

MECHTRON 2MD3

Data Structures and Algorithms for Mechatronics

Winter 2022

# 23 Trees

Department of Computing and Software

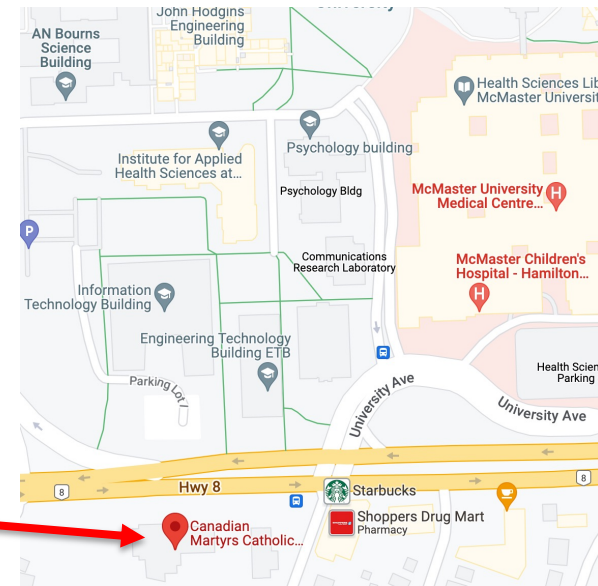
Instructor:

Omid Isfahanialamdari

March 16, 2022

# Admin

- Mid-Term 2:
  - Wednesday 23 March 2022
  - Duration: **1 hour**
  - Location: **MCMST CDN\_MARTYRS**
    - Seems to be here, I am not sure



- Covers: Topics from “Doubly Linked Lists” until the lecture of Wednesday 16 March 2022 (inclusive)

# Iterating through a Container (Review)

- Let **C** be a **container** and **p** be an **iterator** for **C**

for (p = C.begin(); p != C.end(); ++p)

loop\_body

- Example: (with an STL vector)

- Notice how for loop is implemented
- Notice how we access the element of iterator

Output:

Iterator is on data 1  
Iterator is on data 2  
Iterator is on data 3  
Iterator is on data 4  
Iterator is on data 5  
The sum is :15

```
#include <iostream>
#include <vector>
using namespace std;

int main(){
    vector<int> v1;
    for (int i = 1; i <= 5; i++){
        v1.push_back(i);
    }

    typedef vector<int>::iterator Iterator;
    int sum = 0;
    for (Iterator p = v1.begin(); p != v1.end(); ++p){
        cout << "Iterator is on data " << *p << endl;
        sum += *p;
    }
    cout << "The sum is :" << sum << endl;
}
```

# Implementing Iterators (remained from pre. lecture)

- Array-based
  - **Array A of the n elements**
  - **index i** that keeps track of the cursor
  - **begin() = 0**
  - **end() = n** (index following the last element)
- Linked list-based
  - For example: A doubly linked-list **L** storing the elements, with sentinels for header and trailer
  - **pointer to node** containing the current element
  - **begin()** = front node (immediately after the header)
  - **end()** = trailer node (immediately after last node)

# Implementation of Iterator for (Node) List

- This is an implementation of the iterator for our NodeList (aka DLL)

```
class Iterator { // an iterator for the list
public:
    Elem& operator*(); // reference to the element
    bool operator==(const Iterator& p) const; // compare positions
    bool operator!=(const Iterator& p) const;
    Iterator& operator++(); // move to next position
    Iterator& operator--(); // move to previous position
    friend class NodeList; // give NodeList access
private:
    Node* v; // pointer to the node
    Iterator(Node* u); // create from node
};
```

```
NodeList::Iterator::Iterator(Node* u) // constructor from Node*
{ v = u; }
```

```
Elem& NodeList::Iterator::operator*() // reference to the element
{ return v->elem; }
```

```
// compare positions
bool NodeList::Iterator::operator==(const Iterator& p) const
{ return v == p.v; }
```

```
bool NodeList::Iterator::operator!=(const Iterator& p) const
{ return v != p.v; }
```

```
// move to next position
NodeList::Iterator& NodeList::Iterator::operator++()
{ v = v->next; return *this; }
```

```
// move to previous position
NodeList::Iterator& NodeList::Iterator::operator--()
{ v = v->prev; return *this; }
```

```
struct Node {
    Elem elem;
    Node* prev;
    Node* next;
};
```

# (Node) List ADT (duplicate slide)

- This implementation is basically a Doubly Linked-List
- **n**:
  - The number of data elements
- **Iterator** is an implementation of the Position ADT
  - We use iterator instead of a pointer to a specific node in this implementation

```
typedef int Elem; // list base element type
class NodeList { // node-based list
private:
    // insert Node declaration here...
public:
    // insert Iterator declaration here...
public:
    NodeList(); // default constructor
    int size() const; // list size
    bool empty() const; // is the list empty?
    Iterator begin() const; // beginning position
    Iterator end() const; // (just beyond) last position
    void insertFront(const Elem& e); // insert at front
    void insertBack(const Elem& e); // insert at rear
    void insert(const Iterator& p, const Elem& e); // insert e before p
    void eraseFront(); // remove first
    void eraseBack(); // remove last
    void erase(const Iterator& p); // remove p
    // housekeeping functions omitted...
private:
    // data members
    int n; // number of items
    Node* header; // head-of-list sentinel
    Node* trailer; // tail-of-list sentinel
};
```

# STL Iterators in C++

- Each STL container type **C** supports iterators:
  - **C::iterator** – read/write iterator type
  - **C::const\_iterator** – read-only iterator type
  - **C.begin()**, **C.end()** – return start/end iterators
- This iterator-based operators and methods:
  - **\*p**: access current element
  - **++p**, **--p**: advance to next/previous element
  - **C.assign(p, q)**: replace C with contents referenced by the iterator range [p, q) (from p up to, but not including, q)
  - **insert(p, e)**: insert e prior to position p
  - **erase(p)**: remove element at position p
  - **erase(p, q)**: remove elements in the iterator range [p, q)



