



**SIMATS**  
ENGINEERING



**SIMATS**

Saveetha Institute of Medical And Technical Sciences  
(Declared as Deemed to be University under Section 3 of UGC Act 1956)

## **Department of Electronics and Communication Engineering**

### **ECA02 DIGITAL CIRCUITS**

### **LAB MANUAL**

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## 1. STUDY OF LOGIC GATES

DATE:

### AIM

:

To study logic gates and verify their truth tables of all the logic functions.

### APPARATUS REQUIRED:

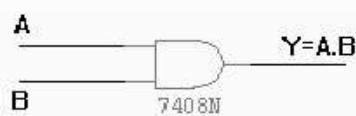
SL No.	COMPONENT	SPECIFICATION	QTY
1.	AND GATE	IC 7408	1
2.	OR GATE	IC 7432	1
3.	NOT GATE	IC 7404	1
4.	NAND GATE 2 I/P	IC 7400	1
5.	NOR GATE	IC 7402	1
6.	X-OR GATE	IC 7486	1
8.	IC TRAINER KIT	-	1
9.	CONNECTING WIRES		

### PROCEDURE:

- (i) Connections are given as per the circuit diagram.
- (ii) Logical inputs are given as per the circuit diagram.
- (iii) Observe the output and verify the truth table.

### AND GATE:

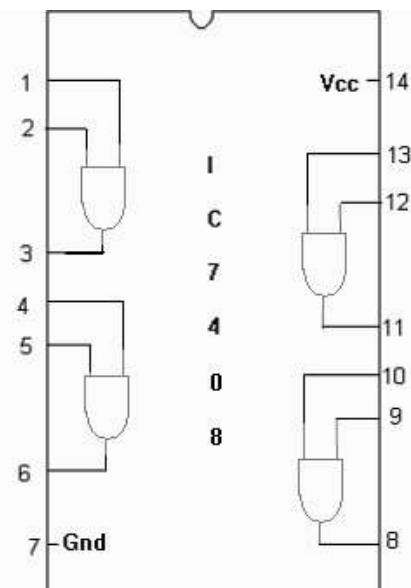
#### SYMBOL:



#### TRUTH TABLE:

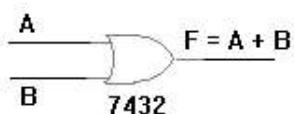
A	B	A.B
0	0	0
0	1	0
1	0	0
1	1	1

#### PIN DIAGRAM:

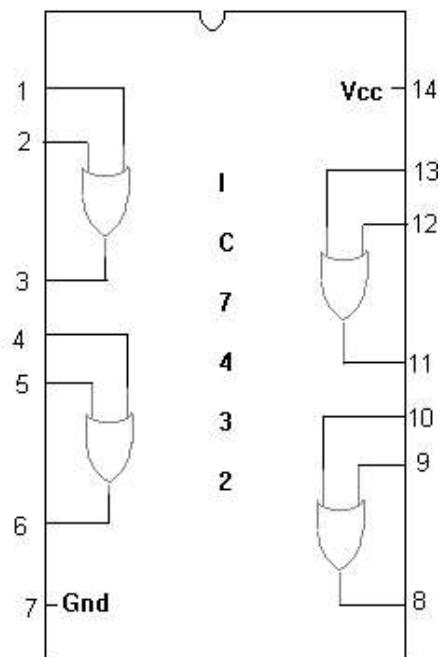


### OR GATE:

SYMBOL :



PIN DIAGRAM :

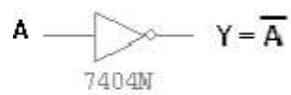


### TRUTH TABLE

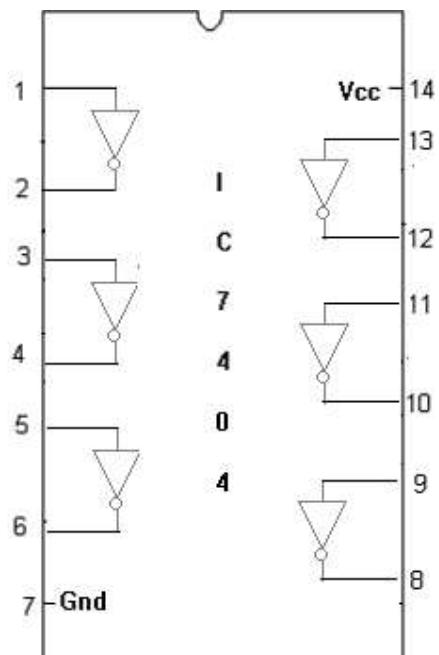
A	B	A+B
0	0	0
0	1	1
1	0	1
1	1	1

### NOT GATE:

SYMBOL:



PIN DIAGRAM:

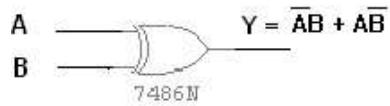


### TRUTH TABLE :

A	$\bar{A}$
0	1
1	0

### X-OR GATE:

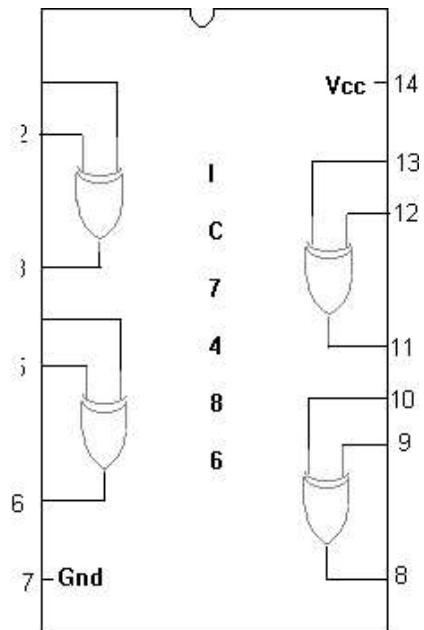
#### SYMBOL:



#### TRUTH TABLE :

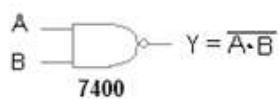
A	B	$\overline{AB} + \overline{A}\overline{B}$
0	0	0
0	1	1
1	0	1
1	1	0

#### PIN DIAGRAM:



### 2-INPUT NAND GATE:

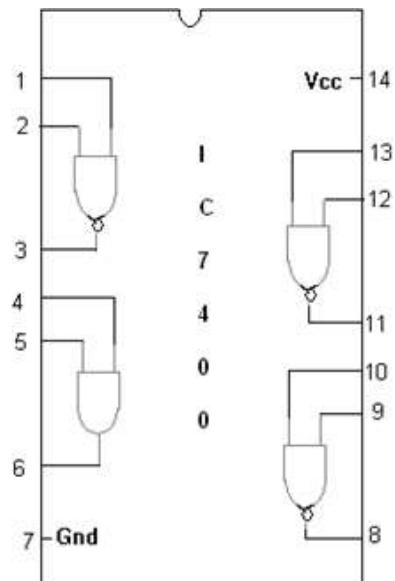
#### SYMBOL:



#### TRUTH TABLE

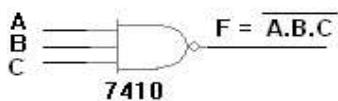
A	B	$\overline{A} \cdot \overline{B}$
0	0	1
0	1	1
1	0	1
1	1	0

#### PIN DIAGRAM:

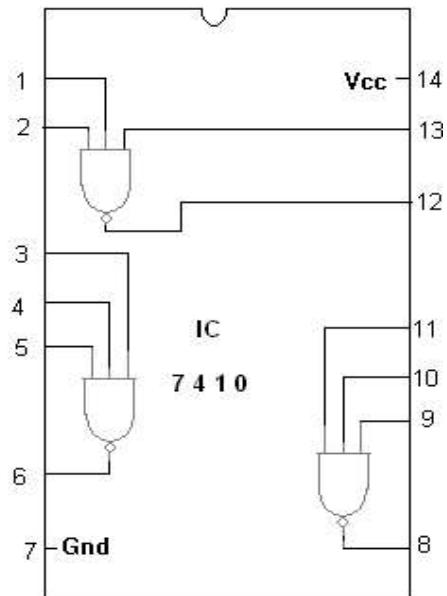


### 3-INPUT NAND GATE :

SYMBOL :



PIN DIAGRAM :

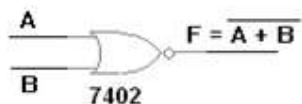


TRUTH TABLE

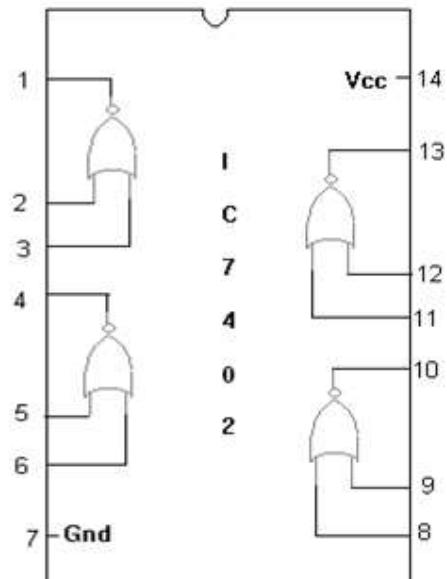
A	B	C	$\overline{A.B.C}$
0	0	0	1
0	0	1	1
0	1	0	1
0	1	1	1
1	0	0	1
1	0	1	1
1	1	0	1
1	1	1	0

### NOR GATE:

SYMBOL :



PIN DIAGRAM :



TRUTH TABLE

A	B	$\overline{A+B}$
0	0	1
0	1	0
1	0	0
1	1	0

### **RESULT:**

Thus, the gates (AND, OR, NOT, NAND, NOR, XOR, and XNOR GATES) are studied and verified using the Truth table.

**Test Case1:**

**Design a 3-bit majority voting circuit that outputs 1 if most of the three inputs are 1**

**Steps:**

- (a) Use three 2-input AND gates to compute  $A \cdot B$ ,  $A \cdot C$ , and  $B \cdot C$
- (b) Use a 3-input OR gate to combine these results

**Truth Table**

A	B	C	$A \cdot B$	$A \cdot C$	$B \cdot C$	F
0	0	0	0	0	0	0
0	0	1	0	0	0	0
0	1	0	0	0	0	0
0	1	1	0	0	1	1
1	0	0	0	0	0	0
1	0	1	0	1	0	1
1	1	0	1	0	0	1
1	1	1	1	1	1	1

**Test Case 2:**

**Design a circuit to check the parity of a 4-bit binary number. The output should be 1 if the number has even parity (even number of 1s)**

Steps: Cascade XOR gates

Compute  $P1 = A \oplus B$

Compute  $P2 = P1 \oplus C$

Compute  $F = P2 \oplus D$

**Truth Table**

A	B	C	D	$A \oplus B$	P2	F
0	0	0	0	0	0	0
0	0	0	1	0	0	1
0	0	1	0	0	1	1
0	0	1	1	0	1	0
1	1	1	1	0	1	0

## 2. VERIFICATION OF BOOLEAN THEOREMS USING LOGIC GATES

DATE:

### AIM

:

To verify Boolean theorems using logic gates

### APPARATUS REQUIRED:

SI.No.	COMPONENT	SPECIFICATION	QTY.
1.	AND GATE	IC 7408	1
2.	NOT GATE	IC 7404	1
3.	OR GATE	IC 7432	1
4.	IC TRAINER KIT	-	1
5.	CONNECTING WIRES		

### PROCEDURE:

- Connections are given as per the logic diagram
- Input is given to the circuit making high '1' i.e., +5 or Vcc Supply to the 14<sup>th</sup> pin, and for low '0' i.e., GND to the 7<sup>th</sup> pin of Gate IC
- Depending upon the truth table, if the LED Glow represents 1 and else it represents '0'
- Verify the truth table as given
- Repeat the procedure steps for different theorems.

### BOOLEAN THEOREM:

POSTULATES	I	II
Identity	$A + 0 = A$	$A \cdot 1 = A$
Commutative	$A + B = B + A$	$AB = BA$
Distributive	$A(B + C) = AB + AC$	$A + BC = (A + B)(A + C)$
Complement	$A + A' = 1$	$A \cdot A' = 0$
Idempotency	$A + A = A$	$A \cdot A = A$
	$A + 1 = 1$	$A \cdot 0 = 0$
Involution		$(A')' = A$
Absorption	$A + AB = A$	$A(A + B) = A$
	$A + A'B = A + B$	$A \cdot (A' + B) = AB$
Associative	$A + (B + C) = (A + B) + C$	$A(BC) = (AB)C$
De Morgan's Law	$(A + B)' = A' \cdot B'$	$(AB)' = A' + B'$

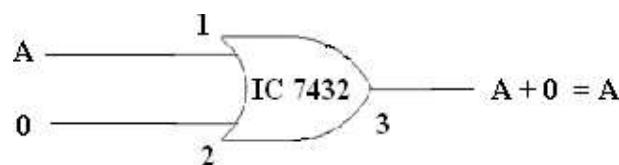
1.

## IDENTITY

### TRUTH TABLE

I/P		O/P
A	0	$A+0=A$
0	0	0
1	0	1

### LOGIC DIAGRAM

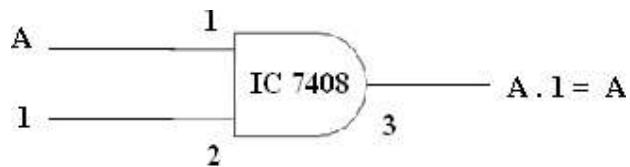


Identity law  $A + 0 = A$

### TRUTH TABLE

I/P		O/P
A	0	$A \cdot 1 = A$
0	0	0
1	0	1

### LOGIC DIAGRAM



Identity law  $A \cdot 1 = A$

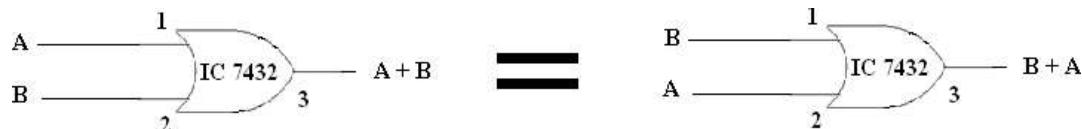
2.

