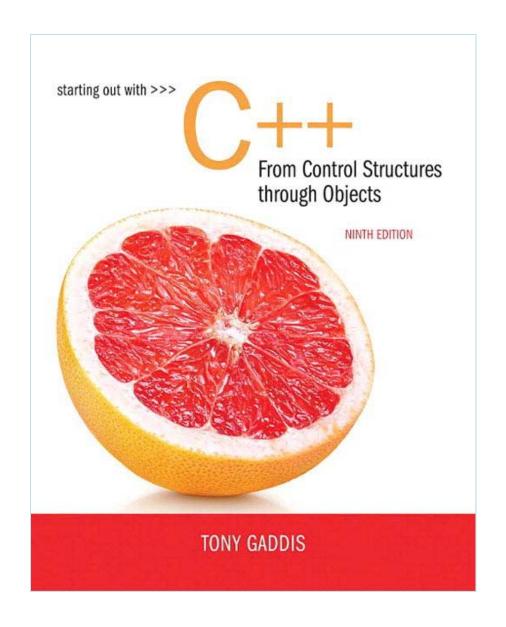
# Chapter 20: Recursion



#### Introduction to Recursion

A <u>recursive function</u> contains a call to itself:

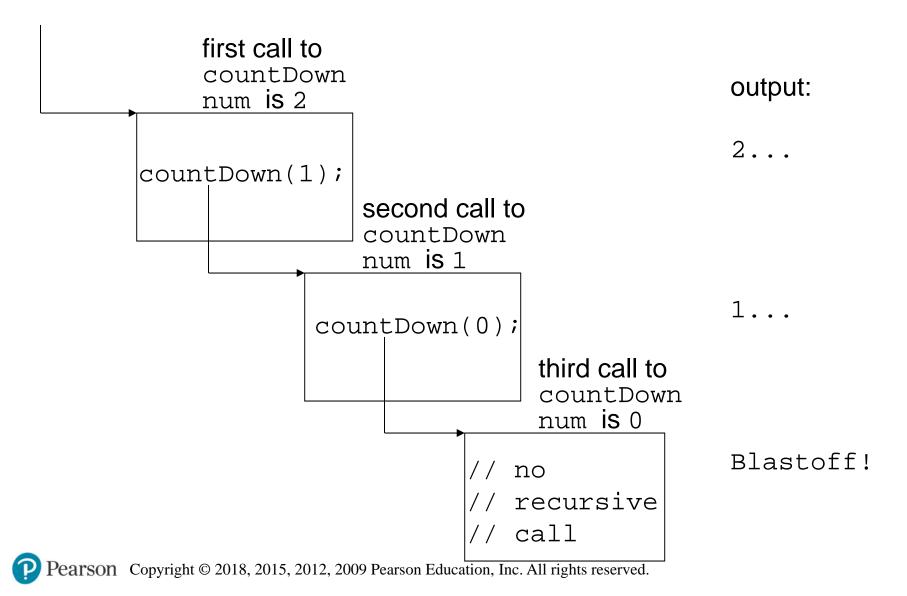
```
void countDown(int num)
{
   if (num == 0)
      cout << "Blastoff!";
   else
   {
      cout << num << "...\n";
      countDown(num-1); // recursive
   }
   }
}</pre>
```

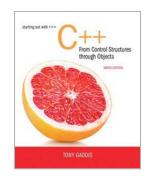
#### What Happens When Called?

If a program contains a line like countDown(2);

- 1. countDown(2) generates the output 2..., then it calls countDown(1)
- 2. countDown(1) generates the output 1..., then it calls countDown(0)
- 3. countDown(0) generates the output Blastoff!, then returns to countDown(1)
- 4. countDown(1) returns to countDown(2)
- 5. countDown(2) returns to the calling function

### What Happens When Called?





20.2

#### Solving Problems with Recursion

#### Recursive Functions - Purpose

- Recursive functions are used to reduce a complex problem to a simpler-to-solve problem.
- The simpler-to-solve problem is known as the <u>base case</u>
- Recursive calls stop when the base case is reached

- A recursive function must always include a test to determine if another recursive call should be made, or if the recursion should stop with this call
- In the sample program, the test is:

```
if (num == 0)
```

```
void countDown(int num)
{
   if (num == 0) // test
     cout << "Blastoff!";
   else
   {
     cout << num << "...\n";
     countDown(num-1); // recursive
   }
} // call
}</pre>
```

