### Investigating 3SAT

(Guide presentation for 380CT Coursework 2)

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#### **Notation**

Let  $x_1, x_2, ..., x_n$  be Boolean **variables**, and let  $\phi$  be a Boolean formula written in 3-cnf (Conjunctive Normal Form)

$$\phi = c_1 \wedge c_2 \wedge \cdots \wedge c_\ell,$$

where each **clause**  $c_m = x_i \lor x_j \lor x_k$ , for some i, j, k = 1, 2, ..., n and  $m = 1, ..., \ell$ .

A **literal** can be  $x_i$  or  $\neg x_i$  for some i = 1, 2, ..., n.

The ratio  $\ell/n$  is important for experiments, and will be denoted by  $\rho$ .

# Definition of the problem

#### **Decisional 3SAT**

Decide if  $\phi$  is satisfiable.

#### Computational/Search 3SAT

If  $\phi$  is satisfiable then find a satisfying assignment.

#### Optimization 3SAT (Max 3SAT)

Find an assignment that minimizes the number of non-satisfying clauses.

# Sampling strategy

General 3SAT instances will be generated by selecting literals from

$$\{x_1, \neg x_1, x_2, \neg x_2, \dots, x_n \neg x_n\}$$

uniformly at random.

For 'yes' instances, a random variable assignment is fixed first, then clauses are randomly constructed making sure each is satisfiable.

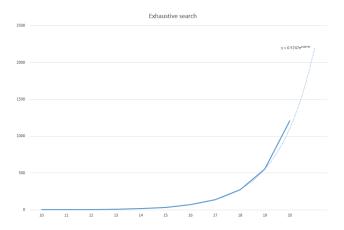
### Exhaustive search – theory

- 1: **for** all possible variable assignments of  $x_1, x_2, ..., x_n$  **do**
- 2: **if**  $\phi(x_1, x_2, \dots, x_n)$  evaluates to True **then**
- 3: **return** True
- 4: end if
- 5: end for
- 6: return False

There are  $2^n$  possible assignments, and each evaluation of  $\phi$  costs  $O(\ell)$ . So this algorithm costs

$$O(\ell 2^n)$$
.

### Exhaustive search – empirical results



Average time in  $100 \times$  seconds [TODO: REDO EXPERIMENT] for randomly generated instances with  $n = \ell$  for  $n = 10, \dots, 20$ . Dotted line: fitted exponential curve.

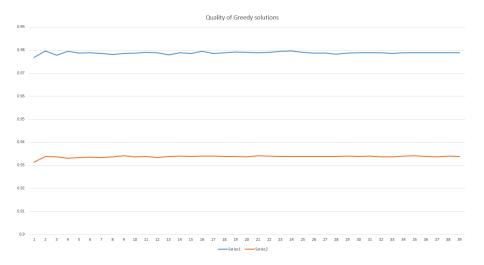
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# Greedy method

Find the variable that appears most often and assign it accordingly to maximize ...

- 1:  $L \leftarrow \emptyset$
- 2: **for**  $w \in \{x_1, \neg x_1, \dots, x_n, \neg x_n\}$  **do**
- 3: Count occurrences of w in  $\phi$
- 4: Append pair  $(w, count of occurrences of w in \phi)$  to L
- 5: end for
- 6: Sort *L* with respect to the second component
- 7: for  $(w, c) \in L$  do
- 8: Set w to True  $\triangleright$  If  $w = \neg x_i$  then set  $x_i$  to False
- 9: end for
- 10: return count of satisfied clauses

Cost:  $O(n \log n)$  assuming the use of an  $O(n \log n)$  sorting algorithm.



Average ratio of clauses satisfied by Greedy. [TODO: WRONG X-AXIS VALUES] Blue when  $\rho=1$  giving about 98%, and orange when  $\rho=10$  dropping to about 93%.

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#### References

- Hoos, H. and Stutzler, T. (2005) Stochastic Local Search: Foundations and Applications. Morgan Kaufmann
- Garey, S. and Johnson, D. (1979) Computers and Intractability: A Guide to the Theory of NP-Completeness. Freeman