A is well-ordered.

**Proof.** Let  $B \subseteq A$  and  $B \neq \phi \Rightarrow B \subseteq S$ 

So B has a min x which means that A is well-ordered by definition.  $\Box$ 

**Example.**  $[1,\infty)$  is not well-ordered because a subset  $(1,\infty)$  does not have a min

**Theorem 8.13.**  $\forall a \in \mathbb{Z}, d \in \mathbb{N}, \exists q, r \in \mathbb{Z} \times \{0, 1, \dots, d-1\} \text{ s.t.}$ 

$$a = dq + r$$

**Proof.** Let  $S = \{a - nd : n \in \mathbb{Z}, a - nd \in \mathbb{N}\}$ 

First we can see that S is non-empty as we can take

$$n = -|a| - 1 \Rightarrow a - nd > 0$$

Now because this is a subset of  $\mathbb N$  it follows the well-ordering principle implying that  $\min S=a-nd=m$ 

 $m \in S \Rightarrow \exists l \in \mathbb{Z} \text{ s.t. } m = a - ld$ 

Case 1: If m > d then

$$a - (l+1)d > 0$$

$$a - (l+1)d \in S \Rightarrow a - (l+1)d < m$$

Which is a contradiction. This means that  $m \ge d$ 

Case 2: m = d

Let q = l + 1, r = 0

$$m = d \Rightarrow a - ld = d \Rightarrow a - (l+1)d = 0$$

Case 3: 0 < m < d

Let  $q = l, r = m \Rightarrow a = dq + r$ 

Now to show uniqueness,

Suppose,  $(q, r), (q', r') \in \mathbb{Z} \times \{0, 1, ..., d - 1\}$  and

$$a = qd + r = q'd + r'$$

We have,

$$(q - q')d = r' - r$$

$$0 - (d - 1) \le r' - r \le d - 1$$

And,

$$d|r' - r \Rightarrow r' - r = 0$$

**Definition 8.14.** Let  $a,b\in\mathbb{N}, d=GCD(a,b)\in N$  if

- (a). d|a and d|b and
- (b). If  $d' \in N$  s.t. d'|a and d'|b then  $d \geq d'$

**Theorem 8.15.**  $\forall a, b \in \mathbb{N}, \exists p, q \in \mathbb{Z} \text{ s.t. } GCD(a, b) = ap + bq$ 

**Proof.** Let  $S = \{a_m + b_n : m, n \in \mathbb{Z}, a_m + b_n \subseteq \mathbb{N}\}$ 

We know S is non-empty as m, n = 1 makes it a + b > 0 as  $a, b \in \mathbb{N}$ So by well-ordering principle we know that  $\exists \min S = d$  and  $p, q \in \mathbb{Z}$  s.t.

$$d = ap + bq$$

If  $d' \in \mathbb{N}$  s.t. d'|a and  $d'|b \Rightarrow d'|ap + bq = d$ So,  $d \in \mathbb{N} \Rightarrow d \geq d'$ 

$$d \in \mathbb{N} \Rightarrow \exists m \in \mathbb{Z}, r \in \{0, \dots, d-1\} \text{ s.t. } a = md + r$$

Which means r = a - md = a - m(ap + bq) = a(1 - mp) + b(-mq)r < d but  $d = \min S \Rightarrow r \notin S \Rightarrow r = 0$ 

So a = md so d|a. Similarly, d|b

This means d is the greatest common divisor.

**Theorem 8.16** (Induction principle). Suppose  $k \in N, S \subseteq N$  satisfy,

- (a).  $k \in S$
- (b). if  $n \in S$  then  $n + 1 \in S$

then  $\{k, k+1, \dots\} \subseteq$ 

Proof. Let  $A = \{n \in \mathbb{N}, n \ge k : n \not \in S\}$ 

Suppose  $A \neq \phi \Rightarrow n_0 = \min A$  exists

Which means  $n_0 \ge k$  but  $k \not\in A$  due to (a). So,  $n_0 > k \Rightarrow n_0 - 1 \ge k$  and  $n_0 - 1 \not\in A$  as  $n_0 = \min A$ 

(b). and  $n_0 - 1 \notin A \Rightarrow n_0 \notin A$  contradictiont which implies that  $A = \phi$ 

**Corollary 8.17.** If a statement  $P(n), n \in \mathbb{N}$  satisfies

- (a) P(k) is true
- (b)  $P(n) \Rightarrow P(n+1)$

Then P(n) is true for all  $n \geq k$ 

Theorem 8.18. 
$$\sum_{k=1}^{n} k = \frac{n+1}{2}n$$

**Proof.** Proof by induction.

If n = 1, then (1) holds.

If n = k then we have,

$$1 + \dots + k = k \frac{k+1}{2}$$

Now we need to show that  $1 + \cdots + k + k + 1 = (k+1)\frac{k+2}{2}$ Using the statement for n = k we can do,

$$k\frac{k+1}{2} + k + 1 = (k+1)\frac{k+2}{2}$$

Simplyfying the left hand side we get,

$$(k+2)\frac{k+1}{2} = (k+1)\frac{k+2}{2}$$

It is trivial to see that this is true.

Hence by induction our statement is true.

**Definition 8.19.**  $C_n^m = \binom{n}{m} = \frac{n!}{m!(n-m)!}$  if  $n \ge m \ge 0$  and 0 otherwise

**Remark.**  $\Gamma(z) = \int_0^\infty r^z e^{-t} dt$  is a stronger definition of factorial

Theorem 8.20. 
$$\binom{n+1}{m} = \binom{n}{m} + \binom{n}{m+1}, \forall n,m \in \mathbb{Z}$$

**Proof.** We prove by cases.

Case 1:  $m \le -2$  or m > n

Case 2: n = m = -1

Case 3: n > m = -1

Case 4:  $m \ge 0$  and  $m \le n$ 

Theorem 8.21.  $\forall n \in N \cup \{0\}, a, b \in \mathbb{R}$ 

$$(a+b)^n = \sum_{k=0}^n \binom{n}{k} a^{n-k} b^k$$