

# **Aula 6 – Sincronismo 2**

# Exclusão Mútua

- Variável de impedimento
- Locks e Rlocks
  - Lock → qualquer thread que tentar adquiri-lo irá bloquear, mesmo se o mesmo segmento em si já estiver segurando o bloqueio.
  - Nesses casos, RLock (bloqueio de reentrada) é usado.

# Exclusão Mutua

- Lock X Rlock

```
1  import threading
2
3  num = 0
4  lock = threading.Lock()
5
6  lock.acquire()
7  num += 1
8  lock.acquire() # This will block.
9  num += 2
10 lock.release()
```

```
1  import threading
2
3  # With RLock, that problem doesn't happen.
4  lock = threading.RLock()
5
6  lock.acquire()
7  num += 3
8  lock.acquire() # This won't block.
9  num += 4
10 lock.release()
11 lock.release()
```

# Exclusão Mútua

- Com alternância estrita

```
while (TRUE) {  
    while (turn != 0)      /* loop */ ;  
    critical_region();  
    turn = 1;  
    noncritical_region();  
}
```

(a)

```
while (TRUE) {  
    while (turn != 1)      /* loop */ ;  
    critical_region();  
    turn = 0;  
    noncritical_region();  
}
```

(b)

# Exclusão Mútua

- Solução de Peterson

```
#define N      2                                /* number of processes */

int turn;                                       /* whose turn is it? */
int interested[N];                             /* all values initially 0 (FALSE) */

void enter_region(int process);                /* process is 0 or 1 */
{
    int other;                                /* number of the other process */

    other = 1 - process;                      /* the opposite of process */
    interested[process] = TRUE;                /* show that you are interested */
    turn = process;                            /* set flag */
    while (turn == process && interested[other] == TRUE) /* null statement */ ;
}

void leave_region(int process)                  /* process: who is leaving */
{
    interested[process] = FALSE;               /* indicate departure from critical region */
}
```

# Exclusão Mútua

- Instrução TSL
  - Auxílio do hardware

enter\_region:

TSL REGISTER, LOCK

CMP REGISTER, #0

JNE enter\_region

RET

| copy lock to register and set lock to 1

| was lock zero?

| if it was nonzero, lock was set, so loop

| return to caller; critical region entered

leave\_region:

MOVE LOCK, #0

RET

| store a 0 in lock

| return to caller

# Exclusão Mutua

- Soluções anteriores
  - com espera ocupada
- Soluções ideais
  - Com sleep e wakeup

# Semáforos

- E.W. Dijkstra (1965) → variável inteira
  - Operação **down**: se maior que zero, decrementa
  - Operação **up**: se menor que o máximo, incrementa
- Up e down são generalizações de acquire e release
- Ações atômicas



# Mutex

- Mutex → Mutual Exclusion
  - Semáforo binário

mutex\_lock:

TSL REGISTER,MUTEX	copy mutex to register and set mutex to 1
CMP REGISTER,#0	was mutex zero?
JZE ok	if it was zero, mutex was unlocked, so return
CALL thread_yield	mutex is busy; schedule another thread
JMP mutex_lock	try again
ok: RET	return to caller; critical region entered

mutex\_unlock:

MOVE MUTEX,#0	store a 0 in mutex
RET	return to caller