# STL List in C++

```
#include <list>
using std::list;           // make list accessible
list<float> myList;        // an empty list of floats
```

**list( $n$ ):** Construct a list with  $n$  elements; if no argument list is given, an empty list is created.

**size():** Return the number of elements in  $L$ .

**empty():** Return true if  $L$  is empty and false otherwise.

**front():** Return a reference to the first element of  $L$ .

**back():** Return a reference to the last element of  $L$ .

**push\_front( $e$ ):** Insert a copy of  $e$  at the beginning of  $L$ .

**push\_back( $e$ ):** Insert a copy of  $e$  at the end of  $L$ .

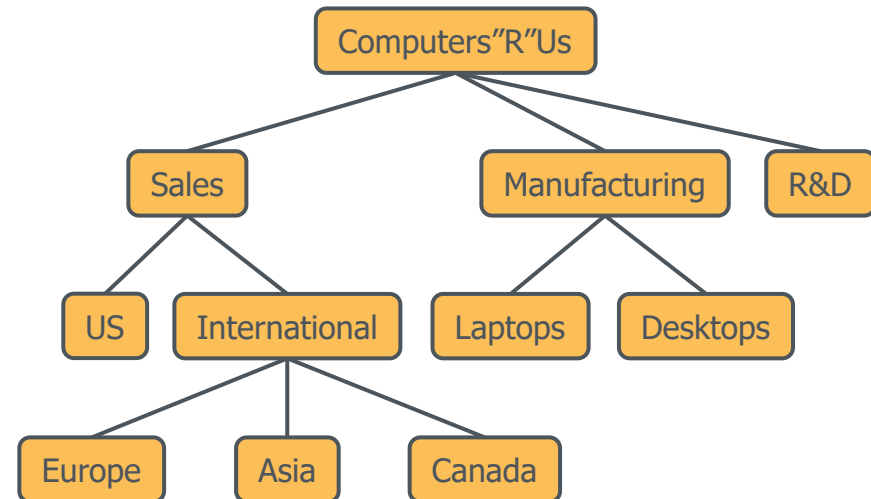
**pop\_front():** Remove the first element of  $L$ .

**pop\_back():** Remove the last element of  $L$ .



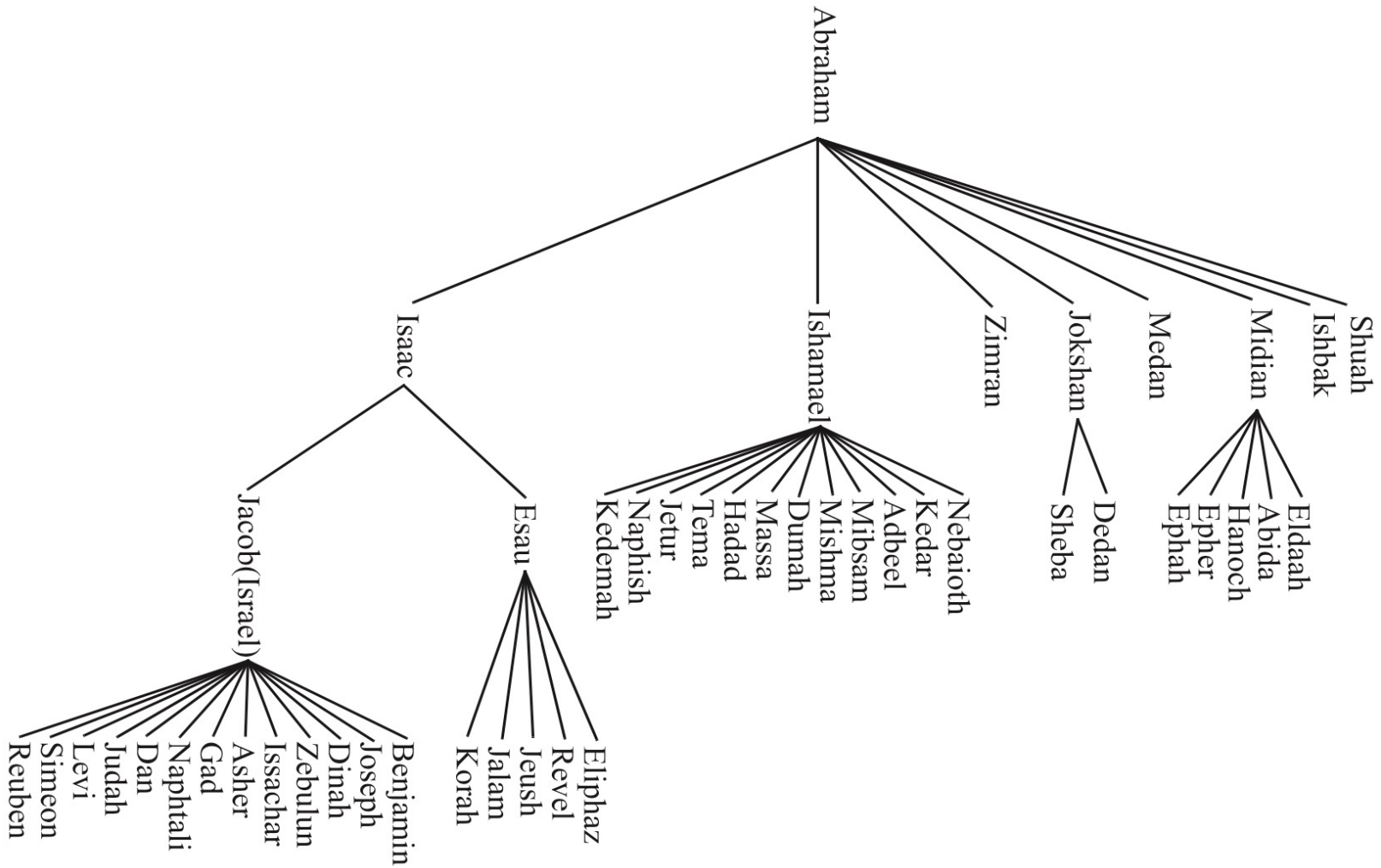
# Trees

- So far you are familiar with:
  - Storing elements in a linear fashion
  - Containers and iterators
- In computer science, a tree is an abstract model of a hierarchical structure
- The relation between elements is non-linear
- A tree consists of nodes with a parent-child relation
- Applications:
  - Organization charts
  - File systems
  - Programming environments



