

Evolving Dynamic Traffic Assignments

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Foreword

May this thesis be the high note with which this challenging but rewarding year at KU Leuven comes to a close.

I would like to express my gratitude to my daily advisors Rutger Claes and Rinde van Lon, for their support and clarity of ideas. My appreciation too, to Willem Himpe who was generous with his time and explained to me the most important traffic concepts and the details of the original simulator.

Finally, I would like to dedicate this work to my family in Spain, always supportive despite the distance and to Judith, my partner in crime, the one with whom I ride the risks of life.

Andrés Arribas

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π the number pi

∞ infinity …

Chapter 1: Introduction

This chapter gives an introduction to the work. The objective is stated and an explanation is given of how it is to be achieved (better known as the theme). If you are not sure what a master’s thesis is, you can always look it up on Wikipedia[2].

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Figure 1 castle

Chapter 2: Traffic Theory[[1]](#footnote-1)

In this chapter basics of traffic theory are explained. The objective is to serve as an introduction to the most important concepts directly used in this thesis. This chapter is intended at readers with little or no background in Traffic theory. Out of scope remains to give a detailed picture of the field.

2.1 Traffic Systems Modelling

Transportation systems can be studied from a variety of angles e.g. economic, social, purely mathematical, etc., and at different scales e.g. regional, local or state wise.

In this thesis we are only concerned with a simplified view of what a transportation system is: we will focus on simulating simple small traffic networks and measuring their efficiency in terms of performance (e.g. travel times) given predefined fictional travel demands.

In order to simulate a traffic system, the first step is to model it.

A traffic model is a mathematical representation of the physical and organisational elements comprising a traffic network, as well as the travel demand and the emerging interactions.

Even the simplest traffic models have to take into consideration the inherent complexity of transportation systems, which arises from the multiplicity of non-linear interactions and feedback cycles.

In order to do so, a traffic model is most commonly defined by means of a demand model plus a combination of a traffic flow model and network model. The former is used to analyse and simulate the performances of the main supply elements, the latter to represent the topological and functional structure of a system.

The subset of traffic flow theory used in this work is a very specific one that requires careful consideration. But first, let us briefly summarize the demand representation to be used and the fundamentals of network/supply models.

2.2 Demand Model

The number of users consuming particular services offered by a transportation system in a given time period is the travel-demand flow. Travel demand flows, result from the aggregation of individual trips made in the area of study during the period of reference. A trip is defined as the act of moving from an origin to a destination.

The spatial characterization of trips is made by grouping them by place of origin and destination; accordingly demand flows can be arranged in tables commonly referred to as O-D matrices, whose rows correspond to the different origin and destination zones, respectively (SEE FIG AND INSERT FIG P 16).

A cell [ins dod] in the matrix gives the number of trips in the reference period from origin o to destination d.

The accumulated number of trips with origin o is therefore:

[FORMULAS IN P15 MAKE IT CONSISTENT WITH THE REST OF THE THESIS DOC]

And the accumulated number of trips attracted by destination d is:

[FORMULAS IN P15 MAKE IT CONSISTENT WITH THE REST OF THE THESIS DOC]

2.2 Supply Model

In this section we look at the foundations of congested network models.

2.2.1 Network Structure

A network structure is represented by a directed graph. A directed graph consists of a set of nodes, N, and a set of connections between pairs of nodes, L, called links, such that L SYMBOL NxN. In a directed graph links are oriented.

A link does not necessarily correspond to a construction in the physical world. Links rather represent phases and/or activities of possible trips between different traffic zones. The core idea behind a link is that its physical and functional characteristics can be assumed to be homogeneous for the whole link. In this sense, links can be seen as the partition of trips into segments, each of which has certain characteristics.

Nodes correspond to significant events delimiting trip phases. A node can represent the same event occurring at different time instants (between two trip phases). For example, the different entry or exit times in a road segment.

A trip is a sequence of several phases/links called a path. A path is defined from an initial node, the origin, to a final node, the destination. Each path is unambiguously associated with one, and only one O-D pair, whereas several paths can connect the same O-D pair.

[ref and pic p 47]

2.2.1 Flows

An instantaneous link flow is the instantaneous number of units using a link (i.e. units in that phase of the trip at that time instant).

