

Algorithms for Programming Contests SS20 - Week 11

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A: Contact List - Sample Solution

- ▶ Build a trie from all possible contact names.
- ▶ Recursively visit the trie, starting at the root node.
- ▶ If a node marks a word end and has children, it has to be renamed.
- ▶ Count nodes that match this criterion.

B: J&J's - Sample Solution

- ▶ Toss J&Js into n glasses standing in a line.
- ▶ i-query: Add v J&Js to all glasses between ℓ and r inclusively.
- ▶ q-query: Return the number of J&Js in glass a .

B: J&J's - Sample Solution

- ▶ Operations needed: range update & point query
- ▶ Solution:
 - ▶ Use a segment tree with lazy propagation.
 - ▶ i-query: Add value v to interval $[\ell, r]$ using lazy propagation.
 - ▶ q-query: Sum up interval $[a, a]$.
 - ▶ Caveat: Without lazy propagation adding J&Js to every glass takes $\mathcal{O}(n \log(n))$.

C: Broken Tetris - Sample Solution

- ▶ Given a broken Tetris game which only considers rectangular pieces, calculate the maximal height at each stage of the game.

C: Broken Tetris - Sample Solution

- ▶ Build a segment tree over the playing field.
- ▶ Use $(\mathbb{N}, \max, 0)$ as monoid.
- ▶ For a new piece $w \times h$ inserted at position p
 - ▶ Determine the maximal height h^* on interval $[p, p + w - 1]$.
 - ▶ Set the maximum on range $[p, p + w - 1]$ to $h^* + h$.
- ▶ Lazy propagation is avoidable as pieces have width ≤ 100 .
(might depend on the language)

D: Treehouse - Sample Solution

Problem

Given a rooted tree, find the shortest path to visit a sequence of nodes

Solution

Let $\text{level}(x)$ be the distance from x to the root.

Then, the shortest path between nodes x, y in a tree is given by

$$\text{level}(x) + \text{level}(y) - 2 * \text{level}(\text{lca}(x, y))$$

- ▶ Precompute the level values.
- ▶ Compute LCA in logarithmic time per query, i.e. by using Segment Trees.

E: Ghost - Sample Solution

- ▶ Given a dictionary of words, find a strategy to win the n -th round of *Ghost*

E: Ghost - Sample Solution

- ▶ Main idea: Build a trie from the dictionary.
- ▶ Minimax algorithm:
 - ▶ Starting from the leaves, mark nodes if Lea can win or lose if she starts.
- ▶ Improvement: Ignore words where a prefix is already in the trie.

E: Ghost - Sample Solution

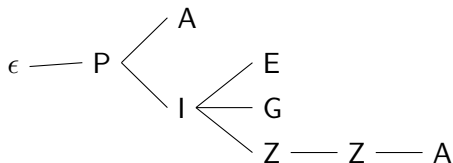
- ▶ Consider scenario 1:
 - ▶ Lea begins in round 1, the winner of a round begins the next.
 - ▶ Then Lea can win the last round iff either
 - ▶ she can win from the root node or
 - ▶ she can lose from the root node and n (the amount of played rounds) is even.
- ▶ Consider scenario 2:
 - ▶ Lea begins, but the loser of a round begins the next.
 - ▶ Then Lea can win the last round iff either
 - ▶ she can win and lose from the root node or
 - ▶ she can win from the root node and n (the amount of played rounds) is odd.
- ▶ Scenarios 3 and 4 are dual to 1 and 2.

E: Ghost - Sample Solution

- ▶ A node can be won if it
 - ▶ is a leaf node or
 - ▶ has a child that can not be won.
- ▶ A node can be lost if it
 - ▶ is not a leaf node and
 - ▶ has a child that can not be lost.
- ▶ Check for all nodes if they can be won or lost.
- ▶ Check what happens:
 - ▶ If root can be won and lost, Lea can force any outcome.
 - ▶ If root can not be won, but can be lost, Lea can win if n is even.
 - ▶ If root can be lost, but can not be won, Bea can win if n is odd.
 - ▶ If root can not be won and not be lost, Bea can force any outcome.

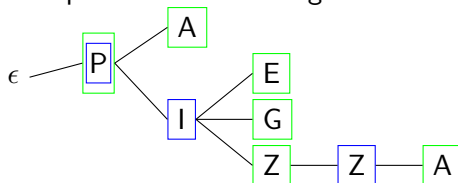
E: Ghost - Sample Solution

► Example trie:



E: Ghost - Sample Solution

- ▶ Example trie with labelling:



- ▶ Green labels: Can be won.
- ▶ Blue labels: Can be lost.
- ▶ Here, Bea can force any outcome.

E: Ghost - Sample Solution

Common Difficulties

- ▶ Organizing the trie through individual structs for each node has some overhead and bad cache efficiency.
- ▶ One of the words in the dictionary might be the prefix of another word. In this case, the game ends already when the shorter word is completed.

F: Catmon Go - Sample Solution

- ▶ Given spawn, despawn and catch events
- ▶ Find the sum of all Catmon in range of a catch event

F: Catmon Go - Sample Solution

- ▶ Build a segment tree over the locations, keeping track of the number of Catmon in a certain range.
- ▶ Spawn event: $+1$ at location a .
- ▶ Despawn event: -1 at location a if number is greater than 0. First recurse down the tree and then propagate change upwards (if there is any).
- ▶ Catch event: Sum in interval $[l, r]$, then set all numbers in this interval to 0. Again propagate change from full intervals upwards and use lazy propagation to later apply change to lower intervals.
- ▶ Caveat: Use fast input

F: Catmon Go - Sample Solution

- ▶ Without segment tree:
 - ▶ Keep a multiset of all Catmons ordered by their positions.
 - ▶ Insert on spawn and remove on despawn events.
 - ▶ Remove every single Catmon in the range on catch events one after another.
 - ▶ Since every Catmon in this range must have been inserted at some point, the amortized runtime of catch events is bounded.