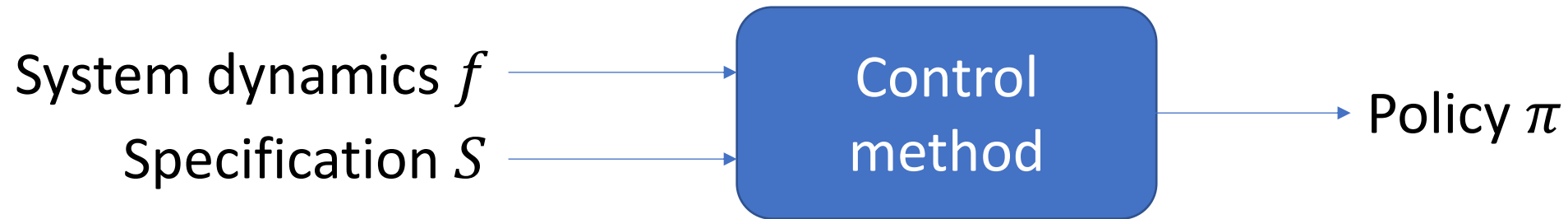


# A Simulation Preorder for Koopman-like Lifted Control Systems

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# Goal

Given a system and a specification, find a policy such that the closed-loop system satisfies the specification.



$$f \models_{\pi} S$$

# Systems and specifications

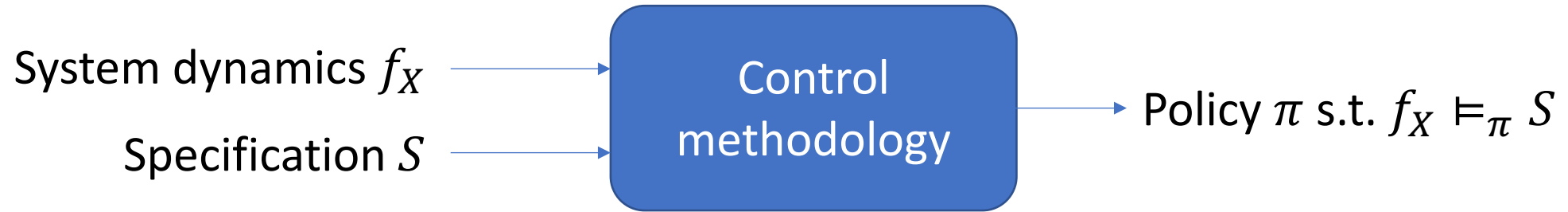
- A discrete-time system

$$x(t + 1) = f_X(x(t), u(t))$$

with state  $x(t) \in X$  and  $u(t) \in U$

- A **specification** is a set of sequences of  $(x, u)$  pairs:  $S \subseteq (X \times U)^\infty$
- The **behavior** of a system under a policy  $\pi$  is the set
$$\mathbf{B}_\pi[f_X] := \{(\mathbf{x}, \mathbf{u}) \mid (\mathbf{x}, \mathbf{u}) \text{ is a max. solution \& } u(t) = \pi(x(0), \dots, x(t))\}$$
- The specification  $S$  is **satisfied** by the system  $f_X$  under the policy  $\pi$  if
$$\mathbf{B}_\pi[f_X] \subseteq S.$$
 This is written  $f_X \models_\pi S$

