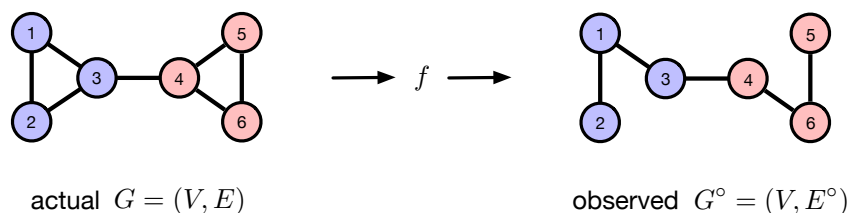


## 1 Predicting Missing Links in Networks

Recall that most network data sets are *incomplete* in some way, either because of the way we measured them or because they are dynamic objects that change over time and we only get to observe them for part of their lifetime. If edges are the things that are missing, our goal is to predict the missing links (or, future connections) based on their correlation with the links we observe.

In this setting, we assume that we have a complete account of a set of nodes  $V$ , but an incomplete account of the edges among them. Letting  $E$  denote the set of actual edges, and  $E^\circ \subset E$  be the subset of them that we observe, we can define the observed network as  $G^\circ = (V, E^\circ)$  and the actual network as  $G = (V, E)$ . Just as with missing node attributes, there is some *missingness function*  $f$  that chooses which subset of edges  $E^\circ$  we actually observe. If there are some patterns within the observed edges  $E^\circ$  that correlate with the actual edges  $E$ , then we may be able to predict (better than baseline guessing) which of the observed unconnected pairs  $Y = V \times V - E^\circ$  are in fact missing links  $X = E - E^\circ$ .

Here is a small example of this process of a missingness function  $f$  being applied to an actual graph  $G$  (on the left) to produce the observed network  $G^\circ$  (on the right), where the removed lines are the missing links.



Mathematically, predicting missing links is a much harder problem than predicting missing attributes of either nodes or edges. With missing attributes, we start off knowing which attributes are missing  $x_i = \emptyset$ , but with missing links, we do not. Instead, our task is to sort among all the 0s of the adjacency matrix, which are the “non-edges” or unconnected pairs in the observed graph  $G^\circ$ , to find those that should be 1s. Because most real-world networks are sparse, this task is like searching for  $O(n)$  needles in a  $\Theta(n^2)$  haystack. Hence, the baseline accuracy for guessing will be  $O(1/n)$ , making it extremely unlikely to guess correctly just by chance.

Link prediction methods operate by defining a *score* function over unconnected pairs  $i, j \in X$ , and the better the predictor, the more likely it will assign a higher score to a pair  $i, j$  that is actually a missing link. The idea of guessing at random provides a simple baseline algorithm—if we can do better than this baseline, then we’re getting somewhere—and this algorithm is equally likely to assign any particular score to each of the unconnected pairs.

**The baseline predictor.** In the absence of any information about  $f$  or about any patterns within  $E^\circ$  and their relationship to the missing links in  $Y$ , the simplest **baseline prediction** algorithm assigns a value that is independent of the graph  $G$ , meaning that every input pair  $i, j$  is equally likely to receive a particular score. We can formalize this notion mathematically, as follows:

$$\text{if } i, j \in X \quad \text{score}(i, j) = \text{Uniform}(0, 1) \quad . \quad (1)$$

That is, the baseline assigns a uniformly random score between 0 and 1 to each unconnected pair. If we apply this algorithm to the above example, and then *sort* the candidate pairs  $i, j$  by their  $\text{score}(i, j)$  values, we obtain a score table:

| $i$ | $j$ | $\text{score}(i, j)$ |
|-----|-----|----------------------|
| 1   | 5   | $r$                  |
| 1   | 4   | $r$                  |
| 1   | 6   | $r$                  |
| 2   | 3   | $r$                  |
| 2   | 6   | $r$                  |
| 2   | 4   | $r$                  |
| 2   | 5   | $r$                  |
| 3   | 5   | $r$                  |
| 3   | 6   | $r$                  |
| 4   | 5   | $r$                  |

where  $r$  is a stand-in for  $\text{Uniform}(0, 1)$ , and the ordering is arbitrary among tied scores. (Note that in practice, ties must be broken randomly. Do you see why?) You can see the two missing edges  $Y = \{(2, 3), (4, 5)\}$  in the list, but there's nothing about their scores to suggest that they are any more or less likely to be the missing links than another pair.