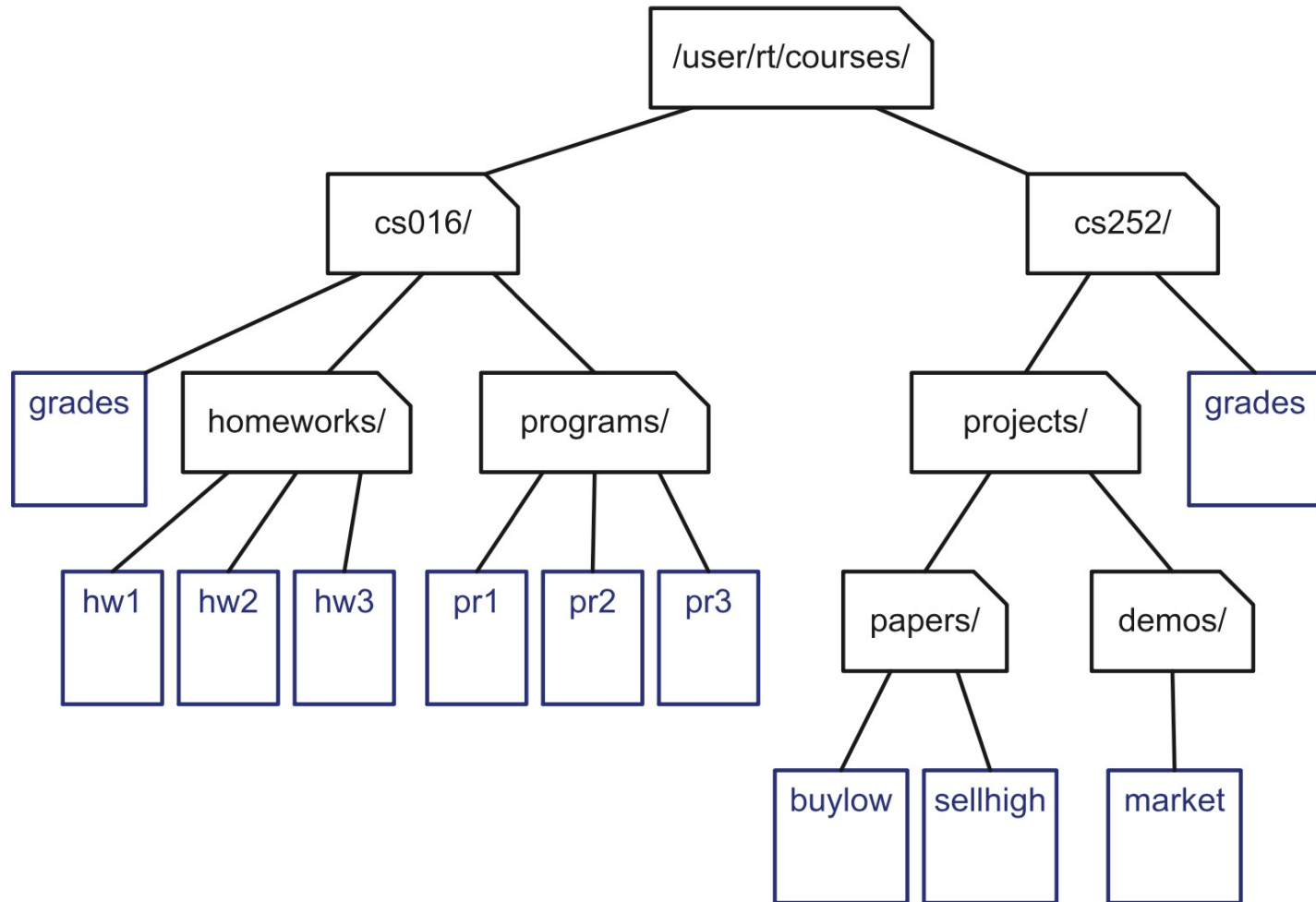
# Tree Examples

- Family Tree



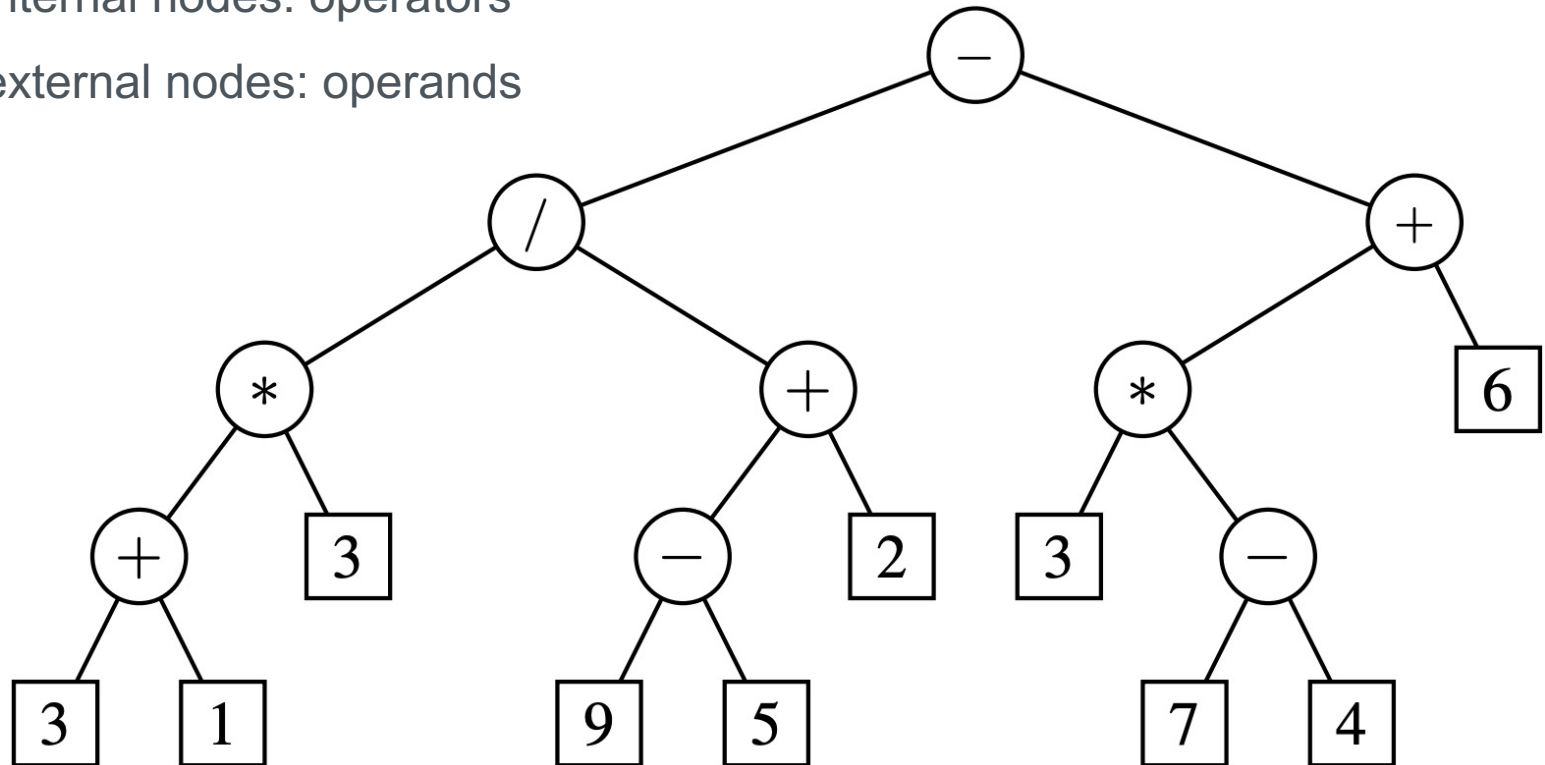
# Tree Examples

- File System Tree



