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// Reference: My (forked) ICPC codebook: https://github com/aayush9/ICPC-CodeBook/blob/master/code/Dinic.cc	
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```
Edge(int u, int v, long long cap): u(u), v(v), cap(
     cap), flow(0) {}
};
int N;
vector<Edge> E;
vector<vector<int>> q;
vector<int> d, pt;
Dinic(int N): N(N), E(0), g(N), d(N), pt(N) {}
void AddEdge(int u, int v, long long cap) {
  if (u != v) {
    E.emplace_back(Edge(u, v, cap));
    g[u].emplace_back(E.size() - 1);
    E.emplace_back(Edge(v, u, 0));
    g[v].emplace_back(E.size() - 1);
bool BFS(int S, int T) {
  queue<int> q({S});
  fill(d.begin(), d.end(), N + 1);
  d[S] = 0;
  while(!q.empty()) {
    int u = q.front(); q.pop();
    if (u == T) break;
    for (int k: g[u]) {
      Edge &e = \tilde{E}[k];
      if (e.flow < e.cap && d[e.v] > d[e.u] + 1) {
   d[e.v] = d[e.u] + 1;
        q.emplace(e.v);
  return d[T] != N + 1;
long long DFS (int u, int T, long long flow = -1) {
  if (u == T || flow == 0) return flow;
  for (int &i = pt[u]; i < g[u].size(); ++i) {</pre>
    Edge &e = E[g[u][i]];
    Edge &oe = E[g[u][i]^1];
    if (d[e.v] == d[e.u] + 1) {
      long long amt = e.cap - e.flow;
      if (flow' != -1 \&\& amt > flow) amt = flow;
      if (long long pushed = DFS(e.v, T, amt)) {
        e.flow += pushed;
        oe.flow -= pushed;
        return pushed;
    }
  return 0;
long long MaxFlow(int S, int T) {
  long long total = 0;
 while (BFS(S, T)) {
    fill(pt.begin(), pt.end(), 0);
    while (long long flow = DFS(S, T))
      total += flow;
  return total;
```

} **;**

1.2 Min-cost max-flow

```
// Implementation of min cost max flow algorithm using
   adjacency
// matrix (Edmonds and Karp 1972). This implementation
   keeps track of
// forward and reverse edges separately (so you can set
   cap[i][j] !=
// cap[j][i]). For a regular max flow, set all edge
   costs to 0.
// Running time, O(|V|^2) cost per augmentation
      max flow:
                         O(|V|^3) augmentations
      min cost max flow: O(|V|^4 * MAX EDGE COST)
   augmentations
// INPUT:
      - graph, constructed using AddEdge()
       - source
       - sink
// OUTPUT:
      - (maximum flow value, minimum cost value)
      - To obtain the actual flow, look at positive
   values only.
#include <cmath>
#include <vector>
#include <iostream>
using namespace std;
typedef vector<int> VI;
typedef vector<VI> VVI;
typedef long long L;
typedef vector<L> VL;
typedef vector<VL> VVL;
typedef pair<int, int> PII;
typedef vector<PII> VPII;
const L INF = numeric limits<L>::max() / 4;
struct MinCostMaxFlow {
  int N;
  VVL cap, flow, cost;
  VI found;
  VL dist, pi, width;
  VPII dad;
  MinCostMaxFlow(int N) :
   N(N), cap(N, VL(N)), flow(N, VL(N)), cost(N, VL(N)),
    found(N), dist(N), pi(N), width(N), dad(N) {}
  void AddEdge(int from, int to, L cap, L cost) {
   this->cap[from][to] = cap;
    this->cost[from][to] = cost;
  void Relax(int s, int k, L cap, L cost, int dir) {
   L val = dist[s] + pi[s] - pi[k] + cost;
```

```
if (cap && val < dist[k]) {
      dist[k] = val;
      dad[k] = make_pair(s, dir);
      width[k] = min(cap, width[s]);
  L Dijkstra(int s, int t) {
    fill(found.begin(), found.end(), false);
    fill(dist.begin(), dist.end(), INF);
    fill(width.begin(), width.end(), 0);
    dist[s] = 0;
    width[s] = INF;
    while (s != -1) {
      int best = -1;
      found(s) = true;
      for (int k = 0; k < N; k++) {
        if (found[k]) continue;
        Relax(s, k, cap[s][k] - flow[s][k], cost[s][k],
        Relax(s, k, flow[k][s], -\cos t[k][s], -1);
        if (best == -1 \mid \mid dist[k] < dist[best]) best = k
      s = best;
    for (int k = 0; k < N; k++)
     pi[k] = min(pi[k] + dist[k], INF);
    return width[t];
  pair<L, L> GetMaxFlow(int s, int t) {
    L \text{ totflow} = 0, totcost = 0;
    while (L amt = Dijkstra(s, t)) {
      totflow += amt;
      for (int x = t; x != s; x = dad[x].first) {
        if (dad[x].second == 1) {
          flow[dad[x].first][x] += amt;
          totcost += amt * cost[dad[x].first][x];
        } else {
          flow[x][dad[x].first] -= amt;
          totcost -= amt * cost[x][dad[x].first];
    return make_pair(totflow, totcost);
};
// BEGIN CUT
// The following code solves UVA problem #10594: Data
   Flow
int main() {
  int N, M;
  while (scanf("%d%d", &N, &M) == 2) {
    VVL v(M, VL(3));
    for (int i = 0; i < M; i++)
      scanf("%Ld%Ld%Ld", &v[i][0], &v[i][1], &v[i][2]);
    scanf("%Ld%Ld", &D, &K);
```

1.3 Edmonds Max Matching

```
/*
GETS:
V->number of vertices
E->number of edges
pair of vertices as edges (vertices are 1..V)
GIVES:
output of edmonds() is the maximum matching
match[i] is matched pair of i (-1 if there isn't a
   matched pair)
#include <bits/stdc++.h>
using namespace std;
const int M=505;
struct struct_edge{int v;struct_edge* n;};
typedef struct edge* edge;
struct_edge pool[M*M*2];
edge top=pool,adj[M];
int V,E,match[M],qh,qt,q[M],father[M],base[M];
bool inq[M],inb[M],ed[M][M];
void add edge(int u,int v)
  top->v=v, top->n=adj[u], adj[u]=top++;
  top \rightarrow v = u, top \rightarrow n = adj[v], adj[v] = top + +;
int LCA(int root,int u,int v)
  static bool inp[M];
  memset(inp,0,sizeof(inp));
  while (1)
      inp[u=base[u]]=true;
      if (u==root) break;
      u=father[match[u]];
  while (1)
```

```
if (inp[v=base[v]]) return v;
      else v=father[match[v]];
void mark blossom(int lca,int u)
  while (base[u]!=lca)
      int v=match[u];
      inb[base[u]]=inb[base[v]]=true;
      u=father[v];
      if (base[u]!=lca) father[u]=v;
void blossom contraction(int s,int u,int v)
  int lca=LCA(s,u,v);
  memset(inb, 0, sizeof(inb));
  mark blossom(lca,u);
  mark_blossom(lca, v);
  if (base[u]!=lca)
    father[u]=v;
  if (base[v]!=lca)
    father[v]=u;
  for (int u=0; u < V; u++)
    if (inb[base[u]])
  base[u]=lca;
  if (!inq[u])
    inq[q[++qt]=u]=true;
int find_augmenting_path(int s)
  memset(inq,0,sizeof(inq));
  memset (father, -1, sizeof (father));
  for (int i=0;i<V;i++) base[i]=i;</pre>
  inq[q[qh=qt=0]=s]=true;
  while (qh<=qt)</pre>
      int u=q[qh++];
      for (edge e=adj[u];e;e=e->n)
    int v=e->v;
    if (base[u]!=base[v]&&match[u]!=v)
      if ((v==s)||(match[v]!=-1 && father[match[v]]!=-1)
        blossom contraction(s,u,v);
      else if (father[v]==-1)
    father[v]=u;
    if (match[v] == -1)
      return v;
    else if (!inq[match[v]])
      inq[q[++qt]=match[v]]=true;
  return -1;
```

```
int augment_path(int s,int t)
  int u=t,v,w;
  while (u!=-1)
      v=father[u];
      w=match[v];
      match[v]=u;
      match[u]=v;
  return t!=-1;
int edmonds()
  int matchc=0;
  memset (match, -1, sizeof (match));
  for (int u=0; u<V; u++)
    if (match[u] == -1)
      matchc+=augment_path(u, find_augmenting_path(u));
  return matchc;
int main()
  int u, v;
  cin>>V>>E;
  while (E^{--})
      cin>>u>>v;
      if (!ed[u-1][v-1])
    add_edge (u-1, v-1);
    ed[u-1][v-1]=ed[v-1][u-1]=true;
  cout << edmonds() << endl;</pre>
  for (int i=0;i<V;i++)</pre>
    if (i<match[i])</pre>
      cout << i+1 << " " << match[i] +1 << endl;
```

1.4 Max bipartite matching

```
// This code performs maximum bipartite matching.
//
Running time: O(|E| |V|) -- often much faster in practice
//
INPUT: w[i][j] = edge between row node i and column node j
// OUTPUT: mr[i] = assignment for row node i, -1 if unassigned
// mc[j] = assignment for column node j, -1 if unassigned
// function returns number of matches made
#include <vector>
using namespace std;
typedef vector<int> VI;
```

