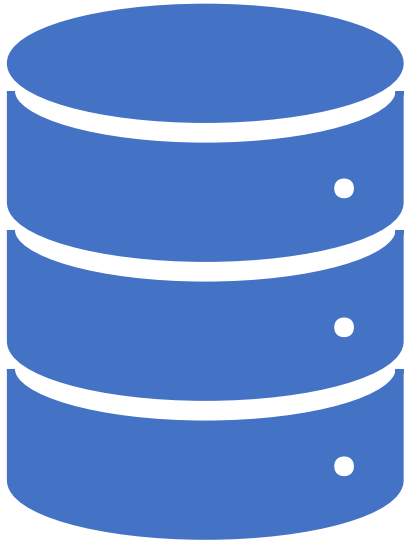


Hashing, Hash Function and Hash Table

What is Hashing



- Hashing is a technique that is used to uniquely identify a specific object from a group of similar objects.
- Examples of how hashing is used in our lives:
 - In universities, each student is assigned a unique roll number that can be used to retrieve information about them.
 - In libraries, each book is assigned a unique number that can be used to determine information about the book
- **Hashing** is the solution that can be used in almost all such situations and performs extremely well compared to above data structures like Array, Linked List, BST in practice

What is Hashing (Cond...)

- With hashing we get $O(1)$ search time on average and $O(n)$ in worst case.
- The idea is to use ***hash function*** that converts a given input number or any other key to a smaller number and uses the small number as index in a table called ***hash table***.
- Hash Function
 - Hash function maps a big number or string to a small integer that can be used as index in hash table
- Hash Table
 - An array that stores pointers to records corresponding to a given input number



0

1

2

3

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Mia

M

77

i

105

a

97

279

4

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0

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2

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Mia

M

77

i

105

a

97

279

4

				Mia						
0	1	2	3	4	5	6	7	8	9	10

Mia	M	77	i	105	a	97	279	4
Tim	T	84	i	105	m	109	298	1

	Tim			Mia						
0	1	2	3	4	5	6	7	8	9	10

Mia	M	77	i	105	a	97	279	4
Tim	T	84	i	105	m	109	298	1
Bea	B	66	e	101	a	97	264	0
Zoe	Z	90	o	111	e	101	302	5
Jan	J	74	a	97	n	110	281	6
Ada	A	65	d	100	a	97	262	9
Leo	L	76	e	101	o	111	288	2

Bea	Tim	Leo		Mia	Zoe	Jan			Ada	
0	1	2	3	4	5	6	7	8	9	10

Mia	M	77	i	105	a	97	279	4
Tim	T	84	i	105	m	109	298	1
Bea	B	66	e	101	a	97	264	0
Zoe	Z	90	o	111	e	101	302	5
Jan	J	74	a	97	n	110	281	6
Ada	A	65	d	100	a	97	262	9
Leo	L	76	e	101	o	111	288	2
Sam	S	83	a	97	m	109	289	3
Lou	L	76	o	111	u	117	304	7
Max	M	77	a	97	x	120	294	8
Ted	T	84	e	101	d	100	285	10

Bea	Tim	Leo	Sam	Mia	Zoe	Jan	Lou	Max	Ada	Ted
0	1	2	3	4	5	6	7	8	9	10

Index number = $\text{sum ASCII codes} \bmod \text{size of array}$

Bea	Tim	Leo	Sam	Mia	Zoe	Jan	Lou	Max	Ada	Ted
0	1	2	3	4	5	6	7	8	9	10

$$\text{Ada} = (65 + 100 + 97) = 262$$

Find Ada

$$262 \text{ Mod } 11 = 9$$

myData = Array(9)

Bea	Tim	Leo	Sam	Mia	Zoe	Jan	Lou	Max	Ada	Ted
0	1	2	3	4	5	6	7	8	9	10

Bea 27/01/1941 English Astronomer	Tim 08/06/1955 English inventor	Leo 31/12/1945 American Mathematician	Sam 27/04/1791 American inventor	Mia 20/02/1986 Russian Space Station	Zoe 19/06/1978 American Actress	Jan 13/02/1956 Polish Logician	Lou 27/12/1822 French Biologist	Max 23/04/1858 German Physicist	Ada 10/12/1815 English Mathematician	Ted 17/06/1937 American Philosopher
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0

