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# 1 Basic

# 1.1 Shell Script

```
g++ -std=c++17 -DABS -Wall -Wextra -Wshadow $1.cpp -o
    $1 && ./$1
for i in {A..J}; do cp tem.cpp $i.cpp; done;
```

#### 1.2 Default Code

```
#include <bits/stdc++.h>
using namespace std;
typedef long long l1;
#define pb push_back
#define pii pair<int, int>
#define all(a) a.begin(), a.end()
#define sz(a) ((int)a.size())
```

#### 1.3 Increase Stack Size

```
const int size = 256 << 20;
register long rsp asm("rsp");
char *p = (char*)malloc(size) + size, *bak = (char*)rsp
;
__asm__("movq %0, %%rsp\n"::"r"(p));
// main
__asm__("movq %0, %%rsp\n"::"r"(bak));</pre>
```

# 1.4 Debug Macro

## 1.5 Stress Test Shell\*

```
#!/usr/bin/env bash
    g++ $1.cpp -o $1
    g++ $2.cpp -o $2
    g++ $3.cpp -o $3
    for i in {1..100}; do
     ./$3 > input.txt
      # st=$(date +%s%N)
      ./$1 < input.txt > output1.txt
      # echo "$((($(date +%s%N) - $st)/1000000))ms"
      ./$2 < input.txt > output2.txt
      if cmp --silent -- "output1.txt" "output2.txt"; then
        continue
      fi
     echo Input:
      cat input.txt
      echo Your Output:
      cat output1.txt
      echo Correct Output:
      cat output2.txt
     exit 1
    done
    echo OK!
20 | ./stress.sh main good gen
```

## 1.6 Pragma / FastIO

```
#pragma GCC optimize("Ofast,inline,unroll-loops")
#pragma GCC target("bmi,bmi2,lzcnt,popcnt,avx2")
#include<unistd.h>
char OB[65536]; int OP;
inline char RC() {
  static char buf[65536], *p = buf, *q = buf;
  return p == q && (q = (p = buf) + read(0, buf, 65536)
      ) == buf ? -1 : *p++;
inline int R() {
  static char c;
  while((c = RC()) < '0'); int a = c ^ '0';
  while((c = RC()) >= '0') a *= 10, a += c ^ '0';
  return a;
inline void W(int n) {
  static char buf[12], p;
 if (n == 0) OB[OP++]='0'; p = 0;
while (n) buf[p++] = '0' + (n % 10), n /= 10;
  for (--p; p >= 0; --p) OB[OP++] = buf[p];
  if (OP > 65520) write(1, OB, OP), OP = 0;
```

#### 1.7 Divide\*

```
ll divdown(ll a, ll b) {
    return a / b - (a < 0 && a % b);
}
ll divup(ll a, ll b) {
    return a / b + (a > 0 && a % b);
}
a / b < x -> divdown(a, b) + 1 <= x
a / b <= x -> divup(a, b) <= x
x < a / b -> x <= divup(a, b) - 1
x <= a / b -> x <= divdown(a, b)</pre>
```

# 2 Data Structure

#### 2.1 Leftist Tree

```
struct node {
  ll rk, data, sz, sum;
  node *1, *r;
  node(11 \ k) : rk(0), data(k), sz(1), l(0), r(0), sum(k)
       ) {}
11 sz(node *p) { return p ? p->sz : 0; }
11 rk(node *p) { return p ? p->rk : -1; }
11 sum(node *p) { return p ? p->sum : 0; }
node *merge(node *a, node *b) {
 if (!a || !b) return a ? a : b;
  if (a->data < b->data) swap(a, b);
  a->r = merge(a->r, b);
  if (rk(a\rightarrow r) \rightarrow rk(a\rightarrow l)) swap(a\rightarrow r, a\rightarrow l);
  a->rk = rk(a->r) + 1, a->sz = sz(a->l) + sz(a->r) +
       1;
  a \rightarrow sum = sum(a \rightarrow 1) + sum(a \rightarrow r) + a \rightarrow data;
  return a;
void pop(node *&o) {
  node *tmp = o;
  o = merge(o->1, o->r);
  delete tmp;
```

#### 2.2 Splay Tree

```
struct Splay {
  int pa[N], ch[N][2], sz[N], rt, _id;
  ll v[N];
  Splay() {}
  void init() {
    rt = 0, pa[0] = ch[0][0] = ch[0][1] = -1;
    sz[0] = 1, v[0] = inf;
  }
  int newnode(int p, int x) {
    int id = _id++;
    v[id] = x, pa[id] = p;
    ch[id][0] = ch[id][1] = -1, sz[id] = 1;
```

```
return id:
  void rotate(int i) {
    int p = pa[i], x = ch[p][1] == i, gp = pa[p], c =
        ch[i][!x];
    sz[p] -= sz[i], sz[i] += sz[p];
    if (~c) sz[p] += sz[c], pa[c] = p;
    ch[p][x] = c, pa[p] = i;
    pa[i] = gp, ch[i][!x] = p;
    if (~gp) ch[gp][ch[gp][1] == p] = i;
  void splay(int i) {
    while (~pa[i]) {
      int p = pa[i];
      if (~pa[p]) rotate(ch[pa[p]][1] == p ^ ch[p][1]
           == i ? i : p);
      rotate(i);
    rt = i;
  }
  int lower_bound(int x) {
    int i = rt, last = -1;
    while (true) {
      if (v[i] == x) return splay(i), i;
      if (v[i] \rightarrow x) {
        last = i;
        if (ch[i][0] == -1) break;
        i = ch[i][0];
      else {
        if (ch[i][1] == -1) break;
        i = ch[i][1];
      }
    splay(i);
    return last; // -1 if not found
  void insert(int x) {
    int i = lower_bound(x);
    if (i == -1) {
      // assert(ch[rt][1] == -1);
      int id = newnode(rt, x);
      ch[rt][1] = id, ++sz[rt];
      splay(id);
    else if (v[i] != x) {
      splay(i);
      int id = newnode(rt, x), c = ch[rt][0];
      ch[rt][0] = id;
ch[id][0] = c;
      if (~c) pa[c] = id, sz[id] += sz[c];
      ++sz[rt];
      splay(id);
  }
};
```

#### 2.3 Link Cut Tree

```
// weighted subtree size, weighted path max
struct LCT {
  int ch[N][2], pa[N], v[N], sz[N], sz2[N], w[N], mx[N
      ], _id;
  // sz := sum of v in splay, sz2 := sum of v in
      virtual subtree
  // mx := max w in splay
  bool rev[N];
  LCT() : _id(1) {}
  int newnode(int _v, int _w) {
    int x = _id++;
ch[x][0] = ch[x][1] = pa[x] = 0;
    v[x] = sz[x] = _v;
    sz2[x] = 0;
    w[x] = mx[x] = w;
    rev[x] = false;
    return x;
  void pull(int i) {
    sz[i] = v[i] + sz2[i];
    mx[i] = w[i];
    if (ch[i][0])
```

node \*1, \*r;

```
sz[i] += sz[ch[i][0]], mx[i] = max(mx[i], mx[ch[i])
                                                                node(int k) : data(k), sz(1), l(0), r(0) {}
                                                                void up() {
    if (ch[i][1])
                                                                  sz = 1;
                                                                  if (1) sz += 1->sz;
      sz[i] += sz[ch[i][1]], mx[i] = max(mx[i], mx[ch[i])
           ][1]]);
                                                                  if (r) sz += r \rightarrow sz;
  void push(int i) {
                                                                void down() {}
    if (rev[i]) reverse(ch[i][0]), reverse(ch[i][1]),
                                                              };
        rev[i] = false;
                                                              int sz(node *a) { return a ? a->sz : 0; }
                                                              node *merge(node *a, node *b) {
                                                                if (!a || !b) return a ? a : b;
  void reverse(int i) {
                                                                if (rand() % (sz(a) + sz(b)) < sz(a))</pre>
    if (!i) return;
    swap(ch[i][0], ch[i][1]);
                                                                  return a \rightarrow down(), a \rightarrow r = merge(a \rightarrow r, b), a \rightarrow up(), a
    rev[i] ^= true;
                                                                return b->down(), b->l = merge(a, b->l), b->up(), b;
  bool isrt(int i) {// rt of splay
    if (!pa[i]) return true;
                                                              void split(node *o, node *&a, node *&b, int k) {
    return ch[pa[i]][0] != i && ch[pa[i]][1] != i;
                                                                if (!o) return a = b = 0, void();
                                                                o->down():
                                                                if (o->data <= k)
  void rotate(int i) {
    int p = pa[i], x = ch[p][1] == i, c = ch[i][!x], gp
                                                                  a = o, split(o \rightarrow r, a \rightarrow r, b, k), <math>a \rightarrow up();
                                                                else b = o, split(o->1, a, b->1, k), b->up();
          = pa[p];
    if (ch[gp][0] == p) ch[gp][0] = i;
    else if (ch[gp][1] == p) ch[gp][1] = i;
                                                              void split2(node *o, node *&a, node *&b, int k) {
    pa[i] = gp, ch[i][!x] = p, pa[p] = i;
                                                                if (sz(o) \le k) return a = o, b = 0, void();
    ch[p][x] = c, pa[c] = p;
                                                                o->down();
    pull(p), pull(i);
                                                                if (sz(o->1) + 1 <= k)
                                                                 a = o, split2(o->r, a->r, b, k - <math>sz(o->l) - 1);
  void splay(int i) {
                                                                else b = o, split2(o->1, a, b->1, k);
    vector<int> anc;
                                                                o->up();
    anc.push_back(i);
    while (!isrt(anc.back())) anc.push_back(pa[anc.back
                                                             node *kth(node *o, int k) {
                                                                if (k <= sz(o->1)) return kth(o->1, k);
         ()1);
                                                                if (k == sz(o->1) + 1) return o;
    while (!anc.empty()) push(anc.back()), anc.pop_back
                                                                return kth(o\rightarrow r, k - sz(o\rightarrow l) - 1);
        ();
    while (!isrt(i)) {
      int p = pa[i];
                                                              int Rank(node *o, int key) {
       if (!isrt(p)) rotate(ch[p][1] == i ^ ch[pa[p]][1]
                                                                if (!o) return 0;
                                                                if (o->data < key)</pre>
           == p ? i : p);
      rotate(i);
                                                                  return sz(o->1) + 1 + Rank(o->r, key);
    }
                                                                else return Rank(o->1, key);
  void access(int i) {
                                                              bool erase(node *&o, int k) {
    int last = 0;
                                                                if (!o) return 0;
    while (i) {
                                                                if (o->data == k) {
                                                                  node *t = o;
      splay(i);
      if (ch[i][1])
                                                                  o->down(), o = merge(o->1, o->r);
        sz2[i] += sz[ch[i][1]];
                                                                  delete t;
      sz2[i] -= sz[last];
                                                                  return 1;
      ch[i][1] = last;
      pull(i), last = i, i = pa[i];
                                                                node *&t = k < o->data ? o->l : o->r;
    }
                                                                return erase(t, k) ? o->up(), 1 : 0;
  void makert(int i) {
                                                              void insert(node *&o, int k) {
                                                                node *a, *b;
    access(i), splay(i), reverse(i);
                                                                o->down(), split(o, a, b, k),
  void link(int i, int j) {
                                                               o = merge(a, merge(new node(k), b));
    // assert(findrt(i) != findrt(j));
                                                                o->up();
    makert(i);
    makert(j);
                                                             void interval(node *&o, int 1, int r) {
                                                                node *a, *b, *c; // [l, r)
    pa[i] = j;
    sz2[j] += sz[i];
                                                                o->down();
                                                                split2(o, a, b, 1), split2(b, b, c, r - 1);
    pull(j);
                                                                // operate
  void cut(int i, int j) {
                                                                o = merge(a, merge(b, c)), o->up();
                                                             }
    makert(i), access(j), splay(i);
    // assert(sz[i] == 2 && ch[i][1] == j);
    ch[i][1] = pa[j] = 0, pull(i);
                                                              2.5 2D Segment Tree*
  int findrt(int i) {
                                                             // 2D range add, range sum in Log^2
    access(i), splay(i);
                                                             struct seg {
    while (ch[i][0]) push(i), i = ch[i][0];
                                                                int 1, r;
    splay(i);
                                                                11 sum, 1z;
                                                                seg *ch[2]{};
    return i;
                                                                seg(int _1, int _r) : l(_1), r(_r), sum(0), lz(0) {}
};
                                                                void push()
                                                                  if (lz) ch[0]->add(l, r, lz), ch[1]->modify(l, r,
2.4
      Treap
                                                                      1z), 1z = 0;
struct node {
                                                                void pull() \{sum = ch[0] -> sum + ch[1] -> sum;\}
  int data, sz;
                                                                void add(int _l, int _r, ll d) {
```

if (\_1 <= 1 && r <= \_r) {</pre>

```
sum += d * (r - 1);
      1z += d;
      return;
    if (!ch[0]) ch[0] = new seg(1, 1 + r >> 1), ch[1] =
          new seg(1 + r \gg 1, r);
    push();
    if (_l < l + r >> 1) ch[0]->add(_l, _r, d);
    if (1 + r >> 1 < _r) ch[1]->add(_1, _r, d);
    pull();
  11 qsum(int _1, int _r) {
    if (_1 <= 1 && r <= _r) return sum;</pre>
    if (!ch[0]) return lz * (min(r, _r) - max(l, _l));
    push();
    if (_1 < 1 + r >> 1) res += ch[0]->qsum(_1, _r);
    if (1 + r >> 1 < _r) res += ch[1]->qsum(_1, _r);
    return res:
  }
struct seg2 {
  int 1, r;
  seg v, lz;
  seg2 *ch[2]{};
  seg2(int _1, int _r) : 1(_1), r(_r), v(0, N), lz(0, N
    if (1 < r - 1) ch[0] = new seg2(1, 1 + r >> 1), ch
         [1] = new seg2(1 + r >> 1, r);
  void add(int _1, int _r, int _12, int _r2, l1 d) {
  v.add(_12, _r2, d * (min(r, _r) - max(1, _1)));
    if (_1 <= 1 && r <= _r) {</pre>
      lz.add(_12, _r2, d);
    if (_1 < 1 + r >> 1) ch[0]->add(_1, _r, _12, _r2, d)
    if (l + r >> 1 < _r) ch[1]->add(_l, _r, _l2, _r2, d
  11 qsum(int _1, int _r, int _12, int _r2) {
    11 res = v.qsum(_12, _r2);
    if (_1 <= 1 && r <= _r) return res;</pre>
    res += lz.qsum(_12, _r2) * (min(r, _r) - max(1, _1)
    if (_1 < 1 + r >> 1) res += ch[0]->query(_1, _r,
         _12, _r2);
    if (1 + r >> 1 < _r) res += ch[1]->query(_1, _r,
         _12, _r2);
     return res;
  }
};
```

#### 2.6 Zkw\*

```
11\ \mathsf{mx}[\mathsf{N}\ <<\ 1]\ ,\ \mathsf{sum}[\mathsf{N}\ <<\ 1]\ ,\ 1\mathsf{z}[\mathsf{N}\ <<\ 1]\ ;
void add(int 1, int r, 11 d) { // [l, r), 0-based
  int len = 1, cntl = 0, cntr = 0;
  for (1 += N, r += N + 1; l ^ r ^ 1; l >>= 1, r >>= 1,
       len <<= 1) {
    sum[1] += cnt1 * d, sum[r] += cnt[r] * d;
    if (len > 1) {
      mx[1] = max(mx[1 << 1], mx[1 << 1 | 1]) + lz[1];
      mx[r] = max(mx[r << 1], mx[r << 1 | 1]) + lz[r];
    if (~1 & 1)
      sum[1 ^ 1] += d * len, mx[1 ^ 1] += d, lz[1 ^ 1]
           += d, cntl += len;
    if (r & 1)
      sum[r ^ 1] += d * len, mx[r ^ 1] += d, lz[r ^ 1]
          += d, cntr += len;
  sum[1] += cnt1 * d, sum[r] += cntr * d;
  if (len > 1) {
    mx[1] = max(mx[1 << 1], mx[1 << 1 | 1]) + lz[1];
    mx[r] = max(mx[r << 1], mx[r << 1 | 1]) + lz[r];
  cntl += cntr:
  for (1 >>= 1; 1; 1 >>= 1) {
    sum[1] += cntl * d;
    mx[1] = max(mx[1 << 1], mx[1 << 1 | 1]) + lz[1];
```

```
}
11 qsum(int 1, int r) {
  ll res = 0, len = 1, cntl = 0, cntr = 0;
  for (1 += N, r += N + 1; 1 ^ r ^ 1; 1 >>= 1, r >>= 1,
       len <<= 1) {
    res += cntl * lz[1] + cntr * lz[r];
    if (~l & 1) res += sum[l ^ 1], cntl += len;
    if (r & 1) res += sum[r ^ 1], cntr += len;
  res += cntl * lz[1] + cntr * lz[r];
  cntl += cntr;
  for (1 >>= 1; 1; 1 >>= 1) res += cntl * lz[1];
  return res;
11 qmax(int 1, int r) {
  11 maxl = -INF, maxr = -INF;
for (1 += N, r += N + 1; 1 ^ r ^ 1; 1 >>= 1, r >>= 1)
    \max l += lz[1], \max[r] += lz[r];
    if (~1 & 1) maxl = max(maxl, mx[l ^ 1]);
    if (r & 1) maxr = max(maxr, mx[r ^ 1]);
  maxl = max(maxl + lz[1], maxr + lz[r]);
  for (1 >>= 1; 1; 1 >>= 1) max1 += lz[1];
  return max1;
```

