Contents			7.1 2-SAT
1	Combinatorial optimization  1.1 Dinic's	1 1 2 3 4 4 4	7.3 DP Öptimizations       20         7.4 Convex Hull Trick (Dynamic)       21         7.5 Convex Hull Trick (Static)       21         7.6 BigInt library       21         7.7 Manachers algorithm       23         7.8 Dates       23         7.9 Bitset (Text)       23         7.10 Template       23         7.11 Numerical Integration       24
2	Geometry 2.1 Miscellaneous geometry 2.2 3D geometry 2.3 Convex hull 2.4 Min Enclosing Circle 2.5 Pick's Theorem (Text) 2.6 Slow Delaunay triangulation	5 61 7 71.1	Combinatorial optimization  Dinic's  struct Dinic {     struct Edge {         int u, v;     }
3	Numerical algorithms 3.1 Pollard Rho 3.2 Simplex algorithm 3.3 Reduced row echelon form 3.4 Fast Fourier transform 3.5 Number Theoretic transform 3.6 Discrete Logarithm 3.7 Mobius Inversion (Text) 3.8 Burnside Lemma (Text) 3.9 Number Theory (Modular, CRT, Linear Diophantine)	8 8 9 10 10 11 12 12 12 12	<pre>long long cap, flow; Edge() {} Edge(int u, int v, long long cap): u(u), v(v), cap(</pre>
4	Graph algorithms 4.1 Dynamic Connectivity	13 13 13 14	<pre>E.emplace_back(Edge(u, v, cap)); g[u].emplace_back(E.size() - 1); E.emplace_back(Edge(v, u, 0)); g[v].emplace_back(E.size() - 1); } </pre>
5	String Stuff 5.1 Suffix Automaton 5.2 Suffix array 5.3 Z Algorithm 5.4 KMP 5.5 String Hashing 5.6 Palindrome DSU 5.7 Eertree	15 15	<pre>bool BFS(int S, int T) {    queue<int> q({S});    fill(d.begin(), d.end(), N + 1);    d[S] = 0;    while(!q.empty()) {      int u = q.front(); q.pop();      if (u == T) break;      for (int k: g[u]) {         Edge &amp;e = E[k];         if (e.flow &lt; e.cap &amp;&amp; d[e.v] &gt; d[e.u] + 1) {         d[e.v] = d[e.u] + 1;</int></pre>
6	Data structures6.1 BIT Range Queries6.2 Treaps6.3 Link-Cut Tree	17 17 17 18	<pre>q.emplace(e.v); }  } return d[T] != N + 1; }</pre>
7	Miscellaneous	19	<pre>long long DFS(int u, int T, long long flow = -1) {   if (u == T    flow == 0) return flow;</pre>

```
for (int &i = pt[u]; i < q[u].size(); ++i) {
      Edge &e = E[g[u][i]];
     Edge &oe = E[g[u][i]^1];
      if (d[e.v] == d[e.u] + 1) {
        long long amt = e.cap - e.flow;
        if (flow != -1 \&\& amt > flow) amt = flow;
        if (long long pushed = DFS(e.v, T, amt)) {
          e.flow += pushed;
          oe.flow -= pushed;
         return pushed;
     }
    return 0;
  long long MaxFlow(int S, int T) {
   long long total = 0;
    while (BFS(S, T)) {
     fill(pt.begin(), pt.end(), 0);
     while (long long flow = DFS(S, T))
       total += flow;
   return total;
};
```

#### 1.2 Min-cost Circulation

```
// Runs in O(<max_flow > * log(V * max_edge_cost)) = <math>O(
   V^3 * // \log(V * C)) Really fast in // practice , 3e4
    edges are fine. Operates on integers , // costs are
   multiplied by N!!
template<typename flow_t = int, typename cost_t = int>
struct mcSFlow{
  struct Edge{
    cost_t č;
    flow_t f;
    int to, rev;
    Edge(int _to, cost_t _c, flow_t _f, int _rev):c(_c),
        f(_f), to(_to), rev(_rev){}
  static constexpr cost_t INFCOST = numeric_limits<</pre>
     cost_t>::max()/2;
  cost_t eps;
  int N, S, T;
  vector<vector<Edge> > G;
  vector<unsigned int> isq, cur;
  vector<flow_t> ex;
  vector<cost_t> h;
  mcSFlow(int _N, int _S, int _T):eps(0), N(_N), S(_S),
     T(_T), G(_N) \{ \}
  void add_edge(int a, int b, cost_t cost, flow_t cap) {
    assert (cap>=0);
    assert (a \ge 0 \& a \le N \& b \ge 0 \& b \le N);
    if (a==b) {assert (cost>=0); return; }
    cost*=N;
    eps = max(eps, abs(cost));
    G[a].emplace_back(b, cost, cap, G[b].size());
    G[b].emplace\_back(a, -cost, 0, G[a].size()-1);
  void add_flow(Edge& e, flow_t f) {
```

```
Edge &back = G[e.to][e.rev];
  if (!ex[e.to] && f)
   hs[h[e.to]].push_back(e.to);
  e.f -= f; ex[e.to] += f;
 back.f += f; ex[back.to] -= f;
vector<vector<int> > hs;
vector<int> co;
flow_t max_flow() {
  ex.assign(N, 0);
 h.assign(N, 0); hs.resize(2*N);
  co.assign(2*N, 0); cur.assign(N, 0);
 h[S] = N;
  ex[T] = 1;
  co[0] = N-1;
  for(auto &e:G[S]) add_flow(e, e.f);
  if(hs[0].size())
  for (int hi = 0; hi>=0;) {
    int u = hs[hi].back();
    hs[hi].pop_back();
    while (ex[u] > 0) { // discharge u
      if (cur[u] == G[u].size()) {
        h[u] = 1e9;
        for(unsigned int i=0;i<G[u].size();++i){</pre>
          auto &e = G[u][i];
          if (e.f && h[u] > h[e.to]+1) {
            h[u] = h[e.to] + 1, cur[u] = i;
        if (++\cos[h[u]], !--\cos[hi] \&\& hi < N)
          for (int i=0; i<N; ++i)
            if (hi < h[i] && h[i] < N) {
              --co[h[i]];
              h[i] = N + 1;
        hi = h[u];
      } else if (G[u][cur[u]].f && h[u] == h[G[u][cur[
        add_flow(G[u][cur[u]], min(ex[u], G[u][cur[u
           ] ] . f ) ) ;
      else ++cur[u];
    while (hi>=0 && hs[hi].empty()) --hi;
  return -ex[S];
void push(Edge &e, flow t amt) {
  if(e.f < amt) amt=e.f;</pre>
  e.f-=amt; ex[e.to]+=amt;
  G[e.to][e.rev].f+=amt; ex[G[e.to][e.rev].to]-=amt;
void relabel(int vertex) {
  cost_t newHeight = -INFCOST;
  for(unsigned int i=0;i<G[vertex].size();++i){</pre>
    Edge const&e = G[vertex][i];
    if(e.f && newHeight < h[e.to]-e.c){</pre>
      newHeight = h[e.to] - e.c;
      cur[vertex] = i;
 h[vertex] = newHeight - eps;
```

#### static constexpr int scale=2; pair<flow\_t, cost\_t> minCostMaxFlow() { cost\_t retCost = 0; for (int i=0; i<N; ++i)</pre> for (Edge &e:G[i]) retCost += e.c\*(e.f); //find max-flow flow\_t retFlow = max\_flow(); h.assign(N, 0); ex.assign(N, 0);isq.assign(N, 0); cur.assign(N, 0); queue<int> q; for(;eps;eps>>=scale){ //refine fill(cur.begin(), cur.end(), 0); for (int i=0; i<N; ++i) for(auto &e:G[i]) if(h[i] + e.c - h[e.to] < 0 && e.f) push(e, e.for (int i=0; i < N; ++i) {</pre> **if**(ex[i]>0){ q.push(i); isq[i]=1; // make flow feasible while(!q.empty()){ int u=q.front();q.pop(); isq[u]=0;while (ex[u]>0) { if(cur[u] == G[u].size()) relabel(u); for(unsigned int &i=cur[u], max\_i = G[u].size ();i<max\_i;++i){ Edge &e=G[u][i]; if(h[u] + e.c - h[e.to] < 0)push(e, ex[u]); **if**(ex[e.to]>0 && isq[e.to]==0){ q.push(e.to); isq[e.to]=1; if(ex[u]==0) break; **if**(eps>1 && eps>>scale==0) { eps = 1 << scale;for (int i=0; i<N; ++i) {</pre> **for**(Edge &e:G[i]){ retCost -= e.c\*(e.f); return make\_pair(retFlow, retCost/2/N); flow t getFlow(Edge const &e) { return G[e.to][e.rev].f; };

# 1.3 Edmonds Max Matching

```
#define MAXN 505
vector<int>q[MAXN];
int pa[MAXN], match[MAXN], st[MAXN], S[MAXN], v[MAXN];
inline int lca(int x,int y) {
  for (++t;;swap(x,y)) if (x) {
    if (v[x] == t) return x;
    v[x]=t, x=st[pa[match[x]]];
#define qpush(x) q.push(x), S[x]=0
void flower(int x,int y,int 1,queue<int> &q) {
  while(st[x]!=1){
    pa[x]=y, y=match[x];
    if(S[y] == 1) qpush(y);
    st[x]=st[y]=1, x=pa[y];
bool bfs(int x) {
  for (int i=1; i<=n; ++i) st[i]=i;</pre>
  memset (S+1, -1, sizeof(int)*n);
  queue<int>q;
  qpush(x);
  while(q.size()){
    x=q.front(),q.pop();
for(size_t i=0;i<g[x].size();++i){</pre>
      int y=q[x][i];
      if (S[v] = -1) {
         pa[y]=x, S[y]=1;
         if(!match[v]){
           for (int lst; x; y=lst, x=pa[y]) {
             lst=match[x], match[x]=y, match[y]=x;
           return 1;
         qpush (match[y]);
      } else if(!S[y]&&st[y]!=st[x]){
         int l=lca(y,x);
         flower(y, x, l, q); flower(x, y, l, q);
  return 0;
int blossom() {
  int ans=0;
  for (int i=1; i<=n; ++i)</pre>
    if(!match[i]&&bfs(i))++ans;
  return ans;
int main(){
  while (m--) {
    q[x].push back(y);
    g[y].push_back(x);
  printf("%d\n", blossom());
  for(int i=1;i<=n;i++)printf("%d ",match[i]);</pre>
```

