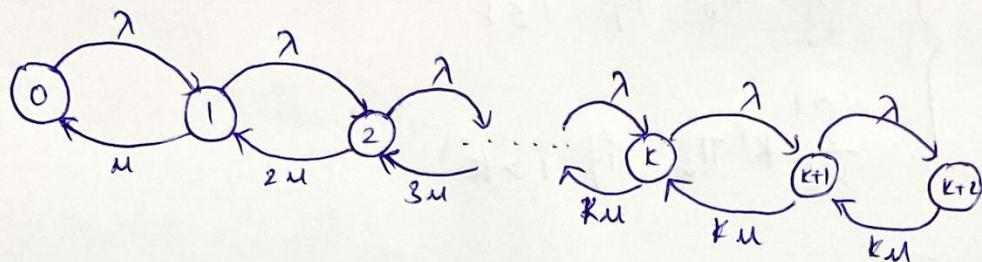


1.1

Performance Modelling for Computer System - Assignment 2
 Aayyan Ajay Sharma - 2022121001.

- The state-space diagram of the given M/M/K/ ∞ queue is:



The time-reversibility equations are as follows:

state

Time-Reversibility Eqn.

Simplified Equation

0

$$\pi_0 \lambda = \pi_1 \mu$$

$$\pi_1 = \frac{\lambda}{\mu} \pi_0$$

1

$$\pi_1 \lambda = \pi_2 (2\mu)$$

$$\pi_2 = \left(\frac{\lambda}{\mu}\right)^2 \frac{1}{2!} \pi_0$$

2

$$\pi_2 \lambda = \pi_3 (3\mu)$$

$$\pi_3 = \left(\frac{\lambda}{\mu}\right)^3 \frac{1}{3!} \pi_0$$

K-1

$$\pi_{K-1} \lambda = \pi_K (K\mu)$$

$$\pi_K = \left(\frac{\lambda}{\mu}\right)^K \frac{1}{K!} \pi_0$$

K

$$\pi_K \lambda = \pi_{K+1} (K\mu)$$

$$\pi_{K+1} = \left(\frac{\lambda}{\mu}\right)^{K+1} \left(\frac{1}{K!}\right) \cdot \frac{1}{K+1} \pi_0$$

K+1

$$\pi_{K+1} \lambda = \pi_{K+2} (K\mu)$$

$$\pi_{K+2} = \left(\frac{\lambda}{\mu}\right)^{K+2} \left(\frac{1}{K!}\right) \cdot \frac{1}{K+2} \pi_0$$

∴

$$\pi_i = \begin{cases} \left(\frac{\lambda}{\mu}\right)^i \frac{1}{i!} \pi_0 & \text{if } i \leq K \\ \left(\frac{\lambda}{\mu}\right)^i \frac{1}{i!} \cdot \left(\frac{1}{K!}\right)^{i-K} \pi_0 & \text{if } i > K \end{cases}$$

1.2

Now, define $\rho = \frac{\lambda}{k\mu}$ as the system utilization.

So, π_i written in terms of ρ will be

$$\pi_i = \begin{cases} \frac{(k\rho)^i}{i!} \pi_0 & \text{if } i \leq k \\ \frac{\rho^i}{k!} \cdot k^k \pi_0 & \text{if } i > k \end{cases}$$

To determine π_0 :

$$\left[\sum_{i=0}^{k-1} \pi_i + \sum_{i=k}^{\infty} \frac{\rho^i}{k!} \cdot k^k \right] = 1$$

$$\Rightarrow \pi_0 \left[\sum_{i=0}^{k-1} \frac{(k\rho)^i}{i!} + \sum_{i=k}^{\infty} \frac{\rho^i}{k!} \cdot k^k \right] = 1$$

$$\Rightarrow \pi_0 \left[\sum_{i=0}^{k-1} \frac{(k\rho)^i}{i!} + \frac{k^k}{k!} \cdot \frac{\rho^k}{1-\rho} \right] = 1$$

$$\therefore \pi_0 = \left[\sum_{i=0}^k \frac{(k\rho)^i}{i!} + \frac{(k\rho)^k}{k!(1-\rho)} \right]^{-1}$$

Next, we find the queuing probability P_Q , and expected probability of number of jobs in the queue $E(N_Q)$, Expected number of time spent in queue $E(T_Q)$ and Expected Response time $E(T)$.