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Hashing Algorithm

- Calculation applied to a key to transform it into an address
- For numeric keys, divide the key by the number of available addresses, n , and take the remainder

$$\text{address} = \text{key} \text{ Mod } n$$

- For alphanumeric keys, divide the sum of ASCII codes in a key by the number of available addresses, n , and take the remainder
- Folding method divides key into equal parts then adds the parts together
 - The telephone number 01452 8345654, becomes $01 + 45 + 28 + 34 + 56 + 54 = 218$
 - Depending on size of table, may then divide by some constant and take remainder

Understand Collision

Mia	M	77	i	105	a	97	279	4
Tim	T	84	i	105	m	109	298	1
Bea	B	66	e	101	a	97	264	0
Zoe	Z	90	o	111	e	101	302	5

Bea	Tim			Mia	Zoe					
0	1	2	3	4	5	6	7	8	9	10

Understand Collision

Mia	M	77	i	105	a	97	279	4
Tim	T	84	i	105	m	109	298	1
Bea	B	66	e	101	a	97	264	0
Zoe	Z	90	o	111	e	101	302	5
Sue	S	83	u	117	e	101	301	4

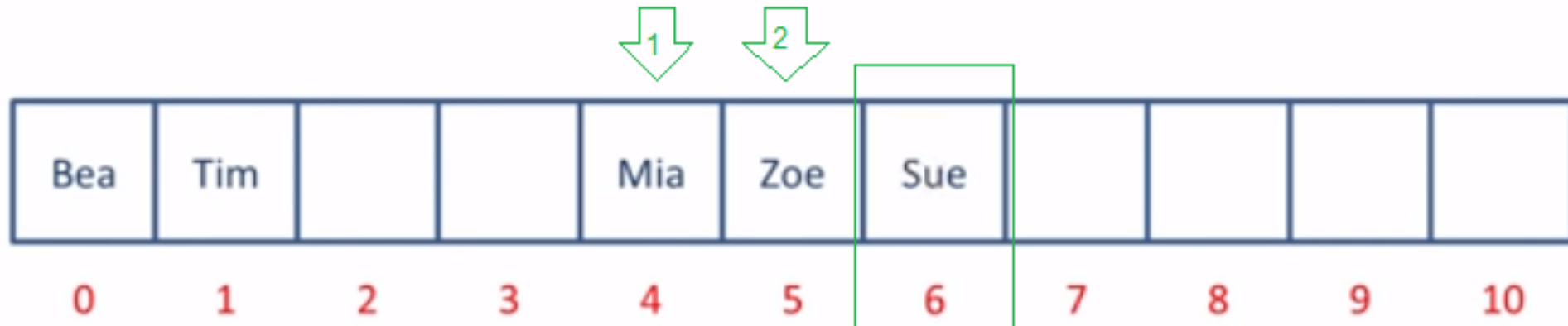
Bea	Tim			Mia	Zoe					
0	1	2	3	4	5	6	7	8	9	10

Collision Handling

- Irrespective of how good a hash function is, collisions are bound to occur.
- Therefore, to maintain the performance of a hash table, it is important to manage collisions through various collision resolution techniques.
 - Open addressing
 - Linear Probing
 - Quadratic Probing
 - Double Hashing
 - Closed addressing

Open Addressing (Linear)

Mia	M	77	i	105	a	97	279	4
Tim	T	84	i	105	m	109	298	1
Bea	B	66	e	101	a	97	264	0
Zoe	Z	90	o	111	e	101	302	5
Sue	S	83	u	117	e	101	301	4