If some correlation exists between the observed edges  $E^\circ$  and the missing edges  $Y$ , then we can likely improve considerably over this baseline. Just like with predicting missing node attributes, there are many ways we could do this. In practice, there are three approaches:

- **Topological predictors** are simple functions of the joint local network structure of the pair of nodes  $i, j$  being scored, e.g., their degrees, measures of overlapping neighborhoods, geodesic distance, and other measures that summarize a network's structure.
- **Model-based predictors** are a broad class of algorithms that rely on models of large-scale network structure, e.g., by decomposing the network into modules or communities, to make predictions about missing links.
- **Embedding-based predictors** are second broad class of algorithms that assign nodes to locations in a  $d$ -dimensional latent space in such a way that nodes are “close” in the embedded space if they are, or are likely to be, connected.

Here, we will focus primarily on topological predictors, and then briefly consider the remaining two in Section 1.4.

## 1.1 Topological predictors and local smoothing

Topological predictors use the structure of the graph itself, from the perspective of nodes  $i, j$  as a way to create a score function that may correlate with the missing links. Each such predictor is based on a specific assumption about how edges form. For instance, we might score a pair by the number of common neighbors (the more, the better), the number of shortest paths between the nodes, the product of their degrees, etc.<sup>1</sup>

Just as node attributes are often assortative (or, homophilous) in networks, edges themselves are often **assortative**, meaning that they tend to cluster together. This pattern is the idea behind the clustering coefficient, which measures the “local” density of edges. And, just as it did for node attributes, this pattern allows us to construct a kind of local smoothing algorithm for predicting missing links. Specifically, we can

*predict a missing link to occur where it would cluster with other edges.*

There are many ways to operationalize such an idea into a specific algorithm for predicting missing links. We will explore two. One uses the **Jaccard coefficient**, which quantifies the degree of overlap among two nodes’ neighborhoods. Using the same notation as in the Lecture on predicting missing node attributes, let  $\nu(i)$  return the set of neighbors of node  $i$ . The Jaccard coefficient is defined as

$$\text{Jaccard}(i, j) = \frac{|\nu(i) \cap \nu(j)|}{|\nu(i) \cup \nu(j)|} , \quad (2)$$

which measures the fraction of neighbors of either node that are neighbors of both nodes, i.e., the density of common neighbors. Two nodes with many common neighbors will have a higher Jaccard coefficient. In that case, adding an edge between  $i, j$  would then “close” each of the open triads  $i, k, j$ , which would increase both their local, and the global, clustering coefficient.

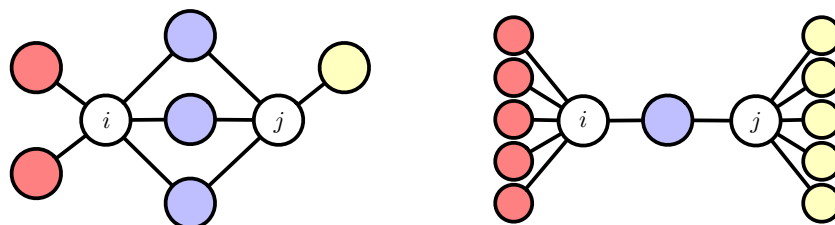
To turn this pairwise network measure into a prediction algorithm, we say

$$\text{score}(i, j) = \text{Jaccard}(i, j) + \text{Uniform}(0, \epsilon) , \quad (3)$$

where  $\text{Uniform}(0, \epsilon)$  adds a small amount of noise in order to break ties randomly without changing the relative ordering outside those ties. (Do you see why this is necessary?)

