## **Contents**

1	Basic												1
-	1.1 Shell Script												1
	1.2 Default Code	•	•		•	•	•	•	•	•	•	•	1
	1.3 Increase Stack Size												1
	1.4 Debug Macro												1
	1.5 Stress Test Shell*	•	•			•	•	•	•	•	•	•	1
	1.6 Pragma / FastIO	•	•			•	•	•	•	•	•	•	2
	1.7 Divide*												2
	1.7 Divide	•	•			•	•	•	•	•	•	•	-
2	Data Structure												2
-	2.1 Leftist Tree												2
	2.2 Splay Tree												2
	2.3 Link Cut Tree												2
	2.4 Treap												3
	2.5 2D Segment Tree*	•	•		•	•	•	•	٠	•	•	•	3
	2.6 Zkw*	•	•	• •	•	•	•	•	•	•	•	•	4
	2.7 Chtholly Tree*	·	•				·	·	Ċ	·	·		4
	2.7 chelory free	•	•	• •	•	•	•	•	•	•	•	•	
3	Flow / Matching												5
	3.1 Dinic												5
	3.2 Min Cost Max Flow												5
	3.3 Kuhn Munkres												5
	3.4 SW Min Cut												e
	3.5 Gomory Hu Tree												6
	3.6 Blossom												6
	3.7 Weighted Blossom												7
	3.8 Flow Model												8
4													8
	4.1 Heavy-Light Decomposition						•	•		•			8
	4.2 Centroid Decomposition		•			•	•	•	•	•	•	•	9
	4.3 Edge BCC												9
	4.4 Block Cut Tree												9
	4.5 SCC / 2SAT		•			•	•	•	•	•	•	•	9
	4.6 Virtual Tree		•			•	•	•	•	•	•	•	9
	4.7 Directed MST												16
	4.8 Dominator Tree	•	•			•	٠	٠	٠	٠	٠	. 1	16
5	String											1	Le
,	5.1 Aho-Corasick Automaton												16
	5.2 KMP Algorithm												11
	5.3 Z Algorithm												L 1
	5.4 Manacher												L1
	5.5 Suffix Array												11
	5.6 SAIS											. 1	L1
	5.7 Suffix Automaton											. 1	12
	5.8 Minimum Rotation											. 1	L 2
	5.9 Palindrome Tree											. 1	12
	5.10Main Lorentz											. 1	12
_	Math												
6													
6	6.1 Fraction*											. 1	13
6	<ul><li>6.1 Fraction*</li></ul>											. 1	13 13 13
6	6.1 Fraction*	:	:					•		:		. 1	13 13
6	6.1 Fraction*	:		 						•	:	. 1 . 1 . 1	13 13 13
6	6.1 Fraction*	•	:				:	:		:	:	. 1 . 1 . 1	13 13 13
6	6.1 Fraction*			  			· · ·					. 1 . 1 . 1	13 13 13
6	6.1 Fraction*			  								. 1 . 1 . 1 . 1	13 13 13 14
6	6.1 Fraction*											. 1 . 1 . 1 . 1 . 1 . 1	13 13 14 14
6	6.1 Fraction* 6.2 Miller Rabin / Pollard Rho 6.3 Ext GCD 6.4 PiCount 6.5 Linear Function Mod Min 6.6 Determinant 6.7 Floor Sum 6.8 Quadratic Residue 6.9 Simplex 6.9 Simplex											. 1 . 1 . 1 . 1 . 1 . 1 . 1	13 13 13 14 14 14
6	6.1 Fraction* 6.2 Miller Rabin / Pollard Rho 6.3 Ext GCD 6.4 PiCount 6.5 Linear Function Mod Min 6.6 Determinant 6.7 Floor Sum 6.8 Quadratic Residue 6.9 Simplex 6.10Berlekamp Massey											. 1 . 1 . 1 . 1 . 1 . 1 . 1 . 1	13 13 14 14 14 14
6	6.1 Fraction* 6.2 Miller Rabin / Pollard Rho 6.3 Ext GCD 6.4 PiCount 6.5 Linear Function Mod Min 6.6 Determinant 6.7 Floor Sum 6.8 Quadratic Residue 6.9 Simplex 6.10Berlekamp Massey 6.11Linear Programming Construction											. 1 . 1 . 1 . 1 . 1 . 1 . 1 . 1 . 1 . 1	13 13 13 14 14 14 15
6	6.1 Fraction* 6.2 Miller Rabin / Pollard Rho 6.3 Ext GCD 6.4 PiCount 6.5 Linear Function Mod Min 6.6 Determinant 6.7 Floor Sum 6.8 Quadratic Residue 6.9 Simplex 6.10Berlekamp Massey											. 1 . 1 . 1 . 1 . 1 . 1 . 1 . 1 . 1 . 1	13 13 14 14 14 15
6	6.1 Fraction* 6.2 Miller Rabin / Pollard Rho 6.3 Ext GCD											. 1 . 1 . 1 . 1 . 1 . 1 . 1 . 1 . 1 . 1	13 13 14 14 14 15 15
6	6.1 Fraction* 6.2 Miller Rabin / Pollard Rho											. 1 . 1 . 1 . 1 . 1 . 1 . 1 . 1	13 13 13 14 14
6	6.1 Fraction* 6.2 Miller Rabin / Pollard Rho 6.3 Ext GCD											. 1 . 1 . 1 . 1 . 1 . 1 . 1 . 1 . 1	13 13 14 14 14 15 15 16
6	6.1 Fraction* 6.2 Miller Rabin / Pollard Rho 6.3 Ext GCD 6.4 PiCount 6.5 Linear Function Mod Min 6.6 Determinant 6.7 Floor Sum 6.8 Quadratic Residue 6.9 Simplex 6.10Berlekamp Massey 6.11Linear Programming Construction 6.12Euclidean 6.13Theorem 6.14Estimation 6.15General Purpose Numbers 6.16Tips for Generating Funtion											. 1 1 . 1 . 1 . 1 . 1 . 1 . 1 . 1 . 1 .	13 13 14 14 14 15 15 16
7	6.1 Fraction* 6.2 Miller Rabin / Pollard Rho 6.3 Ext GCD 6.4 PiCount 6.5 Linear Function Mod Min 6.6 Determinant 6.7 Floor Sum 6.8 Quadratic Residue 6.9 Simplex 6.10Berlekamp Massey 6.11Linear Programming Construction 6.13Theorem 6.14Estimation 6.15General Purpose Numbers 6.16Tips for Generating Funtion  Polynomial											. 1 . 1 . 1 . 1 . 1 . 1 . 1 . 1 . 1 . 1	13 13 14 14 14 15 15 16 16
	6.1 Fraction* 6.2 Miller Rabin / Pollard Rho											. 1 . 1 . 1 . 1 . 1 . 1 . 1 . 1 . 1 . 1	13 13 14 14 14 15 16 16
	6.1 Fraction* 6.2 Miller Rabin / Pollard Rho 6.3 Ext GCD											. 1	13 13 14 14 14 15 15 16 16
	6.1 Fraction* 6.2 Miller Rabin / Pollard Rho 6.3 Ext GCD											. 1	13 13 14 14 14 15 15 16 16 16
	6.1 Fraction* 6.2 Miller Rabin / Pollard Rho 6.3 Ext GCD											. 1 1 . 1 . 1 . 1 . 1 . 1 . 1 . 1 . 1 .	13 13 13 13 13 13 14 14 14 14 14 14 14 14 16 16 16 16 16 16 16 16 16 16 16 16 16
	6.1 Fraction* 6.2 Miller Rabin / Pollard Rho 6.3 Ext GCD											. 1 1 . 1 . 1 . 1 . 1 . 1 . 1 . 1 . 1 .	13 13 14 14 14 15 15 16 16 16
7	6.1 Fraction* 6.2 Miller Rabin / Pollard Rho 6.3 Ext GCD 6.4 PiCount 6.5 Linear Function Mod Min 6.6 Determinant 6.7 Floor Sum 6.8 Quadratic Residue 6.9 Simplex 6.10Berlekamp Massey 6.11Linear Programming Construction 6.12Euclidean 6.13Theorem 6.14Estimation 6.15General Purpose Numbers 6.16Tips for Generating Funtion  Polynomial 7.1 Number Theoretic Transform 7.2 Primes 7.3 Polynomial Operations 7.4 Fast Linear Recursion 7.5 Fast Walsh Transform											. 11. 11. 11. 11. 11. 11. 11. 11. 11. 1	1313 1313 1313 1313 1414 1414 1414 1416 1416
	6.1 Fraction* 6.2 Miller Rabin / Pollard Rho 6.3 Ext GCD 6.4 PiCount 6.5 Linear Function Mod Min 6.6 Determinant 6.7 Floor Sum 6.8 Quadratic Residue 6.9 Simplex 6.10Berlekamp Massey 6.11Linear Programming Construction 6.12Euclidean 6.13Theorem 6.14Estimation 6.15General Purpose Numbers 6.16Tips for Generating Funtion  Polynomial 7.1 Number Theoretic Transform 7.2 Primes 7.3 Polynomial Operations 7.4 Fast Linear Recursion 7.5 Fast Walsh Transform  Geometry											. 11. 11. 11. 11. 11. 11. 11. 11. 11. 1	13131313131444444444444444444444444444
7	6.1 Fraction* 6.2 Miller Rabin / Pollard Rho												13131313131314444444444444444444444444
7	6.1 Fraction* 6.2 Miller Rabin / Pollard Rho												13133133133131313131313131313131313131
7	6.1 Fraction* 6.2 Miller Rabin / Pollard Rho 6.3 Ext GCD 6.4 PiCount 6.5 Linear Function Mod Min 6.6 Determinant 6.7 Floor Sum 6.8 Quadratic Residue 6.9 Simplex 6.10Berlekamp Massey 6.11Linear Programming Construction 6.12Euclidean 6.13Theorem 6.14Estimation 6.15General Purpose Numbers 6.16Tips for Generating Funtion Polynomial 7.1 Number Theoretic Transform 7.2 Primes 7.3 Polynomial Operations 7.4 Fast Linear Recursion 7.5 Fast Walsh Transform  Geometry 8.1 Basic 8.2 Heart 8.3 External Bisector 8.4 Intersection of Seements												13131313131314444444444444444444444444
7	6.1 Fraction* 6.2 Miller Rabin / Pollard Rho 6.3 Ext GCD 6.4 PiCount 6.5 Linear Function Mod Min 6.6 Determinant 6.7 Floor Sum 6.8 Quadratic Residue 6.9 Simplex 6.10Berlekamp Massey 6.11Linear Programming Construction 6.12Euclidean 6.13Theorem 6.14Estimation 6.15General Purpose Numbers 6.16Tips for Generating Funtion Polynomial 7.1 Number Theoretic Transform 7.2 Primes 7.3 Polynomial Operations 7.4 Fast Linear Recursion 7.5 Fast Walsh Transform  Geometry 8.1 Basic 8.2 Heart 8.3 External Bisector 8.4 Intersection of Seements												13131331331313131313131313131313131313
7	6.1 Fraction* 6.2 Miller Rabin / Pollard Rho 6.3 Ext GCD 6.4 PiCount 6.5 Linear Function Mod Min 6.6 Determinant 6.7 Floor Sum 6.8 Quadratic Residue 6.9 Simplex 6.10Berlekamp Massey 6.11Linear Programming Construction 6.12Euclidean 6.13Theorem 6.14Estimation 6.15General Purpose Numbers 6.16Tips for Generating Funtion  Polynomial 7.1 Number Theoretic Transform 7.2 Primes 7.3 Polynomial Operations 7.4 Fast Linear Recursion 7.5 Fast Walsh Transform 7.5 Fast Walsh Transform  Geometry 8.1 Basic 8.2 Heart 8.3 External Bisector 8.4 Intersection of Segments 8.5 Intersection of Circle and Line 8.6 Intersection of Circle and Line												131131311313113113113113113113113113113
7	6.1 Fraction* 6.2 Miller Rabin / Pollard Rho 6.3 Ext GCD 6.4 PiCount 6.5 Linear Function Mod Min 6.6 Determinant 6.7 Floor Sum 6.8 Quadratic Residue 6.9 Simplex 6.10Berlekamp Massey 6.11Linear Programming Construction 6.12Euclidean 6.13Theorem 6.14Estimation 6.15General Purpose Numbers 6.16Tips for Generating Funtion  Polynomial 7.1 Number Theoretic Transform 7.2 Primes 7.3 Polynomial Operations 7.4 Fast Linear Recursion 7.5 Fast Walsh Transform 7.5 Fast Walsh Transform  Geometry 8.1 Basic 8.2 Heart 8.3 External Bisector 8.4 Intersection of Segments 8.5 Intersection of Circle and Line 8.6 Intersection of Circle and Line												13131313131131311313113131131313131313
7	6.1 Fraction* 6.2 Miller Rabin / Pollard Rho 6.3 Ext GCD 6.4 PiCount 6.5 Linear Function Mod Min 6.6 Determinant 6.7 Floor Sum 6.8 Quadratic Residue 6.9 Simplex 6.10Berlekamp Massey 6.11Linear Programming Construction 6.12Euclidean 6.13Theorem 6.14Estimation 6.15General Purpose Numbers 6.16Tips for Generating Funtion Polynomial 7.1 Number Theoretic Transform 7.2 Primes 7.3 Polynomial Operations 7.4 Fast Linear Recursion 7.5 Fast Walsh Transform 6eometry 8.1 Basic 8.2 Heart 8.3 External Bisector 8.4 Intersection of Segments 8.5 Intersection of Circle and Line												13131313131313131313131313131313131313
7	6.1 Fraction* 6.2 Miller Rabin / Pollard Rho 6.3 Ext GCD 6.4 PiCount 6.5 Linear Function Mod Min 6.6 Determinant 6.7 Floor Sum 6.8 Quadratic Residue 6.9 Simplex 6.10Berlekamp Massey 6.11Linear Programming Construction 6.12Euclidean 6.13Theorem 6.14Estimation 6.15General Purpose Numbers 6.16Tips for Generating Funtion  Polynomial 7.1 Number Theoretic Transform 7.2 Primes 7.3 Polynomial Operations 7.4 Fast Linear Recursion 7.5 Fast Walsh Transform 7.5 Fast Walsh Transform  Geometry 8.1 Basic 8.2 Heart 8.3 External Bisector 8.4 Intersection of Segments 8.5 Intersection of Circle and Line 8.6 Intersection of Polygon and Circle 8.7 Intersection of Polygon and Circle												131131131131131131131131131131131131131
7	6.1 Fraction* 6.2 Miller Rabin / Pollard Rho 6.3 Ext GCD 6.4 PiCount 6.5 Linear Function Mod Min 6.6 Determinant 6.7 Floor Sum 6.8 Quadratic Residue 6.9 Simplex 6.9 Simplex 6.10Berlekamp Massey 6.11Linear Programming Construction 6.12Euclidean 6.13Theorem 6.14Estimation 6.15General Purpose Numbers 6.16Tips for Generating Funtion 6.16Tips for Generating Funtion 6.18 Polynomial 7.1 Number Theoretic Transform 7.2 Primes 7.3 Polynomial Operations 7.4 Fast Linear Recursion 7.5 Fast Walsh Transform 6.0 Heart 8.1 Basic 8.2 Heart 8.3 External Bisector 8.4 Intersection of Segments 8.5 Intersection of Circle and Line 8.6 Intersection of Polygon and Circ 8.7 Intersection of Circles 8.7 Intersection of Circle and Poir 8.9 Tangent Lines of Circles 8.10Point In Convex 8.10Point In Convex												13131313131313131313131313131313131313
7	6.1 Fraction* 6.2 Miller Rabin / Pollard Rho 6.3 Ext GCD 6.4 PiCount 6.5 Linear Function Mod Min 6.6 Determinant 6.7 Floor Sum 6.8 Quadratic Residue 6.9 Simplex 6.10Berlekamp Massey 6.11Linear Programming Construction 6.12Euclidean 6.13Theorem 6.14Estimation 6.15General Purpose Numbers 6.16Tips for Generating Funtion Polynomial 7.1 Number Theoretic Transform 7.2 Primes 7.3 Polynomial Operations 7.4 Fast Linear Recursion 7.5 Fast Walsh Transform 7.5 Fast Walsh Transform 6.2 Heart 8.3 External Bisector 8.4 Intersection of Segments 8.5 Intersection of Circle and Line 8.6 Intersection of Circle and Poir 8.7 Tangent Lines of Circle and Poir 8.8 Tangent Lines of Circle and Poir 8.9 Tangent Lines of Circles 9.7 Tangent Lines of Circles 9.7 Tangent Lines of Circle and Poir 8.9 Tangent Lines of Circles 9.7 Tangent Lines of Circle and Poir 8.9 Tangent Lines of Circles 9.7 Tangent Lines of Circles 9.7 Tangent Lines of Circle and Poir 9.8 Tangent Lines of Circle and Poir 9.9 Tangent Lines of Circles 9.7 Tangent Lines of Circles 9.7 Tangent Lines of Circle and Poir 9.8 Tangent Lines of Circle and Poir 9.9 Tangent Lines of Circle and Poir 9.9 Tangent Lines of Circles 9.7 Tangent Lines of Circle and Poir 9.9 Tangent Lines of Circle and Poir												13131313131313131313131313131313131313
7	6.1 Fraction* 6.2 Miller Rabin / Pollard Rho 6.3 Ext GCD 6.4 PiCount 6.5 Linear Function Mod Min 6.6 Determinant 6.7 Floor Sum 6.8 Quadratic Residue 6.9 Simplex 6.10Berlekamp Massey 6.11Linear Programming Construction 6.12Euclidean 6.13Theorem 6.14Estimation 6.15General Purpose Numbers 6.16Tips for Generating Funtion  Polynomial 7.1 Number Theoretic Transform 7.2 Primes 7.3 Polynomial Operations 7.4 Fast Linear Recursion 7.5 Fast Walsh Transform 7.6 Fast Walsh Transform 8.1 Basic 8.2 Heart 8.3 External Bisector 8.4 Intersection of Segments 8.5 Intersection of Circle and Line 8.6 Intersection of Circle and Line 8.6 Intersection of Polygon and Circle 8.7 Intersection of Polygon and Circle 8.8 Tangent Lines of Circle and Poir 8.9 Tangent Lines of Circles 8.10Point In Convex 8.11Point Segment Distance 8.12Convex Hull 8.12Convex Hull 8.2 External Sistance 8.12Convex Hull												131313131311311131113111311131113113113
7	6.1 Fraction* 6.2 Miller Rabin / Pollard Rho												13131313131313131313131313131313131313
7	6.1 Fraction* 6.2 Miller Rabin / Pollard Rho 6.3 Ext GCD 6.4 PiCount 6.5 Linear Function Mod Min 6.6 Determinant 6.7 Floor Sum 6.8 Quadratic Residue 6.9 Simplex 6.9 Simplex 6.10Berlekamp Massey 6.11Linear Programming Construction 6.12Euclidean 6.13Theorem 6.14Estimation 6.15General Purpose Numbers 6.16Tips for Generating Funtion 6.16Tips for Generating Funtion 6.19 Polynomial 7.1 Number Theoretic Transform 7.2 Primes 7.3 Polynomial Operations 7.4 Fast Linear Recursion 7.5 Fast Walsh Transform 6.6 Hassic 8.2 Heart 8.3 External Bisector 8.4 Intersection of Segments 8.5 Intersection of Circle and Line 8.6 Intersection of Circle and Line 8.6 Intersection of Circle and Poir 8.9 Tangent Lines of Circles 8.10Point In Convex 8.11Point Segment Distance 8.14Minimum Enclosing Circle 8.14Minimum Enclosing Circle 8.14Minimum Enclosing Circle 8.14Minimum Enclosing Circle 6.1.1. Convex 8.14Minimum Enclosing Circle												13131313131313131313131313131313131313
7	6.1 Fraction* 6.2 Miller Rabin / Pollard Rho 6.3 Ext GCD 6.4 PiCount 6.5 Linear Function Mod Min 6.6 Determinant 6.7 Floor Sum 6.8 Quadratic Residue 6.9 Simplex 6.9 Simplex 6.10Berlekamp Massey 6.11Linear Programming Construction 6.12Euclidean 6.13Theorem 6.14Estimation 6.15General Purpose Numbers 6.16Tips for Generating Funtion 6.16Tips for Generating Funtion 6.19 Polynomial 7.1 Number Theoretic Transform 7.2 Primes 7.3 Polynomial Operations 7.4 Fast Linear Recursion 7.5 Fast Walsh Transform 6.6 Hassic 8.2 Heart 8.3 External Bisector 8.4 Intersection of Segments 8.5 Intersection of Circle and Line 8.6 Intersection of Circle and Line 8.6 Intersection of Circle and Poir 8.9 Tangent Lines of Circles 8.10Point In Convex 8.11Point Segment Distance 8.14Minimum Enclosing Circle 8.14Minimum Enclosing Circle 8.14Minimum Enclosing Circle 8.14Minimum Enclosing Circle 6.1.1. Convex 8.14Minimum Enclosing Circle												133 133 133 133 133 133 133 133 133 133
7	6.1 Fraction* 6.2 Miller Rabin / Pollard Rho												133 133 133 133 133 133 134 134 134 134