All of these will help us in finding the mean number of jobs in the system, $E(N)$.

1.3

Queuing Probability P_Q :

$$\begin{aligned}
 P_Q &= P \left\{ \text{An arrival finds all servers busy} \right\} \\
 &= P \left\{ \text{An arrival sees } \geq k \text{ jobs in the system} \right\} \\
 &= \text{limiting probability that there are } \geq k \text{ jobs in the system.} \\
 &= \sum_{i=k}^{\infty} \pi_i \\
 &= \frac{k^K}{K!} \pi_0 \sum_{i=k}^{\infty} p^i \\
 &= \frac{(kp)^K}{K!(1-p)} \pi_0, \text{ where } \pi_0 = \left[\sum_{i=0}^{K-1} \frac{(kp)^i}{i!} + \frac{(kp)^K}{K!(1-p)} \right]^{-1}
 \end{aligned}$$

Expected Number in the Queue ($E(N_Q)$):

$$\begin{aligned}
 E(N_Q) &= \sum_{i=k}^{\infty} \pi_i (i-k) \\
 &= \pi_0 \sum_{i=k}^{\infty} \frac{p^i k^k}{k!} (i-k) \\
 &= \pi_0 \cdot \frac{p^k k^k}{k!} \cdot \sum_{i=0}^{\infty} p^i \cdot i \\
 &= \pi_0 \cdot \frac{p^k k^k}{k!} \cdot \frac{p}{(1-p)^2} \\
 &= \frac{P_Q}{\pi_0} \cdot P_Q \cdot \frac{p}{1-p}.
 \end{aligned}$$

1.4

Now,

$$E(T_d) = \frac{1}{\lambda} \cdot E(N_d) = \frac{1}{\lambda} \cdot P_d \cdot \frac{P}{1-P}$$

$$E(T) = E(T_d) + \frac{1}{\mu} = \frac{1}{\mu} + \frac{1}{\lambda} \cdot \frac{P}{1-P} \cdot P_d$$

$$E(N) = \lambda \cdot E(T) = P_d \cdot \frac{P}{1-P} + k_p$$

∴ Expected No. of jobs

$$E(N) = k_p + P_d \frac{P}{1-P} = \text{mean no. of jobs in the system.}$$

we see that as $P \rightarrow 1^-$, $E(N) \rightarrow \infty$

∴ For the system to be stable,

$$P < 1 \quad \text{or} \quad \frac{\lambda}{k\mu} < 1$$

or $\lambda < k\mu$ → stability condition.

2.1

2. $M/M/2/\infty$ queue.

Inter-arrival time rate $\sim \text{Exp}(\lambda)$

Service rate for server 1 $\sim \text{Exp}(\mu_1)$

Service rate for server 2 $\sim \text{Exp}(\mu_2)$

We let our descriptor be: $[n, y, k]$
where

n : State of server 1 (binary: $0 \rightarrow \text{Idle}, 1 \rightarrow \text{Busy}$).

y : State of server 2 (binary: $0 \rightarrow \text{Idle}, 1 \rightarrow \text{busy}$).

k : No. of jobs in the queue [No. of jobs in system -
No. of jobs being served].

We will map our descriptor using the following mapping:

$[n, y, k] \rightarrow \text{state } (2n+y+k)$.

\therefore The map looks as follows:

