# Welcome

Agenda: Log Basics
Quiz
Time Complexity
Big O
TLE
Constraints.

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## Logarithm

$\log_a b = c \implies$ to what value we need to raise 'a' so that its value becomes 'b'

$\log_a b = c \iff a^c = b$

$\log_2 64 = 2^? = 64 = 2^6 = 64$

$\log_3 27 = 3^? = 27 = 3^3 = 27$

$\log_2 2^{10} = 10$

$$2^K = N$$
$$\log_2 N = K$$

$$\log_a a^m = m$$

Q How many times we need to divide 9 by 2 till it reaches 1

$\Rightarrow$ $9/2 = 4$
$4/2 = 2$
$2/2 = 1$

$2^K = 9$
$K = \log_2 9$

Q How many times we need to divide 27 by 2 till it reaches 1

$27/2 = 13$

$13/2 = 6$ Ans = 4

$6/2 = 3$

$3/2 = 1$

---

## No. of Iterations.

```
N>0
i = N
while ( i >1)
{
    i = i/2
}
```

$\Rightarrow \quad i \to i/2 \to i/2^2 \to i/2^3 \cdots$

Ans = $\log_2 N$ $\quad O(\log N)$

Doing

```
for( i=1 ; i<N ; i=i*2)
{
    ....
}
```

| $i_b$ | iterations | $i_a$ |
|---|---|---|
| 1 | 1 | 2 |
| 2 | 2 | $2*2 = 4$ |
| 4 | 3 | $4*2 = 8$ |
| ⋮ | ⋮ | ⋮ |
| $2^{K-1}$ | K | $2^K$ |

$2^K = N \quad \Rightarrow \quad K = \log_2 N$

## Quiz 5

```
for( i=0 ; i<N ; i=i*2)
{
   ....
}
```

# iterations =

| $i_b$ | iterations | $i_a$ |
|---|---|---|
| 0 | 1 | 0 |
| 0 | 2 | 0 |
| 0 | 3 | 0 |
| . | ⋮ | |
| 0 | K | 0 |

## Quiz 6

```
for( i=1 ; i<=10 ; i++)
{
   for( j=1 ; j<=N ; j++)
   {
      ....
   }
}
```

# iteration's = 10N

O(N)

| i | j | # iteration |
|---|---|---|
| 1 | 1→N | N |
| 2 | 1→N | N |
| 3 | 1→N | N |
| ⋮ | ⋮ | |
| 10 | 1→N | N |
| | | 10N |

## Quiz 7

```
for( i=1 ; i<=N ; i++)
{
   for( j=1 ; j<=N ; j++)
   {
      ....
   }
}
```

# iteration = N*N

$= N^2$

$O(N^2)$

**Quiz 8**

```
for( i = 1 ; i <= N ; i++)
{
    for( j = 1 ; j <= N ; j = j * 2)
    {
        ....
    }
}
```

| i | j | # iteration |
|---|---|---|
| 1 | $1 \to \log N$ | $\log N$ |
| 2 | $1 \to \log N$ | $\log N$ |
| 3 | $1 \to \log N$ | $\log N$ |
| ⋮ | | |
| N | $1 \to \log N$ | $\log N$ |
| | | $N \log N$ |

# iterations = $N \log_2 N$ $\quad O(N \log N)$

**Quiz 9**

```
for( i = 1 ; i <= 4 ; i++)
{
    for( j = 1 ; j <= i ; j++)
    {
        ....
    }
}
```

| i | j | # iteration |
|---|---|---|
| 1 | [1, 1] | 1 |
| 2 | [1, 2] | 2 |
| 3 | [1, 3] | 3 |
| 4 | [1, 4] | 4 |
| | | 10 |

# iteration = 10 $\Rightarrow O(1)$

**Quiz 10**

```
for( i = 1 ; i <= N ; i++)
{
    for( j = 1 ; j <= i ; j++)
    {
        ....
    }
}
```

| i | j | # iteration |
|---|---|---|
| 1 | [1, 1] | 1 |
| 2 | [1, 2] | 2 |
| 3 | [1, 3] | 3 |
| 4 | [1, 4] | 4 |
| ⋮ | | ⋮ |
| N | [1, N] | N |
| | | $N(N+1)/2$ |

# iteration = $\frac{N(N+1)}{2}$

$\frac{N^2 + N}{2} \Rightarrow O(N^2)$

Quiz 11

```
for( i=1 ; i<=N ; i++)
{
    for( j=1 ; j<=(2^i); j++)
    {
        ....
    }
}
```

| i | j | # itera$^n$ |
|---|---|---|
| 1 | $[1, 2^1]$ | 2 |
| 2 | $[1, 2^2]$ | 4 |
| 3 | $[1, 2^3]$ | 8 |
| ⋮ | | |
| N | $[1, 2^N]$ | $2^N$ |

$2+4+8+\ldots\ldots 2^N$

10:45

# itera$^n$s = $2 + 4 + 8 + 16 + \ldots\ldots 2^N$

Sum of G.P = $\frac{a(r^N - 1)}{r-1}$

$= \frac{2(2^N - 1)}{2-1} = 2(2^N - 1) \Rightarrow 2^{N+1} - 2$

$\Rightarrow 2 * 2^N$

$O(2^N)$

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comparing itera$^n$s of 2 algo.

