Contents

Contents			<pre>struct Edge { int u, v; lang lang con flow;</pre>
1	Combinatorial optimization1.1 Dinic's	1 1 2 3	<pre>long long cap, flow; Edge() {} Edge(int u, int v, long long cap): u(u), v(v), cap(</pre>
2	Geometry 2.1 Convex hull	4 4 4 6	<pre>vector<vector<int>> g; vector<int> d, pt; Dinic(int N): N(N), E(0), g(N), d(N), pt(N) {} void AddEdge(int u, int v, long long cap) {</int></vector<int></pre>
	Numerical algorithms 3.1 Number theory (modular, Chinese remainder, linear Diophantine)	7 8 8 9 10 10	<pre>if (u != v) { E.emplace_back(Edge(u, v, cap)); g[u].emplace_back(E.size() - 1); E.emplace_back(Edge(v, u, 0)); g[v].emplace_back(E.size() - 1); } bool BFS(int S, int T) { queue<int> q({S}); fill(d.begin(), d.end(), N + 1); d[S] = 0; while(!q.empty()) {</int></pre>
	Graph algorithms 4.1 Bellman-Ford shortest paths with negative edge weights . 4.2 Eulerian path	13 13 13 14 14 15	<pre>int u = q.front(); q.pop(); if (u == T) break; for (int k: g[u]) { Edge &e = E[k]; if (e.flow < e.cap && d[e.v] > d[e.u] + 1) { d[e.v] = d[e.u] + 1; q.emplace(e.v); }</pre>
	Data structures5.1Suffix array5.2KD-tree5.3Merge Sort Tree	16 16 17 19	<pre> } return d[T] != N + 1; } long long DFS(int u, int T, long long flow = -1) { if (u == T flow == 0) return flow; } </pre>
	Miscellaneous 6.1 Miller-Rabin Primality Test 6.2 Pollard-Rho factorization 6.3 Manachers algorithm 6.4 Convex Hull Trick 6.5 Dynamic Programming(DnC) 6.6 Longest increasing subsequence 6.7 Dates 6.8 Knuth-Morris-Pratt 6.9 2-SAT	19 19 19 20 20 21 21 21 22 22	<pre>for (int &i = pt[u]; i < g[u].size(); ++i) { Edge &e = E[g[u][i]]; Edge &oe = E[g[u][i]^1]; if (d[e.v] == d[e.u] + 1) { long long amt = e.cap - e.flow; if (flow != -1 && amt > flow) amt = flow; if (long long pushed = DFS(e.v, T, amt)) { e.flow += pushed; oe.flow -= pushed; return pushed; } } }</pre>
1	Combinatorial optimization		return 0; }
1.1	• • • • • • • • • • • • • • • • • • •		<pre>long long MaxFlow(int S, int T) { long long total = 0; while (BFS(S, T)) { fill(pt.begin(), pt.end(), 0); }</pre>

```
while (long long flow = DFS(S, T))
        total += flow;
    return total;
};
```

1.2 Min-cost max-flow

```
// Implementation of min cost max flow algorithm using
   adjacency
// matrix (Edmonds and Karp 1972). This implementation
   keeps track of
// forward and reverse edges separately (so you can set
   cap[i][j] !=
// cap[j][i]). For a regular max flow, set all edge
   costs to 0.
// Running time, O(|V|^2) cost per augmentation
                           O(|V|^3) augmentations
       max flow:
       min cost max flow: O(|V|^4 * MAX EDGE COST)
   augmentations
// INPUT:
       - graph, constructed using AddEdge()
       - source
       - sink
// OUTPUT:
      - (maximum flow value, minimum cost value)
       - To obtain the actual flow, look at positive
   values only.
#include <cmath>
#include <vector>
#include <iostream>
using namespace std;
typedef vector<int> VI;
typedef vector<VI> VVI;
typedef long long L;
typedef vector<L> VL;
typedef vector<VL> VVL;
typedef pair<int, int> PII;
typedef vector<PII> VPII;
const L INF = numeric_limits<L>::max() / 4;
struct MinCostMaxFlow {
  int N;
  VVL cap, flow, cost;
  VI found;
  VL dist, pi, width;
  VPII dad;
  MinCostMaxFlow(int N) :
    N(N), cap(N, VL(N)), flow(N, VL(N)), cost(N, VL(N)),
    found(N), dist(N), pi(N), width(N), dad(N) {}
  void AddEdge(int from, int to, L cap, L cost) {
   this->cap[from][to] = cap;
    this->cost[from][to] = cost;
```

```
void Relax(int s, int k, L cap, L cost, int dir) {
    L val = dist[s] + pi[s] - pi[k] + cost;
    if (cap && val < dist[k]) {
      dist[k] = val;
      dad[k] = make_pair(s, dir);
      width [k] = min(cap, width [s]);
 L Dijkstra(int s, int t) {
    fill(found.begin(), found.end(), false);
    fill(dist.begin(), dist.end(), INF);
    fill(width.begin(), width.end(), 0);
    dist[s] = 0;
    width[s] = INF;
    while (s != -1) {
      int best = -1;
      found[s] = true;
      for (int k = 0; k < N; k++) {
        if (found[k]) continue;
        Relax(s, k, cap[s][k] - flow[s][k], cost[s][k],
           1);
        Relax(s, k, flow[k][s], -\cos t[k][s], -1);
        if (best == -1 \mid \mid dist[k] < dist[best]) best = k
      s = best;
    for (int k = 0; k < N; k++)
      pi[k] = min(pi[k] + dist[k], INF);
    return width[t];
  pair<L, L> GetMaxFlow(int s, int t) {
    L \text{ totflow} = 0, totcost = 0;
    while (L amt = Dijkstra(s, t)) {
      totflow += amt;
      for (int x = t; x != s; x = dad[x].first) {
        if (dad[x].second == 1) {
          flow[dad[x].first][x] += amt;
          totcost += amt * cost[dad[x].first][x];
        } else {
          flow[x][dad[x].first] -= amt;
          totcost -= amt * cost[x][dad[x].first];
    return make_pair(totflow, totcost);
} ;
// BEGIN CUT
// The following code solves UVA problem #10594: Data
   Flow
int main() {
  int N, M;
 while (scanf("%d%d", &N, &M) == 2) {
    VVL v(M, VL(3));
    for (int i = 0; i < M; i++)
```

```
scanf("%Ld%Ld%Ld", &v[i][0], &v[i][1], &v[i][2]);
   L D. K:
   scanf("%Ld%Ld", &D, &K);
   MinCostMaxFlow mcmf(N+1);
   for (int i = 0; i < M; i++) {
      mcmf.AddEdge(int(v[i][0]), int(v[i][1]), K, v[i
     mcmf.AddEdge(int(v[i][1]), int(v[i][0]), K, v[i
         ][2]);
   mcmf.AddEdge(0, 1, D, 0);
   pair<L, L> res = mcmf.GetMaxFlow(0, N);
   if (res.first == D) {
      printf("%Ld\n", res.second);
    } else
     printf("Impossible.\n");
  return 0;
// END CUT
```

1.3 Edmonds Max Matching

```
/*
Input:
V->number of vertices
E->number of edges
pair of vertices as edges (vertices are 1..V)
Output:
output of edmonds() is the maximum matching
match[i] is matched pair of i (-1 if there isn't a
   matched pair)
#include <bits/stdc++.h>
using namespace std;
const int M=505;
struct struct_edge{int v;struct_edge* n;};
typedef struct edge* edge;
struct_edge pool[M*M*2];
edge top=pool,adj[M];
int V,E,match[M],qh,qt,q[M],father[M],base[M];
bool inq[M], inb[M], ed[M][M];
void add edge(int u,int v)
  top->v=v,top->n=adj[u],adj[u]=top++;
  top->v=u, top->n=adj[v], adj[v]=top++;
int LCA(int root,int u,int v)
  static bool inp[M];
  memset(inp, 0, sizeof(inp));
  while (1)
      inp[u=base[u]]=true;
      if (u==root) break;
      u=father[match[u]];
```

```
while (1)
      if (inp[v=base[v]]) return v;
      else v=father[match[v]];
void mark blossom(int lca,int u)
  while (base[u]!=lca)
      int v=match[u];
      inb[base[u]]=inb[base[v]]=true;
      u=father[v];
      if (base[u]!=lca) father[u]=v;
void blossom contraction(int s,int u,int v)
  int lca=LCA(s,u,v);
  memset(inb,0,sizeof(inb));
  mark_blossom(lca,u);
  mark_blossom(lca, v);
  if (base[u]!=lca)
    father[u]=v;
  if (base[v]!=lca)
    father[v]=u;
  for (int u=0; u<V; u++)
    if (inb[base[u]])
  base[u]=lca;
  if (!inq[u])
    inq[q[++qt]=u]=true;
int find_augmenting_path(int s)
  memset(inq,0,sizeof(inq));
  memset (father, -1, sizeof (father));
  for (int i=0;i<V;i++) base[i]=i;
  inq[q[qh=qt=0]=s]=true;
  while (qh<=qt)</pre>
      int u=q[qh++];
      for (edge e=adj[u];e;e=e->n)
    int v=e->v;
    if (base[u]!=base[v]&&match[u]!=v)
      if ((v==s)||(match[v]!=-1 && father[match[v]]!=-1)
        blossom contraction(s,u,v);
      else if (father[v]==-1)
    father[v]=u;
    if (match[v] == -1)
      return v;
    else if (!ing[match[v]])
      inq[q[++qt]=match[v]]=true;
```

```
return -1;
int augment_path(int s,int t)
  int u=t, v, w;
  while (u!=-1)
      v=father[u];
      w=match[v];
      match[v]=u;
      match[u]=v;
      u=w;
  return t!=-1;
int edmonds()
  int matchc=0;
  memset (match, -1, sizeof (match));
  for (int u=0; u<V; u++)
    if (match[u] == -1)
      matchc+=augment_path(u, find_augmenting_path(u));
  return matchc;
int main()
  int u, v;
  cin>>V>>E;
  while (E^{--})
      cin>>u>>v;
      if (!ed[u-1][v-1])
    add edge (u-1, v-1);
    ed[u-1][v-1]=ed[v-1][u-1]=true;
  cout << edmonds() << endl;
  for (int i=0;i<V;i++)</pre>
    if (i<match[i])</pre>
      cout << i + 1 << " " << match [i] + 1 << endl;
```