8.18Rotating SweepLine . . . . . .

```
8.19Half Plane Intersection .
                                        20
 8.20Minkowski Sum . .
                                        21
 8.22Triangulation Vonoroi . . . . . . . . . . . .
9 Else
                                        22
 9.1 Bit Hack
                                        22
 22
                                        22
                                        22
    9.2.3 Optimal Split Point . . . . . . . . . . .
                                        22
 22
 9.5 Dynamic MST . . . . . . . . . . . . . . . . . .
                                        23
 9.6 ALL LCS . . .
                                        23
 9.7 Hilbert Curve . . . . . . . . . . . . . . . .
                                        23
 24
                                        24
 9.11Matroid Intersection . . . . . . . . . . . . . . . . . .
                                        24
 9.12Python Misc . . . . .
```

## 1 Basic

## 1.1 Shell Script

```
g++ -std=c++17 -DABS -Wall -Wextra -Wshadow $1.cpp -o
    $1 && ./$1
for i in {A..J}; do cp tem.cpp $i.cpp; done;
```

#### 1.2 Default Code

```
#include <bits/stdc++.h>
using namespace std;
typedef long long l1;
#define pb push_back
#define pii pair<int, int>
#define all(a) a.begin(), a.end()
#define sz(a) ((int)a.size())
```

#### 1.3 Increase Stack Size

```
const int size = 256 << 20;
register long rsp asm("rsp");
char *p = (char*)malloc(size) + size, *bak = (char*)rsp
;
__asm__("movq %0, %%rsp\n"::"r"(p));
// main
__asm__("movq %0, %%rsp\n"::"r"(bak));</pre>
```

## 1.4 Debug Macro

## 1.5 Stress Test Shell\*

```
#!/usr/bin/env bash
   g++ $1.cpp -o $1
   g++ $2.cpp -o $2
   g++ $3.cpp -o $3
    for i in {1..100}; do
     ./$3 > input.txt
     # st=$(date +%s%N)
     ./$1 < input.txt > output1.txt
     # echo "$((($(date +%s%N) - $st)/1000000))ms"
      ./$2 < input.txt > output2.txt
     if cmp --silent -- "output1.txt" "output2.txt"; then
       continue
     fi
     echo Input:
     cat input.txt
     echo Your Output:
     cat output1.txt
     echo Correct Output:
     cat output2.txt
     exit 1
   done
   echo OK!
20 |./stress.sh main good gen
```

#### 1.6 Pragma / FastIO

```
#pragma GCC optimize("Ofast,inline,unroll-loops")
#pragma GCC target("bmi,bmi2,lzcnt,popcnt,avx2")
#include<unistd.h>
char OB[65536]; int OP;
inline char RC() {
  static char buf[65536], *p = buf, *q = buf;
  return p == q && (q = (p = buf) + read(0, buf, 65536)
) == buf ? -1 : *p++;
inline int R() {
  static char c;
  while((c = RC()) < '0'); int a = c ^ '0';</pre>
  while((c = RC()) >= '0') a *= 10, a += c ^ '0';
inline void W(int n) {
  static char buf[12], p;
  if (n == 0) OB[OP++]='0'; p = 0;
 while (n) buf[p++] = 0' + (n \% 10), n /= 10;
  for (--p; p >= 0; --p) OB[OP++] = buf[p];
  if (OP > 65520) write(1, OB, OP), OP = 0;
```

## 1.7 Divide\*

```
ll divdown(ll a, ll b) {
  return a / b - (a < 0 && a % b);
}
ll divup(ll a, ll b) {
  return a / b + (a > 0 && a % b);
}
a / b < x -> divdown(a, b) + 1 <= x
a / b <= x -> divup(a, b) <= x
x < a / b -> x <= divup(a, b) - 1
x <= a / b -> x <= divdown(a, b)</pre>
```

# 2 Data Structure

# 2.1 Leftist Tree

```
struct node {
        11 rk, data, sz, sum;
node *1, *r;
          node(11 k) : rk(0), data(k), sz(1), 1(0), r(0), sum(k)
                             ) {}
11 sz(node *p) { return p ? p->sz : 0; }
11 rk(node *p) { return p ? p->rk : -1; }
11 sum(node *p) { return p ? p->sum : 0; }
node *merge(node *a, node *b) {
        if (!a || !b) return a ? a : b;
        if (a->data < b->data) swap(a, b);
        a \rightarrow r = merge(a \rightarrow r, b)
        if (rk(a->r) > rk(a->l)) swap(a->r, a->l);
        a - rk = rk(a - r) + 1, a - rk = sz(a - r) + sz(a 
                          1;
        a \rightarrow sum = sum(a \rightarrow 1) + sum(a \rightarrow r) + a \rightarrow data;
        return a;
void pop(node *&o) {
        node *tmp = o;
         o = merge(o->1, o->r);
         delete tmp;
```