$[0, 0, 0] \rightarrow \text{state 0}$

$[0, 1, 0] \rightarrow \text{state 1}$

$[1, 0, 0] \rightarrow \text{state 2}$

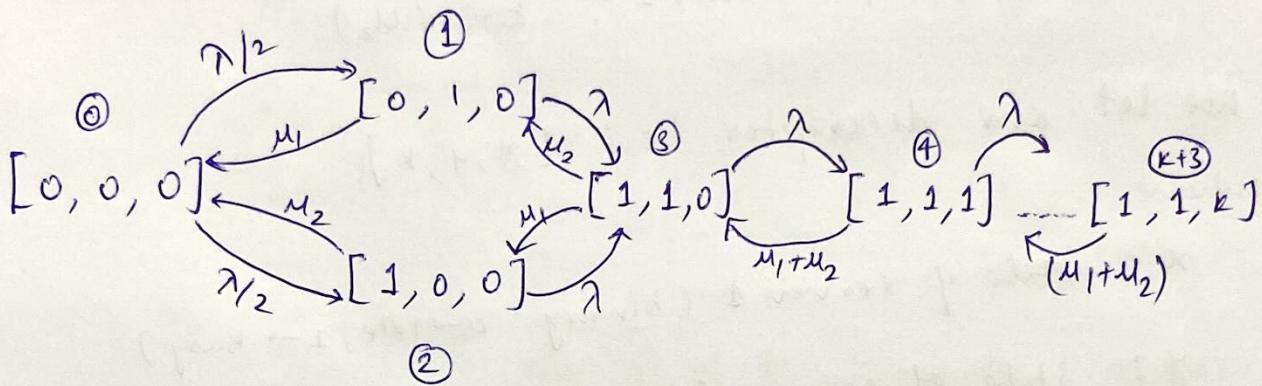
$[1, 1, 0] \rightarrow \text{state 3}$

$[1, 1, 1] \rightarrow \text{state 4}$

$[1, 1, k] \rightarrow \text{state } k$.

2.2

ASSUMING when the first job comes, it gets randomly assigned either to server 1 or 2, the state space diagram of our Markov chain will look as follows:



Looking at the Markov Chain, we get the δ -Matrix as following :

$$\delta = \begin{matrix} & 0 & 1 & 2 & 3 & 4 & \dots \\ 0 & -\lambda & \lambda/2 & \lambda/2 & 0 & 0 & \dots \\ 1 & \mu_1 & -(\mu_1 + \lambda) & 0 & 0 & \lambda & \dots \\ 2 & \mu_2 & 0 & -(\mu_2 + \lambda) & \lambda & 0 & \dots \\ 3 & 0 & \mu_1 & \mu_2 & -(\mu_1 + \mu_2) & \lambda & \dots \\ 4 & & & & & & \dots \\ \vdots & & & & & & \end{matrix}$$

The dimensions of δ are $|1| \times |1| = \sum_{i=0}^k 1$

Now, using Time Reversibility Equations, we get:

State	Time Reversibility Eqn.	Simplified Eqn.
1	$\frac{\lambda}{2}\pi_0 = \mu_1\pi_1$	$\pi_1 = \frac{\lambda}{2\mu_1}\pi_0$
2	$\frac{\lambda}{2}\pi_0 = \mu_2\pi_2$	$\pi_2 = \frac{\lambda}{2\mu_2}\pi_0$
3	$\lambda\pi_1 = \mu_2\pi_3$	$\pi_3 = \frac{\lambda}{\mu_2}\pi_1 = \frac{\lambda^2}{2\mu_1\mu_2}\pi_0$

2.3

From the last equation, we get

$$\pi_n = \left(\frac{\lambda}{\mu_1 + \mu_2} \right)^{n-3} \cdot \left(\frac{\lambda^2}{2\mu_1 \mu_2} \right) \cdot \pi_0 \quad n \geq 3$$

To get π_0 , we have

$$\sum_{i=0}^{\infty} \pi_i = 1$$

$$\Rightarrow \pi_0 + \pi_1 + \pi_2 + \sum_{i=3}^{\infty} \pi_i = 1$$

$$\Rightarrow \pi_0 \left[1 + \frac{\lambda}{2\mu_1} + \frac{\lambda}{2\mu_2} + \sum_{i=3}^{\infty} \left(\frac{\lambda}{\mu_1 + \mu_2} \right)^{i-3} \cdot \frac{\lambda^2}{2\mu_1 \mu_2} \right] = 1$$

$$\Rightarrow \pi_0 \left[1 + \frac{\lambda}{2\mu_1} + \frac{\lambda}{2\mu_2} + \frac{\lambda^2}{2\mu_1 \mu_2} \cdot \frac{1}{1 - \frac{\lambda}{\mu_1 + \mu_2}} \right] = 1$$

$$\Rightarrow \pi_0 \left[1 + \frac{\lambda}{2\mu_1} + \frac{\lambda}{2\mu_2} + \frac{\lambda^2}{2\mu_1 \mu_2} \left[\frac{\mu_1 + \mu_2}{\mu_1 + \mu_2 - \lambda} \right] \right] = 1$$

$$\Rightarrow \pi_0 = \left[1 + \frac{\lambda}{2\mu_1} + \frac{\lambda}{2\mu_2} + \frac{\lambda^2}{2\mu_1 \mu_2} \left(\frac{\mu_1 + \mu_2}{\mu_1 + \mu_2 - \lambda} \right) \right]^{-1}$$

