

# Indexing

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Notes

19<sup>th</sup> Hard day challenge :

1. Assignments + Revision (MCQs)
2. Backlog (Assignments of prev. session)
3. Additional Questions



# Intro to Indexing

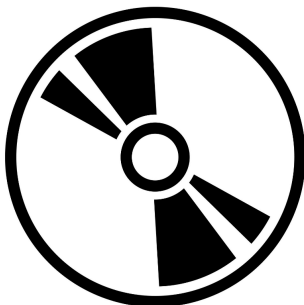
- What is the expected TC for the following queries?

Query-1 :     SELECT     \*  
  
                  FROM       students;

Query-2 :     SELECT     \*  
  
                  FROM       students  
  
                  WHERE     id = 100;

$$Tc = O(N)$$

- Where is all the data getting stored?



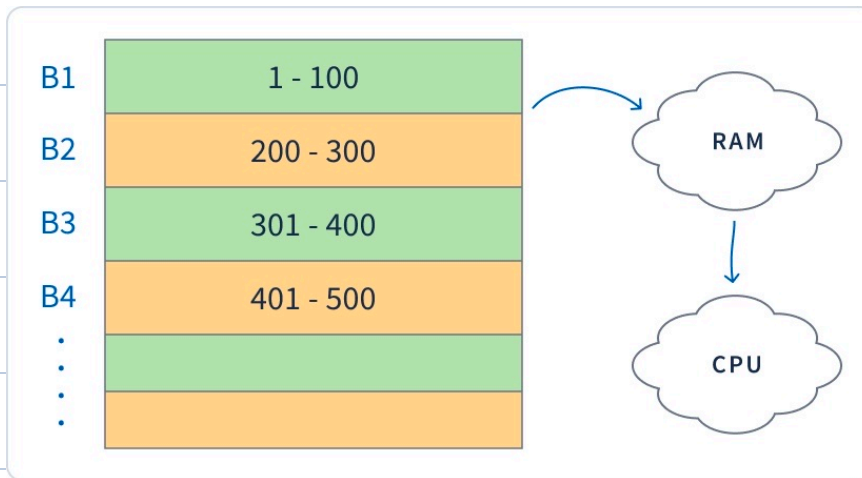


**Query-3 :**      **SELECT**      \*

**FROM**      students

**WHERE**      id = 500;

- The data of id = 500 is being stored in disk in a memory block as shown below :





- Hash Map

What is the TC here?

→ We can store indexes using Hashmap which takes  $O(1)$  to search.

id	block_address
1	B1
2	B1
⋮	⋮
500	B4
⋮	⋮
1050	B10

- How do we create indexing for following query cases?

Query :    **SELECT**        \*

**FROM**        students

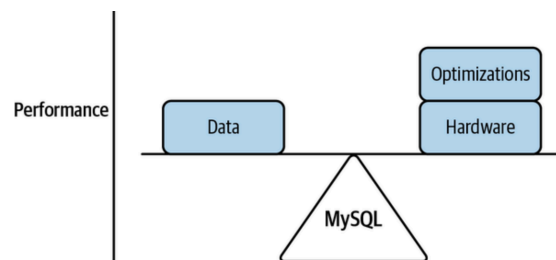
**WHERE**        psp = 80;

psp	memory block
1	[ B1, B5... ]
⋮	⋮
80	[ B2, B4... ]

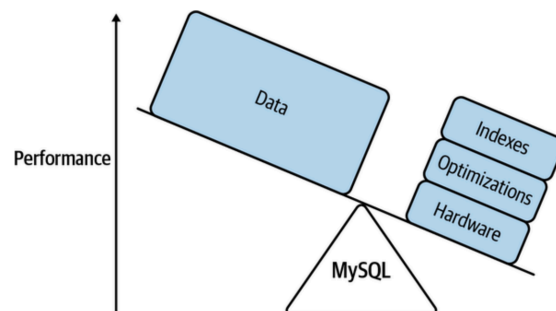


- Indexes makes query lightning fast

- Using indexes we can increase performance of our DB



Performance of SQL queries with indexing:





# Range Queries

```
SELECT      *  
FROM        students  
WHERE       psp between 40.2 and 60.5;
```

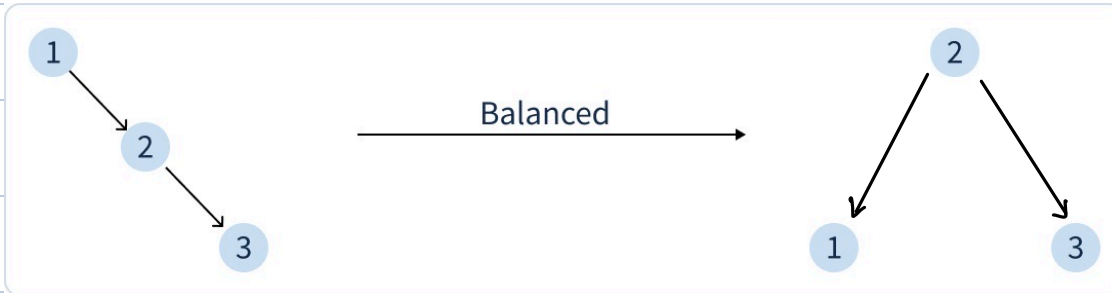


*' Can hash map work on range queries like above one?*

- If we use hash map to store indexing then we might need to exclusively check for every possible values in a range.