#### 2.2 Splay Tree

```
| struct Splay {
   int pa[N], ch[N][2], sz[N], rt, _id;
   ll v[N];
   Splay() {}
   void init() {
     rt = 0, pa[0] = ch[0][0] = ch[0][1] = -1;
     sz[0] = 1, v[0] = inf;
   }
   int newnode(int p, int x) {
     int id = _id++;
     v[id] = x, pa[id] = p;
```

```
ch[id][0] = ch[id][1] = -1, sz[id] = 1;
    return id;
  void rotate(int i) {
    int p = pa[i], x = ch[p][1] == i, gp = pa[p], c =
        ch[i][!x];
    sz[p] -= sz[i], sz[i] += sz[p];
    if (~c) sz[p] += sz[c], pa[c] = p;
    ch[p][x] = c, pa[p] = i;
    pa[i] = gp, ch[i][!x] = p;
    if (~gp) ch[gp][ch[gp][1] == p] = i;
  void splay(int i) {
    while (~pa[i]) {
      int p = pa[i];
      if (~pa[p]) rotate(ch[pa[p]][1] == p ^ ch[p][1]
           == i ? i : p);
      rotate(i);
    rt = i;
  int lower_bound(int x) {
    int i = rt, last = -1;
    while (true) {
      if (v[i] == x) return splay(i), i;
      if (v[i] > x) {
        last = i;
        if (ch[i][0] == -1) break;
        i = ch[i][0];
      }
      else {
        if (ch[i][1] == -1) break;
        i = ch[i][1];
      }
    splay(i);
    return last; // -1 if not found
  void insert(int x) {
    int i = lower_bound(x);
    if (i == -1) {
      // assert(ch[rt][1] == -1);
      int id = newnode(rt, x);
      ch[rt][1] = id, ++sz[rt];
      splay(id);
    else if (v[i] != x) {
      splay(i);
      int id = newnode(rt, x), c = ch[rt][0];
      ch[rt][0] = id;
      ch[id][0] = c;
      if (~c) pa[c] = id, sz[id] += sz[c];
      ++sz[rt];
      splay(id);
  }
};
```

# 2.3 Link Cut Tree

```
// weighted subtree size, weighted path max
struct LCT {
  int ch[N][2], pa[N], v[N], sz[N], sz2[N], w[N], mx[N
      ], _id;
  // sz := sum of v in splay, sz2 := sum of v in
      virtual subtree
  // mx := max w in splay
  bool rev[N];
  LCT() : _id(1) {}
  int newnode(int _v, int _w) {
    int x = _id++;
ch[x][0] = ch[x][1] = pa[x] = 0;
    v[x] = sz[x] = _v;
    sz2[x] = 0;
    w[x] = mx[x] = w;
    rev[x] = false;
    return x;
  void pull(int i) {
    sz[i] = v[i] + sz2[i];
    mx[i] = w[i];
    if (ch[i][0])
```

node \*1, \*r;

```
sz[i] += sz[ch[i][0]], mx[i] = max(mx[i], mx[ch[i])
                                                                node(int k) : data(k), sz(1), l(0), r(0) {}
                                                                void up() {
    if (ch[i][1])
                                                                  sz = 1;
                                                                  if (1) sz += 1->sz;
      sz[i] += sz[ch[i][1]], mx[i] = max(mx[i], mx[ch[i])
           ][1]]);
                                                                  if (r) sz += r \rightarrow sz;
  void push(int i) {
                                                                void down() {}
    if (rev[i]) reverse(ch[i][0]), reverse(ch[i][1]),
                                                              };
        rev[i] = false;
                                                              int sz(node *a) { return a ? a->sz : 0; }
                                                              node *merge(node *a, node *b) {
                                                                if (!a || !b) return a ? a : b;
  void reverse(int i) {
                                                                if (rand() % (sz(a) + sz(b)) < sz(a))</pre>
    if (!i) return;
    swap(ch[i][0], ch[i][1]);
                                                                  return a \rightarrow down(), a \rightarrow r = merge(a \rightarrow r, b), a \rightarrow up(), a
    rev[i] ^= true;
                                                                return b->down(), b->l = merge(a, b->l), b->up(), b;
  bool isrt(int i) {// rt of splay
    if (!pa[i]) return true;
                                                              void split(node *o, node *&a, node *&b, int k) {
    return ch[pa[i]][0] != i && ch[pa[i]][1] != i;
                                                                if (!o) return a = b = 0, void();
                                                                o->down():
                                                                if (o->data <= k)
  void rotate(int i) {
    int p = pa[i], x = ch[p][1] == i, c = ch[i][!x], gp
                                                                  a = o, split(o \rightarrow r, a \rightarrow r, b, k), <math>a \rightarrow up();
                                                                else b = o, split(o->1, a, b->1, k), b->up();
          = pa[p];
    if (ch[gp][0] == p) ch[gp][0] = i;
    else if (ch[gp][1] == p) ch[gp][1] = i;
                                                              void split2(node *o, node *&a, node *&b, int k) {
    pa[i] = gp, ch[i][!x] = p, pa[p] = i;
                                                                if (sz(o) \le k) return a = o, b = 0, void();
    ch[p][x] = c, pa[c] = p;
                                                                o->down();
    pull(p), pull(i);
                                                                if (sz(o->1) + 1 <= k)
                                                                 a = o, split2(o->r, a->r, b, k - <math>sz(o->l) - 1);
  void splay(int i) {
                                                                else b = o, split2(o->1, a, b->1, k);
    vector<int> anc;
                                                                o->up();
    anc.push_back(i);
    while (!isrt(anc.back())) anc.push_back(pa[anc.back
                                                             node *kth(node *o, int k) {
                                                                if (k <= sz(o->1)) return kth(o->1, k);
         ()1);
                                                                if (k == sz(o->1) + 1) return o;
    while (!anc.empty()) push(anc.back()), anc.pop_back
                                                                return kth(o\rightarrow r, k - sz(o\rightarrow l) - 1);
        ();
    while (!isrt(i)) {
      int p = pa[i];
                                                              int Rank(node *o, int key) {
       if (!isrt(p)) rotate(ch[p][1] == i ^ ch[pa[p]][1]
                                                                if (!o) return 0;
                                                                if (o->data < key)</pre>
           == p ? i : p);
      rotate(i);
                                                                  return sz(o->1) + 1 + Rank(o->r, key);
    }
                                                                else return Rank(o->1, key);
  void access(int i) {
                                                              bool erase(node *&o, int k) {
    int last = 0;
                                                                if (!o) return 0;
    while (i) {
                                                                if (o->data == k) {
                                                                  node *t = o;
      splay(i);
      if (ch[i][1])
                                                                  o->down(), o = merge(o->1, o->r);
        sz2[i] += sz[ch[i][1]];
                                                                  delete t;
      sz2[i] -= sz[last];
                                                                  return 1;
      ch[i][1] = last;
      pull(i), last = i, i = pa[i];
                                                                node *&t = k < o->data ? o->l : o->r;
    }
                                                                return erase(t, k) ? o->up(), 1 : 0;
  void makert(int i) {
                                                              void insert(node *&o, int k) {
                                                                node *a, *b;
    access(i), splay(i), reverse(i);
                                                                o->down(), split(o, a, b, k),
  void link(int i, int j) {
                                                               o = merge(a, merge(new node(k), b));
    // assert(findrt(i) != findrt(j));
                                                                o->up();
    makert(i);
    makert(j);
                                                             void interval(node *&o, int 1, int r) {
                                                                node *a, *b, *c; // [l, r)
    pa[i] = j;
    sz2[j] += sz[i];
                                                                o->down();
                                                                split2(o, a, b, 1), split2(b, b, c, r - 1);
    pull(j);
                                                                // operate
  void cut(int i, int j) {
                                                                o = merge(a, merge(b, c)), o->up();
                                                             }
    makert(i), access(j), splay(i);
    // assert(sz[i] == 2 && ch[i][1] == j);
    ch[i][1] = pa[j] = 0, pull(i);
                                                              2.5 2D Segment Tree*
  int findrt(int i) {
                                                             // 2D range add, range sum in Log^2
    access(i), splay(i);
                                                             struct seg {
    while (ch[i][0]) push(i), i = ch[i][0];
                                                                int 1, r;
    splay(i);
                                                                11 sum, 1z;
                                                                seg *ch[2]{};
    return i;
                                                                seg(int _1, int _r) : l(_1), r(_r), sum(0), lz(0) {}
};
                                                                void push()
                                                                  if (lz) ch[0]->add(l, r, lz), ch[1]->modify(l, r,
2.4
      Treap
                                                                      1z), 1z = 0;
struct node {
                                                                void pull() \{sum = ch[0] -> sum + ch[1] -> sum;\}
  int data, sz;
                                                                void add(int _l, int _r, ll d) {
```

if (\_1 <= 1 && r <= \_r) {</pre>

```
sum += d * (r - 1);
      1z += d;
      return;
    if (!ch[0]) ch[0] = new seg(1, 1 + r >> 1), ch[1] =
          new seg(1 + r \gg 1, r);
    push();
    if (_l < l + r >> 1) ch[0]->add(_l, _r, d);
    if (1 + r >> 1 < _r) ch[1]->add(_1, _r, d);
    pull();
  11 qsum(int _1, int _r) {
    if (_1 <= 1 && r <= _r) return sum;</pre>
    if (!ch[0]) return lz * (min(r, _r) - max(l, _l));
    push();
    if (_1 < 1 + r >> 1) res += ch[0]->qsum(_1, _r);
    if (1 + r >> 1 < _r) res += ch[1]->qsum(_1, _r);
    return res:
  }
struct seg2 {
  int 1, r;
  seg v, lz;
  seg2 *ch[2]{};
  seg2(int _1, int _r) : 1(_1), r(_r), v(0, N), lz(0, N
    if (1 < r - 1) ch[0] = new seg2(1, 1 + r >> 1), ch
         [1] = new seg2(1 + r >> 1, r);
  void add(int _1, int _r, int _12, int _r2, l1 d) {
  v.add(_12, _r2, d * (min(r, _r) - max(1, _1)));
    if (_1 <= 1 && r <= _r) {</pre>
      lz.add(_12, _r2, d);
    if (_1 < 1 + r >> 1) ch[0]->add(_1, _r, _12, _r2, d)
    if (l + r >> 1 < _r) ch[1]->add(_l, _r, _l2, _r2, d
  11 qsum(int _1, int _r, int _12, int _r2) {
    11 res = v.qsum(_12, _r2);
    if (_1 <= 1 && r <= _r) return res;</pre>
    res += lz.qsum(_12, _r2) * (min(r, _r) - max(1, _1)
    if (_1 < 1 + r >> 1) res += ch[0]->query(_1, _r,
         _12, _r2);
    if (1 + r >> 1 < _r) res += ch[1]->query(_1, _r,
         _12, _r2);
     return res;
  }
};
```

#### 2.6 Zkw\*

```
11\ \mathsf{mx}[\mathsf{N}\ <<\ 1]\ ,\ \mathsf{sum}[\mathsf{N}\ <<\ 1]\ ,\ 1\mathsf{z}[\mathsf{N}\ <<\ 1]\ ;
void add(int 1, int r, 11 d) { // [l, r), 0-based
  int len = 1, cntl = 0, cntr = 0;
  for (1 += N, r += N + 1; l ^ r ^ 1; l >>= 1, r >>= 1,
       len <<= 1) {
    sum[1] += cnt1 * d, sum[r] += cnt[r] * d;
    if (len > 1) {
      mx[1] = max(mx[1 << 1], mx[1 << 1 | 1]) + lz[1];
      mx[r] = max(mx[r << 1], mx[r << 1 | 1]) + lz[r];
    if (~1 & 1)
      sum[1 ^ 1] += d * len, mx[1 ^ 1] += d, lz[1 ^ 1]
           += d, cntl += len;
    if (r & 1)
      sum[r ^ 1] += d * len, mx[r ^ 1] += d, lz[r ^ 1]
          += d, cntr += len;
  sum[1] += cnt1 * d, sum[r] += cntr * d;
  if (len > 1) {
    mx[1] = max(mx[1 << 1], mx[1 << 1 | 1]) + lz[1];
    mx[r] = max(mx[r << 1], mx[r << 1 | 1]) + lz[r];
  cntl += cntr:
  for (1 >>= 1; 1; 1 >>= 1) {
    sum[1] += cnt1 * d;
    mx[1] = max(mx[1 << 1], mx[1 << 1 | 1]) + lz[1];
```

```
}
11 qsum(int 1, int r) {
  ll res = 0, len = 1, cntl = 0, cntr = 0;
  for (1 += N, r += N + 1; 1 ^ r ^ 1; 1 >>= 1, r >>= 1,
       len <<= 1) {
    res += cntl * lz[1] + cntr * lz[r];
    if (~l & 1) res += sum[l ^ 1], cntl += len;
    if (r & 1) res += sum[r ^ 1], cntr += len;
  res += cntl * lz[1] + cntr * lz[r];
  cntl += cntr;
  for (1 >>= 1; 1; 1 >>= 1) res += cntl * lz[1];
  return res;
11 qmax(int 1, int r) {
  11 maxl = -INF, maxr = -INF;
for (1 += N, r += N + 1; 1 ^ r ^ 1; 1 >>= 1, r >>= 1)
    \max l += lz[1], \max[r] += lz[r];
    if (~1 & 1) maxl = max(maxl, mx[l ^ 1]);
    if (r & 1) maxr = max(maxr, mx[r ^ 1]);
  maxl = max(maxl + lz[1], maxr + lz[r]);
  for (1 >>= 1; 1; 1 >>= 1) max1 += lz[1];
  return max1;
```

