

**Inference, Models and Simulation for Complex Systems**  
**CSCI 7000-003, Fall 2010**  
**Prof. Aaron Clauset**  
**Problem Set 3, due 10/13**

1. Erdős-Rényi random graphs

(**step 1**) Implement an algorithm that generates an instance of an Erdős-Rényi random graph  $G(n, p)$ , where  $n$  is the number of nodes and  $p$  is the probability that a given edge  $(i, j)$  exists.

(**step 2**) Implement an algorithm (such as those discussed in Lecture 8) that takes as input an unweighted, undirected simple network and gives as output a vector  $\vec{c}$  where  $c_i$  is the index of the component to which the  $i$ th vertex belongs.

- (a) Using both your  $G(n, p)$  generator and components algorithms from above, conduct a numerical experiment to demonstrate the phase transition in the size of the largest component, as a function of  $c$ , the average degree. The product of the experiment should be (i) a brief description of what you did, (ii) a figure showing both your numerical results for the size of the largest component  $S$ , measured as a fraction of  $n$ , and the theoretical prediction we derived in Lecture 9, as  $c$  goes from close to zero to well above the critical threshold at  $c = 1$ , and (iii) some comments about the shape and location of the transition, and its agreement with the analytic prediction.

(Hint: choose  $n > 10^3$ , average your results over several graphs at a particular value of  $c$  and choose at least 50 values of  $c$  on the interval  $0 \leq c \leq 5$ .)

- (b) Derive analytically or measure numerically the component size distribution at (i)  $c < 1$ , (ii)  $c = 1$  and (iii)  $c > 1$ . Present and comment on your results. Include a visualization of each type of network for  $n = 100$ .

(Hint: For the visualization, write your own or use an existing implementation of the Fruchterman-Reingold spring-embedder algorithm. The course webpage includes links to several existing pieces of software that can do this.)

- (c) (**optional**) For  $c = 5$ , measure (via numerical simulation) the diameter of a  $G(n, p)$  graph as a function of  $n$ . Show that the estimated diameter grows like  $O(\log n)$  (a straight line on linear-log axes).

(Hint: use a dozen or more values of  $n$  spaced logarithmically over several orders of magnitude, and average over several instances at a given value of  $n$ .)

- (d) **(optional)** Recall that the configuration model constructs a random graph with a particular degree sequence. Using tools from class, study the case of random graphs with power-law degree distributions (for  $2 < \alpha < 3$ ). Characterize the location of the high-degree vertices with respect to other low- or high-degree nodes. Present your results and discuss.  
(Hint: think centrality and assortativity measures.)
- (e) **(optional)** As a function of  $c$ , derive analytically or measure numerically the average time for an unbiased random walker to return to its starting location in a  $G(n, p)$  network for  $c > 1$ . Here, unbiased means the walker chooses its next location uniformly at random from its current neighbors.

## 2. Mathematical exercises

- (a) Consider an undirected (connected) tree of  $n$  vertices. Suppose that a particular vertex in the tree has degree  $k$ , so that its removal would divide the tree into  $k$  disjoint regions, and suppose that the sizes of those regions are  $n_1, \dots, n_k$ . Show that the betweenness centrality  $b$  of the vertex is

$$b = n^2 - \sum_{i=1}^k n_i^2 .$$

- (b) Consider a bipartite network with  $n_1$  vertices of type 1 and  $n_2$  vertices of type 2. Show that the mean degrees  $\langle k \rangle_1$  and  $\langle k \rangle_2$  of the two types are related by

$$\langle k \rangle_2 = \langle k \rangle_1 \frac{n_1}{n_2} .$$

- (c) An alternative definition of the *closeness centrality* is  $c'_i = n / \sum_j d_{ij}$ .

Consider an undirected, unweighted network of  $n$  vertices that contains exactly two subnetworks of size  $n_A$  and  $n_B$ , which are connected by a single edge  $(A, B)$ . Show that the closeness centralities  $c'_A$  and  $c'_B$  of these two vertices are related by

$$\frac{1}{c'_A} + \frac{n_A}{n} = \frac{1}{c'_B} + \frac{n_B}{n} .$$

- (d) What is the time complexity, as a function of the number of vertices  $n$  and the number of edges  $m$ , of the following network operations, if the network in question is stored in an adjacency list format? Give pseudo-code for each solution.
- i. Calculating the mean degree.
  - ii. Calculating the median degree.
  - iii. Calculating the clustering coefficient.