Consider the following small examples. On the left,  $\nu(i)$  returns 5 nodes,  $\nu(j)$  returns 4 nodes,

<sup>1</sup>There are an enormous number of topological predictors for missing links. Examples include the number of common neighbors, shortest-path betweenness centrality, the Leicht–Holme–Newman index, personalized page rank, shortest path length, the mean neighbor entries within a low rank approximation, the Jaccard coefficient, the Adamic–Adar index, the resource allocation index, the dot product of columns  $i$  and  $j$  in a low rank approximation via a singular value decomposition, the local clustering coefficients of  $i$  and  $j$ , the average neighbor degrees, and any number of “centrality” measures.



and there are 3 common neighbors. Hence, the  $\text{Jaccard}(i, j) = 0.50$ , reflecting the relatively high proportion of total neighbors that  $i, j$  have in common. On the right,  $i, j$  have many fewer common neighbors, and their Jaccard coefficient is correspondingly smaller, only 0.091.

A second local topological predictor is the **degree product**, which embodies the idea we learned from our analysis of random graphs that nodes with high degrees are likely themselves to be connected, just by chance. We define this predictor as  $\text{score}(i, j) = k_i k_j + \text{Uniform}(0, \epsilon)$ . How would degree-product scores differ from those of the Jaccard coefficient for the above examples?

Let us return to the example network on the first page of this lecture, from which we removed two edges, and then apply the Jaccard and degree-product predictors to the observed network  $G'$ . Doing so produces the following score tables:

| $i$      | $j$      | Jaccard<br>$\text{score}(i, j)$ | $i$      | $j$      | degree product<br>$\text{score}(i, j)$ |
|----------|----------|---------------------------------|----------|----------|--|
| <b>4</b> | <b>5</b> | $1/2 + r$                       | 1        | 4        | $4 + r$                                |
| <b>2</b> | <b>3</b> | $1/2 + r$                       | 1        | 6        | $4 + r$                                |
| 3        | 6        | $1/3 + r$                       | 3        | 6        | $4 + r$                                |
| 1        | 4        | $1/3 + r$                       | 1        | 5        | $2 + r$                                |
| 1        | 5        | $r$                             | <b>2</b> | <b>3</b> | $2 + r$                                |
| 1        | 6        | $r$                             | 2        | 6        | $2 + r$                                |
| 2        | 6        | $r$                             | 2        | 4        | $2 + r$                                |
| 2        | 4        | $r$                             | 3        | 5        | $2 + r$                                |
| 2        | 5        | $r$                             | <b>4</b> | <b>5</b> | $2 + r$                                |
| 3        | 5        | $r$                             | 2        | 5        | $1 + r$                                |

where the missing links  $Y = \{(2, 3), (4, 5)\}$  are highlighted, and links with the same score are ordered arbitrarily.

On the left, the Jaccard link predictor does very well at assigning these missing links high scores compared to the non-missing links, while on the right, the degree product seems very poor. These different behaviors illustrate a key point about link predictors: each predictor captures different types of correlations between observed and missing edges, and hence each will perform well on some inputs and poorly on others. In this case, the actual network has substantial local edge density, and so the Jaccard function does well. But the degree product predictor would perform better

if the network had much higher variance in its degree structure.

To quantify these differences in performance, compare them to the baseline, or understand how close they might be to an *optimal* predictor, we need a compact way of quantifying their performance.

## 1.2 Measuring performance: the AUC

Predicting missing links is a binary classification problem: every candidate  $i, j \in X$  is either a missing link or not.