```
typedef vector<VI> VVI;
bool FindMatch (int i, const VVI &w, VI &mr, VI &mc, VI &
  for (int j = 0; j < w[i].size(); j++) {</pre>
    if (w[i][j] && !seen[j]) {
      seen[j] = true;
      if (mc[j] < 0 \mid | FindMatch(mc[j], w, mr, mc, seen)
        mr[i] = j;
        mc[j] = i;
        return true;
  return false;
int BipartiteMatching(const VVI &w, VI &mr, VI &mc) {
  mr = VI(w.size(), -1);
  mc = VI(w[0].size(), -1);
  int ct = 0;
  for (int i = 0; i < w.size(); i++) {</pre>
    VI seen(w[0].size());
    if (FindMatch(i, w, mr, mc, seen)) ct++;
  return ct;
```

1.5 Global min-cut

```
// Adjacency matrix implementation of Stoer-Wagner min
   cut algorithm.
// Running time:
       O(|V|^3)
// INPUT:
       - graph, constructed using AddEdge()
// OUTPUT:
       - (min cut value, nodes in half of min cut)
#include <cmath>
#include <vector>
#include <iostream>
using namespace std;
typedef vector<int> VI;
typedef vector<VI> VVI;
const int INF = 1000000000;
pair<int, VI> GetMinCut(VVI &weights) {
  int N = weights.size();
  VI used(N), cut, best_cut;
  int best_weight = -1;
  for (int phase = N-1; phase >= 0; phase--) {
    VI w = weights[0];
    VI added = used;
    int prev, last = 0;
```

```
for (int i = 0; i < phase; i++) {
      prev = last;
      last = -1;
      for (int j = 1; j < N; j++)
        if (!added[j] && (last == -1 || w[j] > w[last]))
            last = j;
      if (i == phase-1) {
        for (int j = 0; j < N; j++) weights[prev][j] +=
           weights[last][j];
        for (int j = 0; j < N; j++) weights[j][prev] =
           weights[prev][j];
        used[last] = true;
        cut.push_back(last);
        if (best_weight == -1 || w[last] < best_weight)</pre>
          best_cut = cut;
          best_weight = w[last];
      } else {
        for (int j = 0; j < N; j++)
          w[j] += weights[last][j];
        added[last] = true;
  return make_pair(best_weight, best_cut);
// BEGIN CUT
// The following code solves UVA problem #10989: Bomb,
   Divide and Conquer
int main() {
  int N;
  cin >> N;
  for (int i = 0; i < N; i++) {
    int n, m;
    cin >> n >> m;
    VVI weights(n, VI(n));
    for (int j = 0; j < m; j++) {
      int a, b, c;
      cin >> a >> b >> c;
      weights [a-1][b-1] = weights [b-1][a-1] = c;
    pair<int, VI> res = GetMinCut(weights);
    cout << "Case #" << i+1 << ": " << res.first << endl
// END CUT
```

2 Geometry

2.1 Convex hull

```
typedef pair<long long, long long> PT;
long double dist(PT a, PT b) {
  return sqrt(pow(a.first-b.first,2)+pow(a.second-b.second,2));
}
```

```
long long cross(PT o, PT a, PT b) {
  PT OA = {a.first-o.first,a.second-o.second};
  PT OB = {b.first-o.first,b.second-o.second};
  return OA.first*OB.second - OA.second*OB.first;
vector<PT> convexhull(){
  vector<PT> hull;
  sort(a,a+n,[](PT i, PT j){
    if(i.second!=j.second)
      return i.second < j.second;</pre>
    return i.first < j.first;</pre>
  for(int i=0; i<n; ++i) {
    while(hull.size()>1 && cross(hull[hull.size()-2],
       hull.back(),a[i]) <= 0
      hull.pop back();
    hull.push_back(a[i]);
  for (int i=n-1, siz = hull.size(); i--;) {
    while(hull.size()>siz && cross(hull[hull.size()-2],
       hull.back(),a[i])<=0)
      hull.pop_back();
    hull.push_back(a[i]);
  return hull;
```

2.2 Miscellaneous geometry

```
double INF = 1e100, EPS = 1e-12;
struct PT {
  double x, y;
  PT() {}
  PT (double x, double y) : x(x), y(y) {}
  PT(const PT &p) : x(p.x), y(p.y)
                                      {}
  PT operator + (const PT &p) const { return PT(x+p.x,
     \sqrt{p.y}; }
  PT operator - (const PT &p) const { return PT(x-p.x,
     y-p.y);
  PT operator * (double c)
                               const { return PT(x*c,
     y*c ); }
  PT operator / (double c)
                               const { return PT(x/c,
     y/c ); }
};
double dot (PT p, PT q)
                           { return p.x*q.x+p.y*q.y; }
double dist2(PT p, PT q)
                         { return dot(p-q,p-q); }
double cross(PT p, PT q) { return p.x*q.y-p.y*q.x; }
ostream & operator << (ostream & os, const PT &p) {
 os << "(" << p.x << "," << p.y << ")";
// rotate a point CCW or CW around the origin
PT RotateCCW90 (PT p)
                     { return PT(-p.y,p.x);
PT RotateCW90(PT p)
                       { return PT(p.y,-p.x); }
PT RotateCCW(PT p, double t) {
  return PT(p.x*cos(t)-p.y*sin(t), p.x*sin(t)+p.y*cos(t)
// project point c onto line through a and b
```

```
// assuming a != b
PT ProjectPointLine(PT a, PT b, PT c) {
     return a + (b-a)*dot(c-a, b-a)/dot(b-a, b-a);
 // project point c onto line segment through a and b
PT ProjectPointSegment (PT a, PT b, PT c) {
     double r = dot(b-a,b-a);
     if (fabs(r) < EPS) return a;</pre>
     r = dot(c-a, b-a)/r;
     if (r < 0) return a;</pre>
     if (r > 1) return b;
     return a + (b-a) *r;
// compute distance from c to segment between a and b
double DistancePointSegment(PT a, PT b, PT c) {
     return sqrt(dist2(c, ProjectPointSegment(a, b, c)));
// compute distance between point (x,y,z) and plane ax+
       bv+cz=d
double DistancePointPlane (double x, double y, double z,
                                                             double a, double b, double c,
                                                                     double d)
     return fabs (a*x+b*y+c*z-d)/sqrt (a*a+b*b+c*c);
// determine if lines from a to b and c to d are
        parallel or collinear
bool LinesParallel (PT a, PT b, PT c, PT d) {
     return fabs(cross(b-a, c-d)) < EPS;
bool LinesCollinear(PT a, PT b, PT c, PT d) {
     return LinesParallel(a, b, c, d)
              && fabs(cross(a-b, a-c)) < EPS
              && fabs(cross(c-d, c-a)) < EPS;
// determine if line segment from a to b intersects with
// line segment from c to d
bool SegmentsIntersect(PT a, PT b, PT c, PT d) {
     if (LinesCollinear(a, b, c, d)) {
         if (dist2(a, c) < EPS || dist2(a, d) < EPS ||
              dist2(b, c) < EPS || dist2(b, d) < EPS) return
         if (dot(c-a, c-b) > 0 \&\& dot(d-a, d-b) > 0 \&\& dot(c-a) > 0 \&\& dot(c-a) > 0 &\& dot(c-a) > 0
                 b, d-b) > 0
              return false;
         return true;
     if (cross(d-a, b-a) * cross(c-a, b-a) > 0) return
            false:
     if (cross(a-c, d-c) * cross(b-c, d-c) > 0) return
             false:
     return true;
// compute intersection of line passing through a and b
// with line passing through c and d, assuming that
// intersection exists; for segment intersection, check
```

```
if
// segments intersect first
PT ComputeLineIntersection(PT a, PT b, PT c, PT d) {
  b=b-a; d=c-d; c=c-a;
  assert (dot (b, b) > EPS && dot (d, d) > EPS);
  return a + b*cross(c, d)/cross(b, d);
// compute center of circle given three points
PT ComputeCircleCenter(PT a, PT b, PT c) {
  b = (a + b) / 2;
  c = (a + c) / 2;
  return ComputeLineIntersection(b, b+RotateCW90(a-b), c
      , c+RotateCW90(a-c);
// determine if point is in a possibly non-convex
   polygon (by William
// Randolph Franklin); returns 1 for strictly interior
   points, 0 for
// strictly exterior points, and 0 or 1 for the
   remaining points.
// Note that it is possible to convert this into an \star
   exact* test using
// integer arithmetic by taking care of the division
   appropriately
// (making sure to deal with signs properly) and then by
    writing exact
// tests for checking point on polygon boundary
bool PointInPolygon(const vector<PT> &p, PT q) {
  bool c = 0;
  for (int i = 0; i < p.size(); i++) {</pre>
    int j = (i+1) p.size();
    if ((p[i].y <= q.y && q.y < p[j].y ||</pre>
      p[j].y \le q.y \&\& q.y < p[i].y) \&\&
      q.x < p[i].x + (p[j].x - p[i].x) * (q.y - p[i].y)
         /(p[j].y - p[i].y))
      c = !c;
  return c;
// determine if point is on the boundary of a polygon
bool PointOnPolygon(const vector<PT> &p, PT q) {
  for (int i = 0; i < p.size(); i++)</pre>
    if (dist2(ProjectPointSegment(p[i], p[(i+1)%p.size()
      ], q), q) < EPS) return true;
    return false;
// compute intersection of line through points a and b
   with
// circle centered at c with radius r > 0
vector<PT> CircleLineIntersection(PT a, PT b, PT c,
   double r) {
  vector<PT> ret;
  b = b-a;
  a = a-c;
  double A = dot(b, b);
  double B = dot(a, b);
  double C = dot(a, a) - r*r;
  double D = B*B - A*C;
  if (D < -EPS) return ret;</pre>
```