PB with the previous and time [AAN]

[REV POSSIBLE FORMULAS]

Typically, when computing the link flow, the difference between types of units has to be taken into account, however, in this work, we consider all units to be of the same homogenous type.

A path flow is obtained as the sum of the link flows for all links in the path, considering in the flow calculations only those units matching the origin and destination of the path.

2.2.1 Performance Variables and Transportation Costs

Some variables perceived by the users can be associated to individual trip phases (hence links). Examples of such variables are travel times, monetary cost, and discomfort. These variables are referred as level-of-service or performance attributes. In general, performance variables correspond to disutilities or costs for the users.

Given our limited scope and simplified approach we focus on travel times as the only performance variable. Accordingly, from now onwards, “travel time”, “cost” and “performance variable/indicator” may be used indistinctly.

Remark that a cost will reflect not only the assigned physical characteristics of the link at hand (e.g. length, free speed), but also the interaction of the units in the link. In particular, the phenomenon of congestion, which we define precisely in the next section, will have a significant impact on the cost.

Similarly to the flows, a path cost can be computed as the sum of the cost of all links in the path.

2.2 Traffic Flow Model

Usually traffic flow models can be split in stationary models, where there is no variance in time of variables, and dynamic models, where there is variance in time.

Stationarity could only be observed in the real work if demand, path choices and supply system remained constant for a very long period of time. In less than a day of time, this is not the case and the behaviour of the traffic system does not only depend on its characteristics but also on its history.

Hence, we are interested in dynamic models, that, contrary to stationary models, allow the effective, realistic, simulation of important phenomena, especially the creation, propagation and dissipation of queues.

There are several approaches to dynamic traffic modelling (also called “within-day dynamics) depending on the scale at which units are considered e.g. continuous fluid analogy, packet based, or individual units. We are concerned with so called “Macroscopic” models.

Macroscopic models are based on the analogy with fluid dynamics. In such a models the individual vehicles are treated as a continuous[[2]](#footnote-2) (one-dimensional) fluid for which variables such as flow, density and velocity are defined at each point in space and time. Macroscopic models comprise equations for the conservation of mass (vehicles) and the relation between flow and density.

More specifically, we are interested in “Space-discrete” Macroscopic models for which it is assumed that basic variables affecting link performance, such as density or speed, do not vary along the link.

A fundamental aspect to take into account given dynamicity, is that the number of users within a link depends now on the travel times required to reach that link, which, in turn, depends on the number of users encountered on the network links in the previous instants. This circular dependency between flow and link performance, illustrated in fig [FIG p 435 Cascetta], does not allow the resolution of the two models (flow propagation an link performance) in a sequential fashion. This problem is known as Dynamic Network Loading.

2.2.1 Dynamic Network Loading

The DNL model is responsible for propagating path demands through the network. The traffic flows are calculated in time and space, congestion may form, delays may be encountered and the travel times on the network are determined.

The DNL we use is the Link Transmission Model (LTM)

Whereas a mastery of the LTM or even in depth knowledge of it are not a requirement for this Master Thesis it is desirable to look at a brief introduction to the LTM which will also serve as a means of covering some key aspects of traffic flow theory worth keeping in mind when assessing the following sections and the final results.

2.2.1 Link Transmission Model

The Link Transmission Model is a DNL model proposed by Yperman at KU Leuven in 2007 [BIB].

Amongst its most significant characteristics we may count a high realism of representation of congestion formation (whenever the traffic in a link surpasses its free flow capacity) and spillback (backwards propagation of the effect of congestion) as well as computational efficiency.

We will now discuss firstly, the concept of Cumulative Vehicle Numbers, which forms the basis of the calculations in the LTM, and secondly, the involved basics of traffic propagation.

Many details of traffic theory of no interest to this work have been omitted.

2.2.1 Cumulative Vehicle Numbers

The aggregation of vehicles having passed a location x by time t is the Cumulative Vehicle Number, or CVN, denoted by *N(x,t).*

In the LTM units are assumed to follow a predefined route between origin and destination creating different streams of traffic given different origins and destinations.

For a particular link, CVN in the LTM are calculated at all upstream and downstream link boundaries. Through time the CVN, given that the locations of measurement are fixed, will draw a shape through time. Additionally, CVN curves are also computed for complete streams at origin and destination.

