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## **Assignment 6 Design**

### **Description of Program**

#### **Deliverables**

- encode.c
  - This contains the implementation of the Huffman encoder
- decode.c
  - This contains the implementation of the Huffman decoder.
- defines.h
  - This declares the macros necessary for completing the assignment. This is given in the resources repository by Professor Long
- header.h
  - This contains the struct definition for a file header
  - This is given by Professor Long in the resources repository.
- node.h
  - This file contains the interface for the node ADT.
  - This is given by Professor Long in the resources repository
- pq.h
  - This contains the interface to the ADT queue.
  - This is given by Professor Long in the resources repository
- code.h
  - This file will contain the code ADT interface.
  - This is given by Professor Long in the resources repository
- io.h
  - This file will contain the I/O module interface.
  - Given in the resources repository by Professor Long.
- stack.h
  - This file will contain the stack ADT interface.
  - Given in the resources repository by Professor Long.
- huffman.h
  - This file will contain the Huffman coding module interface. This file will be provided.
  - This code is given in the resources repository by Professor Long.
- node.c
  - This file will contain the implementation of the node ADT/
- pq.c
  - This file will contain the implementation of the priority queue ADT.
- code.c
  - This file will contain the implementation of the code ADT.
- io.c

- This file will contain the implementation of the input and output module.
- **huffman.c**
  - This file will have the implementation of the Huffman coding module interface.
- **Makefile**
  - Make the above program.
  - Links the above files
- **README.md**
  - A document in markdown syntax on the usage of the program.
- **DESIGN.pdf**
  - The pdf being viewed at the moment (this document).

## Notes

- Public key cryptography by RSA method.
  - First choose two large primes  $p$  and  $q$ .
  - Then calculate product,  $n = pq$
  - Very very computationally hard to do largest prime factorization
  - Computationally infeasible for  $n$  factor  $p$  and  $q$
  - Calculate the totient of  $n$ , which is  $(p-1)*(q-1)$
  - coprime aka relative prime mean gcd is 1
  - Then choose the public exponent
    - This  $e$  is some exponent
    - Calculate inverse of  $e$  which is  $e^{-1}$  modulo totient  $n$
    - This means  $d*e$  is equal to 1 modulo totient( $n$ )
    - 5 mod 13, then inverse of 5 is 8, 5 inverse mod 13 congruent 8 mod 13
    - This is because  $5*8 = 40 \text{ mod } 13$ , which is 1 mod 13.
  - Using this RSA is computed as
    - $d$  is private key
    - $n$  and  $e$  are public key
    - This is RSA
  - The encryption of a message  $m$  is  $m^e \text{ mod } (n)$
  - decryption of said encryption is encryption raised to  $d \text{ mod } n$ .
- Keygen makes the keys
- Encryptor encrypts files
- decryptor decrypts files
- The stuff should be gigantic
- c numbers small we need big so we use gmp
- We need to make a random state
- This is step one
- **randstate.c**
  - Using a seed we make a random number using mersenne twister algo
  - Then set seed using passed in seed

- then we can clear the random state
- use gmp lib funcs
- **Modular exponentiation**
  - pow mod
  - compute a base raised to exponent power modulo modulus, store computed result in out
- **Testing primality**
  - don't do it the normal deterministic way
  - Use a randomized algorithm with high probability of something being prime
  - Miller Rabin primality test
  - first find an s and r such that  $n-1 = 2^s \cdot r$  and r is odd
    - Basically, while not found, divide n-1 by  $2^s$  until r is odd, then when it is like that use s and 3.
  - take a k number of iterations and do it and return if.
  - is prime? take a number, check if it is prime using miller rabin
  - make\_prime, generates a prime that is at least bits number of bits long
- **GCD**
  - Euclidian algorithm of calculating gcd
  - We also need a function to compute modulo inverse of something and mod n
  - If no inverse found set i to 0
- **RSA lib**
  - Creates parts of a new RSA pub key, two large primes p and q, their product n and the public exponent e
    - First make primes p and q using make\_prime such that  $\log_2(n) \geq \text{nbits}$
    - Let the number of bits for p be a random num in the range nbits/4 to  $\frac{3}{4}$  nbits, the remaining bits will go to q, the number of miller rabin iterations is specified by iters.
    - next to the totient of p and q
    - We need a suitable public exponent e, generate random nums of around nbits using mpz\_urandomb() compute the gcd of each rand num and the computed totient stop the loop when found a coprime with totient
  - Write public key
    - Write the key to a pbfile for public file
  - Read pub key
    - Read key from file
  - make a private key and also wr the priv key
  - **RSA encrypt**
    - Take m to the e mod n and put that into c,
    - encrypt file
    - read in bytes of a file