## 1.4 Hungarian Min-cost

```
double MinCostMatching(const VVD &cost, VI &Lmate, VI &
  int n = int(cost.size());
  // construct dual feasible solution
  VD u(n);
  VD v(n);
  for (int i = 0; i < n; i++) {
    u[i] = cost[i][0];
    for (int j = 1; j < n; j++) u[i] = min(u[i], cost[i]
       ][ | ] );
 for (int j = 0; j < n; j++) {
  v[j] = cost[0][j] - u[0];</pre>
    for (int i = 1; i < n; i++) v[j] = min(v[j], cost[i
       ][j] - u[i]);
  // construct primal solution satisfying complementary
  Lmate = VI(n, -1);
  Rmate = VI(n, -1);
  int mated = 0;
  for (int i = 0; i < n; i++) {
    for (int j = 0; j < n; j++) {
      if (Rmate[j] != -1) continue;
      if (fabs(cost[i][j] - u[i] - v[j]) < 1e-10) {</pre>
        Lmate[i] = j;
        Rmate[j] = i;
        mated++;
        break;
  VD dist(n);
  VI dad(n);
  VI seen(n);
  // repeat until primal solution is feasible
  while (mated < n) {</pre>
    // find an unmatched left node
    int s = 0;
    while (Lmate[s] !=-1) s++;
    // initialize Dijkstra
    fill(dad.begin(), dad.end(), -1);
    fill(seen.begin(), seen.end(), 0);
    for (int k = 0; k < n; k++)
      dist[k] = cost[s][k] - u[s] - v[k];
    int i = 0;
    while (true) {
      // find closest
      i = -1;
      for (int k = 0; k < n; k++) {
        if (seen[k]) continue;
        if (j == -1 \mid | dist[k] < dist[j]) j = k;
      seen[j] = 1;
      // termination condition
      if (Rmate[j] == -1) break;
      // relax neighbors
      const int i = Rmate[j];
```

```
for (int k = 0; k < n; k++) {
      if (seen[k]) continue;
     const double new_dist = dist[j] + cost[i][k] - u
         [i] - v[k];
      if (dist[k] > new_dist) {
        dist[k] = new_dist;
        dad[k] = j;
  // update dual variables
  for (int k = 0; k < n; k++) {
    if (k == j || !seen[k]) continue;
    const int i = Rmate[k];
    v[k] += dist[k] - dist[j];
    u[i] -= dist[k] - dist[j];
 u[s] += dist[j];
  // augment along path
 while (dad[j] >= 0) {
    const int d = dad[j];
    Rmate[j] = Rmate[d];
    Lmate[Rmate[j]] = j;
    i = d;
 Rmate[j] = s;
 Lmate[s] = i;
 mated++:
double value = 0;
for (int i = 0; i < n; i++)
 value += cost[i][Lmate[i]];
return value;
```

### 1.5 Konig's Theorem (Text)

In any bipartite graph, the number of edges in a maximum matching equals the number of vertices in a minimum vertex cover. To exhibit the vertex cover:

1. Find a maximum matching

2. Change each edge **used** in the matching into a directed edge from **right to left** 

3. Change each edge **not used** in the matching into a directed edge from **left to right** 

4. Compute the set *T* of all vertices reachable from unmatched vertices on the left (including themselves)

5. The vertex cover consists of all vertices on the right that are **in** *T*, and all vertices on the left that are **not in** *T* 

# 1.6 Minimum Edge Cover (Text)

If a minimum edge cover contains C edges, and a maximum matching contains M edges, then C + M = |V|. To obtain the edge cover, start with a maximum matching, and then, for every vertex not matched, just select some edge incident upon it and add it to the edge set.

# 2 Geometry

### 2.1 Miscellaneous geometry

```
double INF = 1e100, EPS = 1e-12;
struct PT {
  double x, y;
  PT() {}
  PT (double x, double y) : x(x), y(y) {}
  PT(const PT \&p) : x(p.x), y(p.y)
  PT operator + (const PT &p) const { return PT(x+p.x,
     y+p.y);
  PT operator - (const PT &p) const { return PT(x-p.x,
     y-p.y); }
  PT operator * (double c)
                                const { return PT(x*c,
     V*C ); }
  PT operator / (double c)
                                const { return PT(x/c)
     y/c ); }
double dot (PT p, PT q)
                            { return p.x*q.x+p.y*q.y; }
double dist2(PT p, PT q) { return dot(p-q,p-q); }
double cross(PT p, PT q) { return p.x*q.y-p.y*q.x; }
ostream & operator << (ostream & os, const PT & p) {
 os << "(" << p.x << "," << p.y << ")";
// rotate a point CCW or CW around the origin
PT RotateCCW90(PT p) { return PT(-p.y,p.x); } PT RotateCW90(PT p) { return PT(p.y,-p.x); }
PT RotateCCW(PT p, double t) {
  return PT(p.x*cos(t)-p.y*sin(t), p.x*sin(t)+p.y*cos(t)
// project point c onto line through a and b
// assuming a != b
PT ProjectPointLine(PT a, PT b, PT c) {
  return a + (b-a)*dot(c-a, b-a)/dot(b-a, b-a);
// project point c onto line segment through a and b
PT ProjectPointSegment(PT a, PT b, PT c) {
  double r = dot(b-a, b-a);
  if (fabs(r) < EPS) return a;</pre>
  r = dot(c-a, b-a)/r;
  if (r < 0) return a;</pre>
  if (r > 1) return b;
  return a + (b-a) *r;
// compute distance from c to segment between a and b
double DistancePointSegment(PT a, PT b, PT c) {
  return sqrt(dist2(c, ProjectPointSegment(a, b, c)));
// compute distance between point (x,y,z) and plane ax+
double DistancePointPlane(double x, double y, double z,
                           double a, double b, double c,
                               double d)
  return fabs(a*x+b*y+c*z-d)/sqrt(a*a+b*b+c*c);
```

```
// determine if lines from a to b and c to d are
       parallel or collinear
bool LinesParallel(PT a, PT b, PT c, PT d) {
    return fabs(cross(b-a, c-d)) < EPS;</pre>
bool LinesCollinear(PT a, PT b, PT c, PT d) {
    return LinesParallel(a, b, c, d)
             && fabs(cross(a-b, a-c)) < EPS
             && fabs(cross(c-d, c-a)) < EPS;
// determine if line segment from a to b intersects with
 // line segment from c to d
bool SegmentsIntersect (PT a, PT b, PT c, PT d) {
    if (LinesCollinear(a, b, c, d)) {
         if (dist2(a, c) < EPS || dist2(a, d) < EPS ||
             dist2(b, c) < EPS || dist2(b, d) < EPS) return
         if (dot(c-a, c-b) > 0 \&\& dot(d-a, d-b) > 0 \&\& dot(c-a) > 0 \&\& dot(c-a) > 0 && dot(c-a) && dot(c-a) > 0 && dot(c-a) && 
               b, d-b) > 0
             return false;
         return true;
    if (cross(d-a, b-a) * cross(c-a, b-a) > 0) return
    if (cross(a-c, d-c) * cross(b-c, d-c) > 0) return
            false:
    return true;
// compute intersection of line passing through a and b
// with line passing through c and d, assuming that
 // intersection exists; for segment intersection, check
 // segments intersect first
PT ComputeLineIntersection(PT a, PT b, PT c, PT d) {
    b=b-a; d=c-d; c=c-a;
    assert (dot (b, b) > EPS && dot (d, d) > EPS);
    return a + b*cross(c, d)/cross(b, d);
 // compute center of circle given three points
PT ComputeCircleCenter(PT a, PT b, PT c) {
    b = (a+b)/2;
    c = (a + c) / 2;
    return ComputeLineIntersection(b, b+RotateCW90(a-b), c
            , c+RotateCW90(a-c);
// determine if point is in a possibly non-convex
       polygon (by William
// Randolph Franklin); returns 1 for strictly interior
       points, 0 for
// strictly exterior points, and 0 or 1 for the
       remaining points.
// Note that it is possible to convert this into an *
       exact* test using
// integer arithmetic by taking care of the division
       appropriately
// (making sure to deal with signs properly) and then by
         writing exact
```

```
// tests for checking point on polygon boundary
bool PointInPolygon(const vector<PT> &p, PT q) {
  bool c = 0;
  for (int i = 0; i < p.size(); i++) {</pre>
    int j = (i+1) %p.size();
    if ((p[i].y \le q.y \&\& q.y < p[j].y ||
      p[j].y \le q.y \&\& q.y < p[i].y) \&\&
      q.x < p[i].x + (p[j].x - p[i].x) * (q.y - p[i].y)
         / (p[j].y - p[i].y))
      c = !c;
  return c;
// determine if point is on the boundary of a polygon
bool PointOnPolygon(const vector<PT> &p, PT q) {
  for (int i = 0; i < p.size(); i++)</pre>
    if (dist2(ProjectPointSegment(p[i], p[(i+1)%p.size()
       ], q), q) < EPS)
      return true;
    return false;
// compute intersection of line through points a and b
   with
// circle centered at c with radius r > 0
vector<PT> CircleLineIntersection(PT a, PT b, PT c,
   double r) {
  vector<PT> ret;
  b = b-a;
  a = a-c;
  double A = dot(b, b);
  double B = dot(a, b);
  double C = dot(a, a) - r*r;
  double D = B*B - A*C;
  if (D < -EPS) return ret;</pre>
  ret.push back(c+a+b*(-B+sgrt(D+EPS))/A);
  if (D > EPS)
    ret.push_back(c+a+b*(-B-sqrt(D))/A);
  return ret;
// compute intersection of circle centered at a with
   radius r
// with circle centered at b with radius R
vector<PT> CircleCircleIntersection(PT a, PT b, double r
   , double R) {
  vector<PT> ret;
  double d = sqrt(dist2(a, b));
  if (d > r+R \mid | d+min(r, R) < max(r, R)) return ret;
  double x = (d*d-R*R+r*r)/(2*d);
  double y = sqrt(r*r-x*x);
  PT v = (b-a)/d;
  ret.push_back(a+v*x + RotateCCW90(v)*y);
  if (y > 0)
   ret.push_back(a+v*x - RotateCCW90(v)*y);
  return ret;
// This code computes the area or centroid of a (
   possibly nonconvex)
// polygon, assuming that the coordinates are listed in
   a clockwise or
```

```
// counterclockwise fashion. Note that the centroid is
   often known as
// the "center of gravity" or "center of mass".
double ComputeSignedArea(const vector<PT> &p) {
  double area = 0;
  for(int i = 0; i < p.size(); i++) {</pre>
    int j = (i+1) % p.size();
    area += p[i].x*p[j].y - p[j].x*p[i].y;
  return area / 2.0;
double ComputeArea(const vector<PT> &p) {
  return fabs(ComputeSignedArea(p));
PT ComputeCentroid(const vector<PT> &p) {
  PT c(0,0);
  double scale = 6.0 * ComputeSignedArea(p);
  for (int i = 0; i < p.size(); i++) {</pre>
    int i = (i+1) % p.size();
    c = c + (p[i]+p[j])*(p[i].x*p[j].y - p[j].x*p[i].y);
  return c / scale;
// tests whether or not a given polygon (in CW or CCW
   order) is simple
bool IsSimple(const vector<PT> &p) {
  for (int i = 0; i < p.size(); i++) {</pre>
    for (int k = i+1; k < p.size(); k++) {
      int i = (i+1) % p.size();
      int 1 = (k+1) % p.size();
      if (i == 1 \mid | i == k) continue;
      if (SegmentsIntersect(p[i], p[j], p[k], p[l]))
        return false;
  return true;
```

### 2.2 3D geometry

```
// distance from point (px, py, pz) to line (x1, y1,
   z1) - (x2, y2, z2)
// (or ray, or segment; in the case of the ray, the
   endpoint is the
// first point)
public static final int LINE = 0;
public static final int SEGMENT = 1;
public static final int RAY = 2;
public static double ptLineDistSq(double x1, double y1
   , double z1,
   double x2, double y2, double z2, double px, double
        py, double pz,
    int type) {
 double pd2 = (x1-x2) * (x1-x2) + (y1-y2) * (y1-y2) + (z1)
     -z2)*(z1-z2);
 double x, y, z;
 if (pd2 == 0) {
   x = x1;
   y = y1;
    z = z1;
  } else {
    double u = ((px-x1)*(x2-x1) + (py-y1)*(y2-y1) + (
       pz-z1)*(z2-z1)) / pd2;
   x = x1 + u * (x2 - x1);
   y = y1 + u * (y2 - y1);
   z = z1 + u * (z2 - z1);
   if (type != LINE && u < 0) {
     x = x1;
     y = y1;
     z = z1;
   if (type == SEGMENT && u > 1.0) {
     x = x2;
     y = y2;
     z = z2;
 return (x-px)*(x-px) + (y-py)*(y-py) + (z-pz)*(z-pz)
public static double ptLineDist(double x1, double y1,
   double z1,
   double x2, double y2, double z2, double px, double
        py, double pz,
    int type) {
 return Math.sqrt(ptLineDistSq(x1, y1, z1, x2, y2, z2
     , px, py, pz, type));
```