Once all the curves have being obtained different variables can be derived:

* Link travel times can be derived as the horizontal distance between the curve at the upstream boundary N(x0,t), and the curve at the downstream boundary N(xL,t). Analogously, total travel time for a path is the horizontal difference between CVN at origin and CVN ay destination.
* The total number of vehicles in a link is the vertical distance between upstream and downstream CVN curves.
* The flow q (veh/h) is simply the slope of the CVN curve.
* The density k (veh/km) on a link can easily be calculated by dividing the number of vehicles N by the length L.

2.2.1 Traffic Propagation

In the LTM, the “fundamental triangular diagram” is considered an accurate representation of the propagation of traffic in unidirectional links. The diagram itself is just a plot of the flow in relation to the density. What is fundamental is the set of assumptions and definitions behind the diagram.

Below the critical density kc, which corresponds to the link’s capacity C, vehicles are assumed to travel with a fixed free speed vf (km/h) characteristic of the link.

Above the critical density kc , hence for congested traffic, the speed in the link v is given by q/k.

The maximum density is defined as the jam density, kjam, and corresponds to vehicles standing still.

Different traffic states are separated by shock waves that may propagate up- or downstream with a speed ws, that can be derived as the slope of the connecting line between the adjoining traffic states.

The maximum negative shockwave speed w is the fastest possible speed with which congestion may spill back in the upstream direction.

[intro in the text references –as in the figure-]

[INSERT FIG OF FUNDAMENTAL DIAGRAM]

In this thesis it suffices to know that all the above can be computed for a particular time step from CVN curves (including curves produced in previous time steps).

Finally, note that the above is a description of the propagation of traffic in a link.

The LTM may use different models to describe intersections i.e. propagation of incoming flows to outgoing links. These are out of scope; suffice to say, nevertheless, that once the flows are computed the CVN per link can be easily updated.

2.2.1 Dynamic Traffic Assignment

But, does DNL account for a realistic simulation of traffic? No, it does not on its own. DNL serves as a means of computing the model response for an a-priori assignment of traffic. That is, given that the path each vehicle will follow is known, we may use DNL to compute flow propagation and link performance. From the DNL point of view, it is not relevant whether the assignment reflects or does not reflect appropriate route choices.

It is the task of so called Dynamic Traffic Assignment (DTA) Models to, in addition, consider assignment accuracy and henceforth reproduce more adequately system dynamics within the period of simulation.

2.2.1 Introduction

Dynamic Traffic Assignments are used to study user route choice and propagation of traffic in time, including the reaction to changes and the impact on other users of congestion and driver’s decisions.

DTA models are mainly made of four sub-models, the demand model, the supply model, the DNL and the route choice model. We have already seen the first three; we will focus on the route choic.

Route choice, also called traffic assignment and traffic swapping, determines the division of traffic to alternative routes for every origin-destination pair, by entry time to the network. The final aim of the route choice is to attain user equilibrium, commonly called Dynamic User Equilibrium (DUE), stating that the cost of any route being used is not higher than the travel cost of any alternative route. In the dynamic case this condition must hold for all times. Analogously, no user can improve his situation (in terms of cost) by transferring to an alternate route, under the assumption that the users are rational and will chose the shortest route to their destination.

The general formulation of the DUE requires consideration of continuous time [INSERT T]. However, we are interested in practical applications, which require discretization of time into intervals of the form [INSERT FORM].

The cost or travel time is designated by [INSERT]. The instantaneous route demand (veh/h) at a time [insert t] at the entry of route [insert r] is denoted by [drt]. The total instantaneous demand assigned to any alternative route from origin [p] to destination [q] at time [t] is denoted by [Tpqt], hence:

[insert formula]

where [zpq] is the set of alternative routes between origin [p] and destination [q]. The general DUE condition in continuous time is:

[ins formula]

where [C\*] indicates the minimal travel cost from origin [p] to destination [q] a time t. That is to say, at every point in time the travel cost for each route in use [crt], is equal to the minimal cost of the origin-destination pair, and not higher than the travel cost of any alternative route.