# Tree Examples

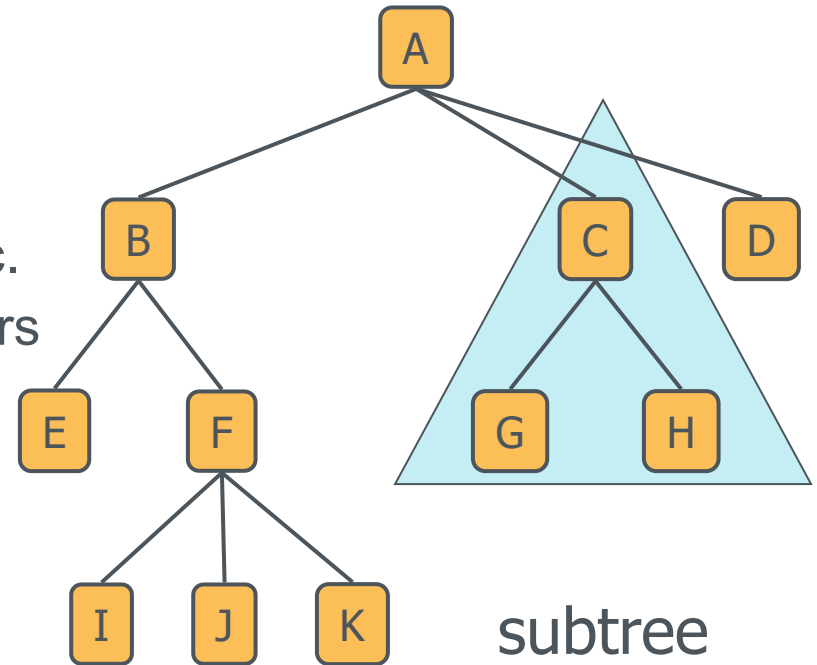
- Expression Tree
  - Binary tree associated with an arithmetic expression
    - internal nodes: operators
    - external nodes: operands



