MECHTRON 2MD3

Data Structures and Algorithms for Mechatronics Winter 2022

19 Queues

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Stack Example - Parentheses Matching Problem

Each "(", "{", or "[" must be paired with a matching ")", "}", or "["

```
    correct: ( )(( )){([( )])}
    correct: ((( )(( )){([( )])}
    incorrect: )(( )){([( )])}
    incorrect: ({[ ])}
    incorrect: (
```

Stack Example - Parentheses Matching Problem

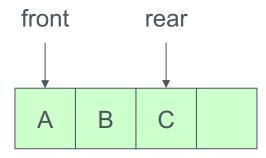
Each "(", "{", or "[" must be paired with a matching ")", "}", or "["

```
correct: ( )(( )){([( )])}
correct: ((( )(( )){([( )])}
incorrect: )(( )){([( )])}
incorrect: ({[ ])}
incorrect: (
```

```
Algorithm ParenMatch(X, n):
   Input: An array X of n tokens, each of which is either a grouping symbol, a
      variable, an arithmetic operator, or a number
   Output: true if and only if all the grouping symbols in X match
    Let S be an empty stack
    for i \leftarrow 0 to n-1 do
      if X[i] is an opening grouping symbol then
         S.\mathsf{push}(X[i])
      else if X[i] is a closing grouping symbol then
         if S.empty() then
           return false
                                {nothing to match with}
         if S.top() does not match the type of X[i] then
                                {wrong type}
           return false
         S.pop()
    if S.empty() then
      return true
                          {every symbol matched}
    else
                           {some symbols were never matched}
      return false
```

Queue

- A queue is a container of elements that are inserted and removed according to the First-In First-Out (FIFO) principle.
- Elements enter queue at the **rear** and *are removed* from the **front** "push" adds a new item on top of the stack.



Queue and its Applications

- Applications
 - Waiting lists, bureaucracy
 - Access to shared resources (e.g., printer)
 - Multiprogramming
 - As a data structure for algorithms to solve problems
 - BFS(Breadth First Search) for graphs
 - Component of other data structures



The Queue ADT

- The Queue ADT stores arbitrary objects
- Main stack operations:
 - enqueue(e): Insert element e at the rear of the queue
 - dequeue(): Remove element at the front of the queue; an error occurs if the queue is empty.
 - front(): Return, but do not remove, a reference to the front element in the queue; an error occurs if the queue is empty.
 - size(): Return the number of elements in the queue
 - empty(): Return true if the queue is empty and false otherwise.

```
template <typename E>
class Queue {
    public:
    int size() const;
    bool empty() const;
    const E& front() const throw(QueueEmpty);
    void enqueue (const E& e);
    void dequeue() throw(QueueEmpty);
};

// an interface for a queue
// number of items in queue
// is the queue empty?
// the front element
// enqueue element at rear
// dequeue element at front
};
```

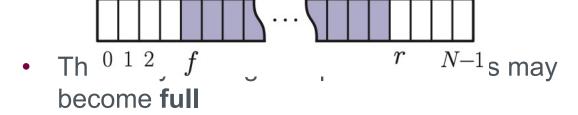
Exceptions for Queue ADT

- In the Queue ADT, operations dequeue and front cannot be performed if the Queue is empty
- Attempting dequeue or front on an empty queue throws a QueueEmpty exception

```
class QueueEmpty : public RuntimeException {
public:
   QueueEmpty(const string& err) : RuntimeException(err) { }
};
```

Array-Based Queue

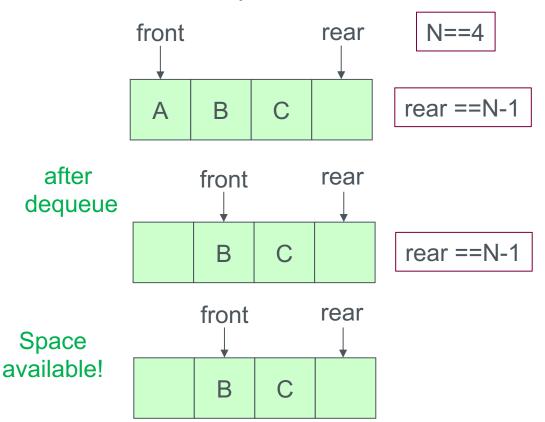
- An array with capacity N stores elements
- Two variables keeps track of the index of the front and rear element;
- f: index of front element
- r: index of element following rear element