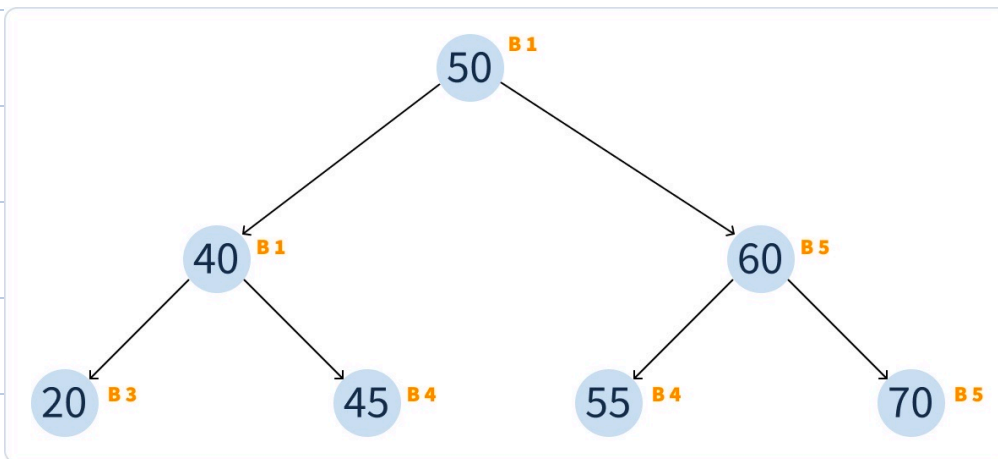
# Tree Map ( Self Balancing ) (BBST)



' What is TC here? '

$$TC = O(\log N)$$

## Example

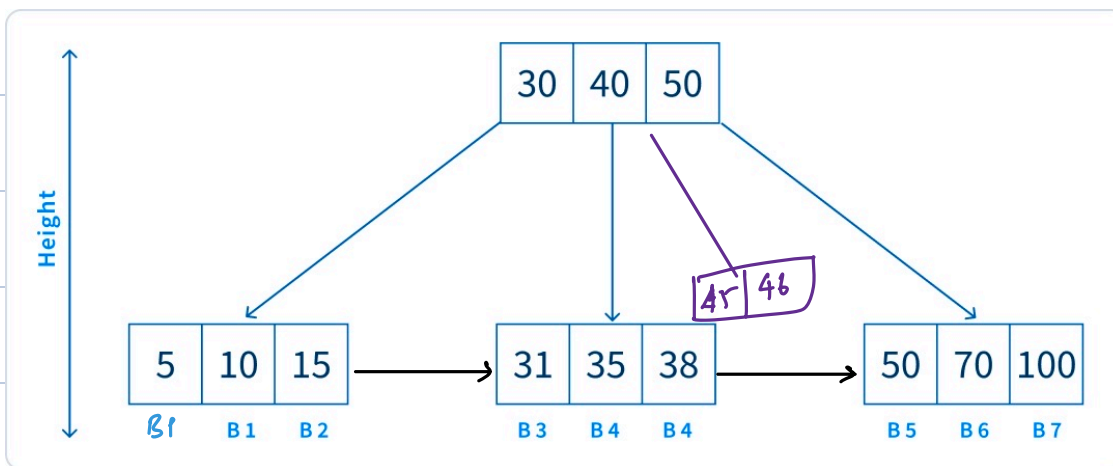


- Using tree maps we can use range queries easily.



## B Trees, B+ Trees

- We can have multiple children for every node.



- This is actual DS used by SQL.
- since we can have multiple children for every node, the overall height

of B+ Trees < Tree maps.

- $TC = O(H) < \text{Tree map.}$

$$(TC = O \log N)$$

→ By default indexing on PK : Clustered Indexing

→ Custom index on any other column : Non-clustered Indexing





# Cons of Indexing

'Where are these indexes getting stored?'



- Extra space required to store them.
- Write write / update / delete operations will be slower.

→ Cardinality : No. of unique values

→ Selectivity :  $\frac{\text{No. of unique values}}{\text{Total no. of values}} \rightarrow [0 - 1]$

→ 1 Billion

id	Name	Drinks
		Tea
		Coffee
		⋮

$$S = \frac{2}{10} \sim 0$$



# Indexing on multiple column

	id		name		psp		attendance	
--	----	--	------	--	-----	--	------------	--



# Indexing on Strings

- Let's say at Scaler we do a lot of queries email\_id column like :

```
SELECT      *  
  
FROM        students  
  
WHERE       email = "Naveen@gmail.com"
```



*' Do we create index on full email ? '*

\* Whenever we create indexing, we create it on a part of string.  
→ 4-6 chr.



"abc@gmail.com"

"Rahul@scaler.com"

"Naveen@gmail.com"

"badboy@gmail.com"

"abc@scaler.com"



' Can we create index on some part of this string ? '

- Creating index on part of a string will reduce space required

"abc@gmail.com"	81
"Rahul@scaler.com"	.
"Naveen@gmail.com"	.
"badboy@gmail.com"	.
"abc@scaler.com"	8500



abc : [81, 8500]



## 1. Clustered Indexing

- Table have default indexing on PK.
- Clustered.

## 2. Non - clustered Indexing

- Indexing created by us.