# 2.7 Chtholly Tree\*

```
struct ChthollyTree {
  struct interval {
    int 1, r;
    11 v;
    interval (int _1, int _r, ll _v) : l(_l), r(_r), v(
  struct cmp {
    bool operator () (const interval &a, const interval
        & b) const {
      return a.l < b.l;</pre>
    }
  };
  set <interval, cmp> s;
  vector <interval> split(int 1, int r) {
    // split into [0, l), [l, r), [r, n) and return [l, r]
    vector <interval> del, ans, re;
    auto it = s.lower_bound(interval(1, -1, 0));
    if (it != s.begin() && (it == s.end() || 1 < it->1)
        ) {
      --it;
      del.pb(*it);
      if (r < it->r) {
        re.pb(interval(it->1, 1, it->v));
        ans.pb(interval(1, r, it->v));
        re.pb(interval(r, it->r, it->v));
      } else ·
        re.pb(interval(it->1, 1, it->v));
        ans.pb(interval(l, it->r, it->v));
      ++it;
    for (; it != s.end() && it->r <= r; ++it) {</pre>
      ans.pb(*it);
      del.pb(*it);
    if (it != s.end() && it->l < r) {</pre>
      del.pb(*it);
      ans.pb(interval(it->l, r, it->v));
      re.pb(interval(r, it->r, it->v));
    for (interval &i : del)
      s.erase(i);
    for (interval &i : re)
     s.insert(i);
    return ans;
  void merge(vector <interval> a) {
    for (interval &i : a)
      s.insert(i);
```

# 3 Flow / Matching

# 3.1 Dinic

**}**;

```
struct Dinic { // 0-base
  struct edge {
    int to, cap, flow, rev;
  vector<edge> adj[N];
  int s, t, dis[N], cur[N], n;
  int dfs(int u, int cap) {
    if (u == t || !cap) return cap;
    for (int &i = cur[u]; i < (int)adj[u].size(); ++i)</pre>
       edge &e = adj[u][i];
      if (dis[e.to] == dis[u] + 1 && e.flow != e.cap) {
        int df = dfs(e.to, min(e.cap - e.flow, cap));
        if (df) {
          e.flow += df;
           adj[e.to][e.rev].flow -= df;
          return df;
      }
    dis[u] = -1;
    return 0;
  bool bfs() {
    fill_n(dis, n, -1);
    queue<int> q;
    q.push(s), dis[s] = 0;
    while (!q.empty()) {
      int tmp = q.front();
      q.pop();
      for (auto &u : adj[tmp])
        if (!~dis[u.to] && u.flow != u.cap) {
           q.push(u.to);
           dis[u.to] = dis[tmp] + 1;
    return dis[t] != -1;
  int maxflow(int _s, int _t) {
    s = _s, t = _t;
    int flow = 0, df;
    while (bfs()) {
      fill_n(cur, n, 0);
      while ((df = dfs(s, INF))) flow += df;
    return flow;
  void init(int _n) {
    n = _n;
for (int i = 0; i < n; ++i) adj[i].clear();</pre>
  void reset() {
    for (int i = 0; i < n; ++i)</pre>
      for (auto &j : adj[i]) j.flow = 0;
  void add_edge(int u, int v, int cap) {
    adj[u].pb(edge{v, cap, 0, (int)adj[v].size()});
    adj[v].pb(edge{u, 0, 0, (int)adj[u].size() - 1});
};
```

#### 3.2 Min Cost Max Flow

```
template <typename T1, typename T2>
struct MCMF { // T1 -> flow, T2 -> cost, 0-based
  const T1 INF1 = 1 << 30;
  const T2 INF2 = 1 << 30;
  struct edge {
    int v; T1 f; T2 c;
  } E[M << 1];
  vector <int> adj[N];
  T2 dis[N], pot[N];
  int rt[N], vis[N], n, m, s, t;
  bool SPFA() {
    fill_n(rt, n, -1), fill_n(dis, n, INF2);
```

```
fill_n(vis, n, false);
     queue <int> q;
     q.push(s), dis[s] = 0, vis[s] = true;
     while (!q.empty()) {
       int v = q.front(); q.pop();
       vis[v] = false;
       for (int id : adj[v]) if (E[id].f > 0 && dis[E[id
           ].v] > dis[v] + E[id].c + pot[v] - pot[E[id].
           v1) {
           dis[E[id].v] = dis[v] + E[id].c + pot[v] -
                pot[E[id].v], rt[E[id].v] = id;
           if (!vis[E[id].v]) vis[E[id].v] = true, q.
               push(E[id].v);
         }
    }
     return dis[t] != INF2;
  bool dijkstra() {
     fill_n(rt, n, -1), fill_n(dis, n, INF2);
    priority_queue <pair <T2, int>, vector <pair <T2,
    int>>, greater <pair <T2, int>>> pq;
     dis[s] = 0, pq.emplace(dis[s], s);
     while (!pq.empty()) {
       auto [d, v] = pq.top(); pq.pop();
       if (dis[v] < d) continue;</pre>
       for (int id : adj[v]) if (E[id].f > 0 && dis[E[id
           ].v] > dis[v] + E[id].c + pot[v] - pot[E[id].
           v1) {
           dis[E[id].v] = dis[v] + E[id].c + pot[v] -
               pot[E[id].v], rt[E[id].v] = id;
           pq.emplace(dis[E[id].v], E[id].v);
     return dis[t] != INF2;
  pair <T1, T2> solve(int _s, int _t) {
     s = _s, t = _t, fill_n(pot, n, 0);
     T1 flow = 0; T2 cost = 0;
     bool fr = true;
     while ((fr ? SPFA() : dijkstra())) {
       for (int i = 0; i < n; i++) {</pre>
         dis[i] += pot[i] - pot[s];
       T1 add = INF1;
       for (int i = t; i != s; i = E[rt[i] ^ 1].v) {
        add = min(add, E[rt[i]].f);
       for (int i = t; i != s; i = E[rt[i] ^ 1].v) {
        E[rt[i]].f -= add, E[rt[i] ^ 1].f += add;
       flow += add, cost += add * dis[t];
       fr = false;
       for (int i = 0; i < n; ++i) swap(dis[i], pot[i]);</pre>
    return make_pair(flow, cost);
  void init(int n) {
    n = _n, m = 0;
for (int i = 0; i < n; ++i) adj[i].clear();</pre>
  void reset() {
    for (int i = 0; i < m; ++i) E[i].f = 0;</pre>
  void add_edge(int u, int v, T1 f, T2 c) {
    adj[u].pb(m), E[m++] = \{v, f, c\};
     adj[v].pb(m), E[m++] = {u, 0, -c};
};
```

# 3.3 Kuhn Munkres

```
void add_edge(int a, int b, int wei) {
    w[a][b] = wei;
  bool check(int x) {
    if (vl[x] = 1, ~fl[x]) return q.push(fl[x]), vr[fl[
        x]] = 1;
    while (\sim x) swap(x, fr[fl[x] = pre[x]]);
    return 0:
  void bfs(int s) {
    fill(slk, slk + n, INF), fill(vl, vl + n, 0), fill(
        vr, vr + n, 0);
    q.push(s), vr[s] = 1;
    while (1) {
      T d;
      while (!q.empty()) {
        int y = q.front(); q.pop();
        for (int x = 0; x < n; ++x)
          if (!vl[x] \&\& slk[x] >= (d = hl[x] + hr[y] -
              w[x][y])
            if (pre[x] = y, d) slk[x] = d;
            else if (!check(x)) return;
      d = INF;
      for (int x = 0; x < n; ++x)
        if (!v1[x] && d > s1k[x]) d = s1k[x];
      for (int x = 0; x < n; ++x) {
        if (v1[x]) h1[x] += d;
        else slk[x] -= d;
        if (vr[x]) hr[x] -= d;
      for (int x = 0; x < n; ++x) if (!v1[x] && !s1k[x]
           && !check(x)) return;
    }
  }
  T solve() {
    fill(fl, fl + n, -1), fill(fr, fr + n, -1), fill(hr
          hr + n, 0);
    for (int i = 0; i < n; ++i) hl[i] = *max_element(w[</pre>
        i], w[i] + n);
    for (int i = 0; i < n; ++i) bfs(i);</pre>
    for (int i = 0; i < n; ++i) res += w[i][fl[i]];</pre>
    return res;
 }
};
```

#### 3.4 SW Min Cut

```
template <typename T>
struct SW { // 0-based
  T g[N][N], sum[N]; int n;
  bool vis[N], dead[N];
  void init(int _n) {
    n = _n;
for (int i = 0; i < n; ++i) fill(g[i], g[i] + n, 0)</pre>
    fill(dead, dead + n, false);
  void add_edge(int u, int v, T w) {
    g[u][v] += w, g[v][u] += w;
  T solve() {
    T ans = 1 << 30;
    for (int round = 0; round + 1 < n; ++round) {</pre>
      fill(vis, vis + n, false), fill(sum, sum + n, 0);
      int num = 0, s = -1, t = -1;
      while (num < n - round) {</pre>
        int now = -1;
        for (int i = 0; i < n; ++i) if (!vis[i] && !</pre>
             dead[i]) {
             if (now == -1 || sum[now] < sum[i]) now = i</pre>
        s = t, t = now;
        vis[now] = true, num++;
        for (int i = 0; i < n; ++i) if (!vis[i] && !</pre>
             dead[i]) {
             sum[i] += g[now][i];
      }
```

```
ans = min(ans, sum[t]);
  for (int i = 0; i < n; ++i) {
    g[i][s] += g[i][t];
    g[s][i] += g[t][i];
  }
  dead[t] = true;
}
return ans;
};</pre>
```

# 3.5 Gomory Hu Tree

```
vector <array <int, 3>> GomoryHu(vector <vector <pii>>>
    adj, int n) {
  Tree edge min -> mincut (0-based)
  Dinic flow(n);
  for (int i = 0; i < n; ++i) for (auto [j, w] : adj[i</pre>
      1)
      flow.add_edge(i, j, w);
  flow.record();
  vector <array <int, 3>> ans;
  vector <int> rt(n);
  for (int i = 0; i < n; ++i) rt[i] = 0;</pre>
  for (int i = 1; i < n; ++i) {</pre>
    int t = rt[i];
    flow.reset(); // clear flows on all edge
    ans.push_back({i, t, flow.solve(i, t)});
    flow.runbfs(i);
    for (int j = i + 1; j < n; ++j) if (rt[j] == t &&</pre>
        flow.vis[j]) {
        rt[j] = i;
  }
  return ans;
}
```

#### 3.6 Blossom

```
struct Matching { // 0-based
  int fa[N], pre[N], match[N], s[N], v[N], n, tk;
  vector <int> g[N];
  queue <int> q;
  Matching (int _n) : n(_n), tk(0) {
    for (int i = 0; i <= n; ++i) match[i] = pre[i] = n;</pre>
    for (int i = 0; i < n; ++i) g[i].clear();</pre>
  void add_edge(int u, int v) {
    g[u].push_back(v), g[v].push_back(u);
  int Find(int u) {
    return u == fa[u] ? u : fa[u] = Find(fa[u]);
  int lca(int x, int y) {
    tk++;
    x = Find(x), y = Find(y);
    for (; ; swap(x, y)) {
      if (x != n) {
        if (v[x] == tk) return x;
        v[x] = tk;
        x = Find(pre[match[x]]);
      }
    }
  void blossom(int x, int y, int 1) {
    while (Find(x) != 1) {
      pre[x] = y, y = match[x];
      if (s[y] == 1) q.push(y), s[y] = 0;
      if (fa[x] == x) fa[x] = 1;
      if (fa[y] == y) fa[y] = 1;
      x = pre[y];
    }
  bool bfs(int r) {
    for (int i = 0; i <= n; ++i) fa[i] = i, s[i] = -1;</pre>
    while (!q.empty()) q.pop();
    q.push(r);
    s[r] = 0;
    while (!q.empty()) {
      int x = q.front(); q.pop();
      for (int u : g[x]) {
        if (s[u] == -1) {
```

```
pre[u] = x, s[u] = 1;
          if (match[u] == n) {
            for (int a = u, b = x, last; b != n; a =
                last, b = pre[a])
              last = match[b], match[b] = a, match[a] =
            return true;
          q.push(match[u]);
          s[match[u]] = 0;
        } else if (!s[u] && Find(u) != Find(x)) {
          int 1 = lca(u, x);
          blossom(x, u, 1);
          blossom(u, x, 1);
     }
    return false;
  int solve() {
    int res = 0;
    for (int x = 0; x < n; ++x) {
      if (match[x] == n) res += bfs(x);
    return res:
 }
};
```

