## **ASSIGNMENT 3 REPORT**

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```
anguy102@ocelot:~/cop4610-os/assignment3 120% ./main
I'm thread1, I did 2000000 updates and I got the bonus for 19441 times, counter = 2080293
I'm thread2, I did 2000000 updates, counter = 4000000
from parent counter = 4000000
anguy102@ocelot:~/cop4610-os/assignment3 121%
```

Figure 1: Screenshot

## Task 1

Done.

## Task 2

1. Section annotations:

```
void *thread1() {
  int bonus_ok = 0;
  int bonus_times = 0;
  int updates = 0;
  // Remainder section
 while (updates < MAX_UPDATES) {</pre>
    // Entry section
    if (pthread_mutex_lock(&mutex) == 0) {
      // Critical section
      bonus_ok = MAX_UPDATES - updates >= 100 && counter->value % 100 == 0;
      if (bonus_ok) {
        counter->value += 100;
      } else {
        counter->value += 1;
      // Exit section
      pthread_mutex_unlock(&mutex);
    // Remainder section
    if (bonus_ok) {
      bonus_times += 1;
      updates += 100;
    } else
```

```
updates += 1;
  // Remainder section
 printf("I'm thread1, I did %d updates and I got the bonus for %d times, "
         "counter = %d\n",
         updates, bonus_times, counter->value);
 return NULL;
}
void *thread2() {
  int updates = 0;
  // Remainder section
 while (updates < MAX_UPDATES) {</pre>
    // Entry section
    if (pthread_mutex_lock(&mutex) == 0) {
      // Critical section
      counter->value += 1;
      // Exit section
      pthread_mutex_unlock(&mutex);
    // Remainder section
    updates += 1;
 }
  // Remainder section
 printf("I'm thread2, I did %d updates, counter = %d\n", updates,
         counter->value);
  return NULL;
}
```

- 2. Critical section problem analysis:
- Mutual Exclusion: Met, through the use of pthread\_mutex\_lock(&mutex) and pthread\_mutex\_unlock(&mutex). When a thread enters the critical section to modify counter->value, the mutex ensures that no other thread can simultaneously access the critical section.
- **Progress**: Met, the implementation guarantees progress because when the mutex is available, any thread attempting to enter the critical section can do so. Threads are not indefinitely blocked outside the critical section as long as they seek to acquire the lock once it's released.
- Bounded Waiting: Met, benefiting from the operating system's (Linux, in this case) effort to manage threads fairly. This includes mechanisms that indirectly support the bounded waiting condition. Linux scheduling

and pthreads ensure fairness and prevent starvation.