### 2.3 Convex hull

```
typedef pair<long long, long long> PT;
inline long long cross(PT o, PT a, PT b) {
  PT OA = {a.first-o.first,a.second-o.second};
  PT OB = {b.first-o.first,b.second-o.second};
  return OA.first*OB.second - OA.second*OB.first;
}
inline long double dist(PT a, PT b) {
```

```
return sqrt (pow (a.first-b.first, 2) +pow (a.second-b.
     second, 2));
vector<PT> convexhull(vector<PT> a) {
  vector<PT> hull;
  sort(a.begin(), a.end(), [](PT i, PT j){
    if(i.second!=j.second)
      return i.second < j.second;</pre>
    return i.first < j.first;
  });
  for (int i=0; i < a.size(); ++i) {</pre>
    while(hull.size()>1 && cross(hull[hull.size()-2],
       hull.back(),a[i])<=0)
      hull.pop_back();
    hull.push_back(a[i]);
  for(int i=a.size()-1, siz = hull.size();i--;){
    while(hull.size()>siz && cross(hull[hull.size()-2],
       hull.back(),a[i]) <= 0
      hull.pop_back();
    hull.push_back(a[i]);
  return hull;
```

# **2.4** Min Enclosing Circle

```
// Minimum enclosing circle, Welzl's algorithm
// Expected linear time.
// If there are any duplicate points in the input, be
   sure to remove them first.
struct point {
        double x;
        double y;
struct circle {
        double x;
        double y;
        double r;
        circle() {}
        circle(double x, double y, double r): x(x), y(y)
           r(r) \{ \}
circle b md(vector<point> R) {
        if (R.size() == 0) {
                return circle(0, 0, -1);
        } else if (R.size() == 1) {
                return circle(R[0].x, R[0].y, 0);
        } else if (R.size() == 2) {
                return circle((R[0].x+R[1].x)/2.0,
                               (R[0].y+R[1].y)/2.0,
                                           hypot (R[0].x-R
                                               [1].x, R
                                               [0].y-R[1].
                                               y)/2.0);
        } else {
                double D = (R[0].x - R[2].x)*(R[1].y - R
                    [2].y) - (R[1].x - R[2].x)*(R[0].y -
                    R[2].y);
                double p0 = (((R[0].x - R[2].x) * (R[0].x)
                    + R[2].x) + (R[0].y - R[2].y) * (R[0].y
```

```
+ R[2].y)) / 2 * (R[1].y - R[2].y) -
                     ((R[1].x - R[2].x)*(R[1].x + R[2].x)
                     + (R[1].y - R[2].y) * (R[1].y + R[2].y
                   )) /2 * (R[0].y - R[2].y))/D;
                double p1 = (((R[1].x - R[2].x) * (R[1].x)
                    + R[2].x) + (R[1].y - R[2].y) * (R[1].y
                    + R[2].y)) / 2 * (R[0].x - R[2].x) -
                     ((R[0].x - R[2].x)*(R[0].x + R[2].x)
                     + (R[0].y - R[2].y) * (R[0].y + R[2].y
                    )) /2 * (R[1].x - R[2].x))/D;
                return circle(p0, p1, hypot(R[0].x - p0,
                     R[0].y - p1));
circle b_minidisk(vector<point>& P, int i, vector<point>
    R)
        if (i == P.size() || R.size() == 3) {
                return b_md(R);
        } else {
                circle D = b_{minidisk}(P, i+1, R);
                if (hypot(P[i].x-D.x, P[i].y-D.y) > D.r)
                         R.push back(P[i]);
                         D = b \min idisk(P, i+1, R);
                return D;
// Call this function.
circle minidisk(vector<point> P) {
        random_shuffle(P.begin(), P.end());
        return b_minidisk(P, 0, vector<point>());
```

# 2.5 Pick's Theorem (Text)

For a polygon with all vertices on lattice points, A = i + b/2 - 1, where Akon is the area, i is the number of lattice points strictly within the polygon, and b is the number of lattice points on the boundary of the polygon. (Note, 3 there is no generalization to higher dimensions)

# 2.6 Slow Delaunay triangulation

```
// Slow but simple Delaunay triangulation. Does not
// handle degenerate cases (from O'Rourke)
//
// Running time: O(n^4)
// INPUT: x[] = x-coordinates
// OUTPUT: triples = a vector containing m triples of
// indices corresponding to triangle vertices
typedef double T;
struct triple {
  int i, j, k;
  triple() {}
  triple(int i, int j, int k) : i(i), j(j), k(k) {}
};
```

```
vector<triple> delaunayTriangulation(vector<T>& x,
   vector<T>& y) {
  int n = x.size();
  vector < T > z(n);
  vector<triple> ret;
  for (int i = 0; i < n; i++)
    z[i] = x[i] * x[i] + y[i] * y[i];
  for (int i = 0; i < n-2; i++) {
    for (int j = i+1; j < n; j++) {
      for (int k = i+1; k < n; k++) {
        if ( == k) continue;
        double xn = (y[j]-y[i])*(z[k]-z[i]) - (y[k]-y[i])
            ]) \star (z[j]-z[i]);
        double yn = (x[k]-x[i])*(z[j]-z[i]) - (x[j]-x[i])
            ])*(z[k]-z[i]);
        double zn = (x[j]-x[i])*(y[k]-y[i]) - (x[k]-x[i])
            ])*(y[j]-y[i]);
        bool flag = zn < 0;
        for (int m = 0; flag && m < n; m++)</pre>
          flag = flag && ((x[m]-x[i])*xn + (y[m]-y[i])*
              yn + (z[m]-z[i])*zn <= 0);
        if (flag) ret.push_back(triple(i, j, k));
    }
  return ret;
int main(){
  T \times s[] = \{0, 0, 1, 0.9\};
  T ys[] = \{0, 1, 0, 0.9\};
  vector<T> x(\&xs[0], \&xs[4]), y(\&ys[0], \&ys[4]);
  vector<triple> tri = delaunayTriangulation(x, y);
  //expected: 0 1 3
  for (int i = 0; i < tri.size(); i++)</pre>
    printf("%d %d %d\n", tri[i].i, tri[i].j, tri[i].k);
```

# Numerical algorithms

### 3.1 Pollard Rho

```
typedef long long unsigned int llui;
typedef long long int 111;
typedef long double float 64;
llui mul mod(llui a, llui b, llui m) {
  llui y = (llui)((float64)a*(float64)b/m+(float64)1/2);
  v = v^* * m;
  llui x = a * b;
  llui r = x - y;
  if ((11i)r < 0)
    r = r + m; y = y - 1;
  return r;
llui C,a,b;
llui gcd() {
  llui c:
  if(a>b) {
    c = a; a = b; b = c;
```

```
while (1) {
    if(a == 1LL) return 1LL;
    if(a == 0 || a == b) return b;
    c = a; a = b%a;
    b = c;
llui f(llui a, llui b){
  llui tmp;
  tmp = mul mod(a,a,b);
  tmp+=C; tmp%=b;
  return tmp;
llui pollard(llui n) {
  if(!(n&1)) return 2;
  C=0:
  llui iteracoes = 0;
  while(iteracoes <= 1000) {</pre>
    llui x,y,d;
    x = y = 2; d = 1;
    while (d == 1) {
       x = f(x, n);
       y = f(f(y,n),n);
       llui m = (x>y)?(x-y):(y-x);
       a = m; b = n; d = qcd();
    if(d != n)
       return d;
    iteracoes++; C = rand();
  return 1;
llui pot(llui a, llui b, llui c){
  if(b == 0) return 1;
  if(b == 1) return a%c;
  llui resp = pot(a,b>>1,c);
  resp = mul_mod(resp, resp, c);
  if (b&1)
    resp = mul_mod(resp,a,c);
  return resp;
// Rabin-Miller primality testing algorithm
bool isPrime(llui n) {
  llui d = n-1;
  llui s = 0;
  if (n \le 3 \mid \mid n == 5) return true;
  if(!(n&1)) return false;
  while(!(d&1)){ s++; d>>=1; }
  for(llui i = 0;i<32;i++) {
    llui a = rand();
    a <<=32;
    a + = rand();
    a\%=(n-3); a+=2;
    llui x = pot(a,d,n);
    if (x == 1 \mid | x == n-1) continue;
    for(llui j = 1; j<= s-1; j++) {
      x = mul\_mod(x, x, n);
      if(x == 1) return false;
      if (x == n-1) break;
```

```
if (x != n-1) return false;
  return true;
map<llui,int> factors;
// Precondition: factors is an empty map, n is a
   positive integer
// Postcondition: factors[p] is the exponent of p in
   prime factorization of n
void fact(llui n) {
  if(!isPrime(n)){
    llui fac = pollard(n);
    fact (n/fac); fact (fac);
  }else{
    map<llui,int>::iterator it;
    it = factors.find(n);
    if(it != factors.end()){
      (*it).second++;
    }else{
      factors[n] = 1;
```

### 3.2 Simplex algorithm

```
// Two-phase simplex algorithm for solving linear
   programs of the form
       maximize
                   C^T X
       subject to Ax \le b ; x \ge 0
// INPUT: Ā -- an m x n matrix
          b -- an m-dimensional vector
          c -- an n-dimensional vector
          x -- a vector where the optimal solution will
   be stored
// OUTPUT: value of the optimal solution (infinity if
          unbounded above, nan if infeasible)
// To use this code, create an LPSolver object with
// A, b, and c as arguments. Then, call Solve(x).
typedef long double DOUBLE;
typedef vector<DOUBLE> VD;
typedef vector<VD> VVD;
typedef vector<int> VI;
const DOUBLE EPS = 1e-9;
struct LPSolver {
  int m, n;
  VI B, N;
  LPSolver(const VVD &A, const VD &b, const VD &c) :
   m(b.size()), n(c.size()), N(n + 1), B(m), D(m + 2,
       VD(n + 2))
    for (int i = 0; i < m; i++) for (int j = 0; j < n; j
       ++) D[i][j] = A[i][j];
    for (int i = 0; i < m; i++) { B[i] = n + i; D[i][n]</pre>
       = -1; D[i][n + 1] = b[i]; 
    for (int j = 0; j < n; j++) { N[j] = j; D[m][j] = -c
       [j]; }
   N[n] = -1; D[m + 1][n] = 1;
  void Pivot(int r, int s)
    double inv = 1.0 / D[r][s];
```

```
for (int i = 0; i < m + 2; i++) if (i != r)
      for (int j = 0; j < n + 2; j++) if (j != s)
        D[i][j] = D[r][j] * D[i][s] * inv;
    for (int j = 0; j < n + 2; j++) if (j != s) D[r][j]
       *=inv;
    for (int i = 0; i < m + 2; i++) if (i != r) D[i][s]
       *=-inv;
    D[r][s] = inv;
    swap(B[r], N[s]);
  bool Simplex(int phase) {
    int x = phase == 1 ? m + 1 : m;
    while (true) {
      int s = -1;
      for (int j = 0; j \le n; j++) {
        if (phase == 2 \&\& N[i] == -1) continue;
        if (s == -1 || D[x][\dot{\eta}] < D[x][s] || D[x][\dot{\eta}] == D
            [x][s] \&\& N[j] < N[s]) s = j;
      if (D[x][s] > -EPS) return true;
      int r = -1;
      for (int i = 0; i < m; i++) {
        if (D[i][s] < EPS) continue;</pre>
        if (r == -1 \mid | D[i][n + 1] / D[i][s] < D[r][n +
           1) / D[r][s] ||
           (D[i][n + 1] / D[i][s]) == (D[r][n + 1] / D[r]
             [s]) && B[i] < B[r]) r = i;
      if (r == -1) return false;
      Pivot(r, s);
  DOUBLE Solve(VD &x) {
    int r = 0;
    for (int i = 1; i < m; i++) if (D[i][n + 1] < D[r][n
        + 11) r = i;
    if (D[r][n + 1] < -EPS) {
      Pivot(r, n);
      if (!Simplex(1) || D[m + 1][n + 1] < -EPS) return
          -numeric limits<DOUBLE>::infinity();
      for (int i = 0; i < m; i++) if (B[i] == -1) {
        int s = -1;
        for (int j = 0; j \le n; j++)
          if (s == -1 || D[i][j] < D[i][s] || D[i][j] ==</pre>
               D[i][s] \&\& N[j] < N[s]) s = j;
        Pivot(i, s);
    if (!Simplex(2)) return numeric limits<DOUBLE>::
       infinity();
    x = VD(n);
    for (int i = 0; i < m; i++) if (B[i] < n) \times[B[i]] =
       D[i][n + 1];
    return D[m][n + 1];
};
int main() {
  const int m = 4, n = 3;
  DOUBLE A[m][n] = {
    \{ 6, -1, 0 \},
```

```
{ -1, -5, 0 },
{ 1, 5, 1 },
{ -1, -5, -1 }
};
DOUBLE _b[m] = { 10, -4, 5, -5 }, _c[n] = { 1, -1, 0 };
VVD A(m);
VD b(_b, _b + m), c(_c, _c + n);
for (int i = 0; i < m; i++) A[i] = VD(_A[i], _A[i] + n );
LPSolver solver(A, b, c);
VD x;
DOUBLE value = solver.Solve(x);
cerr << "VALUE: " << value << endl; // VALUE: 1.29032
cerr << "SOLUTION:"; // SOLUTION: 1.74194 0.451613 1
for (size_t i = 0; i < x.size(); i++) cerr << " " << x
[i];</pre>
```