In the discrete version we assume that the demands are constant within each interval, so … [copy phrase from p3 article].Therefore,

[insert updated formula]

Given that travel times are not really constant a key dilemma in the formulation of the DUE is which travel time should be associated with an interval of departure times. The most computationally stable option has been proven to associate with each interval the travel time at the end of the interval i.e [intro c= c. ].

In addition, route choice is often modelled as being the product of user experience. User experience is limited by the users perception of the traffic situation. In order to model the sensitivity of the user to traffic changes, the route choice algorithm may consider the users impervious to changes to the traffic cost for a period of time [def of t prime] This leads us to an alternative formulation of the DUE condition:

[insert my versoin of eq 4.]

The excess cost of a single route, denoted by λr(τ), is defined as the difference between the travel cost on route r and the minimal travel cost from origin p to destination q. So that,

ECr(τ)=λr(τ)=Cr(τ)−C∗pq(τ), ∀r∈Zpq,∀pq,∀τ.

(5) Total excess cost of traffic travelling on the route r, denoted by ρr(τ), is

TECr(τ) = ρr(τ) = dr(τ)·ECr(τ), ∀r ∈ Zpq,∀pq,∀τ. (6)

In perfect equilibrium the TEC of all routes should be zero. Convergence is mea- sured by the Global Excess Cost (GEC), which is the summation of the total excess travel costs for all departure times, in all the alternative routes and for all origin- destination pairs.

GEC=∑ ∑ ∑TECr(τ).

Route Choice Algorithms

The route choice algorithms we are concerned with follow an iterative process in which the assigned demand is updated between iterations according to a correction factor.

This can described as follows:

[insert func]

where F is a correction function suggesting a change to the demand flow matrix [d] containing the demands for all routes and all departure intervals. The suggested change is dependent on the demand flow matrix [dj-1] of the previous iteration. The solution [dj] for iteration j is the weighted average of the proposed solution [Fdj] in iteration j and of the solution [dj-1] attained in iteration [j-1] , where [aj] represents the step size in iteration j.

Analogously, one can directly consider the correction function that expresses the proposed change matrix [lj] in iteration j. This correction matches the difference between the proposed solution in iteration j and the actual solution in iteration j-1.

[insert func]

In traffic dynamics applying a change to a demand assignment for a route may affect, not only traffic departing at later times on the same route, but also already departed traffic on different routes. As a consequence, attaining the DUE requires the iterative process to compute a full DNL solution per iteration.

Notice that as a requisite to the usage of the DNL lays the need to specify the paths in a network. For small networks all possible paths may be enumerated a-priori, but with bigger networks and more realistic scenarios this becomes impractical very fast. Henceforth a pathfinder component is required.

Path finding is very costly computationally and as such is often applied incrementally. When computing inputs to the DNL for the first time, potential shortest routes are computed by the pathfinder based on costs a priori (assuming no congestion). Once a DNL solution has been obtained the new shortest routes according to the new costs are computed (and the GEC is calculated), if these differ from the existing ones they are added to the pool of potential paths.

The route choice algorithm uses the complete pool of paths in order to assign the demand to different routes with the DUE as an objective.

Given comparable DNL models and pathfinders, the route choice algorithm is the component that influences the most the speed of convergence of the DTA. Finding an efficient route choice algorithm is therefore a key challenge. We will know look at two of the existing simple methods.

[FIG TO BE ADDED]

Fig [] depicts the complete iterative process.

MSA

MSA, or method of successive averages, is a very popular route choice numerical method applied to the DUE problem.

MSA works by proposing an all or nothing assignment then…

DEC

Conclusion

Genetic Programming

A central idea of genetic algorithms is that they are not based on the traditional correct, consistent, justifiable, certain, orderly, parsimonious, and decisive principles of the vast majority of techniques in machine learning. Rather than seeking an exact mathematical reasoning that leads us to the solution, genetic algorithms are based on entirely new principles, namely on applying the same kind of uncertain and non-deterministic evolutionary process observed in nature, where diversity is often wanted and therefore inconsistency and contradiction are considered of significant value.

Genetic programming (GP) is in fact a type of genetic algorithm (GA), which focuses on the evolution of computer programs. In the subsequent sections we will have a look at the basic of what a genetic algorithm is, then, we will re-interpret the genetic algorithm paradigm under the light of genetic programming, and finally we will consider some of the details and practicalities of genetic programming itself.