∴ The stationary distribution is:

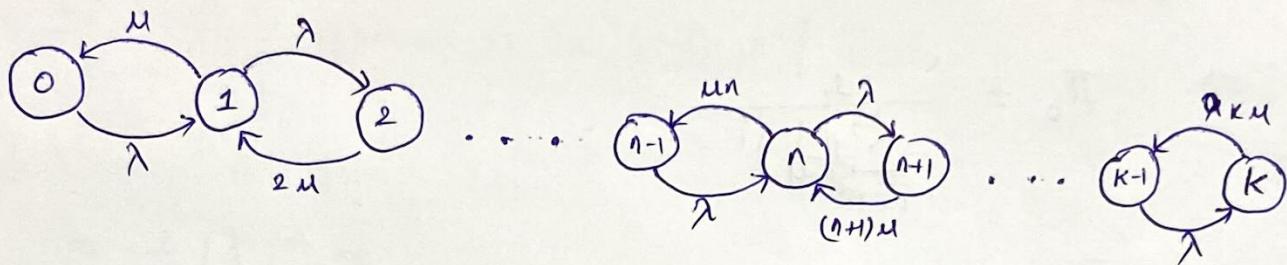
$$\pi_i = \left(\frac{\lambda}{\mu_1 + \mu_2} \right)^{i-3} \cdot \frac{\lambda^2}{2\mu_1 \mu_2} \cdot \pi_0, \quad i \geq 3$$

$$\pi_1 = \frac{\lambda}{2\mu_1} \pi_0, \quad \pi_2 = \frac{\lambda}{2\mu_2} \pi_0$$

$$\pi_3 = \frac{\lambda^2}{2\mu_1 \mu_2} \cdot \underline{\underline{Ans}}$$

3.1

3. The state space diagram for an M/M/k/k queue is as follows:



The Global balance equation for M/M/k/k will be:

$$\pi_n(\lambda + n\mu) = \pi_{n-1} \cdot \lambda + \pi_{n+1} \cdot (n+1)\mu, \text{ where } 0 < n < k \rightarrow ①$$

The local balance equation will be (for state 0):

$$\pi_0 \circ \lambda = \pi_1 \circ \mu.$$

$$\therefore \pi_1 = \frac{\lambda}{\mu} \pi_0. \rightarrow ②$$

$$\pi_1(\lambda + \mu) = \pi_0 \lambda + \pi_2(2\mu) \quad [n=1 \text{ in } ①]$$

$$\pi_1(\lambda + \mu) = \pi_1 \mu + \pi_2(2\mu) \quad [\cancel{\text{in}}]$$

$$\pi_1 \circ \lambda = \pi_2(2\mu)$$

$$\therefore \pi_2 = \frac{\lambda}{2\mu} \circ \pi_1.$$

Similarly, for π_i , we find

$$\pi_i = \frac{\lambda}{i\mu} \circ \pi_{i-1}. \quad i \leq k.$$

$$\therefore \pi_i = \left(\frac{\lambda}{\mu}\right)^i \cdot \frac{1}{i \cdot (i-1) \cdots 1} \cdot \pi_0 = \frac{1}{i!} \left(\frac{\lambda}{\mu}\right)^i \pi_0.$$

Now,

$$\sum_{i=0}^K \pi_i^o = 1$$

$$\Rightarrow \sum_{i=0}^K \frac{1}{i!} \cdot \left(\frac{\lambda}{\mu}\right)^i \pi_0^o = 1$$

$$\Rightarrow \pi_0^o = \frac{1}{\sum_{i=0}^K \frac{1}{i!} \left(\frac{\lambda}{\mu}\right)^i}$$