#### 3.3 Reduced row echelon form

```
// Reduced row echelon form via Gauss-Jordan elimination
// with partial pivoting. This can be used for
// computing the rank of a matrix.
// Running time: O(n^3)
// INPUT:
             a[][] = an nxm matrix
// OUTPUT:
             rref[][] = an nxm matrix (stored in a[][])
             returns rank of a[][]
const double EPSILON = 1e-10;
typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;
int rref(VVT &a) {
  int n = a.size(), m = a[0].size(), r = 0;
  for (int c = 0; c < m \&\& r < n; c++) {
    int j = r;
    for (int i = r + 1; i < n; i++)
      if (fabs(a[i][c]) > fabs(a[j][c])) j = i;
    if (fabs(a[j][c]) < EPSILON) continue;</pre>
    swap(a[j], a[r]);
    T s = 1.0 / a[r][c];
    for (int j = 0; j < m; j++) a[r][j] *= s;
    for (int i = 0; i < n; i++) if (i != r) {
      T t = a[i][c];
      for (int j = 0; j < m; j++) a[i][j] -= t * a[r][j</pre>
    r++;
  return r;
```

### 3.4 Fast Fourier transform

```
template<typename fpt>
struct fft_wrap {
  using cpx_t = complex<fpt>;
  const fpt two_pi = 4 * acosl(0);

  vector<cpx_t> roots; //stores the N-th roots of unity.
  int N;
```

```
fft_wrap(int N) : roots(N), N(N) {
    for (int i = 0; i < N; ++i) {
      roots[i] = EXP(two_pi * i / fpt(N));
  cpx_t EXP(fpt theta) {
   return {cos(theta), sin(theta)};
  void fft(cpx t *in, cpx t *out, int size, int dir){
    bit_reverse(in, out, size);
    for (int s = 0; (1 << s) < size; ++s) {
      int s_{-} = s + 1;
      for (int k = 0; k < size; k += (1 << s_{-})) {
        for (int j = 0; j < (1 << s); ++j) {
          int id = (N + dir * (N >> s) * †) & (N - 1);
          cpx_t w = roots[id];
          cpx_t t t = w * out[k + j + (1 << s)];
          cpx_t u = out[k + j];
          out[k + j] = u + t;
          out [k + j + (1 << s)] = u - t;
  void bit_reverse(cpx_t *in, cpx_t *out, int size){
    for (int i = 0; i < size; ++i) {</pre>
      int rev = 0, i_copy = i;
      for (int j = 0; (1 << j) < size; ++j) {
        rev = (rev << 1) + (i_copy & 1);
        i_copy >>= 1;
      out[rev] = in[i];
};
int main(){
 typedef complex<double> cpx_t;
  fft_wrap<double> fft_wrapper(2048);
 0.0}, out (8);
 fft_wrapper.fft(&in[0], &out[0], 8, 1);
fft_wrapper.fft(&out[0], &in[0], 8, -1);
 for (int i = 0; i < 8; ++i) {
  cout << in[i].real() << ' ' << in[i].imag() << endl;</pre>
```

# 3.5 Number Theoretic transform

```
// gen should be 5^((p - 1) / fft_wrapper_size)
template<int P>
struct field_t { ... };
template<typename fpt, int gen>
struct fft_wrap {
  vector<fpt> roots; //stores the N-th roots of unity.
  int N;
```

```
fft_wrap(int _N) : roots(_N), N(_N) {
    roots[0] = 1;
    for (int i = 1; i < N; ++i) {
      roots[i] = roots[i - 1] * qen;
      if (i != N - 1)
        assert(roots[i].v != 1);
    }
  void fft(fpt *in, fpt *out, int size, int dir) {
    bit_reverse(in, out, size);
    for (int s = 0; (1 << s) < size; ++s) {
      int s = s + 1;
      for (int k = 0; k < size; k += (1 << s_{-})) {
        for (int j = 0; j < (1 << s); ++j) {
          int id = (N + dir * (N >> s_{\underline{}}) * j) & (N - 1);
          fpt w = roots[id];
          fpt t = w * out[k + j + (1 << s)];
          fpt u = out[k + j];
          out[k + j] = u + t;
          out[k + j + (1 << s)] = u - t;
  void bit_reverse(fpt *in, fpt *out, int size) {
    for (int i = 0; i < size; ++i) {</pre>
      int rev = 0, i_copy = i;
for (int j = 0; (1 << j) < size; ++j) {</pre>
        rev = (rev << 1) + (i copy & 1);
        i copy >>= 1;
      out[rev] = in[i];
};
using fpt = field t<mod>;
fft_wrap<fpt, gen> fft_wrapper(1 << 18);</pre>
vector<fpt> polymul(vector<fpt> P1, vector<fpt> P2) {
  int fsize = P1.size() + P2.size() - 1;
  int N = max(P1.size(), P2.size());
  while (N & (N - 1))
    ++N:
  N \star = 2;
  P1.resize(N); P2.resize(N);
  vector<fpt> temp(N);
  fft_wrapper.fft(&P1[0], &temp[0], N, 1);
  fft_wrapper.fft(&P2[0], &P1[0], N, 1);
  for (int i = 0; i < N; ++i) {
    P1[i] *= temp[i];
  fft_wrapper.fft(&P1[0], &temp[0], N, -1);
  field_t<mod> inv(N);
  inv = inv.modexp(mod - 2);
  for (int i = 0; i < N; ++i) {
    temp[i] *= inv;
  temp.resize(fsize);
  return temp;
```

3.6 Discrete Logarithm

```
// Calculates x such that g^x % md == h
int baby_giant(int g, int h, int md) {
  unordered_map<int,int> mp;
  int sq = ceil(sqrtl(md));
  for(int i=0,now=1;i<sq;++i)
    mp[now] = i, now = (long long) now*g % md;
  for(int i=0,jmp=power(g,md-1-sq);i<sq;++i) {
    if(mp.find(h)!=mp.end()) return i*sq+mp[h];
    h = (long long) h*jmp % md;
  }
  return -1;
}</pre>
```

### 3.7 Mobius Inversion (Text)

```
\mu(n) = \begin{cases} 0 & n \text{ not squarefree} \\ 1 & n \text{ squarefree w/ even no. of prime factors} \\ 1 & n \text{ squarefree w/ odd no. of prime factors} \end{cases}
Note that \mu(a)\mu(b) = \mu(ab) for a, b relatively prime
\text{Also } \sum_{d|n} \mu(d) = \begin{cases} 1 & \text{if } n = 1 \\ 0 & \text{otherwise} \end{cases}
Möbius Inversion If a(n) = \sum_{l} \mu(d) for all n > l
```

**Möbius Inversion** If  $g(n) = \sum_{d|n} f(d)$  for all  $n \ge 1$ , then  $f(n) = \sum_{d|n} \mu(d)g(n/d)$  for all  $n \ge 1$ .

# 3.8 Burnside Lemma (Text)

The number of orbits of a set X under the group action G equals the average number of elements of X fixed by the elements of G.

Here's an example. Consider a square of 2n times 2n cells. How many ways are there to color it into X colors, up to rotations and/or reflections? Here, the group has only 8 elements (rotations by 0, 90, 180 and 270 degrees, reflections over two diagonals, over a vertical line and over a horizontal line). Every coloring stays itself after rotating by 0 degrees, so that rotation has  $X^{4n^2}$  fixed points. Rotation by 180 degrees and reflections over a horizonal/vertical line split all cells in pairs that must be of the same color for a coloring to be unaffected by such rotation/reflection, thus there exist  $X^{2n^2}$  such colorings for each of them. Rotations by 90 and 270 degrees split cells in groups of four, thus yielding  $X^{n^2}$  fixed colorings. Reflections over diagonals split cells into 2n groups of 1 (the diagonal itself) and  $2n^2 - n$  groups of 2 (all remaining cells), thus yielding  $X^{2n^2-n+2n} = X^{2n^2+n}$  unaffected colorings. So, the answer is  $(X^{4n^2} + 3X^{2n^2} + 2X^{n^2} + 2X^{2n^2+n})/8$ .

