# L6 - 21/08/2024

## Support of Fnn

For sets X & Y, we can define  $y \times to be the set of all maps <math>f: X \rightarrow Y$ 

Consider  $Y = \{0,1\}$ Then,  $\{0,1\}^{\times}$  is the set of all maps  $f: X \rightarrow \{0,1\}$ 

Given such a map, we can define a subset of X as

Support of f

### Power Set

Let Y be a set.

The Power Set P(Y) is defined to be the set of all subsets of Y.

$$y = \{a, b\}$$
 $\Rightarrow P(y) = \{\emptyset, \{a\}, \{b\}, \{a, b\}\}$ 

#### Thm

There is a natural bijection  $\varphi: \{0,13^X \rightarrow P(X)\}$ 

$$\{0,1\}^{\times} \longrightarrow P(\times)$$

$$f \longmapsto S_{f}$$

Pf - grijectivity

Given an 
$$f: X \rightarrow \{0,1\}$$
,

consider an  $g: X \rightarrow \{0,1\}$  s.t

 $Sa = Sa$ 

Sg = Sf $\Rightarrow g(n) = \{1, n \in S_f : 0, otherwise$ 

$$S_f = \{ x \in X \mid f(x) = 1 \}$$

 $:: S_f = \{ x \in X \mid f(x) = 1 \}$ 

$$\Rightarrow g = f \qquad \left( \begin{array}{c} \vdots \text{ Value of } f, g \\ \text{is same } \forall x \in X \end{array} \right)$$

$$S_f = S_g \Rightarrow g = f$$

(i.e. Sf completely determines f)

 $... S_f = S_g \implies g = f$ 

... φ is injective.

Surjectivity

Let 
$$T \subset X$$
 (or equiv.  $T \in P(X)$ )

Consider  $X_T : X \to \{0,1\}$  s.t

Characteristic

$$\chi_{T}(x) = \begin{cases} 1, & x \in T \\ 0, & \text{otherwise} \end{cases}$$

$$S_{\chi_{T}} = \{ x \in X \mid \chi_{T}(x) = 1 \}$$
$$= \{ x \in X \mid x \in T \}$$

$$\therefore S_{x_{T}} = T \Rightarrow \varphi(x_{T}) = T$$

$$\therefore \forall T \in P(x), \exists x_{T} \in \{0,1\}^{X} \land t$$

$$\varphi(x_{T}) = T$$

Rel" on a set X Subset of XXX RC XXX . Equivalence Reen RC XXX s.t it is - Reflexive Yn, (n,n) ER (n,y)∈R => (y,x)ER - Symmetric  $(x,y) \in \mathbb{R}$  and  $(y,z) \in \mathbb{R}$ - Transitive  $\Rightarrow (x, z) \in \mathbb{R}$ eg- het RCZXZ Ra = { (a,6): d(a-b) de Z,0} RVSVTV >> Rd is equiv. run on Z

NOTE - (Informal) We can denote a rein R in the following manner. 9f (2,y) CR > So, if ~ is a run on X, is eq. run if it satisfies ∀n, n~n S - n~y > y~n T - n~y and y~z >> n~z

#### · Equivalence class

For an equivalence rel<sup>n</sup> R on X, we can define subsets of X. called equivalence classes as follows.

Given nex, Ec(x) c X  $EC(\alpha) = \{ y \in X : (\alpha, y) \in R \}$ 

i. L

For  $x, y \in X$ , we say  $p \in X$   $x \sim y$  if we can join XQ. Let XCR2 ne e y by a cont. path. gs ~ an eq. rel"? A. R / S / T / > Eq. rel"

This eg. run has exactly 2 eq. classes (namely, the 2 discs)

EC(P) = Disc 1EC(Q) = Disc 2

Q. Let ~ be on eq. reen on X. for x, y ex, PT ! Ec(n) n Ec(y) = Ø Of Ec(n) = Ec(y) 2. Ec(n) = Ec(y) ⇔ n~y Pf -1. Consider ZEEC(x) NEC(y) >> 2~2 and y~2 To show EC(n) = EC(y), we first show EC(n) C EC(y) I.I EC(x) G EC(y) Consider tEEC(x) = x~t ⇒ t~n (S)  $t \sim \pi$  and  $x \sim z \Rightarrow t \sim z$  (T)  $t\sim z$  and  $z\sim y \Rightarrow t\sim y$  (T) >> y~t (s)

- : te Ec(n) > te Ec(y)
- $\therefore$  EC(x)  $\subseteq$  EC(y)

Similarly, we can show that  $EC(y) \subseteq EC(x)$ 

- 2.1 EC(x) = EC(y) ⇒ n~y
  - EC(x) = EC(y)=> 3 ZEEC(x) NEC(y)
  - or na zama yaz
    - $\Rightarrow$   $x \sim z$  and  $z \sim y$  (S)
    - > x~y (T)

2.2 
$$n \sim y \Rightarrow EC(n) = EC(y)$$

Let  $z \in EC(n)$ 
 $\Rightarrow x \sim z$ 
 $\Rightarrow z \sim n$  (S)

 $z \sim n$  and  $n \sim y$ 
 $\Rightarrow z \sim y$  (T)

 $\Rightarrow y \sim z$  (S)

 $\Rightarrow z \in EC(y)$  (S)

 $\therefore EC(n) \subseteq EC(y)$ 

Similarly, we can prove that

 $EC(y) \subseteq EC(n)$ 

 $\therefore$  EC(y) = EC(n)