Now, there is no closed form for the expression

$$\sum_{i=0}^K \frac{1}{i!} \left(\frac{\lambda}{\mu}\right)^i, \text{ but we can express it as follows:}$$

$$\sum_{i=0}^K \frac{1}{i!} \left(\frac{\lambda}{\mu}\right)^i = \frac{\Gamma(K+1, \lambda/\mu)}{\Gamma(K+1)} \cdot e^{\lambda/\mu}$$

where $\Gamma(n, y)$ is the incomplete Gamma function

and $\Gamma(n, y) = \int_y^\infty e^{-t} t^{n-1} dt$. and $\Gamma(n)$ is the

classical Gamma function where $\Gamma(n) = \int_0^\infty e^{-t} \cdot t^{n-1} dt$.

and $\Gamma(n) = (n-1)!$ when $n \in \mathbb{N} \cup \{0\}$.

$$\therefore \pi_0^o = \frac{\Gamma(K+1)}{\Gamma(K+1, \lambda/\mu)} e^{-\lambda/\mu}$$

\therefore The stationary distribution for M/M/K/X is

$$\pi_i^o = \frac{1}{i!} \left(\frac{\lambda}{\mu}\right)^i \cdot \pi_0^o \quad \forall i \leq K.$$

3.3

The blocking probability = Probability that the arriving jobs
get blocked

= Probability that our system has K servers busy

= $P \{$ system is in state $K \}$

= π_K

$$= \frac{1}{K!} \left(\frac{\lambda}{\mu}\right)^K \cdot \pi_0.$$

4.1

4. Given that $N(t) = \min(A(t) - D(t), 0)$,

To prove: $N(t)$ is a Markov chain, where $A(t)$ = Arrivals till time t & $D(t)$ = Departures till time t .

Consider the quantity / Probability:

$$P\left\{N(t_{k-1}+\varepsilon) = j \mid N(0), N(t_1), \dots, N(t_{k-1}) = l\right\}$$

for some $\varepsilon > 0$ and $\forall i \quad 0 < t_i < \infty \quad \& \quad i = 0, 1, 2, \dots, k-1$.

Case 1: If $j = l-1 \Rightarrow$ in duration ε , the job currently receiving service finished server being serviced.

$$\therefore P(\text{Remaining service for current job} \leq \varepsilon)$$

$$= P(\text{Service time of a job} \leq \varepsilon)$$

This is the memoryless property of, which satisfied only by exponential distribution.

$$\therefore P(\text{Service time of a job} \leq \varepsilon) = \underset{\cancel{u=\text{rate}}}{\int_0^{\varepsilon} u e^{-ut} dt}$$

Case 2: If $j = l$,

This means there were no arrivals & departure in ε time

$$\therefore P\left\{\text{no arrival \& no departure in } \varepsilon\right\} = \int_0^{\varepsilon} (1-\lambda t) e^{-\mu t} dt$$

Case 3: If $j = l+1$, then in ε time there is one arrival,

$$\therefore P\{\text{one arrival in time } \varepsilon\} = \lambda \varepsilon$$

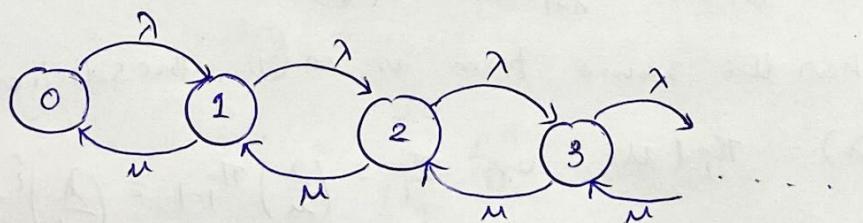
4.2

$$\therefore P \left\{ N(t_{k-1} + \varepsilon) = j \mid N(t_0), N(t_1), \dots, N(t_{k-1}) = l \right\}$$

$$= \begin{cases} \int_0^{\varepsilon} \mu e^{-\mu t} dt & \text{if } j = l-1 \\ \lambda \varepsilon & \text{if } j = l+1 \\ \int_0^{\varepsilon} (1-\lambda t) \cdot e^{\lambda t} dt & \text{if } j = l \end{cases}$$

$N(t)$ is a Markovian Descriptor for $M/M/1/\infty$ queue.