2 Geometry

2.1 Convex hull

```
typedef pair<long long, long long> PT;
long double dist(PT a, PT b) {
   return sqrt(pow(a.first-b.first,2)+pow(a.second-b.second,2));
}
long long cross(PT o, PT a, PT b) {
   PT OA = {a.first-o.first,a.second-o.second};
   PT OB = {b.first-o.first,b.second-o.second};
   return OA.first*OB.second - OA.second*OB.first;
}
vector<PT> convexhull() {
   vector<PT> hull;
   sort(a,a+n,[](PT i, PT j) {
```

2.2 Miscellaneous geometry

```
double INF = 1e100, EPS = 1e-12;
struct PT {
  double x, y;
  PT() {}
  PT (double x, double y) : x(x), y(y) {}
  PT (const PT &p) : x(\bar{p}.x), y(p.y)
  PT operator + (const PT &p) const { return PT(x+p.x,
     y+p.y);
  PT operator - (const PT &p) const { return PT(x-p.x,
     y-p.y); }
  PT operator * (double c)
                                const { return PT(x*c,
     V*C ); }
  PT operator / (double c)
                                const { return PT(x/c,
     y/c ); }
};
double dot(PT p, PT q)
                            { return p.x*q.x+p.y*q.y; }
double dist2(PT p, PT q)
                          { return dot(p-q,p-q); }
double cross(PT p, PT q) { return p.x*q.y-p.y*q.x; }
ostream & operator << (ostream & os, const PT & p) {
  os << "(" << p.x << "," << p.y << ")";
// rotate a point CCW or CW around the origin
PT RotateCCW90(PT p)
                      { return PT(-p.y,p.x); }
PT RotateCW90 (PT p)
                        { return PT(p.y,-p.x); }
PT RotateCCW(PT p, double t) {
  return PT(p.x*cos(t)-p.y*sin(t), p.x*sin(t)+p.y*cos(t)
// project point c onto line through a and b
// assuming a != b
PT ProjectPointLine(PT a, PT b, PT c) {
  return a + (b-a) *dot (c-a, b-a) /dot (b-a, b-a);
// project point c onto line segment through a and b
PT ProjectPointSegment (PT a, PT b, PT c) {
  double r = dot(b-a,b-a);
  if (fabs(r) < EPS) return a;</pre>
```

```
r = dot(c-a, b-a)/r;
  if (r < 0) return a;
  if (r > 1) return b;
  return a + (b-a) *r;
// compute distance from c to segment between a and b
double DistancePointSegment(PT a, PT b, PT c) {
  return sqrt(dist2(c, ProjectPointSegment(a, b, c)));
// compute distance between point (x,y,z) and plane ax+
   bv+cz=d
double DistancePointPlane(double x, double y, double z,
                          double a, double b, double c,
                              double d)
  return fabs (a*x+b*y+c*z-d) /sqrt (a*a+b*b+c*c);
// determine if lines from a to b and c to d are
   parallel or collinear
bool LinesParallel(PT a, PT b, PT c, PT d) {
  return fabs(cross(b-a, c-d)) < EPS;</pre>
bool LinesCollinear(PT a, PT b, PT c, PT d) {
  return LinesParallel(a, b, c, d)
      && fabs(cross(a-b, a-c)) < EPS
      && fabs(cross(c-d, c-a)) < EPS;
// determine if line segment from a to b intersects with
// line segment from c to d
bool SegmentsIntersect (PT a, PT b, PT c, PT d) {
  if (LinesCollinear(a, b, c, d)) {
    if (dist2(a, c) < EPS || dist2(a, d) < EPS ||
      dist2(b, c) < EPS || dist2(b, d) < EPS) return
         true;
    if (dot(c-a, c-b) > 0 && dot(d-a, d-b) > 0 && dot(c-
       b, d-b) > 0
      return false;
    return true;
  if (cross(d-a, b-a) * cross(c-a, b-a) > 0) return
  if (cross(a-c, d-c) * cross(b-c, d-c) > 0) return
     false:
  return true;
// compute intersection of line passing through a and b
// with line passing through c and d, assuming that
// intersection exists; for segment intersection, check
   i f
// segments intersect first
PT ComputeLineIntersection(PT a, PT b, PT c, PT d) {
  b=b-a; d=c-d; c=c-a;
  assert (dot (b, b) > EPS && dot (d, d) > EPS);
  return a + b*cross(c, d)/cross(b, d);
// compute center of circle given three points
```

```
PT ComputeCircleCenter(PT a, PT b, PT c) {
  b = (a+b)/2;
  c = (a+c)/2;
  return ComputeLineIntersection(b, b+RotateCW90(a-b), c
     , c+RotateCW90(a-c);
// determine if point is in a possibly non-convex
   polygon (by William
// Randolph Franklin); returns 1 for strictly interior
   points, 0 for
// strictly exterior points, and 0 or 1 for the
   remaining points.
// Note that it is possible to convert this into an \star
   exact* test using
// integer arithmetic by taking care of the division
   appropriately
// (making sure to deal with signs properly) and then by
    writing exact
// tests for checking point on polygon boundary
bool PointInPolygon(const vector<PT> &p, PT q) {
 bool c = 0;
  for (int i = 0; i < p.size(); i++) {</pre>
    int j = (i+1) %p.size();
    if ((p[i].y <= q.y && q.y < p[j].y ||</pre>
      p[j].y \le q.y \&\& q.y < p[i].y) \&\&
      q.x < p[i].x + (p[j].x - p[i].x) * (q.y - p[i].y)
         / (p[j].y - p[i].y))
  return c;
// determine if point is on the boundary of a polygon
bool PointOnPolygon(const vector<PT> &p, PT q) {
  for (int i = \bar{0}; i < p.size(); i++)
    if (dist2(ProjectPointSegment(p[i], p[(i+1)%p.size()
       ], q), q) < EPS)
      return true;
    return false:
// compute intersection of line through points a and b
   with
// circle centered at c with radius r > 0
vector<PT> CircleLineIntersection(PT a, PT b, PT c,
   double r) {
  vector<PT> ret;
  b = b-a;
  a = a-c;
  double A = dot(b, b);
  double B = dot(a, b);
  double C = dot(a, a) - r*r;
  double D = B*B - A*C;
  if (D < -EPS) return ret;</pre>
  ret.push back(c+a+b*(-B+sqrt(D+EPS))/A);
  if (D > EPS)
    ret.push_back(c+a+b*(-B-sqrt(D))/A);
  return ret;
// compute intersection of circle centered at a with
   radius r
// with circle centered at b with radius R
```

```
vector<PT> CircleCircleIntersection(PT a, PT b, double r
   , double R) {
  vector<PT> ret:
  double d = sqrt(dist2(a, b));
  if (d > r+R \mid | d+min(r, R) < max(r, R)) return ret;
  double x = (d*d-R*R+r*r)/(2*d);
  double y = sqrt(r*r-x*x);
  PT v = (b-a)/d;
  ret.push_back(a+v*x + RotateCCW90(v)*y);
  if (v > 0)
    ret.push_back(a+v*x - RotateCCW90(v)*y);
  return ret;
// This code computes the area or centroid of a (
   possibly nonconvex)
// polygon, assuming that the coordinates are listed in
   a clockwise or
// counterclockwise fashion. Note that the centroid is
   often known as
// the "center of gravity" or "center of mass".
double ComputeSignedArea(const vector<PT> &p) {
  double area = 0;
  for(int i = 0; i < p.size(); i++) {</pre>
    int j = (i+1) % p.size();
    area += p[i].x*p[j].y - p[j].x*p[i].y;
  return area / 2.0;
double ComputeArea(const vector<PT> &p) {
  return fabs(ComputeSignedArea(p));
PT ComputeCentroid(const vector<PT> &p) {
  PT c(0,0);
  double scale = 6.0 * ComputeSignedArea(p);
  for (int i = 0; i < p.size(); i++) {
    int j = (i+1) % p.size();
    c = c + (p[i]+p[j])*(p[i].x*p[j].y - p[j].x*p[i].y);
  return c / scale;
// tests whether or not a given polygon (in CW or CCW
   order) is simple
bool IsSimple(const vector<PT> &p) {
  for (int i = 0; i < p.size(); i++) {</pre>
    for (int k = i+1; k < p.size(); k++) {</pre>
      int j = (i+1) % p.size();
      int 1 = (k+1) % p.size();
      if (i == 1 \mid | i == k) continue;
      if (SegmentsIntersect(p[i], p[j], p[k], p[l]))
        return false;
    }
  return true;
```