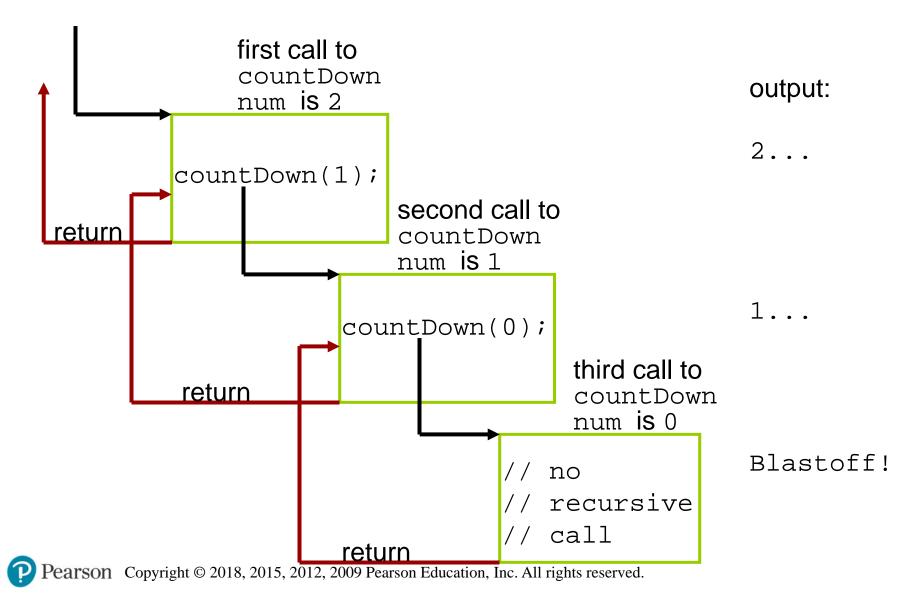
- Recursion uses a process of breaking a problem down into smaller problems until the problem can be solved
- In the countDown function, a different value is passed to the function each time it is called
- Eventually, the parameter reaches the value in the test, and the recursion stops

```
void countDown(int num)
  if (num == 0)
     cout << "Blastoff!";</pre>
  else
      cout << num << "...\n";
     countDown(num-1);// note that the value
                        // passed to recursive
                        // calls decreases by
                        // one for each call
```

### What Happens When Called?

- Each time a recursive function is called, a new copy of the function runs, with new instances of parameters and local variables created
- As each copy finishes executing, it returns to the copy of the function that called it
- When the initial copy finishes executing, it returns to the part of the program that made the initial call to the function

### What Happens When Called?



#### Types of Recursion

- Direct
  - a function calls itself
- Indirect
  - function A calls function B, and function B calls function A
  - function A calls function B, which calls ..., which calls function A

# The Recursive Factorial Function

The factorial function:

```
n! = n*(n-1)*(n-2)*...*3*2*1 if n > 0

n! = 1 if n = 0
```

Can compute factorial of n if the factorial of (n-1) is known:

```
n! = n * (n-1)!
```

on = 0 is the base case

# The Recursive Factorial Function

```
int factorial (int num)
{
  if (num > 0)
    return num * factorial(num - 1);
  else
    return 1;
}
```

#### Program 20-3

```
1  // This program demonstrates a recursive function to
2  // calculate the factorial of a number.
3  #include <iostream>
4  using namespace std;
5
6  // Function prototype
7  int factorial(int);
8
9  int main()
10  {
11   int number;
12
```

(program continues)

#### Program 20-3 (Continued)

```
// Get a number from the user.
        cout << "Enter an integer value and I will display\n";</pre>
14
15
        cout << "its factorial: ";</pre>
16
        cin >> number;
17
18
        // Display the factorial of the number.
        cout << "The factorial of " << number << " is ";</pre>
19
        cout << factorial(number) << endl;</pre>
20
21
         return 0:
22
23
24
    // Definition of factorial. A recursive function to calculate *
26
    // the factorial of the parameter n.
27
28
29
   int factorial(int n)
30
31
        if (n == 0)
32
             return 1:
                                            // Base case
33
         else
             return n * factorial(n - 1); // Recursive case
34
35 }
Program Output with Example Input Shown in Bold
Enter an integer value and I will display
its factorial: 4 Enter
The factorial of 4 is 24
```

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20.3

#### The Recursive gcd Function

## The Recursive gcd Function

- Greatest common divisor (gcd) is the largest factor that two integers have in common
- Computed using Euclid's algorithm:

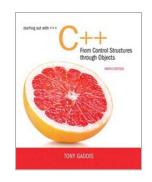
```
gcd(x, y) = y \text{ if } y \text{ divides } x \text{ evenly}

gcd(x, y) = gcd(y, x % y) \text{ otherwise}
```

ogcd(x, y) = y is the base case

## The Recursive gcd Function

```
int gcd(int x, int y)
{
   if (x % y == 0)
      return y;
   else
      return gcd(y, x % y);
}
```



