


Square Roots

Theorem Let (F, P) be an ordered field.

Let $a \in F$. (0 is always its only square root.)

Define, $S_a := \{x \in F \mid x^2 \leq a\}$

Then, $S_a \neq \emptyset$ (as 0 is in the set)

S_a is bounded above (by $a+1$)

now we must show

(3) $\sup(S_a)$ exists $\iff x^2 = a$ has a solution
in F and so $x^2 = \sup(S_a)$

PF] \Leftarrow

$x^2 = a$ has a solution in F . Let $b \in F$ be one such solution. As $(-b)^2 = b^2$, wlog we can say $b \in F$. (note $b \neq 0$ as $a > 0$).

Note: $s \in S_a$ implies

$$s^2 \leq a = b^2 \text{ so } s^2 \leq b^2 \text{ by hwz}$$

$$s \leq |s| \leq b$$

That is $\forall s \in S_a, s \leq b$, so b is an upper bound for S_a .

Furthermore, $b \in S_a$ as $b^2 = a \leq a$

So, b is a maximal element of S_a .

So, $b = \sup(S_a)$.



\Rightarrow

assume $\sup(S_a)$ exists and let $b = \sup(S_a)$
we want to show $b^2 = a$ why?

What can we say about b ?

As $0 \in S_a$ (since $0 < a^2 \nexists a \in P$)

So, $b \geq 0$.

We will show $b \neq 0$.

Claim: if $b = 0$, then

$S_a = \{0\}$ \rightarrow if another element exists in S_a it is positive so $\sup(S_a)$ is positive. $\rightarrow (-b)^2 = b^2$
lets now show $S_a \neq \{0\}$.

Recall] $a > 0 \wedge$ if $a > 1$ then $a^2 > a > 1 = 1^2$

so, $1 \in S_a$

if $a \leq 1$, then $a^2 \leq a$ so, $a \in S_a$.

In either case $S_a \neq \{0\}$

so $b \neq 0$ and $b > 0$.

We want to show $b^2 = a$, by dichotomy exactly 1 of the 3 holds:

(I) $b^2 = a$ (II) $b^2 > a$ (III) $b^2 < a$

We want to exclude (II) and (III).

Assume (III) holds, suppose $b^2 < a$

Set $\epsilon_0 := a - b^2 > 0$

set $\epsilon := \min(1, \frac{\epsilon_0}{2+4b}) \in P$ as $b \in P$

set $b' := b + \epsilon > 0$

Since $\epsilon > 0$ $b' > b$

Plan] Show $(b')^2 < a$

If we do this $b' \in A$. But, $b' > b$ and b is defined to be the supremum so, oops.

$$(b')^2 = (b^2 + 2b\varepsilon + \varepsilon^2)$$

$$b^2 = (a - \varepsilon_0) \quad (\text{by our definition of } \varepsilon_0)$$

$$(b')^2 = (a - \varepsilon_0) + 2b\varepsilon + \varepsilon^2$$

$$\text{So } (b')^2 \leq a \text{ iff } -\varepsilon_0 + 2b\varepsilon + \varepsilon^2 \leq 0, \text{ so}$$

lets show that this holds.

$$\underline{\varepsilon^2 + 2b\varepsilon} \leq \varepsilon_0$$

$\Leftrightarrow \varepsilon(2b + \varepsilon) \text{ as } \varepsilon \leq 1$ This is less than or equal

$$\text{to } \varepsilon(2b + 1) \leq \frac{\varepsilon_0(2b+1)}{2+4b} \text{ as } \varepsilon \leq \frac{\varepsilon_0}{2+4b}$$

$$= \frac{\varepsilon_0}{2} \text{ which is certainly less than } \varepsilon_0.$$

$\therefore (b')^2 < a$ And we get a contradiction
so III doesn't hold.