2.3 3D geometry

```
public class Geom3D {
```

```
// distance from point (x, y, z) to plane aX + bY + cZ
public static double ptPlaneDist(double x, double y,
   double z,
    double a, double b, double c, double d) {
  return Math.abs(a*x + b*y + c*z + d) / Math.sqrt(a*a
      + b*b + c*c);
// distance between parallel planes aX + bY + cZ + d1
   = 0 and
// aX + bY + cZ + d2 = 0
public static double planePlaneDist(double a, double b
   , double c,
    double d1, double d2) {
  return Math.abs(d1 - d2) / Math.sqrt(a*a + b*b + c*c
     );
// distance from point (px, py, pz) to line (x1, y1,
   z1) - (x2, y2, z2)
// (or ray, or segment; in the case of the ray, the
   endpoint is the
// first point)
public static final int LINE = 0;
public static final int SEGMENT = 1;
public static final int RAY = 2;
public static double ptLineDistSq(double x1, double y1
   , double z1,
    double x2, double y2, double z2, double px, double
        py, double pz,
    int type) {
  double pd2 = (x1-x2)*(x1-x2) + (y1-y2)*(y1-y2) + (z1
     -z2)*(z1-z2);
  double x, y, z;
  if (pd2 == 0) {
    x = x1;
    v = y1;
    z = z1:
  } else {
    double u = ((px-x1)*(x2-x1) + (py-y1)*(y2-y1) + (
       pz-z1)*(z2-z1)) / pd2;
    x = x1 + u * (x2 - x1);
    y = y1 + u * (y2 - y1);
    z = z1 + u * (z2 - z1);
    if (type != LINE && u < 0) {
     x = x1;
      y = y1;
      z = z1;
    if (type == SEGMENT && u > 1.0) {
     x = x2;
     y = y2;
      z = z2;
  return (x-px)*(x-px) + (y-py)*(y-py) + (z-pz)*(z-pz)
public static double ptLineDist(double x1, double y1,
   double z1,
```

```
double x2, double y2, double z2, double px, double
    py, double pz,
    int type) {
    return Math.sqrt(ptLineDistSq(x1, y1, z1, x2, y2, z2
        , px, py, pz, type));
    }
}
```

3 Numerical algorithms

3.1 Number theory (modular, Chinese remainder, linear Diophantine)

```
// All algorithms described here work on nonnegative
   integers.
// return a % b (positive value)
int mod(int a, int b) {
  return ((a%b) + b) % b;
// computes lcm(a,b)
int lcm(int a, int b) {
  return a / __gcd(a, b) *b;
// (a^b) mod m via successive squaring
int powermod(int a, int b, int m) {
  return b?powermod(a*a%m,b/2,m)*(b%2?a:1)%m:1;
// returns q = qcd(a, b); finds x, y such that d = ax + b
int extended_euclid(int a, int b, int &x, int &y) {
  int xx = y = 0;
  int yy = x = 1;
  while (b) {
    int q = a / b;
    int t = b; b = a%b; a = t;
    t = xx; xx = x - q*xx; x = t;
    t = yy; yy = y - q*yy; y = t;
  return a;
// finds all solutions to ax = b \pmod{n}
VI modular_linear_equation_solver(int a, int b, int n) {
  int x, y;
  VI ret;
  int g = extended_euclid(a, n, x, y);
  if (!(b%q))
   x = mod(x*(b / g), n);
    for (int i = 0; i < q; i++)
      ret.push_back(mod(x + i*(n / q), n));
  return ret;
// computes b such that ab = 1 \pmod{n}, returns -1 on
   failure
int mod inverse(int a, int n) {
  int x, y;
  int g = extended_euclid(a, n, x, y);
```

```
if (q > 1) return -1;
  return mod(x, n);
// Chinese remainder theorem (special case): find z such
// z % m1 = r1, z % m2 = r2. Here, z is unique modulo M
    = 1cm(m1, m2).
// Return (z, M). On failure, M = -1.
PII chinese_remainder_theorem(int m1, int r1, int m2,
  int s, t;
  int g = extended_euclid(m1, m2, s, t);
  if (r1%g != r2%g) return make_pair(0, -1);
  return make_pair(mod(s*r2*m1 + t*r1*m2, m1*m2) / g, m1
     *m2 / q);
// Chinese remainder theorem: find z such that
//z % m[i] = r[i] for all i. Note that the solution is
// unique modulo M = lcm_i (m[i]). Return (z, M). On
// failure, M = -1. Note that we do not require the a[i
// to be relatively prime.
PII chinese remainder theorem (const VI &m, const VI &r)
  PII ret = make_pair(r[0], m[0]);
  for (int i = 1; i < m.size(); i++) {</pre>
    ret = chinese_remainder_theorem(ret.second, ret.
       first, m[i], r[i]);
    if (ret.second == -1) break;
  return ret;
// computes x and y such that ax + by = c
// returns whether the solution exists
bool linear_diophantine(int a, int b, int c, int &x, int
    &y) {
  if (!a && !b) {
    if (c) return false;
    x = 0; y = 0;
    return true;
  if (!a) {
    if (c % b) return false;
    x = 0; y = c / b;
    return true;
  if (!b) {
    if (c % a) return false;
    x = c / a; y = 0;
    return true;
  int g = \underline{gcd}(a, b);
  if (c % q) return false;
  x = c / g * mod_inverse(a / g, b / g);
  v = (c - a*x) / b;
  return true;
```

3.2 Systems of linear equations, matrix inverse, determinant

```
// Gauss-Jordan elimination with full pivoting.
// Uses:
   (1) solving systems of linear equations (AX=B)
     (2) inverting matrices (AX=I)
     (3) computing determinants of square matrices
// Running time: O(n^3)
   INPUT:
             a[][] = an nxn matrix
b[][] = an nxm matrix
             X = an nxm matrix (stored in b[][])
// OUTPUT:
             A^{-1} = an \ nxn \ matrix \ (stored in a[][])
             returns determinant of a[][]
#include <iostream>
#include <vector>
#include <cmath>
using namespace std;
const double EPS = 1e-10;
typedef vector<int> VI;
typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;
T GaussJordan (VVT &a, VVT &b) {
  const int n = a.size();
  const int m = b[0].size();
  VI irow(n), icol(n), ipiv(n);
  T \det = 1;
  for (int i = 0; i < n; i++) {
    int p_{1} = -1, p_{k} = -1;
    for (int j = 0; j < n; j++) if (!ipiv[j])</pre>
      for (int k = 0; k < n; k++) if (!ipiv[k])
        if (pj == -1 || fabs(a[j][k]) > fabs(a[pj][pk]))
            \{ pj = j; pk = k; \}
    if (fabs(a[pj][pk]) < EPS) { cerr << "Matrix is</pre>
        singular." << endl; exit(0); }</pre>
    ipiv[pk]++;
    swap(a[pj], a[pk]);
    swap(b[pj], b[pk]);
    if (pj != pk) det *= -1;
    irow[i] = pj;
    icol[i] = pk;
    T c = 1.0 / a[pk][pk];
    det *= a[pk][pk];
    a[pk][pk] = 1.0;
    for (int p = 0; p < n; p++) a[pk][p] *= c;
    for (int p = 0; p < m; p++) b[pk][p] *= c;
    for (int p = 0; p < n; p++) if (p != pk) {
      c = a[p][pk];
      a[p][pk] = 0;
      for (int q = 0; q < n; q++) a[p][q] -= a[pk][q] *
```

```
for (int q = 0; q < m; q++) b[p][q] -= b[pk][q] *
  for (int p = n-1; p >= 0; p--) if (irow[p] != icol[p])
    for (int k = 0; k < n; k++) swap(a[k][irow[p]], a[k
       ][icol[p]]);
  return det:
int main() {
  const int n = 4;
  const int m = 2;
  double A[n][n] = {
     \{1,2,3,4\},\{1,0,1,0\},\{5,3,2,4\},\{6,1,4,6\}\};
  double B[n][m] = \{\{1,2\},\{4,3\},\{5,6\},\{8,7\}\}\};
  VVT a(n), b(n);
  for (int i = 0; i < n; i++) {
   a[i] = VT(A[i], A[i] + n);
   b[i] = VT(B[i], B[i] + m);
  double det = GaussJordan(a, b);
  // expected: 60
  cout << "Determinant: " << det << endl;</pre>
  // expected: -0.233333 0.166667 0.133333 0.0666667
               0.166667 0.166667 0.333333 -0.333333
               0.233333 0.833333 -0.133333 -0.0666667
               0.05 - 0.75 - 0.1 0.2
  cout << "Inverse: " << endl;</pre>
  for (int i = 0; i < n; i++) {
    for (int j = 0; j < n; j++)
      cout << a[i][j] << '';
    cout << endl;
  // expected: 1.63333 1.3
               -0.166667 0.5
               2.36667 1.7
               -1.85 - 1.35
  cout << "Solution: " << endl;</pre>
  for (int i = 0; i < n; i++) {
    for (int j = 0; j < m; j++)
      cout << b[i][j] << ' ';
    cout << endl;
```

3.3 Reduced row echelon form, matrix rank

```
// Reduced row echelon form via Gauss-Jordan elimination
// with partial pivoting. This can be used for
    computing
// the rank of a matrix.
//
// Running time: O(n^3)
//
// INPUT: a[][] = an nxm matrix
```

rref[][] = an nxm matrix (stored in a[][]) returns rank of a[][] **#include** <iostream> #include <vector> #include <cmath> using namespace std; const double EPSILON = 1e-10; typedef double T; typedef vector<T> VT; typedef vector<VT> VVT; int rref(VVT &a) { int n = a.size();int m = a[0].size();int r = 0; for (int c = 0; c < m && r < n; c++) { int j = r; for (int i = r + 1; i < n; i++) **if** (fabs(a[i][c]) > fabs(a[j][c])) j = i; if (fabs(a[j][c]) < EPSILON) continue;</pre> swap(a[j], a[r]); T s = 1.0 / a[r][c];for (int j = 0; j < m; j++) a[r][j] *= s; for (int i = 0; i < n; i++) if (i != r) { T t = a[i][c];**for** (**int** j = 0; j < m; j++) a[i][j] -= t * a[r][j |; <u>ŕ</u>++; return r; int main() { const int n = 5, m = 4; double $A[n][m] = {$ {16, 2, 3, 13}, { 5, 11, 10, 8}, 9, 7, 6, 12}, 4, 14, 15, 1}, {13, 21, 21, 13}}; VVT a(n);for (int i = 0; i < n; i++) a[i] = VT(A[i], A[i] + m);int rank = rref(a); // expected: 3 cout << "Rank: " << rank << endl;</pre> // expected: 1 0 0 1 // 0 1 0 3 // $0 \ 0 \ 1 \ -3$