20.4

## Solving Recursively Defined Problems

# Solving Recursively Defined Problems

- The natural definition of some problems leads to a recursive solution
- Example: Fibonacci numbers:

```
0, 1, 1, 2, 3, 5, 8, 13, 21, ...
```

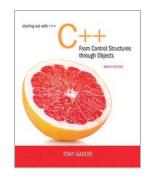
- After the starting 0, 1, each number is the sum of the two preceding numbers
- Recursive solution:

```
fib(n) = fib(n - 1) + fib(n - 2);
```

Base cases: n <= 0, n == 1</pre>

# Solving Recursively Defined Problems

```
int fib(int n)
{
   if (n <= 0)
      return 0;
   else if (n == 1)
      return 1;
   else
      return fib(n - 1) + fib(n - 2);
}</pre>
```



20.5

#### Recursive Linked List Operations

# Recursive Linked List Operations

- Recursive functions can be members of a linked list class
- Some applications:
  - Compute the size of (number of nodes in) a list
  - Traverse the list in reverse order

# Counting the Nodes in a Linked List

- Uses a pointer to visit each node
- Algorithm:
  - pointer starts at head of list
  - If pointer is null pointer, return 0 (base case)
     else, return 1 + number of nodes in the list pointed to by current node
- See the NumberList class in Chapter 19

## The countNodes function, a private member function

The countNodes function is executed by the public numNodes function:

# Contents of a List in Reverse Order

- Algorithm:
  - pointer starts at head of list
  - If the pointer is null pointer, return (base case)
  - If the pointer is not null pointer, advance to next node
  - Upon returning from recursive call, display contents of current node

# The showReverse function, a private member function

The showReverse function is executed by the public displayBackwards function:



20.6

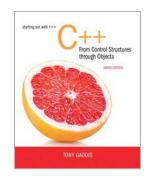
## A Recursive Binary Search Function

# A Recursive Binary Search Function

- Binary search algorithm can easily be written to use recursion
- Base cases: desired value is found, or no more array elements to search
- Algorithm (array in ascending order):
  - If middle element of array segment is desired value, then done
  - Else, if the middle element is too large, repeat binary search in first half of array segment
  - Else, if the middle element is too small, repeat binary search on the second half of array segment

#### A Recursive Binary Search Function (Continued)

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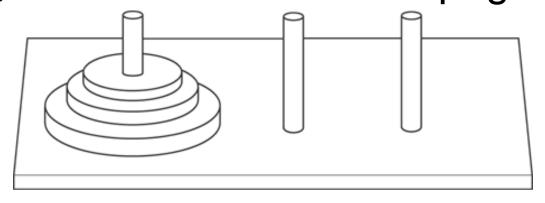


20.7

#### The Towers of Hanoi

#### The Towers of Hanoi

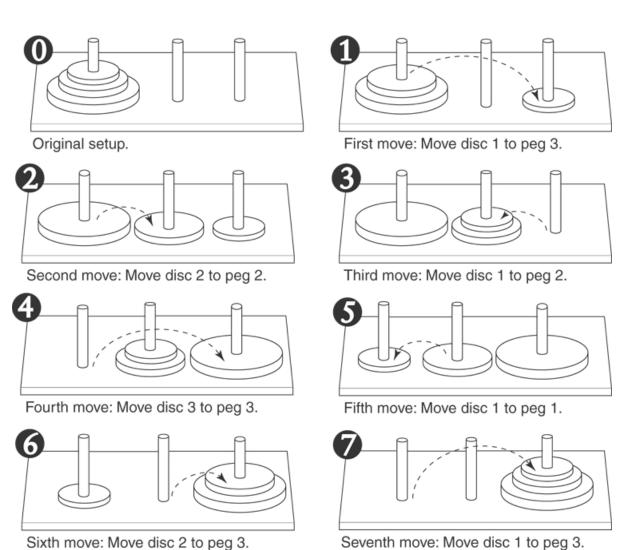
- The Towers of Hanoi is a mathematical game that is often used to demonstrate the power of recursion.
- The game uses three pegs and a set of discs, stacked on one of the pegs.

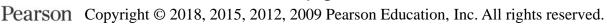


#### The Towers of Hanoi

- The object of the game is to move the discs from the first peg to the third peg. Here are the rules:
  - Only one disc may be moved at a time.
  - A disc cannot be placed on top of a smaller disc.
  - All discs must be stored on a peg except while being moved.

