

# Stochastic Modeling and Optimization Problemset 3

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## Problem 1

Problem variables

$P_l$  = probability of large demand

$P_s$  = probability of small demand

$x_k$  = current inventory level

$w_k$  = demand drawn from  $P_l$  or  $P_s$

$q$  = apriori probability of large demand

In the standard inventory control problem, for every overage or underage in stock given demand, there are associated holding and shortage costs:

$$r(x) = p(\max(0, -x)) + h(\max(0, x))$$

If there is a non-negative  $x$ ,  $h$  denotes the unit holding cost. If there  $x$  is negative, there is an opportunity cost of  $p$  per unit of demand not met.

The minimized cost function is:

$$\min_{u \geq 0} \left\{ cu + \mathbb{E}_w \{ p \max(0, w - x - u) + h \max(0, x + u - w) \} \right\}$$

where  $c$  is the cost of ordering level  $u$ . E.g. it is optimized by the  $u$  minimizing the holding cost plus the overage cost and  $y$  is the indicator in  $\{1, 2\}$  indicating the realization of large or small demand.

### Part 1 Solution

Let  $y_i \in \{1, 2\}$  denote the realization of large or small demand and  $D^i$  the level of demand. The single-period optimal choice is the same for either large or small demand:

$$\mu^*(x, y) = \begin{cases} D^i - x, & \text{if } y = i \text{ and } x < D^i, \\ 0, & \text{otherwise.} \end{cases}$$

### Part 2 Solution

The augmented system for the multi-period problem is:

$$x_{k+1} = x_k + u_k - w_k$$

$$y_{k+1} = \psi_k$$

Where  $\psi_k$  takes on values 1 and 2 indicating a realization of demand from  $P_l$  and  $P_s$  respectively and taking on value 1 with probability  $q$  and 2 with probability  $1 - q$

The DP algorithm is:

1. Terminal cost is:

$$J_N(x_N, y_N) = 0$$

2. Choose the  $u_k$  minimizing the cost function at time  $k$  given by:

$$\begin{aligned} \text{cost function} &= cu_k + \{p\max(0, w_k - x_k - u_k) + h\max(0, x_k + u_k - w_k)\} \\ J_k(x_k, y_k) &= \min_{u \geq 0} \{ \mathbb{E}_{w_k} \left\{ \text{cost function} + qJ_{k+1}(x_k + u_k - w_k, 1) + (1 - q)J_{k+1}(x_k + u_k - w_k, 2) \right\} | y_k \} \end{aligned}$$

$$\mu_k^*(x_k, y_k) = \begin{cases} D_k^i - x_k, & \text{if } y_k = i, x_k < D_k^i, \\ 0, & \text{otherwise.} \end{cases}$$

## Problem 5

$p_i$  : probability of answering question  $i$  correctly  
 $r_i$  : reward for answering question  $i$  correctly  
 $F_i$  : cost for answering question  $i$  incorrectly

The expected value of question  $i$ :

$$R_i = \mathbb{E}[q_i] = p_i R_i - (1 - p_i) F_i$$

*Part 1 Solution*

Claim the optimal ordering of the set of questions is  $L$  and questions  $i$  and  $j$  are the  $k^{th}$  and  $(k + 1)^{st}$  questions:

$$L = \{i_0, \dots, i_{k-1}, i, j, i_{k+2}, \dots, i_{N-1}\}$$

If we consider the list with  $i$  and  $j$  interchanged:

$$L' = \{i_0, \dots, i_{k-1}, j, i, i_{k+2}, \dots, i_{N-1}\}$$

The expectation of the reward for  $L$  is greater than  $L'$  by definition, so it follows that:

$$p_i R_i - (1 - p_i) F_i / (1 - p_i) \geq p_j R_j - (1 - p_j) F_j / (1 - p_j)$$

*Part 2 Solution*

If there is a no cost option to stop answering questions, the contestant will stop answering questions whenever the expected value of answering the next question is negative.

The game stops at period  $k - 1$ , with  $k$  satisfying:

$$p_k R_k < (1 - p_k) F_k$$

That is, the game stops whenever the expected value of answering the next question becomes negative.