Genetic Algorithms

In nature, given a population of self-reproducing entities, diversity amongst them, and a connection between the entities diversity and their survival ability, we observe a tendency for the fittest individuals to reproduce more frequently, as well as a higher rate of survival of their genetic information. Through time this process reshapes the structure of the population towards a population dominated by fitter individuals. We say that the population has then evolved.

John Holland proposed in his book “Adaptation in Natural and Artificial Systems” to extend to artificial systems this sort of evolutionary process ruled by natural selection. The kind of algorithms that solve adaptation problems formulated in genetic terms is what we call a “genetic algorithm”.

A genetic algorithm is typically a highly parallel algorithm, which will transform a population of mathematical objects that we call individuals, each with an associated fitness value, into a new population using operations based on the principles of survival of the fitness and sexual reproduction.

Briefly put and in plain terms a typical genetic algorithm proceeds as follows:

It begins with an effort to learn something about the environment by testing a number of M randomly chosen points in the search space, equivalent to generating M individuals randomly. This is generation 0.

Per generation the fitness of each individual is assessed against the environment. Depending on the fitness of individuals and a probabilistic parameter fixed a-priori [pr], a pool of individuals is created by reproducing (copying) individuals based on a criterion that favours the fittest ones. Individuals in this pool are then recombined into new individuals (mixing their parts to produce 2 individuals), this operation, which also favours fittest individuals and uses a new parameter [pc], is called crossover and it allows exploring new points in the search space by replacing parents with their fitter offspring. An additional operation of mutation may uniformly computed per individual with parameter [pm].

The iterative repetition of the previous steps has the potential to improve the average fitness of the population. The best individual encountered, often labelled best-so-far is our solution. Once the optimal solution has been reached or a maximum number of generations has been reached we have attained the termination criterion and the evolutionary process can be stopped.

More formally:

1) The preparation of a genetic algorithm requires:

The selection of a representation scheme

A fitness measure

The parameters and variables controlling the algorithm

A way of designating the result

A termination criterion

The representation scheme is a mapping that allows expressing each possible point in the search space. A representation that facilitates the search for a solution by means of the genetic algorithm requires significant insight into the problem domain, for example, in the travelling salesman problem, or TSP, a possible solution or path may be represented by enumeration of all nodes: 1-3-4-6-5.

The fitness measure must ensure the evaluation of all possible individuals. It is often inherent to the problem, for example, in the TSP the fitness may be the length of the path.

The primary parameters for controlling a genetic algorithm are the population size (M) and the maximum number of generations to be run (G). Secondary parameters are [pr], [pc] and [pm]. Other control variables are often required depending on the specifics of the GA.

The most common way of designating the result is obviously to take the best individual so far.

A run may be terminated according to different criteria. The most frequent criteria are encountering the best possible result and reaching the maximum number of generations.

2) Once the preparation is complete the genetic algorithm can be run as follows:

Randomly crate an initial population of individuals.

Iteratively perform the following substeps on the population until the termination criterion has been satisfied

Evaluate the fitness of each individual in the population

Create a new population by applying at least the first two of the following operations, on individuals chose with a probability based on the fitness:

Copy existing individuals to the new population

Create two new individuals by genetically recombining two existing individuals

Create a new individual by randomly mutating an existing individual

The best-so-far individual is then designated as the result of the genetic algorithm for the run.

This process is illustrated in figure [ins num]

[INSERT FLOWCHART]

There are numerous variations on the basic genetic algorithm but this is the most frequent setting. It is important to understand that domain specific knowledge may help substantially into taking the best choices during the preparation step, which will in turn improve the performance of the genetic algorithm. However, the main point is that the genetic algorithm is a domain-independent approach to rapidly searching an unknown search space for good solutions.

For further explanations, particularly regarding why genetic algorithms work, the reader is referred to [insert Koza book and chapter 3.2].

Overview of Genetic Programming

Genetic programming arises from the observation that numerous problems in machine learning, artificial intelligence and symbolic processing can be interpreted as requiring the discovery of a computer program that provides a certain output given a certain input.