# 3.9 Number Theory (Modular, CRT, Linear Diophantine)

```
// All algorithms described here work on nonnegative
   integers.
int mod(int a, int b) {
  return ((a%b) + b) % b;
int lcm(int a, int b) {
  return a / __gcd(a, b) *b;
int powermod(int a, int b, int m) {
  return b?powermod(a*a%m,b/2,m)*(b%2?a:1)%m:1;
^{\prime}// returns q=qcd(a,\ b); finds x, y such that d=ax+
int extended_euclid(int a, int b, int &x, int &y) {
  int xx = y = 0;
  int yy = x = 1;
  while (b) {
    int q = a / b;
    int t = b; b = a%b; a = t;
   t = xx; xx = x - q*xx; x = t;
    t = yy; yy = y - q*yy; y = t;
  return a;
// finds all solutions to ax = b (mod n)
VI modular_linear_equation_solver(int a, int b, int n) {
  int x, y;
  VI ret;
  int g = extended_euclid(a, n, x, y);
  if (!(b%g))
    x = mod(x*(b / g), n);
    for (int i = 0; i < q; i++)
      ret.push back(mod(x + i*(n / q), n));
  return ret;
// computes b such that ab = 1 \pmod{n}, returns -1 on
int mod_inverse(int a, int n) {
  int g = extended_euclid(a, n, x, y);
  if (q > 1) return -1;
  return mod(x, n);
// Chinese remainder theorem (special case): find z such
// that z % m1 = r1, z % m2 = r2.
// Here, z is unique modulo M = lcm(m1, m2).
// Return (z, M). On failure, M = -1.
PII chinese_remainder_theorem(int m1, int r1, int m2,
   int r2) {
  int s, t;
  int g = extended_euclid(m1, m2, s, t);
  if (r1%g != r2%g) return make_pair(0, -1);
  return make pair(mod(s*r2*m1 + t*r1*m2, m1*m2) / q, m1
     *m2 / q);
// Find z such that z % m[i] = r[i] for all i.
// The solution is unique modulo M = lcm_i \ (m[i]) .
// Return (z, M). On failure, M = -1.
// We don't require a[i]'s to be relatively prime.
PII chinese remainder theorem (const VI &m, const VI &r) {
```

```
PII ret = make_pair(r[0], m[0]);
  for (int i = 1; i < m.size(); i++) {</pre>
    ret = chinese_remainder_theorem(ret.second, ret.
       first, m[i], r[i]);
    if (ret.second == -1) break;
  return ret;
// computes x and y such that ax + by = c
// returns whether the solution exists
bool linear_diophantine(int a, int b, int c, int &x, int
    &y) {
  if (!a && !b) {
    if (c) return false;
    x = 0; v = 0;
    return true;
  if (!a) {
    if (c % b) return false;
    x = 0; y = c / b;
    return true;
  if (!b) {
    if (c % a) return false;
    x = c / a; y = 0;
    return true;
  int g = \underline{gcd(a, b)};
  if (c % q) return false;
  x = c / q * mod_inverse(a / g, b / g);
  v = (c - a*x) / b;
  return true;
```

# 4 Graph algorithms

# 4.1 Dynamic Connectivity

```
struct UnionFind {
  int n, comp;
  vector<int> uf,si,c;
  UnionFind(int n=0):n(n),comp(n),uf(n),si(n,1){
    for (int i=0; i< n; ++i)
      uf[i]=i;
  int find(int x) {return x==uf[x]?x:find(uf[x]);}
  bool join(int x, int y) {
    if((x=find(x)) == (y=find(y))) return false;
    if(si[x] < si[y]) swap(x,y);
    si[x] + = si[y]; uf[y] = x; comp--;
    c.push back(y);
    return true;
  int snap() {return c.size();}
  void rollback(int snap) {
    while(c.size()>snap) {
      int x=c.back(); c.pop_back();
      si[uf[x]]-=si[x];uf[x]=x;comp++;
};
```

```
enum {ADD, DEL, QUERY};
struct Query {int type, x, y; };
struct DynCon {
  vector<Query> q;
  UnionFind dsu;
  vector<int> mt;
  map<pair<int,int>,int> last;
  DynCon(int n):dsu(n){}
  void add(int x, int y) {
    if(x>y) swap(x,y);
    q.push_back((Query){ADD,x,y}),mt.push_back(-1);
    last [make_pair(x,y)] = q.size()-1;
  void remove(int x, int y) {
    if(x>y) swap(x,y);
    q.push_back((Query) {DEL,x,y});
    int pr=last[make_pair(x,y)];
    mt[pr]=q.size()-1;
    mt.push_back(pr);
  void query(int x, int y) {
    q.push_back((Query) {QUERY, x, y});
    mt.push_back(-1);
  void process() { // answers all queries in order
    if(!q.size()) return;
    for (int i=0; i < q. size(); ++i)</pre>
      if(q[i].type==ADD&&mt[i]<0)
        mt[i]=q.size();
    go(0,q.size());
  void go(int 1, int r) {
    if(1+1==r) {
      if(q[1].type==QUERY) // answer query using DSU
        puts (dsu.find(q[1].x) == dsu.find(q[1].y)? "YES":"
      return;
    int s=dsu.snap(), m=(1+r)/2;
    for (int i=r-1; i>=m; --i)
      if (mt [i] >= 0 \& \&mt [i] < 1)
        dsu.join(q[i].x,q[i].y);
    qo(1,m);
    dsu.rollback(s);
    for (int i=m-1;i>=1;--i) if (mt[i]>=r) dsu.join(q[i].x,q
        [i].y);
    go(m,r);
    dsu.rollback(s);
} ;
```

## 4.2 Bridges

```
// Finds bridges and cut vertices
//
// Receives:
// N: number of vertices
// 1: adjacency list
//
// Gives:
// vis, seen, par (used to find cut vertices)
```

```
// ap - 1 if it is a cut vertex, 0 otherwise
// brid - vector of pairs containing the bridges
typedef pair<int, int> PII;
int N;
vector <int> 1[MAX];
vector <PII> brid;
int vis[MAX], seen[MAX], par[MAX], ap[MAX];
int cnt, root;
void dfs(int x) {
  if(vis[x] != -1)
    return;
  vis[x] = seen[x] = cnt++;
  int adj = 0;
  for(int i = 0; i < (int)1[x].size(); i++){</pre>
    int v = l[x][i];
    if(par[x] == v)
      continue;
    if(vis[v] == -1){
      adj++;
      par[v] = x;
      dfs(v);
seen[x] = min(seen[x], seen[v]);
      if(seen[v] >= vis[x] && x != root)
        ap[x] = 1;
      if(seen[v] == vis[v])
        brid.push_back(make_pair(v, x));
    else{
      seen[x] = min(seen[x], vis[v]);
      seen[v] = min(seen[x], seen[v]);
  if(x == root) ap[x] = (adj>1);
void bridges() {
  brid.clear();
  for (int i = 0; i < N; i++) {
    vis[i] = seen[i] = par[i] = -1;
    ap[i] = 0;
  cnt = 0;
  for (int i = 0; i < N; i++)
    if(vis[i] == -1){
      root = i;
      dfs(i);
```

## 4.3 Strongly connected components

```
struct SCC {
  int V, group_cnt;
  vector<vector<int> > adj, radj;
  vector<int> group_num, vis;
  stack<int> stk;
  // V = number of vertices
  SCC(int V): V(V), group_cnt(0), group_num(V), vis(V),
      adj(V), radj(V) {}
```

```
// Call this to add an edge (0-based)
 void add_edge(int v1, int v2) {
    adj[v1].push_back(v2);
    radj[v2].push_back(v1);
 void fill forward(int x) {
    vis[x] = true;
    for (int i = 0; i < adj[x].size(); i++) {</pre>
      if (!vis[adj[x][i]])
        fill_forward(adj[x][i]);
    stk.push(x);
 void fill backward(int x) {
    vis[x] = false;
    group_num[x] = group_cnt;
    for (int i = 0; i < radj[x].size(); i++) {</pre>
      if (vis[radj[x][i]]) {
        fill_backward(radj[x][i]);
   }
  // Returns number of strongly connected components.
  // After this is called, group num contains component
     assignments (0-based)
  int get_scc() {
    for (int i = 0; i < V; i++) {
      if (!vis[i]) fill_forward(i);
    group\_cnt = 0;
   while (!stk.empty()) {
      if (vis[stk.top()]) {
        fill_backward(stk.top());
        group_cnt++;
      stk.pop();
    return group_cnt;
} ;
```

# 5 String Stuff

### 5.1 Suffix Automaton

```
struct SuffixAutomaton {
  vector<map<char,int>> edges; // edges[i]: the labeled
      edges from node i
  vector<int> link; // link[i] : the parent of i
  vector<int> length; // length[i]: length of longest
      string in ith class
  vector<int> cnt; // No. of times substring occurs
  int last; // index of equivalence class of whole
      string
SuffixAutomaton(string const& s) {
      // add the initial node
      edges.push_back(map<char,int>());
      link.push_back(-1);
      length.push_back(0);
```

```
cnt.push back(0);
last = 0;
for (int i=0; i < s.size(); i++) {</pre>
  // construct r
  edges.push back(map<char,int>());
  length.push_back(i+1);
  link.push back(0);
 cnt.push_back(1);
  int r = edges.size() - 1;
  // add edges to r and find p with link to q
  int p = last;
  while (p \ge 0 \& \& !edges[p].count(s[i])) 
    edges[p][s[i]] = r;
    p = link[p];
  if (p !=-1) {
    int q = edges[p][s[i]];
    if(length[p] + 1 == length[q]) {
      // we do not have to split q, just set the
         correct suffix link
      link[r] = q;
    } else {
      // we have to split, add q'
      edges.push_back(edges[q]); // copy edges of q
      length.push_back(length[p] + 1);
      link.push_back(link[q]); // copy parent of q
      cnt.push_back(0);
      int qq = edges.size()-1;
      // add qq as the new parent of q and r
      link[q] = link[r] = qq; cnt[r] = 1;
      // move short classes pointing to q to point
      while (p >= 0 \& \& edges[p][s[i]] == q)  {
        edges[p][s[i]] = qq;
        p = link[p];
  last = r;
vector<int> ind(length.size());
iota(ind.begin(), ind.end(), 0);
sort(ind.begin(), ind.end(), [&](int i, int j){
  return length[i] > length[j];
for (auto i:ind) if (link[i] >= 0)
 cnt[link[i]] += cnt[i];
```

# 5.2 Suffix array

};

```
vector<int> suffix_array(string &A) {
  int n=A.size(),i=n, *M=new int[5*n];
  int *B=M,*C=M+n,*F=M+2*n,*G=M+3*n,*S=M+4*n;
  for(;i--;S[i]=n-i-1) B[i]=A[i];
  stable_sort(S,S+n,[&](int i,int j) {return A[i]<A[j];})
  for(int L=1,p;L<n;L*=2) {
    for(;++i<n;F[i]=B[S[i]],G[i]=B[S[i]+L/2]);</pre>
```

```
for(;--i;F[i]=F[i]==F[i-1]&&G[i]==G[i-1]&&S[i-1]<n-L
    for (p=B[*S]=0;++i< n;B[S[i]]=p=F[i]?p:i);
    for (fill_n(G,n,0);i--;F[i]=S[i]<L?-1:B[S[i]-L]);
    for (iota(C, C+n, 0); ++i<n; ~F[i]?G[i]=C[F[i]]++:0);</pre>
    for (copy_n(S, n, F); i--; F[i] < L?0:S[G[i]] = F[i]-L);
  vector<int>res(S,S+n);
  delete[] M;
  return res;
vector<int> kasai(string &s, vector<int> &sa){
  int n = s.size();
  vector<int> lcp(n),inv(n);
  for (int i=0; i < n; ++i) inv[sa[i]] = i;</pre>
  for (int i=0, k=0; i < n; ++i) {
    if(k<0) k = 0;
    if(inv[i]==n-1) { k=0; continue; }
    for (int j=sa[inv[i]+1]; max(i,j)+k < n&&s[i+k]==s[j+k]
    lcp[inv[i]] = k--;
  return lcp;
```

# 5.3 Z Algorithm

```
vector<int> compute_Z(string s) {
  int n = s.length();
  vector<int> z(n, 0);
  z[0] = n;
  for (int i = 1, l = 0, r = 0; i < n; ++i) {
    if (r >= i) {
        z[i] = min(z[i - l], r - i + 1);
    }
    while (i+z[i] < n and s[i+z[i]] == s[z[i]]) {
        ++z[i];
    }
  if (i + z[i] - 1 > r) {
        r = i + z[i] - 1;
        l = i;
    }
  }
  return z;
}
```

### **5.4** KMP

```
vector<int> find_prefix(const vector<int> &P) {
   int M = P.size();
   vector<int> pi(M);
   /* pi[i] <- largest prefix P[0..pi[i]] which is a
        suffix of P[0..i]
   * (but not equal to it) */
   pi[0] = -1;
   for (int i = 1, k = -1; i < M; ++i) {
      while(k > -1 && P[k + 1] != P[i])
        k = pi[k];
      if (P[k + 1] == P[i]) ++k;
      pi[i] = k;
   }
}
```

