

# Problem Set 10

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1. Construct—with proof—the *explicit* functions requested below.

(a) A bijection from  $\{x \in \mathbb{R} \mid -1 < x < 1\}$  to  $\{x \in \mathbb{R} \mid -\pi < x < \pi\}$ .

*Proof.* Let  $A := \{x \in \mathbb{R} \mid -1 < x < 1\}$  and let  $B := \{x \in \mathbb{R} \mid -\pi < x < \pi\}$ . Let  $f : A \rightarrow B$  where  $f(a) = \pi \cdot a$ .

- Injective

Let  $a_1, a_2 \in A$  and assume  $f(a_1) = f(a_2)$ . By definition of  $f$ , we get  $\pi \cdot a_1 = \pi \cdot a_2$ . Since  $\pi \neq 0$ , we know by multiplicative cancellation  $a_1 = a_2$ . Thus,  $f$  is injective.

- Surjective

Let  $b \in B$ . Let  $x := \frac{b}{\pi}$ . Observe that  $x > -1$  because  $b > -\pi \implies \frac{b}{\pi} > \frac{-\pi}{\pi} = -1$ . Similarly, observe that  $x < 1$  because  $b < \pi \implies \frac{b}{\pi} < \frac{\pi}{\pi} = 1$ . Since  $x \in \mathbb{R} \wedge -1 < x < 1$ ,  $x \in A$  so  $f(x) = b$ . Therefore,  $f$  is surjective.

As  $f$  is surjective and injective, it is bijective.

Q.E.D.

(b) A surjection from  $\mathbb{N}$  to  $\{p \mid p \text{ is prime}\}$ .

*Proof.* Let  $P := \{p \mid p \text{ is prime}\}$ . Let  $f : \mathbb{N} \rightarrow P$  where

$$f(n) = \begin{cases} n & \text{when } n \text{ is prime} \\ 2 & \text{otherwise} \end{cases}$$

Let  $p \in P$ . Let  $x = p$ . Observe that  $x \in \mathbb{N}$  and  $x$  is prime so  $f(x) = p$ . Thus,  $f$  is a surjection.

Q.E.D.

(c) An injection from  $X$  to  $\mathbb{P}(X)$  for *every* set  $X$ .

*Proof.* Let  $X$  be an arbitrary set. Let  $f : X \rightarrow \mathbb{P}(X)$  where  $f(x) = \{x\}$ . Let  $x_1, x_2 \in X$  and assume  $f(x_1) = f(x_2)$ . Thus, we know that  $\{x_1\} = \{x_2\}$ . By extensionality, we know that  $x_1 \in \{x_1\} \iff x_1 \in \{x_2\}$ . Since  $x_2$  is the only element in  $\{x_2\}$ , we know  $x_1 = x_2$ . Thus,  $f$  is injective.

Q.E.D.

(d) A surjection from  $\mathbb{P}(X)$  to  $X$  for every set  $X \neq \emptyset$ .

*Proof.* Let  $X$  be an arbitrary set and let  $x_0 \in X$ . Let  $f : \mathbb{P}(X) \rightarrow X$  where

$$f(y) = \begin{cases} \cup y & \text{if } |y| = 1 \\ x_0 & \text{otherwise} \end{cases}$$

Let  $x \in X$ . Let  $y = \{x\}$ . Observe that  $y \subseteq X$  and  $|y| = 1$  so  $f(y) = \cup y = x$ . Thus,  $f$  is surjective. Q.E.D.

2. Let  $A$  be an arbitrary finite set of cardinality  $|A| = n$ , where  $n \in \mathbb{N}$ . How many finite strings over  $A$  are there?