The discovery of such a computer program can be reinterpreted as the search for the fittest computer program in a space of computer programs. As we have seen such a search can be achieved through the principles of survival of the fittest, reproduction and crossover.

The leap between basic genetic algorithms and genetic programming lies at the representation level. The type of structure undergoing adaptation is now a hierarchy of functions and terminals of varying size and shape. In addition, such computer programs are now “active” in the sense that they are capable of being executed.

Genetic programming is therefore a domain-independent method that provides a single unified approach to the problem of finding a computer program to solve a problem.

The execution of a genetic algorithm under the genetic programming paradigm becomes:

Randomly crate an initial population of random compositions of the functions and terminals of the problem (computer programs)

Iteratively perform the following substeps on the population until the termination criterion has been satisfied

Execute each program in the population and assign it a fitness value according to how well it solves the problem.

Create a new population by applying at least the first two of the following operations, on computer programs in the population chosen with a probability based on their fitness:

Copy existing programs to the new population

Create new computer programs by genetically recombining two existing individuals

Create a new individual by randomly mutating an existing individual

The best-so-far computer program is then designated as the result of the genetic algorithm for the run.

Genetic Programming Details

As in Koza’s book we will describe in more detail genetic programming focusing on John Holland’s key features common to tall adaptive systems, one at a time.

These key features are:

Structures that undergo adaptation

Initial structures

Fitness measure

Operations that modify the structures

State of the system at each stage

Termination method

Method to designate the result

Parameters that control de process

The following explanations cover all that is needed in order to understand this thesis, not more.

Structures that undergo adaptation

As already introduced the structures that undergo adaptation are computer programs. More precisely hierarchical structures made of functions and terminals.

The functions may be:

Arithmetic operations, e.g. +,-,\*

Mathematical functions, e.g. sin, exp

Boolean operations, e.g. and, not

Conditional operators, e.g. if, then

Functions causing iteration, e.g. loop

Functions causing recursion and

Domain-specific functions that may be defined

The terminals are atomic entities, i.e., variables, constants or functions taking no arguments.

A common representation of these structures is the s-expression used in the LISP language, for example (OR (AND (NOT d0) (NOT d1) ) (AND d0 d1)) which corresponds to the tree diagram in figure [INS FIG REF].

[INSERT FIGURE P 81]

Initial structures

The initial structure in genetic programming consists of the individuals in the initial population of individual computer programs for the problem.

The generation of each initial computer program is done by randomly generating a tree structure made of functions and terminals.

In practice there are three dominant techniques to do this: full method, grow method and the ramped half-and-half.

The full method involves generating trees with a fixed length, in which all branches from root to the endpoint have same length. On the other hand, the grow method produces trees that are variably shaped, but not deeper than a maximum length.

Ramped half-and-half is a method in which the population is split in two parts, often in half but not necessarily. The full method is then applied to one of the parts, whereas the grow method is applied to the other part.

Fitness measure:

The most common approach is to define an explicit measure of fitness for each individual. This is accomplished via an explicit well-defined evaluation procedure producing a scalar fitness value as output.

We are interested in 3 interpretations of the fitness value:

Raw fitness: r, the fitness as normally measured in the problem domain.

Standardized fitness: raw fitness manipulated so as to ensure that a lower fitness coincides with a better individual. In particular if the raw fitness increases for worse individuals s = r.

Adjusted fitness: Defined as [INS EQUATION]. Adjusted fitness lies between 0 and 1. The adjusted fitness is bigger for better individuals in the population.

Operations that modify the structures:

We already know that the primary operations in genetic programming are reproduction and crossover.

Reproduction is the basic engine of survival of the fittest. It consists of two steps. First and individual is selected. Second, the individual is copied, without alteration, from the current population into the new population (i.e., the new generation).

A popular method of selection is tournament selection, in which a specified group of individuals are chosen at random from the current population and the one with the better fitness is then selected. Note that tournament selection is a selection strategy with replacement, therefore the same individual may be selected to become part of the next generation several times. Tournament selection has the particularity of adding additional control to the selection pressure by increasing or decreasing the size of the group of individuals chosen at random.

State of the system at each stage:

The state of the system at any point in time is given by the current population. Control parameters (which may be dynamic in complex implementations) and best-so-far individual are commonly kept in memory too.