```
ret.push_back(c+a+b*(-B+sqrt(D+EPS))/A);
  if (D > EPS)
    ret.push_back(c+a+b*(-B-sqrt(D))/A);
  return ret;
// compute intersection of circle centered at a with
// with circle centered at b with radius R
vector<PT> CircleCircleIntersection(PT a, PT b, double r
   , double R) {
  vector<PT> ret;
  double d = sqrt(dist2(a, b));
  if (d > r+R || d+min(r, R) < max(r, R)) return ret;</pre>
  double x = (d*d-R*R+r*r)/(2*d);
  double y = sqrt(r*r-x*x);
  PT v = (b-a)/d;
  ret.push back(a+v*x + RotateCCW90(v)*y);
  if (v > 0)
   ret.push_back(a+v*x - RotateCCW90(v)*y);
  return ret;
// This code computes the area or centroid of a (
   possibly nonconvex)
// polygon, assuming that the coordinates are listed in
   a clockwise or
// counterclockwise fashion. Note that the centroid is
   often known as
// the "center of gravity" or "center of mass".
double ComputeSignedArea(const vector<PT> &p) {
  double area = 0;
  for(int i = 0; i < p.size(); i++) {</pre>
    int j = (i+1) % p.size();
    area += p[i].x*p[j].y - p[j].x*p[i].y;
  return area / 2.0;
double ComputeArea(const vector<PT> &p) {
  return fabs(ComputeSignedArea(p));
PT ComputeCentroid(const vector<PT> &p) {
  PT c(0,0);
  double scale = 6.0 * ComputeSignedArea(p);
  for (int i = 0; i < p.size(); i++) {
    int j = (i+1) % p.size();
    c = c + (p[i]+p[j])*(p[i].x*p[j].y - p[j].x*p[i].y);
  return c / scale;
// tests whether or not a given polygon (in CW or CCW
   order) is simple
bool IsSimple(const vector<PT> &p) {
  for (int i = 0; i < p.size(); i++) {</pre>
    for (int k = i+1; k < p.size(); k++) {</pre>
      int j = (i+1) % p.size();
      int 1 = (k+1) % p.size();
      if (i == 1 \mid | j == k) continue;
      if (SegmentsIntersect(p[i], p[j], p[k], p[l]))
        return false;
```

```
return true;
int main() {
  cerr << RotateCCW90(PT(2,5)) << endl;</pre>
  cerr << RotateCCW(PT(2,5),M_PI/2) << endl;</pre>
  cerr << ProjectPointLine(PT(-5,-2), PT(10,4), PT(3,7))
       << endl:
  cerr << ProjectPointSegment(PT(-5,-2), PT(10,4), PT</pre>
      (3,7));
  cerr << DistancePointPlane(4, -4, 3, 2, -2, 5, -8) <math><< endl;
  cerr << LinesParallel(PT(1,1), PT(3,5), PT(2,1), PT
      (4,5));
  cerr << LinesCollinear(PT(1,1), PT(3,5), PT(2,1), PT</pre>
      (4,5));
  cerr << SegmentsIntersect(PT(0,0), PT(2,4), PT(3,1),
     PT(-1,3));
  cerr << ComputeLineIntersection(PT(0,0), PT(2,4), PT</pre>
      (3,1), PT(-1,3)) << endl;
  cerr << ComputeCircleCenter(PT(-3,4), PT(6,1), PT(4,5)</pre>
     ) << endl;
  vector<PT> v;
  v.push back (PT(0,0));
  v.push back (PT(5,0));
  v.push_back(PT(5,5));
  v.push_back(PT(0,5));
  cerr << PointInPolygon(v, PT(2,2));</pre>
  cerr << PointOnPolygon(v, PT(2,2))</pre>
  vector<PT> u = CircleLineIntersection(PT(0,6), PT(2,6)
     , PT(1,1), 5);
  for (int i = 0; i < u.size(); i++) cerr << u[i] << " "</pre>
     ; cerr << endl;
  u = CircleCircleIntersection(PT(1,1), PT(4.5,4.5), 5,
     sart(2.0)/2.0);
  for (int i = 0; i < u.size(); i++) cerr << u[i] << " "</pre>
     ; cerr << endl;</pre>
  PT pa[] = { PT(0,0), PT(5,0), PT(1,1), PT(0,5) };
  vector<PT> p(pa, pa+4);
  PT c = ComputeCentroid(p);
  cerr << "Area: " << ComputeArea(p) << endl;</pre>
  cerr << "Centroid: " << c << endl;</pre>
```

2.3 3D geometry

```
double a, double b, double c, double d) {
  return Math.abs(a*x + b*y + c*z + d) / Math.sqrt(a*a
      + b*b + c*c);
// distance between parallel planes ax + by + cz + d1 2.4 Slow Delaunay triangulation
   = 0 and
// aX + bY + cZ + d2 = 0
public static double planePlaneDist(double a, double b
   , double c,
    double d1, double d2) {
  return Math.abs(d1 - d2) / Math.sqrt(a*a + b*b + c*c
// distance from point (px, py, pz) to line (x1, y1,
   z1)-(x2, y2, z2)
// (or ray, or segment; in the case of the ray, the
   endpoint is the
// first point)
public static final int LINE = 0;
public static final int SEGMENT = 1;
public static final int RAY = 2;
public static double ptLineDistSq(double x1, double y1
    double z1,
    double x2, double y2, double z2, double px, double
        py, double pz,
    int type) {
  double pd2 = (x1-x2)*(x1-x2) + (y1-y2)*(y1-y2) + (z1
     -z2)*(z1-z2);
  double x, y, z;
  if (pd2 == 0) {
   x = x1;
   y = y1;
   z = z1;
  } else {
    double u = ((px-x1)*(x2-x1) + (py-y1)*(y2-y1) + (
       pz-z1)*(z2-z1)) / pd2;
   x = x1 + u * (x2 - x1);
   y = y1 + u * (y2 - y1);
    z = z1 + u * (z2 - z1);
    if (type != LINE && u < 0) {
     x = x1;
     y = y1;
      z = z1;
    if (type == SEGMENT && u > 1.0) {
     x = x2;
     y = y2;
      z = z2;
  return (x-px) * (x-px) + (y-py) * (y-py) + (z-pz) * (z-pz)
public static double ptLineDist(double x1, double y1,
   double z1,
   double x2, double y2, double z2, double px, double
        py, double pz,
    int type) {
  return Math.sqrt(ptLineDistSq(x1, y1, z1, x2, y2, z2
```

, px, py, pz, type));

```
// Slow but simple Delaunay triangulation. Does not
// degenerate cases (from O'Rourke, Computational
   Geometry in C)
// Running time: O(n^4)
// INPUT:
             x[] = x-coordinates
             v[] = v-coordinates
// OUTPUT:
             triples = a vector containing m triples of
   indices
                       corresponding to triangle
   vertices
#include<vector>
using namespace std;
typedef double T;
struct triple {
    int i, j, k;
    triple() {}
    triple(int i, int j, int k) : i(i), j(j), k(k) {}
vector<triple> delaunayTriangulation(vector<T>& x,
   vector<T>& y) {
        int n = x.size();
        vector<T> z(n);
        vector<triple> ret;
        for (int i = 0; i < n; i++)
            Z[i] = X[i] * X[i] + Y[i] * Y[i];
        for (int i = 0; i < n-2; i++) {
            for (int j = i+1; j < n; j++) {
                for (int k = i+1; k < n; k++) {
                    if (j == k) continue;
                    double xn = (y[j]-y[i])*(z[k]-z[i])
                        - (y[k]-y[i])*(z[i]-z[i]);
                    double yn = (x[k]-x[i])*(z[j]-z[i])
                        - (x[j]-x[i])*(z[k]-z[i]);
                    double zn = (x[j]-x[i])*(y[k]-y[i])
                        - (x[k]-x[i])*(y[i]-y[i]);
                    bool flag = zn < 0;
                    for (int m = 0; flaq && m < n; m++)</pre>
                         flag = flag && ((x[m]-x[i])*xn +
                                          (y[m]-y[i])*yn +
                                         (z[m]-z[i])*zn
                                             <= 0);
                    if (flag) ret.push_back(triple(i, j,
                         k));
        return ret;
```

```
int main()
{
    T xs[]={0, 0, 1, 0.9};
    T ys[]={0, 1, 0, 0.9};
    vector<T> x(&xs[0], &xs[4]), y(&ys[0], &ys[4]);
    vector<triple> tri = delaunayTriangulation(x, y);
    //expected: 0 1 3
    // 0 3 2

int i;
    for(i = 0; i < tri.size(); i++)
        printf("%d %d %d\n", tri[i].i, tri[i].j, tri[i].
        k);
    return 0;
}</pre>
```

3 Numerical algorithms

3.1 Number theory (modular, Chinese remainder, linear Diophantine)