## 2.7 Chtholly Tree\*

```
struct ChthollyTree {
  struct interval {
    int 1, r;
    11 v;
    interval (int _1, int _r, ll _v) : l(_l), r(_r), v(
  struct cmp {
    bool operator () (const interval &a, const interval
        & b) const {
      return a.l < b.l;</pre>
    }
  };
  set <interval, cmp> s;
  vector <interval> split(int 1, int r) {
    // split into [0, l), [l, r), [r, n) and return [l, r]
    vector <interval> del, ans, re;
    auto it = s.lower_bound(interval(1, -1, 0));
    if (it != s.begin() && (it == s.end() || 1 < it->1)
        ) {
      --it;
      del.pb(*it);
      if (r < it->r) {
        re.pb(interval(it->1, 1, it->v));
        ans.pb(interval(1, r, it->v));
        re.pb(interval(r, it->r, it->v));
      } else ·
        re.pb(interval(it->1, 1, it->v));
        ans.pb(interval(l, it->r, it->v));
      ++it;
    for (; it != s.end() && it->r <= r; ++it) {</pre>
      ans.pb(*it);
      del.pb(*it);
    if (it != s.end() && it->l < r) {</pre>
      del.pb(*it);
      ans.pb(interval(it->l, r, it->v));
      re.pb(interval(r, it->r, it->v));
    for (interval &i : del)
      s.erase(i);
    for (interval &i : re)
     s.insert(i);
    return ans;
  void merge(vector <interval> a) {
    for (interval &i : a)
      s.insert(i);
```

# 3 Flow / Matching

## 3.1 Dinic

**}**;

```
struct Dinic { // 0-base
  struct edge {
    int to, cap, flow, rev;
  vector<edge> adj[N];
  int s, t, dis[N], cur[N], n;
  int dfs(int u, int cap) {
    if (u == t || !cap) return cap;
    for (int &i = cur[u]; i < (int)adj[u].size(); ++i)</pre>
       edge &e = adj[u][i];
      if (dis[e.to] == dis[u] + 1 && e.flow != e.cap) {
        int df = dfs(e.to, min(e.cap - e.flow, cap));
        if (df) {
          e.flow += df;
           adj[e.to][e.rev].flow -= df;
          return df;
      }
    dis[u] = -1;
    return 0;
  bool bfs() {
    fill_n(dis, n, -1);
    queue<int> q;
    q.push(s), dis[s] = 0;
    while (!q.empty()) {
      int tmp = q.front();
      q.pop();
      for (auto &u : adj[tmp])
        if (!~dis[u.to] && u.flow != u.cap) {
           q.push(u.to);
           dis[u.to] = dis[tmp] + 1;
    return dis[t] != -1;
  int maxflow(int _s, int _t) {
    s = _s, t = _t;
    int flow = 0, df;
    while (bfs()) {
      fill_n(cur, n, 0);
      while ((df = dfs(s, INF))) flow += df;
    return flow;
  void init(int _n) {
    n = _n;
for (int i = 0; i < n; ++i) adj[i].clear();</pre>
  void reset() {
    for (int i = 0; i < n; ++i)</pre>
      for (auto &j : adj[i]) j.flow = 0;
  void add_edge(int u, int v, int cap) {
    adj[u].pb(edge{v, cap, 0, (int)adj[v].size()});
    adj[v].pb(edge{u, 0, 0, (int)adj[u].size() - 1});
};
```

#### 3.2 Min Cost Max Flow

```
template <typename T1, typename T2>
struct MCMF { // T1 -> flow, T2 -> cost, 0-based
  const T1 INF1 = 1 << 30;
  const T2 INF2 = 1 << 30;
  struct edge {
    int v; T1 f; T2 c;
  } E[M << 1];
  vector <int> adj[N];
  T2 dis[N], pot[N];
  int rt[N], vis[N], n, m, s, t;
  bool SPFA() {
    fill_n(rt, n, -1), fill_n(dis, n, INF2);
```

```
fill_n(vis, n, false);
     queue <int> q;
     q.push(s), dis[s] = 0, vis[s] = true;
     while (!q.empty()) {
       int v = q.front(); q.pop();
       vis[v] = false;
       for (int id : adj[v]) if (E[id].f > 0 && dis[E[id
           ].v] > dis[v] + E[id].c + pot[v] - pot[E[id].
           v1) {
           dis[E[id].v] = dis[v] + E[id].c + pot[v] -
                pot[E[id].v], rt[E[id].v] = id;
           if (!vis[E[id].v]) vis[E[id].v] = true, q.
               push(E[id].v);
         }
    }
     return dis[t] != INF2;
  bool dijkstra() {
     fill_n(rt, n, -1), fill_n(dis, n, INF2);
    priority_queue <pair <T2, int>, vector <pair <T2,
    int>>, greater <pair <T2, int>>> pq;
     dis[s] = 0, pq.emplace(dis[s], s);
     while (!pq.empty()) {
       auto [d, v] = pq.top(); pq.pop();
       if (dis[v] < d) continue;</pre>
       for (int id : adj[v]) if (E[id].f > 0 && dis[E[id
           ].v] > dis[v] + E[id].c + pot[v] - pot[E[id].
           v1) {
           dis[E[id].v] = dis[v] + E[id].c + pot[v] -
               pot[E[id].v], rt[E[id].v] = id;
           pq.emplace(dis[E[id].v], E[id].v);
     return dis[t] != INF2;
  pair <T1, T2> solve(int _s, int _t) {
     s = _s, t = _t, fill_n(pot, n, 0);
     T1 flow = 0; T2 cost = 0;
     bool fr = true;
     while ((fr ? SPFA() : dijkstra())) {
       for (int i = 0; i < n; i++) {</pre>
         dis[i] += pot[i] - pot[s];
       T1 add = INF1;
       for (int i = t; i != s; i = E[rt[i] ^ 1].v) {
        add = min(add, E[rt[i]].f);
       for (int i = t; i != s; i = E[rt[i] ^ 1].v) {
        E[rt[i]].f -= add, E[rt[i] ^ 1].f += add;
       flow += add, cost += add * dis[t];
       fr = false;
       for (int i = 0; i < n; ++i) swap(dis[i], pot[i]);</pre>
    return make_pair(flow, cost);
  void init(int n) {
    n = _n, m = 0;
for (int i = 0; i < n; ++i) adj[i].clear();</pre>
  void reset() {
    for (int i = 0; i < m; ++i) E[i].f = 0;</pre>
  void add_edge(int u, int v, T1 f, T2 c) {
    adj[u].pb(m), E[m++] = \{v, f, c\};
     adj[v].pb(m), E[m++] = {u, 0, -c};
};
```

## 3.3 Kuhn Munkres

```
void add_edge(int a, int b, int wei) {
    w[a][b] = wei;
  bool check(int x) {
    if (vl[x] = 1, ~fl[x]) return q.push(fl[x]), vr[fl[
        x]] = 1;
    while (\sim x) swap(x, fr[fl[x] = pre[x]]);
    return 0:
  void bfs(int s) {
    fill(slk, slk + n, INF), fill(vl, vl + n, 0), fill(
        vr, vr + n, 0);
    q.push(s), vr[s] = 1;
    while (1) {
      T d;
      while (!q.empty()) {
        int y = q.front(); q.pop();
        for (int x = 0; x < n; ++x)
          if (!vl[x] \&\& slk[x] >= (d = hl[x] + hr[y] -
              w[x][y])
            if (pre[x] = y, d) slk[x] = d;
            else if (!check(x)) return;
      d = INF;
      for (int x = 0; x < n; ++x)
        if (!v1[x] && d > s1k[x]) d = s1k[x];
      for (int x = 0; x < n; ++x) {
        if (v1[x]) h1[x] += d;
        else slk[x] -= d;
        if (vr[x]) hr[x] -= d;
      for (int x = 0; x < n; ++x) if (!v1[x] && !s1k[x]
           && !check(x)) return;
    }
  }
  T solve() {
    fill(fl, fl + n, -1), fill(fr, fr + n, -1), fill(hr
          hr + n, 0);
    for (int i = 0; i < n; ++i) hl[i] = *max_element(w[</pre>
        i], w[i] + n);
    for (int i = 0; i < n; ++i) bfs(i);</pre>
    for (int i = 0; i < n; ++i) res += w[i][fl[i]];</pre>
    return res;
 }
};
```

#### 3.4 SW Min Cut

```
template <typename T>
struct SW { // 0-based
  T g[N][N], sum[N]; int n;
  bool vis[N], dead[N];
  void init(int _n) {
    n = _n;
for (int i = 0; i < n; ++i) fill(g[i], g[i] + n, 0)</pre>
    fill(dead, dead + n, false);
  void add_edge(int u, int v, T w) {
    g[u][v] += w, g[v][u] += w;
  T solve() {
    T ans = 1 << 30;
    for (int round = 0; round + 1 < n; ++round) {</pre>
      fill(vis, vis + n, false), fill(sum, sum + n, 0);
      int num = 0, s = -1, t = -1;
      while (num < n - round) {
        int now = -1;
        for (int i = 0; i < n; ++i) if (!vis[i] && !</pre>
             dead[i]) {
            if (now == -1 || sum[now] < sum[i]) now = i</pre>
        s = t, t = now;
        vis[now] = true, num++;
        for (int i = 0; i < n; ++i) if (!vis[i] && !</pre>
             dead[i]) {
             sum[i] += g[now][i];
      }
```

```
ans = min(ans, sum[t]);
  for (int i = 0; i < n; ++i) {
    g[i][s] += g[i][t];
    g[s][i] += g[t][i];
  }
  dead[t] = true;
}
return ans;
};</pre>
```

## 3.5 Gomory Hu Tree

```
vector <array <int, 3>> GomoryHu(vector <vector <pii>>>
    adj, int n) {
  Tree edge min -> mincut (0-based)
  Dinic flow(n);
  for (int i = 0; i < n; ++i) for (auto [j, w] : adj[i</pre>
      1)
      flow.add_edge(i, j, w);
  flow.record();
  vector <array <int, 3>> ans;
  vector <int> rt(n);
  for (int i = 0; i < n; ++i) rt[i] = 0;</pre>
  for (int i = 1; i < n; ++i) {</pre>
    int t = rt[i];
    flow.reset(); // clear flows on all edge
    ans.push_back({i, t, flow.solve(i, t)});
    flow.runbfs(i);
    for (int j = i + 1; j < n; ++j) if (rt[j] == t &&</pre>
        flow.vis[j]) {
        rt[j] = i;
  }
  return ans;
}
```

#### 3.6 Blossom

```
struct Matching { // 0-based
  int fa[N], pre[N], match[N], s[N], v[N], n, tk;
  vector <int> g[N];
  queue <int> q;
  Matching (int _n) : n(_n), tk(0) {
    for (int i = 0; i <= n; ++i) match[i] = pre[i] = n;</pre>
    for (int i = 0; i < n; ++i) g[i].clear();</pre>
  void add_edge(int u, int v) {
    g[u].push_back(v), g[v].push_back(u);
  int Find(int u) {
    return u == fa[u] ? u : fa[u] = Find(fa[u]);
  int lca(int x, int y) {
    tk++;
    x = Find(x), y = Find(y);
    for (; ; swap(x, y)) {
      if (x != n) {
        if (v[x] == tk) return x;
        v[x] = tk;
        x = Find(pre[match[x]]);
      }
    }
  void blossom(int x, int y, int 1) {
    while (Find(x) != 1) {
      pre[x] = y, y = match[x];
      if (s[y] == 1) q.push(y), s[y] = 0;
      if (fa[x] == x) fa[x] = 1;
      if (fa[y] == y) fa[y] = 1;
      x = pre[y];
    }
  bool bfs(int r) {
    for (int i = 0; i <= n; ++i) fa[i] = i, s[i] = -1;</pre>
    while (!q.empty()) q.pop();
    q.push(r);
    s[r] = 0;
    while (!q.empty()) {
      int x = q.front(); q.pop();
      for (int u : g[x]) {
        if (s[u] == -1) {
```

```
pre[u] = x, s[u] = 1;
          if (match[u] == n) {
            for (int a = u, b = x, last; b != n; a =
                last, b = pre[a])
              last = match[b], match[b] = a, match[a] =
            return true;
          q.push(match[u]);
          s[match[u]] = 0;
        } else if (!s[u] && Find(u) != Find(x)) {
          int 1 = lca(u, x);
          blossom(x, u, 1);
          blossom(u, x, 1);
     }
    return false;
  int solve() {
    int res = 0;
    for (int x = 0; x < n; ++x) {
      if (match[x] == n) res += bfs(x);
    return res:
 }
};
```

