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National University of Computer and Emerging Sciences, Lahore Campus							
SOURCES OF THE PROPERTY OF THE	Course: Program: Duration: Paper Date: Section:	Advanced Database Concepts BS(Computer Science) 60 Minutes Fri 13-Apr-18 CS	Course Code: Semester: Total Marks: Weight Page(s):	CS451 Spring 2018 30 12.5%			

Midterm-II

Instruction/Notes:

Scratch sheet can be used for rough work however, all the questions and steps are to be shown on question paper. You may use backside of paper. No extra/rough sheets should be submitted with question paper. You will not get any credit if you do not show proper working, reasoning and steps as asked in question statements. Calculators are ALLOWED.

Q1. (8 points) Assume a relation R (A, B, C) is given; R is stored as an ordered file (un-spanned) on non-key field C and contains 100,000 records. Attributes A, B and C need 10 bytes of storage each (i.e. record size= 30), and blocks have a size of 512 Bytes. Each A value occurs at an average 5 times in the database, each B value occurs 50 times in the database, and each C value occurs 5000 times in the database. Assume there is no index structure exists. Estimate the number of block fetches needed to compute the following queries (where Ca, Cb, Cc1 and Cc2 are integer constants):

- a) SELECT B, C FROM R WHERE  $A = C_a$ ;
- **b)** SELECT COUNT(\*) FROM R WHERE  $B = C_b$ ;

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- c) SELECT A, B FROM R WHERE  $C = C_{c1}$ ;
- **d)** SELECT A, B FROM R WHERE  $C = C_{c1}$  OR  $C = C_{c2}$ ;

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**Q2.** (1+1+1+1+4=8 points) Consider a disk with block size B=1024 bytes. A block pointer is P=6 bytes long, and a record pointer is  $P_R=7$  bytes long. A file has r=1,000,000 EMPLOYEE records of fixed-length. Record length R is 115 bytes long and DEPTCODE field is 15 bytes long.

Suppose the file is ordered by the non-key field DEPTCODE and we want to construct a clustering index on DEPTCODE that uses block anchors (every new value of DEPTCODE starts at the beginning of a new block). Assume there are *500* distinct values of DEPTCODE, and that the EMPLOYEE records are evenly distributed among these values. Calculate:

- a) The index blocking factor (bfr<sub>i</sub>).
- b) The number of first-level index entries (r1) and the number of first-level index blocks (b1).
- c) The number of levels needed (x) if we make it a multi-level index.
- **d)** The total number of blocks required by the multi-level index (b<sub>i</sub>).
- **e)** The number of block accesses needed to search for and retrieve all records in the file having a specific DEPTCODE value using the clustering index (assume that multiple blocks in a cluster are either contiguous or linked by pointers).

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Q3. (6 points) Suppose you are building an extensible hash index on a table of 100,000 rows. Key values are 8 bytes, a pointer (block/record) to a row is 8 bytes, and a disk block is 2048 bytes. Assume all keys are distinct.  a) What is the (lowest possible) global depth? Provide valid reasons.  b) What is the average occupancy of a bucket, assuming all buckets have a local depth equal to the global depth from part (a)? Justify your answer.					

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<b>Q4.</b> $(2+3+3=8 \text{ points})$ Assume that you have just built a dense 1,000,000 records. The key field for this B <sup>+</sup> -tree index is a 40-byte stri 10-byte values. The size of one disk page is 1000 bytes. The index was at each level were filled up as much as possible.	ng, and it is a candidate key. Pointers (Record/block) are
<ul><li>a) How many levels does the resulting tree have?</li><li>b) For each level of the tree, how many nodes are at that level?</li><li>c) How many levels would the resulting tree have with all pages 70 per</li></ul>	ercent full?

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