## Moving Three Discs





### The Towers of Hanoi

- The following statement describes the overall solution to the problem:
  - Move n discs from peg 1 to peg 3 using peg 2 as a temporary peg.

### The Towers of Hanoi

- Algorithm
  - To move n discs from peg A to peg C, using peg B as a temporary peg:

If n > 0 Then
Move n − 1 discs from peg A to peg B, using peg C as a temporary peg.

Move the remaining disc from the peg A to peg C.

Move n-1 discs from peg B to peg C, using peg A as a temporary peg.

End If

#### Program 20-10

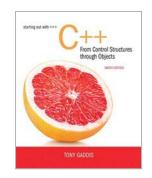
```
// This program displays a solution to the Towers of
   // Hanoi game.
   #include <iostream>
    using namespace std;
 5
    // Function prototype
    void moveDiscs(int, int, int, int);
 8
 9
    int main()
10
11
        const int NUM_DISCS = 3; // Number of discs to move
12
        const int FROM_PEG = 1; // Initial "from" peg
13
        const int TO_PEG = 3;  // Initial "to" peg
        const int TEMP PEG = 2; // Initial "temp" peg
14
15
16
        // Play the game.
        moveDiscs(NUM_DISCS, FROM_PEG, TO_PEG, TEMP_PEG);
17
        cout << "All the pegs are moved!\n";</pre>
18
                                                               (program continues)
```

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#### Program 20-10 (continued) 19 return 0; 20 21 22 // The moveDiscs function displays a disc move in // the Towers of Hanoi game. 24 25 // The parameters are: // num: The number of discs to move. 26 27 // fromPeg: The peg to move from. // toPeg: The peg to move to. 28 // tempPeg: The temporary peg. 29 30 31 32 void moveDiscs(int num, int fromPeg, int toPeg, int tempPeg) 33 { 34 if (num > 0)35 36 moveDiscs(num - 1, fromPeg, tempPeg, toPeg); 37 cout << "Move a disc from peg " << from Peg << " to peg " << toPeg << endl; 38 39 moveDiscs(num - 1, tempPeg, toPeg, fromPeg); 40 41 }

#### Program 20-10 (Continued)

```
Move a disc from peg 1 to peg 3
Move a disc from peg 1 to peg 2
Move a disc from peg 3 to peg 2
Move a disc from peg 1 to peg 3
Move a disc from peg 1 to peg 3
Move a disc from peg 2 to peg 1
Move a disc from peg 2 to peg 3
Move a disc from peg 1 to peg 3
All the pegs are moved!
```

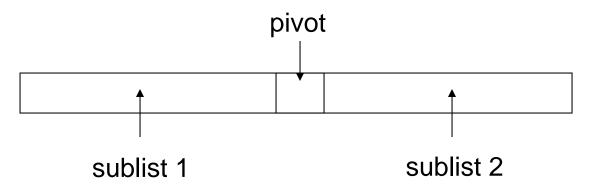


20.8

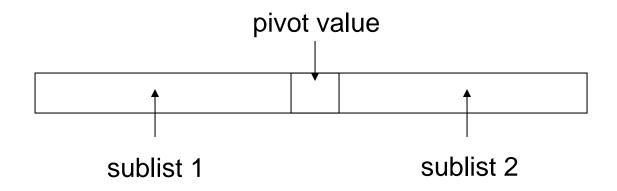
#### The QuickSort Algorithm

## The QuickSort Algorithm

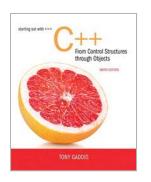
- Recursive algorithm that can sort an array or a linear linked list
- Determines an element/node to use as pivot value:



## The QuickSort Algorithm



- Once pivot value is determined, values are shifted so that elements in sublist1 are < pivot and elements in sublist2 are > pivot
- Algorithm then sorts sublist1 and sublist2
- Base case: sublist has size 1
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20.9

## Exhaustive and Enumeration Algorithms

# Exhaustive and Enumeration Algorithms

- Exhaustive algorithm: search a set of combinations to find an optimal one Example: change for a certain amount of money that uses the fewest coins
- Uses the generation of all possible combinations when determining the optimal one.



20.10

Recursion vs. Iteration

#### Recursion vs. Iteration

- Benefits (+), disadvantages(-) for recursion:
  - + Models certain algorithms most accurately
  - + Results in shorter, simpler functions
  - May not execute very efficiently
- Benefits (+), disadvantages(-) for iteration:
  - + Executes more efficiently than recursion
  - Often is harder to code or understand