Assume III holds, suppose $b^2 > a$.

$$\text{set } \varepsilon_0 := b^2 - a > 0$$

$$\text{set } \varepsilon := \min(b, \frac{\varepsilon_0}{2+4b}) \in \mathbb{P} \text{ for the same reason as earlier } (b > 0).$$

$$b' := b - \varepsilon$$

As $\varepsilon \leq b$ we know

$$b' \geq 0 \text{ and } b' < b.$$

Plan] Show $a < (b')^2$. Then b' is an upper bound for S_a as $S_a := \{x \in F \mid x^2 \leq a\}$

$$\text{fix } y \in S_a \Rightarrow y^2 \leq a < b^2 \Rightarrow y^2 < b'^2 \Rightarrow |y| < |b'| \\ y \leq |y| < |b|$$

if b' is an upper bound we have a problem
as $b' < b$ and b is $\sup(S_a)$. oops.

Let's show $(b')^2 \geq a$

$$(b')^2 = (b - \varepsilon^2) = b^2 - 2b\varepsilon + \varepsilon^2 \\ = (\varepsilon_0 + a) - 2b\varepsilon + \varepsilon^2$$

if $a < (b')^2$ then, $\varepsilon_0 - 2b\varepsilon + \varepsilon^2 > 0$,

$$\varepsilon_0 - 2b\varepsilon + \varepsilon^2 = \varepsilon_0 + \varepsilon(\varepsilon - 2b) \quad \text{as } -1 < \varepsilon \\ > \varepsilon_0 + \varepsilon(-1 - 2b) \\ > \varepsilon_0 - \varepsilon(1 + 2b) \\ \geq \varepsilon_0 - \frac{\varepsilon_0}{2+4b} (1+2b) \\ = \frac{\varepsilon_0}{2}$$

So, garbage $\succ \frac{\varepsilon_0}{2} \succ 0 \therefore$ garbage e.p.

$\therefore a < (b')^2 \rightarrow$ contradiction

So, (ii) doesn't hold. By trichotomy,

$$b^2 = a$$

□

What did we just prove?

In an ordered field, we now know exactly when a positive element has a square root.

It is when $\sup(S_a)$ exists in F .

$$S_a := \{x \in F \mid x^2 \leq a\}$$

Corollary due to completeness all positive elements of \mathbb{R} have square roots in \mathbb{R} .

Weak induction

Def] The $S \subseteq \mathbb{N}$ is said to be weakly inductive if

- ① $1 \in S$
- ② $k \in S \rightarrow k+1 \in S$

Thm] Every weakly inductive $S \subseteq \mathbb{N}$ contains $\mathbb{N} = \mathbb{N}$.

Pf] Clear since every weakly inductive set is inductive and \mathbb{N} is the smallest inductive set so, $\mathbb{N} \subseteq S$. By 2 way containment $S = \mathbb{N}$.

Real Important Lemma

Lemma] $\forall n \in \mathbb{N}, n \geq 1$.

Ie: 1 is the minimal element of \mathbb{N} .

Pf] Define $S := \{m \in \mathbb{N} \mid m \geq 1\}$

Show S is weakly inductive

- $1 \in S$? yes
- fix $x \in S$
- $\Rightarrow x \geq 1$
- $\Rightarrow x+1 \geq 2 > 1 \Rightarrow x+1 > 1 \Rightarrow x+1 \geq 1$
- so $x+1 \in S$

$\therefore S$ is weakly inductive and the property holds for all $n \in \mathbb{N}$.

BIG PROOF
Incoming



Lemma) If $n, m \in \mathbb{N}$ we have

$$m > n \iff m \geq n+1$$

Corollary) $\forall n, m \in \mathbb{N}$, we have

$$n < m \leq n+1 \iff m = n+1$$

Immediate from lemma. \Leftarrow is free

$= \subset$

$$n < m \leq n+1$$

(\supset by lemma) $m \geq n+1$

$$\begin{cases} m \geq n+1 \\ m \leq n+1 \end{cases}$$

$$\begin{cases} m = n+1 \end{cases}$$

Pf) fix $m, n \in \mathbb{N}$

note: \leq is immediate

$$m \geq n+1 > n$$

\implies

Note: Things get messy and we will do a double induction

fix $m, n \in \mathbb{N}$, assume $m > n$.

Def) given $k \in \mathbb{N}$

$$T_k := \{l \in \mathbb{N} \mid l \leq k \text{ or } l = k' + 1 \text{ for some } k' \in \mathbb{N} \text{ w/ } k' \geq k\}$$

Consider T_n . If we show $m \in T_n$, we're done.

because then we know either $m = n+1$ for $n' \geq n$
 $n' \in \mathbb{N}$ or $m < n$ \rightarrow false.