Now, for an $M/M/1/\infty$ queue, the state space diagram is as follows:



looking the diagram, we see that:

$$q_{01} = \lambda, q_{12} = \lambda, q_{23} = \lambda \dots$$

$$q_{10} = \mu, q_{21} = \mu, q_{32} = \mu \dots$$

Assuming Φ is conservative, we get:

$$q_{00} = -\lambda$$

$$q_{ii} = -(\lambda + \mu) \text{ for } i > 0 \text{ & } i \in \mathbb{Z}$$

$$q_{11} = -(\lambda + \mu)$$

$$q_{22} = -(\lambda + \mu)$$

4.83

$$\therefore Q = \begin{matrix} & 0 & 1 & 2 & 3 & 4 & \dots \\ 0 & -\lambda & \lambda & 0 & 0 & 0 & \\ 1 & \mu & -(\lambda+\mu) & \lambda & 0 & 0 & \\ 2 & 0 & \mu & -(\lambda+\mu) & \lambda & 0 & \\ 3 & 0 & 0 & \mu & -(\lambda+\mu) & \lambda & \\ 4 & 0 & 0 & 0 & \mu & -(\lambda+\mu) & \\ \vdots & & & & & & \end{matrix}$$

where the dimensions of Q are $|N \times N| = N$.

To find stationary distribution, we find the balance equations for the first few states :

$$\pi_0(\lambda) = \pi_1(\mu) \rightarrow \textcircled{i}$$

$$\pi_1(\lambda+\mu) = \pi_0(\lambda) + \pi_2(\mu) \rightarrow \textcircled{ii}.$$

$$\Rightarrow \pi_1\lambda = \pi_2(\mu) \rightarrow \{ \text{using } \textcircled{i} \text{ & } \textcircled{ii} \}.$$

The Global Balance equation is as follows :

$$\pi_n(\lambda+\mu) = \pi_{n-1}(\lambda) + \pi_{n+1}(\mu) \rightarrow \textcircled{iii}$$

and since it has the same form as ~~\textcircled{ii}~~, the solution will be:

$$\pi_{i-1}(\lambda) = \pi_i(\mu) \text{ or } \pi_i = \left(\frac{\lambda}{\mu}\right) \pi_{i-1} = \left(\frac{\lambda}{\mu}\right)^i \pi_0.$$

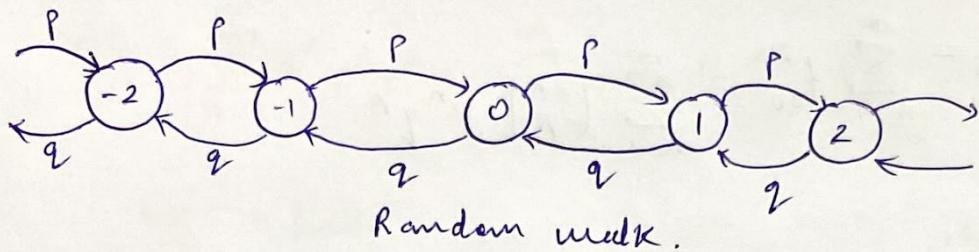
Now, since $\sum_{i=0}^{\infty} \pi_i = 1$

$$\begin{aligned} \sum_{i=0}^{\infty} \pi_i &= \sum_{i=0}^{\infty} \left(\frac{\lambda}{\mu}\right)^i \pi_0 = 1 \\ &= \frac{1}{1 - \frac{\lambda}{\mu}} \pi_0 = 1 \Rightarrow \pi_0 = 1 - \frac{\lambda}{\mu}. \end{aligned}$$

$$\therefore \pi_i = \left(\frac{\lambda}{\mu}\right)^i \left(1 - \frac{\lambda}{\mu}\right).$$

5.1

S. The state diagram for the described random walk is:



In the diagram, $q = 1 - p$.

Since all the states are accessible to each other / they communicate, either all states are recurrent or all states are transient.