```
// This is a collection of useful code for solving
   problems that
// involve modular linear equations. Note that all of
// algorithms described here work on nonnegative
   integers.
#include <iostream>
#include <vector>
#include <algorithm>
using namespace std;
typedef vector<int> VI;
typedef pair<int, int> PII;
// return a % b (positive value)
int mod(int a, int b) {
        return ((a%b) + b) % b;
// computes gcd(a,b)
int gcd(int a, int b) {
        while (b) { int t = a%b; a = b; b = t; }
        return a;
// computes lcm(a,b)
int lcm(int a, int b) {
        return a / gcd(a, b) *b;
// (a^b) mod m via successive squaring
int powermod(int a, int b, int m)
        int ret = 1;
        while (b)
```

```
a = mod(a*a, m);
                b >>= 1;
        return ret;
// returns q = qcd(a, b); finds x, y such that d = ax + b
int extended euclid(int a, int b, int &x, int &y) {
        int xx = y = 0;
        int yy = x = 1;
        while (b) {
                int q = a / b;
                int t = b; b = a%b; a = t;
                t = xx; xx = x - q*xx; x = t;
                t = yy; yy = y - q*yy; y = t;
        return a;
// finds all solutions to ax = b (mod n)
VI modular_linear_equation_solver(int a, int b, int n) {
        int x, y;
        VI ret;
        int g = extended_euclid(a, n, x, y);
        if (!(b%q)) {
                x = mod(x*(b / g), n);
                for (int i = 0; i < g; i++)
                        ret.push back (mod(x + i*(n / q)),
                            n));
        return ret;
// computes b such that ab = 1 \pmod{n}, returns -1 on
   failure
int mod_inverse(int a, int n) {
        int x, y;
        int g = extended_euclid(a, n, x, y);
        if (q > 1) return -1;
        return mod(x, n);
// Chinese remainder theorem (special case): find z such
//z % m1 = r1, z % m2 = r2. Here, z is unique modulo M
    = 1cm(m1, m2).
// Return (z, M). On failure, M = -1.
PII chinese remainder theorem(int m1, int r1, int m2,
   int r2) {
        int s, t;
        int g = extended_euclid(m1, m2, s, t);
        if (r1%g != r2%g) return make_pair(0, -1);
        return make_pair(mod(s*r2*m1 + t*r1*m2, m1*m2) /
            q, m1*m2 / g);
// Chinese remainder theorem: find z such that
//z % m[i] = r[i] for all i. Note that the solution is
// unique modulo M = lcm_i (m[i]). Return (z, M). On
// failure, M = -1. Note that we do not require the a[i
   1's
// to be relatively prime.
```

if (b & 1) ret = mod(ret*a, m);

```
PII chinese_remainder_theorem(const VI &m, const VI &r)
        PII ret = make pair(r[0], m[0]);
        for (int i = 1; i < m.size(); i++) {</pre>
                ret = chinese_remainder_theorem(ret.
                    second, ret.first, m[i], r[i]);
                if (ret.second == -1) break;
        return ret;
// computes x and y such that ax + by = c
// returns whether the solution exists
bool linear diophantine (int a, int b, int c, int &x, int
        if (!a && !b)
                if (c) return false;
                x = 0; y = 0;
                return true;
        if (!a)
                if (c % b) return false;
                x = 0; y = c / b;
                return true;
        if (!b)
                if (c % a) return false;
                x = c / a; y = 0;
                return true;
        int g = gcd(a, b);
        if (c % q) return false;
        x = c / q * mod_inverse(a / g, b / g);
        y = (c - a*x) / b;
        return true;
int main() {
        // expected: 2
        cout << gcd(14, 30) << endl;
        // expected: 2 -2 1
        int x, y;
        int g = extended_euclid(14, 30, x, y);
        cout << q << " " << x << " " << y << endl;
        // expected: 95 451
        VI sols = modular_linear_equation_solver(14, 30,
        for (int i = 0; i < sols.size(); i++) cout <<</pre>
           sols[i] << " ";
        cout << endl;</pre>
        // expected: 8
        cout << mod_inverse(8, 9) << endl;</pre>
        // expected: 23 105
        PII ret = chinese_remainder_theorem(VI({ 3, 5, 7
            }), VI({ 2, 3, 2 }));
        cout << ret.first << " " << ret.second << endl;</pre>
```

```
ret = chinese_remainder_theorem(VI({ 4, 6 }), VI
      ({ 3, 5 }));
cout << ret.first << " " << ret.second << endl;

// expected: 5 -15
if (!linear_diophantine(7, 2, 5, x, y)) cout <<
      "ERROR" << endl;
cout << x << " " << y << endl;
return 0;</pre>
```

3.2 Systems of linear equations, matrix inverse, determinant

```
// Gauss-Jordan elimination with full pivoting.
// Uses:
   (1) solving systems of linear equations (AX=B)
     (2) inverting matrices (AX=I)
     (3) computing determinants of square matrices
// Running time: O(n^3)
// INPUT:
             a[][] = an nxn matrix
             b[][] = an nxm matrix
// OUTPUT: X = an nxm matrix (stored in b[][])
             A^{-1} = an nxn matrix (stored in a[][])
             returns determinant of a[][]
#include <iostream>
#include <vector>
#include <cmath>
using namespace std;
const double EPS = 1e-10;
typedef vector<int> VI;
typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;
T GaussJordan (VVT &a, VVT &b) {
  const int n = a.size();
  const int m = b[0].size();
  VI irow(n), icol(n), ipiv(n);
  T \det = 1;
  for (int i = 0; i < n; i++) {</pre>
    int p_{1} = -1, p_{k} = -1;
    for (int j = 0; j < n; j++) if (!ipiv[j])</pre>
      for (int k = 0; k < n; k++) if (!ipiv[k])
        if (pj == -1 || fabs(a[j][k]) > fabs(a[pj][pk]))
    { pj = j; pk = k; }

if (fabs(a[pj][pk]) < EPS) { cerr << "Matrix is"
        singular." << endl; exit(0); }</pre>
    ipiv[pk]++;
    swap(a[pj], a[pk]);
    swap(b[pj], b[pk]);
    if (pj != pk) det *= -1;
    irow[i] = pj;
    icol[i] = pk;
```

T c = 1.0 / a[pk][pk];det *= a[pk][pk];a[pk][pk] = 1.0;for (int p = 0; p < n; p++) a[pk][p] *= c; for (int p = 0; p < m; p++) b[pk][p] *= c; for (int p = 0; p < n; p++) if (p != pk) { c = a[p][pk];a[p][pk] = 0;for (int q = 0; q < n; q++) a[p][q] -= a[pk][q] *for (int q = 0; q < m; q++) b[p][q] -= b[pk][q] *for (int p = n-1; p >= 0; p--) if (irow[p] != icol[p]) for (int k = 0; k < n; k++) swap(a[k][irow[p]], a[k][icol[p]]); return det; int main() { const int n = 4; const int m = 2; double $A[n][n] = {$ $\{1,2,3,4\},\{1,0,1,0\},\{5,3,2,4\},\{6,1,4,6\}\};$ **double** $B[n][m] = \{ \{1,2\}, \{4,3\}, \{5,6\}, \{8,7\} \};$ VVT a(n), b(n);for (int i = 0; i < n; i++) { a[i] = VT(A[i], A[i] + n);b[i] = VT(B[i], B[i] + m);double det = GaussJordan(a, b); // expected: 60 cout << "Determinant: " << det << endl;</pre> // expected: -0.233333 0.166667 0.133333 0.0666667 0.166667 0.166667 0.333333 -0.333333 0.05 - 0.75 - 0.1 0.2cout << "Inverse: " << endl;</pre> for (int i = 0; i < n; i++) { for (int j = 0; j < n; j++) cout << a[i][j] << ' ';</pre> cout << endl;</pre> // expected: 1.63333 1.3 -0.166667 0.5 2.36667 1.7 -1.85 - 1.35cout << "Solution: " << endl;</pre> for (int i = 0; i < n; i++) { for (int j = 0; j < m; j++) cout << b[i][j] << ' ';</pre> cout << endl;</pre>

3.3 Reduced row echelon form, matrix rank

```
// Reduced row echelon form via Gauss-Jordan elimination
// with partial pivoting. This can be used for
   computing
// the rank of a matrix.
// Running time: O(n^3)
// INPUT:
             a[][] = an nxm matrix
  OUTPUT: rref[][] = an nxm matrix (stored in a[][])
             returns rank of a[][]
#include <iostream>
#include <vector>
#include <cmath>
using namespace std;
const double EPSILON = 1e-10;
typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;
int rref(VVT &a) {
  int n = a.size();
  int m = a[0].size();
  int r = 0;
  for (int c = 0; c < m && r < n; c++) {</pre>
    int j = r;
    for (int i = r + 1; i < n; i++)
      if (fabs(a[i][c]) > fabs(a[j][c])) j = i;
    if (fabs(a[j][c]) < EPSILON) continue;</pre>
    swap(a[j], a[r]);
    T s = 1.0 / a[r][c];
    for (int j = 0; j < m; j++) a[r][j] *= s;
    for (int i = 0; i < n; i++) if (i != r) {
      T t = a[i][c];
      for (int j = 0; j < m; j++) a[i][j] -= t * a[r][j]
         ];
    r++;
  return r;
int main() {
  const int n = 5, m = 4;
  double A[n][m] = {
    {16, 2, 3, 13},
    { 5, 11, 10, 8},
    { 9, 7, 6, 12}, { 4, 14, 15, 1},
    {13, 21, 21, 13}};
  VVT a(n);
  for (int i = 0; i < n; i++)
    a[i] = VT(A[i], A[i] + m);
  int rank = rref(a);
```

