## Keys and Null Values

If there are nulls, keys do not work that well:

- They do not guarantee unique identification;
- They do not help in establishing correspondences between data in different relations

RegNum	Surname	FirstName	BirthDate	DegreeProg	
NULL	Smith	John	NULL	Computing	
587614	Smith	Lucy	01/05/61	Engineering	
934856	Black	Lucy	NULL	NULL	
NULL	Black	Lucy	05/03/58	Engineering	

- Are the third and fourth tuple the same?
- How do we access the first tuple?

## Primary Keys

- The presence of nulls in keys has to be limited.
- Each relation must have a primary key on which nulls are not allowed (in any attribute)
- Notation: the attributes of the primary key are <u>underlined</u>
- References between relations are realized through primary keys

RegNum	Surname	FirstName	BirthDate	DegreeProg
643976	Smith	John	NULL	Computing
587614	Smith	Lucy	01/05/61	Engineering
934856	Black	Lucy	NULL	NULL
735591	Black	Lucy	05/03/58	Engineering

### Do we Always Have Primary Keys?

- In most cases, we do have reasonable primary keys (e.g., student number, SIN)
- There may be multiple keys, one of which is designated as primary.

## Recap

- $\square$  A set of fields is a <u>key</u> for a relation if:
  - 1. No two distinct tuples can have same values in all key fields, and
  - 2. This is not true for any subset of the key.
- ☐ If #2 false, then a *superkey*.
- □ If there's >1 key for a relation, one of the keys is chosen to be the *primary key*.
- E.g., sid is a key for Students. (What about name?) The set {sid, gpa} is a superkey.

### Primary and Candidate Keys

- 1. "For a given student and course, there is a single grade." vs.
- 2. "Students can take only one course, and receive a single grade for that course; further, no two students in a course receive the same grade."
- Be careful to define Integrity Constraints (ICs) correctly at design time.
- ICS are checked when data is updated.

Enrolled(sid, cid, grade)

Enrolled(sid, cid, grade)

Enrolled(sid, cid, grade)

• key (cid, grade)

## Foreign Keys

- Pieces of data in different relations are correlated by means of values of primary keys.
- Referential integrity constraints are imposed in order to guarantee that the values refer to existing tuples in the referenced relation.
- A foreign key requires that the values on a set X of attributes of a relation R<sub>1</sub> must appear as values for the primary key of another relation R<sub>2</sub>.
  - In other words, set of attributes in one relation that is used to `refer' to a tuple in another relation. (Must correspond to primary key of the second relation.) Like a `logical pointer'.

# Referential Integrity

- □ E.g. *sid* is a foreign key referring to Students:
  - Enrolled(sid: string, cid: string, grade: string)
  - If all foreign key constraints are enforced, <u>referential</u> <u>integrity</u> is achieved, i.e., no dangling references.

# Referential Integrity (cont'd)

 Only students listed in the Students relation should be allowed to enroll for courses.

#### Enrolled

sid	cid	grade	ade					
53666	Carnatic101	C		sid	name	login	age	gpa
	Reggae203	B -		53666	Jones	jones@cs	18	3.4
	Topology112	Δ	7	53688	Smith	smith@eecs	18	3.2
	History 105	B	<b>/</b>	53650	Smith	smith@math	19	3.8
33000	1115tO1 y 105	י ע						

Studente