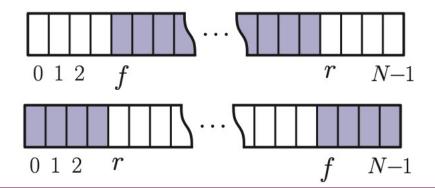
 A enqueue operation will then throw a QueueFull exception

Array-Based Queue - An Issue

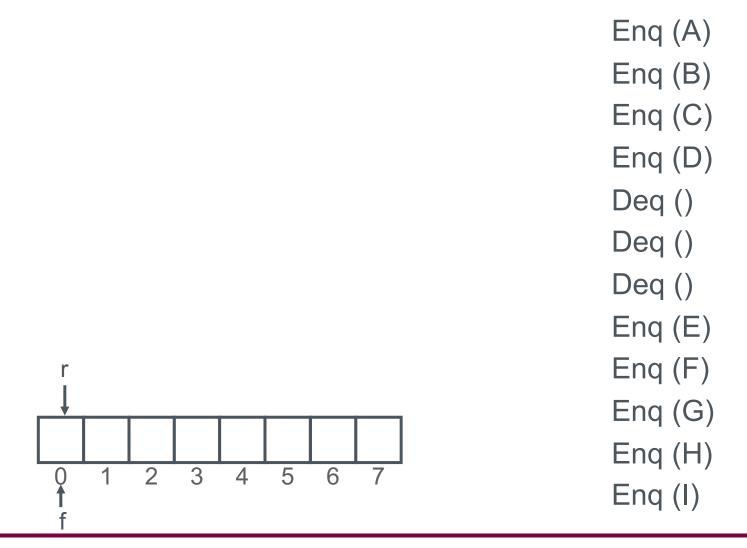
- **f**: index of front element
- r: index of element following rear element
- We can run into problem

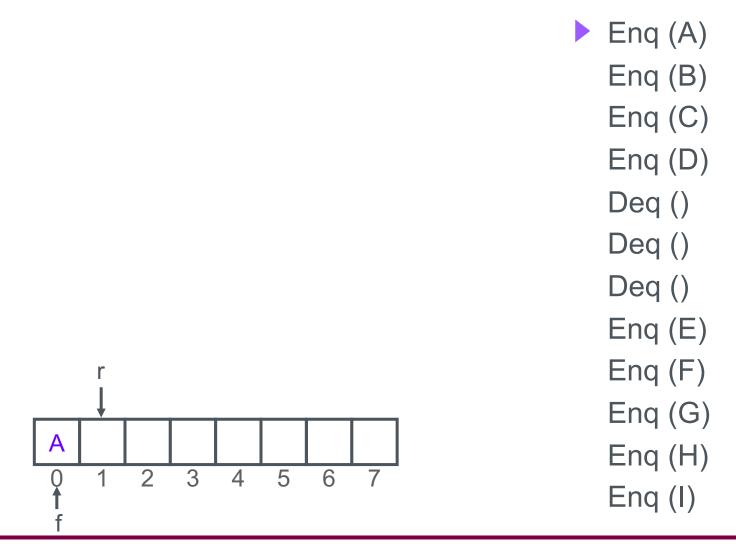


- An array with capacity N stores elements
- f: index of front element
- r: index of element following rear element
- Each time we increment f or r, we simply need to compute this increment as:
 - 。 (f + 1) mod N
 - 。 (r + 1) mod N
- modulo operator in C++ is %