### 5.5 String Hashing

```
struct hasher{
  int hashes[MAXN+5];
  int *pow, *inv;
  int mod;
  int n;
  void init(string &str, int *p, int *i, int m) {
    pow = p;
    inv = i;
    mod = m;
    n = str.size();
    int last = 0;
    for (int i = 0; i < n; i++) {
      int c = str[i] - 'a' + 1;
      last = (last + 1ll*c*pow[i]) % mod;
      hashes[i] = last;
  int getHash(int 1, int r){
    if (r >= n || 1 < 0)
      return -1;
    int curr = hashes[r] - (1-1 >= 0 ? hashes[1-1] : 0);
    curr = ((curr % mod) + mod) % mod;
    curr = (1ll*curr*inv[1]) % mod;
    return curr;
} A, B, C;
```

### **5.6** Palindrome DSU

```
// given an unknown string s and Q ranges that
// are know to be palindromes, this computes the
    characters
// that have to be equal in O(Q + n log n)
struct Palindrome{
    Palindrome() {}
    Palindrome(int n_):n(n_), m(3+__lg(n)), qs(m), p(2*n) {
        iota(p.begin(), p.end(), 0);
    }
    int f(int i) {
        return p[i] == i ? i : p[i] = f(p[i]);
    }
    void u(int a, int b) {
```

```
assert(0 \le a \&\& a < 2*n);
  assert (0 <= b \& \& b < 2*n);
  // union with splicing is a bit faster than
  // just path compression also guarantees p[i] <= i
  while(p[a] != p[b]){
    if(p[a] < p[b]) swap(a, b);
    if(p[a] == a)
      p[a] = b;
      return;
    int tmp = p[a];
    p[a] = p[b];
    a = p[tmp];
int components(){
  int ret = 0;
  for (int i=0; i < (int) p.size(); ++i)</pre>
    if(p[i] == i) ++ret;
  return ret;
// call this after adding all queries
void compute() {
  vector<int> p2(2*n);
  for(int l=m-1; l>=0; --1) {
    const int s = 1 << 1;
    for (int i=0; i<2*n; ++i) p2[i] = f(i);
    for (int i=0; i+s<2*n; ++i) {
      const int j = p2[i];
      if (j+s < 2*n) u (i+s, j+s);
    for(auto const&e:qs[1])
      u(e.first, e.second);
  // link point with mirror-image
  for (int i=0; i< n; ++i)
    u(i, 2*n-1 - i);
// force [l, r] to be a palindrome
void add_q(int 1, int r) {
  assert (0 \le 1 \&\& 1 \le r \&\& r \le n);
  if(l==r) return;
  const int range = r-l+1;
  const int k = __lg(range);
  qs[k].emplace back(1, 2*n-1 - r);
int n, m;
vector<vector<pair<int, int> > > qs;
vector<int> p;
```

### 5.7 Eertree

```
/*
Palindrome tree. Useful structure to deal with
   palindromes in strings. O(N)
This code counts no. of palindrome substrings of string.
   Based on problem 1750 from informatics.mccme.ru:
http://informatics.mccme.ru/moodle/mod/statements/view.
   php?chapterid =1750
*/
```

```
const int MAXN = 105000;
struct node
    int next[26];
    int len;
    int sufflink;
    int num;
int len;
char s[MAXN];
node tree[MAXN];
                    // node 1 - root with len -1, node 2
    - root with len 0
int suff;
                    // max suffix palindrome
long long ans;
bool addLetter(int pos) {
    int cur = suff, curlen = 0;
    int let = s[pos] - 'a';
    while (true) {
        curlen = tree[cur].len;
        if (pos - 1 - curlen) = 0 && s[pos - 1 - curlen]
            == s[pos])
            break:
        cur = tree[cur].sufflink;
    if (tree[cur].next[let]) {
        suff = tree[cur].next[let];
        return false;
    num++;
    suff = num;
    tree[num].len = tree[cur].len + 2;
    tree[cur].next[let] = num;
    if (tree[num].len == 1) {
        tree[num].sufflink = 2;
        tree[num].num = 1;
        return true;
    while (true) {
        cur = tree[cur].sufflink;
        curlen = tree[cur].len;
        if (pos - 1 - curlen >= 0 && s[pos - 1 - curlen]
            == s[pos]) {
            tree[num].sufflink = tree[cur].next[let];
            break;
    tree[num].num = 1 + tree[tree[num].sufflink].num;
    return true;
void initTree() {
    num = 2; suff = 2;
    tree[1].len = -1; tree[1].sufflink = 1;
    tree[2].len = 0; tree[2].sufflink = 1;
int main() {
    gets(s);
    len = strlen(s)
    initTree()
    for (int i = 0; i < len; i++) {</pre>
        addLetter(i);
```

```
ans += tree[suff].num;
}
cout << ans << endl;</pre>
```

# **6** Data structures

# **6.1 BIT Range Queries**

```
struct BIT {
  int n;
 vector<int> slope;
  vector<int> intercept;
  // BIT can be thought of as having entries f[1], ...,
     f[n] which are 0-initialized
  BIT(int n): n(n), slope(n+1), intercept(n+1) {}
  // returns f[1] + \dots + f[idx-1]
  // precondition idx <= n+1
  int query(int idx) {
    int m = 0, b = 0;
    for (int i = idx-1; i > 0; i -= i\&-i) {
     m += slope[i];
     b += intercept[i];
   return m*idx + b;
 // adds amt to f[i] for i in [idx1, idx2)
 // precondition 1 <= idx1 <= idx2 <= n+1 (you can't
     update element 0)
 void update(int idx1, int idx2, int amt) {
    for (int i = idx1; i \le n; i += i\&-i) {
      slope[i] += amt;
      intercept[i] -= idx1*amt;
    for (int i = idx2; i \le n; i += i\&-i) {
      slope[i] -= amt;
      intercept[i] += idx2*amt;
} ;
```

# 6.2 Treaps

```
typedef struct node{
  int prior, size;
  int val; //value stored in the array
  int sum; //whatever info you want to maintain in
      segtree for each node
  int lazy; //whatever lazy update you want to do
  struct node *1,*r;
} node;
struct Treap {
  typedef node* pnode;
  int sz(pnode t) {
    return t?t->size:0;
}
```

```
void upd sz(pnode t) {
  if (t) t - size = sz(t - sl) + 1 + sz(t - sr);
void lazy(pnode t) {
  if(!t || !t->lazy)return;
  t->val+=t->lazy;//operation of lazy
  t \rightarrow sum + = t \rightarrow lazy * sz(t);
  if(t->1)t->1->lazy+=t->lazy;//propagate lazy
  if(t->r)t->r->lazy+=t->lazy;
  t \rightarrow lazy=0;
void reset(pnode t){
  if(t)t->sum = t->val;//no need to reset lazy coz
      when we call this lazy would itself be propagated
void combine(pnode& t,pnode l,pnode r){//combining two
    ranges of segtree
  if(!l || !r) return void(t = 1?l:r);
  t \rightarrow sum = 1 \rightarrow sum + r \rightarrow sum;
void operation(pnode t){//operation of segtree
  if(!t)return;
  reset(t); //reset the value of current node assuming
     it now represents a single element of the array
  lazy(t->1);lazy(t->r);//imp:propagate lazy before
     combining t->1,t->r;
  combine (t, t->1, t);
  combine (t, t, t->r);
void split(pnode t,pnode &1,pnode &r,int pos,int add
   =0) {
  if(!t)return void(l=r=NULL);
  lazy(t);
  int curr_pos = add + sz(t->1);
  if(curr_pos<=pos)//element at pos goes to left</pre>
     subtree(1)
    split(t->r,t->r,r,pos,curr_pos+1),l=t;
  else
    split(t->1,1,t->1,pos,add),r=t;
  upd sz(t);
  operation(t);
void merge(pnode &t,pnode 1,pnode r) { //1->leftarray,r
   ->rightarray,t->resulting array
  lazy(1); lazy(r);
  if(!l || !r) t = l?l:r;
  else if(l->prior>r->prior)merge(l->r,l->r,r),t=1;
          merge(r->1,1,r->1), t=r;
  upd sz(t);
  operation(t);
pnode init(int val){
  pnode ret = (pnode)malloc(sizeof(node));
  ret->prior=rand(); ret->size=1;
  ret->val=val;
  ret->sum=val; ret->lazy=0;
  return ret;
int range_query(pnode t,int 1,int r) {//[1,r]
  pnode L, mid, R;
```

```
split(t,L,mid,l-1);
split(mid,t,R,r-1);//note: r-1!!
int ans = t->sum;
merge(mid,L,t);
merge(t,mid,R);
return ans;
}
void range_update(pnode t,int l,int r,int val){//[l,r]
pnode L,mid,R;
split(t,L,mid,l-1);
split(mid,t,R,r-1);//note: r-1!!
t->lazy+=val; //lazy_update
merge(mid,L,t);
merge(t,mid,R);
}
};
```

#### 6.3 Link-Cut Tree

```
const int MXN = 100005, MEM = 100005;
struct Splay {
  static Splay nil, mem[MEM], *pmem;
  Splay *ch[2], *f;
  int val, rev, size;
  Splay (int val=-1): val(val), rev(0), size(1)
  \{f = ch[0] = ch[1] = &nil; \}
  bool isr()
  { return f->ch[0] != this && f->ch[1] != this; }
  int dir()
  { return f->ch[0] == this ? 0 : 1; }
  void setCh(Splay *c, int d){
    ch[d] = c;
    if (c != &nil) c->f = this;
    pull();
  void push() {
    if( !rev ) return;
    swap(ch[0], ch[1]);
    if (ch[0] != &nil) ch[0]->rev ^= 1;
    if (ch[1] != &nil) ch[1]->rev ^= 1;
    rev=0;
  void pull(){
    size = ch[0] -> size + ch[1] -> size + 1;
    if (ch[0] != &nil) ch[0]->f = this;
    if (ch[1] != &nil) ch[1] -> f = this;
} Splay::nil, Splay::mem[MEM], *Splay::pmem = Splay::mem
Splay *nil = &Splay::nil;
void rotate(Splay *x) {
  Splay *p = x -> f;
  int d = x - > dir();
  if (!p->isr()) p->f->setCh(x, p->dir());
  else x->f = p->f;
        p\rightarrow setCh(x\rightarrow ch[!d], d); x\rightarrow setCh(p, !d); p\rightarrow pull
            (); x->pull();
vector<Splay*> splayVec;
void splay (Splay *x) {
  splayVec.clear();
```

```
for (Splay *q=x;; q=q->f) {
    splayVec.push back(q);
    if (q->isr()) break;
  reverse(begin(splayVec), end(splayVec));
  for (auto it : splayVec) it->push();
  while (!x->isr())
    if (x->f->isr()) rotate(x);
    else if (x->dir()==x->f->dir()) rotate(x->f), rotate(
    else rotate(x), rotate(x);
int id(Splay *x) { return x - Splay::mem + 1; }
Splay* access(Splay *x){
  Splay *q = nil;
  for (;x!=nil;x=x->f) {splay(x); x->setCh(q, 1); q = x;
  return q;
void chroot(Splay *x){
  access(x); splay(x); x\rightarrow rev ^= 1; x\rightarrow push(); x\rightarrow pull()
void link(Splay *x, Splay *y){
  access(x); splay(x);
  chroot(y); x->setCh(y, 1);
void cut_p(Splay *y) {
  access(y); splay(y); y->push(); y->ch[0] = y->ch[0]->f
void cut(Splay *x, Splay *y){
  chroot(x); cut p(y);
Splay* get_root(Splay *x) {
  access(x); splay(x);
  for (; x \rightarrow ch[0] != nil; x = x \rightarrow ch[0]) x \rightarrow push();
  splay(x); return x;
bool conn(Splay *x, Splay *y)
  return (x = get\_root(x)) == (y = get\_root(y));
Splay* lca(Splay *x, Splay *y) {
  access(x); access(y); splay(x);
  if (x->f == nil) return x;
  else return x->f;
```