Termination Method:

As we have already discussed the most common termination criterion is to attain the maximum number of generations or to encounter a result that matches a given condition. Other criteria such as the loss of diversity (stagnation) may be part of the termination criteria.

Result:

Again, as already discussed the designated result is typically the best-so-far. An alternative is to focus on the best individual in the last generation. However, there is no guarantee that the best individual in the last generation is the overall best encountered result.

Parameters that control the process:

FINISH GP TOMO – MAYBE ALREADY DO DEC AND MSA!

The Eiffel Tower has three floors:

the first;

the second;

the third.

But do this:

The Eiffel Tower has three floors: the first, the second and the third.

2.2.1 Flows

blabla

2.2.1 Performance Variables

2.3 Conclusion of this chapter

If you have reached important findings or conclusions in this chapter, it is only logical that you should end the chapter by summarising them. This is not necessary for chapters such as the introduction or list of the cited literature.

|  |  |  |
| --- | --- | --- |
|  | 1 | 2 |
| A | A1 | A2 |
| B | B1 | B2 |
| C | C1 | C2 |

Table 1 This is the first table containing data

Text

|  |  |  |
| --- | --- | --- |
|  | 1 | 2 |
| 1 | 11 | 12 |

Table 2 A second table

Chapter 3: A new chapter

A chapter contains a cohesive[[3]](#footnote-3) whole that stands, more or less, on its own. It is therefore only logical that it should start with an introduction, i.e. that part of the text which you are now reading.

3.1 First subject in this chapter

Information introducing the subject.

3.1.1 An item

Text is never presented on its own. This means that references are bound to be needed. Reference can be made to online documents[2] or books[3].

3.2 Second subject in this chapter

A chapter will contain several subjects. Let us assume that this one is the last.

3.3 Conclusion of this chapter

If you have reached important findings or conclusions in this chapter, it is only logical that you should end the chapter by summarising them. This is not necessary for chapters such as the introduction or list of the cited literature.



Figure 2 Airplane

Chapter 4: The final chapter

A chapter contains a cohesive whole of information that stands, more or less, on its own. It is therefore only logical that it should start with an introduction, i.e. that part of the text which you are now reading.

4.1 First subject in this chapter

Information introducing the subject.

4.1.1 An item

The accompanying text. Remember to keep paragraphs long enough, but make sure the sentences are not too long.

A paragraph contains a train of thought and so will always contain a couple of sentences. Do not write a paragraph which consists of only one line.



Figure 3 www etc

4.2 Second subject in this chapter

A chapter will contain several subjects. Let us assume that this one is the last.

4.3 Conclusion of this chapter

If you have reached important findings or conclusions in this chapter, it is only logical that you should end the chapter by summarising them. This is not necessary for chapters such as the introduction or list of the cited literature.

Chapter 5: Conclusion

The master’s thesis is brought to a close with a chapter summarising all the conclusions once again. This is also the place to include suggestions on further use of the results, in both industrial applications and further research.

Appendices

Appendix A: The first appendix

The appendices contain information that is likely to be useful to the reader, but not essential to a sound understanding of the argument in the normal body of text. Examples include source files, configuration information, lengthy mathematical deductions, etc.

Needless to say, an appendix may be further divided into sections, or contain figures and references[1].

Appendix B: The final appendix

The appendices contain information that is likely to be useful to the reader, but not essential to a sound understanding of the argument in the normal body of text. Examples include source files, configuration information, lengthy mathematical deductions, etc.

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Master’s Thesis file

Students: First name Surname

Title: Title of master’s thesis

English title: English translation of the title

UDC:

Content in brief: (maximum 500 words)

(Title of the article: Title of the article)

(English translation: English translation of the title)

Thesis submitted to obtain the degree of Master in Engineering: ….....................................

Supervisor(s):

Assessors:

Mentor(s):

1. This chapter is essentially based on reformulated excerpts from Transportation Systems Analysis by Ennio Cascetta [↑](#footnote-ref-1)
2. Notice that despite the continuous fluid analogy the simulation and resolution of in Macroscopic models requires the discretization of time. [↑](#footnote-ref-2)
3. Insert footnote via References 🡪 insert footnote [↑](#footnote-ref-3)