∴ To determine whether state 0 is transient or recurrent it should suffice to determine whether the chain as a whole is recurrent or not.

For that, we need to calculate

$$F_{ii} = \sum_{n=1}^{\infty} f_{ii}^n \quad \text{for } i=0$$

where $f_{ii}^n \neq f_{ii}^{(n)} = \text{Probability of starting in } i \text{ and returning to state } i \text{ for the first time in } n \text{-steps}$

$$\therefore F_{00} = \sum_{n=1}^{\infty} f_{00}^n$$

→ ~~the concept~~ If F_{00} is finite, then the chain is transient otherwise recurrent.

5.2

since we cannot go from state 0 to itself in odd number of steps, it follows that

$$F_{00} = \sum_{n=1}^{\infty} t_{00}^{2n} = \sum_{n=1}^{\infty} \binom{2n}{n} \cdot p^n q^n.$$

Consider,

$$\sum_{k=0}^{2n} \binom{2n}{k} = (1+1)^{2n} = 2^{2n} = 4^n.$$

since the last term cannot be greater than its sum, we have

$$\binom{2n}{n} < 4^n$$

and since the last term always greater than the average we have

$$\frac{4^n}{2n+1} < \binom{2n}{n}.$$

$$\therefore \frac{4^n}{2n+1} < \binom{2n}{n} < 4^n.$$

This is Misha/Lavrov's lemma.

$$\therefore \sum_{n=0}^{\infty} \frac{4^n}{(2n+1)} p^n q^n < F_{00} < \sum_{n=0}^{\infty} 4^n p^n q^n.$$

when $p=q=1/2$, LHS = $\sum_{n=1}^{\infty} \frac{1}{2n+1}$ = Harmonic series = ∞

But if $p \neq q$ \therefore when $p=q=1/2$, $F_{00} = \infty$ &

the chain is recurrent.

But if $p \neq q$, then $\frac{p+q}{2} < \sqrt{pq} \Rightarrow \frac{1}{2} < \sqrt{pq}$ or $pq < \frac{1}{4}$.

$$\therefore \text{RHS} = \sum_{n=1}^{\infty} 4^n p^n q^n = \sum_{n=1}^{\infty} (4pq)^n = \frac{4pq}{1-4pq} \text{ since } 4pq < 1.$$

$\therefore \text{RHS} = \text{finite.}$

\therefore when $p \neq q$, F_{00} is finite and therefore the chain would be transient.

The proof of the mean recurrence time of state 0 is given at the end, on page 5.4.

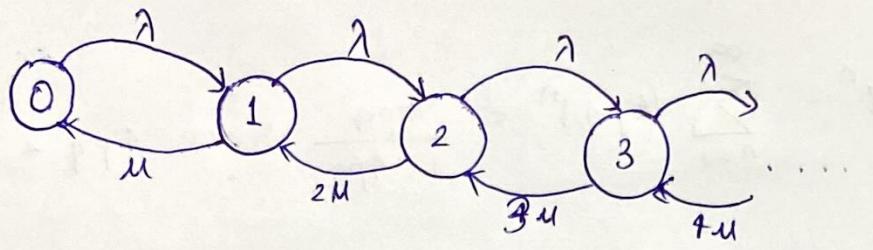
Now, since the mean recurrence time of state 0 is infinite, this means that state 0 is null recurrent. And since state 0 is null recurrent, and since every state communicates with each other, they too are null recurrent. (Theorem 9.22 Mor Balter, pg. 162)

This implies $\forall i \in \mathbb{Z}$, $M_{ii}^0 = \text{mean recurrence time of state } i = \infty$.

\therefore The chain is null recurrent and hence the mean recurrence time for the chain is infinite.

