

p_1, p_2, \dots, p_n

1 2 3 4 ... $(n-1), n \rightarrow p_1 \rightarrow 0$

1 2 3 4 ... $(n), (n-1) \rightarrow p_2 \rightarrow 1$

1 2 3 4 ... $(n), (n-2), (n-1) \rightarrow p_3 \rightarrow 2$

1 2 3 4 ... $(n), (n-1), (n-2) \rightarrow p_4 \rightarrow 3$

1 2 3 4 ... $(n), (n-3), (n-2), (n-1) \rightarrow p_5 \rightarrow 3$

$(n), (n-5), (n-4), (n-2) \rightarrow p_6 \rightarrow 4$

$(n), (n-2), (n-1), (n-3) \rightarrow p_6 \rightarrow$

$E(V) =$

$$V_i^* = \begin{cases} 1 & \text{if present after process} \\ 0 & \text{o/w} \end{cases}$$

$$\Pr(V_i^* = 1) = \frac{1}{d}$$

$$\Pr(V_i^* = 0) = 1 - \frac{1}{d}$$

$$E(V_i^*) = \frac{1}{d}$$

$E(V) =$ Expected no. of vertices after process

$$= \sum_{i=1}^n E(V_i^*) = \sum_{i=1}^n \frac{1}{d} = \frac{n}{d}$$

$n(1 - \frac{1}{d}) \rightarrow$ Expected no. of vertices deleted

A vertex was connected to d edges in expectation.

$$\sum_{i=1}^n ND_i = 2E$$

$$\sum_{i=1}^n ND_i = nd$$

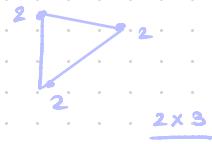
$$E(ND) = d$$

$nd \times (1 - \frac{1}{d})$ edges deleted } wrong

$$\frac{nd}{2} - nd \left(1 - \frac{1}{d}\right)$$

$$nd \left[\frac{1}{2} - 1 + \frac{1}{d}\right]$$

$$nd \left[\frac{1}{d} - \frac{1}{2}\right]$$



$$\frac{2E}{N} =$$

A edge survives if and only if two of its vertices survive.

$$\therefore \text{Expected no. of edges to survive} = \frac{nd}{2} \cdot \left(\frac{1}{d}\right)^2 \\ = \frac{n}{2d}$$

Algorithm to get Independent set

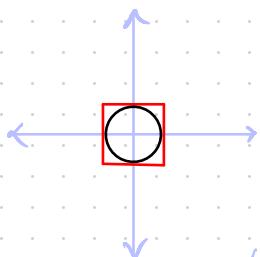
let m_p be the no. of edges after



Theorem: Let $G = (V, E)$ be a graph on n vertices with $\frac{nd}{2}$ edges. Then G has an independent set with at least $\frac{n}{2d}$ vertices.

$$\frac{nd}{2} \geq \frac{n}{2}$$

$$\frac{n^2}{4 \cdot \frac{nd}{2}} = \frac{n^2}{4m}$$



$$e^{-\frac{\pi d \pi}{4} \cdot \frac{16 \cdot \varepsilon^2}{(m \pi)^2}}$$

$$e^{-\frac{4 \varepsilon^2}{m \pi}}$$

$$E(z) = \frac{\pi}{4}$$

$$\Pr\left(|W - \frac{m\pi}{4}| \geq \varepsilon\right)$$

$$E(W) = \frac{m\pi}{4}$$

$$\leq e^{-\frac{m\pi}{4} \cdot \left(\frac{4\varepsilon}{m\pi}\right)^2}$$

$$E(z) = \frac{\pi}{4}$$

$$\Pr\left(|\frac{4W}{m} - \pi| \geq \varepsilon\right) = \Pr\left(|\frac{4W}{m} - \pi| \geq \frac{\delta}{\varepsilon \cdot \pi}\right) \quad E(z') = \frac{\pi}{4}$$

$$\leq 2e^{-\pi \cdot \left(\frac{\varepsilon}{\pi}\right)^2 / 3}$$

$$= 2e^{-\varepsilon^2 / 3\pi}$$

standard deviation σ of Z will be

$$\sqrt{E(z'^2) - (E(z))^2} = \sqrt{\frac{\pi}{4} - \left(\frac{\pi}{4}\right)^2} =$$

$$\Rightarrow \sqrt{\left(\frac{\pi}{4}\right)\left(1-\frac{\pi}{4}\right)} = \sigma =$$

$$\text{Var}\left(\frac{Z_1 + Z_2 + \dots + Z_m}{m}\right) = \frac{1}{m^2} \cdot m \cdot \left(\frac{\pi}{4}\right)\left(1-\frac{\pi}{4}\right)$$
$$= \frac{\left(\frac{\pi}{4}\right)\left(1-\frac{\pi}{4}\right)}{m} = \varepsilon^2$$

$$m = \frac{4\left(\frac{\pi}{4}\right)\left(1-\frac{\pi}{4}\right)}{\varepsilon^2}$$

0.001

for our estimation to be 0.001-accurate, we would need approximately 2,696,766 trials.

