

Mechatronics/Software Engineering 4AA4 Midterm Exam

Monday, October 17, 2022
McMaster University (CAS)

Duration: 80 Minutes
Instructor: Dr. Wenbo He

Student Name: Peng Cui

ID: 400231347

This test paper includes 11 pages. You are responsible for ensuring that your copy of the test is complete. Bring any discrepancy to the attention of your invigilator.

Special Instructions:

1. Put all answers on the test paper. Try to keep your answers brief. Use point form if need be.
2. It is an open-note exam.
3. The burden of communication is upon you. Solutions not properly explained, will not be considered correct. If we cannot decode "what you wrote, we cannot grade it as a correct answer.
4. For any re-read (re-mark) to occur, answers must be in pen.
5. You may need the following Table of the value of $n(2^{1/n} - 1)$

N	$n(2^{1/n} - 1)$
1	1
2	0.828
3	0.780
4	0.757
5	0.743
6	0.735
7	0.729
8	0.724
9	0.721
10	0.718
∞	0.693

Grading

Section	Grade
Question #1 – Multiple Choice	20 points <u>18.5</u>
Question #2 – Short Answer Questions	10 points <u>10</u>
Total:	30points <u>28.5</u>

Student Name: _____

ID: _____

Question 1: Multiple Choice Questions (20 points)

A 1) Given the task set: $T_1(4, 1)$; $T_2(5, 1)$; $T_3(10, 2)$, and you are asked to find a Cyclic Executive (CE) schedule using flow graph. Assume that the frame size is 2. Which of the following is true?

- A. There are 11 job nodes and 10 frame nodes in the flow graph.
- B. There are 14 job nodes and 11 frame nodes in the flow graph.
- C. Using the network flow to model CE scheduling cannot handle the task split in CE.
- D. The algorithms finding the solution to network flow problem is NP-complete.

B 2) The _____ scheduling algorithm schedules periodic tasks using a static priority policy with preemption.

- A. earliest deadline first *EDF dynamic*
- B. rate monotonic
- C. first come first served *X*
- D. priority *X*

C 3) If a set of processes cannot be successfully scheduled by rate monotonic scheduling algorithm, then:

- A. they can be scheduled by EDF algorithm *X*
- B. they cannot be scheduled by EDF algorithm
- C. they cannot be scheduled by any other algorithm
- D. None of the above

EDF

RM

A 4) Consider the task set $T = \{(8, 4), (10, 2), (12, 3)\}$, which of the following is true:

- A. T is both RM and EDF schedulable.
- B. T is not RM schedulable but EDF schedulable.
- C. T is neither RM nor EDF schedulable.
- D. T is not EDF schedulable but RM schedulable.

Student Name: _____

ID: _____

A 5) A Task set consists of n pre-emptive and periodic tasks. If the task is NOT RM schedulable which of the following can be inferred?

- A. The CPU utilization is over $0.693 \times$
- B. The task set is DM schedulable \times
- C. The task set is EDF schedulable \times
- ☒ D. The CPU utilization is over 1

D 6) What is the output of g value if the following code is given:

```
#include <stdio.h>
```

```
#include <pthread.h>
```

```
int g = 0;
```

```
void *aThread()
```

```
{
```

```
    g++;
```

```
    pthread_exit(NULL);
```

```
}
```

```
int main (int argc, char *argv[])
```

```
{
```

```
    int i;
```

```
    pthread_t thread[3];
```

```
    for (i=0; i<3; i++)
```

```
    {
```

```
        if( pthread_create( thread+i, NULL, aThread, NULL) )
```

```
        {
```

```
            printf("ERROR; return code from pthread_create()\n");
```

```
            return -1;
```

```
        }
```

```
    }
```

```
    printf("The value of  $g$  is %d\n", g);
```

```
    return 0;
```

```
}
```

A. The value of g is 1

B. The value of g is 2

C. The value of g is 3

D. There is a race condition, so the value of g cannot be determined.

C 7) Which of the following is a drawback of thread programming

A. Context switch time among different threads can be large.

B. Tasks implemented by threads are less efficient than tasks implemented by processes.

C. Without a synchronization, race condition on shared variables can be disastrous.

D. It does not share code among threads in the same process.

ID:

2X

```
#include <stdio.h>
#include <sys/types.h>
int main()
{
    fork();
    fork();
    fork();
    printf("hello\n");
    return 0;
}
```

A hand-drawn diagram consisting of two rows of four rectangles each. The rectangles are arranged in a 2x4 grid. The top row has four rectangles, and the bottom row has four rectangles. The rectangles are drawn with simple lines, and the entire diagram is enclosed in a larger rectangular frame.

- ```
void *even(void *arg)
{
 printf("This is even thread()\n");
```