### COMMUTATIVE

#### TRUTH TABLE

I/P		LHS (O/P)	RHS (O/P)
A	B	$A+B$	$B+A$
0	0	0	0
0	1	1	1
1	0	1	1
1	1	1	1

### LOGIC DIAGRAM

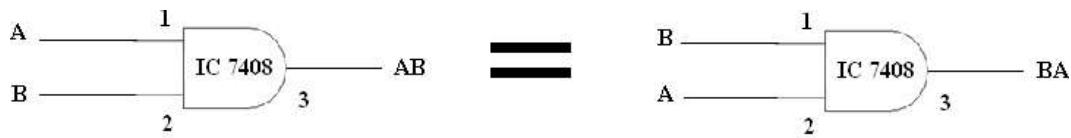


Commutative law  $A + B = B + A$

### TRUTH TABLE

I/P		LHS (O/P)	RHS (O/P)
A	B	$A \cdot B$	$B \cdot A$
0	0	0	0
0	1	0	0
1	0	0	0
1	1	1	1

### LOGIC DIAGRAM



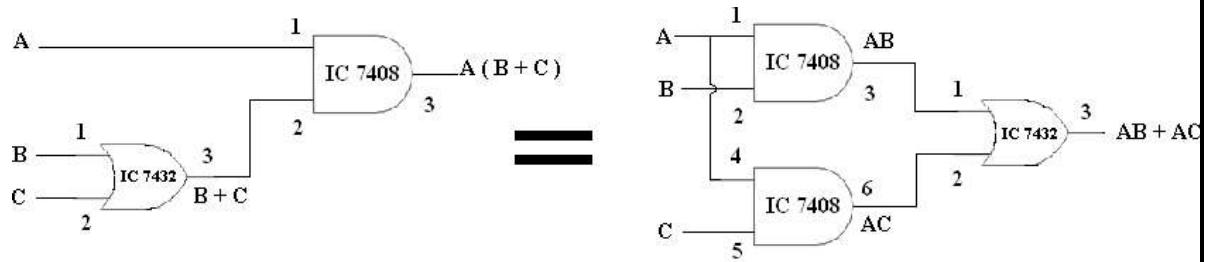
Commutative law  $AB = BA$

### 3. DISTRIBUTIVE

### TRUTH TABLE

I/P				LHS (O/P)			RHS (O/P)
A	B	C	$B+C$	$A(B+C)$	AB	AC	$AB+A C$
0	0	0	0	0	0	0	0
0	0	1	1	0	0	0	0
0	1	0	1	0	0	0	0
0	1	1	1	0	0	0	0
1	0	0	0	0	0	0	0
1	0	1	1	1	0	1	1
1	1	0	1	1	1	0	1
1	1	1	1	1	1	1	1

### LOGIC DIAGRAM

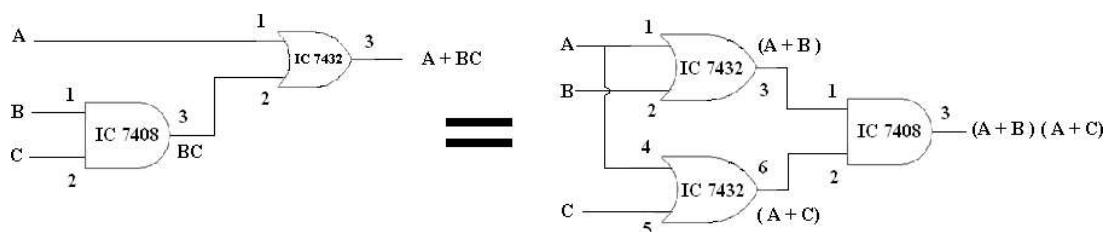


$$\text{Distributive 1 } A(B+C) = AB + AC$$

### TRUTH TABLE

I/P				LHS (O/P)			RHS (O/P)
A	B	C	BC	A+BC	A+B	A+C	(A+B)(A+C)
0	0	0	0	0	0	0	0
0	0	1	0	0	0	1	0
0	1	0	0	0	1	0	0
0	1	1	1	1	1	1	1
1	0	0	0	1	1	1	1
1	0	1	0	1	1	1	1
1	1	0	0	1	1	1	1
1	1	1	1	1	1	1	1

### LOGIC DIAGRAM



$$\text{Distributive 2 } A + BC = (A + B)(A + C)$$

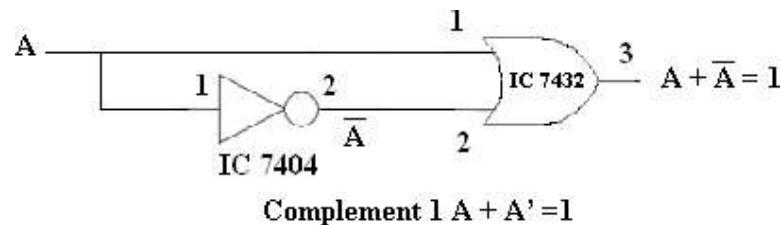
4.

### COMPLEMENT

#### TRUTH TABLE

I/P		O/P
A	A'	A+A'
0	1	1
1	0	1

### LOGIC DIAGRAM

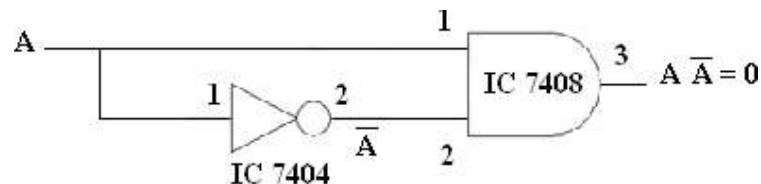


Complement 1  $A + A' = 1$

### TRUTH TABLE

I/P		O/P
A	$A'$	$AA'$
0	1	0
1	0	0

### LOGIC DIAGRAM



Complement 2  $A \cdot A' = 0$

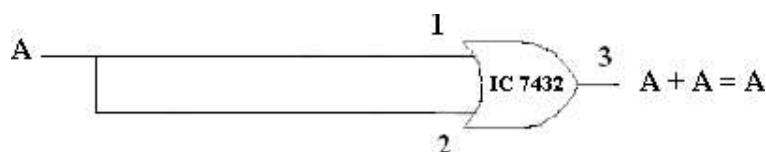
5.

### IDEMPOTENCY

### TRUTH TABLE

I/P		O/P
A	$A$	$A+A$
0	0	0
1	1	1

### LOGIC DIAGRAM

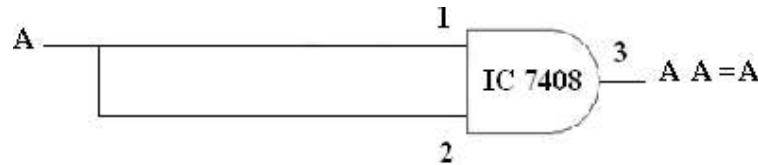


Idempotency 1  $A + A = A$

### TRUTH TABLE

I/P		O/P
A	A	AA
0	0	0
1	1	1

### LOGIC DIAGRAM

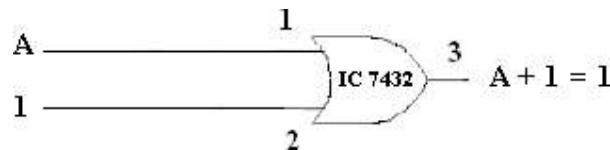


$$\text{Idempotency 2 } A \cdot A = A$$

### TRUTH TABLE

I/P		O/P
A	1	A+1
0	1	1
1	1	1

### LOGIC DIAGRAM

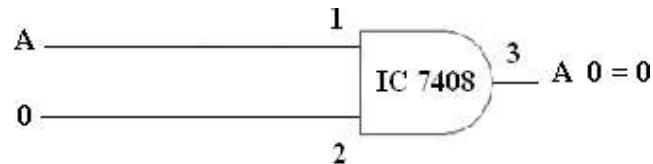


$$\text{Idempotency 3 } A + 1 = 1$$

### TRUTH TABLE

I/P		O/P
A	0	A.0
0	0	0
1	0	0

### LOGIC DIAGRAM



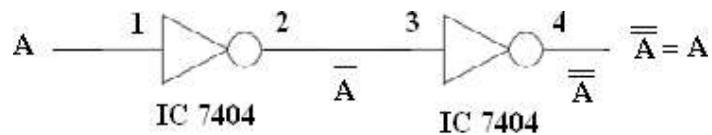
**Idempotency 4  $A \cdot 0 = 0$**

### 6. INVOLUTION

#### TRUTH TABLE

I/P		O/P
A	A'	$(A')'$
0	1	0
1	0	1

#### LOGIC DIAGRAM



**Involution  $(A')' = A$**

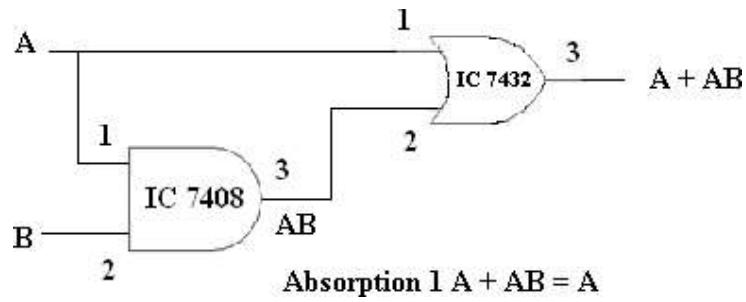
7.

### ABSORPTION

#### TRUTH TABLE

I/P			O/P
A	B	AB	$A + AB$
0	0	0	0
0	1	0	0
1	0	0	1
1	1	1	1

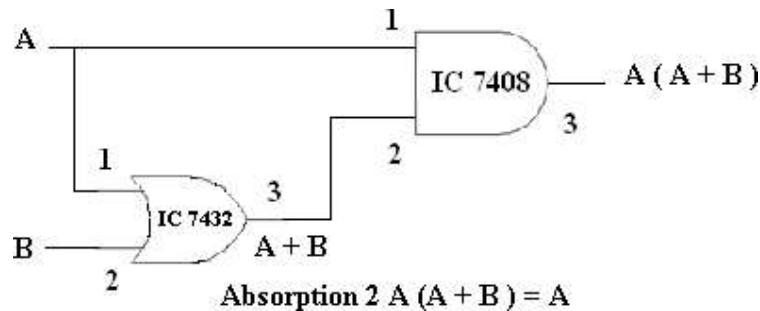
### LOGIC DIAGRAM



### TRUTH TABLE

I/P		O/P	
A	B	$A+B$	$A(A+B)$
0	0	0	0
0	1	1	0
1	0	1	1
1	1	1	1

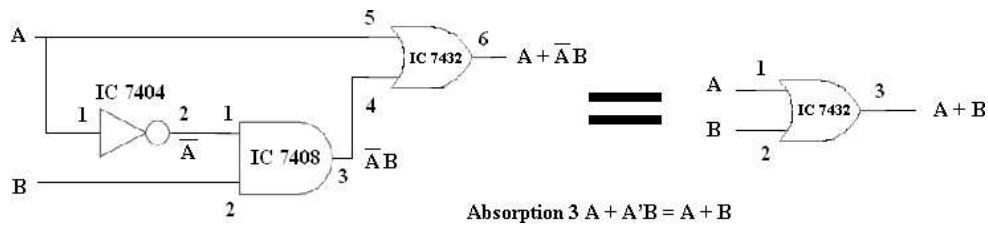
### LOGIC DIAGRAM



### TRUTH TABLE

I/P				LHS (O/P)	RHS (O/P)
A	B	$A'$	$A'B$	$A + (A'B)$	$A + B$
0	0	1	0	0	0
0	1	1	1	1	1
1	0	0	0	1	1
1	1	0	0	1	1

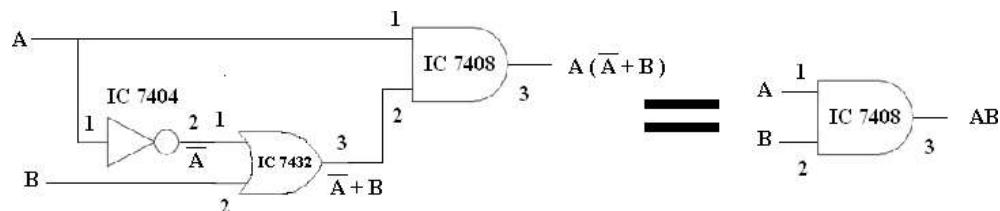
## LOGIC DIAGRAM



## TRUTH TABLE

I/P				LHS (O/P)	RHS (O/P)
A	B	A'	A' + B	$A (A' + B)$	$AB$
0	0	1	1	0	0
0	1	1	1	0	0
1	0	0	0	0	0
1	1	0	1	1	1

## LOGIC DIAGRAM



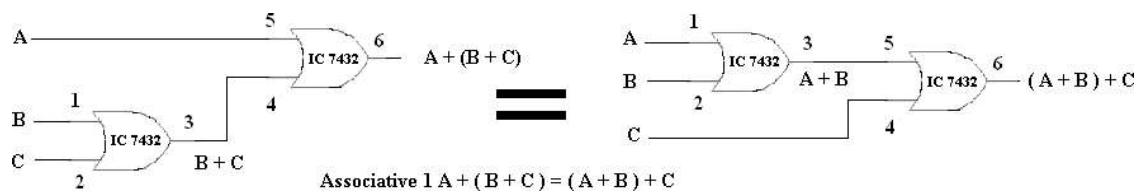
8.

## ASSOCIATIVE

### TRUTH TABLE

I/P				LHS (O/P)		RHS (O/P)
A	B	C	$B+C$	$A + (B+C)$	$A + B$	$(A + B) + C$
0	0	0	0	0	0	0
0	0	1	1	1	0	1
0	1	0	1	1	1	1
0	1	1	1	1	1	1
1	0	0	0	1	1	1
1	0	1	1	1	1	1
1	1	0	1	1	1	1
1	1	1	1	1	1	1

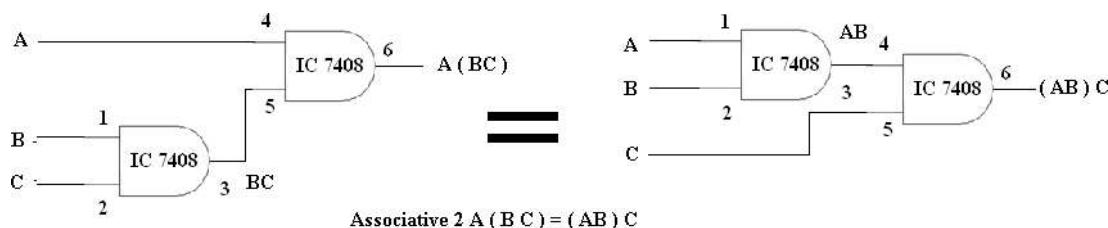
### LOGIC DIAGRAM



### TRUTH TABLE

I/P				LHS (O/P)		RHS (O/P)
A	B	C	BC	A (BC)	AB	(AB) C
0	0	0	0	0	0	0
0	0	1	0	0	0	0
0	1	0	0	0	0	0
0	1	1	1	0	0	0
1	0	0	0	0	0	0
1	0	1	0	0	0	0
1	1	0	0	0	1	0
1	1	1	1	1	1	1

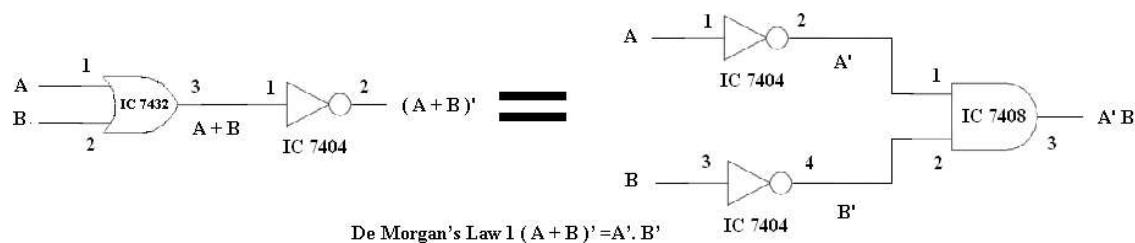
### LOGIC DIAGRAM



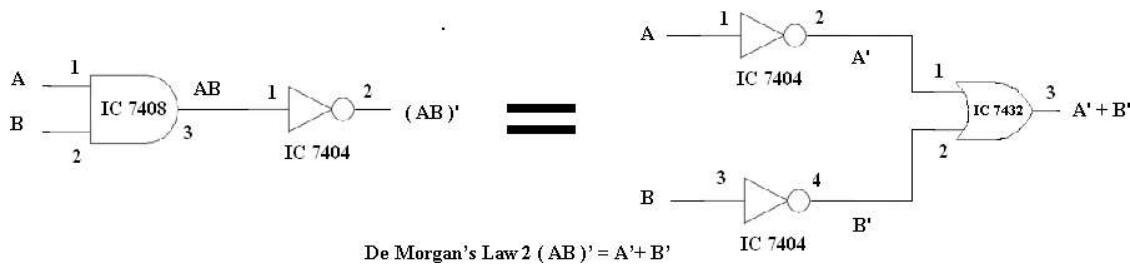
9.