```
// expected: 3
cout << "Rank: " << rank << endl;
// expected: 1 0 0 1
// 0 1 0 3
// 0 0 1 -3
// 0 0 0 3.10862e-15
// 0 0 0 2.22045e-15
cout << "rref: " << endl;
for (int i = 0; i < 5; i++) {
   for (int j = 0; j < 4; j++)
      cout << endl;
}
cout << endl;
}</pre>
```

3.4 Simplex algorithm

```
// Two-phase simplex algorithm for solving linear
   programs of the form
       maximize
       subject to Ax <= b
                    x >= 0
  INPUT: A -- an m x n matrix
          b -- an m-dimensional vector
          c -- an n-dimensional vector
          x -- a vector where the optimal solution will
   be stored
// OUTPUT: value of the optimal solution (infinity if
           above, nan if infeasible)
// To use this code, create an LPSolver object with A, b
   , and c as
// arguments. Then, call Solve(x).
#include <iostream>
#include <iomanip>
#include <vector>
#include <cmath>
#include <limits>
using namespace std;
typedef long double DOUBLE;
typedef vector<DOUBLE> VD;
typedef vector<VD> VVD;
typedef vector<int> VI;
const DOUBLE EPS = 1e-9;
struct LPSolver {
  int m, n;
  VI B, N;
  VVD D;
  LPSolver (const VVD &A, const VD &b, const VD &c) :
   m(b.size()), n(c.size()), N(n + 1), B(m), D(m + 2,
       VD(n + 2) {
    for (int i = 0; i < m; i++) for (int j = 0; j < n; j
```

```
++) D[i][j] = A[i][j];
  for (int i = 0; i < m; i++) { B[i] = n + i; D[i][n]
     = -1; D[i][n + 1] = b[i];
  for (int j = 0; j < n; j++) { N[j] = j; D[m][j] = -c
     [ † ]; }
 N[n] = -1; D[m + 1][n] = 1;
void Pivot(int r, int s)
  double inv = 1.0 / D[r][s];
  for (int i = 0; i < m + 2; i++) if (i != r)
    for (int j = 0; j < n + 2; j++) if (j != s)
      D[i][j] = D[r][j] * D[i][s] * inv;
  for (int j = 0; j < n + 2; j++) if (j != s) D[r][j]
  for (int i = 0; i < m + 2; i++) if (i != r) D[i][s]
     *=-inv;
  D[r][s] = inv;
  swap(B[r], N[s]);
bool Simplex(int phase) {
  int x = phase == 1 ? m + 1 : m;
  while (true) {
    int s = -1;
    for (int j = 0; j <= n; j++) {
      if (phase == 2 \&\& N[\dot{j}] == -1) continue;
      if (s == -1 || D[x][\dot{\eta}] < D[x][s] || D[x][\dot{\eta}] == D
          [x][s] \&\& N[j] < N[s]) s = j;
    if (D[x][s] > -EPS) return true;
    int r = -1;
    for (int i = 0; i < m; i++) {
      if (D[i][s] < EPS) continue;</pre>
      if (r == -1 \mid | D[i][n + 1] / D[i][s] < D[r][n +
         1] / D[r][s] ||
        (D[i][n + 1] / D[i][s]) == (D[r][n + 1] / D[r]
           [s] && B[i] < B[r]) r = i;
    if (r == -1) return false;
    Pivot(r, s);
DOUBLE Solve(VD &x) {
  int r = 0;
  for (int i = 1; i < m; i++) if (D[i][n + 1] < D[r][n
      + 1]) r = i;
  if (D[r][n + 1] < -EPS) {
    Pivot(r, n);
    if (!Simplex(1) || D[m + 1][n + 1] < -EPS) return
       -numeric limits<DOUBLE>::infinity();
    for (int i = 0; i < m; i++) if (B[i] == -1) {
      int s = -1;
      for (int j = 0; j \le n; j++)
        if (s == -1 || D[i][j] < D[i][s] || D[i][j] ==
            D[i][s] \&\& N[j] < N[s]) s = j;
      Pivot(i, s);
  if (!Simplex(2)) return numeric_limits<DOUBLE>::
```

```
infinity();
    x = VD(n);
    for (int i = 0; i < m; i++) if (B[i] < n) x[B[i]] =
        D[i][n + 1];
    return D[m][n + 1];
};
int main() {
  const int m = 4;
  const int n = 3;
  DOUBLE A[m][n] = {
    \{ 6, -1, 0 \},
      -1, -5, 0},
      1, 5, 1 },
    \{-1, -5, -1\}
  DOUBLE _b[m] = { 10, -4, 5, -5 };
  DOUBLE _{c[n]} = \{ 1, -1, 0 \};
  VVD A(m);
  VD b(\underline{b}, \underline{b} + m);
  VD c(\underline{c}, \underline{c} + n);
  for (int i = 0; i < m; i++) A[i] = VD(_A[i], _A[i] + n
     );
  LPSolver solver(A, b, c);
  DOUBLE value = solver.Solve(x);
  cerr << "VALUE: " << value << endl; // VALUE: 1.29032</pre>
  cerr << "SOLUTION:"; // SOLUTION: 1.74194 0.451613 1</pre>
  for (size t i = 0; i < x.size(); i++) cerr << " " << x</pre>
  cerr << endl;
  return 0;
```

3.5 Fast Fourier transform (C++)

```
// Source: My (forked) ICPC codebook: https://github.com
   /aayush9/ICPC-CodeBook/blob/master/code/FFT.cc
namespace FFT{
  vector<complex<long double>> A;
  int n, L;
  int ReverseBits(int k) {
    int ret = 0;
    for (int i=0; i++<L; k/=2)
      ret = (ret << 1) | (k&1);
    return ret;
  void BitReverseCopy(vector<complex<long double>> a) {
    for (n=1, L=0; n<a.size(); n<<=1, L++);
    A.resize(n);
    for (int k = 0; k < n; k++)
      A[ReverseBits(k)] = a[k];
  vector<complex<long double>> DFT(vector<complex<long</pre>
     double>> a, bool inverse) {
    BitReverseCopy(a);
    for (int s = 1; s <= L; s++) {
```

```
int m = 1 << s;</pre>
    complex<long double> wm = exp(complex<long double</pre>
       >(0, 2.0 * M_PI / m));
    if (inverse) wm = complex<long double>(1, 0) / wm;
    for (int k = 0; k < n; k += m) {
      complex<long double> w = 1;
      for (int j = 0; j < m/2; j++) {
        complex < long double > t = w * A[k + j + m/2];
        complex<long double> u = A[k + j];
        A[k + j] = u + t;
        A[k + j + m/2] = u - t;

w = w * wm;
    if (inverse) for (int i = 0; i < n; i++) A[i] /= n</pre>
    return A:
vector<long double> Convolution(vector<long double> a,
   vector<long double> b) {
  int L = 1;
  while ((1 << L) < a.size()) L++;</pre>
  while ((1 << L) < b.size()) L++;</pre>
  int n = 1 << (L+1);
  vector<complex<long double>> aa, bb;
  for (size t i = 0; i < n; i++) aa.push back(i < a.
     size() ? complex<long double>(a[i], 0) : 0);
  for (size_t i = 0; i < n; i++) bb.push_back(i < b.</pre>
     size() ? complex<long double>(b[i], 0) : 0);
  vector<complex<long double>> AA = DFT(aa, false);
  vector<complex<long double>> BB = DFT(bb, false);
  vector<complex<long double>> CC;
  for (size t i = 0; i < AA.size(); i++) CC.push back(
     AA[i] * BB[i]);
  vector<complex<long double>> cc = DFT(CC, true);
  vector<long double> c;
  for (int i=0;i<a.size()+b.size()-1;++i)</pre>
    c.push_back(cc[i].real());
  return c;
```

4 Graph algorithms

4.1 Bellman-Ford shortest paths with negative edge weights

```
// This function runs the Bellman-Ford algorithm for
    single source
// shortest paths with negative edge weights. The
    function returns
// false if a negative weight cycle is detected.
    Otherwise, the
// function returns true and dist[i] is the length of
    the shortest
// path from start to i.
```

```
// Running time: O(|V|^3)
              start, w[i][j] = cost of edge from i to j
     OUTPUT: dist[i] = min weight path from start to i
              prev[i] = previous node on the best path
   from the
                         start node
#include <iostream>
#include <queue>
#include <cmath>
#include <vector>
using namespace std;
typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;
typedef vector<int> VI;
typedef vector<VI> VVI;
bool BellmanFord (const VVT &w, VT &dist, VI &prev, int
   start){
  int n = w.size();
  prev = VI(n, -1);
  dist = VT(n, 1000000000);
  dist[start] = 0;
  for (int k = 0; k < n; k++) {
    for (int i = 0; i < n; i++) {
      for (int j = 0; j < n; j++) {
        if (dist[j] > dist[i] + w[i][j]){
          if (k == n-1) return false;
          dist[j] = dist[i] + w[i][j];
          prev[j] = i;
  return true;
```

4.2 Strongly connected components