Mia	M	77	i	105	a	97	279	4
Tim	T	84	i	105	m	109	298	1
Bea	B	66	e	101	a	97	264	0
Zoe	Z	90	o	111	e	101	302	5
Sue	S	83	u	117	e	101	301	4
Len	L	76	e	101	n	110	287	1



Mia	M	77	i	105	a	97	279	4
Tim	T	84	i	105	m	109	298	1
Bea	B	66	e	101	a	97	264	0
Zoe	Z	90	o	111	e	101	302	5
Sue	S	83	u	117	e	101	301	4
Len	L	76	e	101	n	110	287	1
Moe	M	77	o	111	e	101	289	3
Lou	L	76	o	111	u	117	304	7
Rae	R	82	a	97	e	101	280	5
Max	M	77	a	97	x	120	294	8
Tod	T	84	o	111	d	100	295	9

Bea	Tim	Len	Moe	Mia	Zoe	Sue	Lou	Rae	Max	Tod
0	1	2	3	4	5	6	7	8	9	10

Find Rae

$$\text{Rae} = (82 + 97 + 101) = 280$$

$$280 \text{ Mod } 11 = 5$$

myData = Array(5)

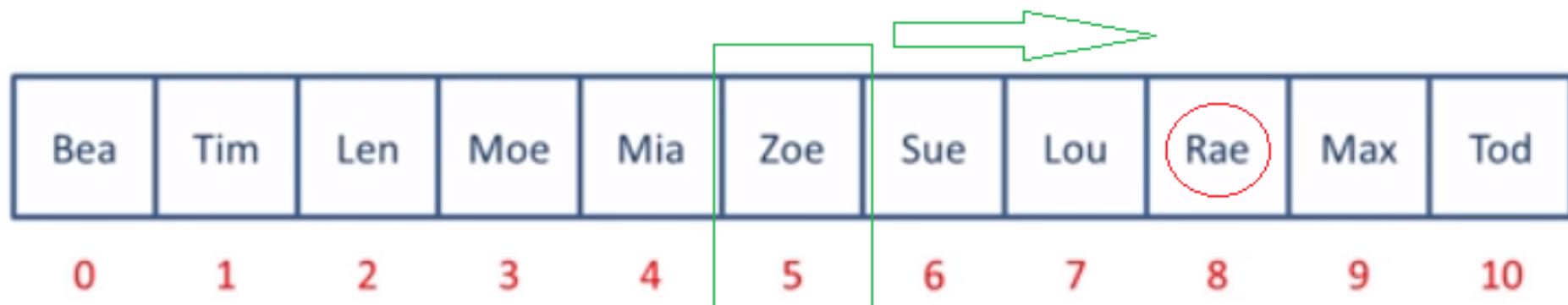
Bea	Tim	Len	Moe	Mia	Zoe	Sue	Lou	Rae	Max	Tod
0	1	2	3	4	5	6	7	8	9	10

Find Rae

$$\text{Rae} = (82 + 97 + 101) = 280$$

$$280 \text{ Mod } 11 = 5$$

myData = Array(5)