# Few words on simulation relations



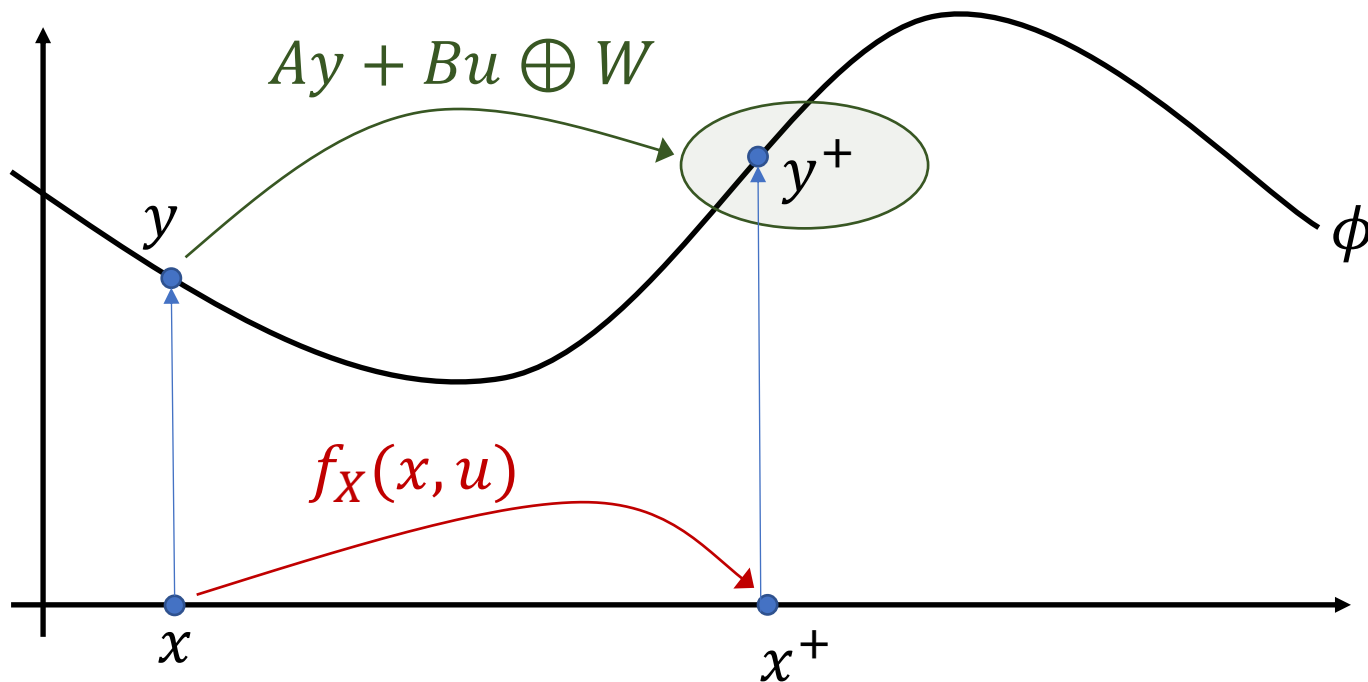
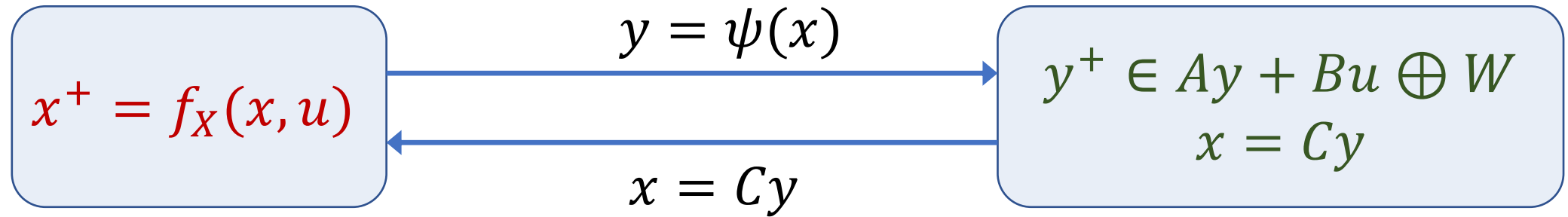
Simulation-based approach:

1. Define a notion of simulation between systems  $f_X \preccurlyeq f_Y$
2. Prove that  $f_X \preccurlyeq f_Y$  implies  $\mathbf{B}_{\pi}[f_X] \subseteq \mathbf{B}_{\pi}[f_Y]$

Why is it useful?

1. Replace the system  $f_X$  by a simpler (e.g., affine) one  $f_Y$  s.t.  $f_X \preccurlyeq f_Y$
2. Synthesize a policy  $\pi$  s.t.  $f_Y \models_{\pi} S$
3. It follows that  $f_X \models_{\pi} S$

# Koopman approach



$$\psi(f_X(x, u)) \in A\psi(x) + Bu \oplus W$$

$$\psi(x) = \begin{bmatrix} x \\ \phi(x) \end{bmatrix}$$

$$C = \begin{bmatrix} I & 0 \end{bmatrix}$$

Question: How to choose  $\psi$  ?