\hookrightarrow if true $m = n'+1 \geq n+1$ as $n' \geq n$.

We want $m \in T_n$. It will be enough to show
 $T_n = \mathbb{N}$ as $m \in \mathbb{N}$.

Show $T_n = \mathbb{N}$. Here it goes ...

Consider $\bar{T} := \{s \in \mathbb{N} \mid T_s = \mathbb{N}\}$

If we show $\bar{T} = \mathbb{N}$ we win

Claim: \bar{T} is weakly inductive

- $1 \in \bar{T}$? Let's check. We must show T_1 is weakly inductive

$$T_1 := \{l \in \mathbb{N} \mid l \leq 1 \text{ or } l = n' + 1 \text{ for } n' \in \mathbb{N}, n' \geq 1\}$$

$\rightarrow 1 \in T_1$? yes

$\rightarrow x \in T_1 \rightarrow x+1 \in T_1$.

If $x \in T_1$, we know $x \leq 1$ or

$$\underline{x > 1.}$$

$x > 1$ and $n \in T_1$,

$$\exists n' \in \mathbb{N} \quad x = n' + 1 \quad n' \geq 1$$

↳ by our lemma

$x = 1$ so $x+1 = 2$
which is in T_1

Then $x+1 = n'+1 + 1$

So is $(x+1) \in T_1$? Yes as $n'+1 > n \geq 1$

$\Rightarrow T_1$ is weakly inductive

- Suppose $k \in \bar{T}$ show $k+1 \in \bar{T}$

Since $k \in \bar{T}$ we know $T_k = \mathbb{N}$

$$T_k := \{l \in \mathbb{N} \mid l \leq k \text{ or } l = n' + 1 \text{ for } n' \in \mathbb{N}, n' \geq k\}$$

We need to get $T_{k+1} = \mathbb{N}$

We need to get $T_{k+1} = \mathbb{N}$ by 2 way containment.

By definition $T_{k+1} \subseteq \mathbb{N}$. So lets show $\mathbb{N} \subseteq T_{k+1}$

Let $r \in \mathbb{N}$. Show $r \in T_{k+1}$

Since $r \in \mathbb{N}$ either $r \leq k+1$ or $r > k+1$

If $r \leq k+1$ $r \in T_{k+1}$

If $r > k+1$ we need to show $\exists n'' \in \mathbb{N}$ with

$$n'' \geq k+1 \text{ s.t. } r = n'' + 1$$

Since $T_1 = \mathbb{N}$ and $r \in \mathbb{N}$ we know $r \in T_1$, this means that $r=1$ or $\exists n' \in \mathbb{N}$ s.t. $n' \geq 1$

$\rightarrow r = n'+1$ since $r > k+1, r \neq 0$.

Using this, we proceed.

Because $k \in T$, we know $n' \in T_k$ so,

$$n' \leq k \quad \text{or} \quad n' = n''' + 1 \text{ where } n''' \geq k$$

$$\begin{aligned} r &> k+1 \text{ by def} \\ r &= n' + 1 \end{aligned}$$

$$n' + 1 > k+1$$

$$\Rightarrow n' > k$$

so this is true

Putting everything together

$$r = n' + 1, n' \geq 1$$

$$n' = n''' + 1, n''' \geq k$$

$$r = (n''' + 1) + 1 \text{ where } n''' + 1 \geq k+1$$

$$\text{So } r \in T_{k+1}$$

$$\therefore T_{k+1} = \mathbb{N}$$

So T is inductive and $T = \mathbb{N}$

$$\text{So, } \forall n, m \in \mathbb{N} \quad m > n \iff m \geq n+1$$

Defn] let (F, \leq) be an ordered field the subset $V \subseteq F$ is said to be well ordered P.t every nonempty $V \subseteq V$ has a minimal element.

Thm] \mathbb{N} is well ordered as a subset of \mathbb{R} .

Pf] Let $V \subseteq \mathbb{N}$ be non-empty

Task : Show V has a minimal element.