To summarize the performance of a link prediction algorithm, we can use the **AUC** statistic,<sup>2 3</sup> a context-agnostic measure of its ability to distinguish a missing link  $i, j \in Y$  (a true positive, or TP) from a non-edge  $X - Y$  (a true negative, or TN). The AUC has two other attractive properties:

1. *scale invariance*, meaning it is not dependent on the scores themselves, and instead only depends on their relative values, and
2. *threshold invariance*, meaning it is a general measure of accuracy that doesn't require choosing a threshold on the scores for making predictions.

Mathematically, the AUC can be defined as

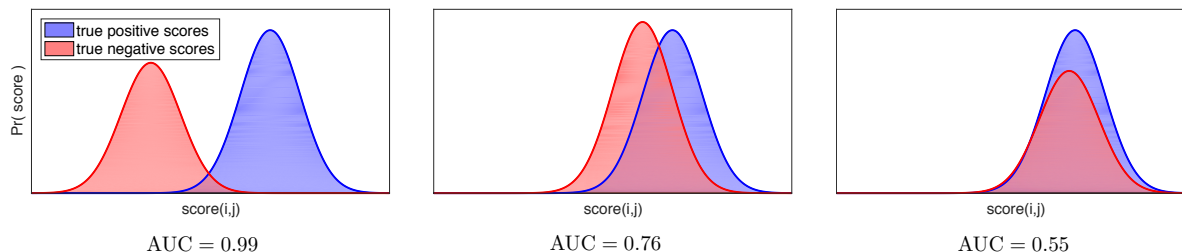
$$\text{AUC} = \Pr[\text{score}(\text{TP}) > \text{score}(\text{TN})] . \quad (4)$$

That is, if we choose a uniformly random missing edge  $a, b \in Y$  (TP) and a uniformly random non-edge  $c, d \in X - Y$  (TN), the AUC is the probability that  $\text{score}(a, b)$  is higher than  $\text{score}(c, d)$ . The figure below illustrates this idea, for three hypothetical predictors. On the left, the predictor assigns scores in a way that the score distribution for true positives is well-separated from the score distribution of true negatives, which produces a high AUC value. In the middle example, the predictor assigns scores so that the distributions overlap more, which lowers the AUC. And, on the right, the predictor assigns scores so that the distributions are nearly the same, leading to a very low AUC.

If the algorithm is no better than guessing at random, then  $\text{AUC} = 0.5$  and it is as equally likely that  $\text{score}(\text{TP}) < \text{score}(\text{TN})$  as it is that  $\text{score}(\text{TP}) > \text{score}(\text{TN})$ . If that's the case, the probability of the latter is  $1/2$ , and hence that the  $\text{AUC} = 0.5$ . The baseline predictor of Eq. (1), which

<sup>2</sup>AUC is short for “Area Under the Curve,” where the “curve” here is the Receiver Operating Characteristic (ROC) curve. ROC curves were invented during World War 2 as a method for assessing the performance of radar at detecting enemy aircraft. Neato. See [https://en.wikipedia.org/wiki/Receiver\\_operating\\_characteristic](https://en.wikipedia.org/wiki/Receiver_operating_characteristic)

<sup>3</sup>There are other ways to quantify accuracy in this setting, each of which has a slightly different way of weighting the cost of different types of errors. Common choices include the F1-measure, which is the harmonic mean of the precision and recall, or the entire precision-recall curve. These alternatives are less context agnostic than the AUC, and hence may provide a better guide to performance in particular settings.



assigns a uniformly random value to every candidate, has exactly this behavior and thus has an  $\text{AUC} = 0.5$ . Hence, any algorithm with  $\text{AUC} > 0.5$  does better than chance (the baseline). The maximum value of  $\text{AUC} = 1.0$  is attained only when the algorithm assigns a higher score to every missing link (true positive) than it does to any non-edge (true negative).

