Zero-Knowledge Proofs - II

CS 601.642/442 Modern Cryptography

Fall 2019

Zero-Knowledge Proofs for **NP**

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- Think: How to prove the theorem?
- Construct ZK proof for every **NP** language?
- Not efficient!

Zero-Knowledge Proofs for **NP** (contd.)

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Step 1: Construct a ZK proof for an **NP**-complete language. We will consider *Graph 3-Coloring*: language of all graphs whose vertices can be colored using only three colors s.t. no two connected vertices have the same color

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- Given instance x and witness w, P and V reduce x into an instance x' of Graph 3-coloring using Cook's (deterministic) reduction

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- Given instance x and witness w, P and V reduce x into an instance x' of Graph 3-coloring using Cook's (deterministic) reduction
- P also applies the reduction to witness w to obtain witness w' for x'
- Now, P and V can run the ZK proof from Step 1 on common input x'

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- **Intuition for Soundness:** In each iteration, cheating prover is caught with probability $\frac{1}{|E|}$
- **Intuition for ZK:** In each iteration, V only sees something it knew before – two random (but different) colors

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- Need two properties from digital locked boxes:
  - **Hiding**: V should not be able to see the content inside a locked box
  - **Binding**: P should not be able to modify the content inside a box once its locked

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- Two phases:
  Commit phase: Sender locks a value v inside a box
  Open phase: Sender unlocks the box and reveals v
- Can be implemented using interactive protocols, but we will consider non-interactive case. Both commit and reveal phases will consist of single messages

Commitment Schemes: Definition

Definition (Commitment)

A randomized polynomial-time algorithm Com is called a *commitment scheme* for n-bit strings if it satisfies the following properties:

- **Binding:** For all $v_0, v_1 \in \{0,1\}^n$ and $r_0, r_1 \in \{0,1\}^n$, it holds that $\mathsf{Com}(v_0; r_0) \neq \mathsf{Com}(v_1; r_1)$

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- **Corollary:** One-bit commitment implies string commitment

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- Binding follows from construction since f is a permutation
- Hiding follows in the same manner as IND-CPA security of public-key encryption scheme constructed from trapdoor permutations

ZK Proof for Graph 3-Coloring

Common Input: $G = (V, E)$**, where** $|V| = n$

P**'s witness: Colors** $\mathsf{color}_1, \ldots, \mathsf{color}_n \in \{1, 2, 3\}$

Protocol (P, V)**:** Repeat the following procedure $n|E|$ times *using fresh randomness*

$P \to V$: P chooses a random permutation π over $\{1, 2, 3\}$. For every $i \in [n]$, it computes $C_i = \mathsf{Com}(\widetilde{\mathsf{color}}_i)$ where $\widetilde{\mathsf{color}}_i = \pi(\mathsf{color}_i)$. It sends $(C_1, \ldots, C_n)$ to V

$V \to P$: V chooses a random edge $(i, j) \in E$ and sends it to P

$P \to V$: Prover opens C_i and C_j to reveal $(\widetilde{\mathsf{color}}_i, \widetilde{\mathsf{color}}_j)$

V: If the openings of C_i, C_j are valid and $\widetilde{\mathsf{color}}_i \neq \widetilde{\mathsf{color}}_j$, then V accepts the proof. Otherwise, it rejects.

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- Then, with probability $\frac{1}{|E|}$, V chooses $i = i^*, j = j^*$ and catches P
- In $n|E|$ independent repetitions, P successfully cheats in all repetitions with probability at most

$$\left(1 - \frac{1}{|E|}\right)^{n|E|} \approx e^{-n}$$

Proving Zero Knowledge: Strategy

- Will prove that a single iteration of (P, V) is zero knowledge
- For the full protocol, use the following (read proof online):

Theorem

Sequential repetition of any ZK protocol is also ZK

- To prove that a single iteration of (P, V) is ZK, we need to do the following:
  - Construct a Simulator S for every PPT V^*
  - Prove that expected runtime of S is polynomial
  - Prove that the output distribution of S is correct (i.e., indistinguishable from real execution)
- Intuition for proving ZK for a single iteration: V only sees two random colors. Hiding property of Com guarantees that everything else remains hidden from V.

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- If $(i, j) = (i', j')$, then feed the openings of C_i, C_j to V^* and output its view. Otherwise, restart the above procedure, at most $n|E|$ times

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- If $(i, j) = (i', j')$, then feed the openings of C_i, C_j to V^* and output its view. Otherwise, restart the above procedure, at most $n|E|$ times
- If simulation has not succeeded after $n|E|$ attempts, then output `fail`

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- H_2: Simulator S

Correctness of Simulation (contd.)

- $H_0 \approx H_1$: If S' does not output `fail`, then H_0 and H_1 are identical. Since (i, j) and (i', j') are independently chosen, S' fails with probability at most:

$$\left(1 - \frac{1}{|E|}\right)^{n|E|} \approx e^{-n}$$

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- $H_1 \approx H_2$: The only difference between H_1 and H_2 is that for all $k \in [n] \setminus \{i', j'\}$, C_k is a commitment to $\pi(\mathsf{color}_k)$ in H_1 and a commitment to 1 in H_2. Then, from the multi-value hiding property of Com, it follows that $H_1 \approx H_2$

Additional Reading

- Zero-knowledge Proofs for Nuclear Disarmament [Glaser-Barak-Goldston'14]

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- Concurrent Composition of Zero-Knowledge Proofs [Dwork-Naor-Sahai'98, Richardson-Kilian'99, Kilian-Petrank'01,Prabhakaran-Rosen-Sahai'02]

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- Non-malleable Commitments and ZK Proofs [Dolev-Dwork-Naor'91]