0 0 0 3.10862e-15

0 0 0 2.22045e-15

cout << "rref: " << endl;</pre>

cout << endl;

for (int i = 0; i < 5; i++) {

for (int j = 0; j < 4; j++)
 cout << a[i][j] << ' ';</pre>

3.4 Simplex algorithm

```
// Two-phase simplex algorithm for solving linear
   programs of the form
       maximize
       subject to Ax \le b
                    x >= 0
  INPUT: A -- an m x n matrix
        b -- an m-dimensional vector
          c -- an n-dimensional vector
         x -- a vector where the optimal solution will
   be stored
// OUTPUT: value of the optimal solution (infinity if
           above, nan if infeasible)
// To use this code, create an LPSolver object with A, b
   , and c as
// arguments. Then, call Solve(x).
#include <iostream>
#include <iomanip>
#include <vector>
#include <cmath>
#include <limits>
using namespace std;
typedef long double DOUBLE;
typedef vector<DOUBLE> VD;
typedef vector<VD> VVD;
typedef vector<int> VI;
const DOUBLE EPS = 1e-9;
struct LPSolver {
  int m, n;
  VI B, N;
  VVD D;
  LPSolver(const VVD &A, const VD &b, const VD &c) :
   m(b.size()), n(c.size()), N(n + 1), B(m), D(m + 2,
       VD(n + 2)
    for (int i = 0; i < m; i++) for (int j = 0; j < n; j
       ++) D[i][j] = A[i][j];
    for (int i = 0; i < m; i++) { B[i] = n + i; D[i][n]
       = -1; D[i][n + 1] = b[i]; 
    for (int j = 0; j < n; j++) { N[j] = j; D[m][j] = -c
       [ j ]; }
   N[n] = -1; D[m + 1][n] = 1;
  void Pivot(int r, int s) {
    double inv = 1.0 / D[r][s];
    for (int i = 0; i < m + 2; i++) if (i != r)
      for (int j = 0; j < n + 2; j++) if (j != s)
        D[i][j] = D[r][j] * D[i][s] * inv;
    for (int j = 0; j < n + 2; j++) if (j != s) D[r][j]
       *=inv;
```

```
for (int i = 0; i < m + 2; i++) if (i != r) D[i][s]
       \star = -inv;
    D[r][s] = inv;
    swap(B[r], N[s]);
  bool Simplex(int phase) {
    int x = phase = 1 ? m + 1 : m;
    while (true) {
      int s = -1;
      for (int j = 0; j \le n; j++) {
        if (phase == 2 \&\& N[j] == -1) continue;
        if (s == -1 || D[x][\dot{\eta}] < D[x][s] || D[x][\dot{\eta}] == D
            [x][s] \&\& N[j] < N[s]) s = j;
      if (D[x][s] > -EPS) return true;
      int r = -1;
      for (int i = 0; i < m; i++) {
        if (D[i][s] < EPS) continue;</pre>
        if (r == -1 || D[i][n + 1] / D[i][s] < D[r][n +
           1] / D[r][s] ||
          (D[i][n + 1] / D[i][s]) == (D[r][n + 1] / D[r]
             [s] && B[i] < B[r]) r = i;
      if (r == -1) return false;
      Pivot(r, s);
  DOUBLE Solve (VD &x) {
    int r = 0;
    for (int i = 1; i < m; i++) if (D[i][n + 1] < D[r][n
        + 1) r = i;
    if (D[r][n + 1] < -EPS) {
      Pivot(r, n);
      if (!Simplex(1) || D[m + 1][n + 1] < -EPS) return
         -numeric limits<DOUBLE>::infinity();
      for (int i = 0; i < m; i++) if (B[i] == -1) {
        int s = -1;
        for (int j = 0; j <= n; j++)
          if (s == -1 || D[i][j] < D[i][s] || D[i][j] ==
              D[i][s] \&\& N[j] < N[s]) s = j;
        Pivot(i, s);
    if (!Simplex(2)) return numeric_limits<DOUBLE>::
       infinity();
   x = VD(n);
    for (int i = 0; i < m; i++) if (B[i] < n) x[B[i]] =
       D[i][n + 1];
    return D[m][n + 1];
};
int main() {
  const int m = 4;
  const int n = 3;
  DOUBLE A[m][n] = {
    \{ 6, -1, 0 \},
      -1, -5, 0 },
    { 1, 5, 1 },
```

```
\{-1, -5, -1\}
DOUBLE _b[m] = { 10, -4, 5, -5 };
DOUBLE _{c[n]} = \{ 1, -1, 0 \};
VVD A(m);
VD b (\_b, \_b + m);
VD c(c, c+n);
for (int i = 0; i < m; i++) A[i] = VD(A[i], A[i] + n
LPSolver solver (A, b, c);
VD x;
DOUBLE value = solver.Solve(x):
cerr << "VALUE: " << value << endl; // VALUE: 1.29032</pre>
cerr << "SOLUTION:"; // SOLUTION: 1.74194 0.451613 1</pre>
for (size t i = 0; i < x.size(); i++) cerr << " " << x</pre>
   [i];
cerr << endl;
return 0;
```

3.5 Fast Fourier transform

```
auto FFT = [](vector<long double>a, vector<long double>b)
  auto DFT = [](vector<complex<long double>>&a, bool inv
     ) {
    int L=31-__builtin_clz(a.size()), n=1<<L;</pre>
    vector<complex<long double>> A(n);
    for (int k=0, r, i; k<n; A[r] =a[k++])</pre>
      for (i=r=0; i<L; (r<<=1) | = (k>>i++) & 1);
    complex<long double> w,wm,t;
    for(int m=2, 1, k; m<=n; m<<=1)
      for (w=\{0, 2*acos(-1)/m\}, wm=exp(inv?-w:w), k=0; k<n; k
        for (j=0, w=1; j < m/2; ++j, w*=wm)
           t=w*A[k+j+m/2], A[k+j+m/2]=A[k+j]-t, A[k+j]+=t;
    return A:
  };
  int n=4<<31-__builtin_clz(max(a.size(),b.size()));</pre>
  vector<complex<long double>> A(n), B(n), CC(n);
  for (int i=0; i< n; ++i)
    A[i]=i < a.size()?a[i]:0, B[i]=i < b.size()?b[i]:0;
  vector<complex<long double>> AA=DFT(A,0), BB=DFT(B,0);
  for (int i=0;i<n;++i) CC[i]=AA[i]*BB[i];</pre>
  vector<long double> c;
  for(auto i:DFT(CC,1)) if(c.size() < a.size() + b.size() - 1)</pre>
    c.push_back(i.real()/n+1e-5);
  return c;
};
```