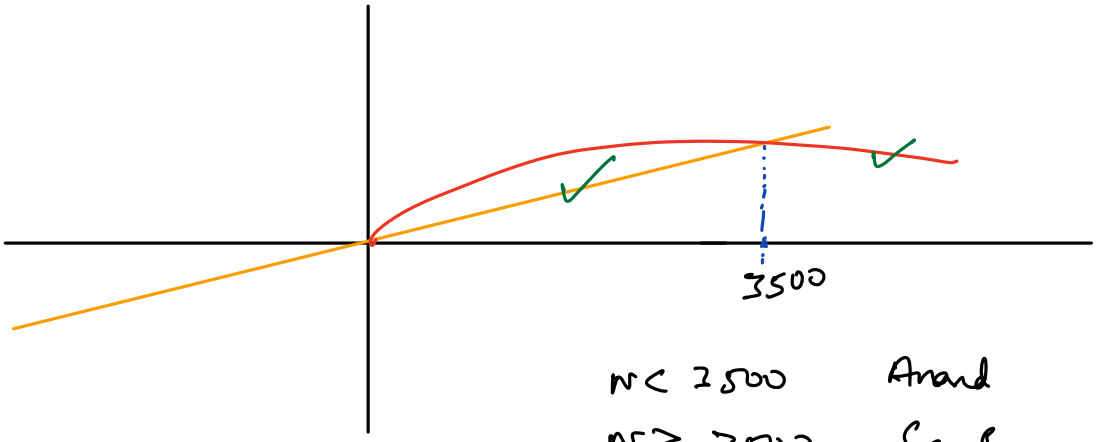
Sanjay

$100 \log_2 N$ (winner)

large inputs.

Anand.

$N/10$




$N < 3500$ Anand

$N > 3500$ Sanjay

obs → in real life, inputs can be very large.

# Asymptotic Analysis of Algorithm.

→ Big (O)

⇒ analysing performance of algo. for larger i/ps

⇒ Big (O)

Sanjay (winner) — Anand.

$100 \log_2 N$ ⇒ $N/10$

To calculate Big (O)

1) calculate iterations
2) Take higher order terms (neglect small order terms)
3) Ignore constant coefficients.

$100 \log_2 N \Rightarrow \log_2 N$

$N/10 \Rightarrow N$

Big O

$N > \log N$

eg: $4N^2 + 3N + 1$

$\Rightarrow N^2 \Rightarrow O(N^2)$

Quiz 12

$F(N) = 4N + 3N \log N + 1 \Rightarrow O(N \log N)$

Quiz 13

$F(N) = 4N \log N + 3N \text{ Sqrt}(N) + 10^6$

$\Rightarrow O(N \sqrt{N})$

$\log N < \text{sqrt}(N) < N < N \log N < N \text{ sqrt}(N) < N^2 < N^3 < 2^N < N! < N^N$

Neglect lower order terms. ?

Aniket's algo → $N^2 + 10N$

| I/P | # iteration's | % contribution of lower order terms. |
|---|---|---|
| $N = 10$ | 200 | 50% |
| $N \geq 100$ | $10^4 + 10^3$ | $\frac{10^3}{10^4 + 10^3} \simeq 9\%$ |
| $N = 10^4$ | $10^8 + 10^5$ | $\frac{10^5}{10^8 + 10^5} = 0.1\%$ |

Obs: as input size increases, contri$^n$ of lower order terms decreases.

Neglect constant coefficients ?

| Saketh | Aman | For larger i/p. |
|---|---|---|
| $10 \log N$ | $N$ | Saket |
| $100 \log N$ | $N$ | Saket |
| $10^3 \log N$ | $N$ | Saket |
| $10N$ | $N^2/10$ | Saket |
| $N \log N$ | $100N$ | Aman. |

---

Drawbacks of Big O

1) $N \log N$ $\quad$ $10^5 N$

⇒ When coefficient is very large, it can affect iteration's. But since we neglect coefficient while calculating Big O, it can sometime give incorrect comparison.

2)

Biswa

$2N^2 + 4N$

$O(N^2)$

Ramiz

$3N^2$

$O(N^2)$

Obs: Since Big(O) is same for both algo, we are not able to compare these 2 algo.

⇒ Sometimes, we cannot compare 2 algo. using Big(O)

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Time Limit Exceeded (TLE)

Ayush → online amazon interview

↳ reads → idea → code → TLE

Online Editors

1 GHz ⇒ $10^9$ instruc$^n$s / sec

Obs At max our code can have $10^9$ instruc$^n$

Pseudo code

```
bool func^n (int N)
{
    int c = 0 → 1
    i = 1 → 1 ; i <= N → 1 ; i++ → 1
    {
        if ( ____ ) → 1
            c++ → 1
    }
    return c → 1
}
```

6 - 7 instruc^n s

6N - 7N instruc^n s

Approach 1

Our code I itera^n → 10 instruc^n s.

Atmax ⟹ $10^9$ instruc^n s

$10 \times 10^8$ "

$10^8$ iterations

Approach 2

Our code I itera^n → 100 instruc^n s.

Atmax ⟹ $10^9$ instruc^n s

$100 \times 10^7$ "

$10^7$ iterations

⇒ In general, your code can have $10^6, 10^7 \sim 10^8$ iterations.

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Constraints

$$1 \leq N \leq 10^5$$

Manskant ⇒ T.C ⇒ $O(N^2)$

$N = 10^5$ ⇒ $10^{10}$ iterations. ⇒ no need to code.
→ optimize further.

Shreeram ⇒

$$1 \leq N \leq 10^3$$

T.C ⇒ $O(N^2)$

# iterations = $10^6$ ✓

Rohith

$$1 \leq N \leq 10^4$$

T.C $O(N^2)$

# iterations = $10^8$ ⇐ grey area
→ need to check.
→ not sure if I will get TLE.

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