# 3.7 Weighted Blossom

```
struct WeightGraph { // 1-based
  static const int inf = INT_MAX;
  static const int maxn = 514;
  struct edge {
    int u, v, w;
    edge(){}
    edge(int u, int v, int w): u(u), v(v), w(w) {}
  };
  int n, n_x;
  edge g[maxn * 2][maxn * 2];
  int lab[maxn * 2];
  int match[maxn * 2], slack[maxn * 2], st[maxn * 2],
      pa[maxn * 2];
 int flo_from[maxn * 2][maxn + 1], S[maxn * 2], vis[
      maxn * 2];
  vector<int> flo[maxn * 2];
  queue<int> q;
  int e_delta(const edge &e) { return lab[e.u] + lab[e.
  v] - g[e.u][e.v].w * 2; }
void update_slack(int u, int x) { if (!slack[x] ||
      e_{delta}(g[u][x]) < e_{delta}(g[slack[x]][x])) slack
      [x] = u; 
  void set_slack(int x) {
    slack[x] = 0;
    for (int u = 1; u <= n; ++u)</pre>
      if (g[u][x].w > 0 && st[u] != x && S[st[u]] == 0)
        update_slack(u, x);
  void q_push(int x) {
    if (x \le n) q.push(x);
    else for (size_t i = 0; i < flo[x].size(); i++)</pre>
        q_push(flo[x][i]);
  void set_st(int x, int b) {
    st[x] = b;
    if (x > n) for (size_t i = 0; i < flo[x].size(); ++</pre>
        i) set_st(flo[x][i], b);
  int get_pr(int b, int xr) {
    int pr = find(flo[b].begin(), flo[b].end(), xr) -
        flo[b].begin();
    if (pr % 2 == 1) {
      reverse(flo[b].begin() + 1, flo[b].end());
      return (int)flo[b].size() - pr;
    return pr;
  void set_match(int u, int v) {
    match[u] = g[u][v].v;
    if (u <= n) return;</pre>
    edge e = g[u][v];
    int xr = flo_from[u][e.u], pr = get_pr(u, xr);
```

```
for (int i = 0; i < pr; ++i) set_match(flo[u][i],</pre>
      flo[u][i ^ 1]);
  set_match(xr, v);
  rotate(flo[u].begin(), flo[u].begin() + pr, flo[u].
      end());
void augment(int u, int v) {
  for (; ; ) {
    int xnv = st[match[u]];
    set_match(u, v);
    if (!xnv) return;
    set_match(xnv, st[pa[xnv]]);
    u = st[pa[xnv]], v = xnv;
int get_lca(int u, int v) {
 static int t = 0:
  for (++t; u || v; swap(u, v)) {
    if (u == 0) continue;
    if (vis[u] == t) return u;
    vis[u] = t;
    u = st[match[u]];
   if (u) u = st[pa[u]];
 return 0:
void add_blossom(int u, int lca, int v) {
 int b = n + 1;
  while (b <= n_x && st[b]) ++b;</pre>
  if (b > n_x) ++n_x;
  lab[b] = 0, S[b] = 0;
  match[b] = match[lca];
  flo[b].clear();
  flo[b].push_back(lca);
  for (int x = u, y; x != lca; x = st[pa[y]])
    flo[b].push_back(x), flo[b].push_back(y = st[
        match[x]]), q_push(y);
  reverse(flo[b].begin() + 1, flo[b].end());
  for (int x = v, y; x != lca; x = st[pa[y]])
    flo[b].push_back(x), flo[b].push_back(y = st[
        match[x]]), q_push(y);
  set_st(b, b);
  for (int x = 1; x \le n_x; ++x) g[b][x].w = g[x][b].
      w = 0;
  for (int x = 1; x <= n; ++x) flo_from[b][x] = 0;
  for (size_t i = 0; i < flo[b].size(); ++i) {</pre>
    int xs = flo[b][i];
    for (int x = 1; x <= n_x; ++x)
      if (g[b][x].w == 0 | e_delta(g[xs][x]) <</pre>
          e_delta(g[b][x]))
        g[b][x] = g[xs][x], g[x][b] = g[x][xs];
    for (int x = 1; x <= n; ++x)
      if (flo_from[xs][x]) flo_from[b][x] = xs;
  set_slack(b);
void expand blossom(int b) {
  for (size_t i = 0; i < flo[b].size(); ++i)</pre>
    set_st(flo[b][i], flo[b][i]);
  int xr = flo_from[b][g[b][pa[b]].u], pr = get_pr(b)
       xr);
  for (int i = 0; i < pr; i += 2) {</pre>
    int xs = flo[b][i], xns = flo[b][i + 1];
    pa[xs] = g[xns][xs].u;
    S[xs] = 1, S[xns] = 0;
    slack[xs] = 0, set_slack(xns);
    q_push(xns);
  S[xr] = 1, pa[xr] = pa[b];
  for (size_t i = pr + 1; i < flo[b].size(); ++i) {</pre>
   int xs = flo[b][i];
    S[xs] = -1, set_slack(xs);
 st[b] = 0;
bool on_found_edge(const edge &e) {
  int u = st[e.u], v = st[e.v];
  if (S[v] == -1) {
    pa[v] = e.u, S[v] = 1;
    int nu = st[match[v]];
    slack[v] = slack[nu] = 0;
    S[nu] = 0, q_push(nu);
```

```
} else if (S[v] == 0) {
    int lca = get_lca(u, v);
    if (!lca) return augment(u,v), augment(v,u), true
    else add_blossom(u, lca, v);
  return false;
bool matching() {
  memset(S + 1, -1, sizeof(int) * n_x);
  memset(slack + 1, 0, sizeof(int) * n_x);
  q = queue<int>();
  for (int x = 1; x <= n_x; ++x)
    if (st[x] == x \&\& !match[x]) pa[x] = 0, S[x] = 0,
         q_push(x);
  if (q.empty()) return false;
  for (; ; ) {
    while (q.size()) {
      int u = q.front(); q.pop();
      if (S[st[u]] == 1) continue;
      for (int v = 1; v <= n; ++v)
        if (g[u][v].w > 0 && st[u] != st[v]) {
          if (e_delta(g[u][v]) == 0) {
            if (on_found_edge(g[u][v])) return true;
          } else update_slack(u, st[v]);
    int d = inf;
    for (int b = n + 1; b <= n_x; ++b)</pre>
      if (st[b] == b && S[b] == 1) d = min(d, lab[b]
          / 2);
    for (int x = 1; x <= n_x; ++x)
      if (st[x] == x && slack[x]) {
        if (S[x] == -1) d = min(d, e_delta(g[slack[x
            ]][x]));
        else if (S[x] == 0) d = min(d, e_delta(g[
            slack[x]][x]) / 2);
    for (int u = 1; u <= n; ++u) {</pre>
      if (S[st[u]] == 0) {
        if (lab[u] <= d) return 0;</pre>
        lab[u] -= d;
      } else if (S[st[u]] == 1) lab[u] += d;
    for (int b = n + 1; b \le n_x; ++b)
      if (st[b] == b) {
        if (S[st[b]] == 0) lab[b] += d * 2;
        else if (S[st[b]] == 1) lab[b] -= d * 2;
    q = queue<int>();
    for (int x = 1; x <= n_x; ++x)</pre>
      if (st[x] == x && slack[x] && st[slack[x]] != x
           && e_delta(g[slack[x]][x]) == 0)
        if (on_found_edge(g[slack[x]][x])) return
            true;
    for (int b = n + 1; b \le n_x; ++b)
      if (st[b] == b && S[b] == 1 && lab[b] == 0)
          expand_blossom(b);
  return false;
}
pair<long long, int> solve() {
  memset(match + 1, 0, sizeof(int) * n);
  n_x = n;
  int n_matches = 0;
  long long tot_weight = 0;
  for (int u = 0; u <= n; ++u) st[u] = u, flo[u].
      clear();
  int w_max = 0;
  for (int u = 1; u <= n; ++u)</pre>
    for (int v = 1; v <= n; ++v) {
      flo_from[u][v] = (u == v ? u : 0);
      w_max = max(w_max, g[u][v].w);
  for (int u = 1; u <= n; ++u) lab[u] = w_max;</pre>
  while (matching()) ++n_matches;
  for (int u = 1; u <= n; ++u)</pre>
    if (match[u] && match[u] < u)</pre>
      tot_weight += g[u][match[u]].w;
  return make_pair(tot_weight, n_matches);
void add_edge(int ui, int vi, int wi) { g[ui][vi].w =
```

```
g[vi][ui].w = wi; }
  void init(int _n) {
    n = _n;
    for (int u = 1; u <= n; ++u)</pre>
      for (int v = 1; v <= n; ++v)
        g[u][v] = edge(u, v, 0);
  }
};
```

# 3.8 Flow Model

- Maximum/Minimum flow with lower bound / Circulation problem
  - 1. Construct super source S and sink T.
  - 2. For each edge (x,y,l,u), connect x o y with capacity u-l.
  - 3. For each vertex v, denote by in(v) the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.
  - 4. If in(v)>0, connect  $S\to v$  with capacity in(v), otherwise, connect  $v\to T$  with capacity -in(v).
    - To maximize, connect t 
      ightarrow s with capacity  $\infty$  (skip this in circulation problem), and let f be the maximum flow from S to T. If  $f \neq \sum_{v \in V, in(v) > 0} in(v)$ , there's no solution. Otherwise, the maximum flow from s to t is the answer
    - To minimize, let f be the maximum flow from S to T. Connect  $t \to s$  with capacity  $\infty$  and let the flow from S to T be f'. If  $f+f' \neq \sum_{v \in V, in(v)>0} in(v)$ , there's no solution. Otherwise, f' is the answer.
  - 5. The solution of each edge e is  $l_e+f_e$ , where  $f_e$  corresponds to the flow of edge  $\boldsymbol{e}$  on the graph.
- ullet Construct minimum vertex cover from maximum matching M on bipar- ${\rm tite\ graph\ }(X,Y)$ 
  - 1. Redirect every edge:  $y \to x$  if  $(x,y) \in M$ ,  $x \to y$  otherwise.

  - 2. DFS from unmatched vertices in X. 3.  $x \in X$  is chosen iff x is unvisited. 4.  $y \in Y$  is chosen iff y is visited.
- Maximum density induced subgraph
  - 1. Binary search on answer, suppose we're checking answer  ${\cal T}$
  - 2. Construct a max flow model, let K be the sum of all weights 3. Connect source  $s\to v$ ,  $v\in G$  with capacity K

  - 4. For each edge (u,v,w) in G, connect  $u \to v$  and  $v \to u$  with capacity  $\boldsymbol{w}$
  - 5. For  $v \in G$ , connect it with sink  $v \to t$  with capacity  $K + 2T (\sum_{e \in E(v)} w(e)) 2w(v)$
  - 6. T is a valid answer if the maximum flow  $f < K \lvert V \rvert$
- Minimum weight edge cover 1. For each  $v \in V$  create a copy v' , and connect  $u' \to v'$  with
  - weight w(u,v). 2. Connect  $v \to v'$  with weight  $2\mu(v)$ , where  $\mu(v)$  is the cost of the cheapest edge incident to v .
  - 3. Find the minimum weight perfect matching on  $G^{\prime}$ .
- Project selection problem
  - 1. If  $p_v > 0$ , create edge (s,v) with capacity  $p_v$ ; otherwise,
  - create edge (v,t) with capacity  $-p_v$ . 2. Create edge (u,v) with capacity w with w being the cost of choosing u without choosing v.
  - 3. The mincut is equivalent to the maximum profit of a subset
- 0/1 quadratic programming

$$\sum_{x} c_{x}x + \sum_{y} c_{y}\bar{y} + \sum_{xy} c_{xy}x\bar{y} + \sum_{xyx'y'} c_{xyx'y'}(x\bar{y} + x'\bar{y'})$$

can be minimized by the mincut of the following graph:

- 1. Create edge (x,t) with capacity  $c_x$  and create edge (s,y) with
- capacity  $c_y$  . Create edge (x,y) with capacity  $c_{xy}$
- 3. Create edge (x,y) and edge (x',y') with capacity  $c_{xyx'y'}$ .

# 4 Graph

## 4.1 Heavy-Light Decomposition

```
vector<int> dep, pa, sz, ch, hd, id;
int _id;
void dfs(int i, int p) {
  dep[i] = \sim p ? dep[p] + 1 : 0;
  pa[i] = p, sz[i] = 1, ch[i] = -1;
  for (int j : g[i])
    if (j != p) {
      dfs(j, i);
      if (ch[i] == -1 || sz[ch[i]] < sz[j]) ch[i] = j;</pre>
      sz[i] += sz[j];
```

```
void hld(int i, int p, int h) {
  hd[i] = h;
  id[i] = _id++;
  if (~ch[i]) hld(ch[i], i, h);
  for (int j : g[i]) if (j != p && j != ch[i])
     hld(j, i, j);
}
void query(int i, int j) {
  while (hd[i] != hd[j]) {
    if (dep[hd[i]] < dep[hd[j]]) swap(i, j);
     query2(id[hd[i]], id[i] + 1), i = pa[hd[i]];
  }
  if (dep[i] < dep[j]) swap(i, j);
  query2(id[j], id[i] + 1);
}</pre>
```

# 4.2 Centroid Decomposition

```
vector<vector<int>> dis; // dis[n][logn]
vector<int> pa, sz, dep;
vector<bool> vis;
void dfs_sz(int i, int p) {
  sz[i] = 1;
  for (int j : g[i]) if (j != p && !vis[j])
    dfs_sz(j, i), sz[i] += sz[j];
int cen(int i, int p, int _n) {
  for (int j : g[i]) if (j != p && !vis[j] && sz[j] >
      _n / 2)
    return cen(j, i, _n);
  return i;
void dfs_dis(int i, int p, int d) { // from i to
    ancestor with depth d
  dis[i][d] = \sim p ? dis[p][d] + 1 : 0;
  for (int j : g[i]) if (j != p && !vis[j])
    dfs_dis(j, i, d);
void cd(int i, int p, int d) {
  dfs_sz(i, -1), i = cen(i, -1, sz[i]);
  vis[i] = true, pa[i] = p, dep[i] = d;
  dfs_dis(i, -1, d);
  for (int j : g[i]) if (!vis[j])
    cd(j, i, d + 1);
}
```

# 4.3 Edge BCC

```
vector<int> low, dep, bcc id, stk;
vector<bool> vis;
int _id;
void dfs(int i, int p) {
  low[i] = dep[i] = \sim p ? dep[p] + 1 : 0;
  stk.push_back(i);
  vis[i] = true;
  for (int j : g[i])
    if (j != p) {
      if (!vis[j])
        dfs(j, i), low[i] = min(low[i], low[j]);
       else
        low[i] = min(low[i], dep[j]);
  if (low[i] == dep[i]) {
    int id = _id++;
    while (stk.back() != i) {
      int x = stk.back();
       stk.pop_back();
      bcc_id[x] = id;
    stk.pop_back();
    bcc_id[i] = id;
}
```

#### 4.4 Block Cut Tree

```
vector<vector<int>> g, _g;
vector<int>> dep, low, stk;
void dfs(int i, int p) {
  dep[i] = low[i] = ~p ? dep[p] + 1 : 0;
  stk.push_back(i);
  for (int j : g[i]) if (j != p) {
```

```
if (dep[j] == -1) {
    dfs(j, i), low[i] = min(low[i], low[j]);
    if (low[j] >= dep[i]) {
        int id = _g.size();
        _g.emplace_back();
        while (stk.back() != j) {
            int x = stk.back();
            stk.pop_back();
            _g[x].push_back(id), _g[id].push_back(x);
        }
        stk.pop_back();
        _g[j].push_back(id), _g[id].push_back(j);
        _g[j].push_back(id), _g[id].push_back(i);
    }
    } else low[i] = min(low[i], dep[j]);
}
```

#### 4.5 SCC / 2SAT

```
struct SAT {
  vector<vector<int>> g;
   vector<int> dep, low, scc_id;
   vector<bool> is;
   vector<int> stk;
   int n, _id, _t;
   SAT() {}
   void init(int _n) {
     n = _n, _id = _t = 0;
     g.assign(2 * n, vector<int>());
dep.assign(2 * n, -1), low.assign(2 * n, -1);
scc_id.assign(2 * n, -1), is.assign(2 * n, false);
     stk.clear();
   void add_edge(int x, int y) {g[x].push_back(y);}
   int rev(int i) {return i < n ? i + n : i - n;}</pre>
   void add_ifthen(int x, int y) {add_clause(rev(x), y)
       ;}
   void add_clause(int x, int y) {
     add_edge(rev(x), y);
     add_edge(rev(y), x);
   void dfs(int i) {
     dep[i] = low[i] = _t++;
     stk.push_back(i);
     for (int j : g[i])
       if (scc_id[j] == -1) {
         if (dep[j] == -1)
           dfs(j);
         low[i] = min(low[i], low[j]);
     if (low[i] == dep[i]) {
       int id = _id++;
       while (stk.back() != i) {
         int x = stk.back();
          stk.pop_back();
         scc id[x] = id;
       }
       stk.pop_back();
       scc_id[i] = id;
   bool solve() {
     for (int i = 0; i < 2 * n; ++i)
       if (dep[i] == -1)
         dfs(i);
     for (int i = 0; i < n; ++i) {</pre>
       if (scc_id[i] == scc_id[i + n]) return false;
       if (scc_id[i] < scc_id[i + n])</pre>
         is[i] = true;
       else
         is[i + n] = true;
     return true;
|};
```

#### 4.6 Virtual Tree

```
vector<vector<int>> _g;
vector<int> st, ed, stk;
void solve(vector<int> v) {
```

template <typename T> struct DMST { // 1-based

#### 4.7 Directed MST

```
T g[maxn][maxn], fw[maxn];
int n, fr[maxn];
bool vis[maxn], inc[maxn];
void clear() {
  for (int i = 0; i < maxn; ++i) {</pre>
    for (int j = 0; j < maxn; ++j) g[i][j] = inf;</pre>
    vis[i] = inc[i] = false;
  }
}
void addedge(int u, int v, T w) {
  g[u][v] = min(g[u][v], w);
T query(int root, int _n) {
  n = n;
  if (dfs(root) != n) return -1;
  T ans = 0;
  while (true) {
    for (int i = 1; i <= n; ++i) fw[i] = inf, fr[i] =</pre>
    for (int i = 1; i <= n; ++i) if (!inc[i]) {</pre>
      for (int j = 1; j <= n; ++j) {</pre>
        if (!inc[j] && i != j && g[j][i] < fw[i]) {</pre>
          fw[i] = g[j][i];
           fr[i] = j;
        }
      }
    }
    int x = -1;
    for (int i = 1; i <= n; ++i) if (i != root && !</pre>
        inc[i]) {
      int j = i, c = 0;
      while (j != root && fr[j] != i && c <= n) ++c,</pre>
           j = fr[j];
      if (j == root || c > n) continue;
      else { x = i; break; }
    if (!~x) {
      for (int i = 1; i <= n; ++i) if (i != root && !</pre>
           inc[i]) ans += fw[i];
      return ans;
    for (int i = 1; i <= n; ++i) vis[i] = false;</pre>
    do { ans += fw[y]; y = fr[y]; vis[y] = inc[y] =
         true; } while (y != x);
    inc[x] = false;
    for (int k = 1; k <= n; ++k) if (vis[k]) {</pre>
      for (int j = 1; j <= n; ++j) if (!vis[j]) {</pre>
           if (g[x][j] > g[k][j]) g[x][j] = g[k][j];
           if (g[j][k] < inf \&\& g[j][k] - fw[k] < g[j]
               ][x]) g[j][x] = g[j][k] - fw[k];
        }
    }
  return ans;
int dfs(int now) {
  int r = 1:
  vis[now] = true;
  for (int i = 1; i <= n; ++i) if (g[now][i] < inf &&</pre>
        !vis[i]) r += dfs(i);
```

```
return r;
}
};
```

# 4.8 Dominator Tree

```
struct Dominator_tree {
  int n, id;
  vector <vector <int>> adj, radj, bucket;
  vector <int> sdom, dom, vis, rev, par, rt, mn;
  Dominator_tree (int _n) : n(_n), id(0) {
    adj.resize(n), radj.resize(n), bucket.resize(n);
    sdom.resize(n), dom.resize(n, -1), vis.resize(n,
    rev.resize(n), rt.resize(n), mn.resize(n), par.
        resize(n);
  void add_edge(int u, int v) {adj[u].pb(v);}
  int query(int v, bool x) {
    if (rt[v] == v) return x ? -1 : v;
    int p = query(rt[v], true);
    if (p == -1) return x ? rt[v] : mn[v];
    if (sdom[mn[v]] > sdom[mn[rt[v]]]) mn[v] = mn[rt[v
        ]];
    rt[v] = p;
    return x ? p : mn[v];
  void dfs(int v) {
    vis[v] = id, rev[id] = v;
    rt[id] = mn[id] = sdom[id] = id, id++;
    for (int u : adj[v]) {
      if (vis[u] == -1) dfs(u), par[vis[u]] = vis[v];
      radj[vis[u]].pb(vis[v]);
    }
  void build(int s) {
    dfs(s);
    for (int i = id - 1; ~i; --i) {
      for (int u : radj[i]) {
        sdom[i] = min(sdom[i], sdom[query(u, false)]);
      if (i) bucket[sdom[i]].pb(i);
      for (int u : bucket[i]) {
        int p = query(u, false);
        dom[u] = sdom[p] == i ? i : p;
      if (i) rt[i] = par[i];
    }
    vector <int> res(n, -1);
    for (int i = 1; i < id; ++i) {</pre>
      if (dom[i] != sdom[i]) dom[i] = dom[dom[i]];
    for (int i = 1; i < id; ++i) res[rev[i]] = rev[dom[</pre>
        i]];
    res[s] = s;
    dom = res;
};
```