3.6 BigInt library

```
struct bigint {
  const int base = 1000000000, base_digits = 9;
  vector<int> a;
  int sign;
  bigint(): sign(1) {}
  bigint(long long v) {
    *this = v;
}
```

```
"(cur), "c"(base));
bigint (const string &s) {
 read(s);
                                                                trim();
void operator=(const bigint &v) {
 sign = v.sign;
a = v.a;
                                                              bigint operator*(int v) const {
                                                                bigint res = *this;
                                                                reš *= v;
                                                                return res;
void operator=(long long v) {
  sign = 1;
                                                              friend pair < bigint, bigint > divmod (const bigint &al,
  if (v < 0) sign = -1, v = -v;
 for (; v > 0; v = v / base)
                                                                  const bigint &b1) {
                                                                int norm = a1.base / (b1.a.back() + 1);
    a.push back(v % base);
                                                                bigint a = a1.abs() * norm;
                                                                bigint b = b1.abs() * norm;
bigint operator+(const bigint &v) const {
 if (sign == v.sign) {
                                                                bigint q, r;
   bigint res = v_i
                                                                q.a.resize(a.a.size());
                                                                for (int i = a.a.size() - 1; i >= 0; i--) {
    for (int i = 0, carry = 0; i < (int) \max(a.size(),
                                                                  r *= a1.base;
        v.a.size()) || carry; ++i) {
                                                                  r += a.a[i];
      if (i == (int) res.a.size())
                                                                  int s1 = r.a.size() \le b.a.size() ? 0 : r.a[b.a.
        res.a.push_back(0);
      res.a[i] += carry + (i < (int) a.size() ? a[i] :
                                                                  int s2 = r.a.size() \le b.a.size() - 1 ? 0 : r.a[b.
                                                                      a.size() - 1];
      carry = res.a[i] >= base;
                                                                  int d = ((long long) al.base * s1 + s2) / b.a.back
      if (carry)
                                                                     ();
        res.a[i] -= base;
                                                                  r -= b * d;
                                                                  while (r < 0)
    return res;
                                                                   r += b, --d;
                                                                  q.a[i] = d;
  return *this - (-v);
                                                                q.sign = al.sign * bl.sign;
bigint operator-(const bigint &v) const {
                                                                r.sign = al.sign;
  if (sign == v.sign) {
                                                                q.trim();
    if (abs() >= v.abs()) {
                                                                r.trim();
      bigint res = *this;
                                                                return make_pair(q, r / norm);
      for (int i = 0, carry = 0; i < (int) v.a.size()</pre>
                                                              bigint operator/(const bigint &v) const {
         || carry; ++i) {
                                                                return divmod(*this, v).first;
        res.a[i] \rightarrow carry + (i < (int) v.a.size() ? v.
           a[i] : 0);
                                                              bigint operator% (const bigint &v) const {
        carry = res.a[i] < 0;
                                                                return divmod(*this, v).second;
        if (carry)
         res.a[i] += base;
                                                              void operator/=(int v) {
                                                                if (v < 0) sign = -sign, v = -v;
      res.trim();
      return res;
                                                                for (int i = (int) a.size() - 1, rem = 0; i >= 0; --
    return - (v - *this);
                                                                  long long cur = a[i] + rem * (long long) base;
                                                                  a[i] = (int) (cur / v);
                                                                  rem = (int) (cur % v);
  return *this + (-v);
                                                                trim();
void operator*=(int v) {
  if (v < 0)
                                                              bigint operator/(int v) const {
    sign = -sign, v = -v;
                                                                bigint res = *this;
  for (int i = 0, carry = 0; i < (int) a.size() ||</pre>
                                                                res /= v;
     carry; ++i) {
                                                                return res;
    if (i == (int) a.size())
      a.push_back(0);
                                                              int operator%(int v) const {
    long long cur = a[i] * (long long) v + carry;
                                                                if (v < 0)
    carry = (int) (cur / base);
    a[i] = (int) (cur % base);
                                                                int m = 0;
    //asm("divl %%ecx" : "=a"(carry), "=d"(a[i]) : "A
                                                                for (int i = a.size() - 1; i >= 0; --i)
```

```
m = (a[i] + m * (long long) base) % v;
  return m * sign;
void operator+=(const bigint &v) {
  *this = *this + \vee;
void operator = (const bigint &v) {
  *this = *this - v;
void operator*=(const bigint &v) {
  *this = *this * v;
void operator/=(const bigint &v) {
  *this = *this / v;
bool operator<(const bigint &v) const {</pre>
  if (sign != v.sign)
    return sign < v.sign;</pre>
  if (a.size() != v.a.size())
    return a.size() * sign < v.a.size() * v.sign;</pre>
  for (int i = a.size() - 1; i >= 0; i--)
    if (a[i] != v.a[i])
      return a[i] * sign < v.a[i] * sign;</pre>
 return false;
bool operator>(const bigint &v) const {
  return v < *this;
bool operator<=(const bigint &v) const {</pre>
  return ! (v < *this);
bool operator>=(const bigint &v) const {
  return ! (*this < v);
bool operator==(const bigint &v) const {
  return ! (*this < v) && ! (v < *this);
bool operator!=(const bigint &v) const {
  return *this < v || v < *this;
void trim() {
 while (!a.empty() && !a.back())
    a.pop_back();
  if (a.empty())
    sign = 1;
bool isZero() const {
  return a.empty() || (a.size() == 1 && !a[0]);
bigint operator-() const {
  bigint res = *this;
  res.sign = -sign;
  return res;
bigint abs() const {
 bigint res = *this;
  res.sign *= res.sign;
  return res;
long long longValue() const {
  long long res = 0;
```

```
for (int i = a.size() - 1; i >= 0; i--)
    res = res * base + a[i];
  return res * sign;
friend bigint gcd(const bigint &a, const bigint &b) {
  return b.isZero() ? a : gcd(b, a % b);
friend bigint lcm(const bigint &a, const bigint &b) {
  return a / gcd(a, b) * b;
void read(const string &s) {
  sign = 1;
  a.clear();
  int pos = 0;
  while (pos < (int) s.size() && (s[pos] == '-' || s[
     pos] == '+')) {
    if (s[pos] == '-') sign = -sign;
  for (int i = s.size() - 1; i >= pos; i -=
     base digits) {
    int x = 0;
    for (int j = max(pos, i - base_digits + 1); j <= i</pre>
       ; j++)
      x = x * 10 + s[j] - '0';
    a.push back(x);
  trim();
friend istream& operator>>(istream &stream, bigint &v)
  string s;
  stream >> s;
  v.read(s);
  return stream;
friend ostream& operator<<(ostream &stream, const</pre>
   bigint &v) {
  if (v.sign == -1) stream << '-';
  stream << (v.a.empty() ? 0 : v.a.back());</pre>
  for (int i = (int) v.a.size() - 2; i >= 0; --i)
    stream << setw(v.base_digits) << setfill('0') << v</pre>
        .a[i];
  return stream;
static vector<int> convert base(const vector<int> &a,
   int old_digits, int new_digits) {
  vector<long long> p(max(old_digits, new_digits) + 1)
  p[0] = 1;
  for (int i = 1; i < (int) p.size(); i++)</pre>
    p[i] = p[i - 1] * 10;
  vector<int> res;
  long long cur = 0;
  int cur digits = 0;
  for (int i = 0; i < (int) a.size(); i++) {</pre>
    cur += a[i] * p[cur digits];
    cur_digits += old_digits;
    while (cur_digits >= new_digits) {
      res.push_back(int(cur % p[new_digits]));
      cur /= p[new digits];
```

```
cur digits -= new digits;
  res.push back((int) cur);
  while (!res.empty() && !res.back())
    res.pop_back();
  return res;
typedef vector<long long> vll;
static vll karatsubaMultiply(const vll &a, const vll & 1
  int n = a.size();
  vll res(n + n);
  if (n \le 32) {
    for (int i = 0; i < n; i++)</pre>
      for (int j = 0; j < n; j++)
        res[i + j] += a[i] * b[j];
    return res;
  int k = n \gg 1;
 vll al(a.begin(), a.begin() + k);
 vll a2(a.begin() + k, a.end());
  vll b1(b.begin(), b.begin() + k);
  vll b2(b.begin() + k, b.end());
  vll a1b1 = karatsubaMultiply(a1, b1);
  vll a2b2 = karatsubaMultiply(a2, b2);
  for (int i = 0; i < k; i++)
    a2[i] += a1[i];
  for (int i = 0; i < k; i++)
    b2[i] += b1[i];
  vll r = karatsubaMultiply(a2, b2);
  for (int i = 0; i < (int) alb1.size(); i++)</pre>
    r[i] -= a1b1[i];
  for (int i = 0; i < (int) a2b2.size(); i++)
    r[i] -= a2b2[i];
  for (int i = 0; i < (int) r.size(); i++)</pre>
    res[i + k] += r[i];
  for (int i = 0; i < (int) alb1.size(); i++)</pre>
    res[i] += a1b1[i];
  for (int i = 0; i < (int) a2b2.size(); i++)</pre>
    res[i + n] += a2b2[i];
  return res;
bigint operator*(const bigint &v) const {
  vector<int> a6 = convert_base(this->a, base_digits,
  vector<int> b6 = convert_base(v.a, base_digits, 6);
  vll a(a6.begin(), a6.end());
  vll b(b6.begin(), b6.end());
  while (a.size() < b.size())</pre>
    a.push_back(0);
  while (b.size() < a.size())</pre>
    b.push_back(0);
  while (a.size() & (a.size() - 1))
    a.push_back(0), b.push_back(0);
  vll c = karatsubaMultiply(a, b);
  bigint res;
  res.sign = sign * v.sign;
  for (int i = 0, carry = 0; i < (int) c.size(); i++)</pre>
    long long cur = c[i] + carry;
```

```
res.a.push_back((int) (cur % 1000000));
    carry = (int) (cur / 1000000);
}
res.a = convert_base(res.a, 6, base_digits);
res.trim();
return res;
}
};
```

4 Graph algorithms

4.1 Bellman-Ford shortest paths with negative edge weights

```
// Single source shortest paths with negative edge
   weights.
// Returns false if a negative weight cycle is detected.
// Running time: O(|V|^3)
              start, w[i][j] = cost of edge from i to j
     INPUT:
     OUTPUT: dist[i] = min weight path from start to i
              dad[i] = prevector<int>ous node on the
   best path from the
                        start node
vector<int> dad;
vector<double> dist;
bool BellmanFord(int start, vector<vector<double>> &w) {
  int n = w.size();
  dad = vector < int > (n, -1);
  dist = vector<double>(n, 1000000000);
  dist[start] = 0;
  for (int k = 0; k < n; k++)
    for (int i = 0; i < n; i++)
      for (int j = 0; j < n; j++)
        if (dist[j] > dist[i] + w[i][j]){
          if (k == n-1) return false;
          else dist[j] = dist[i] + w[i][j], dad[j] = i;
  return true;
int main(){}
```

4.2 Eulerian path

```
vector<int> path;
void find_path(int v)
        while (adj[v].size() > 0)
                int vn = adj[v].front().next_vertex;
                adj[vn].erase(adj[v].front().
                    reverse_edge);
                adj[v].pop front();
                find_path(vn);
        path.push back(v);
void add_edge(int a, int b)
        adj[a].push_front(Edge(b));
        iter ita = adj[a].begin();
        adj[b].push_front(Edge(a));
        iter itb = adj[b].begin();
        ita->reverse_edge = itb;
        itb->reverse_edge = ita;
```

4.3 Minimum spanning trees

```
// This function runs Prim's algorithm for constructing
   minimum
  weight spanning trees.
// Running time: O(|V|^2)
     INPUT:
            w[i][j] = cost of edge from i to j
              NOTE: Make sure that w[i][j] is
   nonnegative and
              symmetric. Missing edges should be given
              weight.
     OUTPUT: edges = list of pair<int, int> in minimum
   spanning tree
              return total weight of tree
#include <iostream>
#include <queue>
#include <cmath>
#include <vector>
using namespace std;
typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;
typedef vector<int> VI;
typedef vector<VI> VVI;
typedef pair<int,int> PII;
typedef vector<PII> VPII;
T Prim (const VVT &w, VPII &edges) {
  int n = w.size();
  VI found (n);
```

```
VI prev (n, -1);
  VT dist (n, 1000000000);
  int here = 0;
  dist[here] = 0;
  while (here !=-1) {
    found[here] = true;
    int best = -1;
    for (int k = 0; k < n; k++) if (!found[k]) {</pre>
      if (w[here][k] != -1 && dist[k] > w[here][k]){
        dist[k] = w[here][k];
        prev[k] = here;
      if (best == -1 || dist[k] < dist[best]) best = k;</pre>
    here = best;
  T tot_weight = 0;
  for (int i = 0; i < n; i++) if (prev[i] != -1) {</pre>
    edges.push_back (make_pair (prev[i], i));
    tot_weight += w[prev[i]][i];
  return tot_weight;
int main(){
  int ww[5][5] = {
    {0, 400, 400, 300, 600},
    {400, 0, 3, -1, 7}, {400, 3, 0, 2, 0},
    {300, -1, 2, 0, 5}, {600, 7, 0, 5, 0}
  };
  VVT w(5, VT(5));
  for (int i = 0; i < 5; i++)
    for (int j = 0; j < 5; j++)
      w[i][j] = ww[i][j];
  // expected: 305
                3 2
                0 3
                2 4
  VPII edges;
  cout << Prim (w, edges) << endl;</pre>
  for (int i = 0; i < edges.size(); i++)</pre>
    cout << edges[i].first << " " << edges[i].second <<</pre>
        endl;
```