Tree represents:  $((((3 + 1) * 3) / ((9 - 5) + 2)) - ((3 * (7 - 4)) + 6))$

# Trees

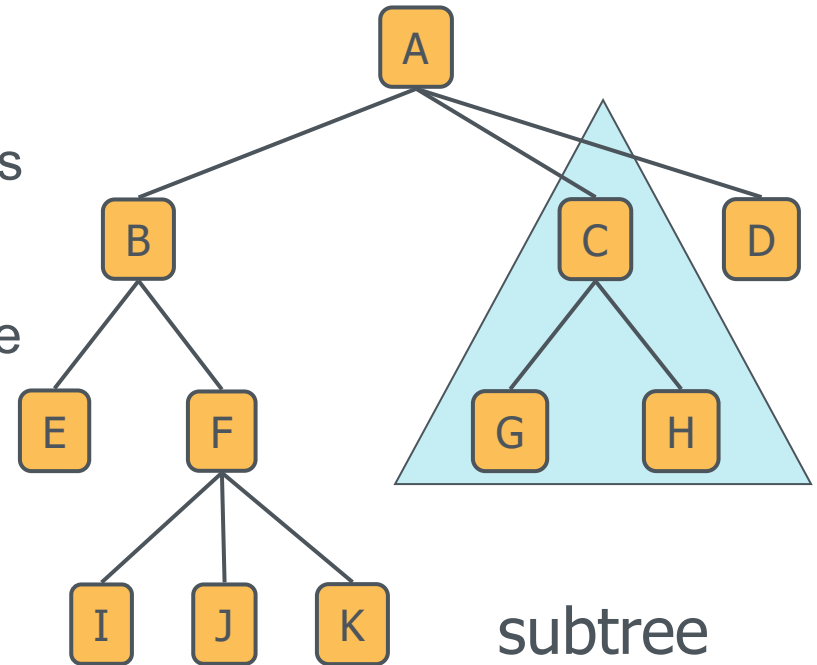
- Root: node without parent (A)
- Internal node: node with at least one child (A, B, C, F)
- External node (a.k.a. leaf ): node without children (E, I, J, K, G, H, D)
- Ancestors of a node: parent, grandparent, grand-grandparent, etc.
- Depth of a node: number of ancestors
- Height of a tree: maximum depth of any node (3)
- Descendant of a node: child, grandchild, grand-grandchild, etc.



- Subtree: tree consisting of a node and its descendants

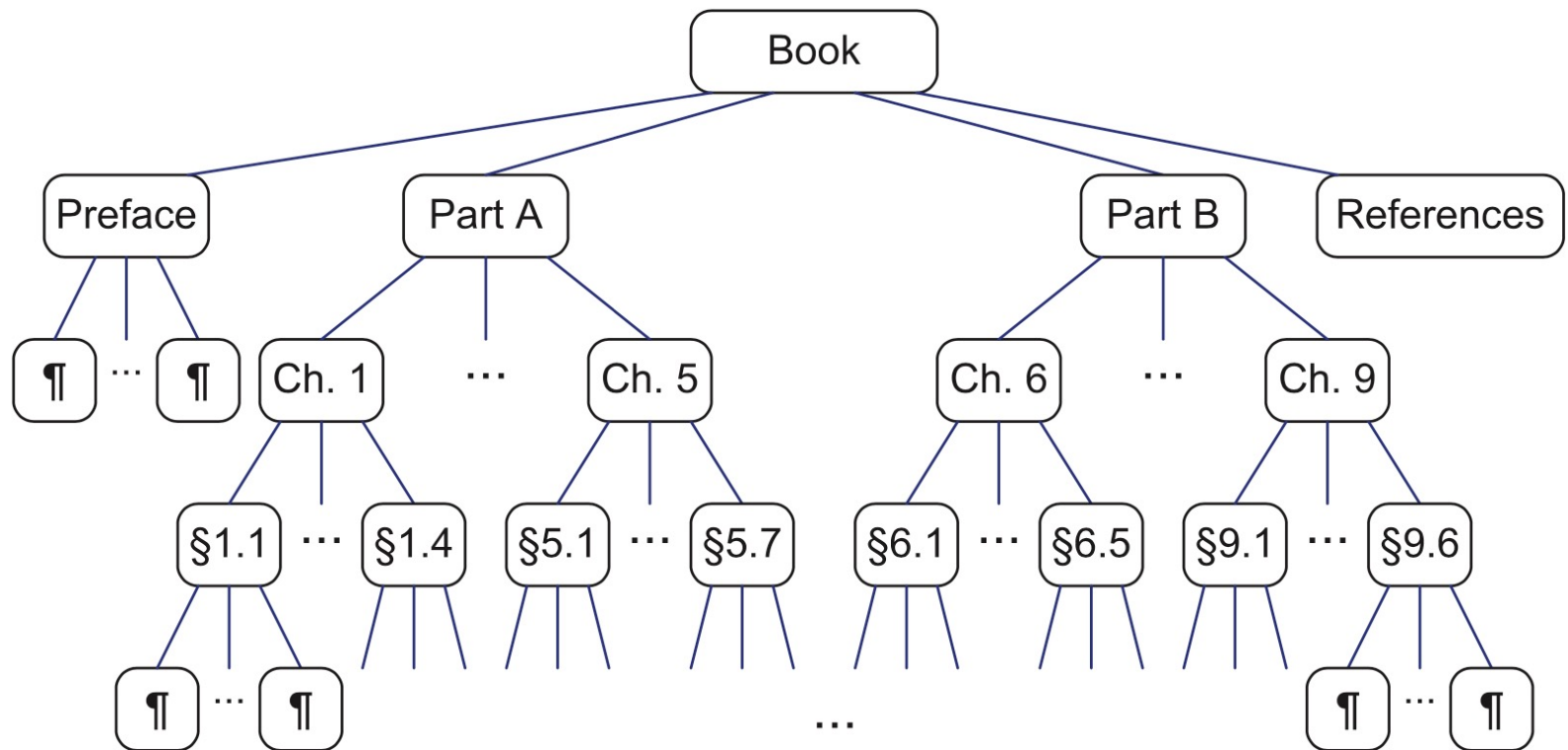
# Formal Definition of Tree

- we define tree **T** to be a **set of nodes** storing elements in a **parent-child** relationship with the following properties:
  - If **T** is *nonempty*, it has a special node, called the **root** of **T**, that has no parent.
  - Each node  $v$  of **T** different from the root has a unique **parent** node  $w$ ; every node with parent  $w$  is a **child** of  $w$ .
- A tree can be empty!
- So, with the above definition we can define it recursively:
  - a tree **T** is either empty or consists of a node **r**, called the root of **T**, and a (possibly empty) set of trees whose roots are the children of **r**



# Ordered Tree Examples

- Ordered Tree: A tree is ordered if there is a linear ordering defined for the children of each node
  - linear order relationship existing between siblings
- Book Organization Tree