## 3.7 Weighted Blossom

```
struct WeightGraph { // 1-based
  static const int inf = INT_MAX;
  static const int maxn = 514;
  struct edge {
    int u, v, w;
    edge(){}
    edge(int u, int v, int w): u(u), v(v), w(w) {}
  };
  int n, n_x;
  edge g[maxn * 2][maxn * 2];
  int lab[maxn * 2];
  int match[maxn * 2], slack[maxn * 2], st[maxn * 2],
      pa[maxn * 2];
 int flo_from[maxn * 2][maxn + 1], S[maxn * 2], vis[
      maxn * 2];
  vector<int> flo[maxn * 2];
  queue<int> q;
  int e_delta(const edge &e) { return lab[e.u] + lab[e.
  v] - g[e.u][e.v].w * 2; }
void update_slack(int u, int x) { if (!slack[x] ||
      e_{delta}(g[u][x]) < e_{delta}(g[slack[x]][x])) slack
      [x] = u; 
  void set_slack(int x) {
    slack[x] = 0;
    for (int u = 1; u <= n; ++u)</pre>
      if (g[u][x].w > 0 && st[u] != x && S[st[u]] == 0)
        update_slack(u, x);
  void q_push(int x) {
    if (x \le n) q.push(x);
    else for (size_t i = 0; i < flo[x].size(); i++)</pre>
        q_push(flo[x][i]);
  void set_st(int x, int b) {
    st[x] = b;
    if (x > n) for (size_t i = 0; i < flo[x].size(); ++</pre>
        i) set_st(flo[x][i], b);
  int get_pr(int b, int xr) {
    int pr = find(flo[b].begin(), flo[b].end(), xr) -
        flo[b].begin();
    if (pr % 2 == 1) {
      reverse(flo[b].begin() + 1, flo[b].end());
      return (int)flo[b].size() - pr;
    return pr;
  void set_match(int u, int v) {
    match[u] = g[u][v].v;
    if (u <= n) return;</pre>
    edge e = g[u][v];
    int xr = flo_from[u][e.u], pr = get_pr(u, xr);
```

```
for (int i = 0; i < pr; ++i) set_match(flo[u][i],</pre>
      flo[u][i ^ 1]);
  set_match(xr, v);
  rotate(flo[u].begin(), flo[u].begin() + pr, flo[u].
      end());
void augment(int u, int v) {
  for (; ; ) {
    int xnv = st[match[u]];
    set_match(u, v);
    if (!xnv) return;
    set_match(xnv, st[pa[xnv]]);
    u = st[pa[xnv]], v = xnv;
int get_lca(int u, int v) {
 static int t = 0:
  for (++t; u || v; swap(u, v)) {
    if (u == 0) continue;
    if (vis[u] == t) return u;
    vis[u] = t;
    u = st[match[u]];
   if (u) u = st[pa[u]];
 return 0:
void add_blossom(int u, int lca, int v) {
 int b = n + 1;
  while (b <= n_x && st[b]) ++b;</pre>
  if (b > n_x) ++n_x;
  lab[b] = 0, S[b] = 0;
  match[b] = match[lca];
  flo[b].clear();
  flo[b].push_back(lca);
  for (int x = u, y; x != lca; x = st[pa[y]])
    flo[b].push_back(x), flo[b].push_back(y = st[
        match[x]]), q_push(y);
  reverse(flo[b].begin() + 1, flo[b].end());
  for (int x = v, y; x != lca; x = st[pa[y]])
    flo[b].push_back(x), flo[b].push_back(y = st[
        match[x]]), q_push(y);
  set_st(b, b);
  for (int x = 1; x \le n_x; ++x) g[b][x].w = g[x][b].
      w = 0;
  for (int x = 1; x <= n; ++x) flo_from[b][x] = 0;
  for (size_t i = 0; i < flo[b].size(); ++i) {</pre>
    int xs = flo[b][i];
    for (int x = 1; x <= n_x; ++x)
      if (g[b][x].w == 0 | e_delta(g[xs][x]) <</pre>
          e_delta(g[b][x]))
        g[b][x] = g[xs][x], g[x][b] = g[x][xs];
    for (int x = 1; x <= n; ++x)
      if (flo_from[xs][x]) flo_from[b][x] = xs;
  set_slack(b);
void expand blossom(int b) {
  for (size_t i = 0; i < flo[b].size(); ++i)</pre>
    set_st(flo[b][i], flo[b][i]);
  int xr = flo_from[b][g[b][pa[b]].u], pr = get_pr(b)
       xr);
  for (int i = 0; i < pr; i += 2) {
    int xs = flo[b][i], xns = flo[b][i + 1];
    pa[xs] = g[xns][xs].u;
    S[xs] = 1, S[xns] = 0;
    slack[xs] = 0, set_slack(xns);
    q_push(xns);
  S[xr] = 1, pa[xr] = pa[b];
  for (size_t i = pr + 1; i < flo[b].size(); ++i) {</pre>
   int xs = flo[b][i];
    S[xs] = -1, set_slack(xs);
 st[b] = 0;
bool on_found_edge(const edge &e) {
  int u = st[e.u], v = st[e.v];
  if (S[v] == -1) {
    pa[v] = e.u, S[v] = 1;
    int nu = st[match[v]];
    slack[v] = slack[nu] = 0;
    S[nu] = 0, q_push(nu);
```

```
} else if (S[v] == 0) {
    int lca = get_lca(u, v);
    if (!lca) return augment(u,v), augment(v,u), true
    else add_blossom(u, lca, v);
  return false;
bool matching() {
  memset(S + 1, -1, sizeof(int) * n_x);
  memset(slack + 1, 0, sizeof(int) * n_x);
  q = queue<int>();
  for (int x = 1; x <= n_x; ++x)
    if (st[x] == x \&\& !match[x]) pa[x] = 0, S[x] = 0,
         q_push(x);
  if (q.empty()) return false;
  for (; ; ) {
    while (q.size()) {
      int u = q.front(); q.pop();
      if (S[st[u]] == 1) continue;
      for (int v = 1; v <= n; ++v)
        if (g[u][v].w > 0 && st[u] != st[v]) {
          if (e_delta(g[u][v]) == 0) {
            if (on_found_edge(g[u][v])) return true;
          } else update_slack(u, st[v]);
    int d = inf;
    for (int b = n + 1; b <= n_x; ++b)
      if (st[b] == b && S[b] == 1) d = min(d, lab[b]
          / 2);
    for (int x = 1; x <= n_x; ++x)
      if (st[x] == x && slack[x]) {
        if (S[x] == -1) d = min(d, e_delta(g[slack[x
            ]][x]));
        else if (S[x] == 0) d = min(d, e_delta(g[
            slack[x]][x]) / 2);
    for (int u = 1; u <= n; ++u) {</pre>
      if (S[st[u]] == 0) {
        if (lab[u] <= d) return 0;</pre>
        lab[u] -= d;
      } else if (S[st[u]] == 1) lab[u] += d;
    for (int b = n + 1; b \le n_x; ++b)
      if (st[b] == b) {
        if (S[st[b]] == 0) lab[b] += d * 2;
        else if (S[st[b]] == 1) lab[b] -= d * 2;
    q = queue<int>();
    for (int x = 1; x <= n_x; ++x)</pre>
      if (st[x] == x && slack[x] && st[slack[x]] != x
           && e_delta(g[slack[x]][x]) == 0)
        if (on_found_edge(g[slack[x]][x])) return
            true;
    for (int b = n + 1; b \le n_x; ++b)
      if (st[b] == b && S[b] == 1 && lab[b] == 0)
          expand_blossom(b);
  return false;
}
pair<long long, int> solve() {
  memset(match + 1, 0, sizeof(int) * n);
  n_x = n;
  int n_matches = 0;
  long long tot_weight = 0;
  for (int u = 0; u \le n; ++u) st[u] = u, flo[u].
      clear();
  int w_max = 0;
  for (int u = 1; u <= n; ++u)</pre>
    for (int v = 1; v <= n; ++v) {
      flo_from[u][v] = (u == v ? u : 0);
      w_max = max(w_max, g[u][v].w);
  for (int u = 1; u <= n; ++u) lab[u] = w_max;</pre>
  while (matching()) ++n_matches;
  for (int u = 1; u <= n; ++u)</pre>
    if (match[u] && match[u] < u)</pre>
      tot_weight += g[u][match[u]].w;
  return make_pair(tot_weight, n_matches);
void add_edge(int ui, int vi, int wi) { g[ui][vi].w =
```

```
g[vi][ui].w = wi; }
  void init(int _n) {
    n = _n;
    for (int u = 1; u <= n; ++u)</pre>
      for (int v = 1; v <= n; ++v)
        g[u][v] = edge(u, v, 0);
  }
};
```

## 3.8 Flow Model

- Maximum/Minimum flow with lower bound / Circulation problem
  - 1. Construct super source S and sink T.
  - 2. For each edge (x,y,l,u), connect x o y with capacity u-l.
  - 3. For each vertex v, denote by in(v) the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.
  - 4. If in(v)>0, connect  $S\to v$  with capacity in(v), otherwise, connect  $v\to T$  with capacity -in(v).
    - To maximize, connect t 
      ightarrow s with capacity  $\infty$  (skip this in circulation problem), and let f be the maximum flow from S to T. If  $f \neq \sum_{v \in V, in(v) > 0} in(v)$ , there's no solution. Otherwise, the maximum flow from s to t is the answer
    - To minimize, let f be the maximum flow from S to T. Connect  $t \to s$  with capacity  $\infty$  and let the flow from S to T be f'. If  $f+f' \neq \sum_{v \in V, in(v)>0} in(v)$ , there's no solution. Otherwise, f' is the answer.
  - 5. The solution of each edge e is  $l_e+f_e$ , where  $f_e$  corresponds to the flow of edge  $\boldsymbol{e}$  on the graph.
- ullet Construct minimum vertex cover from maximum matching M on bipar- ${\rm tite\ graph\ }(X,Y)$ 
  - 1. Redirect every edge:  $y \to x$  if  $(x,y) \in M$ ,  $x \to y$  otherwise.

  - 2. DFS from unmatched vertices in X. 3.  $x \in X$  is chosen iff x is unvisited. 4.  $y \in Y$  is chosen iff y is visited.
- Maximum density induced subgraph
  - 1. Binary search on answer, suppose we're checking answer  ${\cal T}$
  - 2. Construct a max flow model, let K be the sum of all weights 3. Connect source  $s\to v$ ,  $v\in G$  with capacity K

  - 4. For each edge (u,v,w) in G, connect  $u \to v$  and  $v \to u$  with capacity  $\boldsymbol{w}$
  - 5. For  $v \in G$ , connect it with sink  $v \to t$  with capacity  $K + 2T (\sum_{e \in E(v)} w(e)) 2w(v)$
  - 6. T is a valid answer if the maximum flow  $f < K \lvert V \rvert$
- Minimum weight edge cover 1. For each  $v \in V$  create a copy v' , and connect  $u' \to v'$  with
  - weight w(u,v). 2. Connect  $v \to v'$  with weight  $2\mu(v)$ , where  $\mu(v)$  is the cost of the cheapest edge incident to v .
  - 3. Find the minimum weight perfect matching on  $G^{\prime}$ .
- Project selection problem
  - 1. If  $p_v > 0$ , create edge (s,v) with capacity  $p_v$ ; otherwise,
  - create edge (v,t) with capacity  $-p_v$ . 2. Create edge (u,v) with capacity w with w being the cost of choosing u without choosing v.
  - 3. The mincut is equivalent to the maximum profit of a subset
- 0/1 quadratic programming

$$\sum_{x} c_{x}x + \sum_{y} c_{y}\bar{y} + \sum_{xy} c_{xy}x\bar{y} + \sum_{xyx'y'} c_{xyx'y'}(x\bar{y} + x'\bar{y'})$$

can be minimized by the mincut of the following graph:

- 1. Create edge (x,t) with capacity  $c_x$  and create edge (s,y) with
- capacity  $c_y$  . Create edge (x,y) with capacity  $c_{xy}$
- 3. Create edge (x,y) and edge (x',y') with capacity  $c_{xyx'y'}$ .

# 4 Graph

## 4.1 Heavy-Light Decomposition

```
vector<int> dep, pa, sz, ch, hd, id;
int _id;
void dfs(int i, int p) {
  dep[i] = \sim p ? dep[p] + 1 : 0;
  pa[i] = p, sz[i] = 1, ch[i] = -1;
  for (int j : g[i])
    if (j != p) {
      dfs(j, i);
      if (ch[i] == -1 || sz[ch[i]] < sz[j]) ch[i] = j;</pre>
      sz[i] += sz[j];
```

```
void hld(int i, int p, int h) {
  hd[i] = h;
  id[i] = _id++;
  if (~ch[i]) hld(ch[i], i, h);
  for (int j : g[i]) if (j != p && j != ch[i])
     hld(j, i, j);
}
void query(int i, int j) {
  while (hd[i] != hd[j]) {
    if (dep[hd[i]] < dep[hd[j]]) swap(i, j);
     query2(id[hd[i]], id[i] + 1), i = pa[hd[i]];
  }
  if (dep[i] < dep[j]) swap(i, j);
  query2(id[j], id[i] + 1);
}</pre>
```

## 4.2 Centroid Decomposition

```
vector<vector<int>> dis; // dis[n][logn]
vector<int> pa, sz, dep;
vector<bool> vis;
void dfs_sz(int i, int p) {
  sz[i] = 1;
  for (int j : g[i]) if (j != p && !vis[j])
    dfs_sz(j, i), sz[i] += sz[j];
int cen(int i, int p, int _n) {
  for (int j : g[i]) if (j != p && !vis[j] && sz[j] >
      _n / 2)
    return cen(j, i, _n);
  return i;
void dfs_dis(int i, int p, int d) { // from i to
    ancestor with depth d
  dis[i][d] = \sim p ? dis[p][d] + 1 : 0;
  for (int j : g[i]) if (j != p && !vis[j])
    dfs_dis(j, i, d);
void cd(int i, int p, int d) {
  dfs_sz(i, -1), i = cen(i, -1, sz[i]);
  vis[i] = true, pa[i] = p, dep[i] = d;
  dfs_dis(i, -1, d);
  for (int j : g[i]) if (!vis[j])
    cd(j, i, d + 1);
}
```

## 4.3 Edge BCC

```
vector<int> low, dep, bcc id, stk;
vector<bool> vis;
int _id;
void dfs(int i, int p) {
  low[i] = dep[i] = \sim p ? dep[p] + 1 : 0;
  stk.push_back(i);
  vis[i] = true;
  for (int j : g[i])
    if (j != p) {
      if (!vis[j])
        dfs(j, i), low[i] = min(low[i], low[j]);
       else
        low[i] = min(low[i], dep[j]);
  if (low[i] == dep[i]) {
    int id = _id++;
    while (stk.back() != i) {
      int x = stk.back();
       stk.pop_back();
      bcc_id[x] = id;
    stk.pop_back();
    bcc_id[i] = id;
}
```