**DE MORGAN'S LAW****TRUTH TABLE**

I/P			LHS (O/P)			RHS (O/P)
A	B	(A + B)	(A + B)'	A'	B'	A'.B'
0	0	0	1	1	1	1
0	1	1	0	1	0	0
1	0	1	0	0	1	0
1	1	1	0	0	0	0

**LOGIC DIAGRAM****TRUTH TABLE**

I/P			LHS (O/P)			RHS (O/P)
A	B	(AB)	(AB)'	A'	B'	A' + B'
0	0	0	1	1	1	1
0	1	0	1	1	0	1
1	0	0	1	0	1	1
1	1	1	0	0	0	0

**LOGIC DIAGRAM****RESULT:**

The different theorems of Boolean algebra and Demorgan's theorem are designed and verified using logic gates.

**Test Case 1:****Cascading De Morgan's Theorem - Verify the equivalence of the two expressions**

1.  $((A \cdot B)' + (C \cdot D)')'$
2.  $A \cdot B \cdot C' + D'$

**Stepwise calculation - Expression 1:  $((A \cdot B)' + (C \cdot D)')'$** 

1. Compute  $A \cdot B$  and  $C \cdot D$
2. Compute  $(A \cdot B)'$  and  $(C \cdot D)'$
3. Compute  $(A \cdot B)' + (C \cdot D)'$
4. Finally, negate the result:  $((A \cdot B)' + (C \cdot D)')'$

**Expression 2:  $A \cdot B \cdot (C' + D')$** 

1. Compute  $C'$  and  $D'$
2. Compute  $C' + D'$
3. Compute  $A \cdot B \cdot (C' + D')$

**Conclusion: Two expressions are not logically equivalent.****Test Case 2:****Design a digital circuit that demonstrates the practical application of the Complement Law in a multi-variable logic system. Specifically, simplify and verify the expression  $F = A \cdot A \cdot B + (A + A) \cdot C$** 

Step 1: Simplify the Expression

A	B	C	$A \cdot A'$	$A + A'$	$(A + A') \cdot C$	$F = C$
0	0	0	0	1	0	0
0	0	1	0	1	1	1
0	1	0	0	1	0	0
0	1	1	0	1	1	1
1	0	0	0	1	0	0
1	0	1	0	1	1	1
1	1	0	0	1	0	0
1	1	1	0	1	1	1

**Simplified Expression:  $F = C$**

### 3. IMPLEMENTATION OF COMBINATION CIRCUIT FOR AN ARBITRARY FUNCTION

DATE:

#### AIM:

To Realize the Boolean Expression  $AB + BC + AC$  using Logic gates and verify its performance with its truth table and the Digital Trainer kit.  $Y = AC + AB + BC$

#### APPARATUS REQUIRED:

SI.No.	COMPONENT	SPECIFICATION	QTY.
1.	AND GATE	IC 7408	1
2.	OR GATE	IC 7432	1
3.	IC TRAINER KIT	-	1
4.	CONNECTING WIRES	-	few

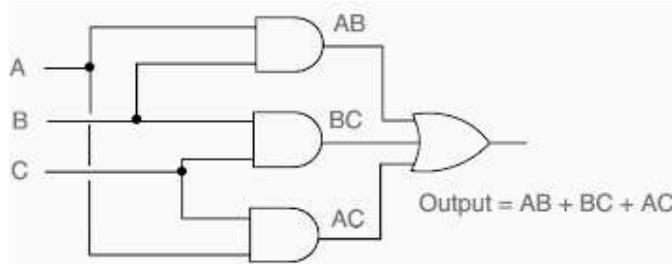
#### THEORY:

Combinational logic circuits are circuits in which the output at any time depends upon the combination of input signals present at that instant only, and does not depend on any past conditions. The combinational circuit block can be considered as a network of logic gates that accept signals from inputs and generate signals to outputs

#### PROCEDURE:

- a. Connections are given as per the logic diagram
- b. The input is given to the circuit, making high '1', i.e., +5 or Vcc Supply to the 14th pin, and low '0', i.e., GND to the 7<sup>th</sup> pin of the Gate IC.
- c. Depending upon the truth table, if the LED Glows, it represents 1, and otherwise, it represents 0.
- d. Verify the truth table as given

#### LOGIC DIAGRAM:



**TRUTH TABLE:**

I/P						O/P
A	B	C	AB	BC	AC	<b>AB + BC + AC</b>
0	0	0	0	0	0	0
0	0	1	0	0	0	0
0	1	0	0	0	0	0
0	1	1	0	1	0	1
1	0	0	0	0	0	0
1	0	1	0	0	1	1
1	1	0	1	0	0	1
1	1	1	1	1	1	1

**RESULT:**

Thus, the given Arbitrary function  $AB+BC=AC$  was verified using truth table.

**Test case:**

- Set the values of A, B, and C using the input switches.
- Observe the output LED.
- Compare the LED output with the expected result from the truth table

**Test Case 1: A=0, B=0, C=0**

- Inputs: A = 0, B = 0, C = 0
- Expected Output: AB+BC+AC=0
- Reasoning: All terms AB, BC, and AC are 0, so the final output should be 0.

**Test Case 2: A=0, B=0, C=1**

- Inputs: A = 0, B = 0, C = 1
- Expected Output: AB+BC+AC=0
- Reasoning: All terms AB, BC, and AC are 0, so the final output should be 0.

**Test Case 3: A=0, B=1, C=0**

- Inputs: A = 0, B = 1, C = 0
- Expected Output: AB+BC+AC=0
- Reasoning: All terms AB, BC, and AC are 0, so the final output should be 0.

**Test Case 4: A=0, B=1, C=1**

- Inputs: A = 0, B = 1, C = 1
- Expected Output: AB+BC+AC=1
- Reasoning: AB=0, BC=1, AC=0. The OR gate will output 1.

**Test Case 5: A=1, B=0, C=1**

- Inputs: A = 1, B = 0, C = 1
- Expected Output: AB+BC+AC=1
- Reasoning: AB=0, BC=0, AC=1. The OR gate will output 1.

#### 4. DESIGN AND IMPLEMENTATION OF CODE CONVERTERS USING LOGIC GATES

DATE:

**AIM:** To design and implement a 4-bit code converter for the following

- a. Binary to gray code converter
- b. Gray to binary code converter

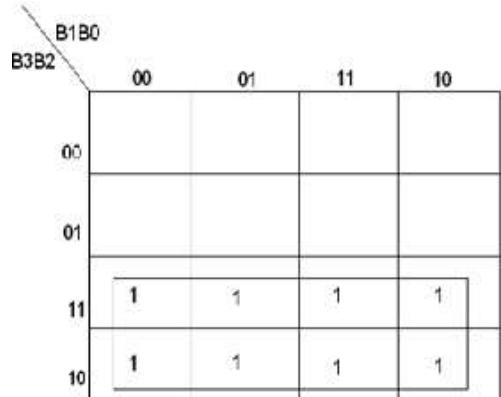
**APPARATUS REQUIRED:**

Sl.No.	COMPONENT	SPECIFICATION	QTY
1.	X-OR GATE	IC 7486	1
5.	IC TRAINER KIT	-	1
6.	CONNECTING WIRES	-	35

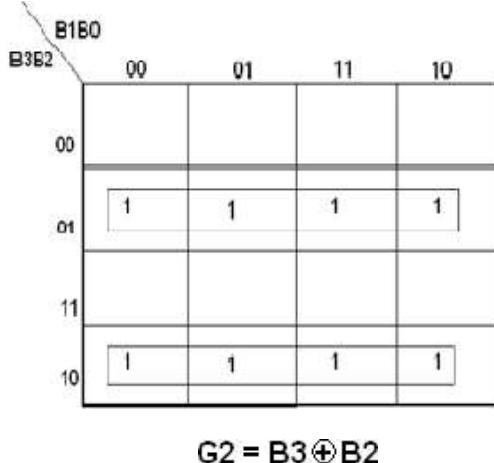
**DESIGN FOR BINARY TO GRAY CODE CONVERTERTRUTH TABLE:**

Binary input				Gray code output			
B3	B2	B1	B0	G3	G2	G1	G0
0	0	0	0	0	0	0	0
0	0	0	1	0	0	0	1
0	0	1	0	0	0	1	1
0	0	1	1	0	0	1	0
0	1	0	0	0	1	1	0
0	1	0	1	0	1	1	1
0	1	1	0	0	1	0	1
0	1	1	1	0	1	0	0
1	0	0	0	1	1	0	0
1	0	0	1	1	1	0	1
1	0	1	0	1	1	1	1
1	0	1	1	1	1	1	0
1	1	0	0	1	0	1	0
1	1	0	1	1	0	1	1
1	1	1	0	1	0	0	1
1	1	1	1	1	0	0	0

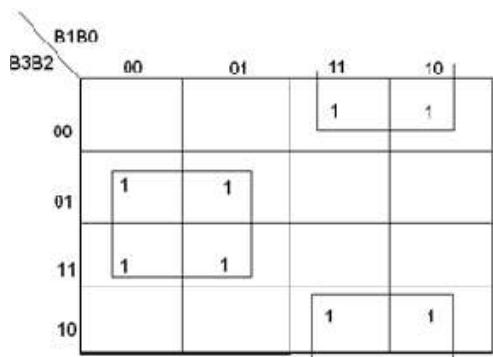
K-Map for G3:



K-Map for G2:

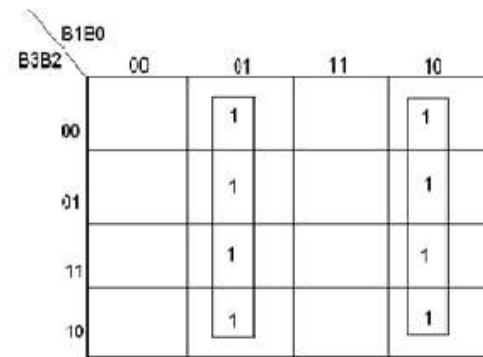


K-Map for G1:



$$G1 = B_1 \oplus B_2$$

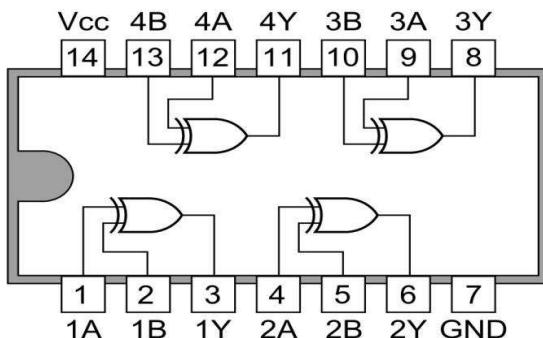
K-Map for G0:



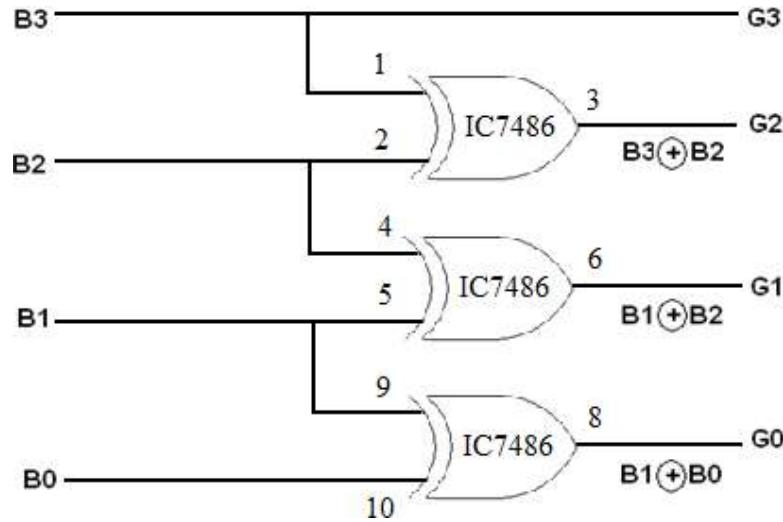
$$G0 = B_1 \oplus B_0$$

#### PIN DIAGRAM:

7486 Quad 2-input ExOR Gates



**LOGIC DIAGRAM FOR BINARY TO GRAY CODE CONVERTOR:**



**DESIGN FOR GRAY TO BINARY CODE CONVERTER TRUTH TABLE:**

Gray Code				Binary Code			
G3	G2	G1	G0	B3	B2	B1	B0
0	0	0	0	0	0	0	0
0	0	0	1	0	0	0	1
0	0	1	1	0	0	1	0
0	0	1	0	0	0	1	1
0	1	1	0	0	1	0	0
0	1	1	1	0	1	0	1
0	1	0	1	0	1	1	0
0	1	0	0	0	1	1	1
1	1	0	0	1	0	0	0
1	1	0	1	1	0	0	1
1	1	1	1	1	0	1	0
1	1	1	0	1	0	1	1
1	0	1	0	1	1	0	0
1	0	1	1	1	1	0	1
1	0	0	1	1	1	1	0
1	0	0	0	1	1	1	1