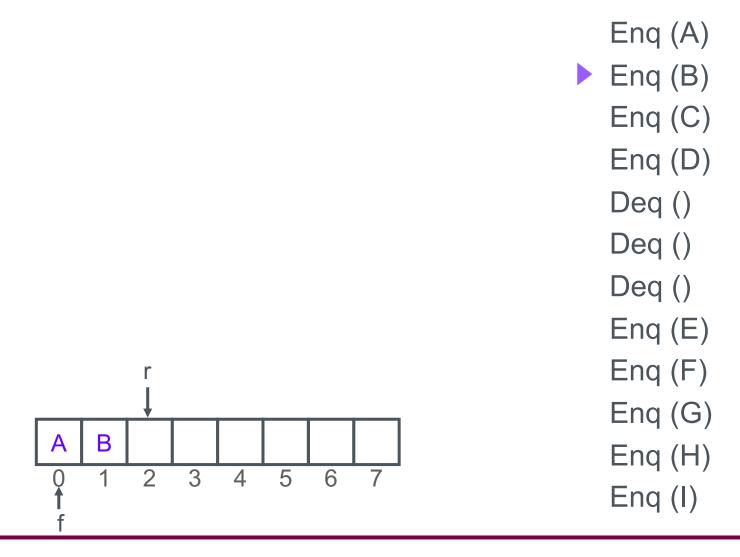


```
Algorithm size():
    return n
Algorithm empty():
    return (n = 0)
Algorithm front():
    if empty() then
      throw QueueEmpty exception
    return Q[f]
Algorithm dequeue():
    if empty() then
      throw QueueEmpty exception
    f \leftarrow (f+1) \bmod N
   n = n - 1
Algorithm enqueue(e):
    if size() = N then
      throw QueueFull exception
    Q[r] \leftarrow e
    r \leftarrow (r+1) \bmod N
   n = n + 1
```