4.4 Centroid decomposition

```
set < int > v[100005];
map < int, int > mp[100005];
int n, up[100005][17], lvl[100005], par[100005], CNT, siz
      [100005], tin[100005], tout[100005];
void dfspre(int u, int dad=1, int depth = 0) {
    static int clk = 0;
    tin[u] = clk++;
    up[u][0] = dad;
    lvl[u] = depth;
```

```
for (int i=1; i<17; ++i)
    up[u][i] = up[up[u][i-1]][i-1];
  for(int i:v[u]) if(i!=dad)
    dfspre(i,u,depth+1);
  tout[u]=clk++;
int dfs(int u, int dad){
  siz[u] = 1;
  for(int i:v[u]) if(i!=dad)
    siz[u] += dfs(i,u);
  return siz[u];
int centroid(int u, int dad){
  for(int i:v[u]) if(i!=dad && siz[i]>CNT)
    return centroid(i,u);
  return u;
void decompose(int u, int dad) {
  CNT = dfs(u, dad)/2;
  int centre = centroid(u, dad);
  par[centre] = dad;
  for(int i:v[centre]) if(i!=dad){
    v[i].erase(centre);
    decompose (i, centre);
  v[centre].clear();
int lca(int u, int v) {
  if(|v|[u]>|v|[v]) swap(u,v);
  if(tin[u] <=tin[v] && tout[v] <=tout[u]) return u;</pre>
  for(int i=17;i--;)
    if(!(tin[up[u][i]]<=tin[v] && tout[v]<=tout[up[u][i</pre>
       11))
      u = up[u][i];
  return up[u][0];
void update(int u) {
  for(int node = u;u;u = par[u])
    ++mp[u][lvl[node]+lvl[u] - 2*lvl[lca(u, node)]];
int get(int u){
  int ans = INT_MAX;
  for(int node = u; u; u = par[u])
    ans = min(ans, lvl[u]+lvl[node]-2*lvl[lca(u, node)]+(*
       mp[u].begin()).first);
  return ans;
```

4.5 Heavy-Light decomposition

```
#include <bits/stdc++.h>
using namespace std;
int a[500005],sz[500005],lv1[500005];
int seg_id[500005],pos_id[500005],parent[500005];
int up[500005][19],tin[500005],tout[500005],clk,CNT;
vector<int> chain[500005];
template <typename T>
class SegmentTree{
    vector<T> segtree,lazy;
    public:
    SegmentTree(int size){
```

```
seatree.resize(4*size,0);
    lazy.resize(4*size,0);
  void update(int u, int a, int b, int i, int j, T x){
    if(lazy[u]) {
      segtree[u] += (b-a+1) * lazy[u];
      if (a!=b)
        lazy[u*2] += lazy[u], lazy[2*u+1] += lazy[u];
      lazy[u]=0;
    if(j<a || i>b || a>b) return;
    if(i) = b \&\& i <= a)
      segtree [u] += (b-a+1) *x;
      if (a!=b) lazy[u*2]+=x, lazy[2*u+1]+=x;
      return;
    update (u*2,a,(a+b)/2,i,j,x); update (u*2+1,(a+b)/2+1,
    segtree [u] = segtree [u*2] + segtree [u*2+1];
  void update(int i, T x) {
    update(1, 0, seqtree.size()/4-1, i, i, x);
  void update(int i, int j, T x) {
    update(1, 0, seqtree.size()/4-1, i, j, x);
  T query(int u, int a, int b, int i, int j){
    if(j<a || i>b || a>b) return 0;
    if(lazy[u]){
      segtree[u] += (b-a+1) * lazy[u];
      if (a!=b)
        lazy[u*2] += lazy[u], lazy[2*u+1] += lazy[u];
      lazy[u]=0;
    if(j>=b && i<=a) return segtree[u];</pre>
    return query (u*2,a,(a+b)/2,i,j) +query (u*2+1,1+(a+b)
       /2, b, i, j);
  T query(int i, int j){
    return query (1, 0, segtree.size()/4-1,i,j);
};
map<int,int> v[500005];
int dfs(int u, int dad=1, int depth=1, int last edge =
  tin[u]=clk++;
  up[u][0] = dad;
  lvl[u] = depth;
  a[u] = last edge;
  for (int i=1; i<19; ++i)
    up[u][i] = up[up[u][i-1]][i-1];
  lvl[u] = depth, sz[u] = 1;
  for(auto i:v[u]) if(i.first!=dad)
    sz[u] += dfs(i.first,u,depth+1, i.second);
  tout[u] = clk++;
  return sz[u];
void hld(int u, int dad = 1, int chain no = 0, int
   chain_parent = 0) {
  seq_id[\bar{u}] = chain_no;
  pos_id[u] = chain[chain_no].size();
```

```
parent[u] = chain parent;
  chain [chain no].push back(u);
  int max_sz = 0, heavy_child = -1;
  for(auto i:v[u]) if(i.first!=dad && max sz<sz[i.first</pre>
    tie (max_sz, heavy_child) = make_pair(sz[i.first], i.
        first);
  if (heavy_child!=-1)
    hld(heavy_child, u, chain_no, chain_parent);
  for(auto i:v[u]) if(i.first!=dad && i.first!=
     heavy child)
    hld(i.first,u,++CNT, u);
int lca(int u, int v) {
  if(lvl[u]>lvl[v]) swap(u,v);
  if(tin[u]<=tin[v] && tout[v]<=tout[u]) return u;</pre>
  for (int i=19;i--;)
    if(!(tin[up[u][i]]<=tin[v]&&tout[v]<=tout[up[u][i]])</pre>
      u = up[u][i];
  return up[u][0];
vector<SegmentTree> ST;
bool get(int x, int y) {
  if(seg_id[x] == seg_id[y]){
    if(pos_id[x]>pos_id[y]) swap(x,y);
    return ST[seg_id[x]].query(pos_id[x],pos_id[y]);
  if (seg_id[x]>seg_id[y]) swap(x,y);
  return get(x,parent[y]) ^ ST[seg id[y]].guery(0,pos id
      [y]);
int u[100005], T, q, n;
vector<pair<int,int>,int>> edges(1);
int main(){
  ios base::sync with stdio(0);
  cin.tie(0); cout.tie(0);
  for (int i=1, x, y, p; i < n; ++i) {</pre>
    cin>>x>>y>>p;
    p%=2;
    v[x][y] = p;
    v[v][x] = p;
    edges.push_back(\{\{x,y\},p\});
  dfs(1); hld(1);
  ST.clear();
  for (int i=0; i <= CNT; ++i) {</pre>
    ST.push back(SegmentTree(chain[i].size()));
    for(auto u:chain[i])
      ST[i].update(pos_id[u],a[u]);
  for(cin>>q;q--;){
    int type,x,y;
    cin>>type>>x>>y;
    if(type==1) {
      int 1 = lca(x, y);
      bool sum = qet(x,1) ^ qet(y,1);
      cout << (sum? "WE NEED BLACK PANDA\n": "JAKANDA
          FOREVER\n");
    } else {
```

5 Data structures

5.1 Suffix array

```
#include <bits/stdc++.h>
using namespace std;
vector<int> suffix_array(string &s){
  int n = s.size();
  vector<int> sa(n), buckets(n);
  for(int i=0;i<n;++i) sa[i] = n-i-1;</pre>
  stable_sort(sa.begin(),sa.end(),[&](int i, int j){
      return s[i] < s[j]; });
  for (int i=0; i < n; ++i) buckets[i] = s[i];</pre>
  for (int len=1;len<n;len*=2) {</pre>
    vector<int> b(buckets), cnt(n), s(sa);
    for (int i=0; i < n; ++i)
      buckets[sa[i]]=i\&\&b[sa[i-1]]==b[sa[i]]\&\&sa[i-1]+
          len < n\&\&b[sa[i-1]+len/2] == b[sa[i]+len/2]?buckets
          [sa[i-1]]:i;
    iota(cnt.begin(), cnt.end(),0);
    for (int i=0; i< n; ++i) if (s[i]>=len)
      sa[cnt[buckets[s[i]-len]]++]=s[i]-len;
  return sa;
vector<int> kasai(string &s, vector<int> &sa){
  int n = s.size();
  vector<int> lcp(n),inv(n);
  for (int i=0; i < n; ++i) inv[sa[i]] = i;</pre>
  for (int i=0, k=0; i < n; ++i) {
    if (k<0) k = 0;
    if(inv[i]==n-1) { k=0; continue; }
    for (int j=sa[inv[i]+1]; max(i,j)+k < n&&s[i+k]==s[j+k]
        ];++k];
    lcp[inv[i]] = k--;
  return lcp;
int main(){
  ios base::sync with stdio(0);
  cin.tie(0);
  string a.s:
  int K = 0;
  for(;cin>>a;++K) s += a + char(4+K);
  vector<int> color(s.size()), col(s.size());
  for (int i=0, cnt=0; i < s.size(); ++i)</pre>
    col[i]=cnt, cnt+=s[i]<20;
  auto sa = suffix_array(s);
  auto lcp = kasai(s,sa);
```

```
for(int i=0;i<lcp.size();++i) color[i]=col[sa[i]];</pre>
int freq[11] = {};
deque<int> v, mq;
multiset<int> ms;
int ans=0;
for (int i=1,COL=0;i<lcp.size();++i) {</pre>
  if(++freq[color[i]]==1) ++COL;
  while(v.size() && freq[color[v[0]]]>1){
    if (mq[0] == v[0])
      mq.pop front();
    --freq[color[v[0]]];
    v.pop_front();
  if(COL==K) ans = max(ans,lcp[mq[0]]);
  v.push_back(i);
  while (mq.size() && lcp[mq[0]]>lcp[i])
    mq.pop_front();
 mq.push_back(i);
cout << ans << '\n';
return 0:
```