#### 4.4 Block Cut Tree

```
vector<vector<int>> g, _g;
vector<int>> dep, low, stk;
void dfs(int i, int p) {
  dep[i] = low[i] = ~p ? dep[p] + 1 : 0;
  stk.push_back(i);
  for (int j : g[i]) if (j != p) {
```

```
if (dep[j] == -1) {
    dfs(j, i), low[i] = min(low[i], low[j]);
    if (low[j] >= dep[i]) {
        int id = _g.size();
        _g.emplace_back();
        while (stk.back() != j) {
            int x = stk.back();
            stk.pop_back();
            _g[x].push_back(id), _g[id].push_back(x);
        }
        stk.pop_back();
        _g[j].push_back(id), _g[id].push_back(j);
        _g[j].push_back(id), _g[id].push_back(i);
    }
    } else low[i] = min(low[i], dep[j]);
}
```

#### 4.5 SCC / 2SAT

```
struct SAT {
  vector<vector<int>> g;
   vector<int> dep, low, scc_id;
   vector<bool> is;
   vector<int> stk;
   int n, _id, _t;
   SAT() {}
   void init(int _n) {
     n = _n, _id = _t = 0;
     g.assign(2 * n, vector<int>());
dep.assign(2 * n, -1), low.assign(2 * n, -1);
scc_id.assign(2 * n, -1), is.assign(2 * n, false);
     stk.clear();
   void add_edge(int x, int y) {g[x].push_back(y);}
   int rev(int i) {return i < n ? i + n : i - n;}</pre>
   void add_ifthen(int x, int y) {add_clause(rev(x), y)
       ;}
   void add_clause(int x, int y) {
     add_edge(rev(x), y);
     add_edge(rev(y), x);
   void dfs(int i) {
     dep[i] = low[i] = _t++;
     stk.push_back(i);
     for (int j : g[i])
       if (scc_id[j] == -1) {
         if (dep[j] == -1)
           dfs(j);
         low[i] = min(low[i], low[j]);
     if (low[i] == dep[i]) {
       int id = _id++;
       while (stk.back() != i) {
         int x = stk.back();
          stk.pop_back();
         scc id[x] = id;
       }
       stk.pop_back();
       scc_id[i] = id;
   bool solve() {
     for (int i = 0; i < 2 * n; ++i)
       if (dep[i] == -1)
         dfs(i);
     for (int i = 0; i < n; ++i) {</pre>
       if (scc_id[i] == scc_id[i + n]) return false;
       if (scc_id[i] < scc_id[i + n])</pre>
         is[i] = true;
       else
         is[i + n] = true;
     return true;
|};
```

#### 4.6 Virtual Tree

```
vector<vector<int>> _g;
vector<int> st, ed, stk;
void solve(vector<int> v) {
```

template <typename T> struct DMST { // 1-based

### 4.7 Directed MST

```
T g[maxn][maxn], fw[maxn];
int n, fr[maxn];
bool vis[maxn], inc[maxn];
void clear() {
  for (int i = 0; i < maxn; ++i) {</pre>
    for (int j = 0; j < maxn; ++j) g[i][j] = inf;</pre>
    vis[i] = inc[i] = false;
  }
}
void addedge(int u, int v, T w) {
  g[u][v] = min(g[u][v], w);
T query(int root, int _n) {
  n = n;
  if (dfs(root) != n) return -1;
  T ans = 0;
  while (true) {
    for (int i = 1; i <= n; ++i) fw[i] = inf, fr[i] =</pre>
    for (int i = 1; i <= n; ++i) if (!inc[i]) {</pre>
      for (int j = 1; j <= n; ++j) {</pre>
        if (!inc[j] && i != j && g[j][i] < fw[i]) {</pre>
          fw[i] = g[j][i];
           fr[i] = j;
        }
      }
    }
    int x = -1;
    for (int i = 1; i <= n; ++i) if (i != root && !</pre>
        inc[i]) {
      int j = i, c = 0;
      while (j != root && fr[j] != i && c <= n) ++c,</pre>
           j = fr[j];
      if (j == root || c > n) continue;
      else { x = i; break; }
    if (!~x) {
      for (int i = 1; i <= n; ++i) if (i != root && !</pre>
           inc[i]) ans += fw[i];
      return ans;
    for (int i = 1; i <= n; ++i) vis[i] = false;</pre>
    do { ans += fw[y]; y = fr[y]; vis[y] = inc[y] =
         true; } while (y != x);
    inc[x] = false;
    for (int k = 1; k <= n; ++k) if (vis[k]) {</pre>
      for (int j = 1; j <= n; ++j) if (!vis[j]) {</pre>
           if (g[x][j] > g[k][j]) g[x][j] = g[k][j];
           if (g[j][k] < inf \&\& g[j][k] - fw[k] < g[j]
               ][x]) g[j][x] = g[j][k] - fw[k];
        }
    }
  return ans;
int dfs(int now) {
  int r = 1:
  vis[now] = true;
  for (int i = 1; i <= n; ++i) if (g[now][i] < inf &&</pre>
        !vis[i]) r += dfs(i);
```

```
return r;
}
};
```

## 4.8 Dominator Tree

```
struct Dominator_tree {
  int n, id;
  vector <vector <int>> adj, radj, bucket;
  vector <int> sdom, dom, vis, rev, par, rt, mn;
  Dominator_tree (int _n) : n(_n), id(0) {
    adj.resize(n), radj.resize(n), bucket.resize(n);
    sdom.resize(n), dom.resize(n, -1), vis.resize(n,
    rev.resize(n), rt.resize(n), mn.resize(n), par.
        resize(n);
  void add_edge(int u, int v) {adj[u].pb(v);}
  int query(int v, bool x) {
    if (rt[v] == v) return x ? -1 : v;
    int p = query(rt[v], true);
    if (p == -1) return x ? rt[v] : mn[v];
    if (sdom[mn[v]] > sdom[mn[rt[v]]]) mn[v] = mn[rt[v
        ]];
    rt[v] = p;
    return x ? p : mn[v];
  void dfs(int v) {
    vis[v] = id, rev[id] = v;
    rt[id] = mn[id] = sdom[id] = id, id++;
    for (int u : adj[v]) {
      if (vis[u] == -1) dfs(u), par[vis[u]] = vis[v];
      radj[vis[u]].pb(vis[v]);
    }
  void build(int s) {
    dfs(s);
    for (int i = id - 1; ~i; --i) {
      for (int u : radj[i]) {
        sdom[i] = min(sdom[i], sdom[query(u, false)]);
      if (i) bucket[sdom[i]].pb(i);
      for (int u : bucket[i]) {
        int p = query(u, false);
        dom[u] = sdom[p] == i ? i : p;
      if (i) rt[i] = par[i];
    }
    vector <int> res(n, -1);
    for (int i = 1; i < id; ++i) {</pre>
      if (dom[i] != sdom[i]) dom[i] = dom[dom[i]];
    for (int i = 1; i < id; ++i) res[rev[i]] = rev[dom[</pre>
        i]];
    res[s] = s;
    dom = res;
};
```

# 5 String

### 5.1 Aho-Corasick Automaton

```
struct AC {
  int ch[N][26], to[N][26], fail[N], sz;
  vector <int> g[N];
  int cnt[N];
  AC () {sz = 0, extend();}
  void extend() {fill(ch[sz], ch[sz] + 26, 0), sz++;}
  int nxt(int u, int v) {
    if (!ch[u][v]) ch[u][v] = sz, extend();
    return ch[u][v];
  int insert(string s) {
    int now = 0;
    for (char c : s) now = nxt(now, c - 'a');
    cnt[now]++;
    return now:
  void build_fail() {
    queue <int> q;
```

```
for (int i = 0; i < 26; ++i) if (ch[0][i]) {</pre>
       q.push(ch[0][i]);
       g[0].push_back(ch[0][i]);
     while (!q.empty()) {
       int v = q.front(); q.pop();
       for (int j = 0; j < 26; ++j) {</pre>
         to[v][j] = ch[v][j] ? v : to[fail[v]][j];
       for (int i = 0; i < 26; ++i) if (ch[v][i]) {</pre>
         int u = ch[v][i], k = fail[v];
         while (k && !ch[k][i]) k = fail[k];
         if (ch[k][i]) k = ch[k][i];
         fail[u] = k;
         cnt[u] += cnt[k], g[k].push_back(u);
         q.push(u);
       }
    }
  int match(string &s) {
     int now = 0, ans = 0;
     for (char c : s) {
       now = to[now][c - 'a'];
if (ch[now][c - 'a']) now = ch[now][c - 'a'];
       ans += cnt[now]:
     return ans;
  }
};
```

# 5.2 KMP Algorithm

```
vector <int> build fail(string s) {
  vector <int> f(s.length() + 1, 0);
  int k = 0;
  for (int i = 1; i < s.length(); ++i) {
  while (k && s[k] != s[i]) k = f[k];</pre>
    if (s[k] == s[i]) k++;
    f[i + 1] = k;
  return f:
int match(string s, string t) {
  vector <int> f = build_fail(t);
  int k = 0, ans = 0;
  for (int i = 0; i < s.length(); ++i) {</pre>
    while (k && s[i] != t[k]) k = f[k];
    if (s[i] == t[k]) k++;
    if (k == t.length()) ans++, k = f[k];
  return ans;
}
```

# 5.3 Z Algorithm

## 5.4 Manacher

```
vector <int> manacher(string &s) {
   string t = "^#";
   for (char c : s) t += c, t += '#';
   t += '&';
   int n = t.length();
   vector <int> r(n, 0);
   int C = 0, R = 0;
   for (int i = 1; i < n - 1; ++i) {
     int mirror = 2 * C - i;
     r[i] = (i < R ? min(r[mirror], R - i) : 0);
   while (t[i - 1 - r[i]] == t[i + 1 + r[i]]) r[i]++;</pre>
```

```
if (i + r[i] > R) R = i + r[i], C = i;
}
return r;
}
```

## 5.5 Suffix Array

```
int sa[N], tmp[2][N], c[N], rk[N], lcp[N];
void buildSA(string s) {
  int *x = tmp[0], *y = tmp[1], m = 256, n = s.length()
  for (int i = 0; i < m; ++i) c[i] = 0;</pre>
  for (int i = 0; i < n; ++i) c[x[i] = s[i]]++;</pre>
  for (int i = 1; i < m; ++i) c[i] += c[i - 1];
  for (int i = n - 1; ~i; --i) sa[--c[x[i]]] = i;
  for (int k = 1; k < n; k <<= 1) {
    for (int i = 0; i < m; ++i) c[i] = 0;</pre>
    for (int i = 0; i < n; ++i) c[x[i]]++;</pre>
    for (int i = 1; i < m; ++i) c[i] += c[i - 1];</pre>
    int p = 0;
    for (int i = n - k; i < n; ++i) y[p++] = i;
    for (int i = 0; i < n; ++i) if (sa[i] >= k) y[p++]
         = sa[i] - k;
    for (int i = n - 1; ~i; --i) sa[--c[x[y[i]]]] = y[i
    y[sa[0]] = p = 0;
    for (int i = 1; i < n; ++i) {
      int a = sa[i], b = sa[i - 1];
      if (!(x[a] == x[b] \&\& a + k < n \&\& b + k < n \&\& x)
           [a + k] == x[b + k]) p++;
      y[sa[i]] = p;
    if (n == p + 1) break;
    swap(x, y), m = p + 1;
void buildLCP(string s) {
   // Lcp[i] = LCP(sa[i - 1], sa[i])
  // lcp(i, j) = min(lcp[rk[i] + 1], lcp[rk[i] + 2],
      ..., lcp[rk[j]])
  int n = s.length(), val = 0;
  for (int i = 0; i < n; ++i) rk[sa[i]] = i;</pre>
  for (int i = 0; i < n; ++i) {</pre>
    if (!rk[i]) lcp[rk[i]] = 0;
    else {
      if (val) val--
      int p = sa[rk[i] - 1];
      while (val + i < n && val + p < n && s[val + i]
           == s[val + p]) val++;
      lcp[rk[i]] = val;
  }
}
```