# Lifted systems

A **lifted system** is a tuple  $LS_Y = (X, U, \psi_Y, f_Y, g_Y)$

- $X$  is a set of outputs
- $U$  is a set of inputs
- $\psi_Y: \mathbb{R}^{n_X} \rightarrow \mathbb{R}^{n_Y}$  is a lifting function
- $f_Y: \mathbb{R}^{n_Y} \times U \rightrightarrows \mathbb{R}^{n_Y}$  is a set-valued dynamics
- $g_Y: \mathbb{R}^{n_Y} \rightarrow \mathbb{R}^{n_X}$  is an output map such that  $g_Y(\psi_Y(x)) = x$

A **solution** of  $LS_Y$  is a tuple  $(\mathbf{x}, \mathbf{u}, \mathbf{y}) \in X^{[0,T[} \times U^{[0,T[} \times (\mathbb{R}^{n_Y})^{[0,T[}$  s.t.

- $y(0) = \psi_Y(x(0))$
- $y(t+1) \in f_Y(y(t), u(t))$
- $x(t) = g_Y(y(t))$

The **behavior** of  $LS_Y$  under a policy  $\pi$  is the set

$$\mathbf{B}_\pi[LS_Y] = \{ (\mathbf{x}, \mathbf{u}) \mid \exists \mathbf{y} \text{ s.t. } (\mathbf{x}, \mathbf{u}, \mathbf{y}) \text{ is a max. solution \& } u(t) = \pi(x(0), \dots, x(t)) \}$$

# Important classes of lifted systems

- **Unlifted** (i.e., “classical”) systems

$$x(t + 1) = f_X(x(t), u(t))$$

are lifted systems with  $n_Y = n_X$  and  $\psi_Y = g_Y = id$

- **Affine** lifted systems

$$\begin{aligned} y(t + 1) &\in Ay(t) + Bu(t) \oplus W \\ x(t) &= Cy(t) \end{aligned}$$

for which linear control methods can be used

- **Picewise affine** lifted systems

# Simulation relation for lifted systems

$LS_Y$  is **simulated** by  $LS_Z$  (denoted  $LS_Y \preccurlyeq LS_Z$ ) if there exists a set-valued map  $\rho: \mathbb{R}^{n_Z} \rightrightarrows \mathbb{R}^{n_Y}$  such that

- $\forall x \in X: \psi_Y(x) \in \rho(\psi_Z(x))$  *(Relation between liftings)*
- $\forall (z, u) \in \mathbb{R}^{n_Z} \times U: f_Y(\rho(z), u) \subseteq \rho(f_Z(z, u))$  *(between dynamics)*
- $\forall z \in \mathbb{R}^{n_Z}: g_Y(\rho(z)) \subseteq \{g_Z(z)\}$  *(between outputs)*



# Simulation implies behavioral inclusion

## THEOREM

Given two lifted systems  $LS_Y$  and  $LS_Z$  and a policy  $\pi$ , if  $LS_Y$  is simulated by  $LS_Z$ , then the closed loop behavior of  $LS_Y$  under  $\pi$  is included in the closed loop behavior of  $LS_Z$  under  $\pi$ , i.e.,

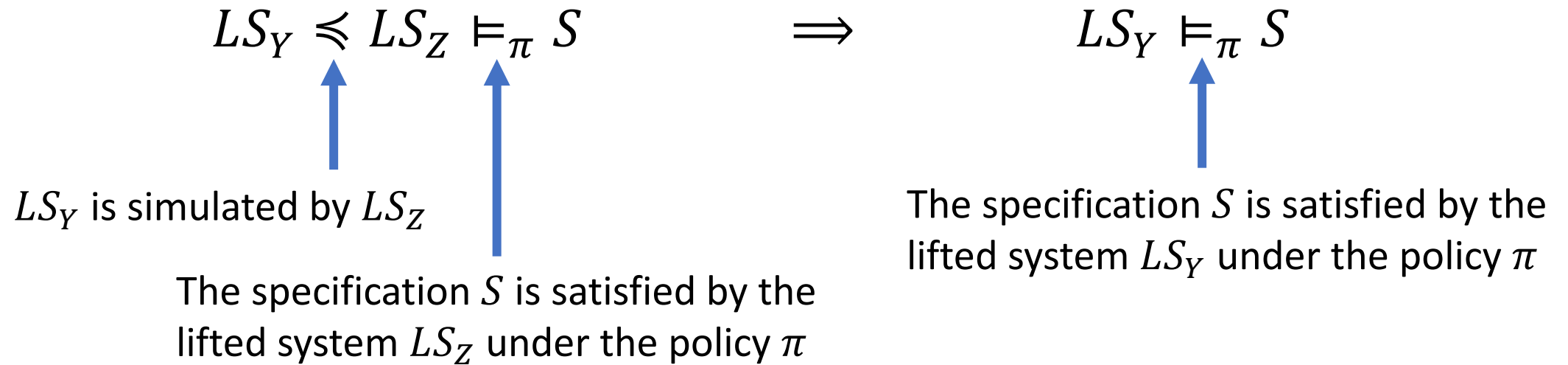
$$LS_Y \preceq LS_Z \implies \mathbf{B}_\pi[LS_Y] \subseteq \mathbf{B}_\pi[LS_Z]$$

