



*Assignment 25: Dijkstra Verification*  
*Due in class Thursday, 3/28*

March 25, 2019  
CS: DS&A  
PROOF SCHOOL

In this short assignment, you'll give a careful proof that Dijkstra's algorithm works. (The next assignment will be a longer coding assignment.) Recall the algorithm:

Start with  $A$ , and first relax along edges coming out of  $A$ . Then, pick an unprocessed vertex with the smallest current  $d$ -value. Call that node processed, and relax along edges coming out of it. Keep going.

**Question 1.** First, give a quick argument that every vertex  $v$  gets relaxed at some point; i.e. no vertex ends up with  $d(v) = \infty$ . (Hint: Suggested first line: We know  $G$  is connected, so there is some path from  $A$  to  $v$ .)

**Question 2.** Now, here's the setup for the main part of the argument. Suppose we're in the middle of the algorithm; let  $P$  be the set of processed nodes (so  $V \setminus P$  is the set of unprocessed nodes.) By induction on stages, we can assume that  $d(x) = D(x)$  for all nodes  $x \in P$ . (When we use  $d$  here, we mean the  $d$ -value at this particular stage.) We've just chosen  $v \in V \setminus P$  to process. By question 1,  $v$  has already been relaxed at some point; suppose it was relaxed most recently by  $u \in P$  (i.e. suppose the current predecessor of  $v$  is  $u$ ). We claim that at this point  $d(v) = D(v)$ . We prove this by contradiction: suppose that  $D(v) < d(v)$ .

Your task for this assignment is to finish the argument. But to make things more clear, *I want you to do this in the style of a two-column argument. Every single statement must be stated separately, and given a reason, unless it's just establishing notation.* For example, "Let  $w$  be the weight of the edge from  $u$  to  $v$ " doesn't need a reason. Something like " $D(v') = d(v')$ " definitely needs a reason, though.

You may use the fact, discussed in class, that if  $A - v_1 - \dots - v_k - v$  is a true path from  $A$  to  $v$ , then  $A - v_1 - \dots - v_i$  is a true path from  $A$  to  $v_i$  for each  $i \leq k$ .