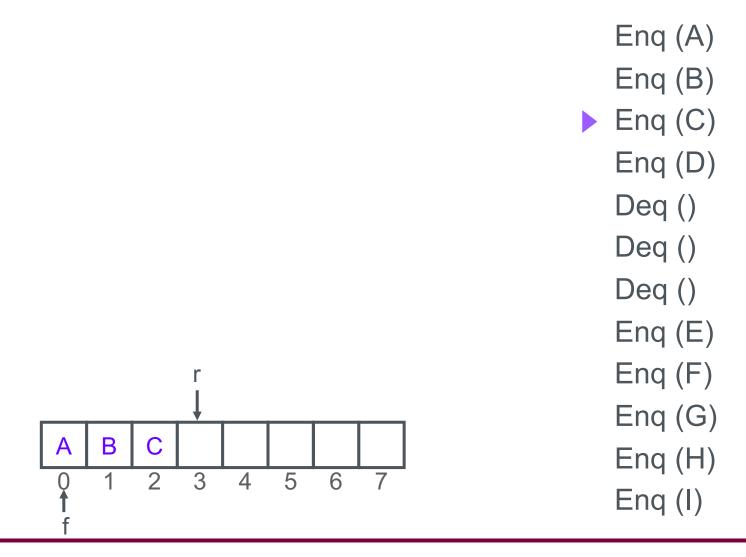




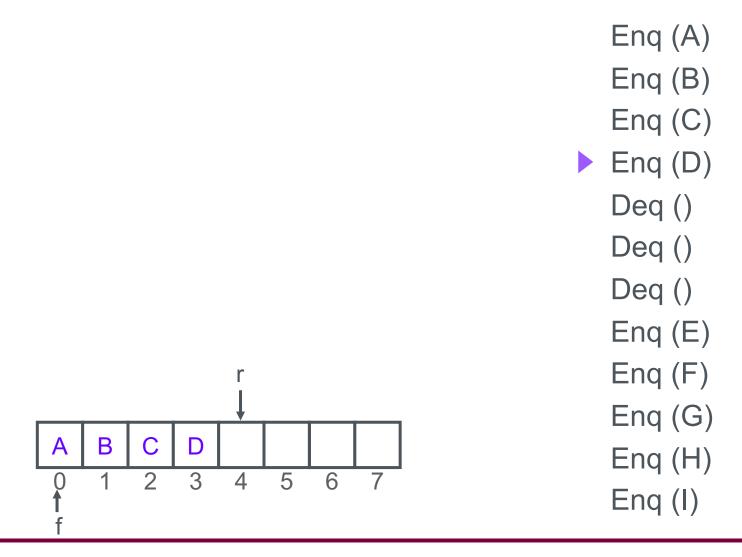




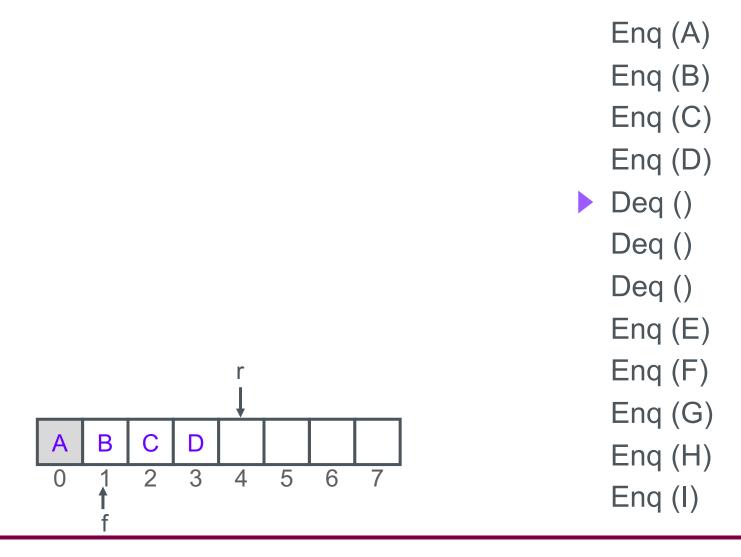




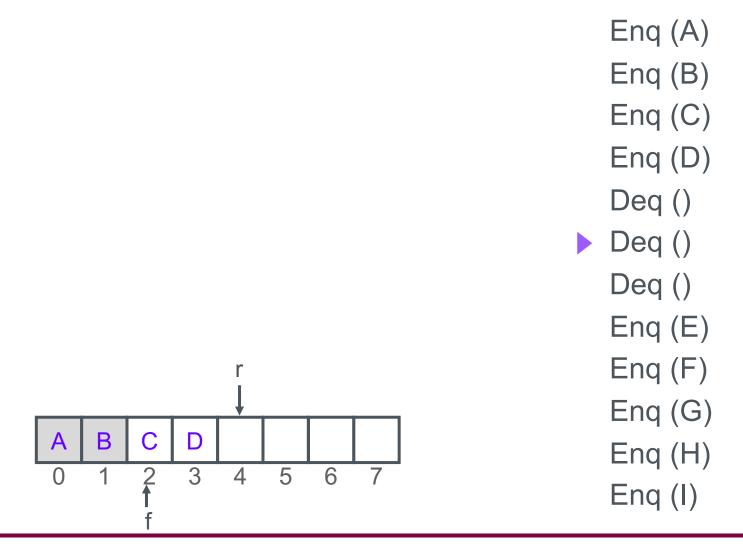




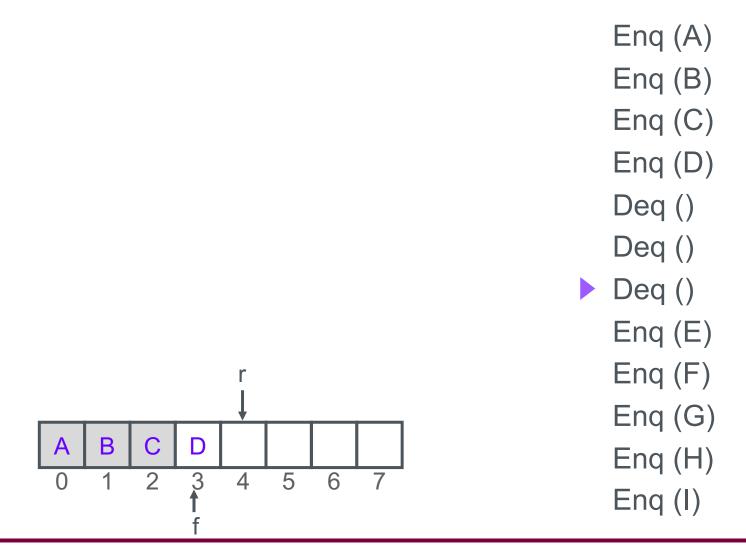




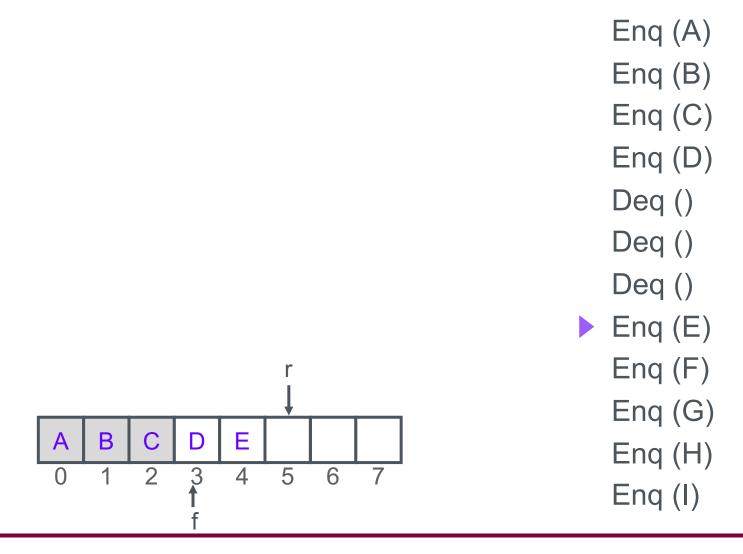




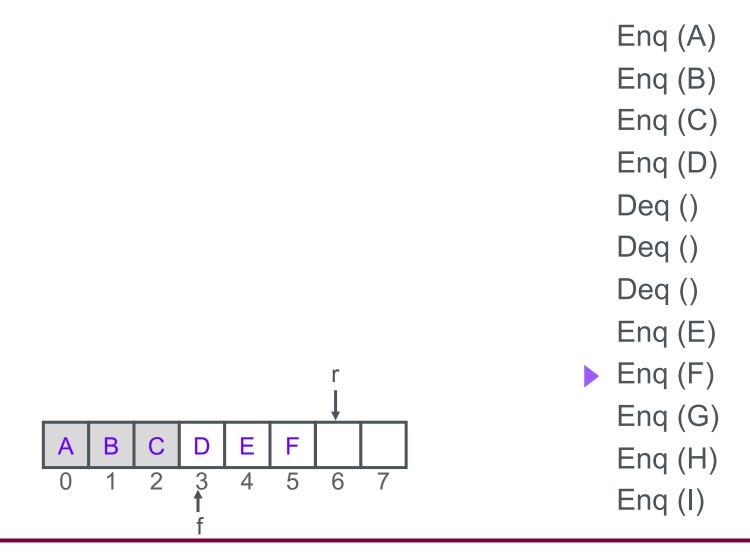


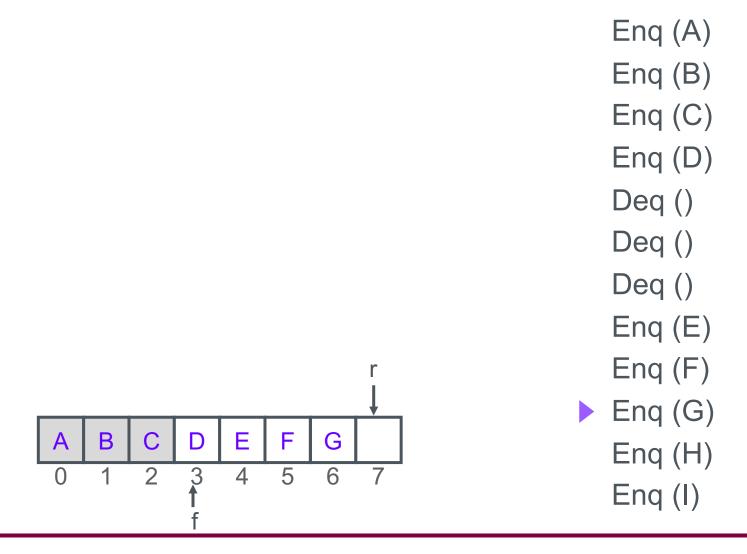




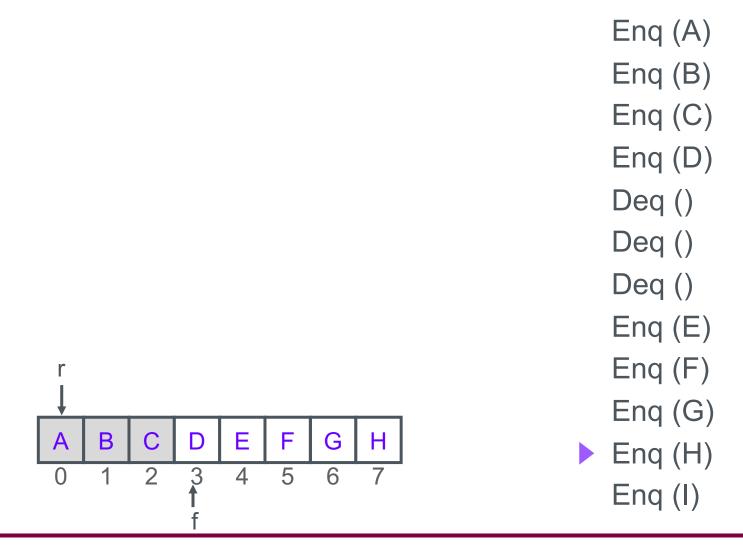




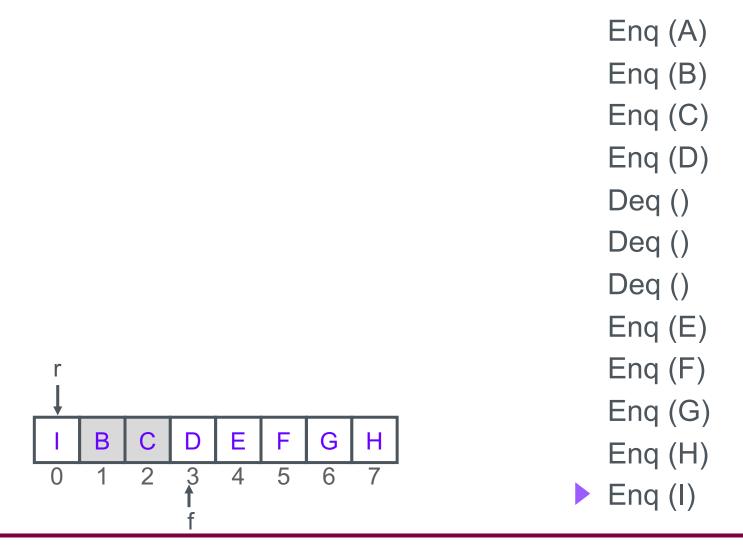












Array-Based Queue - Performance

- Performance:
 - Let n be the number of elements in the Queue
 - The space used is O(N)
 - this is independent of the number of elements n
 - Each operation runs in time O(1)

- Limitations:
 - The maximum size of the Queue must be defined beforehand and cannot be changed
 - Trying to enqueue a new element into a full queue causes an implementation-specific exception (QueueFull Exception)



- We use that CircleList class
 - C stores the Queue elements
 - n stores the number of elements on Queue

```
typedef string Elem;
                                           // queue element type
                                           // queue as doubly linked list
class LinkedQueue {
public:
 LinkedQueue();
                                           // constructor
 int size() const;