```
void fill backward(int x)
  int i;
  v[x] = false;
  group_num[x]=group_cnt;
  for(i=spr[x];i;i=er[i].nxt) if(v[er[i].e])
     fill backward(er[i].e);
void add edge(int v1, int v2) //add edge v1->v2
  e [++E].e=v2; e [E].nxt=sp [v1]; sp [v1]=E;
  er[ E].e=v1; er[E].nxt=spr[v2]; spr[v2]=E;
void SCC()
  int i;
  stk[0]=0;
  memset(v, false, sizeof(v));
  for (i=1; i <= V; i++) if (!v[i]) fill_forward(i);</pre>
  group_cnt=0;
  for(i=stk[0];i>=1;i--) if(v[stk[i]]) {group cnt++;
     fill_backward(stk[i]);}
```

4.3 Eulerian path

```
struct Edge;
typedef list<Edge>::iterator iter;
struct Edge
        int next vertex;
        iter reverse_edge;
        Edge (int next vertex)
                :next_vertex(next_vertex)
};
const int max_vertices = ;
int num vertices;
list<Edge> adj[max_vertices];
                                        // adjacency
   list
vector<int> path;
void find path(int v)
        while (adj[v].size() > 0)
                int vn = adj[v].front().next_vertex;
                adj[vn].erase(adj[v].front().
                    reverse edge);
                adj[v].pop_front();
                find_path(vn);
        path.push back(v);
void add_edge(int a, int b)
        adj[a].push_front(Edge(b));
```

```
iter ita = adj[a].begin();
adj[b].push_front(Edge(a));
iter itb = adj[b].begin();
ita->reverse_edge = itb;
itb->reverse_edge = ita;
```

4.4 Minimum spanning trees

```
// This function runs Prim's algorithm for constructing
// weight spanning trees.
// Running time: O(|V|^2)
     INPUT: w[i][j] = cost \ of \ edge \ from \ i \ to \ j
              NOTE: Make sure that w[i][j] is
   nonnegative and
              symmetric. Missing edges should be given
              weight.
     OUTPUT: edges = list of pair<int,int> in minimum
   spanning tree
              return total weight of tree
#include <iostream>
#include <queue>
#include <cmath>
#include <vector>
using namespace std;
typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;
typedef vector<int> VI;
typedef vector<VI> VVI;
typedef pair<int, int> PII;
typedef vector<PII> VPII;
T Prim (const VVT &w, VPII &edges) {
  int n = w.size();
  VI found (n);
  VI prev (n, -1);
  VT dist (n, 1000000000);
  int here = 0;
  dist[here] = 0;
  while (here !=-1) {
    found[here] = true;
    int best = -1;
    for (int k = 0; k < n; k++) if (!found[k]) {</pre>
      if (w[here][k] != -1 \&\& dist[k] > w[here][k]) {
        dist[k] = w[here][k];
        prev[k] = here;
      if (best == -1 \mid \mid dist[k] < dist[best]) best = k;
    here = best;
```

```
T tot weight = 0;
  for (int i = 0; i < n; i++) if (prev[i] != -1) {
    edges.push_back (make_pair (prev[i], i));
    tot weight += w[prev[i]][i];
  return tot_weight;
int main(){
  int ww[5][5] = {
    \{0, 400, 400, 300, 600\},\
    \{400, 0, 3, -1, 7\},\
    {400, 3, 0, 2, 0},
{300, -1, 2, 0, 5},
{600, 7, 0, 5, 0}
  VVT w(5, VT(5));
  for (int i = 0; i < 5; i++)
    for (int j = 0; j < 5; j++)
      w[i][j] = ww[i][j];
  // expected: 305
                 2 1
                 3 2
0 3
                 2 4
  VPII edges;
  cout << Prim (w, edges) << endl;</pre>
  for (int i = 0; i < edges.size(); i++)
    cout << edges[i].first << " " << edges[i].second <<</pre>
        endl:
```

5 Data structures

5.1 Suffix array

```
// Suffix array construction in O(L log^2 L) time.
   Routine for
// computing the length of the longest common prefix of
// suffixes in O(log L) time.
// INPUT: string s
// OUTPUT: array suffix[] such that suffix[i] = index (
   from 0 to L-1)
            of substring s[i...L-1] in the list of
   sorted suffixes.
            That is, if we take the inverse of the
   permutation suffix[],
            we get the actual suffix array.
#include <vector>
#include <iostream>
#include <string>
using namespace std;
```

```
struct SuffixArray {
  const int L;
  string s;
  vector<vector<int> > P;
  vector<pair<int,int>,int> > M;
  SuffixArray(const string &s) : L(s.length()), s(s), P
     (1, vector<int>(L, 0)), M(L) {
    for (int i = 0; i < L; i++) P[0][i] = int(s[i]);
    for (int skip = 1, level = 1; skip < L; skip \star= 2,
       level++) {
      P.push back(vector<int>(L, 0));
      for (int i = 0; i < L; i++)
        M[i] = make_pair(make_pair(P[level-1][i], i +
           skip < L ? P[level-1][i + skip] : -1000), i);
      sort(M.begin(), M.end());
      for (int i = 0; i < L; i++)
        P[level][M[i].second] = (i > 0 && M[i].first ==
           M[i-1].first) ? P[level][M[i-1].second] : i;
  vector<int> GetSuffixArray() { return P.back(); }
  // returns the length of the longest common prefix of
     s[i...L-1] and s[j...L-1]
  int LongestCommonPrefix(int i, int j) {
    int len = 0;
    if (i == j) return L - i;
    for (int k = P.size() - 1; k \ge 0 && i < L && j < L;
      if (P[k][i] == P[k][j]) {
        i += 1 << k;
        \dagger += 1 << k;
        len += 1 << k;
    return len;
};
// BEGIN CUT
// BEGIN CUT
// The following code solves UVA problem 11512: GATTACA. \overline{\textbf{5.2}} KD-tree
#ifdef TESTING
int main() {
  int T;
  cin >> T;
  for (int caseno = 0; caseno < T; caseno++) {</pre>
    string s;
    cin >> s;
    SuffixArray array(s);
    vector<int> v = array.GetSuffixArray();
    int bestlen = -1, bestpos = -1, bestcount = 0;
    for (int i = 0; i < s.length(); i++) {</pre>
      int len = 0, count = 0;
      for (int j = i+1; j < s.length(); j++) {</pre>
        int 1 = array.LongestCommonPrefix(i, j);
        if (1 >= len) {
          if (1 > len) count = 2; else count++;
          len = 1;
```

```
if (len > bestlen || len == bestlen && s.substr(
         bestpos, bestlen) > s.substr(i, len)) {
        bestlen = len;
        bestcount = count;
        bestpos = i;
    if (bestlen == 0) {
      cout << "No repetitions found!" << endl;</pre>
    } else {
      cout << s.substr(bestpos, bestlen) << " " <<</pre>
         bestcount << endl;</pre>
#else
// END CUT
int main() {
  // bobocel is the O'th suffix
  // obocel is the 5'th suffix
      bocel is the 1'st suffix
       ocel is the 6'th suffix
        cel is the 2'nd suffix
          el is the 3'rd suffix
          l is the 4'th suffix
  SuffixArray suffix("bobocel");
 vector<int> v = suffix.GetSuffixArray();
  // Expected output: 0 5 1 6 2 3 4
  for (int i = 0; i < v.size(); i++) cout << v[i] << " "</pre>
  cout << endl:
  cout << suffix.LongestCommonPrefix(0, 2) << endl;</pre>
// BEGIN CUT
#endif
// END CUT
```

```
// A straightforward, but probably sub-optimal KD-tree
   implmentation
// that's probably good enough for most things (current
   it's a
// 2D-tree)
//
  - constructs from n points in O(n lg^2 n) time
// - handles nearest-neighbor query in O(lg n) if
  points are well
   distributed
// - worst case for nearest-neighbor may be linear in
  pathological
     case
// Sonny Chan, Stanford University, April 2009
#include <iostream>
#include <vector>
```

```
#include <limits>
#include <cstdlib>
using namespace std;
// number type for coordinates, and its maximum value
typedef long long ntype;
const ntype sentry = numeric_limits<ntype>::max();
// point structure for 2D-tree, can be extended to 3D
struct point {
   ntype x, y;
    point (ntype xx = 0, ntype yy = 0) : x(xx), y(yy) {}
};
bool operator==(const point &a, const point &b) {
    return a.x == b.x && a.v == b.v;
// sorts points on x-coordinate
bool on_x(const point &a, const point &b) {
    return a.x < b.x;</pre>
// sorts points on y-coordinate
bool on_y (const point &a, const point &b) {
    return a.y < b.y;</pre>
// squared distance between points
ntype pdist2(const point &a, const point &b) {
    ntype dx = a.x-b.x, dy = a.y-b.y;
    return dx*dx + dy*dy;
// bounding box for a set of points
struct bbox{
    ntype x0, x1, y0, y1;
    bbox(): x0(sentry), x1(-sentry), y0(sentry), y1(-
       sentry) {}
    // computes bounding box from a bunch of points
    void compute(const vector<point> &v) {
        for (int i = 0; i < v.size(); ++i) {</pre>
            x0 = \min(x0, v[i].x);
                                   x1 = max(x1, v[i].x)
            y0 = min(y0, v[i].y); y1 = max(y1, v[i].y)
    // squared distance between a point and this bbox, 0
        if inside
    ntype distance(const point &p) {
        if (p.x < x0) {
            if (p.y < y0)
                                 return pdist2(point(x0,
               y0), p);
            else if (p.y > y1)
                                return pdist2(point(x0,
               y1), p);
            else
                                 return pdist2(point(x0,
               p.y), p);
        else if (p.x > x1) {
            if (p.y < y0)
                                 return pdist2(point(x1,
               y0), p);
            else if (p.y > y1) return pdist2(point(x1,
```

