The OpenCL program to perform multiply and multiply and add was carried out on "rabbit machine (Intel Xeon)", which is running Cent-OS and has NVIDIA-Titan Black graphics card installed in it. The experimental results of Multiply and Multiply and Add are tabulated as shown below. (Note: Arrow horizontal in tables indicate local size and Vertical arrow pointing downwards indicate Global size)

Multiply:

4096000

2.497

3.895

Table 1: Performance parameters of Multiply with respect varying Global and Local work size. All Performance Parameters in GigaMults/Sec

		8	16	32	64	128	256	512
10	000	0.017	0.032	0.027	0.033	0.032	0.034	0.036
20	000	0.035	0.035	0.054	0.061	0.056	0.051	0.036
40	000	0.05	0.051	0.046	0.146	0.109	0.108	0.133
80	000	0.112	0.099	0.099	0.109	0.205	0.243	0.212
160	000	0.306	0.209	0.231	0.257	0.233	0.428	0.422
320	000	0.418	0.552	0.466	0.559	0.436	0.446	1.013
640	000	0.664	0.817	1.109	1.192	1.175	0.815	1.151
1280	000	0.399	0.446	0.542	0.433	0.52	0.541	0.543
2560	000	0.79	0.258	1.175	1.16	0.881	0.941	0.919
5120	000	1.194	1.356	1.801	1.801	1.695	1.417	1.778
10240	000	1.752	0.932	2.789	3.037	3.236	2.654	3.15
20480	000	2.157	3.151	4.2	6.05	5.164	5.951	5.12

5.903

7.732

7.81

7.689

7.943

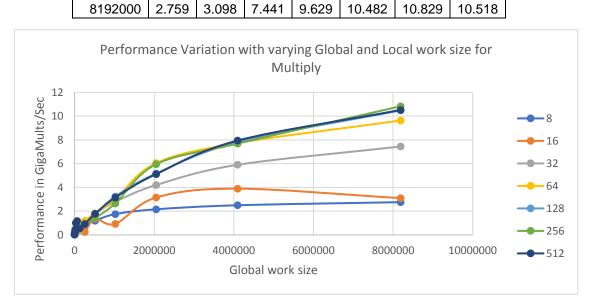


Figure 1: Graphical Representation of Performance measured for multiply and store using OpenCL, with constant local size

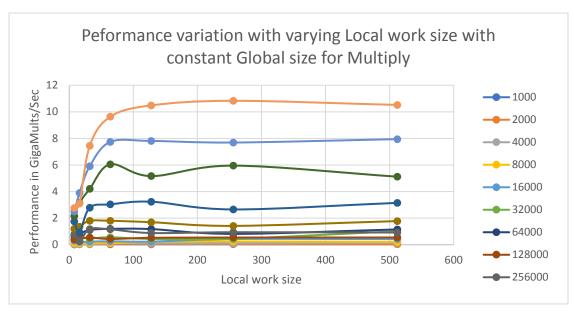


Figure 2: Graphical Representation of Performance measured for multiply and store using OpenCL, with constant global size

From the table 1 it is evident that the keeping global size as constant and varying the local work sizes given better performance than its predecessor, for instance for global size of 2048000 and local size of 32 the performance observed is 4.2 GigaMults/Sec, same simulation carried out with local size of 256 gives 5.951 GigaMults/Sec. This is mainly because the number of computational threads needed at GPU is very less for local work size of greater number this makes it more efficient than its predecessor local work sizes. Graphical representation of the Table 1 is shown in Figure 1.

Multiply and Add:

The performance measurement of multiply and add is like that of multiply only addition that is being made is d=a*b+c, unlike just a*b is prior measurement. The wall clock measurements start and stop time are still the same as that of multiply only. The performance increases with increasing global size, and much better performance is observed for varying local size with keeping constant global size.

Table 2 : Performance parameters of Multiply and Add with respect varying Global and Local work size in GigaMults/Sec



	8	16	32	64	128	256	512
1000	0.014	0.032	0.035	0.036	0.037	0.035	0.035
2000	0.05	0.041	0.055	0.054	0.077	0.072	0.075
4000	0.072	0.067	0.095	0.16	0.159	0.153	0.158
8000	0.179	0.201	0.216	0.203	0.349	0.227	0.367
16000	0.418	0.472	0.446	0.314	0.326	0.537	0.422
32000	0.644	0.737	0.78	0.779	0.79	0.811	1.223

64000	0.901	1.241	1.075	1.541	1.265	1.045	1.37
128000	1.255	2.079	2.487	1.903	2.921	2.077	2.932
256000	1.844	3.139	3.43	3.726	4.662	4.516	4.824
512000	2.172	4.056	5.776	6.014	7.071	7.216	5.983
1024000	2.543	4.728	7.295	9.199	8.673	9.41	8.429
2048000	2.753	5.072	8.186	10.722	11.587	11.328	11.392
4096000	2.842	5.303	8.583	11.73	12.907	12.812	12.897
8192000	2.883	5.544	9.296	12.655	13.875	13.352	13.224