# 5 String

### 5.1 Aho-Corasick Automaton

```
struct AC {
  int ch[N][26], to[N][26], fail[N], sz;
  vector <int> g[N];
  int cnt[N];
  AC () {sz = 0, extend();}
  void extend() {fill(ch[sz], ch[sz] + 26, 0), sz++;}
  int nxt(int u, int v) {
    if (!ch[u][v]) ch[u][v] = sz, extend();
    return ch[u][v];
  int insert(string s) {
    int now = 0;
    for (char c : s) now = nxt(now, c - 'a');
    cnt[now]++;
    return now:
  void build_fail() {
    queue <int> q;
```

```
for (int i = 0; i < 26; ++i) if (ch[0][i]) {</pre>
       q.push(ch[0][i]);
       g[0].push_back(ch[0][i]);
     while (!q.empty()) {
       int v = q.front(); q.pop();
       for (int j = 0; j < 26; ++j) {</pre>
         to[v][j] = ch[v][j] ? v : to[fail[v]][j];
       for (int i = 0; i < 26; ++i) if (ch[v][i]) {</pre>
         int u = ch[v][i], k = fail[v];
         while (k && !ch[k][i]) k = fail[k];
         if (ch[k][i]) k = ch[k][i];
         fail[u] = k;
         cnt[u] += cnt[k], g[k].push_back(u);
         q.push(u);
       }
    }
  int match(string &s) {
     int now = 0, ans = 0;
     for (char c : s) {
       now = to[now][c - 'a'];
if (ch[now][c - 'a']) now = ch[now][c - 'a'];
       ans += cnt[now]:
     return ans;
  }
};
```

# 5.2 KMP Algorithm

```
vector <int> build fail(string s) {
  vector <int> f(s.length() + 1, 0);
  int k = 0;
  for (int i = 1; i < s.length(); ++i) {
  while (k && s[k] != s[i]) k = f[k];</pre>
    if (s[k] == s[i]) k++;
    f[i + 1] = k;
  return f:
int match(string s, string t) {
  vector <int> f = build_fail(t);
  int k = 0, ans = 0;
  for (int i = 0; i < s.length(); ++i) {</pre>
    while (k && s[i] != t[k]) k = f[k];
    if (s[i] == t[k]) k++;
    if (k == t.length()) ans++, k = f[k];
  return ans;
}
```

# 5.3 Z Algorithm

# 5.4 Manacher

```
vector <int> manacher(string &s) {
   string t = "^#";
   for (char c : s) t += c, t += '#';
   t += '&';
   int n = t.length();
   vector <int> r(n, 0);
   int C = 0, R = 0;
   for (int i = 1; i < n - 1; ++i) {
     int mirror = 2 * C - i;
     r[i] = (i < R ? min(r[mirror], R - i) : 0);
   while (t[i - 1 - r[i]] == t[i + 1 + r[i]]) r[i]++;</pre>
```

```
if (i + r[i] > R) R = i + r[i], C = i;
}
return r;
}
```

# 5.5 Suffix Array

```
int sa[N], tmp[2][N], c[N], rk[N], lcp[N];
void buildSA(string s) {
  int *x = tmp[0], *y = tmp[1], m = 256, n = s.length()
  for (int i = 0; i < m; ++i) c[i] = 0;</pre>
  for (int i = 0; i < n; ++i) c[x[i] = s[i]]++;</pre>
  for (int i = 1; i < m; ++i) c[i] += c[i - 1];
  for (int i = n - 1; ~i; --i) sa[--c[x[i]]] = i;
  for (int k = 1; k < n; k <<= 1) {
    for (int i = 0; i < m; ++i) c[i] = 0;</pre>
    for (int i = 0; i < n; ++i) c[x[i]]++;</pre>
    for (int i = 1; i < m; ++i) c[i] += c[i - 1];</pre>
    int p = 0;
    for (int i = n - k; i < n; ++i) y[p++] = i;
    for (int i = 0; i < n; ++i) if (sa[i] >= k) y[p++]
         = sa[i] - k;
    for (int i = n - 1; ~i; --i) sa[--c[x[y[i]]]] = y[i
    y[sa[0]] = p = 0;
    for (int i = 1; i < n; ++i) {
      int a = sa[i], b = sa[i - 1];
      if (!(x[a] == x[b] \&\& a + k < n \&\& b + k < n \&\& x)
           [a + k] == x[b + k]) p++;
      y[sa[i]] = p;
    if (n == p + 1) break;
    swap(x, y), m = p + 1;
void buildLCP(string s) {
   // Lcp[i] = LCP(sa[i - 1], sa[i])
  // lcp(i, j) = min(lcp[rk[i] + 1], lcp[rk[i] + 2],
      ..., lcp[rk[j]])
  int n = s.length(), val = 0;
  for (int i = 0; i < n; ++i) rk[sa[i]] = i;</pre>
  for (int i = 0; i < n; ++i) {</pre>
    if (!rk[i]) lcp[rk[i]] = 0;
    else {
      if (val) val--
      int p = sa[rk[i] - 1];
      while (val + i < n && val + p < n && s[val + i]
           == s[val + p]) val++;
      lcp[rk[i]] = val;
  }
}
```

# 5.6 SAIS

```
namespace sfx {
bool _t[N * 2];
int SA[N * 2], H[N], RA[N];
int _s[N * 2], _c[N * 2], x[N], _p[N],
                                         _q[N * 2];
void pre(int *sa, int *c, int n, int z) {
  fill_n(sa, n, 0), copy_n(c, z, x);
void induce(int *sa, int *c, int *s, bool *t, int n,
    int z) {
  copy_n(c, z - 1, x + 1);
  for (int i = 0; i < n; ++i) if (sa[i] && !t[sa[i] -</pre>
      1]) sa[x[s[sa[i] - 1]]++] = sa[i] - 1;
  copy_n(c, z, x);
for (int i = n - 1; i >= 0; --i) if (sa[i] && t[sa[i]
        - 1]) sa[--x[s[sa[i] - 1]]] = sa[i] - 1;
void sais(int *s, int *sa, int *p, int *q, bool *t, int
     *c, int n, int z) {
  bool uniq = t[n - 1] = true;
  int nn = 0, nmxz = -1, *nsa = sa + n, *ns = s + n,
      last = -1;
  fill_n(c, z, 0);
  for (int i = 0; i < n; ++i) uniq &= ++c[s[i]] < 2;</pre>
  partial_sum(c, c + z, c);
  if (uniq) {
    for (int i = 0; i < n; ++i) sa[--c[s[i]]] = i;</pre>
```

```
return:
  for (int i = n - 2; i >= 0; --i)
    t[i] = (s[i] == s[i + 1] ? t[i + 1] : s[i] < s[i +
         1]);
  pre(sa, c, n, z);
  for (int i = 1; i <= n - 1; ++i)</pre>
    if (t[i] && !t[i - 1])
      sa[--x[s[i]]] = p[q[i] = nn++] = i;
  induce(sa, c, s, t, n, z);
for (int i = 0; i < n; ++i)
    if (sa[i] && t[sa[i]] && !t[sa[i] - 1]) {
      bool neq = last < 0 \mid | !equal(s + sa[i], s + p[q[
           sa[i]] + 1], s + last);
      ns[q[last = sa[i]]] = nmxz += neq;
  sais(ns, nsa, p + nn, q + n, t + n, c + z, nn, nmxz + m
        1);
  pre(sa, c, n, z);
  for (int i = nn - 1; i >= 0; --i)
    sa[--x[s[p[nsa[i]]]]] = p[nsa[i]];
  induce(sa, c, s, t, n, z);
vector<int> build(int *s, int n) {
 copy_n(s, n, _s), _s[n] = 0;
sais(_s, SA, _p, _q, _t, _c, n + 1, 256);
vector <int> sa(n);
  for (int i = 0; i < n; ++i)</pre>
    sa[i] = SA[i + 1];
  return sa;
```

#### 5.7 Suffix Automaton

```
struct SAM {
  int ch[N][26], len[N], link[N], pos[N], cnt[N], sz;
  // node -> strings with the same endpos set
  // length in range [len(link) + 1, len]
  // node's endpos set -> pos in the subtree of node
  // link -> longest suffix with different endpos set
  // len -> longest suffix
  // pos -> end position
// cnt -> size of endpos set
  SAM () \{len[0] = 0, link[0] = -1, pos[0] = 0, cnt[0] \}
      = 0, sz = 1;
  void build(string s) {
    int last = 0;
    for (int i = 0; i < s.length(); ++i) {</pre>
      char c = s[i];
      int cur = sz++;
      len[cur] = len[last] + 1, pos[cur] = i + 1;
      int p = last;
      while (\sim p \&\& !ch[p][c - 'a']) ch[p][c - 'a'] =
           cur, p = link[p];
      if (p == -1) {
        link[cur] = 0;
      } else {
        int q = ch[p][c - 'a'];
        if (len[p] + 1 == len[q]) {
          link[cur] = q;
        } else {
          int nxt = sz++;
          len[nxt] = len[p] + 1, link[nxt] = link[q],
               pos[nxt] = 0;
          for (int j = 0; j < 26; ++j) ch[nxt][j] = ch[
               q][j];
          while (~p && ch[p][c - 'a'] == q) ch[p][c - '
    a'] = nxt, p = link[p];
          link[q] = link[cur] = nxt;
        }
      cnt[cur]++;
      last = cur;
    vector <int> p(sz);
    iota(all(p), 0);
    sort(all(p), [&](int i, int j) {return len[i] > len
         [j];});
    for (int i = 0; i < sz; ++i) cnt[link[p[i]]] += cnt</pre>
         [p[i]];
  }
```

# 5.8 Minimum Rotation

} sam;

```
string rotate(const string &s) {
  int n = s.length();
  string t = s + s;
  int i = 0, j = 1;
  while (i < n && j < n) {
   int k = 0;
  while (k < n && t[i + k] == t[j + k]) ++k;
   if (t[i + k] <= t[j + k]) j += k + 1;
   else i += k + 1;
   if (i == j) ++j;
  }
  int pos = (i < n ? i : j);
  return t.substr(pos, n);
}</pre>
```

# 5.9 Palindrome Tree

```
struct PAM {
  int ch[N][26], cnt[N], fail[N], len[N], sz;
  // 0 -> even root, 1 -> odd root
  PAM (string _s) : s(_s) {
    sz = 0;
    extend(), extend();
    len[0] = 0, fail[0] = 1, len[1] = -1;
    int lst = 1;
    for (int i = 0; i < s.length(); ++i) {</pre>
      while (s[i - len[lst] - 1] != s[i]) lst = fail[
           lst1:
      if (!ch[lst][s[i] - 'a']) {
        int idx = extend();
        len[idx] = len[lst] + 2;
        int now = fail[lst];
        while (s[i - len[now] - 1] != s[i]) now = fail[
            now];
        fail[idx] = ch[now][s[i] - 'a'];
        ch[lst][s[i] - 'a'] = idx;
      lst = ch[lst][s[i] - 'a'], cnt[lst]++;
    }
  void build_count() {
    for (int i = sz - 1; i > 1; --i)
      cnt[fail[i]] += cnt[i];
  int extend() {
    fill(ch[sz], ch[sz] + 26, 0), sz++;
    return sz - 1;
};
```

#### 5.10 Main Lorentz

```
int to_left[N], to_right[N];
vector <array <int, 3>> rep; // l, r, len.
// substr(l ~ r, len * 2) are tandem
void findRep(string &s, int 1, int r) {
  if (r - 1 == 1) return;
  int m = 1 + r >> 1;
  findRep(s, l, m), findRep(s, m, r);
  string sl = s.substr(1, m - 1), sr = s.substr(m, r - 1)
      m);
  vector <int> Z = buildZ(sr + "#" + sl);
  for (int i = 1; i < m; ++i) to_right[i] = Z[r - m + 1</pre>
       + i - 1];
  reverse(all(sl));
  Z = buildZ(sl);
  for (int i = 1; i < m; ++i) to_left[i] = Z[m - i -</pre>
      1];
  reverse(all(sl));
  for (int i = 1; i + 1 < m; ++i) {</pre>
    int k1 = to_left[i], k2 = to_right[i + 1], len = m
         - i - 1;
    if (k1 < 1 || k2 < 1 || len < 2) continue;</pre>
    int tl = max(1, len - k2), tr = min(len - 1, k1);
    if (tl <= tr) rep.pb({i + 1 - tr, i + 1 - tl, len})</pre>
```

```
Z = buildZ(sr);
   for (int i = m; i < r; ++i) to_right[i] = Z[i - m];</pre>
  reverse(all(sl)), reverse(all(sr));
Z = buildZ(sl + "#" + sr);
   for (int i = m; i < r; ++i) to_left[i] = Z[m - l + 1</pre>
       + r - i - 1];
  reverse(all(sl)), reverse(all(sr));
   for (int i = m; i + 1 < r; ++i) {
     int k1 = to_left[i], k2 = to_right[i + 1], len = i
          - m + 1;
     if (k1 < 1 || k2 < 1 || len < 2) continue;</pre>
     int tl = max(len - k2, 1), tr = min(len - 1, k1);
     if (tl <= tr) rep.pb({i + 1 - len - tr, i + 1 - len
           - tl, len});
  Z = buildZ(sr + "#" + sl);
  for (int i = 1; i < m; ++i) {
  if (Z[r - m + 1 + i - 1] >= m - i) {
       rep.pb({i, i, m - i});
  }
}
```

# 6 Math

#### 6.1 Fraction\*

```
struct fraction {
  11 n, d;
  fraction(const ll _n=0, const ll _d=1): n(_n), d(_d)
    11 t = gcd(n, d);
    n /= t, d /= t;
    if (d < 0) n = -n, d = -d;
  fraction operator-() const
  { return fraction(-n, d); }
  fraction operator+(const fraction &b) const
  { return fraction(n * b.d + b.n * d, d * b.d); }
  fraction operator-(const fraction &b) const
  { return fraction(n * b.d - b.n * d, d * b.d); }
  fraction operator*(const fraction &b) const
  { return fraction(n * b.n, d * b.d); }
  fraction operator/(const fraction &b) const
  { return fraction(n * b.d, d * b.n); }
  void print() {
    cout << n;
    if (d != 1) cout << "/" << d;</pre>
};
```

#### 6.2 Miller Rabin / Pollard Rho

```
11 mul(11 x, 11 y, 11 p) {return (x * y - (11)((long
double)x / p * y) * p + p) % p;}
vector<11> chk = {2, 325, 9375, 28178, 450775, 9780504,
     1795265022};
ll Pow(ll a, ll b, ll n) {ll res = 1; for (; b; b >>=
    1, a = mul(a, a, n)) if (b \& 1) res = mul(res, a, n)
    ); return res;}
bool check(ll a, ll d, int s, ll n) {
  a = Pow(a, d, n);
  if (a <= 1) return 1;</pre>
  for (int i = 0; i < s; ++i, a = mul(a, a, n)) {</pre>
    if (a == 1) return 0;
    if (a == n - 1) return 1;
  }
  return 0;
bool IsPrime(ll n) {
  if (n < 2) return 0;
  if (n % 2 == 0) return n == 2;
  \hat{d} = n - 1, s = 0;
  while (d % 2 == 0) d >>= 1, ++s;
  for (ll i : chk) if (!check(i, d, s, n)) return 0;
  return 1;
const vector<ll> small = {2, 3, 5, 7, 11, 13, 17, 19};
11 FindFactor(ll n) {
  if (IsPrime(n)) return 1;
  for (ll p : small) if (n % p == 0) return p;
```

```
11 x, y = 2, d, t = 1;
auto f = [&](11 a) {return (mul(a, a, n) + t) % n;};
  for (int 1 = 2; ; 1 <<= 1) {
    x = y;
    int m = min(1, 32);
    for (int i = 0; i < 1; i += m) {</pre>
       d = 1;
       for (int j = 0; j < m; ++j) {
         y = f(y), d = mul(d, abs(x - y), n);
       il g = __gcd(d, n);
if (g == n) {
         1 = 1, y = 2, ++t;
         break;
       if (g != 1) return g;
    }
  }
map <11, int> res;
void PollardRho(ll n) {
  if (n == 1) return;
  if (IsPrime(n)) return ++res[n], void(0);
  11 d = FindFactor(n);
  PollardRho(n / d), PollardRho(d);
}
```