## 5.6 SAIS

```
namespace sfx {
bool _t[N * 2];
int SA[N * 2], H[N], RA[N];
int _s[N * 2], _c[N * 2], x[N], _p[N],
                                         _q[N * 2];
void pre(int *sa, int *c, int n, int z) {
  fill_n(sa, n, 0), copy_n(c, z, x);
void induce(int *sa, int *c, int *s, bool *t, int n,
    int z) {
  copy_n(c, z - 1, x + 1);
  for (int i = 0; i < n; ++i) if (sa[i] && !t[sa[i] -</pre>
      1]) sa[x[s[sa[i] - 1]]++] = sa[i] - 1;
  copy_n(c, z, x);
for (int i = n - 1; i >= 0; --i) if (sa[i] && t[sa[i]
        - 1]) sa[--x[s[sa[i] - 1]]] = sa[i] - 1;
void sais(int *s, int *sa, int *p, int *q, bool *t, int
     *c, int n, int z) {
  bool uniq = t[n - 1] = true;
  int nn = 0, nmxz = -1, *nsa = sa + n, *ns = s + n,
      last = -1;
  fill_n(c, z, 0);
  for (int i = 0; i < n; ++i) uniq &= ++c[s[i]] < 2;</pre>
  partial_sum(c, c + z, c);
  if (uniq) {
    for (int i = 0; i < n; ++i) sa[--c[s[i]]] = i;</pre>
```

```
return:
  for (int i = n - 2; i >= 0; --i)
    t[i] = (s[i] == s[i + 1] ? t[i + 1] : s[i] < s[i +
         1]);
  pre(sa, c, n, z);
  for (int i = 1; i <= n - 1; ++i)</pre>
    if (t[i] && !t[i - 1])
      sa[--x[s[i]]] = p[q[i] = nn++] = i;
  induce(sa, c, s, t, n, z);
for (int i = 0; i < n; ++i)
    if (sa[i] && t[sa[i]] && !t[sa[i] - 1]) {
      bool neq = last < 0 \mid | !equal(s + sa[i], s + p[q[
           sa[i]] + 1], s + last);
      ns[q[last = sa[i]]] = nmxz += neq;
  sais(ns, nsa, p + nn, q + n, t + n, c + z, nn, nmxz + m
        1);
  pre(sa, c, n, z);
  for (int i = nn - 1; i >= 0; --i)
    sa[--x[s[p[nsa[i]]]]] = p[nsa[i]];
  induce(sa, c, s, t, n, z);
vector<int> build(int *s, int n) {
 copy_n(s, n, _s), _s[n] = 0;
sais(_s, SA, _p, _q, _t, _c, n + 1, 256);
vector <int> sa(n);
  for (int i = 0; i < n; ++i)</pre>
    sa[i] = SA[i + 1];
  return sa;
```

#### 5.7 Suffix Automaton

```
struct SAM {
  int ch[N][26], len[N], link[N], pos[N], cnt[N], sz;
  // node -> strings with the same endpos set
  // length in range [len(link) + 1, len]
  // node's endpos set -> pos in the subtree of node
  // link -> longest suffix with different endpos set
  // len -> longest suffix
  // pos -> end position
// cnt -> size of endpos set
  SAM () \{len[0] = 0, link[0] = -1, pos[0] = 0, cnt[0] \}
      = 0, sz = 1;
  void build(string s) {
    int last = 0;
    for (int i = 0; i < s.length(); ++i) {</pre>
      char c = s[i];
      int cur = sz++;
      len[cur] = len[last] + 1, pos[cur] = i + 1;
      int p = last;
      while (\sim p \&\& !ch[p][c - 'a']) ch[p][c - 'a'] =
           cur, p = link[p];
      if (p == -1) {
        link[cur] = 0;
      } else {
        int q = ch[p][c - 'a'];
        if (len[p] + 1 == len[q]) {
          link[cur] = q;
        } else {
          int nxt = sz++;
          len[nxt] = len[p] + 1, link[nxt] = link[q],
               pos[nxt] = 0;
          for (int j = 0; j < 26; ++j) ch[nxt][j] = ch[
               q][j];
          while (~p && ch[p][c - 'a'] == q) ch[p][c - '
    a'] = nxt, p = link[p];
          link[q] = link[cur] = nxt;
        }
      cnt[cur]++;
      last = cur;
    vector <int> p(sz);
    iota(all(p), 0);
    sort(all(p), [&](int i, int j) {return len[i] > len
         [j];});
    for (int i = 0; i < sz; ++i) cnt[link[p[i]]] += cnt</pre>
         [p[i]];
  }
```

# 5.8 Minimum Rotation

} sam;

```
string rotate(const string &s) {
  int n = s.length();
  string t = s + s;
  int i = 0, j = 1;
  while (i < n && j < n) {
   int k = 0;
  while (k < n && t[i + k] == t[j + k]) ++k;
   if (t[i + k] <= t[j + k]) j += k + 1;
   else i += k + 1;
   if (i == j) ++j;
  }
  int pos = (i < n ? i : j);
  return t.substr(pos, n);
}</pre>
```

## 5.9 Palindrome Tree

```
struct PAM {
  int ch[N][26], cnt[N], fail[N], len[N], sz;
  // 0 -> even root, 1 -> odd root
  PAM (string _s) : s(_s) {
    sz = 0;
    extend(), extend();
    len[0] = 0, fail[0] = 1, len[1] = -1;
    int lst = 1;
    for (int i = 0; i < s.length(); ++i) {</pre>
      while (s[i - len[lst] - 1] != s[i]) lst = fail[
           lst1:
      if (!ch[lst][s[i] - 'a']) {
        int idx = extend();
        len[idx] = len[lst] + 2;
        int now = fail[lst];
        while (s[i - len[now] - 1] != s[i]) now = fail[
            now];
        fail[idx] = ch[now][s[i] - 'a'];
        ch[lst][s[i] - 'a'] = idx;
      lst = ch[lst][s[i] - 'a'], cnt[lst]++;
    }
  void build_count() {
    for (int i = sz - 1; i > 1; --i)
      cnt[fail[i]] += cnt[i];
  int extend() {
    fill(ch[sz], ch[sz] + 26, 0), sz++;
    return sz - 1;
};
```

#### 5.10 Main Lorentz

```
int to_left[N], to_right[N];
vector <array <int, 3>> rep; // l, r, len.
// substr(l ~ r, len * 2) are tandem
void findRep(string &s, int 1, int r) {
  if (r - 1 == 1) return;
  int m = 1 + r >> 1;
  findRep(s, l, m), findRep(s, m, r);
  string sl = s.substr(1, m - 1), sr = s.substr(m, r - 1)
      m);
  vector <int> Z = buildZ(sr + "#" + sl);
  for (int i = 1; i < m; ++i) to_right[i] = Z[r - m + 1</pre>
       + i - 1];
  reverse(all(sl));
  Z = buildZ(sl);
  for (int i = 1; i < m; ++i) to_left[i] = Z[m - i -</pre>
      1];
  reverse(all(sl));
  for (int i = 1; i + 1 < m; ++i) {</pre>
    int k1 = to_left[i], k2 = to_right[i + 1], len = m
         - i - 1;
    if (k1 < 1 || k2 < 1 || len < 2) continue;</pre>
    int tl = max(1, len - k2), tr = min(len - 1, k1);
    if (tl <= tr) rep.pb({i + 1 - tr, i + 1 - tl, len})</pre>
```

```
Z = buildZ(sr);
   for (int i = m; i < r; ++i) to_right[i] = Z[i - m];</pre>
  reverse(all(sl)), reverse(all(sr));
Z = buildZ(sl + "#" + sr);
   for (int i = m; i < r; ++i) to_left[i] = Z[m - l + 1</pre>
       +r-i-1];
  reverse(all(sl)), reverse(all(sr));
   for (int i = m; i + 1 < r; ++i) {
     int k1 = to_left[i], k2 = to_right[i + 1], len = i
          - m + 1;
     if (k1 < 1 || k2 < 1 || len < 2) continue;</pre>
     int tl = max(len - k2, 1), tr = min(len - 1, k1);
     if (tl <= tr) rep.pb({i + 1 - len - tr, i + 1 - len
           - tl, len});
  Z = buildZ(sr + "#" + sl);
  for (int i = 1; i < m; ++i) {
  if (Z[r - m + 1 + i - 1] >= m - i) {
       rep.pb({i, i, m - i});
  }
}
```

## 6 Math

#### 6.1 Fraction\*

```
struct fraction {
  11 n, d;
  fraction(const ll _n=0, const ll _d=1): n(_n), d(_d)
    11 t = gcd(n, d);
    n /= t, d /= t;
    if (d < 0) n = -n, d = -d;
  fraction operator-() const
  { return fraction(-n, d); }
  fraction operator+(const fraction &b) const
  { return fraction(n * b.d + b.n * d, d * b.d); }
  fraction operator-(const fraction &b) const
  { return fraction(n * b.d - b.n * d, d * b.d); }
  fraction operator*(const fraction &b) const
  { return fraction(n * b.n, d * b.d); }
  fraction operator/(const fraction &b) const
  { return fraction(n * b.d, d * b.n); }
  void print() {
    cout << n;
    if (d != 1) cout << "/" << d;</pre>
};
```

#### 6.2 Miller Rabin / Pollard Rho

```
11 mul(11 x, 11 y, 11 p) {return (x * y - (11)((long
double)x / p * y) * p + p) % p;}
vector<11> chk = {2, 325, 9375, 28178, 450775, 9780504,
     1795265022};
ll Pow(ll a, ll b, ll n) {ll res = 1; for (; b; b >>=
    1, a = mul(a, a, n)) if (b \& 1) res = mul(res, a, n)
    ); return res;}
bool check(ll a, ll d, int s, ll n) {
  a = Pow(a, d, n);
  if (a <= 1) return 1;</pre>
  for (int i = 0; i < s; ++i, a = mul(a, a, n)) {</pre>
    if (a == 1) return 0;
    if (a == n - 1) return 1;
  }
  return 0;
bool IsPrime(ll n) {
  if (n < 2) return 0;
  if (n % 2 == 0) return n == 2;
  \hat{d} = n - 1, s = 0;
  while (d % 2 == 0) d >>= 1, ++s;
  for (ll i : chk) if (!check(i, d, s, n)) return 0;
  return 1;
const vector<ll> small = {2, 3, 5, 7, 11, 13, 17, 19};
11 FindFactor(ll n) {
  if (IsPrime(n)) return 1;
  for (ll p : small) if (n % p == 0) return p;
```

```
11 x, y = 2, d, t = 1;
auto f = [&](11 a) {return (mul(a, a, n) + t) % n;};
  for (int 1 = 2; ; 1 <<= 1) {
    x = y;
    int m = min(1, 32);
    for (int i = 0; i < 1; i += m) {</pre>
       d = 1;
       for (int j = 0; j < m; ++j) {
         y = f(y), d = mul(d, abs(x - y), n);
       11 g = __gcd(d, n);
if (g == n) {
         1 = 1, y = 2, ++t;
         break;
       if (g != 1) return g;
    }
  }
map <11, int> res;
void PollardRho(ll n) {
  if (n == 1) return;
  if (IsPrime(n)) return ++res[n], void(0);
  11 d = FindFactor(n);
  PollardRho(n / d), PollardRho(d);
}
```

#### 6.3 Ext GCD

```
//a * p.first + b * p.second = gcd(a, b)
pair<11, 11> extgcd(11 a, 11 b) {
    pair<11, 11> res;
    if (a < 0) {
       res = extgcd(-a, b);
       res.first *= -1;
       return res;
    }
    if (b < 0) {
       res = extgcd(a, -b);
       res.second *= -1;
       return res;
    }
    if (b == 0) return {1, 0};
    res = extgcd(b, a % b);
    return {res.second, res.first - res.second * (a / b)
       };
}</pre>
```

#### 6.4 PiCount

```
const int V = 10000000, N = 100, M = 100000;
vector<int> primes;
bool isp[V];
 int small_pi[V], dp[N][M];
 void sieve(int x){
   for(int i = 2; i < x; ++i) isp[i] = true;</pre>
   isp[0] = isp[1] = false;
   for(int i = 2; i * i < x; ++i) if(isp[i]) for(int j =</pre>
   i * i; j < x; j += i) isp[j] = false;
for(int i = 2; i < x; ++i) if(isp[i]) primes.</pre>
       push_back(i);
void init(){
   sieve(V);
   small_pi[0] = 0;
   for(int i = 1; i < V; ++i) small_pi[i] = small_pi[i -</pre>
        1] + isp[i];
   for(int i = 0; i < M; ++i) dp[0][i] = i;</pre>
   for(int i = 1; i < N; ++i) for(int j = 0; j < M; ++j)</pre>
        dp[i][j] = dp[i - 1][j] - dp[i - 1][j / primes[i]
         - 111:
11 phi(ll n, int a){
  if(!a) return n;
   if(n < M && a < N) return dp[a][n];</pre>
   if(primes[a - 1] > n) return 1;
   if(((11)primes[a - 1]) * primes[a - 1] >= n && n < V)
        return small_pi[n] - a + 1;
   11 de = phi(n, a - 1) - phi(n / primes[a - 1], a - 1)
   return de;
}
```

```
1l PiCount(1l n){
    if(n < V) return small_pi[n];
    int s = sqrt(n + 0.5), y = cbrt(n + 0.5), a =
        small_pi[y];
    ll res = phi(n, a) + a - 1;
    for(; primes[a] <= s; ++a) res -= max(PiCount(n /
        primes[a]) - PiCount(primes[a]) + 1, 0ll);
    return res;
}</pre>
```

#### 6.5 Linear Function Mod Min

```
11 topos(11 x, 11 m) {x %= m; if (x < 0) x += m; return
//min value of ax + b \pmod{m} for x \in [0, n - 1]. O(
    Log m)
11 min_rem(ll n, ll m, ll a, ll b) {
  a = topos(a, m), b = topos(b, m);
  for (ll g = __gcd(a, m); g > 1;) return g * min_rem(n
        m / g, a / g, b / g) + (b % g);
  for (11 nn, nm, na, nb; a; n = nn, m = nm, a = na, b
      = nb) {
    if (a <= m - a) {
    nn = (a * (n - 1) + b) / m;</pre>
      if (!nn) break;
      nn += (b < a);
      nm = a, na = topos(-m, a);
      nb = b < a ? b : topos(b - m, a);
    } else {
      ll lst = b - (n - 1) * (m - a);
      if (lst >= 0) {b = lst; break;}
      nn = -(1st / m) + (1st % m < -a) + 1;
      nm = m - a, na = m % (m - a), nb = b % (m - a);
    }
  }
  return b;
//min value of ax + b \pmod{m} for x \in [0, n - 1],
    also return \min x to get the value. O(\log m)
//{value, x}
pair<ll, ll> min_rem_pos(ll n, ll m, ll a, ll b) {
  a = topos(a, m), b = topos(b, m);
  11 mn = min_rem(n, m, a, b), g = __gcd(a, m);
  //ax = (mn - b) \pmod{m}
  11 x = (extgcd(a, m).first + m) * ((mn - b + m) / g)
      % (m / g);
  return {mn, x};
}
```