Table 3: Performance measurement in GigaMults/Sec for Mutiply and Add

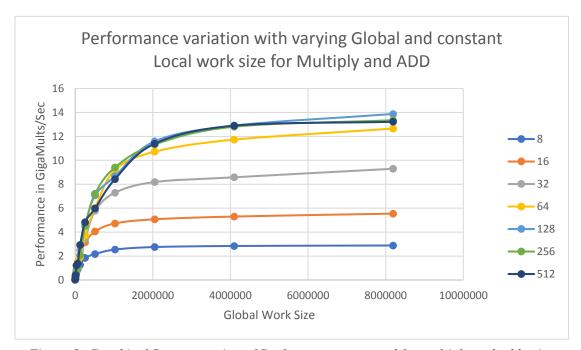


Figure 3: Graphical Representation of Performance measured for multiply and add using OpenCL, with constant local size

Multiply and Reduction:

The simulation to perform multiply and sum i.e. sum+=a*b, was run on the same machine as that of prior simulation. The strange behavior observed in this case was that, there is a saturation point for parallelizing code which much less than that of multiply and multiply and add. For a local size of 64 and global size of 64000 we hit *Wait: clWaitForEvents failed* which is observed immediately after enqueuing the buffer. This is observed for any combinations higher

than this. I could simulate for local work sizes of 8,16,32 for global size varying from 1K-8M and results are presented in Table 3 and graphical representation of the same in Figure 3. From the graph, it is evident that the varying global and local work sizes till saturation point can be done to get better performance out the of the system. The maximum performance observed in my simulation for multiply and reduction is 4.6 GigaMults/Sec with a global work size of 8M and local work size of 32

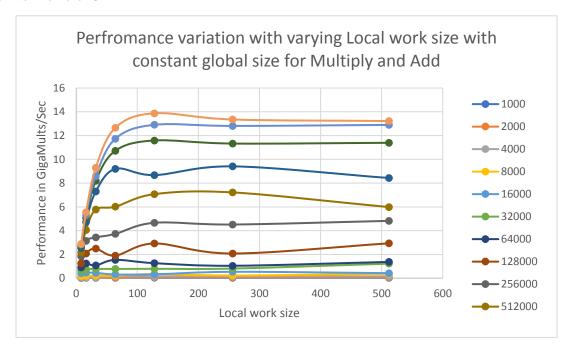


Figure 4: Graphical Representation of Performance measured for multiply and add using OpenCL, with constant global size

Table 3: Performance measurement in GigaMults/Sec for Mutiply and Add



	8	16	32
1000	0.004	0.029	0.035
2000	0.01	0.01	0.069
4000	0.02	0.017	0.019
8000	0.036	0.041	0.039
16000	0.075	0.056	0.078
32000	0.131	0.119	0.136
64000	0.236	0.261	0.237
128000	0.42	0.436	0.52
256000	0.648	0.807	0.877
512000	1.003	1.213	1.448
1024000	1.365	1.665	2.467

2048000	1.49	2.243	3.215
4096000	1.606	2.528	3.826
8192000	1.718	2.712	4.274

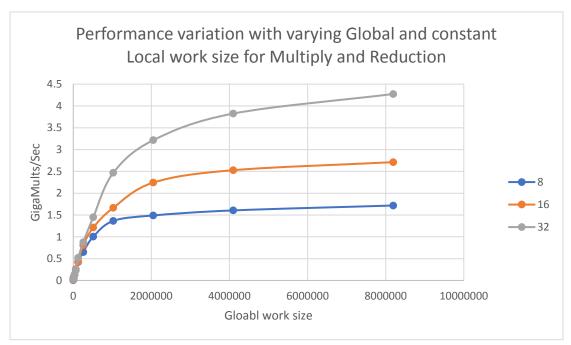


Figure 5: Graphical Representation of Performance measured for multiply and reduction using OpenCL, with constant local size

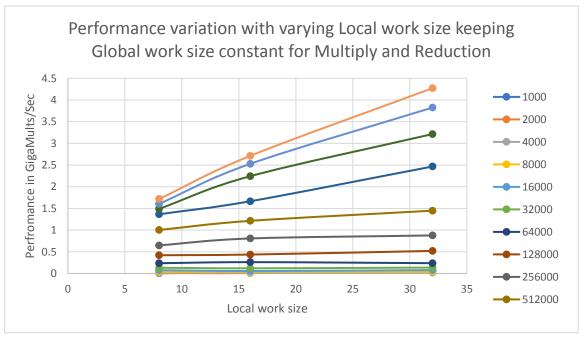


Figure 6: Graphical Representation of Performance measured for multiply and reduction using OpenCL, with constant global size