- Read a block of bytes and turn the bytes into a number
  - We need to guarantee that  $m$  is less than  $n$
  - Calculate block size, that is  $\log_2(n) - 1/8$
  - Dynamically allocate an array of that block size
  - Pad a 0xFF to the beginning of the block
- Decrypt
  - Computing message  $m$  by decrypting ciphertext using priv key and public modulus  $n$
  - `rsa_decrypt` file
    - Decrypts the contents of infile writing the contents to the outfile
    - Decrypt in blocks, array of block size
    - scan in hexstring saving it as a new mpz
    - use `mpz_export` to make  $c$  back into bytes and store them in allocated block, let  $j$  be num bytes converted, you want to set the order param for `mpz_export()` to 1 for most significant word for, 1 for endian and 0 for nats
    - write out  $j-1$  bytes starting from index 1 of the block to outfile
    - This is because index 0 is prepended 0xFF don't print the 0xFF.
- `RSA_sign`
  - produce signature  $x$  by signing message with private key and public modulus, signing is  $s(m) = s = m^d \pmod n$
  - proves sender.
- $s$  is going to be  $x$  to the  $d \pmod n$
- $v = y$  to the  $e \pmod n$
- signing  $s = x$  to the  $d \pmod n$
- verify  $r$  is  $y$  to the  $e \pmod n$
- inverse operations of encrypt decrypt sign verify

## Pseudo

- `randstate_init`
  - Take the extern variable and initialize it with `gmp_randinit_mt()`
  - call `gmp_randseed_ui()` to set the seed
- `randstate_clear`
  - `gmp_randclear` the extern var.
- `pow_mod`
  - Does fast modular exponentiation, base raised to the exponent power module modulus and storing the computed result in out.
  - make a temp variable of 1
  - make a placeholder copy of  $a$
  - while  $d$  is greater than 0
    - IF  $2 \pmod 2$  is 1

- $v$  is  $v$  times  $p \bmod n$
    - Then,  $p$  is  $p^2 \bmod n$
    - Half  $d$
  - then return  $v$
  - Because gmp, instead of return, assign the value to the given output variable. This convention will remain true for the rest of the num\_theory functions
- is\_prime
  - find if  $n$  is prime in  $k$  numbers of iters.
  - First, find  $r$  and  $s$  such that  $n-1 = 2^s \cdot r$  and  $r$  is odd.
    - make a temp var for  $n-1$
    - To do this, make a while loop such that until an odd  $r$  is found, divide  $(n-1)$  by 2, starting at division by one.
    - Have a counter variable starting at zero count each iteration.
    - This counter variable is  $s$ .
    - Once odd  $n$  is found keep  $s$  as  $s$  and  $r$  as  $r$ .
  - For  $i$  in range of 1 to  $k$ (inclusive)
    - Generate a random number, modulo  $(n-4)+2$ , this makes a random number between 2 and  $n-2$
    - $y$  equals the power\_mod of  $a$ ,  $r$  and  $n$
    - if  $y$  is not 1 and  $y$  is not  $n-1$ 
      - new variable  $j$  is 1
      - while  $j$  is less than  $s$  minus one, and  $y$  is not equal to  $n-1$ 
        - $y$  is the power modulo of  $y^2 \bmod n$
        - if  $y == 1$ 
          - return false;
      - if  $y$  does not equal  $n-1$  return false
  - return true
- make\_prime
  - generates a new prime number that is at least bits number of bits, using mpz\_librandomb
  - check if it is prime using inters number of generations
  - if not prime, generate and check again.
- gcd
  - take the greatest common denominator of  $a$  and  $b$ , store it in  $d$  and return it.
  - while  $b$  is not zero
    - store  $b$  into  $t$
    - $b$  is  $a \bmod b$
    - store  $t$  into  $a$
  - return  $a$