Consider : $L_V := \{k \in \mathbb{N} \mid k \text{ is a strict lower bound for } V\}$
↳ lower bound and $k \notin V$

What can we say about L_V ? Focus on 1.

① $1 \notin L_V$, this means $1 \in V$. Why, because if $1 \notin V$, then $1 < v \forall v \in V$. In this case 1 is a strict lower bound for V so $1 \in L_V$.

We are happy because in this case 1 is a min elt.

② $1 \in L_V$,

Since $V \neq \emptyset$, we know $L_V \neq \mathbb{N}$ because $L_V \cap V = \emptyset$. This means L_V contains 1 and $L_V \neq \mathbb{N}$.

As $L_V \neq \mathbb{N}$ L_V isn't inductive, so a criteria of inductivity must fail. As $1 \in L_V$, the $x \in L_V \rightarrow x+1 \in L_V$ criteria fail.

So, $\exists k \in L_V$ s.t $k+1 \notin L_V$.

Thus $k < n \wedge n \in V$ as $k \in L_V$

• $\exists n_0 \in V$ s.t $k+1 \geq n_0 \Rightarrow k+1 \notin L_V$

Putting this together we have.

Since $n_0 \in V$ and $k \in L_V$, $k < n_0 \leq k+1 \Rightarrow$ by lemma
 $n_0 = k+1$

Claim $n_0 = \min(V)$

true since $\forall n \in V$, $n < k \Rightarrow n \leq k+1$ (by lemma)
so, $n \leq n_0$

So n_0 is a lower bound for V and $n_0 \in V$. $\therefore n_0 = \min(V)$

Defn] $S \subseteq \mathbb{N}$ is strongly inductive P.f $\forall n \in \mathbb{N}$ if,

$$\{k \in \mathbb{N} \mid k < n\} \subseteq S$$

then $n \in S$

Remark: if $S \subseteq \mathbb{N}$ has this property we automatically get $1 \in S$ as $\emptyset \subseteq S$.

Thm] If $S \subseteq \mathbb{N}$ is strongly inductive $S = \mathbb{N}$.

Pf] Suppose $S \subseteq \mathbb{N}$ is strongly inductive. If $S \neq \mathbb{N}$, then $V = \mathbb{N} \setminus S \neq \emptyset$. Let $v = \min(V)$ (as \mathbb{N} is well ordered). If $n \in \mathbb{N}$, $n < v$, then $n \notin V$, and as $v = s'$ in context of \mathbb{N} , $n \in S$. Thus we have $\forall n < v \ n \in S \Rightarrow v \in S$ because S is strongly inductive. \leftarrow definition.
But this is bad because $v \in V$ and $S \cap V = \emptyset$ by definition. oops.

So $S = \mathbb{N}$.

Recursion: Spivak

Big Deal : Division alg.

$\forall D, d \in \mathbb{N} \ \exists! q, r \in \mathbb{N} \cup \{0\}$ with $r < d$ s.t
 $D = q \cdot d + r$

Pf] Since there are 2 assertions, prove uniqueness first.

fix $d, D \in \mathbb{N}$

suppose $D = q' \cdot d + r' = q \cdot d + r$ for $q, q', r, r' \in \mathbb{N} \cup \{0\}$
 $r, r' < d$.

wlog, suppose $r \geq r' \rightarrow r - r' \geq p$

$$\text{Then, } 0 = \delta(q' - q) + (r' - r)$$

$$0 \leq r - r' = \delta(q' - q)$$

$$r - r' < \delta \rightarrow \text{by definition as } r, r' \leq \delta$$

$$0 \leq \delta(q' - q) < \delta$$

$$0 \leq q' - q < 1$$

by our lemma,

$$q' - q = 0 \Rightarrow q' = q$$

using this its easy to show $r' = r$

Pf of Existence]

Fix $\delta \in \mathbb{N} \rightarrow$ ignore fixing D for now.

Consider $S = \{D \in \mathbb{N} \mid D = q\delta + r \text{ for } q, r \in \mathbb{N} \cup \{0\}$
 $r < \delta\}$

\hookrightarrow the set of all numbers that can be written as a linear combination of δ or $q\delta + r$ (intuitively this is all of \mathbb{N})

Show S is strongly inductive. $\therefore S = \mathbb{N} \therefore D \in S$.