5.2 KD-tree

```
// A straightforward, but probably sub-optimal KD-tree
   implmentation
// that's probably good enough for most things (current
   it's a
// 2D-tree)
   - constructs from n points in O(n lg^2 n) time
// - handles nearest-neighbor query in O(lq n) if
   points are well
    distributed
   - worst case for nearest-neighbor may be linear in
   pathological
     case
// Sonny Chan, Stanford University, April 2009
#include <iostream>
#include <vector>
#include <limits>
#include <cstdlib>
using namespace std;
// number type for coordinates, and its maximum value
typedef long long ntype;
const ntype sentry = numeric limits<ntype>::max();
// point structure for 2D-tree, can be extended to 3D
struct point {
    ntype x, y;
    point(ntype xx = 0, ntype yy = 0) : x(xx), y(yy) {}
};
bool operator==(const point &a, const point &b) {
    return a.x == b.x && a.y == b.y;
// sorts points on x-coordinate
bool on_x(const point &a, const point &b) {
```

```
return a.x < b.x;
// sorts points on y-coordinate
bool on_y(const point &a, const point &b) {
    return a.y < b.y;</pre>
// squared distance between points
ntype pdist2(const point &a, const point &b) {
    ntype dx = a.x-b.x, dy = a.y-b.y;
    return dx*dx + dy*dy;
// bounding box for a set of points
struct bbox{
    ntype x0, x1, y0, y1;
    bbox(): x0(sentry), x1(-sentry), y0(sentry), y1(-
       sentry) {}
    // computes bounding box from a bunch of points
    void compute(const vector<point> &v) {
        for (int i = 0; i < v.size(); ++i) {</pre>
            x0 = min(x0, v[i].x); x1 = max(x1, v[i].x)
            y0 = min(y0, v[i].y); y1 = max(y1, v[i].y)
    // squared distance between a point and this bbox, 0
        if inside
    ntype distance(const point &p) {
        if (p.x < x0) {
            if (p.y < y0)
                                 return pdist2 (point (x0,
               y0), p);
            else if (p.y > y1)
                                return pdist2(point(x0,
               y1), p);
            else
                                 return pdist2 (point (x0,
               p.y), p);
        else if (p.x > x1) {
            if (p.y < y0)
                                 return pdist2 (point (x1,
               y0), p);
            else if (p.y > y1)
                                 return pdist2 (point (x1,
               y1), p);
            else
                                 return pdist2 (point (x1,
               p.y), p);
        else {
            if (p.y < y0)
                                 return pdist2 (point (p.x,
                y0), p);
            else if (p.y > y1)
                                 return pdist2 (point (p.x,
                y1), p);
            else
                                 return 0;
        }
};
// stores a single node of the kd-tree, either internal
   or leaf
struct kdnode {
                    // true if this is a leaf node (has
    bool leaf;
       one point)
```

```
// the single point of this is a
   point pt;
       leaf
   bbox bound;
                   // bounding box for set of points in
        children
   kdnode *first, *second; // two children of this kd-
   kdnode() : leaf(false), first(0), second(0) {}
   ~kdnode() { if (first) delete first; if (second)
       delete second; }
   // intersect a point with this node (returns squared
        distance)
   ntype intersect(const point &p) {
        return bound.distance(p);
   // recursively builds a kd-tree from a given cloud
   void construct(vector<point> &vp) {
        // compute bounding box for points at this node
       bound.compute(vp);
        // if we're down to one point, then we're a leaf
        if (vp.size() == 1) {
            leaf = true;
           pt = vp[0];
       else {
            // split on x if the bbox is wider than high
                 (not best heuristic...)
            if (bound.x1-bound.x0 >= bound.y1-bound.y0)
                sort(vp.begin(), vp.end(), on x);
            // otherwise split on y-coordinate
                sort(vp.begin(), vp.end(), on_y);
            // divide by taking half the array for each
            // (not best performance if many duplicates
               in the middle)
            int half = vp.size()/2;
            vector<point> vl(vp.begin(), vp.begin()+half
            vector<point> vr(vp.begin()+half, vp.end());
            first = new kdnode();
                                  first->construct(vl)
            second = new kdnode(); second->construct(vr
               );
   }
// simple kd-tree class to hold the tree and handle
   queries
struct kdtree{
   kdnode *root;
   // constructs a kd-tree from a points (copied here,
       as it sorts them)
   kdtree(const vector<point> &vp) {
       vector<point> v(vp.begin(), vp.end());
       root = new kdnode();
```

```
root->construct(v);
    ~kdtree() { delete root; }
    // recursive search method returns squared distance
       to nearest point
    ntype search(kdnode *node, const point &p)
        if (node->leaf) {
            // commented special case tells a point not
               to find itself
              if (p == node->pt) return sentry;
              else
                return pdist2(p, node->pt);
        ntype bfirst = node->first->intersect(p);
        ntvpe bsecond = node->second->intersect(p);
        // choose the side with the closest bounding box
             to search first
        // (note that the other side is also searched if
            needed)
        if (bfirst < bsecond) {</pre>
            ntype best = search(node->first, p);
            if (bsecond < best)</pre>
                best = min(best, search(node->second, p)
            return best;
        else {
            ntype best = search(node->second, p);
            if (bfirst < best)</pre>
                best = min(best, search(node->first, p))
            return best;
    // squared distance to the nearest
    ntype nearest(const point &p) {
        return search(root, p);
// some basic test code here
int main(){
    // generate some random points for a kd-tree
    vector<point> vp;
    for (int i = 0; i < 100000; ++i) {
        vp.push_back(point(rand()%100000, rand()%100000)
           );
    kdtree tree(vp);
    // query some points
    for (int i = 0; i < 10; ++i) {
        point q(rand()%100000, rand()%100000);
        cout << "Closest squared distance to (" << q.x</pre>
           << ", " << q.y << ")"
             << " is " << tree.nearest(q) << endl;
```

5.3 Merge Sort Tree

```
#include <bits/stdc++.h>
using namespace std;
int A[1000005];
vector<int> segtree[400005];
void build(int u, int a, int b) {
        if (a==b) {
                 segtree[u].push_back(A[a]);
                 return;
        build (u*2,a,(a+b)/2); build (u*2+1,(a+b)/2+1,b);
        segtree[u].resize(b-a+1);
        merge(segtree[u*2].begin(), segtree[u*2].end(),
            segtree [u*2+1].begin(), segtree [u*2+1].end(),
            segtree[u].begin());
int query(int u, int a, int b, int i, int j, int k){
        if (b<a || j<a || i>b) return 0;
        if(i<=a && b<=j) return lower_bound(segtree[u].</pre>
            begin(), segtree[u].end(),k) - segtree[u].
         return query (u*2,a,(a+b)/2,i,j,k) +query (u*2+1,(a+b)/2,i,j,k)
            +b)/\bar{2}+1,b,i,j,k);
int main(){
         ios_base::sync_with_stdio(0);
         int n,q,low,high,mid,x,y,k;
         for (cin>>n>>q, x=0; x<n; cin>>A[x++]);
         for (build(1,0,n-1);q--;cout<<low-1<<'\n')</pre>
                 for (cin>>x>>y>>k, low=-1e9, high=1e9; low<
                     high;)
                          if (query (1, 0, n-1, x-1, y-1, mid=low)
                              +high>>1)< k) low=mid+1;
                          else high=mid;
```

6 Miscellaneous

6.1 Miller-Rabin Primality Test

```
Error rate: 2^(-TRIAL)
     Almost constant time, srand is needed
int64 t ModMul(int64 t a, int64 t b, int64 t m) {
  int64_t ret=0, c=a;
  for(;b;b>>=1, c=(c+c)%m)
    if(b&1) ret=(ret+c)%m;
  return ret;
int64 t ModExp(int64 t a, int64 t n, int64 t m) {
  return n?ModMul(ModExp(ModMul(a,a,m),n/2,m),(n%2?a:1),
     m):1:
bool Witness(int64_t a, int64_t n) {
  int64_t u=n-1;
    int t=0;
  while (!(u&1))\{u>>=1; t++;\}
  int64_t x0=ModExp(a, u, n), x1;
  for (int i=1; i<=t; i++) {</pre>
```

```
x1=ModMul(x0, x0, n);
if(x1==1&&x0!=1&&x0!=n-1) return true;
x0=x1;
}
if(x0!=1) return true;
return false;
}
bool IsPrimeFast(int64_t n, int TRIAL=15){
if(n<=2) return (n==2);
static random_device rd;
static mt19937_64 g(rd());
while(TRIAL--)
if(Witness(g()/2%(n-2)+1, n))
return false;
return true;
}</pre>
```

