EXP NO: 08 DATE:

GENERATE THREE ADDRESS CODE FOR A SIMPLE PROGRAM USING LEX AND YACC

AIM:

To design and implement a **LEX and YACC** program that generates **three-address code** (**TAC**) for a simple arithmetic expression or program. The program will:

- Recognize expressions like addition, subtraction, multiplication, and division.
- Generate **three-address code** that represents the operations in a way that could be directly translated into assembly code or intermediate code for a compiler.

ALGORITHM:

1. Lexical Analysis (LEX) Phase:

Input: A string containing an arithmetic expression (e.g., a = b + c * d;).

Output: A stream of tokens such as identifiers (variables), numbers (constants), operators, and special characters (like =, ;, (), etc.).

1. Define the Token Patterns:

- o **ID:** Identifiers (variables) are strings starting with a letter and followed by letters or digits (e.g., a, b, result).
- o **NUMBER:** Constants (e.g., 1, 5, 100).
- o **OPERATOR:** Arithmetic operators (+, -, *, /).
- o ASSIGNMENT: Assignment operator (=).
- o PARENTHESIS: Parentheses for grouping ((and)).
- WHITESPACE: Spaces, tabs, and newline characters (which should be ignored).

2. Write Regular Expressions for the Tokens:

- \circ ID -> [a-zA-Z][a-zA-Z0-9]*
- o NUMBER -> [0-9]+
- \circ OPERATOR -> [\+\-*/]
- o ASSIGN -> "="
- PAREN -> [\(\)]
- WHITESPACE -> [\t\n]+ (skip whitespace)

3. Action on Tokens:

- When a token is matched, pass it to YACC using yylval to store the token values.
- 2. Syntax Analysis and TAC Generation (YACC) Phase:

Input: Tokens provided by the **LEX** lexical analyzer.

Output: Three-address code for the given arithmetic expression.

1. Define Grammar Rules:

Assignment:

bash CopyEdit statement: ID '=' expr

This means an expression is assigned to a variable.

o Expressions:

bash CopyEdit expr: expr OPERATOR expr

An expression can be another expression with an operator (+, -, *, /).

bash CopyEdit expr: NUMBER expr: ID expr: '(' expr ')'

2. Three-Address Code Generation:

- For every arithmetic operation, generate a temporary variable (e.g., t1, t2, etc.) to hold intermediate results.
- o For a = b + c, generate:

ini CopyEdit t1 = b + c a = t1

• For a = b * c + d, generate:

ini CopyEdit t1 = b * c t2 = t1 + da = t2

3. Temporary Variable Management:

- Keep a counter (temp_count) for generating unique temporary variable names (t0, t1, t2, ...).
- Each time a new operation is encountered, increment the temp_count to generate a new temporary variable.

4. Rule Actions:

• When a rule is matched (e.g., expr OPERATOR expr), generate the TAC and assign temporary variables for intermediate results.

Detailed Algorithm:

1. Initialize Lexical Analyzer:

o Define the token patterns for ID, NUMBER, OPERATOR, ASSIGN, PAREN, and WHITESPACE.

2. Define the Syntax Grammar:

- o Define grammar rules for:
 - Assignments: ID = expr
 - Expressions: expr -> expr OPERATOR expr, expr -> NUMBER, expr
 -> ID, expr -> (expr)

3. Token Matching:

- o LEX: Match input characters against the defined regular expressions for tokens
- o YACC: Use the tokens to parse and apply grammar rules.

4. TAC Generation:

- o For Assignment:
 - Upon parsing ID = expr, generate a temporary variable for the result of expr and assign it to the variable ID.
- o For Arithmetic Operations:
 - For each operator (e.g., +, -, *, /), generate temporary variables for intermediate calculations.

5. Output TAC:

Print the generated three-address code, with each expression and its intermediate results represented by temporary variables.

PROGRAM:

```
3address.1
%{
#include "3address.tab.h"
#include <string.h>
#include <stdlib.h>
%}
     [a-zA-Z][a-zA-Z0-9]*
NUM [0-9]+
%%
{ID} { yylval.str = strdup(yytext); return ID; }
{NUM} { yylval.str = strdup(yytext); return NUM; }
     { return '='; }
      { return ';'; }
"("
      { return '('; }
      { return ')'; }
      { return '+'; }
```

```
"_"
      { return '-'; }
      { return '*'; }
!!*!!
      { return '/'; }
[\t\n]; // skip whitespace
%%
int yywrap()
  { return 1;
3address.y
%{
#include <stdio.h>
#include <stdlib.h>
#include <string.h>
int tempCount = 0;
char* createTemp() {
  char* temp = (char*)malloc(10);
  sprintf(temp, "t%d", tempCount++);
  return temp;
void yyerror(const char* s);
int yylex();
%}
%union
   { char* str;
%token <str> ID NUM
%type <str> expr
%left '+' '-'
%left '*' '/'
%%
stmt:
    ID '=' expr ';' {
      printf("%s = %s\n", $1, $3);
expr:
    expr '+' expr {
```

```
char* temp = createTemp();
      printf("%s = %s + %s\n", temp, $1, $3);
      $$ = temp;
  expr'-' expr {
      char* temp = createTemp();
      printf("%s = %s - %s \n", temp, $1, $3);
      $ = temp;
  expr '*' expr {
      char* temp = createTemp();
      printf("% = % * % \n", temp, $1, $3);
      $ = temp;
  expr'/' expr {
      char* temp = createTemp();
      printf("%s = %s / %s\n", temp, 1, 3);
      $ = temp;
  | '(' expr ')' {
      $$ = $2;
  | ID {
      $\$ = strdup(\$1);
  | NUM {
      \$\$ = strdup(\$1);
%%
void yyerror(const char* s)
   { printf("Syntax Error: %s\n", s);
int main() {
  printf("Enter an arithmetic expression :\n");
  yyparse();
  return 0;
```

}

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OUTPUT:

```
yacc -d expr.y
lex expr.l
gcc y.tab.c lex.yy.c -o expr_parser
/expr_parser
a = b * c + d;
t0 = b * c
t1 = t0 + d
a = t1
```

Implementation	
Output/Signature	

RESULT:

Thus the process effectively tokenizes the input, parses it according to defined grammar rules, and generates the corresponding Three-Address Code, facilitating further compilation or interpretation stages.