#### 6.3 Ext GCD

```
//a * p.first + b * p.second = gcd(a, b)
pair<11, 11> extgcd(11 a, 11 b) {
    pair<11, 11> res;
    if (a < 0) {
       res = extgcd(-a, b);
       res.first *= -1;
       return res;
    }
    if (b < 0) {
       res = extgcd(a, -b);
       res.second *= -1;
       return res;
    }
    if (b == 0) return {1, 0};
    res = extgcd(b, a % b);
    return {res.second, res.first - res.second * (a / b)
       };
}</pre>
```

#### 6.4 PiCount

```
const int V = 10000000, N = 100, M = 100000;
vector<int> primes;
bool isp[V];
 int small_pi[V], dp[N][M];
 void sieve(int x){
   for(int i = 2; i < x; ++i) isp[i] = true;</pre>
   isp[0] = isp[1] = false;
   for(int i = 2; i * i < x; ++i) if(isp[i]) for(int j =</pre>
   i * i; j < x; j += i) isp[j] = false;
for(int i = 2; i < x; ++i) if(isp[i]) primes.</pre>
       push_back(i);
void init(){
   sieve(V);
   small_pi[0] = 0;
   for(int i = 1; i < V; ++i) small_pi[i] = small_pi[i -</pre>
        1] + isp[i];
   for(int i = 0; i < M; ++i) dp[0][i] = i;</pre>
   for(int i = 1; i < N; ++i) for(int j = 0; j < M; ++j)</pre>
        dp[i][j] = dp[i - 1][j] - dp[i - 1][j / primes[i]
         - 111:
11 phi(ll n, int a){
  if(!a) return n;
   if(n < M && a < N) return dp[a][n];</pre>
   if(primes[a - 1] > n) return 1;
   if(((11)primes[a - 1]) * primes[a - 1] >= n && n < V)
        return small_pi[n] - a + 1;
   11 de = phi(n, a - 1) - phi(n / primes[a - 1], a - 1)
   return de;
}
```

```
1l PiCount(1l n){
    if(n < V) return small_pi[n];
    int s = sqrt(n + 0.5), y = cbrt(n + 0.5), a =
        small_pi[y];
    ll res = phi(n, a) + a - 1;
    for(; primes[a] <= s; ++a) res -= max(PiCount(n /
        primes[a]) - PiCount(primes[a]) + 1, 0ll);
    return res;
}</pre>
```

#### 6.5 Linear Function Mod Min

```
11 topos(11 x, 11 m) {x %= m; if (x < 0) x += m; return
//min value of ax + b \pmod{m} for x \in [0, n - 1]. O(
    Log m)
11 min_rem(ll n, ll m, ll a, ll b) {
  a = topos(a, m), b = topos(b, m);
  for (ll g = __gcd(a, m); g > 1;) return g * min_rem(n
        m / g, a / g, b / g) + (b % g);
  for (11 nn, nm, na, nb; a; n = nn, m = nm, a = na, b
      = nb) {
    if (a <= m - a) {
    nn = (a * (n - 1) + b) / m;</pre>
      if (!nn) break;
      nn += (b < a);
      nm = a, na = topos(-m, a);
      nb = b < a ? b : topos(b - m, a);
    } else {
      ll lst = b - (n - 1) * (m - a);
      if (lst >= 0) {b = lst; break;}
      nn = -(1st / m) + (1st % m < -a) + 1;
      nm = m - a, na = m % (m - a), nb = b % (m - a);
    }
  }
  return b;
//min value of ax + b \pmod{m} for x \in [0, n - 1],
    also return \min x to get the value. O(\log m)
//{value, x}
pair<ll, ll> min_rem_pos(ll n, ll m, ll a, ll b) {
  a = topos(a, m), b = topos(b, m);
  11 mn = min_rem(n, m, a, b), g = __gcd(a, m);
  //ax = (mn - b) \pmod{m}
  11 x = (extgcd(a, m).first + m) * ((mn - b + m) / g)
      % (m / g);
  return {mn, x};
}
```

#### 6.6 Determinant

```
11 Det(vector <vector <11>>> a) {
  int n = a.size();
  ll det = 1;
  for (int i = 0; i < n; ++i) {
    if (!a[i][i]) {
      det = -det;
      if (det < 0) det += mod;</pre>
      for (int j = i + 1; j < n; ++j) if (a[j][i]) {</pre>
        swap(a[j], a[i]);
        break;
      if (!a[i][i]) return 0;
    det = det * a[i][i] % mod;
    ll mul = mpow(a[i][i], mod - 2);
    for (int j = 0; j < n; ++j) a[i][j] = a[i][j] * mul</pre>
          % mod;
    for (int j = 0; j < n; ++j) if (i ^ j) {</pre>
      11 mul = a[j][i];
      for (int k = 0; k < n; ++k) {
        a[j][k] -= a[i][k] * mul % mod;
        if (a[j][k] < 0) a[j][k] += mod;</pre>
      }
    }
  return det;
```

# 6.7 Floor Sum

# 6.8 Quadratic Residue

```
int Jacobi(int a, int m) {
  int s = 1;
  for (; m > 1; ) {
    a %= m;
    if (a == 0) return 0;
    const int r = __builtin_ctz(a);
    if ((r \& 1) \& \& ((m + 2) \& 4)) s = -s;
    a >>= r;
    if (a \& m \& 2) s = -s;
    swap(a, m);
  }
  return s;
int QuadraticResidue(int a, int p) {
  if (p == 2) return a & 1;
  const int jc = Jacobi(a, p);
  if (jc == 0) return 0;
  if (jc == -1) return -1;
  int b, d;
  for (; ; ) {
    b = rand() % p;
d = (1LL * b * b + p - a) % p;
    if (Jacobi(d, p) == -1) break;
  int f0 = b, f1 = 1, g0 = 1, g1 = 0, tmp;
  for (int e = (1LL + p) >> 1; e; e >>= 1) {
    if (e & 1) {
      tmp = (1LL * g0 * f0 + 1LL * d * (1LL * g1 * f1 %
           p)) % p;
      g1 = (1LL * g0 * f1 + 1LL * g1 * f0) % p;
      g0 = tmp;
    tmp = (1LL * f0 * f0 + 1LL * d * (1LL * f1 * f1 % p
        )) % p;
    f1 = (2LL * f0 * f1) % p;
    f0 = tmp;
  }
  return g0;
```

# 6.9 Simplex

```
struct Simplex { // O-based
  using T = long double;
  static const int N = 410, M = 30010;
  const T eps = 1e-7;
  int n, m;
  int Left[M], Down[N];
  // Ax <= b, max c^T x
  // result : v, xi = sol[i]
  T a[M][N], b[M], c[N], v, sol[N];
  bool eq(T a, T b) {return fabs(a - b) < eps;}</pre>
  bool ls(T a, T b) {return a < b && !eq(a, b);}</pre>
  void init(int _n, int _m) {
    n = _n, m = _m, v = 0;
    for (int i = 0; i < m; ++i) for (int j = 0; j < n;
         ++j) a[i][j] = 0;
    for (int i = 0; i < m; ++i) b[i] = 0;</pre>
    for (int i = 0; i < n; ++i) c[i] = sol[i] = 0;</pre>
  void pivot(int x, int y) {
    swap(Left[x], Down[y]);
    T k = a[x][y]; a[x][y] = 1;
    vector <int> nz;
    for (int i = 0; i < n; ++i) {</pre>
      a[x][i] /= k;
      if (!eq(a[x][i], 0)) nz.push_back(i);
```

```
b[x] /= k;
    for (int i = 0; i < m; ++i) {</pre>
      if (i == x || eq(a[i][y], 0)) continue;
      k = a[i][y], a[i][y] = 0;
b[i] -= k * b[x];
      for (int j : nz) a[i][j] -= k * a[x][j];
    if (eq(c[y], 0)) return;
    k = c[y], c[y] = 0, v += k * b[x];
    for (int i : nz) c[i] -= k * a[x][i];
  }
  // 0: found solution, 1: no feasible solution, 2:
       unbounded
  int solve() {
    for (int i = 0; i < n; ++i) Down[i] = i;</pre>
    for (int i = 0; i < m; ++i) Left[i] = n + i;</pre>
    while (1) {
      int x = -1, y = -1;
       for (int i = 0; i < m; ++i) if (ls(b[i], 0) && (x</pre>
            == -1 \mid \mid b[i] < b[x])) x = i;
       if (x == -1) break;
      for (int i = 0; i < n; ++i) if (ls(a[x][i], 0) &&</pre>
             (y == -1 \mid | a[x][i] < a[x][y])) y = i;
       if (y == -1) return 1;
      pivot(x, y);
    while (1) {
       int x = -1, y = -1;
       for (int i = 0; i < n; ++i) if (ls(0, c[i]) && (y
            == -1 || c[i] > c[y])) y = i;
       if (y == -1) break;
      for (int i = 0; i < m; ++i) if (ls(0, a[i][y]) &&
    (x == -1 || b[i] / a[i][y] < b[x] / a[x][y])</pre>
           ) x = i;
      if (x == -1) return 2;
      pivot(x, y);
    for (int i = 0; i < m; ++i) if (Left[i] < n) sol[</pre>
         Left[i]] = b[i];
    return 0;
  }
};
```

#### 6.10 Berlekamp Massey

```
vector <11> BerlekampMassey(vector <11> a) {
 // find min |c| such that a_n = sum c_j * a_{n - j - j}
      1}, 0-based
  // O(N^2), if |c| = k, |a| >= 2k sure correct
  auto f = [&](vector<11> v, 11 c) {
    for (11 &x : v) x = mul(x, c);
 };
 vector <11> c, best;
  int pos = 0, n = a.size();
  for (int i = 0; i < n; ++i) {
    ll error = a[i];
    for (int j = 0; j < c.size(); ++j) error = sub(
        error, mul(c[j], a[i - 1 - j]));
    if (error == 0) continue;
    11 inv = mpow(error, mod - 2);
    if (c.empty()) {
      c.resize(i + 1);
      pos = i;
      best.pb(inv);
    } else {
      vector <1l> fix = f(best, error);
      fix.insert(fix.begin(), i - pos - 1, 0);
      if (fix.size() >= c.size()) {
        best = f(c, sub(0, inv));
        best.insert(best.begin(), inv);
        pos = i:
        c.resize(fix.size());
      for (int j = 0; j < fix.size(); ++j) c[j] = add(c</pre>
          [j], fix[j]);
   }
 }
  return c;
```

# 6.11 Linear Programming Construction

Standard form: maximize  $\mathbf{c}^T\mathbf{x}$  subject to  $A\mathbf{x} \leq \mathbf{b}$  and  $\mathbf{x} \geq 0$ . Dual LP: minimize  $\mathbf{b}^T\mathbf{y}$  subject to  $A^T\mathbf{y} \geq \mathbf{c}$  and  $\mathbf{y} \geq 0$ .  $\bar{\mathbf{x}}$  and  $\bar{\mathbf{y}}$  are optimal if and only if for all  $i \in [1,n]$ , either  $\bar{x}_i = 0$  or  $\sum_{j=1}^m A_{ji}\bar{y}_j = c_i$  holds and for all  $i \in [1,m]$  either  $\bar{y}_i = 0$  or  $\sum_{j=1}^n A_{ij}\bar{x}_j = b_j$  holds.

- 1. In case of minimization, let  $c_i^\prime = -c_i$
- 2.  $\sum_{1 \le i \le n} A_{ji} x_i \ge b_j \to \sum_{1 \le i \le n} A_{ji} x_i \le -b_j$
- 3.  $\sum_{1 \le i \le n}^{-} A_{ji} x_i = b_j$ 
  - $\begin{array}{ll} \bullet & \sum_{1 \leq i \leq n} A_{ji} x_i \leq b_j \\ \bullet & \sum_{1 \leq i \leq n} A_{ji} x_i \geq b_j \end{array}$
- 4. If  $x_i$  has no lower bound, replace  $x_i$  with  $x_i x_i'$

#### 6.12 Euclidean

- $m = \lfloor \frac{an+b}{c} \rfloor$
- Time complexity:  $O(\log n)$

$$\begin{split} f(a,b,c,n) &= \sum_{i=0}^n \lfloor \frac{ai+b}{c} \rfloor \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)}{2} + \lfloor \frac{b}{c} \rfloor \cdot (n+1) \\ + f(a \mod c, b \mod c, c, n), & a \geq c \vee b \geq c \\ 0, & n < 0 \vee a = 0 \\ nm - f(c, c - b - 1, a, m - 1), & \text{otherwise} \end{cases} \end{split}$$

$$\begin{split} g(a,b,c,n) &= \sum_{i=0}^n i \lfloor \frac{ai+b}{c} \rfloor \\ &= \begin{cases} \left\lfloor \frac{a}{c} \right\rfloor \cdot \frac{n(n+1)(2n+1)}{6} + \left\lfloor \frac{b}{c} \right\rfloor \cdot \frac{n(n+1)}{2} \\ +g(a \mod c, b \mod c, c, n), & a \geq c \lor b \geq c \\ 0, & n < 0 \lor a = 0 \\ \frac{1}{2} \cdot (n(n+1)m - f(c, c-b-1, a, m-1) \\ -h(c, c-b-1, a, m-1)), & \text{otherwise} \end{cases} \end{split}$$

$$\begin{split} h(a,b,c,n) &= \sum_{i=0}^n \lfloor \frac{ai+b}{c} \rfloor^2 \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor^2 \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor^2 \cdot (n+1) \\ &+ \lfloor \frac{a}{c} \rfloor \cdot \lfloor \frac{b}{c} \rfloor \cdot n(n+1) \\ &+ h(a \bmod c, b \bmod c, c, n) \\ &+ 2 \lfloor \frac{a}{c} \rfloor \cdot g(a \bmod c, b \bmod c, c, n) \\ &+ 2 \lfloor \frac{b}{c} \rfloor \cdot f(a \bmod c, b \bmod c, c, n), & a \geq c \lor b \geq c \\ 0, & n < 0 \lor a = 0 \\ nm(m+1) - 2g(c, c-b-1, a, m-1) \\ &- 2f(c, c-b-1, a, m-1) - f(a, b, c, n), & \text{otherwise} \end{cases} \end{split}$$

#### 6.13 Theorem

• Kirchhoff's Theorem

Denote L be a  $n\times n$  matrix as the Laplacian matrix of graph G, where  $L_{ii}=d(i)$ ,  $L_{ij}=-c$  where c is the number of edge (i,j) in G.

- The number of undirected spanning in G is  $|\det(\tilde{L}_{11})|$  .
- The number of directed spanning tree rooted at r in G is  $|\det(\tilde{L}_{rr})|$ .
- Tutte's Matrix

Let D be a  $n \times n$  matrix, where  $d_{ij} = x_{ij}$  ( $x_{ij}$  is chosen uniformly at random) if i < j and  $(i,j) \in E$ , otherwise  $d_{ij} = -d_{ji}$ .  $\frac{rank(D)}{2}$  is the maximum matching on G.

- Cayley's Formula
  - Given a degree sequence  $d_1,d_2,\ldots,d_n$  for each  $\mathit{labeled}$  vertices, there are

$$\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$$

spanning trees.

- Let  $T_{n,k}$  be the number of *labeled* forests on n vertices with k components, such that vertex  $1,2,\ldots,k$  belong to different components. Then  $T_{n,k}=kn^{n-k-1}$ .
- Erdős-Gallai Theorem

A sequence of non-negative integers  $d_1 \geq d_2 \geq \ldots \geq d_n$  can be represented as the degree sequence of a finite simple graph on n vertices if and only if  $d_1+d_2+\ldots+d_n$  is even and

$$\sum_{i=1}^k d_i \le k(k-1) + \sum_{i=k+1}^n \min(d_i, k)$$

holds for all  $1 \leq k \leq n$ .

• Burnside's Lemma

Let X be a set and G be a group that acts on X. For  $g \in G$ , denote by  $X^g$  the elements fixed by g:

$$X^g = \{ x \in X \mid gx \in X \}$$

Then

$$|X/G| = \frac{1}{|G|} \sum_{g \in G} |X^g|.$$

• Gale-Ryser theorem

A pair of sequences of nonnegative integers  $a_1 \geq \cdots \geq a_n$  and  $b_1,\dots,b_n$  is bigraphic if and only if  $\sum_{i=1}^n a_i = \sum_{i=1}^n b_i$  and  $\sum_{i=1}^k a_i \leq$ 

$$\sum_{i=1}^n \min(b_i,k)$$
 holds for every  $1 \leq k \leq n$  .

