

String Alignment 1

Dynamic Programming

How similar are these strings?

TGGCAGTCCCAGAAAGGACTCTCCTC

TGCACAGTCCAGACACGGTCTCGTT

How similar are these strings?

TGGCAGTCCCAGAAGGACTCTCCTC

TGCACAGTCCAGACACGGTCTCGTT

- Phylogenetic ancestry
- Mutations due to disease
- Understanding gene function

String Alignment

G	C	A	G	T	C	C	G	A	C
G	C	G	T	C	T	G	A	C	T

#Matches: 2

#Mismatches: 8

String Alignment

G	C	A	G	T	C	C	G	A	C	
G	C		G	T	C	T	G	A	C	T

#Matches: 8

#Mismatches: 3

String Alignment

G	C	A	G	T	C	C	G	A	C	
G	C		G	T	C	T	G	A	C	T

A: Deleted

C->T

T: Inserted

Mutations: Errors in DNA replication? Disease?

String Similarity

G	C	A	G	T	C	C	G	A	C	
G	C		G	T	C	T	G	A	C	T

Minimum #Mismatches among all alignments

String Similarity

G	C	A	G	T	C	C	G	A	C	
G	C		G	T	C	T	G	A	C	T
M	M	D	M	M	M	R	M	M	M	I

Edit Distance = Minimum #
Insertions + Deletions + Replacements

Edit Distance

GCAGTCCGAC

-> GCGTCCGAC

-> GCGTCTGAC

-> GCGTCTGACT

Minimum # Insertions + Deletions + Replacements

Edit Distance

Problem:

Input: Two strings S_1 and S_2

Output: Edit Distance
and

Optimal Edit Transcript/Alignment

Review:

$ED(\text{CLEAR}, \text{READS}) =$

$ED(\text{ACCTGCAA}, \text{CTGCAAG}) =$

$ED(\epsilon, \text{GACCT}) =$

Dynamic Programming

Step 1: Identify subproblems to be solved.

Step 2: Develop a recurrence relation for each subproblem.

Step 3: Solve subproblems in bottom-up/tabular computation.

Step 4: Traceback to get solution.

Step 1: Subproblems

$D(i,j)$ = Edit Distance between $S1[1,...,i]$ and $S2[1,...,j]$

$S1$: TACCGCA $S2$: ACCGTAC

$D(4,3) = ED(TACC, ACC) = 1$

$D(2,2) = ED(TA, AC) = 2$

$D(7,7) = ED(TACCGCA, ACCGTAC)$

Step 2: Recurrence

S₁:TACCGCA S₂: ACCGTAC

Base Cases:

$$D(o,j)=j$$

$$D(i,o)=i$$

Step 2: Recurrence

S1:TACCGCA S2: ACCGTAC

$ED(TAC, ACC) = ED(TA, AC)$

$ED(TACCGC, ACCGT) = 1 + \text{Min}\{$
 $ED(TACCGC, ACCG),$
 $ED(TACCG, ACCGT),$
 $ED(TACCG, ACCG)\}$

Step 2: Recurrence

$$ED(TAC, ACC) = ED(AC, TA)$$

$$\text{If } S1[i] = S2[j], \text{ then } D(i, j) = D(i-1, j-1)$$

Step 2: Recurrence

$$\text{ED}(\text{TACCGC}, \text{ACCGT}) = 1 + \text{Min}\{\text{ED}(\text{TACCGC}, \text{ACCG}), \text{ED}(\text{TACCG}, \text{ACCGT}), \text{ED}(\text{TACCG}, \text{ACCG})\}$$

If $S_1[i] \neq S_2[j]$, then:

$$D(i,j) = 1 + \text{Min}\{D(i,j-1), D(i-1,j), D(i-1,j-1)\}$$

Step 2: Recurrence

If $S_1[i]=S_2[j]$, then

$$D(i,j)=D(i-1,j-1)$$

Else

$$D(i,j)=1+\text{Min}\{D(i,j-1), D(i-1,j), D(i-1,j-1)\}$$

Step 3:
Tabular
Computation

	$D(i-1,j-1)$	$D(i-1,j)$		
	$D(i,j-1)$	$D(i,j)$		

Step 3:
Tabular
Computation

	€	A	C	C	G	T	A	C
€								
T								
A								
C								
C								
G								
C								
A								

Step 3:
Tabular
Computation

	€	A	C	C	G	T	A	C
€	0	1	2	3	4	5	6	7
T	1	1	2	3	4	4	5	6
A	2	1						
C	3	2						
C	4	3						
G	5	4						
C	6	5						
A	7	6						

Step 3:
Tabular
Computation

	€	A	C	C	G	T	A	C
€	0	1	2	3	4	5	6	7
T	1	1	2	3	4	4	5	6
A	2	1	2	3	4	5	4	5
C	3	2	1	2	3	4	5	4
C	4	3	2	1	2	3	4	5
G	5	4	3	2	1	2	3	4
C	6	5	4	3	2	2	3	3
A	7	6	5	4	3	3	2	3

Step 4:
Traceback

	€	A	C	C	G	T	A	C
€	0	1	2	3	4	5	6	7
T	1	1	2	3	4	4	5	6
A	2	1	2	3	4	5	4	5
C	3	2	1	2	3	4	5	4
C	4	3	2	1	2	3	4	5
G	5	4	3	2	1	2	3	4
C	6	5	4	3	2	2	3	3
A	7	6	5	4	3	3	2	3

Step 4:
Traceback

	€	A	C	C	G	T	A	C
€	0	1	2	3	4	5	6	7
T	1	1	2	3	4	4	5	6
A	2	1	2	3	4	5	4	5
C	3	2	1	2	3	4	5	4
C	4	3	2	1	2	3	4	5
G	5	4	3	2	1	2	3	4
C	6	5	4	3	2	2	3	3
A	7	6	5	4	3	3	2	3

Step 4:
Traceback

	€	A	C	C	G	T	A	C
€	0	1	2	3	4	5	6	7
T	1	1	2	3	4	4	5	6
A	2	1	2	3	4	5	4	5
C	3	2	1	2	3	4	5	4
C	4	3	2	1	2	3	4	5
G	5	4	3	2	1	2	3	4
C	6	5	4	3	2	2	3	3
A	7	6	5	4	3	3	2	3

Transcript:
DMMMMRMI

Transcript:
DMMMMRMI

S1: TACCGCA, S2: ACCGTAC

T	A	C	C	G	C	A	€
€	A	C	C	G	T	A	C

Algorithm Sketch

1. $D(i,0)=i$, $D(0,j)=j$ for all i,j
2. Set $row=col=0$
3. While ($row < m$ or $col < n$)

 If ($row < m$) $row=row+1$

 If ($col < n$) $col=col+1$

Algorithm Sketch

```
for (j=row to n)  
    Compute Cell(row,j) and  
    Parent(row,j) and Transcript(row,j)
```

```
for (i=col+1 to m)  
    Compute Cell(i,col) and  
    Parent(i,col) and Transcript(row,j)
```

End While

Algorithm Sketch

4. $\text{Cell} = (m, n)$

$\text{Transcript} = \epsilon$

While ($\text{Parent}(\text{Cell})$ is not in zeroth row/column)

$\text{Transcript} = \text{Transcript} + \text{Transcript}(\text{Cell})$

$\text{Cell} = \text{Parent}(\text{Cell})$

Time and Space Complexity

Space=Size of Table= $O(mn)$

Time for Table Computation: $O(mn)$

Time for Traceback: $O(m+n)$

Space Complexity

Space=Size of Table= $O(mn)$

Can we do better?