6.2 Pollard-Rho factorization

```
typedef long long unsigned int llui;
typedef long long int lli;
typedef long double float 64;
llui mul_mod(llui a, llui b, llui m) {
  llui y = (llui)((float64)a*(float64)b/m+(float64)1/2);
  y = y^* * m;
  llui x = a * b;
  llui r = x - y;
  if ((11i)r < 0){
    r = r + m; y = y - 1;
  return r;
llui C,a,b;
llui gcd() {
  llui c:
  if(a>b){
    c = a; a = b; b = c;
  while(1){
    if(a == 1LL) return 1LL;
    if(a == 0 || a == b) return b;
    c = a; a = b%a;
    b = c:
llui f(llui a, llui b) {
  llui tmp;
  tmp = mul\_mod(a,a,b);
  tmp+=C; tmp%=b;
  return tmp;
llui pollard(llui n) {
  if(!(n&1)) return 2;
  C=0;
  llui iteracoes = 0;
  while(iteracoes <= 1000) {</pre>
    llui x, y, d;
    x = y = 2; d = 1;
    while (d == 1) {
```

```
x = f(x,n);
       y = f(f(y,n),n);
       llui m = (x>y)?(x-y):(y-x);
       a = m; b = n; d = qcd();
    if(d != n)
       return d;
    iteracoes++; C = rand();
  return 1:
llui pot(llui a, llui b, llui c){
  if(b == 0) return 1;
  if(b == 1) return a%c;
  llui resp = pot(a,b>>1,c);
  resp = mul_mod(resp, resp, c);
  if(b&1)
   resp = mul_mod(resp,a,c);
  return resp;
// Rabin-Miller primality testing algorithm
bool isPrime(llui n) {
  llui d = n-1;
  llui s = 0;
  if (n \le 3 \mid \mid n == 5) return true;
  if(!(n&1)) return false;
  while(!(d&1)){ s++; d>>=1; }
  for(llui i = 0;i<32;i++) {
   llui a = rand();
    a <<=32;
    a+=rand();
    a\%=(n-3); a+=2;
    llui x = pot(a,d,n);
    if (x == 1 \mid | x == n-1) continue;
    for(llui j = 1; j<= s-1; j++) {
      x = mul\_mod(x, x, n);
      if(x == 1) return false;
      if (x == n-1) break;
    if (x != n-1) return false;
  return true;
map<llui,int> factors;
// Precondition: factors is an empty map, n is a
   positive integer
// Postcondition: factors[p] is the exponent of p in
   prime factorization of n
void fact(llui n) {
  if(!isPrime(n)){
    llui fac = pollard(n);
    fact (n/fac); fact (fac);
    map<llui,int>::iterator it;
    it = factors.find(n);
    if(it != factors.end()){
```

(*it).second++;

factors[n] = 1;

}else{

6.3 Manachers algorithm

```
// Maximal palindrome lengths centered around each
// position in a string (including positions between
   characters) and returns
// them in left-to-right order of centres. Linear time.
// Ex: "opposes" -> [0, 1, 0, 1, 4, 1, 0, 1, 0, 1, 0, 3,
    0, 1, 01
vector<int> fastLongestPalindromes(string str) {
  int i=0, j, d, s, e, lLen, palLen=0;
  vector<int> res;
  while (i < str.length()) {</pre>
    if (i > palLen && str[i-palLen-1] == str[i]) {
      palLen += 2; i++; continue;
    res.push back(palLen);
    s = res.size()-2;
    e = s-palLen;
    bool b = true;
    for (j=s; j>e; j--) {
      d = i - e - 1;
      if (res[j] == d) { palLen = d; b = false; break; }
      res.push_back(min(d, res[j]));
    if (b) { palLen = 1; i++; }
  res.push_back(palLen);
  lLen = res.size();
  s = 1Len-2;
  e = s-(2*str.length()+1-lLen);
  for (i=s; i>e; i--) { d = i-e-1; res.push\_back(min(d,
     res[i])); }
  return res;
```

6.4 Convex Hull Trick

```
struct Line {
  long long m, b;
 mutable function<const Line*()> succ;
 bool operator<(const Line& rhs) const{</pre>
    if (rhs.b != -(111<<62)) return m > rhs.m; // < for
    const Line* s = succ();
    if (!s) return 0;
    return b-s->b > (s->m-m)*rhs.m; // < for max
};
struct HullDynamic : public multiset<Line> {
 bool bad(iterator y) {
    auto z = next(y);
    if(y==begin()){
      if(z==end())return 0;
      return y->m == z->m && y->b >= z->b; // <= for max
    auto x = prev(v);
    if (z == end()) return y->m == x->m && y->b >= x->b;
         // <= for max
```

6.5 Dynamic Programming(DnC)

```
long long dp[21][100005];
void cost(int x, int y);
void computeDP(int idx,int jleft,int jright,int kleft,
   int kright) {
  if(jleft>jright) return;
  int jmid=(jleft+jright)/2;
  int bestk=jmid;
  for (int k=kleft; k<=min(kright, jmid); ++k) {</pre>
    cost(k, jmid);
    if (dp[idx-1][k-1]+tot<dp[idx][jmid])
      dp[idx][jmid]=dp[idx-1][k-1]+tot,bestk=k;
  computeDP(idx, jleft, jmid-1, kleft, bestk);
  computeDP(idx, jmid+1, jright, bestk, kright);
int main(){
  for (int i=0; i <= k; ++i)
    for (int j=0; j<=n; dp[i][j++]=1e17);</pre>
  dp[0][0]=0;
  for (int i=1; i <=k; ++i)
    computeDP (i, 1, n, 1, n);
  cout < \bar{q}p[k][n];
```

6.6 Longest increasing subsequence

```
// Given a list of numbers of length n, this routine
    extracts a longest increasing subsequence.
// Running time: O(n log n)
// INPUT: a vector of integers
// OUTPUT: a vector containing the longest increasing
    subsequence

typedef vector<int> VI;
typedef pair<int,int> PII;
typedef vector<PII> VPII;
#define STRICTLY_INCREASNG
VI LongestIncreasingSubsequence(VI v) {
    VPII best;
```

```
VI dad(v.size(), -1);
 for (int i = 0; i < v.size(); i++) {</pre>
#ifdef STRICTLY_INCREASNG
    PII item = make_pair(v[i], 0);
   VPII::iterator it = lower bound(best.begin(), best.
       end(), item);
   item.second = i;
#else
    PII item = make_pair(v[i], i);
   VPII::iterator it = upper_bound(best.begin(), best.
       end(), item);
#endif
   if (it == best.end()) {
      dad[i] = (best.size() == 0 ? -1 : best.back().
         second);
      best.push_back(item);
   } else {
      dad[i] = it == best.begin() ? -1 : prev(it)->
         second;
      *it = item;
 VI ret;
 for (int i = best.back().second; i >= 0; i = dad[i])
    ret.push_back(v[i]);
 reverse (ret.begin(), ret.end());
 return ret;
```

$\overline{6.7}$ Dates

```
// Months are expressed as integers from 1 to 12, Days
   are expressed as integers from 1 to 31, and Years are
    expressed as 4-digit integers.
string dayOfWeek[] = { "Mon", "Tue", "Wed", "Thu", "Fri",
    "Sat", "Sun"};
//converts Gregorian date to integer(Julian day number)
int dateToInt (int m, int d, int y) {
  return
    1461 * (v + 4800 + (m - 14) / 12) / 4 +
    367 * (m - 2 - (m - 14) / 12 * 12) / 12 -
    3 * ((y + 4900 + (m - 14) / 12) / 100) / 4 +
    d - 32075;
// converts integer (Julian day number) to Gregorian
   date: month/day/year
void intToDate (int jd, int &m, int &d, int &y) {
  int x, n, i, j;
  x = jd + 68569;
 n = 4 * x / 146097;
 x = (146097 * n + 3) / 4;
  i = (4000 * (x + 1)) / 1461001;
 x = 1461 * i / 4 - 31;
  j = 80 * x / 2447;
 d = x - 2447 * j / 80;
 x = \frac{1}{2} / 11;
 m = \frac{1}{1} + 2 - 12 * x;
 y = 100 * (n - 49) + i + x;
```

```
// converts integer (Julian day number) to day of week
string intToDay (int jd) {
   return dayOfWeek[jd % 7];
}
```

6.8 Knuth-Morris-Pratt

```
typedef vector<int> VI;
void buildPi(string& p, VI& pi){
  pi = VI(p.length());
  int k = -2;
  for(int i = 0; i < p.length(); i++) {</pre>
    while (k \ge -1 \& \& p[k+1] != p[i])
      k = (k == -1) ? -2 : pi[k];
    pi[i] = ++k;
int KMP(string& t, string& p){
  buildPi(p, pi);
  int k = -1;
  for(int i = 0; i < t.length(); i++) {</pre>
    while (k \ge -1 \&\& p[k+1] != t[i])
      k = (k == -1) ? -2 : pi[k];
    k++;
    if(k == p.length() - 1) {
      // p matches t[i-m+1, ..., i]
      cout << "matched at index " << i-k << ": ";</pre>
      cout << t.substr(i-k, p.length()) << endl;</pre>
      k = (k == -1) ? -2 : pi[k];
int main(){
  KMP ("AABAACAADAABAABA", "AABA"); //Mtches at: 0,9,12
```

6.9 2-SAT

```
struct TwoSat {
  int n;
  vector<vector<int> > adj, radj, scc;
  vector<int> sid, vis, val;
  stack<int> stk;
  int scnt;
  // n: number of variables, including negations
  TwoSat(int n): n(n), adj(n), radj(n), sid(n), vis(n),
    val(n, -1) {}
```

```
// adds an implication
  void impl(int x, int y) { adj[x].push_back(y); radj[y
     ].push_back(x); }
  // adds a disjunction
 void vee(int x, int y) { impl(x^1, y); impl(y^1, x); }
  // forces variables to be equal
 // forces variable to be true
 void tru(int x) { impl(x^1, x); }
 void dfs1(int x) {
   if (vis[x]++) return;
    for (int i = 0; i < adj[x].size(); i++)</pre>
     dfs1(adj[x][i]);
    stk.push(x);
 void dfs2(int x) {
   if (!vis[x]) return; vis[x] = 0;
    sid[x] = scnt; scc.back().push back(x);
   for (int i = 0; i < radj[x].size(); i++)</pre>
     dfs2(radj[x][i]);
  // returns true if satisfiable, false otherwise
  // on completion, val[x] is the assigned value of
     variable x (Note: val[x] = 0 implies val[x^1] = 1)
 bool two sat() {
   scnt = 0;
    for (int i = 0; i < n; dfs1(i++));</pre>
   while (!stk.empty()) {
     int v = stk.top(); stk.pop();
     if (vis[v]) {
       scc.push_back(vector<int>());
       dfs2(v);
       scnt++;
    for (int i = 0; i < n; i += 2)
     if (sid[i] == sid[i+1]) return false;
    vector<int> must(scnt);
    for (int i = 0; i < scnt; i++)</pre>
     for (int j = 0; j < scc[i].size(); j++){</pre>
       val[scc[i][j]] = must[i];
       must[sid[scc[i][j]^1]] = !must[i];
   return true;
};
```