Proof idea: If  $(\mathbf{x}, \mathbf{u}, \mathbf{y})$  is a max. sol. of  $LS_Y$ , it exists  $\mathbf{z}$  s.t.  $(\mathbf{x}, \mathbf{u}, \mathbf{z})$  is a max. sol. of  $LS_Z$

- $\forall x \in X: \psi_Y(x) \in \rho(\psi_Z(x))$   
 $\rightarrow y(0) \in \rho(z(0))$
- $\forall (z, u) \in \mathbb{R}^{n_z} \times U: f_Y(\rho(z), u) \subseteq \rho(f_Z(z, u))$   
 $\rightarrow y(t) \in \rho(z(t)) \implies y(t+1) \in \rho(z(t+1))$
- $\forall z \in \mathbb{R}^{n_z}: g_Y(\rho(z)) \subseteq \{g_Z(z)\}$   
 $\rightarrow y(t) \in \rho(z(t)) \implies g_Y(y(t)) = g_Z(z(t))$

# Why is it useful?

## COROLLARY



Given an unlifted system  $LS_X$ , and two affine lifted systems  $LS_Y$  and  $LS_Z$ , if

$$LS_X \preceq LS_Y \preceq LS_Z$$

then  $LS_Y$  is a « better » representation of  $LS_X$  than  $LS_Z$

# Some special cases of $LS_Y \preceq LS_Z$

If  $LS_Y$  is unlifted and

- $LS_Z$  is affine
  - Koopman over-approximation in (Balim, Aspeel, Liu, Ozay, 2023)
- $LS_Z$  is affine and all systems are autonomous (i.e.,  $U = \{0\}$ )
  - approximate immersion in (Wang, Jungers, Ong, 2023)
- $LS_Z$  is piecewise affine and unlifted
  - Hybridization in (Girard, Martin, 2011)

# Computational aspects

Verifying if  $LS_Y \preceq LS_Z$  is a feasibility problem:

Find  $\rho$  s.t.  $\rho: \mathbb{R}^{n_Z} \rightrightarrows \mathbb{R}^{n_Y}$  Infinite number of variables

- $\forall x \in X: \psi_Y(x) \in \rho(\psi_Z(x))$
- $\forall (z, u) \in \mathbb{R}^{n_Z} \times U: f_Y(\rho(z), u) \subseteq \rho(f_Z(z, u))$
- $\forall z \in \mathbb{R}^{n_Z}: g_Y(\rho(z)) \subseteq \{g_Z(z)\}$

Infinite number of constraints

We derived **finite-dimensional** sufficient conditions in two cases:

1.  $LS_Y$  is unlifted and  $LS_Z$  is (piecewise) affine
2.  $LS_Y$  and  $LS_Z$  are both (piecewise) affine

# In practice...

Given a classical (i.e., unlifted) system  $LS_X$

1. Pick  $K$  lifting functions  $\psi_1, \dots, \psi_K$
2. For each, compute an affine lifted system:  $LS_k$  s.t.  $LS_X \preceq LS_k$
3. If  $LS_i \preceq LS_j$ , delete  $LS_j$