Fix $D_0 \in \mathbb{N}$, so that $D \in S$ whenever $D < D_0$.

Show $D_0 \in S$. There are 3 possibilities,

I) $D_0 = \delta$

take $q=1, r=0, D_0 = q\cdot\delta + r = \delta \text{ so } D_0 \in S$

II) $D_0 < \delta$

take $q=0, r=D_0, D_0 = \delta, \text{ so } D_0 \in S$

Z $D_0 > 0$

so, $D_0 - \delta \in \mathbb{N}$

and $D_0 - \delta < D_0$ as δ is non-zero

we have by assumption,

$D_0 - \delta \in S$ \Rightarrow strong induction

$\exists q, r \in \mathbb{N} \cup \{0\}$ with $r < \delta$ s.t

$D_0 - \delta = q\delta + r$ thus,

$D_0 = (q+1)\delta + r$ so $D_0 \in S$

In all 3 cases $D_0 \in S$. So, S is strongly inductive. So, we are done as $D \in S$.

Let's talk about integers !!

$\mathbb{Z} = \{x \in \mathbb{R} \mid \forall n \in \mathbb{N} \cup \{0\}\}$

Define | let $a, b \in \mathbb{N}$. The element a is called a divisor of b pt $\exists m \in \mathbb{N}$ s.t

$$b = a \cdot m$$

in this case $a \mid b \Rightarrow a$ divides b

Define | let $a, b \in \mathbb{N}$. The gcd (a, b) is their greatest common divisor.

$S = \{d \in \mathbb{N} \mid d \mid a \text{ and } d \mid b\}$

S is nonempty as $1 \mid a$ and $1 \mid b$.

Bounded above by $\min(a, b)$.

Does it have a maximal element?

This is a finite set (see HW3).

By HW3, this set has a maximal element.

Bezout lemma

Suppose $a, b \in \mathbb{N}$. Then $\exists m, n \in \mathbb{Z}$ st
 $a \cdot m + b \cdot n = \gcd(a, b)$

Caution: no claim about uniqueness of a and b .

Pf]

fix $a, b \in \mathbb{N}$

let $S = \{am + bn \mid m, n \in \mathbb{Z}, am + bn \in \mathbb{N}\}$

note: (1) $S \neq \emptyset$

$$a = a \cdot 1 + b \cdot 0 \in S$$

(2) $S \subseteq \mathbb{N}$

(1) + (2) imply S has a least element \rightarrow well ordered

lets call this element s' .

so, $s' = a \cdot m' + b \cdot n'$ for some $m', n' \in \mathbb{Z}$

We will show $s' \mid a$ and $s' \mid b$.

But we also need to show that it is the greatest common divisor.

(A) Show $s' \mid a$

by division alg,

$\exists q, r \in \mathbb{N} \cup \{0\}$ with $r < s'$

$$\text{s.t. } a = q \cdot s' + r$$

$$r = a - q \cdot s'$$

$$= a - (q(am' + bn'))$$

$$r = (1 - qm') \cdot a - (qn')b$$

Is this in the set S ? Its of the right form but it is only in S iff it is positive. If so, it implies $r \in S$. But since $r < s'$ (which is $\min(S)$), oops.

Thus, the expression isn't positive. Since $r \in \mathbb{N}$, it can't be negative. \therefore by trichotomy it is 0 and $r=0$.

$$\text{So, } a = q \cdot s'$$

$$s' \mid a$$

(B) Similar argument for B $s' \mid b$

(A+B) \rightarrow s' is a common divisor of a and b .

We are done if $s' = \gcd(a, b)$ as it shows for some element in the set because it satisfied.

Last part Let d be a common factor of a and b . we must show $d \geq s'$.

Since $d \mid a$ and $d \mid b$,

$$d \mid (a \cdot m + b \cdot n)$$

$$\text{so } d \mid s'$$

$$\Rightarrow \exists q \in \mathbb{N} \text{ st}$$

$$s' = q \cdot d, \quad q \geq 1$$

$$dq \geq d$$

$$s' \geq d$$

$\therefore s'$ is the $\gcd(a, b)$ and we win