*Proof.* Let  $k \in \mathbb{N}$  and let  $s : k \rightarrow A$ . Observe that there are  $k^{|A|}$  distinct strings  $s$ . Let  $S_k := \{s | s : k \rightarrow A\}$  so as shown above  $|S_k| = k^{|A|}$ . Also let  $S := \{S_k | k \in \mathbb{N}\}$ . By theorem 8.8, since  $|S| \leq |\mathbb{N}|$  and  $(\forall k \in \mathbb{N})(|S_k| \leq |\mathbb{N}|)$ , we know that  $|\cup S| \leq |\mathbb{N}|$ .

Towards a contradiction, assume that  $|\cup S| = n$  for some  $n \in \mathbb{N}$ . So there must be  $f : n \rightarrow \cup S$  such that  $f$  is surjective. Let  $a \in A$  and let  $h := f(0) \frown f(1) \frown \dots \frown f(n-1) \frown a$ . Observe  $|h| = (\sum_{i=0}^{n-1} |f(i)|) + 1$ . Thus we know that  $(\forall k \in n)(|f(k)| \leq \sum_{i=0}^{n-1} |f(i)| < (\sum_{i=0}^{n-1} |f(i)|) + 1 = |h|)$ . Since  $(\forall k \in n)(|f(k)| < |h|)$  we know that  $(\forall k \in n)(f(k) \neq h)$ . Thus we have achieved a contradiction as  $h$  is a finite string over  $A$  (aka  $h \in \cup S$ ) and  $f$  does not map to  $h$  so  $f$  can not be surjective.

Consequently,  $\cup S \neq n$  for any  $n \in \mathbb{N}$  so  $|\cup S| \geq |\mathbb{N}|$ .

From  $|\cup S| \leq |\mathbb{N}|$  and  $|\cup S| \geq |\mathbb{N}|$ , we know  $|\cup S| = |\mathbb{N}| = \aleph_0$ .

There are  $\aleph_0$  many finite strings over  $A$ .

Q.E.D.

3. The Library

*Proof.* Let  $S :=$  the set of all books in the library. Let  $\beta :=$  the set of all sentences. Observe that there are a countable infinite amount of sentences so  $|\beta| = |\mathbb{N}|$ . Also observe that since  $\beta$  is numbered, we can use  $\mathbb{N}$  to represent the sentences in  $\beta$ . Now, let  $i : \mathbb{N} \rightarrow \mathbb{N}$  (a string of infinite length over alphabet  $\mathbb{N}$  which represents the sentences in  $\beta$ ) represent a book of infinite length. Let  $I := \{i | i : \mathbb{N} \rightarrow \mathbb{N}\}$  (the set of all infinite books).

Towards a contradiction, assume that  $|\mathbb{N}| \geq |I|$ . This means that there exists an  $f : \mathbb{N} \rightarrow I$  where  $f$  is surjective. Let us define  $g : \mathbb{N} \rightarrow \mathbb{N}$  where:

$$g(i) = \begin{cases} 1 & f(i)(i) \neq 1 \\ 2 & f(i)(i) = 1 \end{cases}$$

Notice that since  $f$  is surjective so for some value  $t \in \mathbb{N}$ ,  $f(t) = g$ . As such,  $(\forall i \in \mathbb{N})(f(t)(i) = g(i))$  so  $f(t)(t) = g(t)$ . Note that there are two cases:  $f(t)(t) = 1$  or  $f(t)(t) \neq 1$

$f(t)(t) = 1$  implies that  $g(t) = 2 \neq 1$  so  $f(t)(t) \neq g(t)$ .

$f(t)(t) \neq 1$  implies that  $g(t) = 1$  so  $f(t)(t) \neq g(t)$ .

In both cases, we achieve a contradiction and since  $f$  was arbitrary, we know that there are no injective functions from  $\mathbb{N} \rightarrow I$ . As such,  $|\mathbb{N}| < |I|$ . Notice that  $I \subseteq S$ . Consequently,  $|\mathbb{N}| < |S|$  as well so we know that  $S$  is uncountable. Therefore, as there are an uncountable amount of books, I will not be set free. Q.E.D.