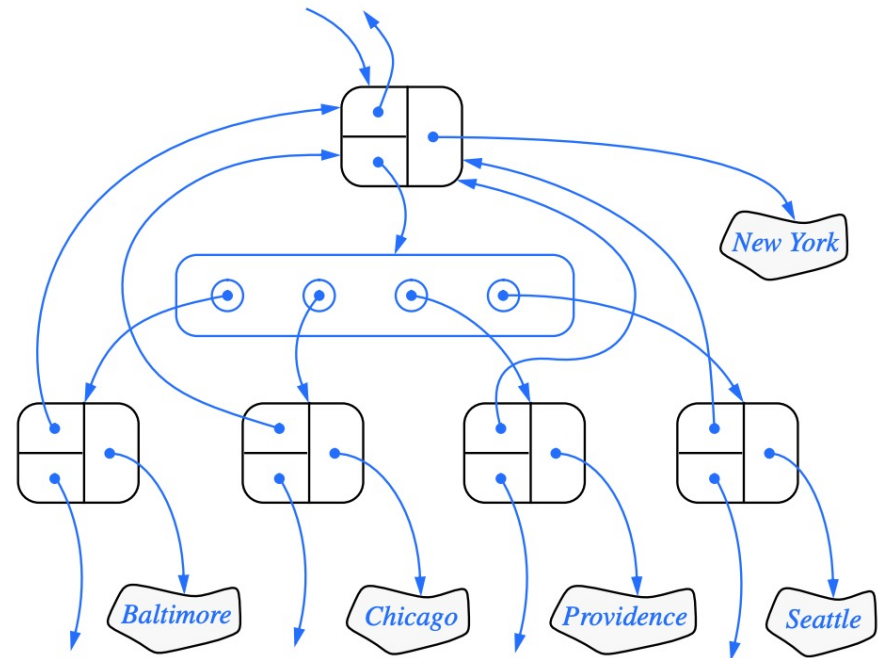
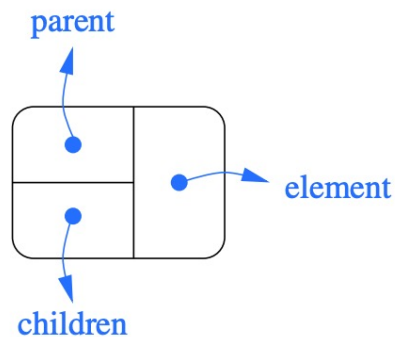
# Tree ADT

- We use **positions** to abstract nodes
- Generic methods:
  - integer **size()**
  - boolean **empty()**
- Accessor methods:
  - position **root()**
  - list<position> **positions()**
- Position-based methods:
  - position p.**parent()**
  - list<position> p.**children()**
- Query methods:
  - boolean p.**isRoot()**
  - boolean p.**isExternal()**
- Additional update methods may be defined by data structures implementing the Tree ADT

# Tree ADT

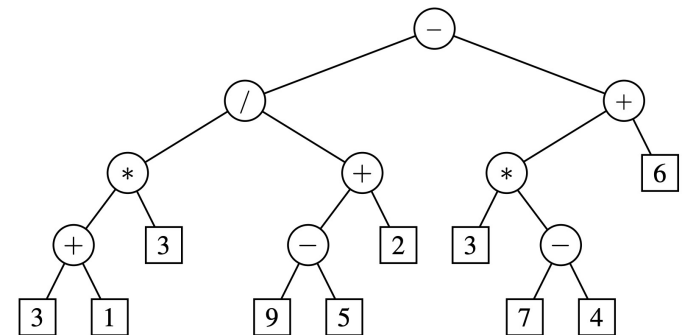
- General Tree Implementation

```
template <typename E>                // base element type
class Position<E> {                  // a node position
public:
    E& operator*();                  // get element
    Position parent() const;         // get parent
    PositionList children() const;   // get node's children
    bool isRoot() const;             // root node?
    bool isExternal() const;         // external node?
};
```



# Tree Traversals

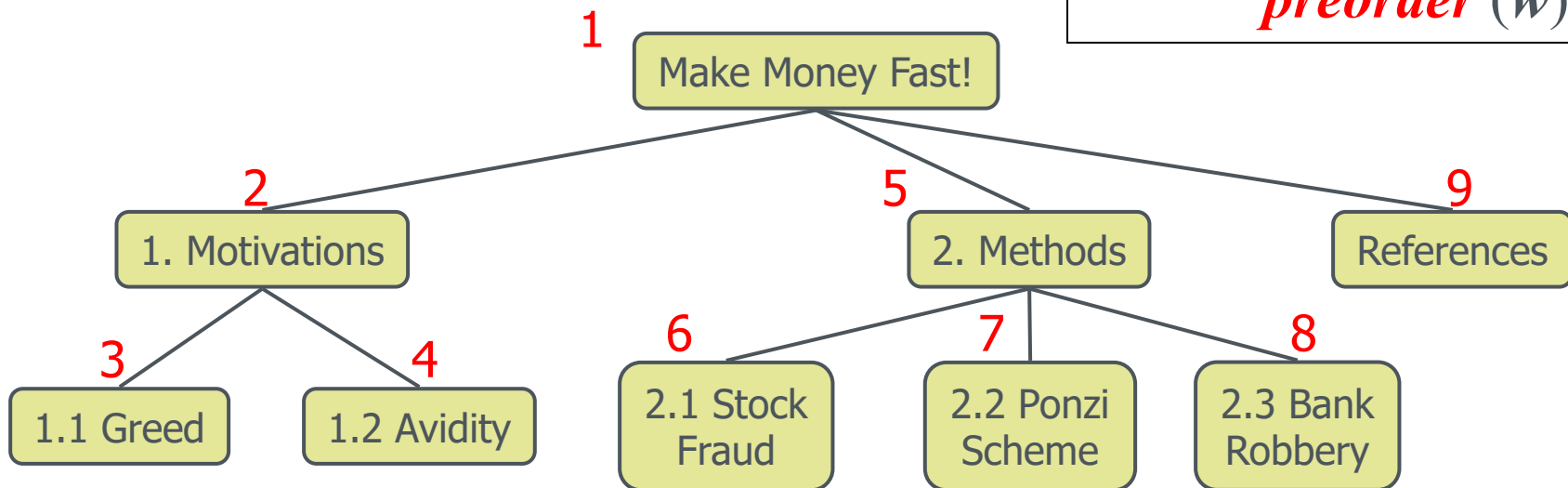
- We were able to easily traverse a linear data structure
  - The traversal of a tree is not trivial
- A **traversal** of a tree **T** is a systematic way of accessing or visiting all the nodes of T.
  - Preorder traversal
  - Inorder traversal
  - Postorder traversal
  - Breadth-first traversal