```
y1), p);
            else
                                return pdist2 (point (x1,
               p.y), p);
        else {
            if (p.y < y0)
                                return pdist2 (point (p.x,
                y0), p);
            else if (p.y > y1)
                                return pdist2 (point (p.x,
                y1), p);
            else
                                return 0;
        }
    }
} ;
// stores a single node of the kd-tree, either internal
   or leaf
struct kdnode {
                    // true if this is a leaf node (has
    bool leaf;
       one point)
    point pt;
                    // the single point of this is a
       leaf
    bbox bound;
                    // bounding box for set of points in
        children
    kdnode *first, *second; // two children of this kd-
       node
    kdnode() : leaf(false), first(0), second(0) {}
    ~kdnode() { if (first) delete first; if (second)
       delete second; }
    // intersect a point with this node (returns squared
        distance)
    ntype intersect(const point &p) {
        return bound.distance(p);
    // recursively builds a kd-tree from a given cloud
       of points
    void construct(vector<point> &vp) {
        // compute bounding box for points at this node
        bound.compute(vp);
        // if we're down to one point, then we're a leaf
            node
        if (vp.size() == 1) {
            leaf = true;
            pt = vp[0];
        else {
            // split on x if the bbox is wider than high
                 (not best heuristic...)
            if (bound.x1-bound.x0 >= bound.y1-bound.y0)
                sort(vp.begin(), vp.end(), on_x);
            // otherwise split on y-coordinate
                sort(vp.begin(), vp.end(), on_y);
            // divide by taking half the array for each
               child
            // (not best performance if many duplicates
               in the middle)
            int half = vp.size()/2;
            vector<point> vl(vp.begin(), vp.begin()+half
```

```
vector<point> vr(vp.begin()+half, vp.end());
            first = new kdnode(); first->construct(v1)
            second = new kdnode(); second->construct(vr
               );
};
// simple kd-tree class to hold the tree and handle
   queries
struct kdtree{
   kdnode *root;
    // constructs a kd-tree from a points (copied here,
       as it sorts them)
    kdtree(const vector<point> &vp) {
       vector<point> v(vp.begin(), vp.end());
        root = new kdnode();
       root->construct(v);
    ~kdtree() { delete root; }
    // recursive search method returns squared distance
       to nearest point
   ntype search(kdnode *node, const point &p)
        if (node->leaf) {
            // commented special case tells a point not
               to find itself
              if (p == node->pt) return sentry;
                return pdist2(p, node->pt);
        ntype bfirst = node->first->intersect(p);
        ntype bsecond = node->second->intersect(p);
        // choose the side with the closest bounding box
            to search first
        // (note that the other side is also searched if
            needed)
        if (bfirst < bsecond) {</pre>
            ntype best = search(node->first, p);
            if (bsecond < best)</pre>
                best = min(best, search(node->second, p)
            return best;
        else {
            ntype best = search(node->second, p);
            if (bfirst < best)</pre>
                best = min(best, search(node->first, p))
            return best;
    // squared distance to the nearest
    ntype nearest (const point &p) {
        return search(root, p);
};
```

5.3 Splay tree

```
#include <cstdio>
#include <algorithm>
using namespace std;
const int N_MAX = 130010;
const int oo = 0x3f3f3f3f;
struct Node
  Node *ch[2], *pre;
  int val, size;
  bool isTurned;
} nodePool[N MAX], *null, *root;
Node *allocNode(int val)
  static int freePos = 0;
  Node *x = &nodePool[freePos ++];
  x->val = val, x->isTurned = false;
  x->ch[0] = x->ch[1] = x->pre = null;
  x->size = 1;
  return x;
inline void update(Node *x)
  x->size = x->ch[0]->size + x->ch[1]->size + 1;
inline void makeTurned(Node *x)
  if(x == null)
    return;
  swap (x->ch[0], x->ch[1]);
  x->isTurned ^= 1;
inline void pushDown(Node *x)
  if(x->isTurned)
    makeTurned(x->ch[0]);
```

```
makeTurned(x->ch[1]);
    x->isTurned ^= 1;
inline void rotate(Node *x, int c)
  Node *y = x - > pre;
  x->pre^{-}=y->pre;
  if(v->pre != null)
    y-pre-ch[y == y-pre-ch[1]] = x;
  y \rightarrow ch[!c] = x \rightarrow ch[c];
  if(x->ch[c] != null)
    x->ch[c]->pre = y;
  x->ch[c] = y, y->pre = x;
  update(y);
  if(y == root)
    root = x;
void splay(Node *x, Node *p)
  while (x->pre != p)
    if(x->pre->pre == p)
      rotate(x, x == x->pre->ch[0]);
    else
      Node *y = x - pre, *z = y - pre;
      if(y == z -> ch[0])
        if(x == y -> ch[0])
           rotate(y, 1), rotate(x, 1);
           rotate (x, 0), rotate (x, 1);
      else
        if(x == y->ch[1])
           rotate(y, 0), rotate(x, 0);
           rotate (x, 1), rotate (x, 0);
  update(x);
void select(int k, Node *fa)
  Node *now = root;
  while (1)
    pushDown(now);
    int tmp = now->ch[0]->size + 1;
    if(tmp == k)
      break;
    else if(tmp < k)</pre>
      now = now -> ch[1], k -= tmp;
    else
      now = now -> ch[0];
  splay(now, fa);
```

```
Node *makeTree(Node *p, int 1, int r)
  if(1 > r)
    return null;
  int mid = (1 + r) / 2;
  Node *x = allocNode(mid);
  x->pre = p;
  x \rightarrow ch[0] = makeTree(x, 1, mid - 1);
  x->ch[1] = makeTree(x, mid + 1, r);
  update(x);
  return x;
int main()
  int n, m;
  null = allocNode(0);
  null->size = 0;
  root = allocNode(0);
  root->ch[1] = allocNode(oo);
  root->ch[1]->pre = root;
  update (root);
  scanf("%d%d", &n, &m);
  root \rightarrow ch[1] \rightarrow ch[0] = makeTree(root \rightarrow ch[1], 1, n);
  splay(root->ch[1]->ch[0], null);
  while (m --)
    int a, b;
    scanf("%d%d", &a, &b);
    a ++, b ++;
    select(a - 1, null);
    select(b + 1, root);
    makeTurned(root->ch[1]->ch[0]);
  for (int i = 1; i \le n; i ++)
    select(i + 1, null);
    printf("%d ", root->val);
```

5.4 Lazy segment tree

```
[2*u+1]+=lazv[u];
                   lazy[u]=0;
         if(j<a || i>b || a>b) return;
         if(j>=b && i<=a){
                   segtree[u] += (b-a+1) *x;
                   if (a!=b) lazy[u*2]+=x, lazy[2*u
                       +1]+=x;
                   return;
         update (u*2,a,(a+b)/2,i,j,x); update (u*2,a,(a+b)/2,i,j,x);
         *2+1, (a+b)/2+1, b, i, j, x);
segtree[u]=segtree[u*2]+segtree[u*2+1];
void update(int i, T x) {
         update(1, 0, segtree.size()/4-1, i, i, x
             );
void update(int i, int j, T x) {
         update(1, 0, segtree.size()/4-1, i, j, x
T query(int u, int a, int b, int i, int j){
         if(j<a || i>b || a>b) return 0;
         if(lazy[u]){
                   segtree [u] += (b-a+1) * lazy [u];
                   if(a!=b)
                            lazy[u*2] += lazy[u], lazy
                                [2*u+1] += lazy[u];
                   lazv[u]=0;
         if(j>=b && i<=a) return segtree[u];</pre>
         return query (u*2,a,(a+b)/2,i,j) +query (u*2,a,(a+b)/2,i,j)
             *2+1,1+(a+b)/2,b,i,j);
T query(int i, int j) {
         return query (1, 0, segtree.size()/4-1,i,j)
```

5.5 Centroid decomposition

};

```
set < int > v[100005];
map<int,int> mp[100005];
int n,up[100005][17],lvl[100005],par[100005],CNT,siz
    [100005], tin[100005], tout[100005];
void dfspre(int u, int dad=1, int depth = 0) {
    static int clk = 0;
    tin[u]=clk++;
    up[u][0] = dad;
    lvl[u] = depth;
    for (int i=1; i<17; ++i)</pre>
        up[u][i] = up[up[u][i-1]][i-1];
    for(int i:v[u]) if(i!=dad)
        dfspre(i,u,depth+1);
    tout[u]=clk++;
int dfs(int u, int dad) {
    siz[u] = 1;
```