• Fulkerson-Chen-Anstee theorem

A sequence  $(a_1,b_1),\ldots,(a_n,b_n)$  of nonnegative integer pairs with  $a_1 \geq \cdots \geq a_n$  is digraphic if and only if  $\sum^n a_i = \sum^n b_i$  and

$$\sum_{i=1}^k a_i \leq \sum_{i=1}^k \min(b_i,k-1) + \sum_{i=k+1}^n \min(b_i,k) \text{ holds for every } 1 \leq k \leq n.$$

• Möbius inversion formula

- 
$$f(n) = \sum_{d|n} g(d) \Leftrightarrow g(n) = \sum_{d|n} \mu(d) f(\frac{n}{d})$$
  
-  $f(n) = \sum_{n|d} g(d) \Leftrightarrow g(n) = \sum_{n|d} \mu(\frac{d}{n}) f(d)$ 

- Spherical cap
  - A portion of a sphere cut off by a plane. r: sphere radius, a: radius of the base of the cap, h: height of the cap,  $\theta$ :  $\arcsin(a/r)$ .  $Volume = \pi h^2(3r-h)/3 = \pi h(3a^2+h^2)/6 = \pi r^3(2+\cos\theta)(1-\sin\theta)$

  - $\cos\theta)^2/3. \label{eq:cos} \text{Area} = 2\pi r h = \pi (a^2 + h^2) = 2\pi r^2 (1 \cos\theta).$
- Chinese Remainder Theorem
  - $x \equiv a_i \pmod{m_i}$
  - $M = \prod m_i, M_i = M/m_i$
  - $t_i M_i \equiv 1 \pmod{m_i}$
  - $x = \sum a_i t_i M_i \pmod{M}$

### 6.14 Estimation

- The number of divisors of n is at most around  $100\ {\rm for}\ n<5e4$  ,
- 500 for n<1e7, 2000 for n<1e10, 200000 for n<1e19. The number of ways of writing n as a sum of positive integers, disregarding the order of the summands. 1,1,2,3,5,7,11,15,22,30for  $n=0\sim 9$ , 627 for n=20,  $\sim 2e5$  for n=50,  $\sim 2e8$  for
- Total number of partitions of n distinct elements: B(n)=1,1,2,5,15,52,203,877,4140,21147,115975,678570,4213597,27644437, 190899322, . . . .

# 6.15 General Purpose Numbers

Bernoulli numbers

$$\begin{array}{l} B_0-1, B_1^{\pm}=\pm\frac{1}{2}, B_2=\frac{1}{6}, B_3=0 \\ \sum_{j=0}^m \binom{m+1}{j} B_j=0 \text{, EGF is } B(x)=\frac{x}{e^x-1}=\sum_{n=0}^\infty B_n \frac{x^n}{n!} \,. \end{array}$$

$$S_m(n) = \sum_{k=1}^n k^m = \frac{1}{m+1} \sum_{k=0}^m {m+1 \choose k} B_k^+ n^{m+1-k}$$

ullet Stirling numbers of the second kind Partitions of n distinct elements into exactly k groups.

$$\begin{split} S(n,k) &= S(n-1,k-1) + kS(n-1,k), S(n,1) = S(n,n) = 1 \\ S(n,k) &= \frac{1}{k!} \sum_{i=0}^k (-1)^{k-i} {k \choose i} i^n \\ x^n &= \sum_{i=0}^n S(n,i)(x)_i \end{split}$$

• Pentagonal number theorem

$$\prod_{n=1}^{\infty} (1-x^n) = 1 + \sum_{k=1}^{\infty} (-1)^k \left( x^{k(3k+1)/2} + x^{k(3k-1)/2} \right)$$

• Catalan numbers

$$C_n^{(k)} = \frac{1}{(k-1)n+1} \binom{kn}{n}$$
 
$$C^{(k)}(x) = 1 + x[C^{(k)}(x)]^k$$

• Eulerian numbers

Number of permutations  $\pi \in S_n$  in which exactly k elements are greater than the previous element. k j:s s.t.  $\pi(j) > \pi(j+1)$ , k+1 j:s s.t.  $\pi(j) \geq j$ , k j:s s.t.  $\pi(j) > j$ .

$$E(n,k) = (n-k)E(n-1,k-1) + (k+1)E(n-1,k)$$

$$E(n,0) = E(n,n-1) = 1$$

$$E(n,k) = \sum_{j=0}^{k} (-1)^{j} {n+1 \choose j} (k+1-j)^{n}$$

# 6.16 Tips for Generating Funtion

• Ordinary Generating Function  $A(x) = \sum_{i \ge 0} a_i x^i$ 

```
- A(rx) \Rightarrow r^n a_n
- A(x) + B(x) \Rightarrow a_n + b_n
- A(x)B(x) \Rightarrow \sum_{i=0}^{n} a_i b_{n-i}
- A(x)^k \Rightarrow \sum_{i_1+i_2+\dots+i_k=n} a_{i_1} a_{i_2} \dots a_{i_k}
- xA(x)' \Rightarrow na_n
- \frac{A(x)}{1-x} \Rightarrow \sum_{i=0}^{n} a_i
```

• Exponential Generating Function  $A(x) = \sum_{i>0} \frac{a_i}{i!} x_i$ 

```
- A(x) + B(x) \Rightarrow a_n + b_n
- A^{(k)}(x) \Rightarrow a_{n+k}

- A(x)B(x) \Rightarrow \sum_{i=0}^{n} {n \choose i} a_i b_{n-i}
- A(x)^k \Rightarrow \sum_{i_1+i_2+\dots+i_k=n} \binom{n}{i_1,i_2,\dots,i_k} a_{i_1} a_{i_2} \dots a_{i_k}
- xA(x) \Rightarrow na_n
```

• Special Generating Function

- 
$$(1+x)^n = \sum_{i\geq 0} \binom{n}{i} x^i$$
  
-  $\frac{1}{(1-x)^n} = \sum_{i\geq 0} \binom{n}{i-1} x^i$ 

# **Polynomial**

# **Number Theoretic Transform**

```
// mul, add, sub, mpow
// ll -> int if too slow
struct NTT {
  11 w[N];
  NTT() {
    ll dw = mpow(G, (mod - 1) / N);
     w[0] = 1;
     for (int i = 1; i < N; ++i) w[i] = mul(w[i - 1], dw
  void operator()(vector<ll>& a, bool inv = false) { //
       0 <= a[i] < P
     int x = 0, n = a.size();
     for (int j = 1; j < n - 1; ++j) {
  for (int k = n >> 1; (x ^= k) < k; k >>= 1);
       if (j < x) swap(a[x], a[j]);</pre>
     for (int L = 2; L <= n; L <<= 1) {
       int dx = N / L, dl = L >> 1;
for (int i = 0; i < n; i += L) {
   for (int j = i, x = 0; j < i + dl; ++j, x += dx</pre>
            ll tmp = mul(a[j + dl], w[x]);
            a[j + dl] = sub(a[j], tmp);
            a[j] = add(a[j], tmp);
       }
     if (inv) {
       reverse(a.begin() + 1, a.end());
       ll invn = mpow(n, mod - 2);
       for (int i = 0; i < n; ++i) a[i] = mul(a[i], invn</pre>
  }
} ntt;
```

# 7.2 Primes

Prime	Root	Prime	Root
7681	17	167772161	3
12289	11	104857601	3
40961	3	985661441	3
65537	3	998244353	3
786433	10	1107296257	10
5767169	3	2013265921	31
7340033	3	2810183681	11
23068673	3	2885681153	3
469762049	3	605028353	3
2061584302081	7	1945555039024054273	5
2748779069441	3	9223372036737335297	3

## 7.3 Polynomial Operations

```
vector <ll> Mul(vector <ll> a, vector <ll> b, int bound
     = N) {
  int m = a.size() + b.size() - 1, n = 1;
  while (n < m) n <<= 1;</pre>
  a.resize(n), b.resize(n);
 ntt(a), ntt(b);
 vector <11> out(n);
  for (int i = 0; i < n; ++i) out[i] = mul(a[i], b[i]);</pre>
 ntt(out, true), out.resize(min(m, bound));
 return out;
vector <ll> Inverse(vector <ll> a) {
 // O(NlogN), a[0] != 0
 int n = a.size();
  vector <11> res(1, mpow(a[0], mod - 2));
  for (int m = 1; m < n; m <<= 1) {</pre>
    if (n < m * 2) a.resize(m * 2);</pre>
    vector <ll> v1(a.begin(), a.begin() + m * 2), v2 =
        res;
    v1.resize(m * 4), v2.resize(m * 4);
    ntt(v1), ntt(v2);
    for (int i = 0; i < m * 4; ++i) v1[i] = mul(mul(v1[</pre>
         i], v2[i]), v2[i]);
    ntt(v1, true);
    res.resize(m * 2);
    for (int i = 0; i < m; ++i) res[i] = add(res[i],</pre>
         res[i]);
    for (int i = 0; i < m * 2; ++i) res[i] = sub(res[i</pre>
         ], v1[i]);
 res.resize(n):
 return res;
pair <vector <ll>, vector <ll>> Divide(vector <ll> a,
    vector <ll> b) {
  // a = bQ + R, O(NLogN), b.back() != 0
 int n = a.size(), m = b.size(), k = n - m + 1;
  if (n < m) return {{0}, a};</pre>
 vector <11> ra = a, rb = b;
 reverse(all(ra)), ra.resize(k);
 reverse(all(rb)), rb.resize(k);
 vector <11> Q = Mul(ra, Inverse(rb), k);
 reverse(all(Q));
 vector <1l> res = Mul(b, Q), R(m - 1);
for (int i = 0; i < m - 1; ++i) R[i] = sub(a[i], res[</pre>
      i]);
 return {Q, R};
vector <1l> SqrtImpl(vector <1l> a) {
  if (a.empty()) return {0};
  int z = QuadraticResidue(a[0], mod), n = a.size();
  if (z == -1) return {-1};
  vector \langle 11 \rangle q(1, z);
  const int inv2 = (mod + 1) / 2;
  for (int m = 1; m < n; m <<= 1) {</pre>
    if (n < m * 2) a.resize(m * 2);</pre>
    q.resize(m * 2);
    vector <11> f2 = Mul(q, q, m * 2);
for (int i = 0; i < m * 2; ++i) f2[i] = sub(f2[i],</pre>
        a[i]);
    f2 = Mul(f2, Inverse(q), m * 2);
    for (int i = 0; i < m * 2; ++i) q[i] = sub(q[i],</pre>
        mul(f2[i], inv2));
  q.resize(n);
  return q;
vector <11> Sqrt(vector <11> a) {
  // O(NlogN), return {-1} if not exists
  int n = a.size(), m = 0;
 while (m < n && a[m] == 0) m++;</pre>
  if (m == n) return vector <11>(n);
  if (m & 1) return {-1};
  vector <ll> s = SqrtImpl(vector <ll>(a.begin() + m, a
      .end()));
  if (s[0] == -1) return {-1};
  vector <11> res(n);
  for (int i = 0; i < s.size(); ++i) res[i + m / 2] = s</pre>
      [i];
  return res;
```

```
vector <11> Derivative(vector <11> a) {
  int n = a.size();
  vector <ll> res(n - 1);
  for (int i = 0; i < n - 1; ++i) res[i] = mul(a[i +</pre>
      1], i + 1);
  return res;
vector <ll> Integral(vector <ll> a) {
  int n = a.size();
  vector \langle 11 \rangle res(n + 1);
  for (int i = 0; i < n; ++i) {
    res[i + 1] = mul(a[i], mpow(i + 1, mod - 2));
  return res;
}
vector <ll> Ln(vector <ll> a) {
  // O(NlogN), a[0] = 1
  int n = a.size();
  if (n == 1) return {0};
  vector <1l> d = Derivative(a);
  a.pop_back();
  return Integral(Mul(d, Inverse(a), n - 1));
vector <11> Exp(vector <11> a) {
  // O(NlogN), a[0] = 0
  int n = a.size();
  vector \langle 11 \rangle q(1, 1);
  a[0] = add(a[0], 1);
  for (int m = 1; m < n; m <<= 1) {
   if (n < m * 2) a.resize(m * 2);</pre>
    vector <ll> g(a.begin(), a.begin() + m * 2), h(all(
         a));
    h.resize(m * 2), h = Ln(h);
    for (int i = 0; i < m * 2; ++i) {</pre>
      g[i] = sub(g[i], h[i]);
    q = Mul(g, q, m * 2);
  }
  q.resize(n);
  return a:
vector <ll> Pow(vector <ll> a, ll k) {
  int n = a.size(), m = 0;
  vector <11> ans(n, 0);
  while (m < n \&\& a[m] == 0) m++;
  if (k \&\& m \&\& (k >= n | | k * m >= n)) return ans;
  if (m == n) return ans[0] = 1, ans;
  ll lead = m * k;
  vector <1l> b(a.begin() + m, a.end());
  11 base = mpow(b[0], k), inv = mpow(b[0], mod - 2);
  for (int i = 0; i < n - m; ++i) b[i] = mul(b[i], inv)</pre>
  b = Ln(b);
  for (int i = 0; i < n - m; ++i) b[i] = mul(b[i], k %</pre>
      mod);
  b = Exp(b);
  for (int i = lead; i < n; ++i) ans[i] = mul(b[i -</pre>
      lead], base);
  return ans;
vector <ll> Evaluate(vector <ll> a, vector <ll> x) {
  if (x.empty()) return {};
  int n = x.size();
  vector <vector <11>> up(n * 2);
  for (int i = 0; i < n; ++i) up[i + n] = {sub(0, x[i])}
      , 1};
  for (int i = n - 1; i > 0; --i) up[i] = Mul(up[i *
      2], up[i * 2 + 1]);
  vector <vector <11>> down(n * 2);
  down[1] = Divide(a, up[1]).second;
  for (int i = 2; i < n * 2; ++i) down[i] = Divide(down</pre>
      [i >> 1], up[i]).second;
  vector <11> y(n);
  for (int i = 0; i < n; ++i) y[i] = down[i + n][0];</pre>
  return v;
vector <1l> Interpolate(vector <1l> x, vector <1l> y) {
  int n = x.size();
  vector <vector <11>> up(n * 2);
  for (int i = 0; i < n; ++i) up[i + n] = {sub(0, x[i])}
      , 1};
```

#### 7.4 Fast Linear Recursion

```
11 FastLinearRecursion(vector <11> a, vector <11> c, 11
      k) {
  // a_n = sigma c_j * a_{n - j - 1}, 0-based
// O(NLogNLogK), |a| = |c|
  int n = a.size();
  if (k < n) return a[k];</pre>
  vector <ll> base(n + 1, 1);
  for (int i = 0; i < n; ++i) base[i] = sub(0, c[n - i</pre>
       - 1]);
  vector <11> poly(n);
  (n == 1 ? poly[0] = c[n - 1] : poly[1] = 1);
  auto calc = [&](vector <ll> p1, vector <ll> p2) {
    // O(n^2) bruteforce or O(nlogn) NTT
    return Divide(Mul(p1, p2), base).second;
  };
  vector \langle 11 \rangle res(n, 0); res[0] = 1;
  for (; k; k >>= 1, poly = calc(poly, poly)) {
    if (k & 1) res = calc(res, poly);
  11 \text{ ans} = 0;
  for (int i = 0; i < n; ++i) {</pre>
    (ans += res[i] * a[i]) %= mod;
  return ans;
```

#### 7.5 Fast Walsh Transform

```
void fwt(vector <int> &a) {
 // \ and : a[j] += x;
  //
         : a[j] -= x;
        : a[j ^ (1 << i)] += y;
         : a[j ^ (1 << i)] -= y;
 //
  // xor : a[j] = x - y, a[j ^ (1 << i)] = x + y;
          : a[j] = (x - y) / 2, a[j ^ (1 << i)] = (x + y)
      ) / 2;
 int n = __lg(a.size());
  for (int i = 0; i < n; ++i) {</pre>
    for (int j = 0; j < 1 << n; ++j) if (j >> i & 1) {
  int x = a[j ^ (1 << i)], y = a[j];</pre>
      // do something
 }
vector<int> subs_conv(vector<int> a, vector<int> b) {
 // c_i = sum_{j \& k = 0, j | k = i} a_j * b_k
 int n = __lg(a.size());
  vector<vector<int>> ha(n + 1, vector<int>(1 << n));</pre>
  vector<vector<int>> hb(n + 1, vector<int>(1 << n));</pre>
  vector<vector<int>> c(n + 1, vector<int>(1 << n));</pre>
  for (int i = 0; i < 1 << n; ++i) {</pre>
    ha[__builtin_popcount(i)][i] = a[i];
    hb[__builtin_popcount(i)][i] = b[i];
  for (int i = 0; i <= n; ++i) fwt(ha[i]),fwt(hb[i]);</pre>
  for (int i = 0; i <= n; ++i)</pre>
    for (int j = 0; i + j <= n; ++j)</pre>
      for (int k = 0; k < 1 << n; ++k)
        // mind overflow
        c[i + j][k] += ha[i][k] * hb[j][k];
  for (int i = 0; i <= n; ++i)</pre>
```

```
fwt(c[i], true);
vector <int> ans(1 << n);
for (int i = 0; i < 1 << n; ++i)
    ans[i] = c[__builtin_popcount(i)][i];
return ans;
}</pre>
```