# Preorder Traversal

- A traversal visits the nodes of a tree in a systematic manner
- In a preorder traversal, a node is visited before its descendants
- Application: print a structured document

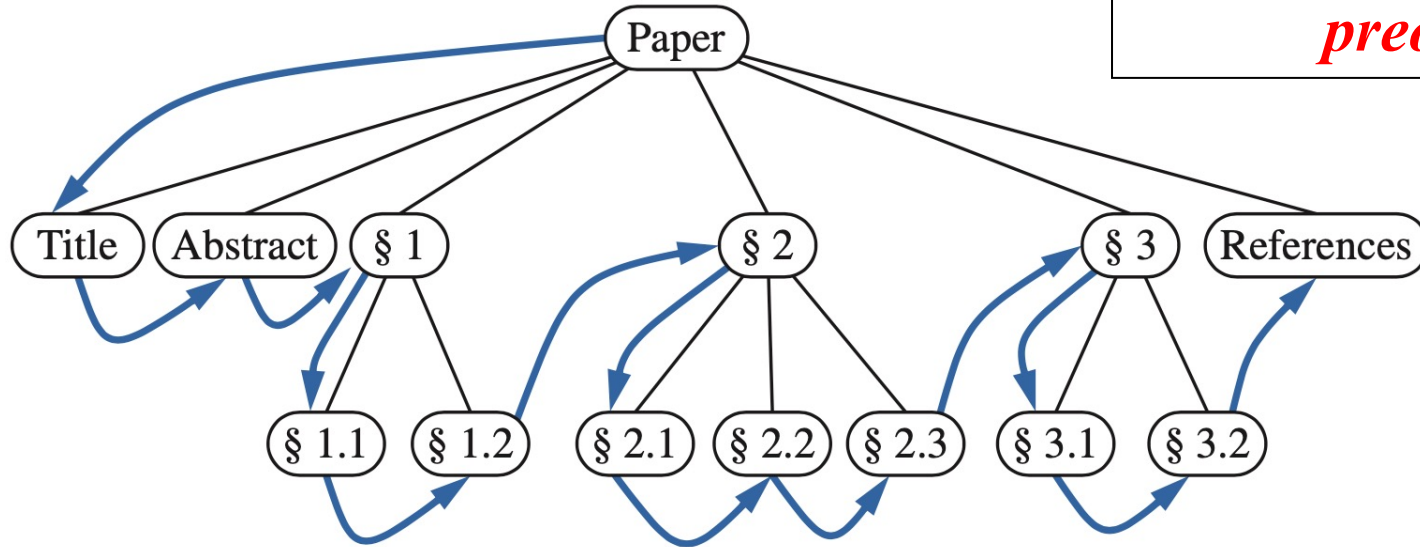
**Algorithm** *preOrder*(*v*)  
*visit*(*v*)  
for each child *w* of *v*  
    *preorder* (*w*)



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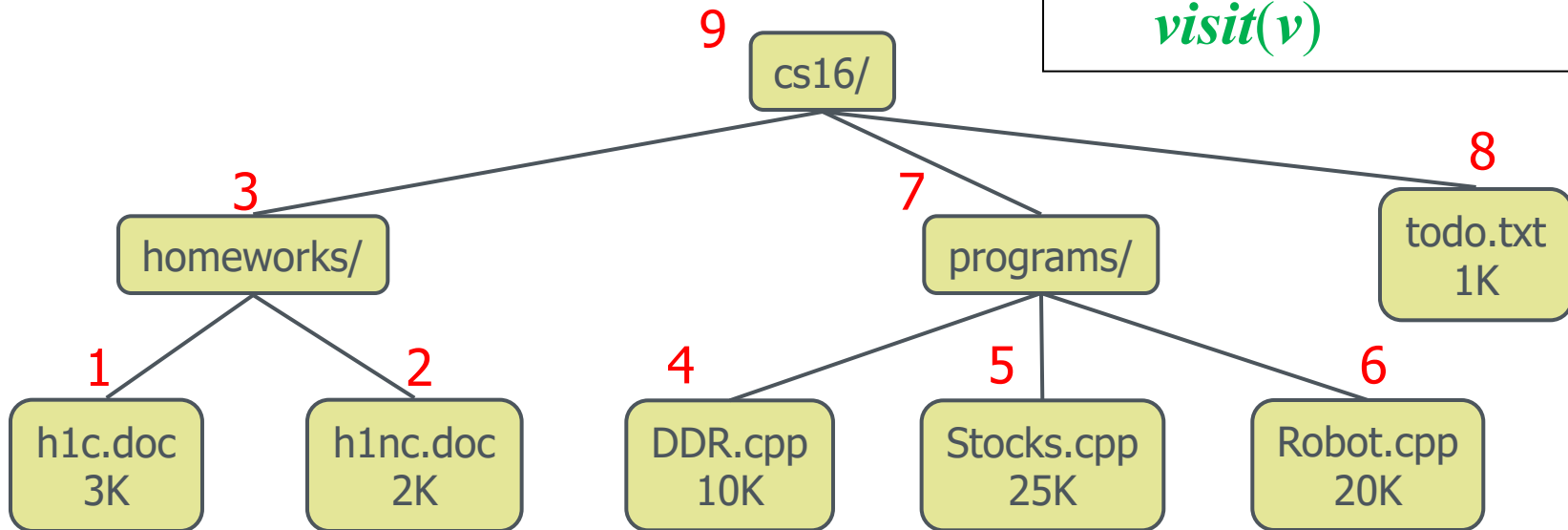
**Algorithm** *preOrder(v)*  
*visit(v)*  
for each child *w* of *v*  
*preorder(w)*

```
void preorderPrint(const Tree& T, const Position& p) {  
    cout << *p;                // print element  
    PositionList ch = p.children(); // list of children  
    for (Iterator q = ch.begin(); q != ch.end(); ++q) {  
        cout << " ";  
        preorderPrint(T, *q);  
    }  
}
```