In practice, we compute a link prediction algorithm's AUC by numerically integrating its corresponding "ROC" curve to get, literally, the area under the curve.<sup>4</sup> Formally, the ROC curve is a parametric plot of the *true positive rate* (TPR) and the *false positive rate* (FPR), as a function of prediction threshold. The idea is straightforward. Imagine drawing a line across one of our score tables just below the  $\ell$ th row, and then sliding it up, or down the table. This line represents a prediction "threshold": all the rows  $\ell$  and above we predict to be missing links, while all rows  $\ell + 1$  and below we predict to be non-missing links. The TPR is the fraction of all the actual missing links that are in the predicted-to-be-missing set (at or above row  $\ell$ ), and the FPR is the fraction of all the actual non-missing links (non-edges) that are in that same set. Plotting  $\text{TPR}(\ell)$  vs.  $\text{FPR}(\ell)$  sweeps out the ROC curve on the unit square, starting at  $(0,0)$  and ending at  $(1,1)$ .

To calculate the AUC, we begin by augmenting the score table with three additional columns: a column  $\tau$  that gives the *actual* status of the  $\ell$ th row (pair  $i,j$ ) as either a true positive ( $\tau_\ell = 1$ ) or true negative ( $\tau_\ell = 0$ ), and then two columns, one for  $\text{TPR}(\ell)$  and one for  $\text{FPR}(\ell)$ . If a predictor performs well, then in the  $\tau$  column, the 1s will tend to be located higher up in the table, while if it is no better than baseline, the 1s will be scattered uniformly at random within this column.<sup>5</sup>

Once the  $\tau$  column has been added, we can fill in the TPR and FPR columns simultaneously in a single scan down the table. Let  $\mathcal{T} = |Y|$  be the total number of true positives (missing links; equal to the column sum of  $\tau$ ), and let  $\mathcal{F} = |X - Y|$  be the total number of true negatives (non-edges;

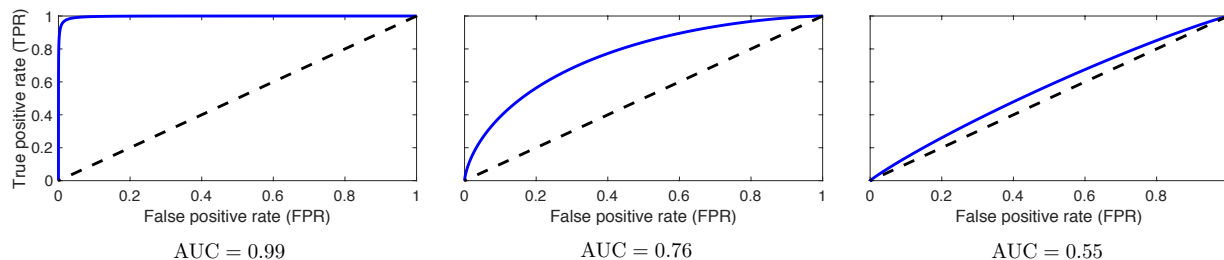
<sup>4</sup>We can also estimate the AUC via Monte Carlo, by evaluating Eq. (4) directly: sample many pairs of true positives and true negatives, and estimate the fraction of draws for which the former is assigned a higher score than the latter.

<sup>5</sup>If the 1s tend to be stacked at the bottom of the  $\tau$  column, then the predictor is, in fact, pretty good at distinguishing 1s from 0s, it just gets it backwards. We can then construct a better-than-baseline predictor out of a worse-than-baseline predictor by simply doing the opposite of what it says (or by inverting its output scores).

equal to the length of the table minus the sum of  $\tau$ ). The  $\text{TPR}(\ell)$  and  $\text{FPR}(\ell)$  are then defined as

$$\text{TPR}(\ell) = \frac{1}{\mathcal{T}} \sum_{k=1}^{\ell} \tau_k \quad \text{FPR}(\ell) = \frac{1}{\mathcal{F}} \sum_{k=1}^{\ell} 1 - \tau_k \quad . \quad (5)$$

Using the same three sets of hypothetical score distributions we saw above, we can now compute the corresponding TPR and FPR functions, and then plot them against each other to obtain the ROC curve for each case. As a reference, each plot also shows the  $\text{TPR} = \text{FPR}$  line, which represents the baseline predictor's performance.