Alternatively, if we run our experiment 200,000 times, our estimate would be 0.003670 accurate.

$$\Pr(X > \mu) < \frac{E(X)}{\mu} = 1$$

$$\Pr(X < \mu) > 0$$

$$n_1 + n_2 \equiv 0 \pmod{3}$$

$$n_1 \quad n_2$$

$$(0,0) \checkmark$$

$$(0,1) \times$$

$$(0,2) \times$$

$$(1,0) \times$$

$$(1,1) \times$$

$$(1,2) \checkmark$$

$$(2,0) \times$$

$$(2,1) \checkmark$$

$$(2,2) \times$$

$$n = 3$$

g

$$P = 2$$

$$f_0(n) = 2n$$

$$f_0(n) \equiv 0 \pmod{2}$$

$$f_1(n) \equiv n \pmod{2} \quad f_1(n) = n$$

$$h_0(n) = g(f_0(n)) \quad n, y = 1, 2$$

$$h_1(n) = g(f_1(n))$$

$$h_0(n) = g(2n) \equiv 2n \pmod{3}$$

$$h_1(n) = g(n) \equiv n \pmod{3}$$

$$H = \{h_0, h_1\}$$

$$(x_1 \vee \bar{x}_2) \wedge (x_2 \vee x_3) \wedge (\bar{x}_3 \vee x_1)$$

$$y_1 + (1-y_2) \geq z_1$$

$$\Pr(X =$$

$$e_1, e_2, \dots, e_m$$

$x_i \rightarrow 1$ if e_i is betⁿ A & B

$$\begin{aligned} E(x_i) &= 1 \left[\Pr([u_i = A \wedge v_i = B] \vee [u_i = B \wedge v_i = A]) \right] \\ &\quad + 0 \left[\Pr([u_i = A \wedge v_i = A] \vee [u_i = B \wedge v_i = B]) \right] \\ &= 1 \cdot \left[\frac{1}{2} \times \frac{1}{2} + \frac{1}{2} \times \frac{1}{2} \right] \\ &= \frac{1}{2} \end{aligned}$$

$$E(x) = E\left[\sum_{i=1}^m x_i\right] = \sum_{i=1}^m E[x_i] = \sum_{i=1}^m \frac{1}{2} = \frac{m}{2}$$

$$\Pr(x \geq \frac{m}{2}) = P = \sum_{i=\lceil m/2 \rceil}^m \Pr(x=i)$$

$$1 - P = \sum_{i=0}^{\lfloor m/2 \rfloor} \Pr(x=i)$$

$$E(X) = \sum_{i=0}^m i \cdot \Pr(X=i) = m/2$$

$$= \sum_{i=0}^{m/2-1} i \cdot \Pr(X=i) + \sum_{i=m/2}^m i \cdot \Pr(X=i)$$

$$\leq \left(\frac{m}{2}-1\right)(1-p) + m p$$

$$\cancel{\frac{m}{2}} \leq p \left[m - \frac{m}{2} + 1\right] + \cancel{\frac{m}{2}} - 1$$

$$p \left(\frac{m}{2} + 1\right) \geq 1$$

$$p \geq \frac{1}{\frac{m}{2} + 1}$$

$\left(\frac{m}{2} + 1\right) \rightarrow \frac{1}{p}$, in expectation, the number of trials to get

To find a cut with at least $m/2$ edges, simply sample $m/2 + 1$ random assignments of vertices to $A \& B$.



$$E(X | x_1) = E(X)$$

$$\begin{aligned} E(X | x_1, \dots, x_k) &= \frac{1}{2} E(X | x_1, \dots, x_k, x_{k+1}=A) \\ &\quad + \frac{1}{2} E(X | x_1, x_2, \dots, x_{k+1}=B) \\ &\leq E(X | x_1, x_2, \dots, x_{k+1}) \end{aligned}$$

To calculate conditioned expectation,

- 1) calculate the # edges contributing to the cut betⁿ A & B by the vertices that are already placed. Call it λ .
- 2) The condition expectation is just the $\lambda + \frac{1}{2} \cdot \# \text{ remaining vertices}$.

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サシスセリ

ラリルレロ

いま二日寺ごろです。

ナニヌネノ

タツ ッテトヲ

山口 甲山 田中

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中山

私 東京

東京はにすんでいますか？

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Are you from Tokyo?

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東京 東 日本
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let $X = \sum_{i=1}^m X_i$, $\mu = E(X)$

and $p_i = \Pr(X_i = 1)$

$$\Pr(|X - \mu| \geq \delta\mu) \leq 2e^{-\mu\delta^2/3}$$

we want to bound the probability of the following

$$\Pr\left(\left|\frac{\sum X_i}{m} - \mu\right| \geq \varepsilon\mu\right)$$

$$\Pr(X_i = 1) = \mu$$

$$E(X_i) = \mu$$

$$\therefore E(X) = \mu m$$

$$\Pr\left(\left|\frac{1}{m}\sum X_i - \mu\right| \geq \varepsilon\mu\right) \leq 2e^{-\varepsilon^2 m \mu / 3} \leq \delta$$

we want

$$\delta \geq 2e^{-\varepsilon^2 m \mu / 3}$$

$$\log\left(\frac{\delta}{2}\right) \geq -\frac{\varepsilon^2 m \mu}{3}$$

$$3\log\left(\frac{2}{\delta}\right) \leq \varepsilon^2 m \mu$$

$$m \geq 3\log\left(\frac{2}{\delta}\right) \cdot \frac{1}{\varepsilon^2 \mu}$$

$$f((x+y) \bmod n) = [f(x) + f(y)] \bmod m$$

n	$F(n)$
1	$f(1)$
2	$f(2)$
3	$f(3)$
\vdots	\vdots
$n-1$	$f(n-1)$

P1 state $\rightarrow x$

$-\ell_1, \dots, -2, -1, 0, 1, 2, 3, \dots, \ell_2$

$$x = x_1 + \dots + x_n$$

$$P(|x|^2 = i^2)$$