# Postorder Traversal

- In a postorder traversal, a node is visited after its descendants
- Application: compute space used by files in a directory and its subdirectories

Algorithm *postOrder*( $v$ )  
for each child  $w$  of  $v$   
    *postorder* ( $w$ )  
visit( $v$ )

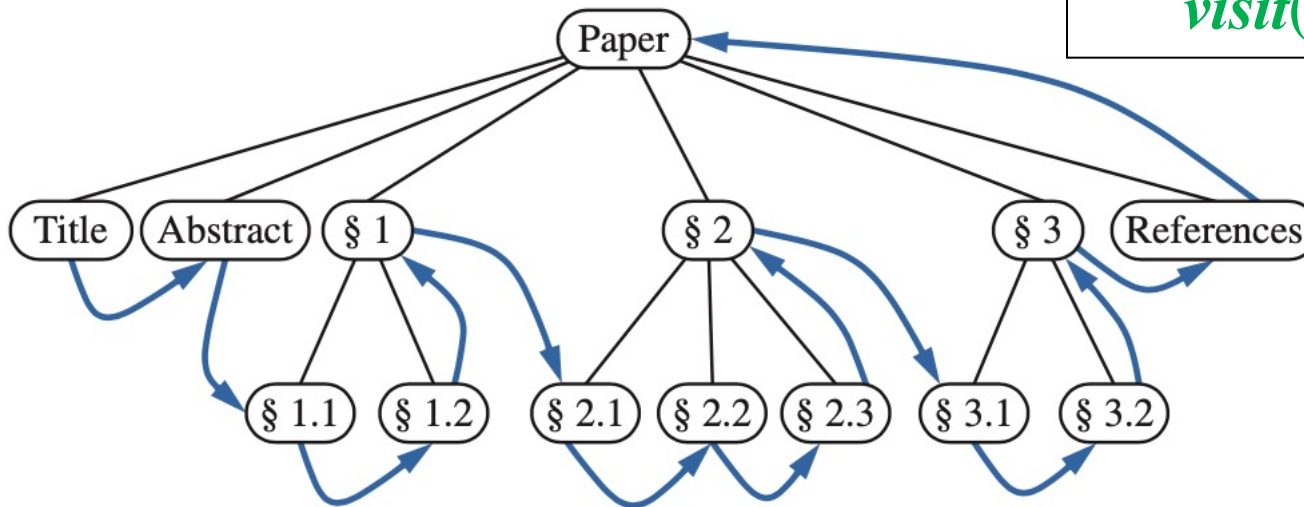




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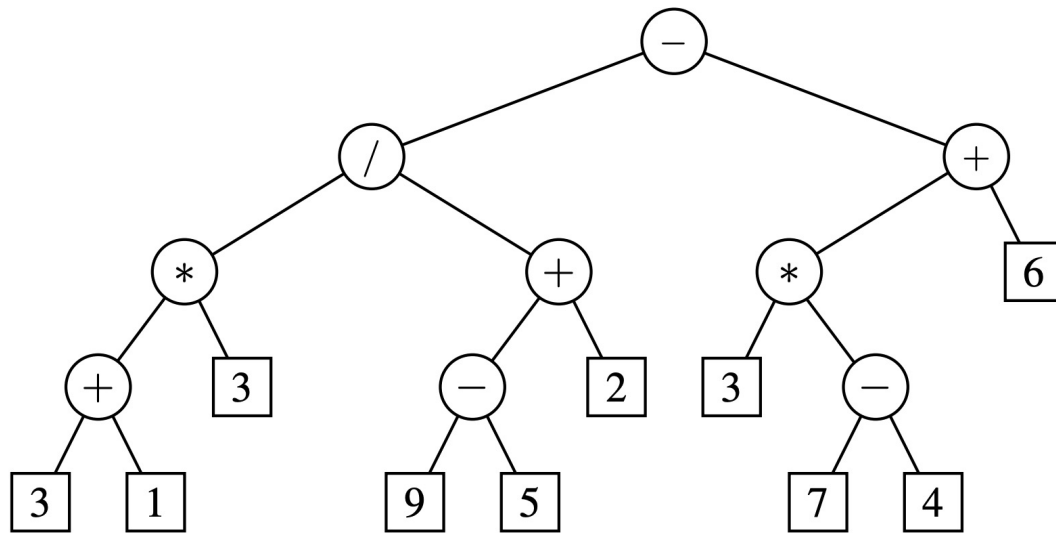
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```
void postorderPrint(const Tree& T, const Position& p) {  
    PositionList ch = p.children();           // list of children  
    for (Iterator q = ch.begin(); q != ch.end(); ++q) {  
        postorderPrint(T, *q);  
        cout << " ";  
    }  
    cout << *p;                               // print element  
}
```

# Postorder Traversal

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$(((((3 + 1) * 3) / ((9 - 5) + 2)) - ((3 * (7 - 4)) + 6)))$

Postorder?

Algorithm *postOrder*(*v*)  
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# Questions?