**K-Map for B3**

		G1G0	G3G2		
		00	01	11	10
00		0	0	0	0
01		0	0	0	0
11		1	1	1	1
10		1	1	1	1

$$B3 = G3$$

**K-Map for B2**

		G1G0	G3G2		
		00	01	11	10
00		0	0	0	0
01		1	1	1	1
11		0	0	0	0
10		1	1	1	1

$$B2 = G3 \oplus G2$$

**K-Map for B1**

		G1G0	G3G2		
		00	01	11	10
00		0	0	1	1
01		1	1	0	0
11		0	0	1	1
10		1	1	0	0

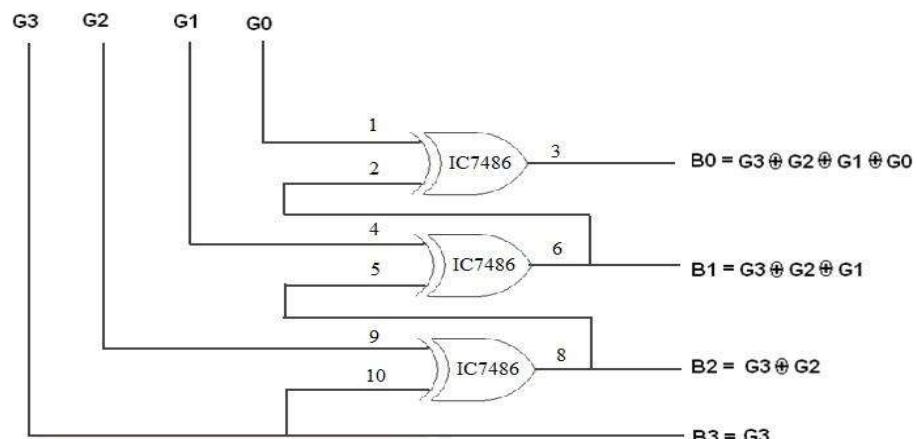
$$B1 = G3 \oplus G2 \oplus G1$$

**K-Map for B0**

		G1G0	G3G2		
		00	01	11	10
00		0	1	0	1
01		1	0	1	0
11		0	1	0	1
10		1	0	1	0

$$B0 = G3 \oplus G2 \oplus G1 \oplus G0$$

#### LOGIC DIAGRAM FOR GRAY TO BINARY CODE CONVERTER:



**RESULT:**

Thus, the design for the code converter was done and verified using logic gates successfully.



**Test Case 1**

Input (Binary): 0000

Expected Output (Gray Code): 0000

Step-by-Step Function Execution:

- $G_3=B_3=0$
- $G_2=B_3 \oplus B_2=0 \oplus 0=0$
- $G_1=B_2 \oplus B_1=0 \oplus 0=0$
- $G_0=B_1 \oplus B_0=0 \oplus 0=0$

Result: Gray Code = 0000

**Test Case 2**

Input (Binary): 0101

Expected Output (Gray Code): 0111

Step-by-Step Function Execution:

- $G_3=B_3=0$
- $G_2=B_3 \oplus B_2=0 \oplus 1=1$
- $G_1=B_2 \oplus B_1=1 \oplus 0=1$
- $G_0=B_1 \oplus B_0=0 \oplus 1=1$

Result: Gray Code = 0111

**Test Case 3**

Input (Binary): 1000

Expected Output (Gray Code): 1100

Step-by-Step Function Execution:

- $G_3=B_3=1$
- $G_2=B_3 \oplus B_2=1 \oplus 0=1$
- $G_1=B_2 \oplus B_1=0 \oplus 0=0$
- $G_0=B_1 \oplus B_0=0 \oplus 0=0$

Result: Gray Code = 1100

**Test Case 4**

Input (Binary): 1110

Expected Output (Gray Code): 1001

Step-by-Step Function Execution:

- $G_3=B_3=1$
- $G_2=B_3 \oplus B_2=1 \oplus 1=0$
- $G_1=B_2 \oplus B_1=1 \oplus 1=0$
- $G_0=B_1 \oplus B_0=1 \oplus 0=1$

Result: Gray Code = 1001

**Test Case 5**

Input (Binary): 0011

Expected Output (Gray Code): 0010

Step-by-Step Function Execution:

- $G_3=B_3=0$
- $G_2=B_3 \oplus B_2=0 \oplus 0=0$
- $G_1=B_2 \oplus B_1=0 \oplus 1=1$
- $G_0=B_1 \oplus B_0=1 \oplus 1=0$

Result: Gray Code = 0010

## 5. DESIGN AND IMPLEMENTATION OF HALF ADDER AND FULL ADDER USING LOGIC GATES

**DATE:**

**AIM:**

To design and construct half adder, and full adder circuits and verify the truth table using logic gates.

**APPARATUS REQUIRED:**

SI.No.	COMPONENT	SPECIFICATION	QTY.
1.	AND GATE	IC 7408	1
2.	X-OR GATE	IC 7486	1
4.	OR GATE	IC 7432	1
5.	IC TRAINER KIT	-	1
6.	CONNECTING WIRES	-	23

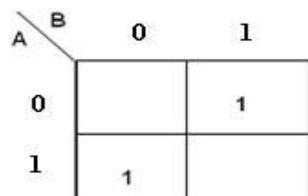
**PROCEDURE:**

- Connections are given as per the logic diagram
- Input are given to the circuit making high '1' i.e. +5 or Vcc Supply to the 14<sup>th</sup> pin and for low '0' i.e. GND to the 7<sup>th</sup> pin of the Gate IC
- Depending upon the truth table, if the LED Glow it represents 1 and else it represents '0'
- Verify the truth table as given

**HALF ADDER: TRUTH TABLE:**

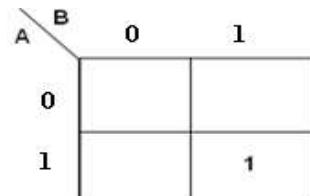
A	B	SUM	CARRY
0	0	0	0
0	1	1	0
1	0	1	0
1	1	0	1

**K-Map for SUM:**



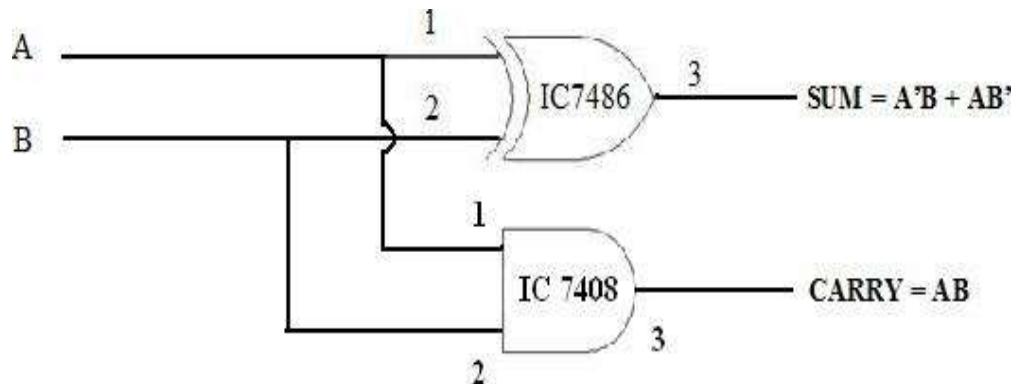
$$\text{SUM} = A'B + AB' = A \oplus B$$

**K-Map for CARRY:**



$$\text{CARRY} = AB$$

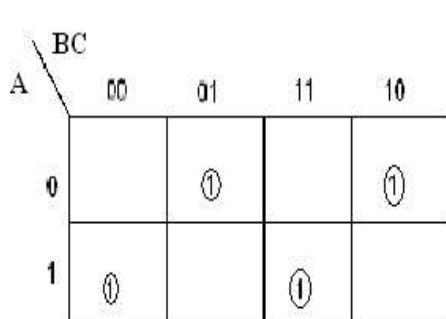
**LOGIC DIAGRAM FOR HALF ADDER:**



**FULL ADDER: TRUTH TABLE:**

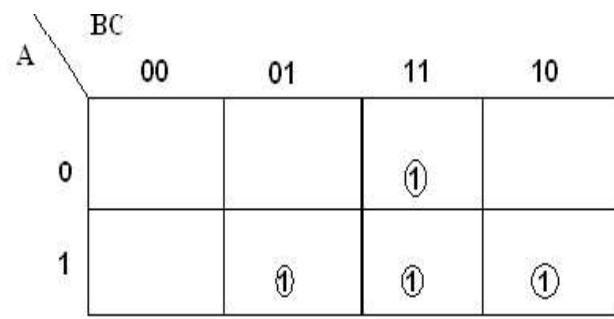
A	B	C	CARRY	SUM
0	0	0	0	0
0	0	1	0	1
0	1	0	0	1
0	1	1	1	0
1	0	0	0	1
1	0	1	1	0
1	1	0	1	0
1	1	1	1	1

K-Map for SUM

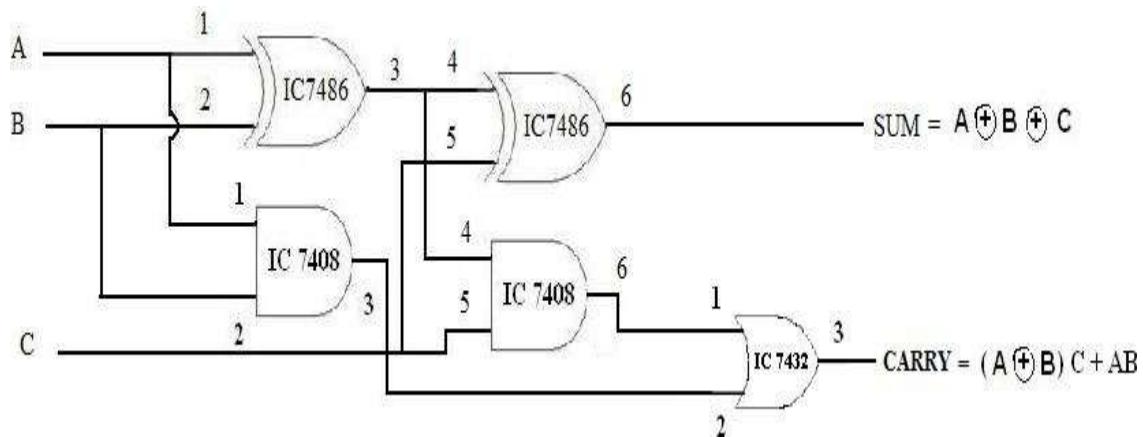


$$\begin{aligned} SUM &= A'B'C + A'BC' + ABC' + ABC \\ &= A \oplus B \oplus C \end{aligned}$$

K-Map for CARRY



$$\begin{aligned} CARRY &= A'B'C + AB'C' + ABC' + ABC \\ &= (A \oplus B)C + AB \end{aligned}$$

**LOGIC DIAGRAM FOR FULL ADDER:****RESULT:**

Thus, the design for half adder and full adder were designed and verified successfully using Truth table.

### **Half Adder Test Case 1: Fault Injection**

**Question:** If the XOR gate in the Half Adder fails and behaves like an OR gate, what will be the output for A = 1 and B = 1?

**Answer:**

Sum = 1 (A OR B = 1 OR 1 = 1).

Carry remains correct = 1 (A AND B = 1 AND 1 = 1).

This creates an incorrect result, indicating XOR gate failure.

**Test Case 2:**

**Question:** While implementing a Half Adder in hardware, your team observes increased complexity due to XOR gates. Is it possible to design a Half Adder without using XOR gates? If so, how?

**Answer:**

Yes, the Sum output can be expressed using only AND, OR, and NOT gates:

Sum = (A AND NOT B) OR (NOT A AND B).

However, this increases the gate count and latency.

**Full Adder Test Case 1: Multi-bit Addition**

**Question:** Implement a 4-bit binary adder using Full Adder circuits. Show the step-by-step calculation for inputs A = 1011 and B = 1101. Include carry propagation.

**Answer:**

Step 1: Add LSBs: A0 = 1, B0 = 1, Cin = 0 → Sum = 0, Carry = 1.

Step 2: A1 = 1, B1 = 0, Cin = 1 → Sum = 0, Carry = 1.

Step 3: A2 = 0, B2 = 1, Cin = 1 → Sum = 0, Carry = 1.

Step 4: A3 = 1, B3 = 1, Cin = 1 → Sum = 1, Carry = 1.

Final result: Sum = 10000.

**Test Case 2: Multi-Bit Binary Addition**

**Question:** You are designing a 4-bit Ripple Carry Adder using Full Adders. Describe the circuit's operation for the binary inputs A = 1011 and B = 1101. Include details about Carry propagation.

**Answer:**

For each bit:

LSB: A0 = 1, B0 = 1 → Sum = 0, Carry = 1.

Second bit: A1 = 1, B1 = 0, Cin = 1 → Sum = 0, Carry = 1.

Third bit: A2 = 0, B2 = 1, Cin = 1 → Sum = 0, Carry = 1.

MSB: A3 = 1, B3 = 1, Cin = 1 → Sum = 1, Carry = 1.

Final output: Sum = 10000.

## **6. DESIGN AND IMPLEMENTATION OF HALF SUBTRACTOR AND FULL SUBTRACTOR USING LOGIC GATES**

**DATE:**

**AIM:**

To design and construct half subtractor and full subtractor circuits and verify the truth table using logic gates.