# 7 Miscellaneous

### 7.1 2-SAT

```
class Two_Sat {
  int N; // number of variables
  vector<int> val; // assignment of x is at val[2x] and
      -x at val[2x+1]
  vector<char> valid; // changes made at time i are kept
      iff valid[i]
  vector<vector<int> > G; // graph of implications G[x][
      i] = y means (x -> y)
```

```
Two_Sat(int N_{-}): N(N_{-}) { // create a formula over N_{-}
     variables (numbered 1 to N)
    G.resize(2*N);
  int add_variable() {
    G.emplace back(); G.emplace back();
    return N++;
private:
  // converts a signed variable index to its position in
      val[] and G[]
  int to_ind(int x) {
    return 2*(abs(x)-1) + (x<0);
  // Add a directed edge to the graph.
  // You most likely do not want to call this yourself!
  void add_edge(int a, int b) {
    G[to_ind(a)].push_back(to_ind(b));
  int time() {
    return valid.size()-1;
  bool dfs(int x) {
    if (valid[abs(val[x])]) return val[x]>0;
    val[x] = time();
    val[x^1] = -time();
    for(int e:G[x])
      if(!dfs(e))
        return false;
    return true;
public:
  // Add the or-clause: (a or b)
  void add_or(int a, int b) {
    add_edge(-a,b); add_edge(-b,a);
  // Add the implication: a -> b
  void add_implication(int a, int b) {
    add or (-a, b);
  // Add condition: x is true
  void add_true(int x) {
    add_or(x,x);
  // At most one with linear number of clauses
  template<typename T>
  void add_at_most_one(T vars) {
    if(vars.begin() == vars.end()) return;
    int last = *vars.begin(), cur = 0;
    for(int const&e:vars) {
      if(e == last) continue;
      if(cur == 0) cur = e;
      else {
        add_or(-cur, -e);
        int new_cur = add_variable();
        cur = add_implication(cur, new_cur);
        add_implication(e, new_cur);
        cur = hew_cur;
    if(cur != 0) add_or(-cur, -last);
```

```
bool solve() {
    val.assign(2*n, 0); valid.assign(1, 0);
    for (int i=0; i < val.size(); i+=2) {</pre>
      if(!valid[abs(val[i])]) {
        valid.push_back(1);
        if(!dfs(i))
          valid.back()=0;
          valid.push back(1);
          if(!dfs(i+1)) return false;
    return true;
};
```

## 7.2 Merge Insertion

```
// Sorting in O(n^2) time with near-optimal number of
   comparisons
// Number of comparisons used is: n lg n - 1.415 n
// The lower bound is: lg(n!) = n lg n - 1.443 n
// Binary search insertion sort would need: n log n - n
struct Merge Insertion Sort{
  template<typename F>
  static void apply_permutation(vector<int> const&p, F
     aet) {
    const int n = p.size();
    vector<int> q(n);
    for (int i=0; i < n; ++i) q[p[i]] = i;
    for (int i=0; i<n; ++i) {</pre>
      while (q[i] != i) {
        swap(get(i), get(q[i]));
        swap(q[i], q[q[i]]);
    }
  // ret.first is the sorted vector ret.second[i] is the
  // index of the i-th smallest element in the original
  // i.e. ret.second is the permutation that was applied
      to sort
  template<typename T, typename F>
 static pair<vector<T>, vector<int> > sort (vector<T> v, 7.3 DP Optimizations
      F comp) {
    const int n = v.size();
    if (n <= 1) return {move(v), {{0}}};
    vector<int> preperm(n);
    iota(preperm.begin(), preperm.end(), 0);
    const int M = n-n/2;
    for (int i=0; i< n/2; ++i) {
      if(comp(v[M+i], v[i])){
        swap(v[M+i], v[i]);
        swap(preperm[M+i], preperm[i]);
    auto ret=sort(vector<T>(v.begin(), v.begin()+n/2),
    apply_permutation(ret.second,[&preperm,M](int const&
       i) ->int&{return preperm[M+i];});
```

```
apply_permutation(ret.second,[&preperm](int const&i)
     ->int&{return preperm[i];});
  apply permutation(ret.second, [&v,M](int const&i)->T
     &{return v[M+i];});
  iota(ret.second.begin(), ret.second.end(), 0);
  // insert one element without comparisons
  ret.first.push_back(v.back());
  ret.second.push_back(n-1);
  // now insert the rest in blocks that optimize the
     binary search sizes
  for (int it=1, r=n-1, s=2; r>n/2; ++it, r-=s, s=(1<<it)
    for (int i=r-s; i<r; ++i) {</pre>
      if (i > = n/2) {
        int a = find(ret.second.begin(), ret.second.
           end(), i-M) - ret.second.begin();
        int b = ret.first.size();
        if(a==b) {
          assert (i==M-1);
          a = -1;
        while (a+1 < b) {
          const int m = a + (b-a)/2;
          if(comp(ret.first[m], v[i])) a = m;
          else b = m;
        ret.first.insert(ret.first.begin()+b, v[i]);
        ret.second.insert(ret.second.begin()+b, i);
  // compose permutations
 apply_permutation(ret.second, [&preperm](int const&i
     ) ->int&{return preperm[i];});
  ret.second.swap(preperm);
  return ret;
template<typename T>
static pair<vector<T>, vector<int> > sort(vector<T> v)
 return sort(move(v), std::less<T>{});
```

};

```
A[i][j]: The smallest k that gives optimal answer
Divide and Conquer:
  dp[i][j] = min(k < j) \{dp[i - 1][k] + C[k][j]\}
  O(kn^2) \rightarrow O(knlog(n))
  Conditions:
    A[i][j] \le A[i][j + 1] OR
    C[a][d] + C[b][c] >= C[a][c] + C[b][d] where a < b <
  Short Description:
    A[i][1] \le A[i][2] \le ... \le A[i][n]
Knuth Optimization:
  dp[i][j] = min(i < k < j) \{dp[i][k] + dp[k][j]\} + C[i][
     j]
```

```
O(n^3) -> O(n^2)
Conditions:
    A[i, j - 1] <= A[i, j] <= A[i + 1, j] OR
    C[a][d] + C[b][c] >= C[a][c] + C[b][d] AND
    C[b][c] <= C[a][d] where a <= b <= c <= d
Short Description:
    For dp[i][j], loop k from A[i][j - 1] to A[i + 1][j]</pre>
```

### 7.4 Convex Hull Trick (Dynamic)

```
struct Line {
  long long m, b;
  mutable function<const Line*()> succ;
  bool operator<(const Line& rhs) const{</pre>
    if(rhs.b!=-(111<<62)) return m>rhs.m; // < for max
    const Line* s = succ();
    if (!s) return 0;
    return b-s->b > (s->m-m)*rhs.m; // < for max
};
struct HullDynamic : public multiset<Line> {
  bool bad(iterator y) {
    auto z = next(y);
    if(y==begin()){
      if(z==end())return 0;
      return y->m == z->m && y->b >= z->b; // <= for max
    auto x = prev(y);
    if (z==end()) return y->m == x->m && y->b >= x->b;
        // <= for max
    return (x-b-y-b)*1.0*(z-m-y-m) >= (y-b-z)
       ->b) *1.0* (y->m - x->m);
  void insert_line(long long m, long long b) {
    auto y = insert({ m, b });
    y->succ = [=] {return next(y) ==end()? 0:&*next(y);};
    if(bad(y)) { erase(y); return; }
    while (\text{next}(y)!=\text{end}() \&\& \text{bad}(\text{next}(y))) = \text{rase}(\text{next}(y));
    while(y!=begin() && bad(prev(y)))erase(prev(y));
  long long eval(long long x) {
    auto 1 = *lower_bound((Line) \{x, -(111 << 62)\});
    return 1.m * x + 1.b;
};
```

# 7.5 Convex Hull Trick (Static)

```
struct ConvexHullTrick {
   typedef long long LL;
   vector<LL> M;
   vector<LL> B;
   vector<double> left;
   ConvexHullTrick() {}
   bool bad(LL m1, LL b1, LL m2, LL b2, LL m3, LL b3) {
      // Careful, this may overflow
      return (b3-b1)*(m1-m2) < (b2-b1)*(m1-m3);
   }
   // Add a new line to the structure, y = mx + b.
   // Lines must be added in decreasing order of slope.</pre>
```

```
void add(LL m, LL b) {
    while (M.size() >= 2 \&\& bad(M[M.size()-2], B[B.size))
       ()-2], M.back(), B.back(), m, b)) {
      M.pop_back(); B.pop_back(); left.pop_back();
    if (M.size() && M.back() == m) {
      if (B.back() > b) {
       M.pop_back(); B.pop_back(); left.pop_back();
        return:
    if (M.size() == 0) {
      left.push_back(-numeric_limits<double>::infinity()
         );
    } else {
      left.push back((double)(b - B.back())/(M.back() -
    M.push_back(m);
    B.push_back(b);
  // Get the minimum value of mx + b among all lines in
     the structure.
  // There must be at least one line.
  LL query(LL x) {
    int i = upper_bound(left.begin(), left.end(), x) -
       left.begin();
    return M[i-1] *x + B[i-1];
};
```