# 8 Geometry

#### 8.1 Basic

```
const double eps = 1e-8, pi = acos(-1);
int sign(double x) \{return abs(x) \leftarrow eps ? 0 : (x > 0 ?
     1 : -1);}
struct Pt {
  double x, y;
  Pt (double _x, double _y) : x(_x), y(_y) {}
  Pt operator + (Pt o) {return Pt(x + o.x, y + o.y);}
  Pt operator - (Pt o) {return Pt(x - o.x, y - o.y);}
  Pt operator * (double k) {return Pt(x * k, y * k);}
  Pt operator / (double k) {return Pt (x / k, y / k);}
double operator * (Pt o) {return x * o.x + y * o.y;}
  double operator ^ (Pt o) {return x * o.y - y * o.x;}
};
struct Line {
  Pt a, b;
struct Cir {
  Pt o; double r;
double abs2(Pt o) {return o.x * o.x + o.y * o.y;}
double abs(Pt o) {return sqrt(abs2(o));}
int ori(Pt o, Pt a, Pt b) {return sign((o - a) ^ (o - b
    ));}
bool btw(Pt a, Pt b, Pt c) { // c on segment ab?
  return ori(a, b, c) == 0 && sign((c - a) * (c - b))
double area(Pt a, Pt b, Pt c) {return abs((a - b) ^ (a
    - c)) / 2;}
Pt unit(Pt o) {return o / abs(o);}
Pt rot(Pt a, double o) { // CCW
  double c = cos(o), s = sin(o);
  return Pt(c * a.x - s * a.y, s * a.x + c * a.y);
Pt proj_vector(Pt a, Pt b, Pt c) { // vector ac proj to
  return (b - a) * ((c - a) * (b - a)) / ((b - a) * (b
      - a));
Pt proj_pt(Pt a, Pt b, Pt c) { // point c proj to ab
  return proj_vector(a, b, c) + a;
```

# 8.2 Heart

```
Pt circenter(Pt p0, Pt p1, Pt p2) { // radius = abs(
   center)
  p1 = p1 - p0, p2 = p2 - p0;
  double x1 = p1.x, y1 = p1.y, x2 = p2.x, y2 = p2.y;
  double m = 2. * (x1 * y2 - y1 * x2);
  Pt center(0, 0);
  center.x = (x1 * x1 * y2 - x2 * x2 * y1 + y1 * y2 * (
      y1 - y2)) / m;
  center.y = (x1 * x2 * (x2 - x1) - y1 * y1 * x2 + x1 *
       y2 * y2) / m;
  return center + p0;
Pt incenter(Pt p1, Pt p2, Pt p3) { // radius = area / s
  double a = abs(p2 - p3), b = abs(p1 - p3), c = abs(p1
       - p2);
  double s = a + b + c;
  return (p1 * a + p2 * b + p3 * c) / s;
Pt masscenter(Pt p1, Pt p2, Pt p3)
{ return (p1 + p2 + p3) / 3; }
Pt orthocenter(Pt p1, Pt p2, Pt p3)
{ return masscenter(p1, p2, p3) * 3 - circenter(p1, p2,
     p3) * 2; }
```

#### 8.3 External Bisector

```
Pt external_bisector(Pt p1, Pt p2, Pt p3) { //213
Pt L1 = p2 - p1, L2 = p3 - p1;
L2 = L2 * abs(L1) / abs(L2);
return L1 + L2;
```

# 8.4 Intersection of Segments

```
Pt LinesInter(Line a, Line b) {
   double abc = (a.b - a.a) ^ (b.a - a.a);
   double abd = (a.b - a.a) ^ (b.b - a.a);
   if (sign(abc - abd) == 0) return b.b;// no inter
   return (b.b * abc - b.a * abd) / (abc - abd);
}
vector<Pt> SegsInter(Line a, Line b) {
   if (btw(a.a, a.b, b.a)) return {b.a};
   if (btw(a.a, a.b, b.b)) return {b.b};
   if (btw(b.a, b.b, a.a)) return {a.a};
   if (btw(b.a, b.b, a.b)) return {a.b};
   if (ori(a.a, a.b, b.a) * ori(a.a, a.b, b.b) == -1 &&
        ori(b.a, b.b, a.a) * ori(b.a, b.b, a.b) == -1)
        return {LinesInter(a, b)};
   return {};
}
```

## 8.5 Intersection of Circle and Line

#### 8.6 Intersection of Circles

# 8.7 Intersection of Polygon and Circle

```
double _area(Pt pa, Pt pb, double r){
  if (abs(pa) < abs(pb)) swap(pa, pb);</pre>
  if (abs(pb) < eps) return 0;</pre>
  double S, h, theta;
 double a = abs(pb), b = abs(pa), c = abs(pb - pa);
double cosB = pb * (pb - pa) / a / c, B = acos(cosB);
  double cosC = (pa * pb) / a / b, C = acos(cosC);
 if (a > r) {
   S = (C / 2) * r * r;
    h = a * b * sin(C) / c;
    if (h < r \&\& B < pi / 2) S -= (acos(h / r) * r * r
         - h * sqrt(r * r - h * h));
  } else if (b > r) {
    theta = pi - B - asin(sin(B) / r * a);
    S = .5 * a * r * sin(theta) + (C - theta) / 2 * r *
  } else S = .5 * sin(C) * a * b;
double area_poly_circle(vector<Pt> poly, Pt 0, double r
    ) {
  double S = 0; int n = poly.size();
  for (int i = 0; i < n; ++i)</pre>
    S += _area(poly[i] - 0, poly[(i + 1) % n] - 0, r) *
          ori(0, poly[i], poly[(i + 1) % n]);
```

```
return fabs(S);
```

# 8.8 Tangent Lines of Circle and Point

# 8.9 Tangent Lines of Circles

```
vector<Line> tangent(Cir a, Cir b) {
#define Pij \
  Pt i = unit(b.o - a.o) * a.r, j = Pt(i.y, -i.x);\
  z.push_back({a.o + i, a.o + i + j});
#define deo(I,J) \
  double d = abs(a.o - b.o), e = a.r I b.r, o = acos(e
      / d);\
  Pt i = unit(b.o - a.o), j = rot(i, o), k = rot(i, -o)
  z.push_back({a.o + j * a.r, b.o J j * b.r});\
z.push_back({a.o + k * a.r, b.o J k * b.r});
  if (a.r < b.r) swap(a, b);</pre>
  vector<Line> z;
  if (abs(a.o - b.o) + b.r < a.r) return z;</pre>
  else if (sign(abs(a.o - b.o) + b.r - a.r) == 0) { Pij
  else {
    deo(-,+); // inter
    // outer
    if (sign(d - a.r - b.r) == 0) { Pij; }
    else if (d > a.r + b.r) { deo(+,-); }
  return z;
```

#### 8.10 Point In Convex

## 8.11 Point Segment Distance

```
double PointSegDist(Pt q0, Pt q1, Pt p) {
  if (sign(abs(q0 - q1)) == 0) return abs(q0 - p);
  if (sign((q1 - q0) * (p - q0)) >= 0 && sign((q0 - q1)
      * (p - q1)) >= 0)
    return fabs(((q1 - q0) ^ (p - q0)) / abs(q0 - q1));
  return min(abs(p - q0), abs(p - q1));
}
```

# 8.12 Convex Hull

```
vector <Pt> ConvexHull(vector <Pt> pt) {
  int n = pt.size();
  sort(all(pt), [&](Pt a, Pt b) {return a.x == b.x ? a.
       y < b.y : a.x < b.x;});</pre>
```

#### 8.13 Convex Hull Distance

```
double ConvexHullDist(vector<Pt> A, vector<Pt> B) {
  for (auto &p : B) p = Pt(0, 0) - p;
  auto C = Minkowski(A, B); // assert SZ(C) > 0
  if (PointInConvex(C, Pt(0, 0))) return 0;
  double ans = PointSegDist(C.back(), C[0], Pt(0, 0));
  for (int i = 0; i + 1 < C.size(); ++i) {
    ans = min(ans, PointSegDist(C[i], C[i + 1], Pt(0, 0)));
  }
  return ans;
}</pre>
```

## 8.14 Minimum Enclosing Circle

```
Cir min_enclosing(vector<Pt> &p) {
 random_shuffle(p.begin(), p.end());
  double r = 0.0:
  Pt cent = p[0];
  for (int i = 1; i < p.size(); ++i) {</pre>
    if (abs2(cent - p[i]) <= r) continue;</pre>
    cent = p[i];
    r = 0.0;
    for (int j = 0; j < i; ++j) {
      if (abs2(cent - p[j]) <= r) continue;</pre>
      cent = (p[i] + p[j]) / 2;
      r = abs2(p[j] - cent);
      for (int k = 0; k < j; ++k) {
        if (abs2(cent - p[k]) <= r) continue;</pre>
        cent = circenter(p[i], p[j], p[k]);
        r = abs2(p[k] - cent);
   }
  return {cent, sqrt(r)};
```

# 8.15 Union of Circles

```
vector<pair<double, double>> CoverSegment(Cir a, Cir b)
  double d = abs(a.o - b.o);
  vector<pair<double, double>> res;
  if (sign(a.r + b.r - d) == 0);
  else if (d <= abs(a.r - b.r) + eps) {
    if (a.r < b.r) res.emplace_back(0, 2 * pi);</pre>
 } else if (d < abs(a.r + b.r) - eps) {</pre>
    double o = acos((a.r * a.r + d * d - b.r * b.r) /
        (2 * a.r * d)), z = atan2((b.o - a.o).y, (b.o -
         a.o).x);
    if (z < 0) z += 2 * pi;
    double 1 = z - o, r = z + o;
    if (1 < 0) 1 += 2 * pi;</pre>
    if (r > 2 * pi) r -= 2 * pi;
    if (1 > r) res.emplace_back(1, 2 * pi), res.
        emplace_back(0, r);
   else res.emplace_back(1, r);
 }
 return res;
double CircleUnionArea(vector<Cir> c) { // circle
    should be identical
  int n = c.size();
  double a = 0, w;
 for (int i = 0; w = 0, i < n; ++i) {</pre>
   vector<pair<double, double>> s = {{2 * pi, 9}}, z;
```

# 8.16 Polar Angle Sort

# 8.17 Rotating Caliper

## 8.18 Rotating SweepLine

```
void RotatingSweepLine(vector <Pt> &pt) {
  int n = pt.size();
  vector <int> id(n), pos(n);
  vector <pair <int, int>> line;
  sort(line.begin(), line.end(), [&](pair <int, int> i,
      pair <int, int> j) {
    Pt a = pt[i.second] - pt[i.first], b = pt[j.second]
         - pt[j.first];
    return (a.pos() == b.pos() ? sign(a ^ b) > 0 : a.
        pos() < b.pos());
  });
  iota(id.begin(), id.end(), 0);
  sort(id.begin(), id.end(), [&](int i, int j) {
    return (sign(pt[i].y - pt[j].y) == 0 ? pt[i].x < pt</pre>
        [j].x : pt[i].y < pt[j].y);
  });
  for (int i = 0; i < n; ++i)</pre>
   pos[id[i]] = i;
  for (auto [i, j] : line) {
   // point sort by the distance to line(i, j)
    // do something
    tie(pos[i], pos[j], id[pos[i]], id[pos[j]]) =
        make_tuple(pos[j], pos[i], j, i);
}
```

#### 8.19 Half Plane Intersection

```
Pt A = a.second - a.first, B = b.second - b.first;
  if (pos(A) == pos(B)) {
    if (sign(A ^ B) == 0) return sign((b.first - a.
        first) * (b.second - a.first)) > 0;
    return sign(A ^ B) > 0;
  return pos(A) < pos(B);</pre>
});
deque <Pt> inter;
deque <pair <Pt, Pt>> seg;
int n = vec.size();
auto get = [&](pair <Pt, Pt> a, pair <Pt, Pt> b) {
    return intersect(a.first, a.second, b.first, b.
    second);};
for (int i = 0; i < n; ++i) if (!i || vec[i] != vec[i</pre>
     - 1]) {
  while (seg.size() >= 2 && sign((vec[i].second -
    inter.back()) ^ (vec[i].first - inter.back()))
      == 1) seg.pop_back(), inter.pop_back();
  while (seg.size() >= 2 && sign((vec[i].second -
      inter.front()) ^ (vec[i].first - inter.front())
      ) == 1) seg.pop_front(), inter.pop_front();
  seg.push_back(vec[i]);
  if (seg.size() >= 2) inter.pb(get(seg[seg.size() -
      2], seg.back()));
while (seg.size() >= 2 && sign((seg.front().second
    inter.back()) ^ (seg.front().first - inter.back()
    )) == 1) seg.pop_back(), inter.pop_back();
inter.push_back(get(seg.front(), seg.back()));
return vector <Pt>(all(inter));
```

## 8.20 Minkowski Sum

```
void reorder(vector <Pt> &P) {
  rotate(P.begin(), min_element(all(P), [&](Pt a, Pt b)
       { return make_pair(a.y, a.x) < make_pair(b.y, b.
      x); }), P.end());
vector <Pt> Minkowski(vector <Pt> P, vector <Pt> Q) {
 // P, Q: convex polygon
 reorder(P), reorder(Q);
 int n = P.size(), m = Q.size();
 P.pb(P[0]), P.pb(P[1]), Q.pb(Q[0]), Q.pb(Q[1]);
  vector <Pt> ans;
 for (int i = 0, j = 0; i < n || j < m; ) {</pre>
   ans.pb(P[i] + Q[j]);
    auto val = (P[i + 1] - P[i]) ^ (Q[j + 1] - Q[j]);
   if (val >= 0) i++;
    if (val <= 0) j++;</pre>
 }
  return ans;
```

## 8.21 Delaunay Triangulation

```
/* Delaunay Triangulation:
Given a sets of points on 2D plane, find a
triangulation such that no points will strictly
inside circumcircle of any triangle.
find : return a triangle contain given point
add_point : add a point into triangulation
A Triangle is in triangulation iff. its has_chd is 0.
Region of triangle u: iterate each u.edge[i].tri,
each points are u.p[(i+1)\%3], u.p[(i+2)\%3]
Voronoi diagram: for each triangle in triangulation,
the bisector of all its edges will split the region.
nearest point will belong to the triangle containing it
const ll inf = MAXC * MAXC * 100; // Lower_bound
    unknown
struct Tri;
struct Edge {
  Tri* tri; int side;
  Edge(): tri(0), side(0){}
  Edge(Tri* _tri, int _side): tri(_tri), side(_side){}
struct Tri {
  pll p[3];
  Edge edge[3];
  Tri* chd[3];
```

```
Tri() {}
  Tri(const pll& p0, const pll& p1, const pll& p2) {
    p[0] = p0; p[1] = p1; p[2] = p2;
     chd[0] = chd[1] = chd[2] = 0;
  bool has_chd() const { return chd[0] != 0; }
  int num_chd() const +
    return !!chd[0] + !!chd[1] + !!chd[2];
  bool contains(pll const& q) const {
     for (int i = 0; i < 3; ++i)</pre>
       if (ori(p[i], p[(i + 1) % 3], q) < 0)
         return 0;
    return 1;
} pool[N * 10], *tris;
void edge(Edge a, Edge b) {
  if(a.tri) a.tri->edge[a.side] = b;
  if(b.tri) b.tri->edge[b.side] = a;
struct Trig { // Triangulation
  Trig() {
    the_root = // Tri should at least contain all
       new(tris++) Tri(pll(-inf, -inf), pll(inf + inf, -
           inf), pll(-inf, inf + inf));
  Tri* find(pll p) { return find(the_root, p); }
  void add_point(const pll &p) { add_point(find(
       the_root, p), p); }
  Tri* the root;
   static Tri* find(Tri* root, const pll &p) {
     while (1) {
       if (!root->has_chd())
         return root;
       for (int i = 0; i < 3 && root->chd[i]; ++i)
         if (root->chd[i]->contains(p)) {
           root = root->chd[i];
           break:
    }
     assert(0); // "point not found"
  void add_point(Tri* root, pll const& p) {
    Tri* t[3];
     /* split it into three triangles */
     for (int i = 0; i < 3; ++i)
       t[i] = new(tris++) Tri(root->p[i], root->p[(i +
    1) % 3], p);
for (int i = 0; i < 3; ++i)
       edge(Edge(t[i], 0), Edge(t[(i + 1) % 3], 1));
     for (int i = 0; i < 3; ++i)
       edge(Edge(t[i], 2), root->edge[(i + 2) % 3]);
     for (int i = 0; i < 3; ++i)
    root->chd[i] = t[i];
for (int i = 0; i < 3; ++i)
       flip(t[i], 2);
  void flip(Tri* tri, int pi) {
    Tri* trj = tri->edge[pi].tri;
     int pj = tri->edge[pi].side;
     if (!trj) return;
     if (!in_cc(tri->p[0], tri->p[1], tri->p[2], trj->p[
         pj])) return;
     /* flip edge between tri,trj *,
     Tri* trk = new(tris++) Tri(tri->p[(pi + 1) % 3],
         trj->p[pj], tri->p[pi]);
     Tri* trl = new(tris++) Tri(trj->p[(pj + 1) % 3],
         tri->p[pi], trj->p[pj]);
     edge(Edge(trk, 0), Edge(trl, 0));
     edge(Edge(trk, 1), tri->edge[(pi + 2) % 3]);
     edge(Edge(trk, 2), trj->edge[(pj + 1) % 3]);
     edge(Edge(trl, 1), trj->edge[(pj + 2) % 3]);
     edge(Edge(trl, 2), tri->edge[(pi + 1) \% 3]);
     tri->chd[0] = trk; tri->chd[1] = trl; tri->chd[2] =
          0:
     trj \rightarrow chd[0] = trk; trj \rightarrow chd[1] = trl; trj \rightarrow chd[2] =
     flip(trk, 1); flip(trk, 2);
     flip(trl, 1); flip(trl, 2);
};
```

```
vector<Tri*> triang; // vector of all triangle
set<Tri*> vst;
void go(Tri* now) { // store all tri into triang
 if (vst.find(now) != vst.end())
    return;
  vst.insert(now);
 if (!now->has_chd())
    return triang.pb(now);
 for (int i = 0; i < now->num_chd(); ++i)
    go(now->chd[i]);
void build(int n, pll* ps) { // build triangulation
 tris = pool; triang.clear(); vst.clear();
 random\_shuffle(ps, ps + n);
 Trig tri; // the triangulation structure
 for (int i = 0; i < n; ++i)</pre>
   tri.add_point(ps[i]);
 go(tri.the_root);
```