The area under each ROC curve is each hypothetical predictor's AUC, which we calculate by numerically integrating the ROC curve, using a simple box-rule approximation<sup>6</sup> technique:

$$\text{AUC} = \sum_{\ell=1}^{\mathcal{F}} \text{TPR}(\ell) \times [\text{FPR}(\ell) - \text{FPR}(\ell - 1)] \quad , \quad (6)$$

where we define  $\text{FPR}(0) = 0$ .

### 1.3 An example

To see how this works in practice, let's consider again the Jaccard and degree-product link predictors from Section 1.1. The two score tables below now include the three new columns: the actual status  $\tau$ , the corresponding true positive rates (TPR), and the corresponding false positive rates (FPR) for a threshold drawn between that row and the next. (Exercise: verify that the columns were calculated correctly.)

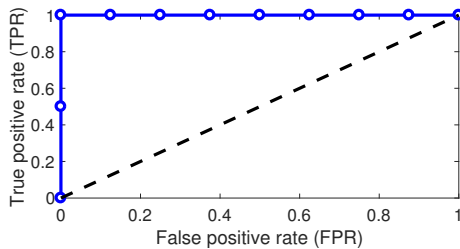
From these, we can plot the corresponding ROC curves, and apply Eq. (6) to calculate the AUC for each predictor. Unsurprisingly, the Jaccard predictor achieves perfect performance, with an  $\text{AUC} = 1.00$ . In this network, the degree-product predictor is in fact *worse* than guessing at

<sup>6</sup>Recall the box rule for numerical integration: given a function  $f(x)$ , a range  $[x_s, x_t]$ , and an increment  $\Delta x$ , the integral  $\int_{x_s}^{x_t} f(x) dx \approx \sum_{x=x_s}^{x_t} f(x) \times \Delta x$ . The trapezoid rule is slightly more accurate, although the difference is negligible when  $\Delta x$  is small, as is the case in link prediction (do you see why?).

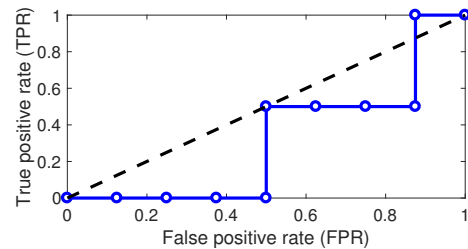
random, producing an  $AUC = 0.31$ . This value is not as precise as it seems, however. Recall that the value  $r$  in the score column is a random variable. This means that the positions of the two true positives in the degree-product table will vary within the block of  $\text{score}(i, j) = 2 + r$  rows, and these lead to slightly different TPR and FPR functions. The right thing to do here is to compute an *average AUC*, when these ties are broken randomly, which produces a more reliable estimate of the AUC. The highest value possible, if the ties break luckily for the missing links is 0.62, while the lowest, is 0.12 (do you see why?). The average is  $AUC = 0.37$ .

| $i$ | $j$ | $\tau_k$ | $TPR_\ell$ | $FPR_\ell$ | Jaccard<br>$\text{score}(i, j)$ |
|-----|-----|----------|------------|------------|---------------------------------|
| 4   | 5   | 1        | 0.5        | 0.0        | $1/2 + r$                       |
| 2   | 3   | 1        | 1.0        | 0.0        | $1/2 + r$                       |
| 3   | 6   | 0        | 1.0        | 0.125      | $1/3 + r$                       |
| 1   | 4   | 0        | 1.0        | 0.250      | $1/3 + r$                       |
| 1   | 5   | 0        | 1.0        | 0.375      | $r$                             |
| 1   | 6   | 0        | 1.0        | 0.500      | $r$                             |
| 2   | 6   | 0        | 1.0        | 0.625      | $r$                             |
| 2   | 4   | 0        | 1.0        | 0.750      | $r$                             |
| 2   | 5   | 0        | 1.0        | 0.875      | $r$                             |
| 3   | 5   | 0        | 1.0        | 1.00       | $r$                             |