## 7.6 BigInt library

```
struct bignum {
 typedef unsigned int uint;
 vector<uint> digits;
 static const uint RADIX = 1000000000;
 bignum(): digits(1, 0) {}
 bignum(const bignum& x): digits(x.digits) {}
 bignum(unsigned long long x) {*this = x;}
 bignum(const char* x) {*this = x;}
 bignum(const string& s) {*this = s;}
 bignum& operator=(const bignum& y)
    {digits = y.digits; return *this;}
 bignum& operator=(unsigned long long x) {
    digits.assign(1, x%RADIX);
    if (x >= RADIX)
     digits.push_back(x/RADIX);
    return *this;
 bignum& operator=(const char* s) {
    int slen=strlen(s),i,l;
    digits.resize((slen+8)/9);
    for (l=0; slen>0; l++, slen-=9) {
     digits[1]=0;
     for (i=slen>9?slen-9:0; i<slen; i++)
       digits[1]=10*digits[1]+s[i]-'0';
   while (digits.size() > 1 && !digits.back()) digits.
       pop_back();
```

```
0);
   return *this;
bignum& operator=(const string& s)
                                                                                                      while (pwr >= 0) {
   {return *this = s.c str();}
                                                                                                          if (dividend.digits.size() > divisor.digits.size()
void add(const bignum& x) {
                                                                                                             unsigned long long q = dividend.digits.back();
   int l = max(digits.size(), x.digits.size());
                                                                                                             q *= RADIX; q += dividend.digits[dividend.digits
   digits.resize(1+1);
                                                                                                                   .size()-2];
   for (int d=0, carry=0; d<=1; d++) {</pre>
                                                                                                             q /= 1+divisor.digits.back();
      uint sum=carry;
                                                                                                             dividend -= divisor*q; digits[pwr] = q;
      if (d<digits.size()) sum+=digits[d];</pre>
                                                                                                             if (dividend >= divisor) { digits[pwr]++;
      if (d<x.digits.size()) sum+=x.digits[d];</pre>
                                                                                                                   dividend -= divisor; }
      digits[d]=sum;
                                                                                                             assert(dividend.digits.size() <= divisor.digits.</pre>
      if (digits[d]>=RADIX)
                                                                                                                   size()); continue;
          digits[d]-=RADIX, carry=1;
      else
                                                                                                          while (dividend.digits.size() == divisor.digits.
         carry=0;
                                                                                                                size()) {
                                                                                                             uint q = dividend.digits.back() / (1+divisor.
  if (!digits.back()) digits.pop_back();
                                                                                                                   digits.back());
                                                                                                             if (q == 0) break;
void sub(const bignum& x) {
                                                                                                             digits[pwr] += q; dividend -= divisor*q;
  // if ((*this) < x) throw; //negative numbers not yet
        supported
                                                                                                          if (dividend >= divisor) { dividend -= divisor;
   for (int d=0, borrow=0; d<digits.size(); d++) {</pre>
                                                                                                               digits[pwr]++; }
      digits[d]-=borrow;
                                                                                                          pwr--; divisor.digits.erase(divisor.digits.begin()
      if (d<x.digits.size()) digits[d]-=x.digits[d];</pre>
      if (digits[d]>>31) { digits[d]+=RADIX; borrow=1; }
             else borrow=0;
                                                                                                      while (digits.size() > 1 && !digits.back()) digits.
                                                                                                            pop_back();
   while (digits.size() > 1 && !digits.back()) digits.
                                                                                                      return dividend;
        pop_back();
                                                                                                   string to_string() const {
void mult(const bignum& x) {
                                                                                                       ostringstream oss;
   vector<uint> res(digits.size() + x.digits.size());
                                                                                                       oss << digits.back();</pre>
   unsigned long long y, z;
                                                                                                       for (int i = digits.size() - 2; i >= 0; i--) {
   for (int i=0; i < digits.size(); i++) {</pre>
                                                                                                          oss << setfill('0') << setw(9) << digits[i];
      for (int j=0; j<x.digits.size(); j++) {</pre>
         unsigned long long y=digits[i]; y*=x.digits[j];
                                                                                                      return oss.str();
         unsigned long long z=v/RADIX;
         res[i+j+1]+=z; res[i+j]+=y-RADIX*z; //mod is
                                                                                                   bignum operator+(const bignum& y) const
                                                                                                       {bignum res(*this); res.add(y); return res;}
         if (res[i+j] >= RADIX) \{ res[i+j] -= RADIX; res
                                                                                                   bignum operator-(const bignum& y) const
               i+j+1]++; }
                                                                                                       {bignum res(*this); res.sub(y); return res;}
         for (int k = i+j+1; res[k] >= RADIX; res[k] -=
                                                                                                   bignum operator*(const bignum& y) const
               RADIX, res[++k]++);
                                                                                                       {bignum res(*this); res.mult(y); return res;}
                                                                                                   bignum operator/(const bignum& y) const
                                                                                                       {bignum res(*this); res.div(y); return res;}
   digits = res;
                                                                                                   bignum operator% (const bignum& y) const
   while (digits.size() > 1 && !digits.back()) digits.
                                                                                                       {bignum res(*this); return res.div(y);}
        pop back();
                                                                                                   bignum& operator+=(const bignum& y)
// returns the remainder
                                                                                                       {add(y); return *this;}
                                                                                                   bignum& operator-=(const bignum& y)
bignum div(const bignum& x) {
                                                                                                       {sub(y); return *this;}
   bignum dividend(*this);
                                                                                                   bignum& operator*=(const bignum& y)
   bignum divisor(x);
                                                                                                       {mult(y); return *this;}
   fill(digits.begin(), digits.end(), 0);
                                                                                                   bignum& operator/=(const bignum& y)
   // shift divisor up
                                                                                                       {div(y); return *this;}
   int pwr = dividend.digits.size() - divisor.digits.
                                                                                                   bignum& operator%=(const bignum& y)
        size();
                                                                                                       \{*this = div(y);\}
   if (pwr > 0)
      divisor.digits.insert(divisor.digits.begin(), pwr,
                                                                                                   bool operator==(const bignum& y)
```

```
{return digits == v.digits;}
  bool operator<(const bignum& y) const {</pre>
    if (digits.size() < y.digits.size()) return true;</pre>
    if (digits.size() > y.digits.size()) return false;
    for (int i = digits.size()-1; i >= 0; i--)
      if (digits[i] < y.digits[i])</pre>
        return true;
      else if (digits[i] > y.digits[i])
        return false:
    return false;
  bool operator>(const bignum& y) const
    {return y<*this;}
  bool operator<=(const bignum& y) const</pre>
    {return ! (y<*this);}
  bool operator>=(const bignum& y) const
    {return ! (*this<v);}
};
```

### 7.7 Manachers algorithm

```
// Ex: "opposes" -> [0,1,0,1,4,1,0,1,0,1,0,3,0,1,0]
vector<int> fastLongestPalindromes(string str) {
  int i=0, j, d, s, e, lLen, palLen=0;
  vector<int> res;
  while (i < str.length()) {</pre>
    if (i > palLen && str[i-palLen-1] == str[i]) {
      palLen += 2; i++; continue;
    res.push back(palLen);
    s = res.size()-2;
    e = s-palLen;
   bool b = true;
    for (j=s; j>e; j--) {
      d = j - e - 1;
      if (res[j] == d) { palLen = d; b = false; break; }
      res.push_back(min(d, res[j]));
    if (b) { palLen = 1; i++; }
  res.push_back(palLen);
  lLen = res.size();
  s = 1Len-2;
  e = s-(2*str.length()+1-lLen);
  for (i=s; i>e; i--) { d = i-e-1; res.push\_back(min(d,
     res[i])); }
  return res;
```

### 7.8 Dates

```
// Months are expressed as integers from 1 to 12, Days
    are expressed as integers from 1 to 31, and Years are
    expressed as 4-digit integers.
string dayOfWeek[] = {"Mon", "Tue", "Wed", "Thu", "Fri",
    "Sat", "Sun"};
//converts Gregorian date to integer(Julian day number)
int dateToInt (int m, int d, int y) {
    return
    1461 * (y + 4800 + (m - 14) / 12) / 4 +
```

```
367 * (m - 2 - (m - 14) / 12 * 12) / 12 -
    3 * ((y + 4900 + (m - 14) / 12) / 100) / 4 +
    d - 32075;
// converts integer (Julian day number) to Gregorian
   date: month/day/year
void intToDate (int jd, int &m, int &d, int &y) {
  int x, n, i, j;
  x = jd + 68569;
 n = 4 * x / 146097;
 x = (146097 * n + 3) / 4;
 i = (4000 * (x + 1)) / 1461001;
 x = 1461 * i / 4 - 31;
 j = 80 * x / 2447;
  d = x - 2447 * j / 80;
 x = \frac{1}{2} / 11;
 m = 1 + 2 - 12 * x;
  y = 100 * (n - 49) + i + x;
// converts integer (Julian day number) to day of week
string intToDay (int jd) {
  return dayOfWeek[jd % 7];
```

### 7.9 Bitset (Text)

Remember \_Find\_first() and \_Find\_next() ex- ist, and run in O(N/W), where W is word size of machine.

## 7.10 Template

```
q++ -std=c++17 -DLOCAL -O2 -Wall -Wshadow -Wextra -
   pedantic -Wfloat-equal -Wlogical-op
#pragma comment(linker, "/stack:200000000")
#pragma GCC optimize("Ofast")
#pragma GCC optimize ("unroll-loops")
#pragma GCC target("sse,sse2,sse3,ssse3,sse4,popcnt,abm,
   mmx, avx, tune=native") // codeforces
//#pragma GCC target("avx,avx2,fma")
//#pragma GCC target("sse,sse2,sse3,ssse3,sse4,popcnt,
   abm, mmx, tune=native") // yandex
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
#include <bits/stdc++.h>
using namespace __gnu_pbds;
using namespace std;
typedef long double ld;
typedef long long 11;
typedef pair <int, int> pii;
typedef pair <ll, ll> pll;
mt19937 rng(std::chrono::duration_cast<std::chrono::</pre>
   nanoseconds>(chrono::high_resolution_clock::now().
   time since epoch()).count());
template<typename has_less>
using ordered set =
tree<has_less,
```

```
null type,
        less<has_less>,
        rb tree tag,
        tree order statistics node update>;
//insert using pref_trie.insert
//get range for prefix using pref_trie.prefix_range
//use iterator from range.first until != range.second
typedef
trie<string,
       null type,
       trie_string_access_traits<>,
       pat_trie_tag,
        trie_prefix_search_node_update>
       pref_trie;
struct chash {
    int operator()(int x) const { return x ^ RANDOM; }
gp_hash_table<key, int, chash> table;
int main(){}
```

## 7.11 Numerical Integration

```
// different schemes for numerical integration
// approximatively ordered by accuracy
// do NOT use integer types for integration range!
struct Integration_Midpoint{
  template<typename Func, typename S>
  static typename result_of<Func(S)>::type
     integrate_step(Func f, S l, S r) {
   S m = (1+r)/2;
   return f(m) * (r-1);
};
struct Integration_Simpson{
 template<typename Func, typename S>
 static typename result_of<Func(S)>::type
     integrate_step(Func f, S l, S r) {
   S m = (1+r)/2;
   return (f(1) + 4*f(m) + f(r))/6 * (r-1);
struct Integration_Gauss_2{
  static constexpr long double A = 1.01/sqrt1(3)/2, x1
     =0.51-A, x2 = 0.51+A;
 template<typename Func, typename S>
  static typename result_of<Func(S)>::type
     integrate_step(Func f, S l, S r) {
   return (f(1*x1 + r*x2) + f(1*x2+r*x1))/2 * (r-1);
struct Integration_NCotes_Open_4{
 template<typename Func, typename S>
  static typename result_of<Func(S)>::type
     integrate_step(Func f, S l, S r) {
    S h = (r-1)/5;
   return (11*f(1+h) + f(1+2*h) + f(r-2*h) + 11*f(r-h))
       /24 * (r-1);
};
```

```
struct Integration_Gauss_3{
  static constexpr long double A = sqrt1(3.01/5.01)/2,
     x1=0.5-A, x^2 = 0.5+A;
  template<typename Func, typename S>
  static typename result_of<Func(S)>::type
     integrate_step(Func f, S l, S r){
    return (5*f(1*x1 + r*x2) + 8*f((1+r)/2) + 5*f(1*x2+r)
       *x1))/18 * (r-1);
};
template<typename Integration Method>
struct Integrator Fixedstep{
  template<typename Func, typename S>
  static typename result_of<Func(S)>::type integrate(
     Func f, S const 1, S const r, size_t const steps) {
    assert(steps>0);
    typename result_of<Func(S)>::type ret(0);
    S cur_1 = 1, cur_r;
    for(size_t i=0;i<steps;++i){</pre>
      cur_r = (1*(steps-i-1) + r*(i+1))/steps;
      ret+=Integration Method::integrate step(f, cur l,
          cur_r);
      cur_l = cur_r;
    return ret;
};
template<typename Integration_Method>
class Integrator_Adaptive{
private:
  template < size_t depth_limit, typename Func, typename S
  static typename result_of<Func(S)>::type integrate(
     Func f, S const l, S const r, typename result_of<</pre>
     Func(S)>::type const val, typename result of<Func(S
     )>::type const eps, const size_t depth) {
    if (depth>=depth limit) {
      return val;
    S const m = (1+r)/2;
    typename result_of<Func(S)>::type val_l =
       Integration_Method::integrate_step(f, l, m);
    typename result_of<Func(S)>::type val_r =
       Integration_Method::integrate_step(f, m, r);
    typename result_of<Func(S)>::type error = abs(val -
       val 1 - val r);
    if(error < eps) {</pre>
      return val 1 + val r;
    return integrate<depth_limit>(f, 1, m, val_1, eps/2,
      + integrate < depth_limit > (f, m, r, val_r, eps/2,
         depth+1);
public:
  template<size_t depth_limit, typename Func, typename S
  static typename result of<Func(S)>::type integrate(
     Func f, S const l, S const r, typename result_of<</pre>
     Func(S)>::type const eps) {
    return integrate<depth_limit>(f, l, r,
```