Student Name: \_\_\_\_\_

ID: \_\_\_\_\_

```
while(count < MAX)
 if(count % 2 == 0)
 printf("%d ", count++);
 pthread_exit(0);
}

void *odd(void *arg)
{
 printf("This is odd thread()\n");
 while(count < MAX)
 if(count % 2 == 1)
 printf("%d ", count++);
 pthread_exit(0);
}

int main()
{
 pthread_t t1;
 pthread_t t2;

 pthread_create(&t1, 0, &even, NULL);
 pthread_create(&t2, 0, &odd, NULL);

 pthread_join(t1, 0);
 pthread_join(t2, 0);

 return 0;
}
```

- A. 1 0 2 3 4 5 6 7 9 8  
B. 1 0 2 4 3 5 6 7 8 9  
C. 0 1 2 3 4 6 5 7 8 9  
D. 1 0 2 4 5 3 6 7 9 8  
D | 2 3 4 5 6 7 8 9

- C 12) In cyclic executive scheduling, which of the following statement is correct.
- A. CE scheduling algorithm is based on static priority of tasks. A task's priority is inverse proportional to its period. X
  - B. A drawback of CE scheduling algorithm is that you have to compute an offline schedule which can be arbitrarily long.
  - C. The frame size of the CE algorithm cannot be too small since we want an instance of a task is completed within a single frame. ✓
  - D. The frame size of the CE algorithm cannot be too large, otherwise the computational complexity is too high. X

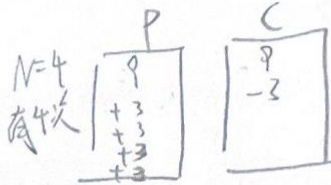
Student Name:

- D 13) What is the Output of the following program? Note that the SIGCHLD signal is sent to the parent of a child process when it exits, is interrupted, or resumes after being interrupted.

```
#include<stdio.h>
#include<signal.h>
#include<sys/wait.h>
```

```
#define N 4
int val = 9;
void handler(sig) {
 val += 3;
 return;
}
```

```
int main() {
 pid_t pid;
 int i;
 signal(SIGCHLD, handler);
 for (i=0; i<N; i++) {
 if ((pid = fork()) == 0) {
 val -= 3;
 exit(0);
 }
 }
 for (i=0; i<N; i++) {
 waitpid(-1, NULL, 0); //suspends execution of the calling process until any child has
 changed state.
 }
 printf("val = %d\n", val);
}
```



- A. 9  
B. 12  
C. 18  
D. 21  
E. None of the above

- D 14) How many times would "Hello World" be printed?

```
int main()
{
 int i;
 for (i = 0; i < 2; i++){
 fork();
 }
}
```

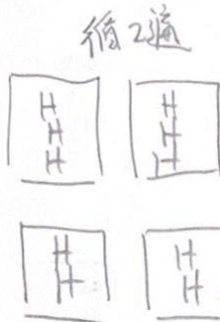


Student Name: \_\_\_\_\_

ID: \_\_\_\_\_

```
printf("Hello World!\n");
}
printf("Hello World!\n");
return 0;
}
```

- A. 4  
B. 6  
C. 8  
D. 10  
E. None of the above



15) Given the task set: T1(4, 1); T2(6, 1); T3(12, 2). What is the largest suitable frame size f?

- A. f=3  
B. f=4  
C. f=6  
D. f=12

$$0414 \quad 0616 \quad 012212$$

$$234612$$

$$24 - \gcd(4, 12) = 24 - 4$$

$$12 - \gcd(4, 6) = 12 - 2$$

$$8 - \gcd(4, 8) = 8 - 4 = 4$$

$$8 - \gcd(6, 4) = 8 - 2 = 6$$

$$8 - \gcd(12, 4) = 8 - 4 = 4$$

16) Real-time systems must have \_\_\_\_\_.

- A. preemptive kernels  
B. non preemptive kernels  
C. memory protection between processes  
D. zero interrupt latency

17) Given 3 periodic tasks T1, T2, and T3. They have the same execution time, but different periods. If the periods of the tasks are 4, 8, and 16. What is the maximum execution time so that the 3 tasks are RM schedulable?

- A. 2      B. 2.28      C. 2.75      D. 3

$$2x + 2x + x \leq 16$$

$$\frac{x}{4} + \frac{x}{8} + \frac{x}{16} \leq 1$$

$$7x \leq 16$$

$$x \leq \frac{16}{7}$$

18) Consider n periodic tasks  $T_i(p_i, e_i)$  for  $1 \leq i \leq n$ , where  $p_i = 2^i$ , and  $e_i = 1$ . What is the value of largest n, so that the task set is (Rate Monotonic) RM schedulable?