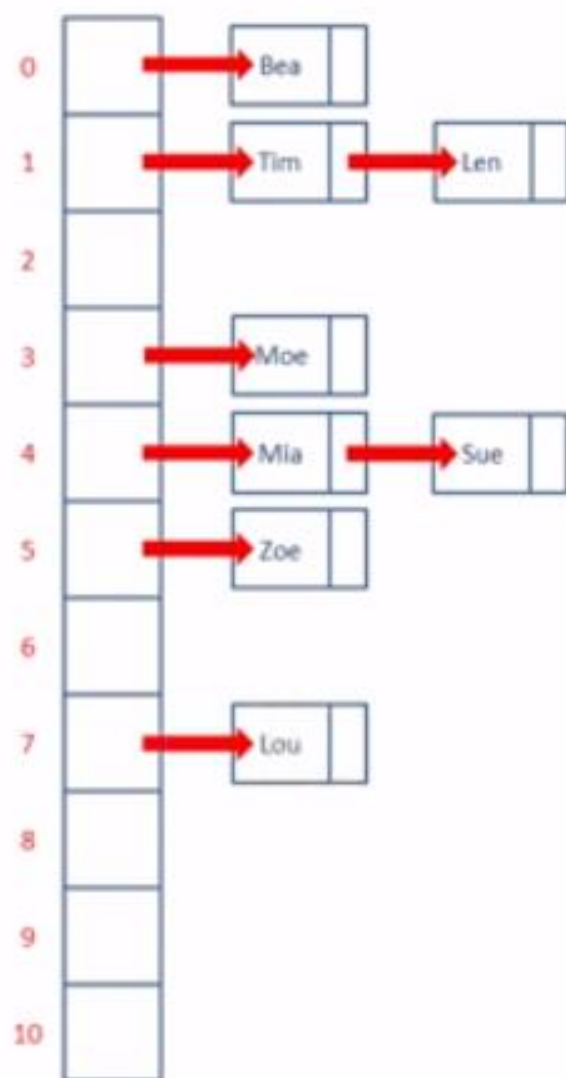


Bea	Tim	Len	Moe	Mia	Zoe	Sue	Lou	Rae	Max	Tod
0	1	2	3	4	5	6	7	8	9	10

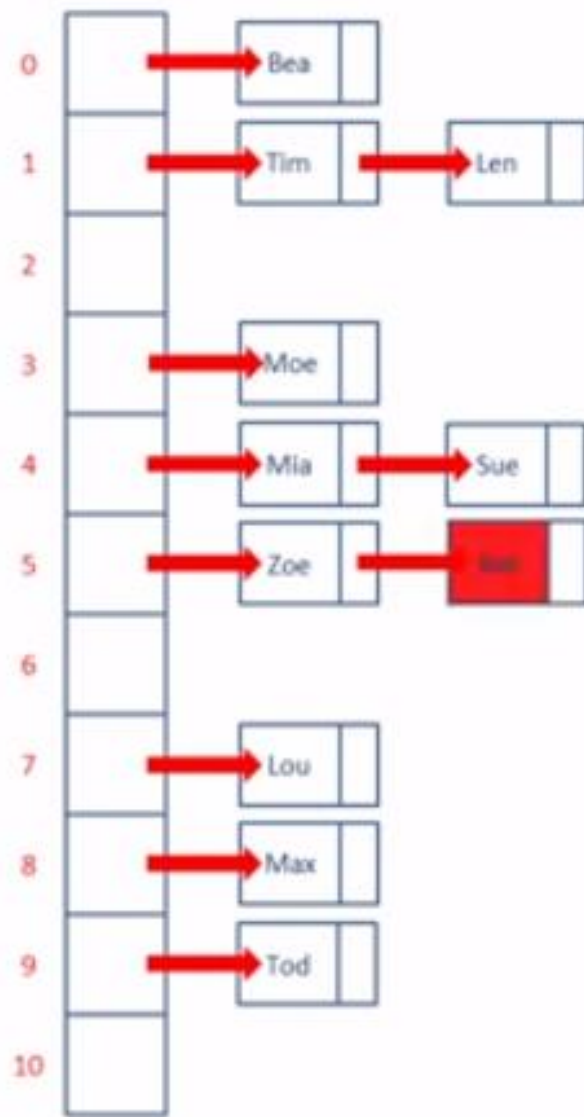
Closed Addressing (Non-Linear)



Mia M 77 i 105 a 97 279 4



Mia	M	77	i	105	a	97	279	4
Tim	T	84	i	105	m	109	298	1
Bea	B	66	e	101	a	97	264	0
Zoe	Z	90	o	111	e	101	302	5
Sue	S	83	u	117	e	101	301	4
Len	L	76	e	101	n	110	287	1
Moe	M	77	o	111	e	101	289	3
Lou	L	76	o	111	u	117	304	7



Find Rae $280 \text{ Mod } 11 = 5$

myData = Array(5)

Objectives of Hash Function

- Minimize collisions
- Uniform distribution of hash values
- Easy to calculate
- Resolve any collisions