```
for(int i:v[u]) if(i!=dad)
        siz[u] += dfs(i,u);
    return siz[u];
int centroid(int u, int dad) {
    for(int i:v[u]) if(i!=dad && siz[i]>CNT)
        return centroid(i,u);
    return u;
void decompose(int u, int dad){
    CNT = dfs(u, dad)/2;
    int centre = centroid(u, dad);
    par[centre] = dad;
    for(int i:v[centre]) if(i!=dad){
        v[i].erase(centre);
        decompose (i, centre);
    v[centre].clear();
int lca(int u, int v) {
    if(lvl[u]>lvl[v]) swap(u,v);
    if(tin[u] <=tin[v] && tout[v] <=tout[u]) return u;</pre>
    for(int i=17;i--;)
        if(!(tin[up[u][i]]<=tin[v] && tout[v]<=tout[up[u]</pre>
            ][i]]))
            u = up[u][i];
    return up[u][0];
void update(int u) {
    for(int node = u;u;u = par[u])
        ++mp[u][lvl[node]+lvl[u] - 2*lvl[lca(u, node)]];
int get(int u) {
    int ans = INT MAX;
    for(int node = u; u; u = par[u])
        ans = min(ans,lvl[u]+lvl[node]-2*lvl[lca(u,node)]
            ] + (*mp[u].begin()).first);
    return ans;
```

5.6 Heavy-Light decomposition

```
#include <bits/stdc++.h>
using namespace std;
int a[100005],sz[100005],lvl[100005];
int seg id[100005],pos id[100005],parent[100005];
int CNT;
int e_to_v[100005];
vector<pair<int,int>> edges(100005);
vector<pair<int,int>> v[100005];
vector<int> chain[100005];
template <typename T>
class SegmentTree{
  vector<T> segtree, lazy;
  public:
  SegmentTree(int size){
    segtree.resize(4*size,0);
    lazy.resize(4*size,0);
  void update(int u, int a, int b, int i, int j, T x) {
```

```
hld(i.first,u,++CNT, u);
    if(lazv[u]) {
      segtree [u] += (b-a+1) * lazy [u];
      if(a!=b)
                                                                int main(){
                                                                  ios base::sync with stdio(0);
        lazy[u*2] += lazy[u], lazy[2*u+1] += lazy[u];
      lazy[u]=0;
                                                                   cin.tie(0); cout.tie(0);
                                                                   int T, n, x, y;
                                                                   for(cin>>T; T-- && cin>>n;) {
    if(j<a || i>b || a>b) return;
    if( j>=b && i<=a) {
                                                                     CNT = 0;
      seatree[u]=x;
                                                                     for(int i=0;i<100005;chain[i].clear(), v[i].clear(),</pre>
      if (a!=b) lazy[u*2] += x, lazy[2*u+1] += x;
      return;
                                                                     for (int i=1, c; i < n; ++i) {</pre>
    update (u*2,a,(a+b)/2,i,j,x); update (u*2+1,(a+b)/2+1,
                                                                       cin>>x>>y>>c;
       b, i, j, x);
                                                                       edges[i] = \{x,y\};
    seqtree [u] =max (segtree [u*2], segtree [u*2+1]);
                                                                       v[x].push_back({y,c});
                                                                       v[y].push_back({x,c});
  void update(int i, T x) {
    update(1, 0, segtree.size()/4-1, i, i, x);
                                                                     dfs(1);
  void update(int i, int j, T x) {
                                                                     for (int i=1; i < n; ++i) {</pre>
    update(1, 0, segtree.size()/4-1, i, j, x);
                                                                       int u = edges[i].first, v = edges[i].second;
                                                                       e to v[i] = (lvl[u] < lvl[v]?v:u);
  T query(int u, int a, int b, int i, int j) {
    if(j<a || i>b || a>b) return 0;
                                                                    hld(1);
    if(lazv[u]) {
      segtree [u] += (b-a+1) * lazy [u];
                                                                     vector<SegmentTree<int>> ST;
                                                                     for (int i=0; i <= CNT; ++i) {</pre>
      if (a!=b)
        lazy[u*2] += lazy[u], lazy[2*u+1] += lazy[u];
                                                                       ST.push back(SegmentTree<int>(chain[i].size()));
      lazy[u]=0;
                                                                       for(auto u:chain[i])
                                                                         ST[i].update(pos_id[u],a[u]);
    if(j>=b && i<=a) return segtree[u];</pre>
    return max (query (u*2, a, (a+b)/2, i, j), query (u*2+1, 1+ (a
                                                                     for(string type;cin>>type && type!="DONE";){
        +b)/2,b,i,i);
                                                                       cin>>x>>y;
                                                                       if(type=="QUERY") {
  T query(int i, int j){
                                                                         int res = 0;
    return query(1,0,segtree.size()/4-1,i,i);
                                                                         while (x!=y) {
                                                                           if(seg_id[x]>seg_id[y]) swap(x,y);
                                                                           else if(seg_id[x] == seg_id[y]){
int dfs(int u, int dad=1, int depth=1) {
                                                                             if (pos_id[x]>pos_id[y]) swap(x,y);
  lvl[u] = depth;
                                                                             res = max(res, ST[seg_id[y]].query(pos_id[x
  sz[u] = 1;
                                                                                 ]+1,pos_id[y]));
  for (auto i:v[u]) if (i.first!=dad)
                                                                             break:
    sz[u] += dfs(i.first,u,depth+1);
  return sz[u];
                                                                           res = max(res, ST[seq_id[y]].query(0,pos_id[y
                                                                               1));
void hld(int u, int dad = 1, int chain no = 0, int
                                                                           y = parent[y];
   chain parent = 0) {
  seq id[\bar{u}] = chain no;
                                                                         cout << res << '\n';
  pos_id[u] = chain[chain_no].size();
  parent[u] = chain_parent;
                                                                       else if(type == "CHANGE") {
  chain[chain_no].push_back(u);
                                                                         int u = e_to_v[x];
  int max_sz = 0, heavy_child = -1;
                                                                         ST[seq_id[u]].update(pos_id[u],y);
  for(auto i:v[u]) if(i.first!=dad){
    a[i.first] = i.second;
                                                                    }
    if(max sz < sz[i.first])</pre>
      max_sz = sz[i.first], heavy_child = i.first;
  if (heavy child!=-1)
    hld(heavy_child, u, chain_no, chain_parent);
  for(auto i:v[u]) if(i.first!=dad && i.first!=
```

heavy_child)

6 Miscellaneous

6.1 Dynamic Programming(DnC)

```
long long dp [21] [100005];
void cost(int x, int y);
void computeDP(int idx,int jleft,int jright,int kleft,
   int kright) {
         if(jleft>jright) return;
         int jmid=(jleft+jright)/2;
         int bestk=jmid;
         for(int k=kleft; k<=min(kright, jmid); ++k) {</pre>
                  cost(k, jmid);
                  if (dp[idx-1][k-1]+tot<dp[idx][jmid])
                           dp[idx][jmid]=dp[idx-1][k-1]+tot
                               ,bestk=k;
         computeDP(idx, jleft, jmid-1, kleft, bestk);
         computeDP (idx, jmid+1, jright, bestk, kright);
int main(){
         for (int i=0; i <= k; ++i)
                  for (int j=0; j<=n; dp[i][j++]=1e17);</pre>
         dp[0][0]=0;
         for(int i=1;i<=k;++i)</pre>
                  computeDP (i, 1, n, 1, n);
         cout << dp[k][n];
```

6.2 Longest increasing subsequence

```
// Given a list of numbers of length n, this routine
   extracts a
// longest increasing subsequence.
// Running time: O(n log n)
     INPUT: a vector of integers
    OUTPUT: a vector containing the longest increasing
   subsequence
typedef vector<int> VI;
typedef pair<int, int> PII;
typedef vector<PII> VPII;
#define STRICTLY INCREASING
VI LongestIncreasingSubsequence(VI v) {
  VPII best;
  VI dad(v.size(), -1);
  for (int i = 0; i < v.size(); i++) {</pre>
#ifdef STRICTLY INCREASING
    PII item = make_pair(v[i], 0);
    VPII::iterator it = lower_bound(best.begin(), best.
       end(), item);
    item.second = i;
```

```
#else
    PII item = make pair(v[i], i);
    VPII::iterator it = upper bound(best.begin(), best.
       end(), item);
#endif
    if (it == best.end()) {
      dad[i] = (best.size() == 0 ? -1 : best.back().
         second);
      best.push back(item);
    } else {
      dad[i] = it == best.begin() ? -1 : prev(it) ->
      *it = item;
 VI ret;
  for (int i = best.back().second; i >= 0; i = dad[i])
    ret.push back(v[i]);
  reverse(ret.begin(), ret.end());
  return ret;
```

6.3 Knuth-Morris-Pratt

```
typedef vector<int> VI;
void buildPi(string& p, VI& pi){
  pi = VI(p.length());
  int k = -2;
  for(int i = 0; i < p.length(); i++) {
    while (k \ge -1 \& \& p[k+1] != p[i])
      k = (k == -1) ? -2 : pi[k];
    pi[i] = ++k;
int KMP(string& t, string& p){
  VI pi;
  buildPi(p, pi);
  int k = -1;
  for(int i = 0; i < t.length(); i++) {</pre>
    while (k \ge -1 \&\& p[k+1] != t[i])
      k = (k == -1) ? -2 : pi[k];
    k++;
    if(k == p.length() - 1) {
      // p matches t[i-m+1, ..., i]
      cout << "matched at index " << i-k << ": ";</pre>
      cout << t.substr(i-k, p.length()) << endl;</pre>
      k = (k == -1) ? -2 : pi[k];
  return 0;
int main(){
  string a = "AABAACAADAABAABA", b = "AABA";
  KMP(a, b); // expected matches at: 0, 9, 12
```