# 8.22 Triangulation Vonoroi

```
vector<Line> ls[N]:
pll arr[N];
Line make_line(pdd p, Line 1) {
  pdd d = 1.Y - 1.X; d = perp(d);
  pdd m = (1.X + 1.Y) / 2;
  l = Line(m, m + d);
  if (ori(1.X, 1.Y, p) < 0)
    l = Line(m + d, m);
  return 1;
double calc_area(int id) {
  // use to calculate the area of point "strictly in
       the convex hull"
  vector<Line> hpi = halfPlaneInter(ls[id]);
  vector<pdd> ps;
  for (int i = 0; i < SZ(hpi); ++i)</pre>
    ps.pb(intersect(hpi[i].X, \, hpi[i].Y, \, hpi[(i\, +\, 1)\, \, \%
         SZ(hpi)].X, hpi[(i + 1) % SZ(hpi)].Y));
  double rt = 0;
  for (int i = 0; i < SZ(ps); ++i)</pre>
    rt += cross(ps[i], ps[(i + 1) % SZ(ps)]);
  return fabs(rt) / 2;
void solve(int n, pii *oarr) {
 map<pll, int> mp;
  for (int i = 0; i < n; ++i)</pre>
    arr[i] = pll(oarr[i].X, oarr[i].Y), mp[arr[i]] = i;
  build(n, arr); // Triangulation
  for (auto *t : triang) {
    vector<int> p;
    for (int i = 0; i < 3; ++i)
       if (mp.find(t->p[i]) != mp.end())
        p.pb(mp[t->p[i]]);
    for (int i = 0; i < SZ(p); ++i)</pre>
       for (int j = i + 1; j < SZ(p); ++j) {</pre>
         Line l(oarr[p[i]], oarr[p[j]]);
         ls[p[i]].pb(make_line(oarr[p[i]], 1));
         ls[p[j]].pb(make_line(oarr[p[j]], 1));
  }
}
```

#### 9 Else

#### 9.1 Bit Hack

# 9.2 Dynamic Programming Condition

## 9.2.1 Totally Monotone (Concave/Convex)

```
\begin{array}{l} \forall i < i', j < j' \text{, } B[i][j] \leq B[i'][j] \implies B[i][j'] \leq B[i'][j'] \\ \forall i < i', j < j' \text{, } B[i][j] \geq B[i'][j] \implies B[i][j'] \geq B[i'][j'] \end{array}
```

#### 9.2.2 Monge Condition (Concave/Convex)

```
\begin{array}{l} \forall i < i', j < j' \text{, } B[i][j] + B[i'][j'] \geq B[i][j'] + B[i'][j] \\ \forall i < i', j < j' \text{, } B[i][j] + B[i'][j'] \leq B[i][j'] + B[i'][j] \end{array}
```

#### 9.2.3 Optimal Split Point

```
If B[i][j] + B[i+1][j+1] \geq B[i][j+1] + B[i+1][j] then H_{i,j-1} \leq H_{i,j} \leq H_{i+1,j}
```

## 9.3 Slope Trick

```
template<typename T>
struct slope_trick_convex {
 T minn = 0, ground_1 = 0, ground_r = 0;
  priority_queue<T, vector<T>, less<T>> left;
  priority_queue<T, vector<T>, greater<T>> right;
  slope_trick_convex() {left.push(numeric_limits<T>::
      min() / 2), right.push(numeric_limits<T>::max() /
       2);}
  void push_left(T x) {left.push(x - ground_l);}
  void push_right(T x) {right.push(x - ground_r);}
  //add a line with slope 1 to the right starting from
  void add_right(T x) {
    T l = left.top() + ground_l;
    if (1 <= x) push_right(x);</pre>
    else push_left(x), push_right(l), left.pop(), minn
        += 1 - x;
  //add a line with slope -1 to the left starting from
  void add_left(T x) {
    T r = right.top() + ground_r;
    if (r >= x) push_left(x);
    else push_right(x), push_left(r), right.pop(), minn
  //val[i]=min(val[j]) for all i-l<=j<=i+r
  void expand(T 1, T r) {ground_1 -= 1, ground_r += r;}
  void shift_up(T x) {minn += x;}
  T get_val(T x) {
    T l = left.top() + ground_l, r = right.top() +
        ground_r;
    if (x >= 1 \&\& x <= r) return minn;
    if (x < 1) {
      vector<T> trash;
      T cur_val = minn, slope = 1, res;
      while (1) {
        trash.push_back(left.top());
        left.pop();
        if (left.top() + ground_l <= x) {</pre>
          res = cur_val + slope * (1 - x);
        }
        cur_val += slope * (1 - (left.top() + ground_1)
        1 = left.top() + ground_l;
        slope += 1;
      for (auto i : trash) left.push(i);
      return res;
    if (x > r) {
      vector<T> trash;
      T cur_val = minn, slope = 1, res;
      while (1) {
        trash.push_back(right.top());
        right.pop();
        if (right.top() + ground_r >= x) {
          res = cur_val + slope *(x - r);
          break:
        cur_val += slope * ((right.top() + ground_r) -
            r);
```

#### 9.4 Manhattan MST

```
void solve(int n) {
 init();
  vector<int> v(n), ds;
  for (int i = 0; i < n; ++i) {</pre>
    v[i] = i:
    ds.push_back(x[i] - y[i]);
  sort(ds.begin(), ds.end());
  ds.resize(unique(ds.begin(), ds.end()) - ds.begin());
  sort(v.begin(), v.end(), [&](int i, int j) { return x
      [i] == x[j] ? y[i] > y[j] : x[i] > x[j]; });
  int j = 0;
  for (int i = 0; i < n; ++i) {
    int p = lower_bound(ds.begin(), ds.end(), x[v[i]] -
         y[v[i]]) - ds.begin() + 1;
    pair<int, int> q = query(p);
    // query return prefix minimum
    if (~q.second) add_edge(v[i], q.second);
    add(p, make_pair(x[v[i]] + y[v[i]], v[i]));
void make_graph() {
 solve(n);
  for (int i = 0; i < n; ++i) swap(x[i], y[i]);</pre>
  solve(n);
 for (int i = 0; i < n; ++i) x[i] = -x[i];
  solve(n);
  for (int i = 0; i < n; ++i) swap(x[i], y[i]);</pre>
  solve(n);
```

# 9.5 Dynamic MST

```
int cnt[maxn], cost[maxn], st[maxn], ed[maxn];
pair<int, int> qr[maxn];
// qr[i].first = id of edge to be changed, qr[i].second
     = weight after operation
// cnt[i] = number of operation on edge i
// call solve(0, q - 1, v, 0), where v contains edges i
     such that cnt[i] == 0
void contract(int 1, int r, vector<int> v, vector<int>
    &x, vector<int> &y) {
  sort(v.begin(), v.end(), [&](int i, int j) {
      if (cost[i] == cost[j]) return i < j;</pre>
      return cost[i] < cost[j];</pre>
      });
  djs.save();
 for (int i = 1; i <= r; ++i) djs.merge(st[qr[i].first</pre>
      ], ed[qr[i].first]);
  for (int i = 0; i < (int)v.size(); ++i) {</pre>
    if (djs.find(st[v[i]]) != djs.find(ed[v[i]])) {
      x.push_back(v[i]);
      djs.merge(st[v[i]], ed[v[i]]);
   }
 dis.undo();
  djs.save();
  for (int i = 0; i < (int)x.size(); ++i) djs.merge(st[</pre>
      x[i], ed[x[i]];
  for (int i = 0; i < (int)v.size(); ++i) {</pre>
   if (djs.find(st[v[i]]) != djs.find(ed[v[i]])) {
      y.push_back(v[i]);
      djs.merge(st[v[i]], ed[v[i]]);
   }
 djs.undo();
void solve(int 1, int r, vector<int> v, long long c) {
 if (1 == r) {
    cost[qr[1].first] = qr[1].second;
```

```
if (st[qr[1].first] == ed[qr[1].first]) {
    printf("%lld\n", c);
    return:
  int minv = qr[1].second;
  for (int i = 0; i < (int)v.size(); ++i) minv = min(</pre>
      minv, cost[v[i]]);
  printf("%lld\n", c + minv);
  return:
int m = (1 + r) >> 1;
vector<int> lv = v, rv = v;
vector<int> x, y;
for (int i = m + 1; i <= r; ++i) {</pre>
  cnt[qr[i].first]--
  if (cnt[qr[i].first] == 0) lv.push_back(qr[i].first
contract(l, m, lv, x, y);
long long lc = c, rc = c;
djs.save();
for (int i = 0; i < (int)x.size(); ++i) {</pre>
  lc += cost[x[i]];
  djs.merge(st[x[i]], ed[x[i]]);
solve(1, m, y, 1c);
djs.undo();
x.clear(), y.clear();
for (int i = m + 1; i <= r; ++i) cnt[qr[i].first]++;</pre>
for (int i = 1; i <= m; ++i) {</pre>
  cnt[qr[i].first]--;
  if (cnt[qr[i].first] == 0) rv.push_back(qr[i].first
      ):
}
contract(m + 1, r, rv, x, y);
djs.save();
for (int i = 0; i < (int)x.size(); ++i) {</pre>
  rc += cost[x[i]];
  djs.merge(st[x[i]], ed[x[i]]);
solve(m + 1, r, y, rc);
djs.undo();
for (int i = 1; i <= m; ++i) cnt[qr[i].first]++;</pre>
```

#### 9.6 ALL LCS

```
void all_lcs(string s, string t) { // 0-base
  vector<int> h(t.size());
  iota(all(h), 0);
  for (int a = 0; a < s.size(); ++a) {
    int v = -1;
    for (int c = 0; c < t.size(); ++c)
        if (s[a] == t[c] || h[c] < v)
            swap(h[c], v);
        // LCS(s[0, a], t[b, c]) =
        // c - b + 1 - sum([h[i] >= b] | i <= c)
        // h[i] might become -1 !!
  }
}</pre>
```

# 9.7 Hilbert Curve

```
long long hilbertOrder(int x, int y, int pow, int
    rotate) {
  if (pow == 0) return 0;
  int hpow = 1 << (pow-1);</pre>
  int seg = (x < hpow) ? ((y < hpow) ? 0 : 3) : ((y < hpow)
      hpow) ? 1 : 2);
  seg = (seg + rotate) & 3;
  const int rotateDelta[4] = {3, 0, 0, 1};
  int nx = x & (x ^ hpow), ny = y & (y ^ hpow);
  int nrot = (rotate + rotateDelta[seg]) & 3;
  long long subSquareSize = 111 << (pow * 2 - 2);</pre>
  long long ans = seg * subSquareSize;
  long long add = hilbertOrder(nx, ny, pow - 1, nrot);
  ans += (seg == 1 || seg == 2) ? add : (subSquareSize
      - add - 1);
  return ans;
```

#### 9.8 Pbds

```
#include <ext/pb_ds/priority_queue.hpp>
#include <ext/pb_ds/assoc_container.hpp>
using namespace __gnu_pbds;
#include <ext/rope>
using namespace __gnu_cxx;
 _gnu_pbds::priority_queue <<mark>int</mark>> pq1, pq2;
pq1.join(pq2); // pq1 += pq2, pq2 = {}
cc_hash_table<int, int> m1;
tree<int, null_type, less<int>, rb_tree_tag,
    tree_order_statistics_node_update> oset;
oset.insert(2), oset.insert(4);
*oset.find_by_order(1), oset.order_of_key(1);// 4 0
bitset <100> BS;
BS.flip(3), BS.flip(5);
BS._Find_first(), BS._Find_next(3); // 3 5
rope <int> rp1, rp2;
rp1.push_back(1), rp1.push_back(3);
rp1.insert(0, 2); // pos, num
rp1.erase(0, 2); // pos, len
rp1.substr(0, 2); // pos, Len
rp2.push_back(4);
rp1 += rp2, rp2 = rp1;
rp2[0], rp2[1]; // 3 4
```

#### 9.9 Random

```
struct custom_hash {
    static uint64_t splitmix64(uint64_t x) {
        x += 0x9e3779b97f4a7c15;
        x = (x ^ (x >> 30)) * 0xbf58476d1ce4e5b9;
        x = (x ^ (x >> 27)) * 0x94d049bb133111eb;
        return x ^ (x >> 31);
    }
    size_t operator()(uint64_t a) const {
        static const uint64_t FIXED_RANDOM = chrono::
            steady_clock::now().time_since_epoch().count();
        return splitmix64(i + FIXED_RANDOM);
    }
};
unordered_map <int, int, custom_hash> m1;
random_device rd; mt19937 rng(rd());
```

#### 9.10 Smawk Algorithm

```
11 query(int 1, int r) {
 // ...
struct SMAWK {
 // Condition:
  // If M[1][0] < M[1][1] then M[0][0] < M[0][1]
 // If M[1][0] == M[1][1] then M[0][0] <= M[0][1]
  // For all i, find r_i s.t. M[i][r_i] is maximum ||
      minimum.
  int ans[N], tmp[N];
  void interpolate(vector <int> 1, vector <int> r) {
    int n = 1.size(), m = r.size();
    vector <int> nl;
    for (int i = 1; i < n; i += 2) {</pre>
      nl.push_back(l[i]);
    run(nl, r);
    for (int i = 1, j = 0; i < n; i += 2) {
      while (j < m && r[j] < ans[l[i]])</pre>
      assert(j < m && ans[l[i]] == r[j]);
      tmp[l[i]] = j;
    for (int i = 0; i < n; i += 2) {</pre>
      int curl = 0, curr = m - 1;
      if (i)
        curl = tmp[l[i - 1]];
      if (i + 1 < n)
        curr = tmp[l[i + 1]];
      11 res = query(l[i], r[curl]);
      ans[l[i]] = r[curl];
      for (int j = curl + 1; j <= curr; ++j) {</pre>
        11 nxt = query(l[i], r[j]);
        if (res < nxt)</pre>
          res = nxt, ans[l[i]] = r[j];
      }
   }
 }
```

```
void reduce(vector <int> 1, vector <int> r) {
    int n = 1.size(), m = r.size();
    vector <int> nr;
    for (int j : r) {
      while (!nr.empty()) {
        int i = nr.size() - 1;
        if (query(l[i], nr.back()) <= query(l[i], j))</pre>
          nr.pop_back();
        else
          break:
      if (nr.size() < n)</pre>
        nr.push_back(j);
    }
    run(1, nr);
  void run(vector <int> 1, vector <int> r) {
    int n = 1.size(), m = r.size();
    if (max(n, m) <= 2) {
      for (int i : 1) {
        ans[i] = r[0];
        if (m > 1) {
          if (query(i, r[0]) < query(i, r[1]))</pre>
            ans[i] = r[1];
    } else if (n >= m) {
      interpolate(l, r);
    } else {
      reduce(1, r);
  }
}:
```

#### 9.11 Matroid Intersection

Start from  $S=\emptyset$ . In each iteration, let

```
• Y_1 = \{x \notin S \mid S \cup \{x\} \in I_1\}
• Y_2 = \{x \notin S \mid S \cup \{x\} \in I_2\}
```

If there exists  $x\in Y_1\cap Y_2$ , insert x into S. Otherwise for each  $x\in S, y\not\in S$ , create edges

```
\begin{array}{ll} \bullet & x \rightarrow y \text{ if } S - \{x\} \cup \{y\} \in I_1\text{.} \\ \bullet & y \rightarrow x \text{ if } S - \{x\} \cup \{y\} \in I_2\text{.} \end{array}
```

Find a shortest path (with BFS) starting from a vertex in  $Y_1$  and ending at a vertex in  $Y_2$  which doesn't pass through any other vertices in  $Y_2$ , and alternate the path. The size of S will be incremented by 1 in each iteration. For the weighted case, assign weight w(x) to vertex x if  $x \in S$  and -w(x) if  $x \not\in S$ . Find the path with the minimum number of edges among all minimum length paths and alternate it.

#### 9.12 Python Misc