| $i$ | $j$ | $\tau_k$ | $TPR_\ell$ | $FPR_\ell$ | degree product<br>$\text{score}(i, j)$ |
|-----|-----|----------|------------|------------|--|
| 1   | 4   | 0        | 0.0        | 0.125      | $4 + r$                                |
| 1   | 6   | 0        | 0.0        | 0.250      | $4 + r$                                |
| 3   | 6   | 0        | 0.0        | 0.375      | $4 + r$                                |
| 1   | 5   | 0        | 0.0        | 0.500      | $2 + r$                                |
| 2   | 3   | 1        | 0.5        | 0.500      | $2 + r$                                |
| 2   | 6   | 0        | 0.5        | 0.625      | $2 + r$                                |
| 2   | 4   | 0        | 0.5        | 0.750      | $2 + r$                                |
| 3   | 5   | 0        | 0.5        | 0.875      | $2 + r$                                |
| 4   | 5   | 1        | 1.0        | 0.875      | $2 + r$                                |
| 2   | 5   | 0        | 1.0        | 1.000      | $1 + r$                                |



Jaccard coefficient,  $AUC = 1.00$



degree product,  $AUC = 0.31$

## 1.4 Other link prediction techniques

There are two other major classes of link prediction algorithms, both of which use *global* information in order to derive a  $\text{score}(i, j)$ , and are often more complicated than most topological predictors.

**Model-based predictors** are a broad class of algorithms that rely on models of large-scale network structure, e.g., by decomposing the network into modules or communities, to make predictions about missing links. The predictions are produced by one of two strategies: likelihood or optimization.



Likelihood-based predictors are basically structured random-graph models, which estimate a probabilistic model of graphs  $\Pr(G \mid \theta)$ , much like the Erdős-Rényi random graph and the configuration model do, but with more elaborate structure. The advantage of these methods is that by learning the model from data, they also learn a model of the existence of individual edges  $\Pr(i \rightarrow j \mid \theta)$ , which we can exploit as a score function for the unconnected pairs. If the model fits the data well, then we can predict that an unconnected pair should be connected if it has a high probability of connection under the model. A very popular approach within this family is the stochastic block model, or one of its many variants.

Optimization-based predictors are based on algorithms that maximize some measure of the quality of the large-scale decomposition of the network. We can convert this objective function into a score function for unconnected pairs by asking how much adding that edge would improve its quality function.

**Embedding-based predictors** are another broad class of algorithms, which are based on the idea that close proximity of an unconnected pair within an “embedded space” is indicative of a missing link. To create this embedded space, the algorithms first assign positions to nodes within some  $\mathbb{R}^d$  space, for  $d$ -dimensions, via some graph embedding algorithm. The way these positions are assigned determines the particular assumptions about structure, and there are many different approaches, e.g., DeepWalk or node2vec; most attempt to preserve “local” structure, meaning that pairs of nodes that are connected in  $G$  have small proximities in the embedded space. Given such an “embedding,” we can then simply sort the unconnected pairs by some distance measure (there are many) within the space.

Both of these families can produce good predictions of missing links, but something called the No Free Lunch Theorem<sup>7</sup> basically guarantees that no one algorithm can be best across all inputs, and some recent work suggests that model-based and topological predictors are good general predictors, especially when used in combination. But, performance also varies across the domain of the network: social networks are generally the easiest, with most predictors performing well, while biological networks are among the hardest.<sup>8</sup>

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<sup>7</sup>See Wolpert and Macready, “No Free Lunch Theorems for optimization.” *IEEE Transactions on Evolutionary Computation* **1**(1), 67–82 (1997)

<sup>8</sup>See Ghasemian et al., “Stacking Models for Nearly Optimal Link Prediction in Complex Networks.” *Proc. Natl. Acad. Sci. USA* **117**(38), 23393–23400 (2020).