- A. 2      B. 4      C. 8      D. infinity

$$\frac{1}{2} + \frac{1}{4} + \frac{1}{8}$$

Student Name: \_\_\_\_\_

ID: \_\_\_\_\_

C 19) Which of the following is most likely to be the output of the following code.

```
#include<stdio.h>
```

```
#include<wait.h>
```

```
#include<signal.h>
```

```
pid_t pid;
```

```
int counter = 0;
```

```
void handler(int sig)
```

```
{
```

```
 counter++;
```

```
}
```

```
int main()
```

```
{
```

```
 pid_t p;
```

```
 int status=0;
```

```
 signal(SIGUSR1, handler);
```

```
 if ((pid = fork()) == 0)
```

```
 {
```

```
 kill(getppid(), SIGUSR1);
```

```
 exit(0);
```

```
 }
```

```
 if ((pid = fork()) == 0)
```

```
 {
```

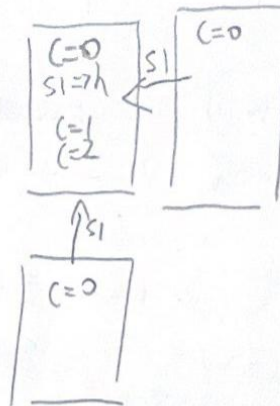
```
 kill(getppid(), SIGUSR1);
```

```
 exit(0);
```

```
 }
```

```
 //Waits for all the child processes
```

```
 while ((p = wait(&status)) > 0);
```





Student Name: \_\_\_\_\_

ID: \_\_\_\_\_

```
sleep(1);
printf("Value of counter is %d", counter);
}
```

- A. Value of counter is 0
- B. Value of counter is 1
- C. Value of counter is 2
- D. Value of counter is 3

- C 20) Which of the following is not a benefit of using kernel module instead of installing all anticipated functionalities into a base kernel?
- A. Save memory space ✓
  - B. No need to rebuilt and reboot the kernel every time ✓
  - C. Allow preemption
  - D. None of the above

Question 2: Short Answer Questions:

1. EDF is an optimal uniprocessor scheduling algorithm, why we still prefer RM instead of EDF in many situations? [1 point]

Because the time behavior of a system scheduled by RM is more predictable. As an engineer, we want more predictable results.

2. Design considerations/objectives for a real time operating system are different from those of a general purpose operating system. Please name two of them. [1 point]

The first consideration is for an OS, the design is for time sharing; but for RTOS, it is for event driven.

What's more, OS requires maximum throughput but RTOS is for good scheduling.

OS wants fast average response time but RTOS wants to meet the deadline of each task.

3. You are required to schedule the following set of independent, preemptable periodic tasks using Rate Monotonic(RM) scheduling algorithm: T1(50; 12); T2(40; 10); T3(30; 10). Use a suitable necessary and sufficient schedulability test to determine if all or part of the tasks can be scheduled. Show your work. [3 points]

By the period  $\bar{T}_3$   $\bar{T}_2$   $\bar{T}_1$

$$\text{Test 1: } \frac{12}{50} + \frac{10}{40} + \frac{10}{30} = 0.823 < 1 \quad \text{OK}$$

Test 3: When  $i=1, j=1, k=1, t=30$

$$W_1(30) = 1 \times 10 = 10 < 30 \quad \text{OK}$$

When  $i=2, j=1, 2, k_1=1, k_2=1, t=30, 40$

$$W_2(30) = 1 \times 10 + 1 \times 10 = 20 < 30 \quad \text{OK}$$

$$W_2(40) = 2 \times 10 + 1 \times 10 = 30 < 40 \quad \text{OK}$$

When  $i=3, j=1, 2, 3, k_1=1, k_2=1, k_3=1, t=30, 40, 50$

$$W_3(30) = 1 \times 10 + 1 \times 10 + 1 \times 12 = 32 > 30$$

$$W_3(40) = 2 \times 10 + 1 \times 10 + 1 \times 12 = 42 > 40$$

$$W_3(50) = 2 \times 10 + 2 \times 10 + 1 \times 12 = 52 > 50$$

$\therefore$  It is not RM schedulable



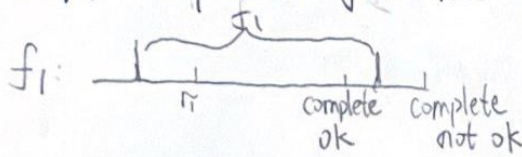
Student Name: \_\_\_\_\_

ID: \_\_\_\_\_

4. In cyclic executive scheduling, why the frame size cannot be too small? and why the frame size cannot be too large? (You can use figure for illustration) [2 points]

If the frame size is too small, a simple instance of a task cannot complete in just one frame

$$f \geq \max_{1 \leq i \leq n} e_i$$



If the frame size is too large, the release time, deadline may be in the same frame, which is impossible for a task to be released and meet deadline in the same frame.

5. You are required to schedule the following set of independent, preemptable periodic tasks using Earliest Deadline First (EDF) scheduling algorithm:  $T_1(50; 15)$ ;  $T_2(40; 10)$ ;  $T_3(30; 10)$ .

- (1) Determine if all or part of the tasks can be scheduled. Show your work. [2 points]  
 (2) Show a feasible schedule using EDF for a period from 0 to 100 seconds if possible. If not, mark when deadline is missed. [1 point]

$$\therefore D_i > P_i \therefore U = \frac{15}{50} + \frac{10}{40} + \frac{10}{30} = 0.88 < 1$$

$\therefore$  all parts are schedulable