6.1

The state space diagram for M/M/ ∞ queue is



The local balance equation for state 0 is:

$$\pi_0(\lambda) = \pi_1(\mu) \rightarrow \textcircled{i}$$

The global balance equation for state i , where $i > 0$ is:

$$\pi_i(\lambda + \mu) = \pi_{i-1} \cdot \lambda + \pi_{i+1}((i+1)\mu) \rightarrow \textcircled{ii}$$

Solving it for $i = 1, 2, \dots$ we get the value of π_i as:

$$\pi_i = \frac{\lambda}{i\mu} \cdot \pi_{i-1} \rightarrow \textcircled{iii}$$

Example: For $i=1$

$$\pi_1(\lambda + \mu) = \cancel{\pi_0} \cdot \lambda + \pi_2 \cdot (2\mu)$$

$$\pi_1(\lambda + \mu) = \pi_1 \mu + \pi_2 (2\mu) \quad [\text{using } \textcircled{i}]$$

$$\lambda \pi_1 \cancel{\pi_1(\mu)} = 2\mu \pi_2$$

$$\Rightarrow \pi_2 = \cancel{2\mu} \frac{\lambda}{2\mu} \pi_1$$

Similarly, we get/do this for all $i = 1, 2, \dots$

$$\therefore \pi_i = \frac{\lambda}{i\mu} \cdot \pi_{i-1}$$

$$\begin{aligned} \therefore \pi_i &= \left(\frac{\lambda}{1\mu} \right) \left(\frac{\lambda}{(i-1)\mu} \right) \cdot \left(\frac{\lambda}{(i-2)\mu} \right) \cdot \left(\frac{\lambda}{(i-3)\mu} \right) \cdots \left(\frac{\lambda}{1\mu} \right) \pi_0 \\ &= \frac{1}{i!} \left(\frac{\lambda}{\mu} \right)^i \cdot \pi_0. \end{aligned}$$

6.2

since $\sum \pi_i = 1$, we get:

$$\sum_{i=0}^{\infty} \pi_i = 1$$

$$\Rightarrow \sum_{i=0}^{\infty} \frac{1}{i!} \left(\frac{\lambda}{\mu}\right)^i \pi_0 = 1$$

$$\Rightarrow \pi_0 \sum_{i=0}^{\infty} \frac{1}{i!} \left(\frac{\lambda}{\mu}\right)^i = 1$$

$$\Rightarrow \pi_0 e^{\lambda/\mu} = 1 \Rightarrow \boxed{\pi_0 = e^{-\lambda/\mu}}$$

$$\therefore \pi_i = \frac{1}{i!} \left(\frac{\lambda}{\mu}\right)^i \cdot e^{-\lambda/\mu}$$

Now, since X is number of jobs at stationarity and

since we know $\pi_i = \frac{1}{i!} \left(\frac{\lambda}{\mu}\right)^i e^{-\lambda/\mu}$, we know

$$X \sim \text{Poisson} \left(\frac{\lambda}{\mu} \right).$$

$$\therefore E(X) = \frac{\lambda}{\mu} = \text{mean no. of jobs at stationarity.}$$

Since expectation of Poisson RV with parameter λ is λ .

S.4 Proof $m_{ii} = \infty$

The mean recurrence time to revisit state i from itself is infinite. i.e. $m_{ii} = \infty \forall i \in \mathbb{Z}$

To prove this, we use the argument used in Mor Balter Book ^{Theorem} on page 163, which is as follows:

Theorem: (9.23) For the symmetric random walk with $p = 1/2$,

$m_{ij}^o = m_{00}^o = \infty$, where $m_{00} =$ mean time between visits to state 0.

Proof by contradiction: Assume $m_{00}^o \neq \infty$.

Then Here, we define $m_{ij}^o =$ mean time until we visit state j from i .

Now,

$$m_{00}^o = \frac{1}{2} m_{1,0} + \frac{1}{2} m_{-1,0} + 1$$

$$\therefore m_{1,0} \neq \infty.$$

$$\begin{aligned} \therefore m_{1,0} &= 1 + \frac{1}{2} \cdot 0 + \frac{1}{2} m_{2,0} \\ &= 1 + \frac{1}{2} m_{2,0}. \end{aligned}$$

$$\text{since } m_{2,1} = m_{1,0} \text{ and } m_{2,0} = m_{2,1} + m_{1,0}$$

$$\therefore m_{2,0} = 2 m_{1,0}$$

$$\therefore m_{1,0} = 1 + \frac{1}{2} (2 m_{2,0})$$

$$m_{1,0} = 1 + m_{1,0}$$

~~writes~~ which is a contradiction.

$\therefore m_{0,0}^o$ is infinite