

326 | Introduction to Automata Theory, Formal Languages and Computation

**Step v :** (Conversion of the string containing terminals and non-terminals to the string of non-terminal. on the RHS) Consider any production of the form

$$NT \rightarrow T_1 T_2 \dots T_n NT_1 \dots NT_n.$$

where  $n \geq 1$  For the part  $T_1 T_2 \dots T_n$ , follow Step III

Check the resultant string by Step II. If not in CNF, follow Step IV.

By this process, the generated new grammar is in CNF.

**Example 6.37**

Convert the following grammar into CNF.

$$\begin{aligned} S &\rightarrow bA/aB \\ A &\rightarrow bAA/aS/a \\ B &\rightarrow aBB/DS/a \end{aligned}$$

**Selation:** The productions  $S \rightarrow bA, S \rightarrow aB, A \rightarrow bAA, A \rightarrow aS, B \rightarrow aBB, B \rightarrow bS$  are not in CNF. So, we have to convert these into CNF. Let us replace terminal 'a' by a non-terminal  $C_a$  and terminal 'b' by a noa-terminal  $C_b$ . Hence, two new productions will be added to the grammar

$$C_a \rightarrow a \text{ and } C_b \rightarrow b$$

By replacing a and b by new non-terminals and including the two productions, the modified grammar will be

$$\begin{aligned} S &\rightarrow C_b A / C_a B \\ A &\rightarrow C_b AA / C_a S / a \\ B &\rightarrow C_a BB / C_b S / a \\ C_a &\rightarrow a \\ C_b &\rightarrow b \end{aligned}$$

In the modified grammar, all the productions are not in CNF. The productions  $A \rightarrow C_b AA$  and  $B \rightarrow C_a BB$  are not in CNF, because they contain more than two noa-terminals at the RHS. Let us replace AA by a new non-terminal D and BB by another new non-terminal E. Hence, two new productions will be added to the grammar  $D \rightarrow AA$  and  $E \rightarrow BB$ . So, the new modified grammar will be

$$\begin{aligned} S &\rightarrow C_b A / C_a B \\ A &\rightarrow C_b D / C_a S / a \\ B &\rightarrow C_b E / C_a S / a \\ D &\rightarrow AA \\ E &\rightarrow BB \\ C_a &\rightarrow a \\ C_b &\rightarrow b \end{aligned}$$

**Example 6.38**

Convert the following grammar into CRF.

$$\begin{aligned} S &\rightarrow ABa \\ A &\rightarrow aab \\ B &\rightarrow Ac \end{aligned}$$