- `mod_inverse`
  - take `i` and `a` and `n`, compute the modular inverse of `a` mod `n`,
  - Pseudo will be written assuming parallel assignment is allowed, every time we do so, in reality a temp variable is used.
  - Set `r` to `n`, `r'` to `a`
  - set `t` to 0 `t'` to 1
  - while `r'` is not 0
    - `q` is `r/r'`
    - parallel assign `r` to `r'`, `r`, to `r - q times r'`
    - `t = t' t' = t - q x t'`
  - if `r` is greater than 1
    - Set `i` to zero for no inverse
  - if `t` is less than 0
    - `t` is `t + n`
  - return `t` through `i`
- `rsa_make_pub`
  - Creates two new primes, `p` and `q`, the bits for `p` will be a random number in the range `nbits/4` and `3nbits/4`.
  - Next compute totient of `pq`
  - While (no prime exp)
    - generate a new random number
    - is the gcd of totient and the random number 1?
      - Then we found the prime exponent
- `rsa_write_pub`
  - Write the public key values, then the username to the outfile
  - Write each with newline after
- `rsa_make_priv`
  - compute the inverse modulo of the totient and `e`
- `rsa_write_priv`
  - Write `n` then `d`, both with trailing newline
- `rsa_read_priv`
  - Reads a private key from `pfile`, the format of a private key should be `n` then `d`, read both of these values
- `rsa_encrypt`
  - Takes a message as a number `m`, then compute the ciphertext as  $c = m^e \pmod{n}$
- `rsa_encrypt file`
  - Encrypts the contents of input file to the output file with mod `n`
  - Step one, set a block size var `k = floor of log2(n-1)/8`
    - do `n - 1`, find the then the size in base 2 aka the size in bits of the number
    - then divide that number by 8 and set it to `k`

- Dynamically allocate an array that can hold k bytes
    - array of type `uint8_t`
  - Set byte one of the array to `0xFF`
  - `fread()` from k -1 bytes from the infile and set j to the number of bytes returned
  - for i in range 1 to block put the bytes into the buffer
  - using `mpz_import`, convert the buffer to a number
  - pass this number to rsa encrypt with n and e
  - Take the encrypted number, write it, then a newline
- RSA decrypt
  - compute message from ciphertext, d, and n
  - powermod of m to the d mod n
  - return the result
- `rsa_decryptor`
  - Calculate block size, same as above
  - Make an array that can hold k bytes to hold the block
  - While the number returned by `mpz_inp_str` >0
    - use `mpz_inp_str` to read the hexstring
    - decrypts the given hexstring with RSA decrypt
    - `mpz_export` the hexstring into the allocated array
    - use `fwrite` to write from the array (minus the first element) to the number of bytes read from `mpz_inp_str` to outfile
- `rsa_sign`
  - Take message s and sign it by doing powermod of the message to the dth mod n
  - return that in s
- `rsa_verify`
  - if t is equal to s to the e mod n, return true, else return false

## KEYGEN

- parse command line opts, same as prev assignment
- `fopen()` the given output files for pub and priv key
- `fchmod()` the private key file permissions to 0600 also know as rw for use only
- using seed from cmd line initialize a random state with `randstate_init()`
- make the public and private keys using respective functions
- `getenv()` the current environment
- Convert the username with `mpz_set_str()` as a base 62
- Then use `rsa_sign` to compute the signature of the username
- Write the computed public and private keys to their respective files
- If verbose is enabled
  - print username, signature, p ,q, n, e, and d
  - Print these in decimal with number of bitse

- Close the files and `randstate_clear`

### **ENCRYPTOR**

- open the public key file using `fopen()`
- Read the entire public key file using `mpz_inp_str()` to read each field `n`, `e`, `s`, and `username`
- Convert the username that was read in to an `mpz_t`
- Verify the signature with `rsa_verify`
- Encrypt the file using `rsa_encrypt` file

### **DECRYPTOR**

- Open the private key file, then read the private key
- Decrypt the file with a call to `rsa_decrypt`