

## Lecture #2

- Two events if and only if the events do not causally affect each other

$e_2^4$  &  $e_3^1 \rightarrow$  concurrent

If it is not causal, then it concurrent.

If it is physically concurrent, then it is definitely concurrent.

- Logical Time Implementation details: 1) scalar time 2) Vector Time  
3) Matrix Time

Partial ordering bet<sup>n</sup> events

scalar logical time cannot have strong consistency.

Concurrenay is not transitive

## Lecture #3

$$\begin{matrix} i & \left[ \begin{matrix} K \\ < P \end{matrix} \right] \\ j & \left[ \begin{matrix} < K \\ P \end{matrix} \right] \end{matrix}$$

$\underbrace{\qquad\qquad}_{\text{concurrent}}$

## lecture #4

$i^{\text{th}}$  row is your vector in matrix clock

$(j, j)$

Home work: Q:1) why should we not take man along column instead of rows?

Q:2) Can we check all if every one has received the broadcast using matrix clock?

If causal system, how do we mark global state?

# Lecture #6

Do we need both ME conditions in RA?

Do you play by the rules or quietly slip into CS as if nothing happened?

Is there disadvantage of breaking rules?

Raymond's Algorithm

Home work:

Will the complexity of the Raymond's Algorithm become  $O(n)$  from  $O(\log n)$ , will the K-ary tree ever become linear chain?

# Lecture # 8

# Lecture #9

- i. Execution of TD algorithm should not indefinitely

Phantom deadlock - detecting deadlock that doesn't exist  
in DS

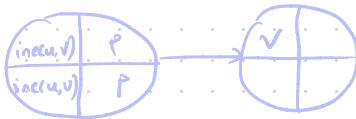


# Lecture #10

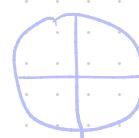
# Lecture #11



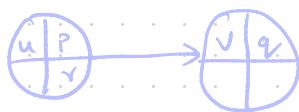
Block



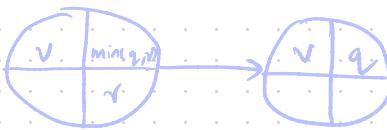
Activate



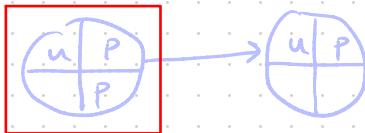
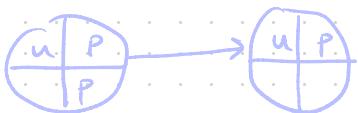
Transmit



$(u < v) \text{ or } (u = v, P > q)$



Detect



$$\begin{array}{c|c} 8 & 3 \\ \hline 8 & 2 \end{array} \quad \xrightarrow{1} \quad \begin{array}{c|c} 8 & 3 \\ \hline 8 & 3 \end{array}$$

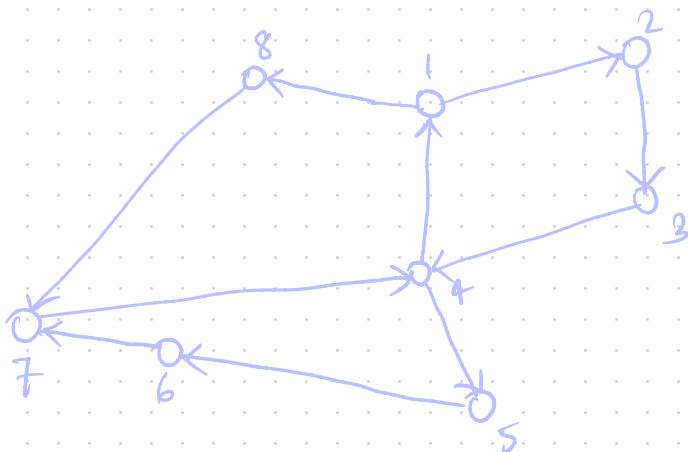
$\uparrow 3$

$$\begin{array}{c|c} 7 & 4 \\ \hline 7 & 4 \end{array} \quad \xleftarrow{2} \quad \begin{array}{c|c} 6 & 1 \\ \hline 6 & 1 \end{array}$$

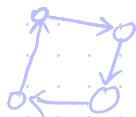
$\downarrow 4$

is it always the last edge that detects that there is a deadlock  
in Mitchell-Merritt algorithm?

# Lecture #12



Phantom deadlock detection  
CMH



# Lecture #13

## Paper Discussion - Mid Semester

1. (a) Mitchell Merritt question  
(b) chandy lamport request queue question  
(c) will it remain deadlock if initiator aborting himself

$s \rightarrow$  AND and OR

(d) Global consistent

2. Termination detection

3. New 2Phase Commit

4. Modify Matrix clock using signal - Kshemkalyani



all non-tail

# Lecture # 14

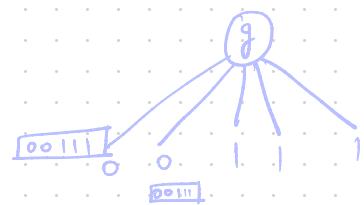
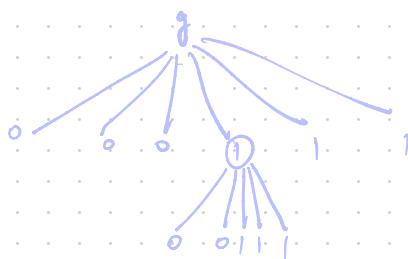
Given  $n$ , if  $n \neq 3m+1$ , then it is not possible to come up with the a solution to Byzantine problem.

H0 - Rama moorthy

4,1

$$N = 7, t = 2$$

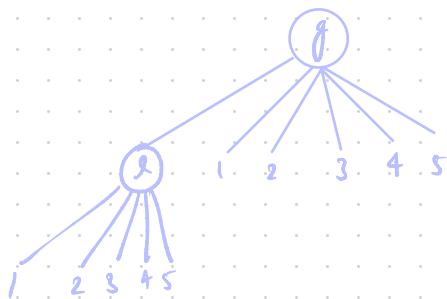
$$N = 7,$$



Slip test question

Find pattern in which is faulty  
or when it's not.

# Lecture #15



# Lecture #16

Slip Test

22/03/24

If all the non-malicious processes prefer V at the start of phase K, then all do at end of phase K.

Each of the non-malicious process receives at least  $n-f$  preferences for V at the end of the first round of phase K.

Now, since  $n > 4f$  it follows that  $n-f > 3f = \frac{n}{2} + f$

∴ each of the non-malicious process still prefer V.

∴ This implies from phase K onwards, all non-mal processes have the same preference.

$$\begin{aligned} n-f &= \frac{n}{2} + \frac{n}{2} - f > \frac{n}{2} + 2f - f \\ &= \frac{n}{2} + f \end{aligned}$$

# Lecture #17

## Minimum Spanning Tree

## Lecture #18

HW: At any level, sends at most one best net leading

HW: Take any MST run normal MST, & run this algo to  
get a clear idea

$$n \rightarrow ?$$

# lecture #19

Merkel Trees

Bitcoin uses Merkle Trees

$$HPI = (H(x), \text{Pointer})$$

B L I

Watermarking of Link Prediction Task in Graph Neural Network using  
backdoors