**APPARATUS REQUIRED:**

SI.No.	COMPONENT	SPECIFICATION	QTY.
1.	AND GATE	IC 7408	1
2.	X-OR GATE	IC 7486	1
3.	NOT GATE	IC 7404	1
4.	OR GATE	IC 7432	1
3.	IC TRAINER KIT	-	1
4.	CONNECTING WIRES	-	23

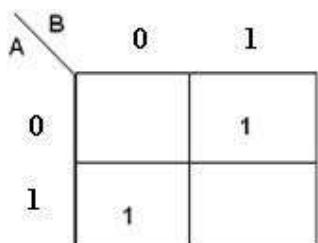
**PROCEDURE:**

- a. Connections are given as per the logic diagram
- b. Input is given to the circuit making high '1' i.e. +5 or Vcc Supply to the 14<sup>th</sup> pin and for low '0' i.e. GND to the 7<sup>th</sup> pin of the Gate IC
- c. Depending upon the truth table, if the LED Glow it represents 1 and else it represents '0'
- d. Verify the truth table as given

**HALF SUBTRACTOR: TRUTH TABLE**

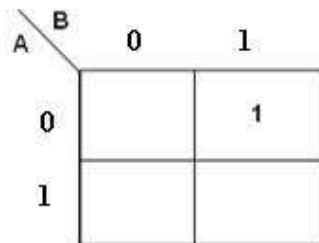
A	B	DIFFERENCE	BORROW
0	0	0	0
0	1	1	1
1	0	1	0
1	1	0	0

K-Map for DIFFERENCE



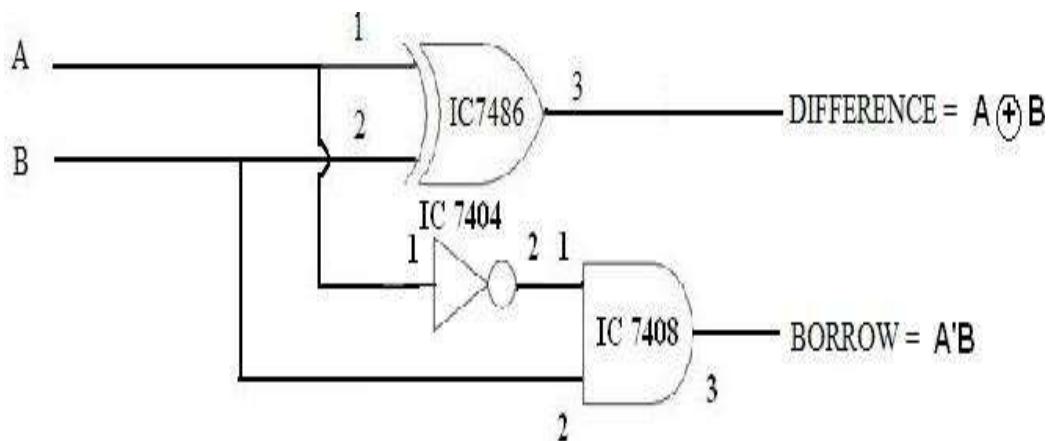
$$\begin{aligned}\text{DIFFERENCE} &= A'B + AB' \\ &= A \oplus B\end{aligned}$$

K-Map for BORROW



$$\text{BORROW} = A'B$$

LOGIC DIAGRAM FOR HALF SUBTRACTOR:



FULL SUBTRACTOR: TRUTH TABLE:

A	B	C	DIFFERENCE	BORROW
0	0	0	0	0
0	0	1	1	1
0	1	0	1	1
0	1	1	0	1
1	0	0	1	0
1	0	1	0	0
1	1	0	0	0
1	1	1	1	1

**K-Map for DIFFERENCE**

		BC	00	01	11	10
		A	0	1	0	1
A	B	00	1			1
		01				
A	B	11		1		
		10				

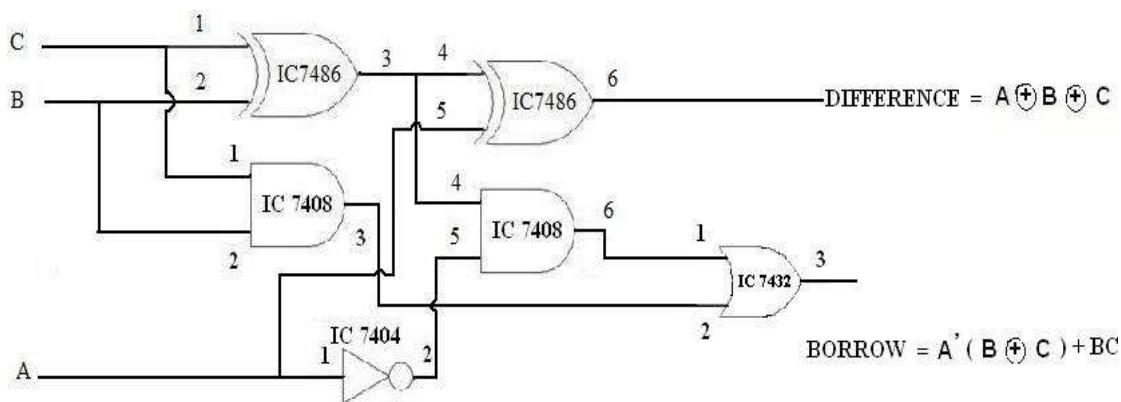
**K-Map for BORROW**

		BC	00	01	11	10
		A	0	1	0	1
A	B	00		1	1	1
		01				
A	B	11				
		10				

$$\begin{aligned}\text{Difference} &= A'B'C + A'BC' + AB'C' + ABC \\ &= A \oplus B \oplus C\end{aligned}$$

$$\begin{aligned}\text{Borrow} &= A'B + BC + A'C \\ &= A'(B \oplus C) + BC\end{aligned}$$

#### LOGIC DIAGRAM FOR FULL SUBTRACTOR:



#### RESULT:

Thus, the design for half subtractor and full subtractor were designed and verified successfully using the Truth table.

### **Half Subtractor Test Case 1: Circuit Design**

**Question:** Explain how the Half Subtractor circuit can be implemented on a digital trainer kit using basic logic gates. Why is the XOR gate used for the Difference output, and how does the AND gate with a NOT gate produce the Borrow output?

**Answer:**

XOR gate computes the Difference:  $A \oplus B$ , as it outputs 1 when the inputs are different.

NOT gate inverts A, and the AND gate computes Borrow:  $\bar{A} \cdot B$ , since Borrow occurs when B is 1 and A is 0.

### **Test Case 2: Fault Detection**

**Question:** During testing, you notice that the Difference output is always 0, irrespective of the inputs. Which component in the Half Subtractor circuit might have failed, and how would you troubleshoot this using the digital trainer kit?

**Answer:**

Likely issue: Faulty XOR gate.

Troubleshooting: Use a logic probe to verify if the XOR gate correctly processes inputs. Check inputs A and B and confirm if  $A \oplus B$  is computed.

### **Full Subtractor Test Case 1: Circuit Design**

**Question:** Design a Full Subtractor circuit using basic logic gates (XOR, AND, OR, NOT). Explain how the Difference and Borrow outputs are computed logically and implement the circuit on the digital trainer kit.

**Expected Answer:**

Logic equations:

Difference =  $A \oplus B \oplus B_{in}$  (using two XOR gates).

Borrow =  $(\bar{A} \cdot B) + (\bar{A} \cdot B_{in}) + (B \cdot B_{in})$  (using AND and OR gates).

Connect gates on the trainer kit and verify the circuit against the truth table.

### **Test Case 2: Validation with Multi-Bit Inputs**

**Question:** Implement a 4-bit binary subtractor using Full Subtractors on the digital trainer kit. Test the circuit for inputs A=1011 and B=1101 with initial Borrow-In = 0. Explain the step-by-step operation and results.

**Answer:**

Perform subtraction bit by bit:

LSB:  $A_0=1, B_0=1, B_{in}=0 \rightarrow$  Difference = 0, Borrow = 0.

Next bit:  $A_1=1, B_1=0, B_{in}=0 \rightarrow$  Difference = 1, Borrow = 0.

Next bit:  $A_2=0, B_2=1, B_{in}=0 \rightarrow$  Difference = 1, Borrow = 1.

MSB:  $A_3=1, B_3=1, B_{in}=1 \rightarrow$  Difference = 1, Borrow = 0.

Final output: Difference = 0010, Borrow = 0.

## 7. DESIGN AND IMPLEMENTATION OF MULTIPLEXERS

DATE:

### AIM:

To Design and implement a 2 to 1 Multiplexer using logic gates.

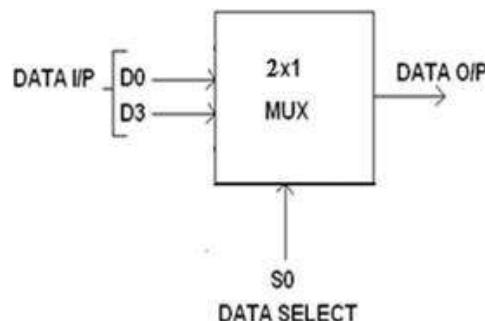
### APPARATUS REQUIRED:

SI.No.	COMPONENT	SPECIFICATION	QTY.
1.	AND GATE	IC 7408	1
2.	OR GATE	IC 7432	1
3.	NOT GATE	IC 7404	1
4.	IC TRAINER KIT	-	1
5.	CONNECTING WIRES	-	32

### PROCEDURE:

- Connections are given as per the logic diagram
- Input is given to the circuit making high '1' i.e. +5 or Vcc Supply to the 16<sup>th</sup> pin and for low '0' i.e. GND to the 8<sup>th</sup> pin of the Gate IC
- Depending upon the truth table, if the LED Glow represents 1 and else it represents '0'
- Verify the truth table as given.

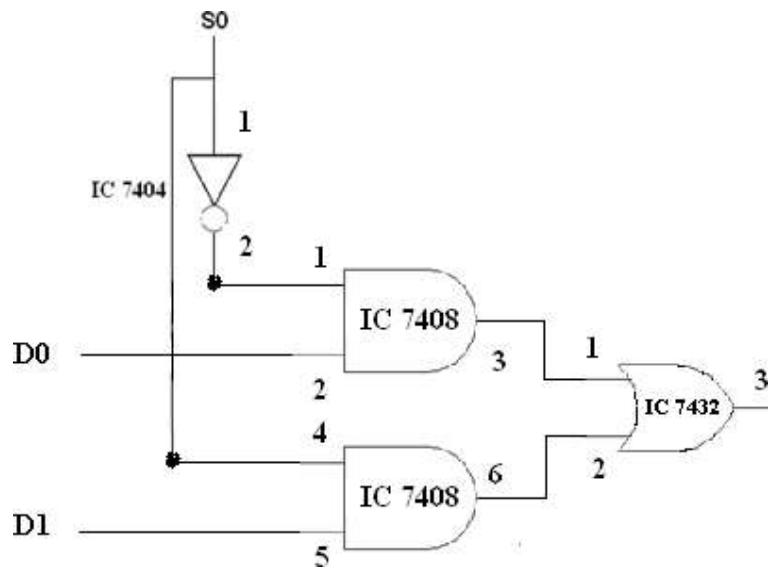
### BLOCK DIAGRAM FOR 2:1 MULTIPLEXER:



### MULTIPLEXER USING GATES:TRUTH TABLE:

S	Y =
0	D0
1	D1

### CIRCUIT DIAGRAM FOR MULTIPLEXER:



### RESULT:

Thus, the 2 to 1 Multiplexer circuit was designed using logic gates, and outputs are verified using the Truth table.

**Test Case 1:**

S = 0, D0 = 0, D1 = 1

Output should be Y = 0 (since S = 0, output should be from D0).

**Test Case 2:**

S = 1, D0 = 1, D1 = 0

Output should be Y = 0 (since S = 1, output should be from D1).

Implementation using Digital Trainer Kit:

Connect the data inputs (D0, D1) to switches or inputs on the trainer kit.

Use another switch to control the select input (S).

Use the trainer's logic gates (AND, OR, NOT) to implement the above logic.

Verify the output using a LED display or an output pin to observe the result.

## 8. DESIGN AND IMPLEMENTATION OF DECODER

DATE:

### AIM:

To design and implement a 2 to 4 Decoder using logic gates.

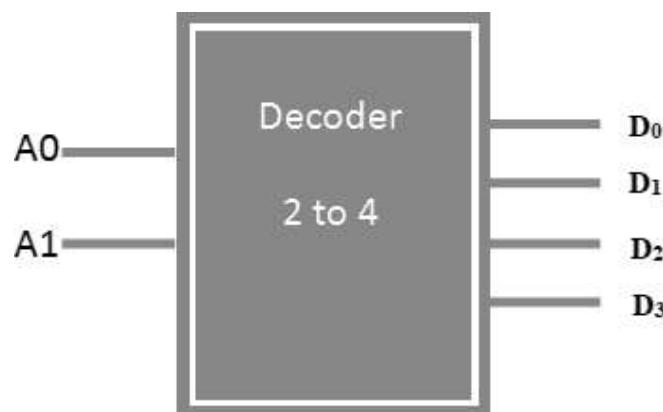
### APPARATUS REQUIRED:

SI.No.	COMPONENT	SPECIFICATION	QTY.
1.	AND GATE	IC 7408	1
2.	NOT GATE	IC7404	1
3.	IC TRAINER KIT	-	1
4.	CONNECTING WIRES	-	few

### PROCEDURE:

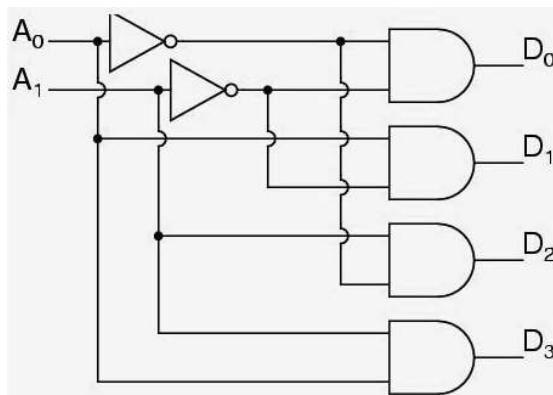
- a. Connections are given as per the logic diagram
- b. Input is given to the circuit making high '1' i.e. +5 or Vcc Supply to the 16<sup>th</sup> pin and for low '0' i.e. GND to the 8<sup>th</sup> pin of the Gate IC
- c. Depending upon the truth table, if the LED Glow it represents 1 and else it represents '0'
- d. Verify the truth table as given

### LOGIC DIAGRAM FOR 2:4 DECODER:



**TRUTH TABLE:**

INPUT		OUTPUT			
A1	A0	D0	D1	D2	D3
0	0	1	0	0	0
0	1	0	1	0	0
1	0	0	0	1	0
1	0	0	0	0	1

**CIRCUIT DIAGRAM FOR DECODER:****RESULT:**

Thus, the 2 to 4 Decoder circuit was designed using logic gates and outputs are verified using Truth table.

**Test Case 1:**

A = 0

Output should be Y0 = 1, Y1 = 0.

**Test Case 2:**

A = 1

Output should be Y0 = 0, Y1 = 1.

Implementation using Digital Trainer Kit:

Connect the input (A) to a switch or input on the trainer kit.

Use the trainer's logic gates (AND, NOT) to implement the above logic.

Observe the outputs using LEDs or output pins to check the values of Y0 and Y1 based on the input A.

## 9. DESIGN AND IMPLEMENTATION OF FLIP-FLOPS

DATE:

### AIM:

To Design and implement Flipflops (D, SR, JK & T) using logic gates.

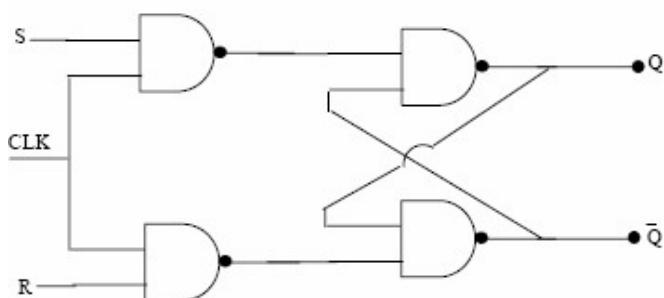
### APPARATUS REQUIRED:

SI.No.	COMPONENT	SPECIFICATION	QTY.
1.	AND GATE	IC 7408	1
2.	OR GATE	IC 7432	1
3.	NOT GATE	IC 7404	1
4.	IC TRAINER KIT	-	1
5.	CONNNECTING WIRES	-	32

### PROCEDURE:

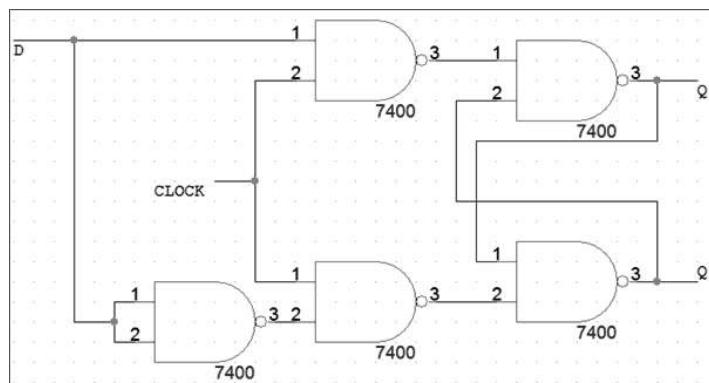
- Connections are given as per the logic diagram
- Input are given to the circuit making high '1' i.e. +5 or Vcc Supply to the 16<sup>th</sup> pin and for low '0' i.e. GND to the 8<sup>th</sup> pin of the Gate IC
- Depending upon the truth table, if the LED Glow it represents 1 and else it represents '0'
- Verify the truth table as given

### SR FLIPFLOP LOGIC DIAGRAM:



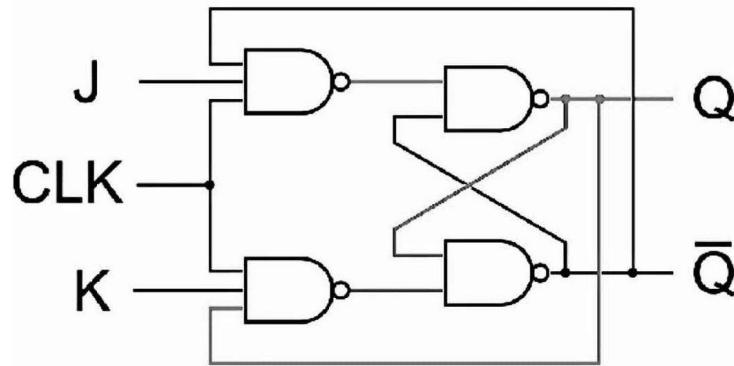
**TRUTH TABLE:**

Inputs			Outputs
R	S	CLK	Qt+1
0	0		Qt (Previous State)
0	1		1
1	0		0
1	1		Indeterminate state

**D FLIPFLOP**  
**LOGIC DIAGRAM:****TRUTH TABLE:**

Outputs		
D	CLK	Q
0		0
1		1

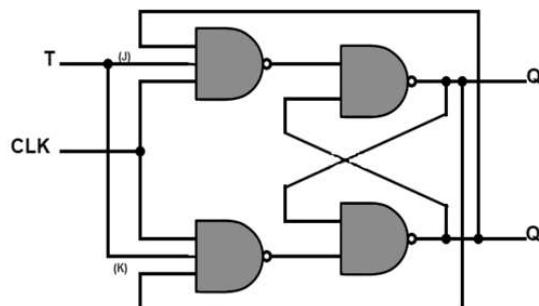
**JK FLIPFLOP****LOGIC DIAGRAM:**



**TRUTH TABLE:**

Truth Table			
Q	J	K	Q(t+1)
0	x	x	Qt (PreviousState)
1	0	0	Qt (PreviousState)
1	0	1	0
1	1	0	1
1	1	1	Qt'(Toggle)

**T FLIPFLOP LOGIC DIAGRAM:**



**TRUTH TABLE:**

Q	T	Q(t+1)
0	x	Qt (Previous State)
1	0	Qt (Previous State)
1	1	Qt'(Toggle)

**RESULT**

Thus, the of Flipflops (D, SR, JK & T) using logic gates were Designed and verified the truth tables.

**Test Case 1:**

D = 1, CLK = 1 (Rising Edge) → Q should become 1, Q\_bar should be 0.  
D = 0, CLK = 1 (Rising Edge) → Q should become 0, Q\_bar should be 1.

**Test Case 2:**

J = 1, K = 0 → Q should be set to 1 on the rising edge of CLK.  
J = 1, K = 1 → Q should toggle on the rising edge of CLK.

**Test Case 3:**

S = 1, R = 0 → Q should be set to 1 on the rising edge of CLK.  
S = 0, R = 1 → Q should be reset to 0 on the rising edge of CLK

**Test Case 4:**

T = 1 → Q should toggle on the rising edge of CLK.  
T = 0 → Q should remain unchanged on the rising edge of CLK.

## 10. DESIGN OF 2 – BIT SYNCHRONOUS COUNTER

DATE:

### AIM:

To Design and implement 2 – bit synchronous counter using Flip-Flops.

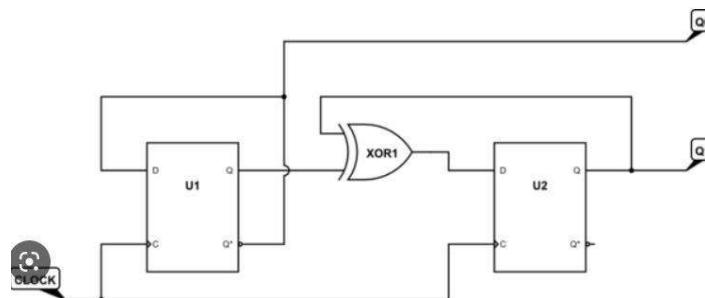
### APPARATUS REQUIRED:

SI. No	COMPONENT	SPECIFICATION	QTY.
1.	D FF	IC 7474	1
2.	OR GATE	IC 7432	1
3.	NOT GATE	IC 7404	1
4.	IC TRAINER KIT	-	1
5.	CONNECTING WIRES	-	few

### PROCEDURE:

- Connections are given as per the logic diagram
- Input is given to the circuit making high '1' i.e. +5 or Vcc Supply to the 16<sup>th</sup> pin and for low '0' i.e. GND to the 8<sup>th</sup> pin of the Gate IC
- Depending upon the truth table, if the LED Glow it represents 1 and else it represents '0'
- Verify the truth table as given

### LOGIC DIAGRAM:



### TRUTH TABLE:

CLK	QA	QB
1	0	0
2	0	1
3	1	0
4	1	1

### RESULT:

The 2-Bit counter sequential circuit was designed, and implemented and truth table was verified.

**Test Case 1:**

Initial state Q1 = 0, Q0 = 0

Clock Cycle 1:

Q0 toggles (Q0 = 1).

Q1 stays at 0 (because Q0 was 0 before).

Result: Q1 = 0, Q0 = 1.

Clock Cycle 2:

Q0 toggles again (Q0 = 0).

Q1 toggles (Q1 = 1, because Q0 transitioned from 1 to 0).

Result: Q1 = 1, Q0 = 0.

Clock Cycle 3:

Q0 toggles (Q0 = 1).

Q1 stays at 1 (because Q0 was 0 before).

Result: Q1 = 1, Q0 = 1.

Clock Cycle 4:

Q0 toggles (Q0 = 0).

Q1 toggles (Q1 = 0, because Q0 transitioned from 1 to 0).

Result: Q1 = 0, Q0 = 0.

Expected Output Sequence:

Q1 Q0 → 00, 01, 10, 11 (and repeat).

**Test Case 2:**

Initial state Q1 = 1, Q0 = 1

Clock Cycle 1:

Q0 toggles (Q0 = 0).

Q1 toggles (Q1 = 0, because Q0 transitioned from 1 to 0).

Result: Q1 = 0, Q0 = 0.

Clock Cycle 2:

Q0 toggles (Q0 = 1).

Q1 stays at 0 (because Q0 was 0 before).

Result: Q1 = 0, Q0 = 1.

Clock Cycle 3:

Q0 toggles (Q0 = 0).

Q1 toggles (Q1 = 1, because Q0 transitioned from 1 to 0).

Result: Q1 = 1, Q0 = 0.

Clock Cycle 4:

Q0 toggles (Q0 = 1).

Q1 stays at 1 (because Q0 was 0 before).

Result: Q1 = 1, Q0 = 1.

Expected Output Sequence:

Q1 Q0 → 11, 10, 01, 00 (and repeat).

Implementation with Digital Trainer Kit

D Flip-Flop Counter:

Connect two D flip-flops to the clock signal.

Set D0 = Q0' to ensure Q0 toggles on every clock pulse.

Set D1 = Q0 AND Q1 to ensure Q1 toggles every time Q0 transitions from 1 to 0.

Connect LEDs or other indicators to observe the states of Q1 and Q0.

## **11. DESIGN AND IMPLEMENTATION OF SHIFT REGISTERS**

**DATE:**

### **AIM:**

To design and implement

- A. Serial in serial out
- B. Serial in parallel out

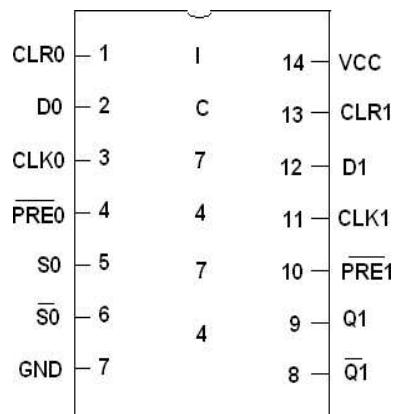
### **APPARATUS REQUIRED:**

SI.No.	COMPONENT	SPECIFICATION	QTY.
1.	D FLIP FLOP	IC 7474	2
2.	IC TRAINER KIT	-	1
3.	CONNECTING WIRES	-	35

### **PROCEDURE:**

- a. Connections are given as per the logic diagram
- b. Input is given to the circuit making high '1' i.e., +5 or Vcc Supply to the 14<sup>th</sup> pin and for low '0' i.e., GND to the 7<sup>th</sup> pin of the Gate IC
- c. Depending upon the truth table, if the LED Glow it represents 1 and else it represents '0'
- d. Verify the truth table as given

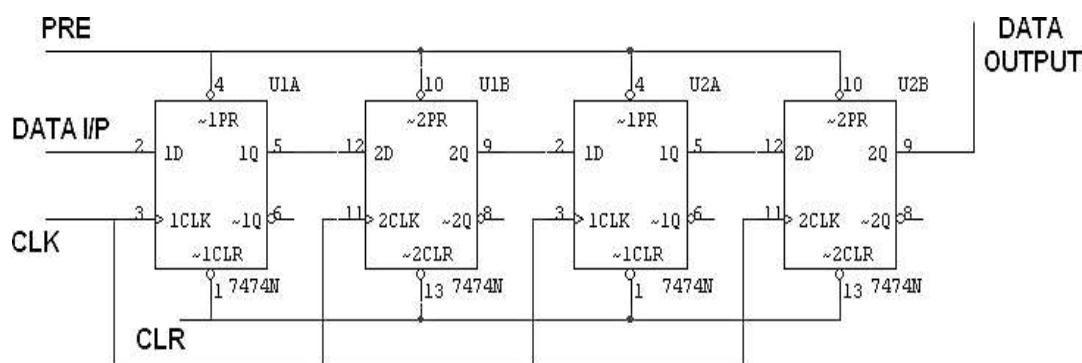
**PIN DIAGRAM:**



**SERIAL IN SERIAL OUTTRUTH TABLE:**

CLK	Serial in	Serial out
1	1	0
2	0	0
3	0	0
4	1	1
5	X	0
6	X	0
7	X	1

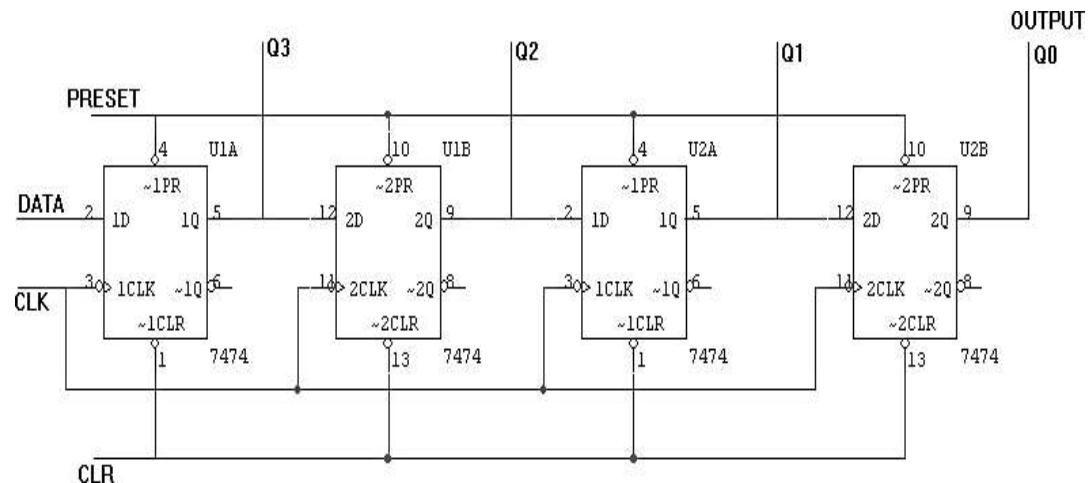
**LOGIC DIAGRAM:**



**SERIAL IN PARALLEL OUTTRUTH TABLE:**

CLK	DATA	OUTPUT			
		QA	QB	QC	QD
1	1	1	0	0	0
2	0	0	1	0	0
3	0	0	0	1	0
4	1	1	0	0	1

**LOGIC DAIGRAM:**



**RESULT:**

Thus, the shift register using D Flip Flop is constructed and the outputs are verified using truth table.

**Test Case 1:**

Initial state: Q1 = 0, Q2 = 0, Q3 = 0 (all flip-flops cleared).

Data input: Serial input = 101 (to be shifted in one bit at a time).

Clock pulses: Apply 6 clock pulses.

After the first clock pulse, Q1 will become 1, and Q2 and Q3 will remain 0. After each subsequent clock pulse, the data shifts right:

After Clock 1: Q1 = 1, Q2 = 0, Q3 = 0

After Clock 2: Q1 = 1, Q2 = 0, Q3 = 0

After Clock 3: Q1 = 1, Q2 = 1, Q3 = 0

After Clock 4: Q1 = 1, Q2 = 1, Q3 = 0

After Clock 5: Q1 = 0, Q2 = 1, Q3 = 1

After Clock 6: Q1 = 1, Q2 = 0, Q3 = 1

Expected Output:

After 6 clock pulses, the output from Q3 will be 1, and the shifting sequence should be 1, 0, 1, 1, 0, 1.

**Test Case 2:**

Initial state: Q1 = 0, Q2 = 0, Q3 = 0.

Data input: Serial input = 110 (to be shifted in).

Clock pulses: Apply 5 clock pulses.

After each clock pulse:

After Clock 1: Q1 = 1, Q2 = 0, Q3 = 0

After Clock 2: Q1 = 1, Q2 = 1, Q3 = 0

After Clock 3: Q1 = 1, Q2 = 1, Q3 = 0

After Clock 4: Q1 = 0, Q2 = 1, Q3 = 1

After Clock 5: Q1 = 1, Q2 = 0, Q3 = 1

Expected Output:

After 5 clock pulses, the parallel output should be Q1 = 1, Q2 = 0, Q3 = 1.

## 12. SIMULATION OF LOGIC GATES USING VHDL

DATE:

### AIM:

To write the source code and the simulation of the logic gates of AND, OR and NOT using VHDL.

### SOFTWARE REQUIRED:

ModelSim Simulator

### PROCEDURE:

- Step1: Define the specifications and initialize the design.
- Step2: Declare the name of the module by using source code.
- Step3: Write the source code in VHDL.
- Step4: Check the syntax and debug the errors if found, obtain the synthesis is report.
- Step5: Verify the output by simulating the source code.

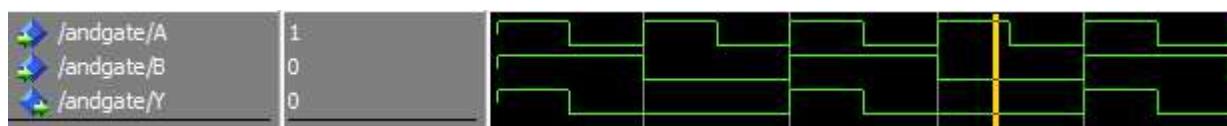
### LOGIC GATES:

#### AND GATE:

##### VHDL CODE:

```
library IEEE;
use IEEE.std_logic_1164.all;
entity andGate is
port(A : in std_logic;      -- AND gate input B : in std_logic;      -- AND gate input Y : out std_logic); -
- AND gate output
end andGate;
architecture andLogic of andGate isbegin
    Y <= A AND B;
end andLogic;
```

##### OUTPUT WAVEFORM:



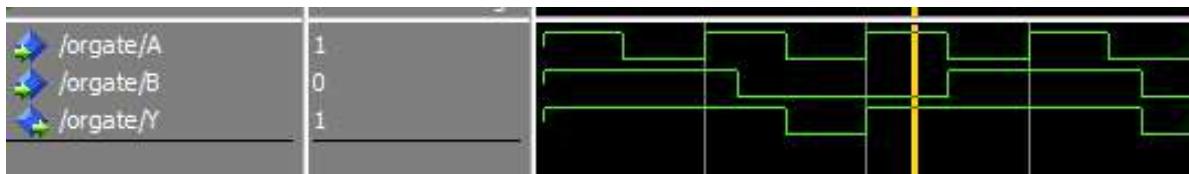
### OR GATE:

#### VHDL CODE:

```
library IEEE;
use IEEE.std_logic_1164.all;
entity orGate is

    port(A : in std_logic;      -- OR gate input B : in std_logic;
          -- OR gate input Y : out std_logic); -- OR gate output
end orGate;
architecture orLogic of orGate is begin
    Y <= A OR B;
end orLogic;
```

#### OUTPUT:

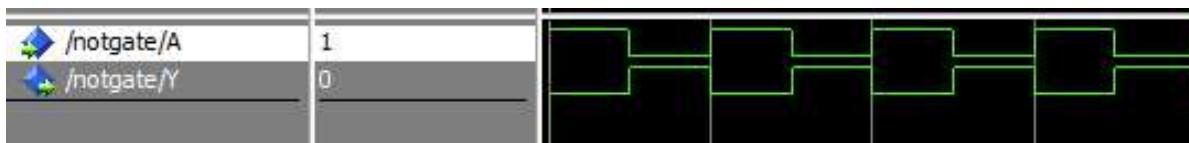


### NOT GATE:

#### VHDL CODE:

```
library IEEE;
use IEEE.std_logic_1164.all;
entity notGate is
    port(A : in std_logic;      -- NOT gate input
         Y : out std_logic); -- NOT gate output
end notGate;
architecture notLogic of notGate isbegin
    Y <= not A;
end notLogic;
```

#### OUTPUT WAVEFORM:



### RESULT

Thus, the VHDL code for logic gates were written, executed and verified the output.

### **Test Case 1: AND Gate**

Test Inputs:

Test case 1: A = 0, B = 0 → Expected Output: Y = 0

Test case 2: A = 1, B = 0 → Expected Output: Y = 0

Test case 3: A = 0, B = 1 → Expected Output: Y = 0

Test case 4: A = 1, B = 1 → Expected Output: Y = 1

Expected Result: The AND gate outputs 1 only when both inputs are 1, otherwise, it outputs 0.

### **Test Case 2: OR Gate**

Test Inputs:

Test case 1: A = 0, B = 0 → Expected Output: Y = 0

Test case 2: A = 1, B = 0 → Expected Output: Y = 1

Test case 3: A = 0, B = 1 → Expected Output: Y = 1

Test case 4: A = 1, B = 1 → Expected Output: Y = 1

Expected Result: The OR gate outputs 1 when at least one of the inputs is 1, otherwise, it outputs 0.

### **Test Case 3: NOT Gate**

Test Inputs:

Test case 1: A = 0 → Expected Output: Y = 1

Test case 2: A = 1 → Expected Output: Y = 0

Expected Result: The NOT gate outputs the inverse of the input value (complement of A).

### **How to Run the Simulation:**

Write the VHDL Code: Write the entity and architecture for each gate (and\_gate.vhdl, or\_gate.vhdl, not\_gate.vhdl).

Create the Testbenches: Create the corresponding testbenches (and\_gate\_tb.vhdl, or\_gate\_tb.vhdl, not\_gate\_tb.vhdl).

Compile the Code: Use an HDL simulator (like ModelSim, Xilinx Vivado, or any other simulator) to compile both the logic gate code and the testbenches.

Run the Simulation: Simulate the design and observe the waveforms to verify the output of each gate.

Verify the Results: Ensure the outputs match the expected results from the truth tables provided above.

### **13. SIMULATION OF ADDER CIRCUITS USING VHDL**

**DATE:**

#### **AIM:**

To write the source code and the simulate the Half adder and Full adder using VHDL.

#### **SOFTWARE REQUIRED:**

ModelSim Simulator

#### **PROCEDURE:**

Step1: Define the specifications and initialize the design.

Step2: Declare the name of the module by using source code.

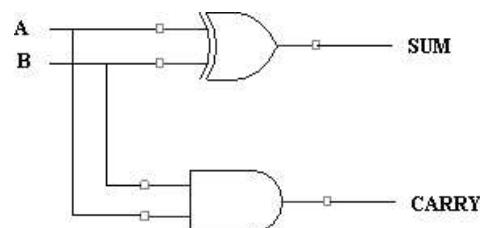
Step3: Write the source code in VHDL.

Step4: Check the syntax and debug the errors if found, obtain the synthesis report.

Step5: Verify the output by simulating the source code.

### HALF ADDER:

#### LOGIC DIAGRAM:



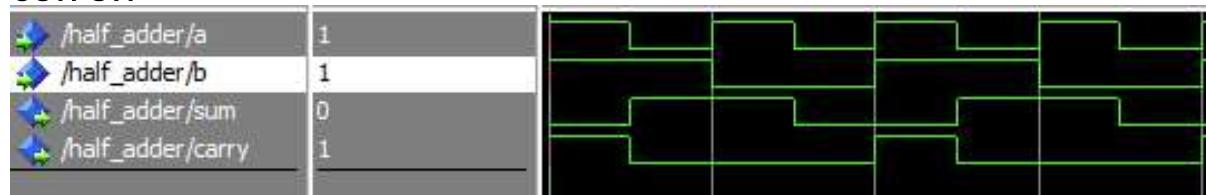
#### TRUTH TABLE:

A	B	SUM	CARRY
0	0	0	0
0	1	1	0

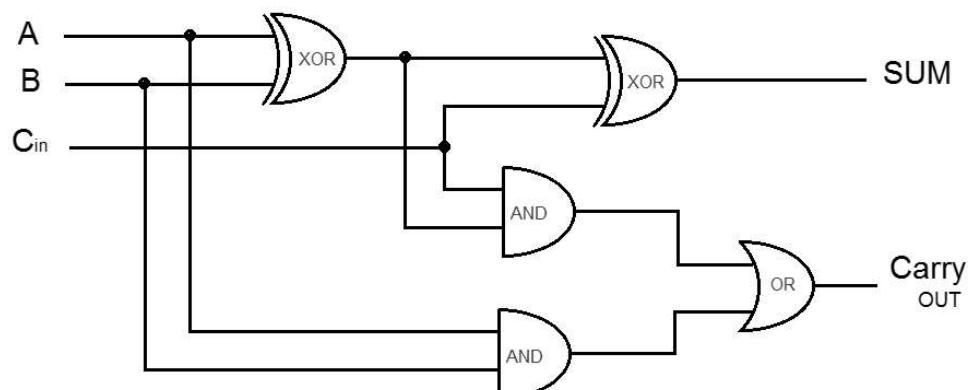
#### VHDL CODE:

```
library ieee;
use ieee.std_logic_1164.all;
entity half_adder is
port(a,b:in bit; sum,carry:out bit);end half_adder;
architecture data of half_adder is begin
sum<= a xor b; carry <= a
and b;
end data;
```

#### OUTPUT:

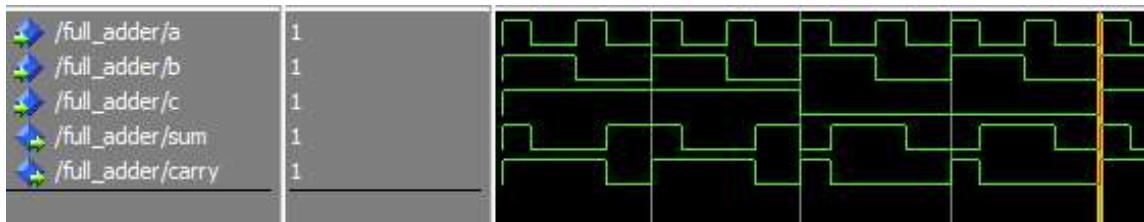


### FULL ADDER Circuit Diagram:



**VHDL CODE:**

```
Library ieee;
use ieee.std_logic_1164.all;
entity full_adder is port (a,b,c:in bit; sum,carry:out bit);end
full_adder;
architecture data of full_adder is begin
sum<= a xor b xor c;
carry <= ((a and b) or (b and c) or (a and c));end data;
```

**Output:****RESULT**

Thus, the VHDL code for half adder and full adder were written, executed and verified the output.

### **Test Case 1: Half Adder**

Inputs:

Test case 1: A = 0, B = 0

Test case 2: A = 0, B = 1

Test case 3: A = 1, B = 0

Test case 4: A = 1, B = 1

Expected Outputs (Sum, Carry):

For A = 0, B = 0 → Sum = 0, Carry = 0

For A = 0, B = 1 → Sum = 1, Carry = 0

For A = 1, B = 0 → Sum = 1, Carry = 0

For A = 1, B = 1 → Sum = 0, Carry = 1

### **Test Case 2: Full Adder**

Inputs:

Test case 1: A = 0, B = 0, Cin = 0

Test case 2: A = 1, B = 1, Cin = 0

Test case 3: A = 1, B = 0, Cin = 1

Test case 4: A = 1, B = 1, Cin = 1

Expected Outputs (Sum, Carry Out):

For A = 0, B = 0, Cin = 0 → Sum = 0, Carry Out = 0

For A = 1, B = 1, Cin = 0 → Sum = 0, Carry Out = 1

For A = 1, B = 0, Cin = 1 → Sum = 0, Carry Out = 1

For A = 1, B = 1, Cin = 1 → Sum = 1, Carry Out = 1

## **How to Simulate the Design**

**Write the VHDL Code:** Write the VHDL code for both the Half Adder and Full Adder (entities and architectures).

**Create the Testbenches:** Create testbenches for each adder.

**Compile the Code:** Use a VHDL simulator like ModelSim, Xilinx Vivado, or Quartus to compile both the adder design and the testbenches.

**Run the Simulation:** Run the simulation and check the output waveform to verify the correctness of the adders.

**Verify the Results:** Ensure that the outputs from the testbenches match the expected values as per the truth tables above.

## **14. SIMULATION OF SUBTRACTOR CIRCUITS USING VHDL**

**DATE:**

### **AIM:**

To write the source code and simulate the Half subtractor and Full subtractor using VHDL.

### **SOFTWARE REQUIRED:**

ModelSim Simulator

### **PROCEDURE:**

Step1: Define the specifications and initialize the design.

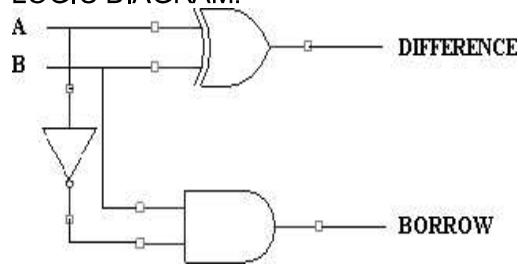
Step2: Declare the name of the module by using source code.

Step3: Write the source code in VHDL.

Step4: Check the syntax and debug the errors if found, obtain the synthesis is report.

### HALF SUBTRACTOR:

#### LOGIC DIAGRAM:



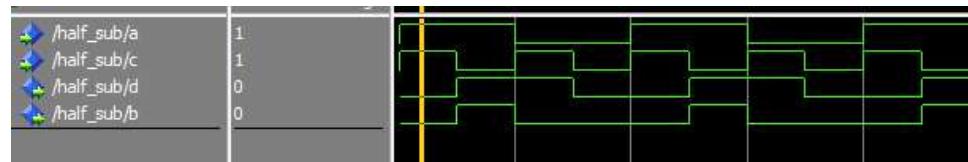
#### TRUTH TABLE:

A	B	DIFFERENCE	BORROW
0	0	0	0
0	1	1	1

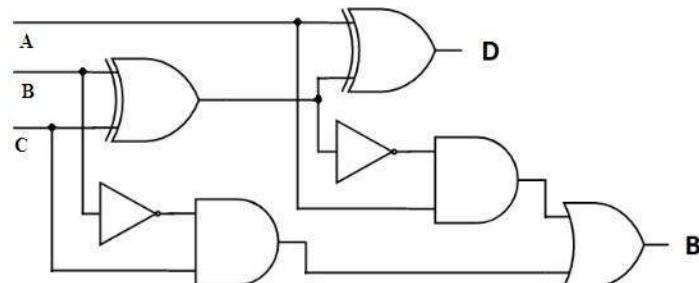
#### VHDL CODE:

```
Library ieee;
use ieee.std_logic_1164.all; entity half_sub
is port(a,c:in bit; d,b:out bit);
end half_sub;
architecture data of half_sub is begin
d<= a xor c;
b<= (a and (not c));end data;
```

#### OUTPUT:



### FULL SUBTRACTOR CIRCUIT DIAGRAM:

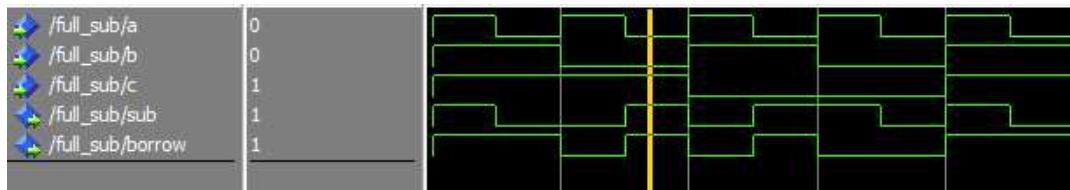


#### TRUTH TABLE

A	B	C	DIFFERENCE	BORROW
0	0	0	0	0
0	0	1	1	1
0	1	0	1	1
0	1	1	0	1
1	0	0	1	0
1	0	1	0	0
1	1	0	0	0
1	1	1	1	1

**VHDL CODE:**

```
Library ieee;
use ieee.std_logic_1164.all;entity
full_sub is
    port (a,b,c:in bit; sub,borrow:out bit);end
full_sub;
architecture data of full_sub isbegin
sub<= a xor b xor c;
borrow <= ((b xor c) and (not a)) or (b and c);end
data;
```

**Output:****RESULT**

Thus, the VHDL code for half subtractor and full subtractor were written, executed and verified the output.

### **Test Case 1: Half Subtractor**

Inputs:

Test case 1: A = 0, B = 0

Test case 2: A = 0, B = 1

Test case 3: A = 1, B = 0

Test case 4: A = 1, B = 1

Expected Outputs (Difference, Borrow):

For A = 0, B = 0 → Difference = 0, Borrow = 0

For A = 0, B = 1 → Difference = 1, Borrow = 1

For A = 1, B = 0 → Difference = 1, Borrow = 0

For A = 1, B = 1 → Difference = 0, Borrow = 0

### **Test Case 2: Full Subtractor**

Inputs:

Test case 1: A = 0, B = 0, Bin = 0

Test case 2: A = 1, B = 1, Bin = 0

Test case 3: A = 1, B = 0, Bin = 1

Test case 4: A = 1, B = 1, Bin = 1

Expected Outputs (Difference, Borrow):

For A = 0, B = 0, Bin = 0 → Difference = 0, Borrow = 0

For A = 1, B = 1, Bin = 0 → Difference = 0, Borrow = 0

For A = 1, B = 0, Bin = 1 → Difference = 0, Borrow = 1

For A = 1, B = 1, Bin = 1 → Difference = 1, Borrow = 1

## **How to Simulate the Design**

Write the VHDL Code: Write the VHDL code for both the Half Subtractor and Full Subtractor (entities and architectures).

Create the Testbenches: Create testbenches for each subtractor.

Compile the Code: Use a VHDL simulator like ModelSim, Xilinx Vivado, or Quartus to compile both the subtractor design and the testbenches.

Run the Simulation: Run the simulation and check the output waveform to verify the correctness of the subtractors.

Verify the Results: Ensure that the outputs from the test benches match the expected values as per the truth tables above.

## **15. SIMULATION OF MUXIPLEXER AND DEMUXIPLEXER USING VHDL**

**DATE**

### **AIM:**

To write the source code and the simulate the 4x1 Multiplexer and 1x4 Demultiplexer using VHDL.

### **SOFTWARE REQUIRED:**

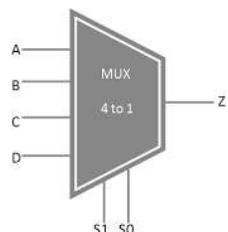
ModelSim Simulator

## PROCEDURE:

- Step1: Define the specifications and initialize the design.
- Step2: Declare the name of the module by using source code.
- Step3: Write the source code in VHDL.
- Step4: Check the syntax and debug the errors if found, obtain the synthesis report.
- Step5: Verify the output by simulating the source code.

## 1x4 MUX:

LOGIC DIAGRAM:



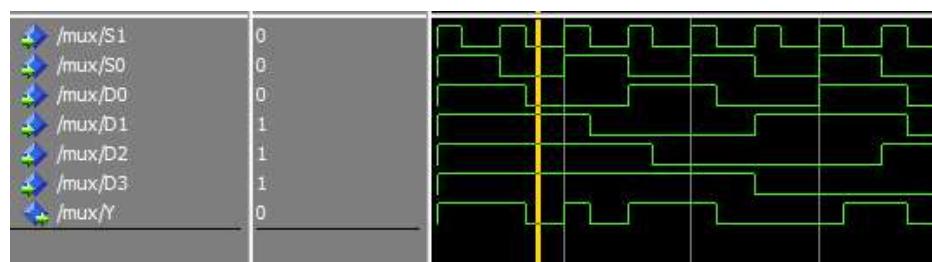
TRUTH TABLE:

S1	S0	Output
0	0	A
0	1	B
-	-	-

## VHDL CODE:

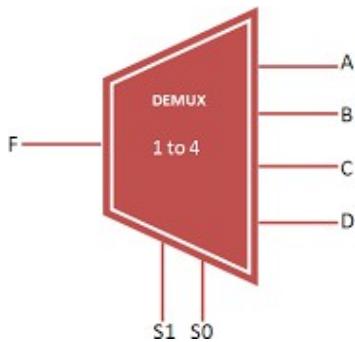
```
library ieee;
use
    ieee.std_logic_1164.all;
entity mux is
    port(S1,S0,D0,D1,D2,D3:in bit; Y:out bit);
end mux;
architecture data of mux
isbegin
    Y<= (not S0 and not S1 and D0) or
        (S0 and not S1 and D1) or
        (not S0 and S1 and D2)
        or(S0 and S1 and D3);
end data;
```

## OUTPUT:



#### 4x1 DEMUTIPLEXER:

##### LOGIC DIAGRAM:

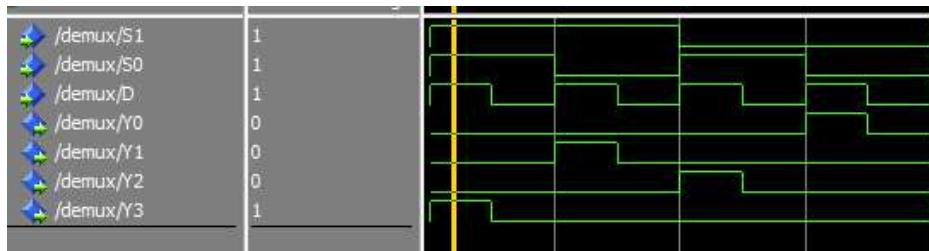


D	INPUT		OUTPUT			
	S0	S1	Y0	Y1	Y2	Y3
1	0	0	1	0	0	0
1	0	1	0	1	0	0
1	1	0	0	0	1	0
1	1	1	0	0	0	1

##### VHDL CODE:

```
Library ieee;
use
  ieee.std_logic_1164.all;
entity demux is
  port(S1,S0,D:in bit; Y0,Y1,Y2,Y3:out bit);
end demux;
architecture data of demux
isbegin
  Y0<= ((Not S0) and (Not S1) and D);
  Y1<= ((Not S0) and S1 and D);
  Y2<= (S0 and (Not S1) and D);
  Y3<= (S0 and S1 and D);
end data;
```

**Output:**



**RESULT**

Thus, the VHDL code for Multiplexer and Demultiplexer were written, executed and verified the output.

### **Test Case 1: 4x1 Multiplexer**

Test Inputs:

Test case 1:  $S_0 = 0, S_1 = 0, I_0 = 1, I_1 = 0, I_2 = 0, I_3 = 0$

Test case 2:  $S_0 = 1, S_1 = 1, I_0 = 0, I_1 = 0, I_2 = 0, I_3 = 1$

Expected Output:

For  $S_0 = 0, S_1 = 0 \rightarrow Y = I_0 = 1$

For  $S_0 = 1, S_1 = 1 \rightarrow Y = I_3 = 1$

### **Test Case 2: 1x4 Demultiplexer**

Test Inputs:

Test case 1:  $D = 1, S_0 = 0, S_1 = 0$

Test case 2:  $D = 1, S_0 = 1, S_1 = 1$

Expected Output:

For  $S_0 = 0, S_1 = 0 \rightarrow O_0 = 1, O_1 = O_2 = O_3 = 0$

For  $S_0 = 1, S_1 = 1 \rightarrow O_3 = 1, O_0 = O_1 = O_2 = 0$

## **How to Simulate the Design**

Write the VHDL Code: Write the VHDL code for both the 4x1 Multiplexer and 1x4 Demultiplexer (entities and architectures).

Create the Testbenches: Create testbenches for both the multiplexer and demultiplexer.

Compile the Code: Use a VHDL simulator such as Model Sim, Xilinx Vivado, or Quartus to compile the designs and test benches.

Run the Simulation: Execute the simulations and observe the waveform to verify that the outputs match the expected results.

Verify the Results: Compare the simulation output against the expected truth tables for each test case.

## 16. SIMULATION OF ENCODER AND DECODER USING VHDL

DATE:

### AIM:

To write the source code and simulate the encoder and decoder using VHDL.

### SOFTWARE REQUIRED:

ModelSim Simulator

### PROCEDURE:

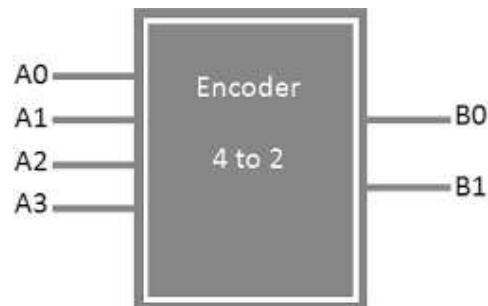
- Step1: Define the specifications and initialize the design.
- Step2: Declare the name of the module by using source code.
- Step3: Write the source code in VHDL.
- Step4: Check the syntax and debug the errors if found, obtain the synthesis report.
- Step5: Verify the output by simulating the source code.

### 4x2 ENCODER:

TRUTH TABLE:

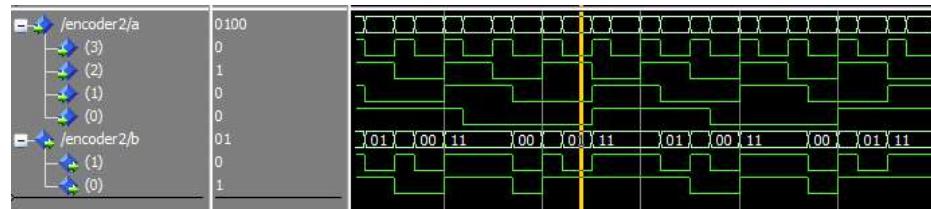
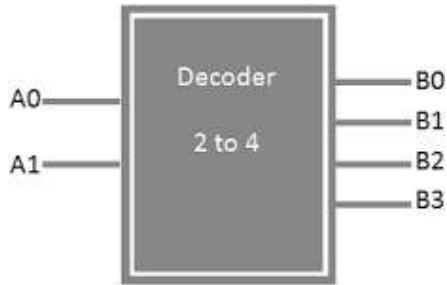
LOGIC DIAGRAM:

INPUT				OUTPUT	
A3	A2	A1	A0	B1	B0
0	0	0	1	0	0
0	0	1	0	0	1
0	1	0	0	1	0
1	0	0	0	1	1



### VHDL CODE:

```
library IEEE;
use IEEE.STD_LOGIC_1164.all;
entity encoder2 is
port(
a : in STD_LOGIC_VECTOR(3 downto 0); b
: out STD_LOGIC_VECTOR(1 downto 0)
);
end encoder2;
architecture bhv of encoder2
isbegin
b(0) <= a(1) or a(2);
b(1) <= a(1) or a(3);
end bhv;
```

**OUTPUT:****4x2 DECODER:****LOGIC DIAGRAM:****TRUTH TABLE:**

INPUT		OUTPUT			
A1	A0	B3	B2	B1	B0
0	0	0	0	0	1
0	1	0	0	1	0
1	0	0	1	0	0
1	1	1	0	0	0

**VHDL CODE:**

```
library IEEE;
use IEEE.STD_LOGIC_1164.all;
entity decoder2 is port( a : in STD_LOGIC_VECTOR(1 downto 0); b : out STD_LOGIC_VECTOR(3 downto 0));
end decoder2;
architecture bhv of decoder2 is
begin
b(0) <= not a(0) and not a(1);
b(1) <= not a(0) and a(1);
b(2) <= a(0) and not a(1);
b(3) <= a(0) and a(1);
end bhv;
```

## OUTPUT:



## RESULT

Thus, the VHDL code for Encoder and Decoder were written, executed and verified the output.

### **Test Case 1: 4x2 Encoder**

Test Inputs:

Test case 1: I<sub>3</sub> = 1, I<sub>2</sub> = 0, I<sub>1</sub> = 0, I<sub>0</sub> = 0

Test case 2: I<sub>3</sub> = 0, I<sub>2</sub> = 1, I<sub>1</sub> = 0, I<sub>0</sub> = 0

Expected Output:

For I<sub>3</sub> = 1, I<sub>2</sub> = 0, I<sub>1</sub> = 0, I<sub>0</sub> = 0 → A = 1, B = 1, V = 1

For I<sub>3</sub> = 0, I<sub>2</sub> = 1, I<sub>1</sub> = 0, I<sub>0</sub> = 0 → A = 1, B = 0, V = 1

### **Test Case 2: 2x4 Decoder**

Test Inputs:

Test case 1: A = 0, B = 0

Test case 2: A = 1, B = 1

Expected Output:

For A = 0, B = 0 → O<sub>0</sub> = 1, O<sub>1</sub> = O<sub>2</sub> = O<sub>3</sub> = 0

For A = 1, B = 1 → O<sub>3</sub> = 1, O<sub>0</sub> = O<sub>1</sub> = O<sub>2</sub> = 0

### **How to Simulate the Design**

Write the VHDL Code: Write the VHDL code for both the 4x2 Encoder and 2x4 Decoder (entities and architectures).

Create the Testbenches: Create testbenches for both the encoder and decoder.

Compile the Code: Use a VHDL simulator such as Model Sim, Xilinx Vivado, or Quartus to compile the designs and testbenches.

Run the Simulation: Execute the simulations and observe the waveform to verify that the outputs match the expected results.

Verify the Results: Compare the simulation output against the expected truth tables for each test case.