# Experiments with backward reachable sets

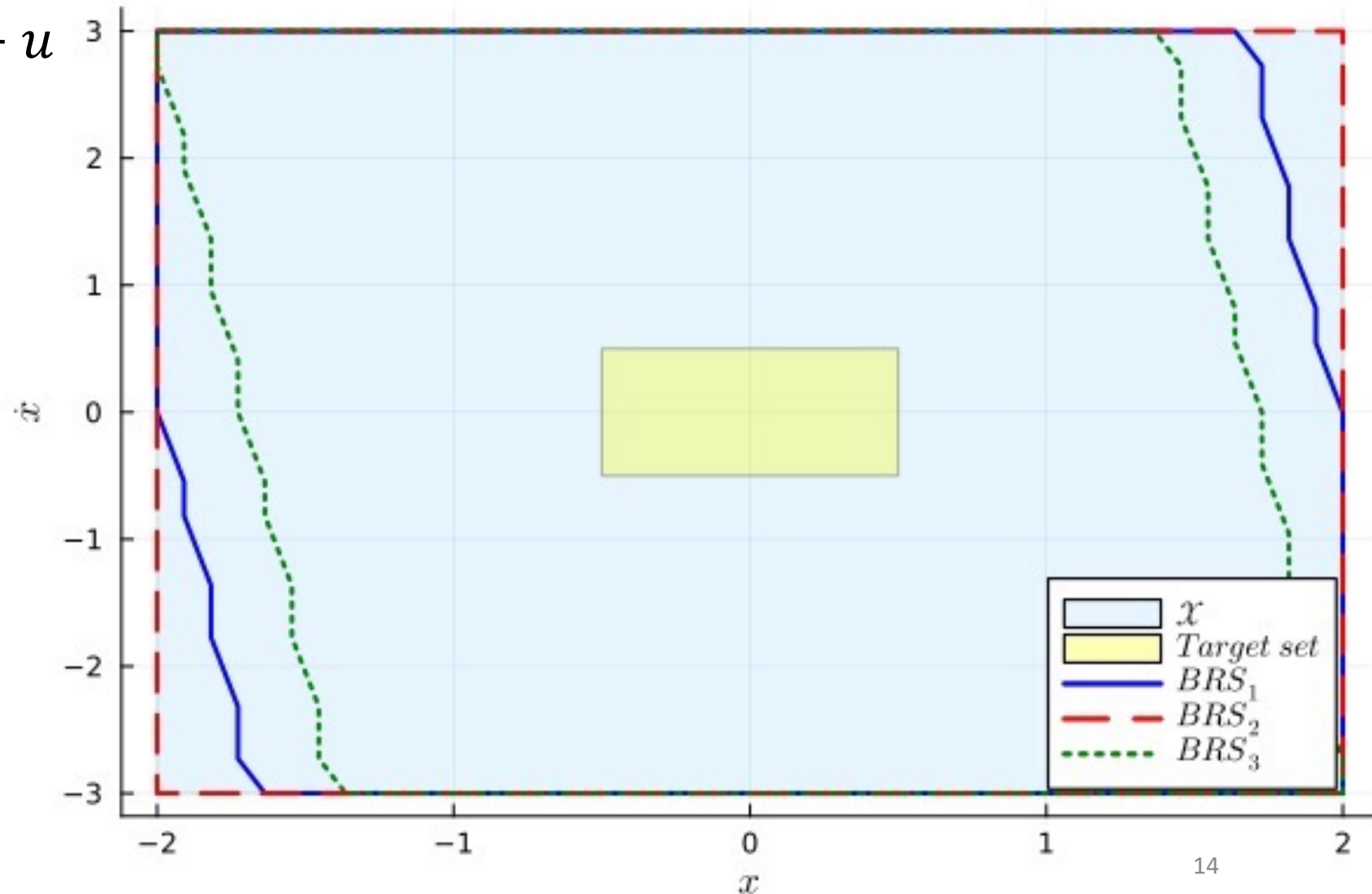
$$\ddot{x} = 2x - 2x^3 - 0.5 \dot{x} + u$$

$$\psi_1(x) = \begin{bmatrix} x \\ \dot{x} \end{bmatrix}$$

$$\psi_2(x) = \begin{bmatrix} x \\ \dot{x} \\ x^3 \end{bmatrix}$$

$$\psi_3(x) = \begin{bmatrix} x \\ \dot{x} \\ x\dot{x} \end{bmatrix}$$

$$LS_1 \preceq LS_3$$



# Take home message

Simulation relations between lifted systems

- Contribution to abstraction/hybridization theory  
→ More general class of simulations
- Contribution to finite-dimensional Koopman theory  
→ Tools to select the lifting function  $\psi$

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