

Updating Probability Values for Improved Bounds

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1 How the current implementation works

Although the theory paper uses Rashtchian's bounds for upper bounding $c_S(w)$, the code calculates the upper bound using the 4-term combinatorial summation expression for $c_S(w)$. The program only uses $l = \lfloor \log |S| \rfloor$. The dispersion index bound is set as $\rho = 1.1$. The expression for the dispersion index is given by

$$\frac{\sigma^2[Cnt_{(S,m)}]}{E[Cnt_{(S,m)}]} \leq \sum_{w=0}^{w=l} c_S(w) \cdot r(w, m) \leq \rho$$

$$\text{where } r(w, m) = \prod_{j=1}^{j=m} \left(\frac{1}{2} + \frac{1}{2}(1 - 2p_j)^w \right) - \frac{1}{2^m}.$$

For a fixed value of m , the p_i 's are started off with a value of 0.5 and are iteratively decreased until the variance bound inequality gets violated. For each term the $c_S(w)$ calculation is done by the 4-term combinatorial expression.

Note: The $c_S(w)$ bound derived in the LICS paper has a small error. The bound should be $c_S(w) \leq \frac{w}{2} \left(\frac{8e\sqrt{nl}}{w} \right)^w$ instead of $c_S(w) \leq 2 \left(\frac{8e\sqrt{nl}}{w} \right)^w$.

2 My First Approach

I used the improved bound I derived in the summer for $l = \lfloor \log |S| \rfloor$ given by

$$c_S(w) \leq \frac{w}{2} \left(\frac{4e\sqrt{nl}}{w} \right)^w$$

The $c_S(w)$ was upper bounded using this analytic expression instead of a combinatorial expression. The p_i 's were iterated in the same manner.

Results: This approach was coded in the file "newvalues1.py" led to higher values of p_i in comparison to the original approach. This was undesirable and probably happened because the improved analytical bound is still weak in comparison to the weaker old combinatorial expression. The program runs much faster than the original code although it gives weaker values.

3 My Second Approach

In the improved bounds for $c_S(w)$ I derived in the summer, instead of the 4-term combinatorial expression, for a fixed $x \in S$ I improved it to

$$c_S(w) \leq \sum_{(b,a) \in U'} e_{(b,a)}(S) \leq \sum_{(b,a) \in U'} \binom{n-|x|}{a} \binom{|x|}{a} \binom{|x|}{b-a} \leq \sum_{(b,a) \in U'} \binom{n-b}{a} \binom{l}{a} \binom{l}{b-a}$$

where $U' = \{(b, a) : b \geq a, b + a = w\}$.

We calculate these $c_S(w)$ values and substitute in the sum and iteratively find the p_i values like in the original approach.

Results: This approach was coded in the file "newvalues2.py" led to slightly higher values of p_i in comparison to the original approach. This is probably happening because of some optimisation included in the

original code that I have not yet incorporated into my code. The program runs significantly faster than the original code and will be able to calculate values for higher row values. I feel that the slightly higher probability values is something we can compromise for the significantly improved complexity.

4 My Third Approach

This approach was just based on using Rashtchian's old analytic bounds but for $l = \min\{\lfloor \log|S| \rfloor, \lfloor \frac{2\log|S|}{\log n - \log \log|S|} \rfloor\}$ instead of just $l = \lfloor \log|S| \rfloor$.

Results:

5 My Fourth Approach

This approach was just based on using Rashtchian's old combinatorial bounds but for $l = \min\{\lfloor \log|S| \rfloor, \lfloor \frac{2\log|S|}{\log n - \log \log|S|} \rfloor\}$ instead of just $l = \lfloor \log|S| \rfloor$. I have used the 4-term combinatorial sum used in the original code except for a changed value of l .

Results: