

Notes for class of 18.8

Practical Vernam cipher: (Stream cipher): A is a field
Algorithm: 1. Exchanging key K

2. Choose a random IV

3. Initial state $x(0) = (K, IV)$

Generate sequence of states $x(k+1) = F(x(k))$
4. Generate output stream

$w(k) = f(x(k))$ in A

5. Encrypt: $c(k) = p(k) + w(k)$

Decryption by same algorithm: $p(k) = c(k) - w(k)$

As a TDWF: $\{c(k)\} = E(k, \{p(k)\})$

1. OWF from $\{p(k)\}$ to $\{c(k)\}$ given k .
2. Easy to invert from $\{c(k)\}$ to $\{p(k)\}$ given k .
3. Given both $\{p(k)\}$ and $\{c(k)\}$ (att. given $w(k)$) difficult to find k .

- o Since Encryption is like a Vernam cipher / OTP property 1 is assured by randomness of $\{w(k)\}$ for any IV.

< Output stream should be indistinguishable from a random stream of alphabets >

- o Since the algorithm to generate output stream from IV is deterministic the output sequence is pseudo-random (PR) stream cipher is thus a PRG.