                                           // number of items in the queue
  bool empty() const;
                                           // is the queue empty?
 const Elem& front() const throw(QueueEmpty); // the front element
 void enqueue(const Elem& e);
                                           // enqueue element at rear
  void dequeue() throw(QueueEmpty);
                                           // dequeue element at front
private:
                                           // member data
                                           // circular list of elements
  CircleList C:
                                           // number of elements
 int n:
```

```
// enqueue element at rear
void LinkedQueue::enqueue(const Elem& e)
  C.add(e);
                                                     // insert after cursor
  C.advance();
                                                     // ...and advance
  n++;
                       (front)
                                              (rear)
                                  MSP
                        LAX
                                     (a)
                                                cursor
                 (front)
                                         (rear)
                  LAX
                             MSP
                                         ATL
                                     (b)
                                                           cursor
                 (front)
                                                    (rear)
                  LAX
                             MSP
                                        ATL
                                                    BOS
                                     (c)
```

```
// dequeue element at front
void LinkedQueue::dequeue() throw(QueueEmpty) {
 if (empty())
   throw QueueEmpty("dequeue of empty queue");
 C.remove();
                                             // remove from list front
 n--;
                                                          cursor
             (front)
                                                  (rear)
             LAX
                       ► MSP
                                      ATL
                                                  BOS
                                  (a)
                                                         cursor
                        (front)
                                                 (rear)
                                                 BOS
                                  (b)
```

Questions?