#### 6.6 Determinant

```
11 Det(vector <vector <11>>> a) {
  int n = a.size();
  ll det = 1;
  for (int i = 0; i < n; ++i) {
    if (!a[i][i]) {
      det = -det;
      if (det < 0) det += mod;</pre>
      for (int j = i + 1; j < n; ++j) if (a[j][i]) {</pre>
        swap(a[j], a[i]);
        break;
      if (!a[i][i]) return 0;
    det = det * a[i][i] % mod;
    ll mul = mpow(a[i][i], mod - 2);
    for (int j = 0; j < n; ++j) a[i][j] = a[i][j] * mul</pre>
          % mod;
    for (int j = 0; j < n; ++j) if (i ^ j) {</pre>
      11 mul = a[j][i];
      for (int k = 0; k < n; ++k) {
        a[j][k] -= a[i][k] * mul % mod;
        if (a[j][k] < 0) a[j][k] += mod;</pre>
      }
    }
  return det;
```

# 6.7 Floor Sum

## 6.8 Quadratic Residue

```
int Jacobi(int a, int m) {
  int s = 1;
  for (; m > 1; ) {
    a %= m;
    if (a == 0) return 0;
    const int r = __builtin_ctz(a);
    if ((r \& 1) \& \& ((m + 2) \& 4)) s = -s;
    a >>= r;
    if (a \& m \& 2) s = -s;
    swap(a, m);
  }
  return s;
int QuadraticResidue(int a, int p) {
  if (p == 2) return a & 1;
  const int jc = Jacobi(a, p);
  if (jc == 0) return 0;
  if (jc == -1) return -1;
  int b, d;
  for (; ; ) {
    b = rand() % p;
d = (1LL * b * b + p - a) % p;
    if (Jacobi(d, p) == -1) break;
  int f0 = b, f1 = 1, g0 = 1, g1 = 0, tmp;
  for (int e = (1LL + p) >> 1; e; e >>= 1) {
    if (e & 1) {
      tmp = (1LL * g0 * f0 + 1LL * d * (1LL * g1 * f1 %
           p)) % p;
      g1 = (1LL * g0 * f1 + 1LL * g1 * f0) % p;
      g0 = tmp;
    tmp = (1LL * f0 * f0 + 1LL * d * (1LL * f1 * f1 % p
        )) % p;
    f1 = (2LL * f0 * f1) % p;
    f0 = tmp;
  }
  return g0;
```

## 6.9 Simplex

```
struct Simplex { // O-based
  using T = long double;
  static const int N = 410, M = 30010;
  const T eps = 1e-7;
  int n, m;
  int Left[M], Down[N];
  // Ax <= b, max c^T x
  // result : v, xi = sol[i]
  T a[M][N], b[M], c[N], v, sol[N];
  bool eq(T a, T b) {return fabs(a - b) < eps;}</pre>
  bool ls(T a, T b) {return a < b && !eq(a, b);}</pre>
  void init(int _n, int _m) {
  n = _n, m = _m, v = 0;
    for (int i = 0; i < m; ++i) for (int j = 0; j < n;
         ++j) a[i][j] = 0;
    for (int i = 0; i < m; ++i) b[i] = 0;</pre>
    for (int i = 0; i < n; ++i) c[i] = sol[i] = 0;</pre>
  void pivot(int x, int y) {
    swap(Left[x], Down[y]);
    T k = a[x][y]; a[x][y] = 1;
    vector <int> nz;
    for (int i = 0; i < n; ++i) {</pre>
      a[x][i] /= k;
      if (!eq(a[x][i], 0)) nz.push_back(i);
```

```
b[x] /= k;
    for (int i = 0; i < m; ++i) {</pre>
      if (i == x || eq(a[i][y], 0)) continue;
      k = a[i][y], a[i][y] = 0;
b[i] -= k * b[x];
      for (int j : nz) a[i][j] -= k * a[x][j];
    if (eq(c[y], 0)) return;
    k = c[y], c[y] = 0, v += k * b[x];
    for (int i : nz) c[i] -= k * a[x][i];
  }
  // 0: found solution, 1: no feasible solution, 2:
       unbounded
  int solve() {
    for (int i = 0; i < n; ++i) Down[i] = i;</pre>
    for (int i = 0; i < m; ++i) Left[i] = n + i;</pre>
    while (1) {
      int x = -1, y = -1;
       for (int i = 0; i < m; ++i) if (ls(b[i], 0) && (x</pre>
            == -1 \mid \mid b[i] < b[x])) x = i;
       if (x == -1) break;
      for (int i = 0; i < n; ++i) if (ls(a[x][i], 0) &&</pre>
             (y == -1 \mid | a[x][i] < a[x][y])) y = i;
       if (y == -1) return 1;
      pivot(x, y);
    while (1) {
       int x = -1, y = -1;
       for (int i = 0; i < n; ++i) if (ls(0, c[i]) && (y
            == -1 || c[i] > c[y])) y = i;
       if (y == -1) break;
      for (int i = 0; i < m; ++i) if (ls(0, a[i][y]) &&
    (x == -1 || b[i] / a[i][y] < b[x] / a[x][y])</pre>
           ) x = i;
      if (x == -1) return 2;
      pivot(x, y);
    for (int i = 0; i < m; ++i) if (Left[i] < n) sol[</pre>
         Left[i]] = b[i];
    return 0;
  }
};
```

#### 6.10 Berlekamp Massey

```
vector <11> BerlekampMassey(vector <11> a) {
 // find min |c| such that a_n = sum c_j * a_{n - j - j}
      1}, 0-based
  // O(N^2), if |c| = k, |a| >= 2k sure correct
  auto f = [&](vector<11> v, 11 c) {
    for (11 &x : v) x = mul(x, c);
 };
 vector <11> c, best;
  int pos = 0, n = a.size();
  for (int i = 0; i < n; ++i) {
    ll error = a[i];
    for (int j = 0; j < c.size(); ++j) error = sub(
        error, mul(c[j], a[i - 1 - j]));
    if (error == 0) continue;
    11 inv = mpow(error, mod - 2);
    if (c.empty()) {
      c.resize(i + 1);
      pos = i;
      best.pb(inv);
    } else {
      vector <1l> fix = f(best, error);
      fix.insert(fix.begin(), i - pos - 1, 0);
      if (fix.size() >= c.size()) {
        best = f(c, sub(0, inv));
        best.insert(best.begin(), inv);
        pos = i:
        c.resize(fix.size());
      for (int j = 0; j < fix.size(); ++j) c[j] = add(c</pre>
          [j], fix[j]);
   }
 }
  return c;
```

# 6.11 Linear Programming Construction

Standard form: maximize  $\mathbf{c}^T\mathbf{x}$  subject to  $A\mathbf{x} \leq \mathbf{b}$  and  $\mathbf{x} \geq 0$ . Dual LP: minimize  $\mathbf{b}^T\mathbf{y}$  subject to  $A^T\mathbf{y} \geq \mathbf{c}$  and  $\mathbf{y} \geq 0$ .  $\bar{\mathbf{x}}$  and  $\bar{\mathbf{y}}$  are optimal if and only if for all  $i \in [1,n]$ , either  $\bar{x}_i = 0$  or  $\sum_{j=1}^m A_{ji}\bar{y}_j = c_i$  holds and for all  $i \in [1,m]$  either  $\bar{y}_i = 0$  or  $\sum_{j=1}^n A_{ij}\bar{x}_j = b_j$  holds.

- 1. In case of minimization, let  $c_i^\prime = -c_i$
- 2.  $\sum_{1 \leq i \leq n} A_{ji} x_i \geq b_j \rightarrow \sum_{1 \leq i \leq n} -A_{ji} x_i \leq -b_j$
- 3.  $\sum_{1 \le i \le n}^{-} A_{ji} x_i = b_j$ 
  - $\begin{array}{ll} \bullet & \sum_{1 \leq i \leq n} A_{ji} x_i \leq b_j \\ \bullet & \sum_{1 \leq i \leq n} A_{ji} x_i \geq b_j \end{array}$
- 4. If  $x_i$  has no lower bound, replace  $x_i$  with  $x_i x_i'$

#### 6.12 Euclidean

- $m = \lfloor \frac{an+b}{c} \rfloor$
- Time complexity:  $O(\log n)$

$$\begin{split} f(a,b,c,n) &= \sum_{i=0}^n \lfloor \frac{ai+b}{c} \rfloor \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)}{2} + \lfloor \frac{b}{c} \rfloor \cdot (n+1) \\ + f(a \mod c, b \mod c, c, n), & a \geq c \vee b \geq c \\ 0, & n < 0 \vee a = 0 \\ nm - f(c, c - b - 1, a, m - 1), & \text{otherwise} \end{cases} \end{split}$$

$$\begin{split} g(a,b,c,n) &= \sum_{i=0}^n i \lfloor \frac{ai+b}{c} \rfloor \\ &= \begin{cases} \left\lfloor \frac{a}{c} \right\rfloor \cdot \frac{n(n+1)(2n+1)}{6} + \left\lfloor \frac{b}{c} \right\rfloor \cdot \frac{n(n+1)}{2} \\ +g(a \mod c, b \mod c, c, n), & a \geq c \lor b \geq c \\ 0, & n < 0 \lor a = 0 \\ \frac{1}{2} \cdot (n(n+1)m - f(c, c-b-1, a, m-1) \\ -h(c, c-b-1, a, m-1)), & \text{otherwise} \end{cases} \end{split}$$

$$\begin{split} h(a,b,c,n) &= \sum_{i=0}^n \lfloor \frac{ai+b}{c} \rfloor^2 \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor^2 \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor^2 \cdot (n+1) \\ &+ \lfloor \frac{a}{c} \rfloor \cdot \lfloor \frac{b}{c} \rfloor \cdot n(n+1) \\ &+ h(a \bmod c, b \bmod c, c, n) \\ &+ 2 \lfloor \frac{a}{c} \rfloor \cdot g(a \bmod c, b \bmod c, c, n) \\ &+ 2 \lfloor \frac{b}{c} \rfloor \cdot f(a \bmod c, b \bmod c, c, n), & a \geq c \lor b \geq c \\ 0, & n < 0 \lor a = 0 \\ nm(m+1) - 2g(c, c-b-1, a, m-1) \\ &- 2f(c, c-b-1, a, m-1) - f(a, b, c, n), & \text{otherwise} \end{cases} \end{split}$$

#### 6.13 Theorem

• Kirchhoff's Theorem

Denote L be a  $n\times n$  matrix as the Laplacian matrix of graph G, where  $L_{ii}=d(i)$ ,  $L_{ij}=-c$  where c is the number of edge (i,j) in G.

- The number of undirected spanning in G is  $|\det(\tilde{L}_{11})|$  .
- The number of directed spanning tree rooted at r in G is  $|\det(\tilde{L}_{rr})|$ .
- Tutte's Matrix

Let D be a  $n \times n$  matrix, where  $d_{ij} = x_{ij}$  ( $x_{ij}$  is chosen uniformly at random) if i < j and  $(i,j) \in E$ , otherwise  $d_{ij} = -d_{ji}$ .  $\frac{rank(D)}{2}$  is the maximum matching on G.

- Cayley's Formula
  - Given a degree sequence  $d_1,d_2,\ldots,d_n$  for each  $\mathit{labeled}$  vertices, there are

$$\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$$

spanning trees.

- Let  $T_{n,k}$  be the number of *labeled* forests on n vertices with k components, such that vertex  $1,2,\ldots,k$  belong to different components. Then  $T_{n,k}=kn^{n-k-1}$ .
- Erdős-Gallai Theorem

A sequence of non-negative integers  $d_1 \geq d_2 \geq \ldots \geq d_n$  can be represented as the degree sequence of a finite simple graph on n vertices if and only if  $d_1+d_2+\ldots+d_n$  is even and

$$\sum_{i=1}^k d_i \le k(k-1) + \sum_{i=k+1}^n \min(d_i, k)$$

holds for all  $1 \leq k \leq n$ .

• Burnside's Lemma

Let X be a set and G be a group that acts on X. For  $g \in G$ , denote by  $X^g$  the elements fixed by g:

$$X^g = \{ x \in X \mid gx \in X \}$$

Then

$$|X/G| = \frac{1}{|G|} \sum_{g \in G} |X^g|.$$

• Gale-Ryser theorem

A pair of sequences of nonnegative integers  $a_1 \geq \cdots \geq a_n$  and  $b_1,\ldots,b_n$  is bigraphic if and only if  $\sum_{i=1}^n a_i = \sum_{i=1}^n b_i$  and  $\sum_{i=1}^k a_i \leq$ 

$$\sum_{i=1}^n \min(b_i,k)$$
 holds for every  $1 \leq k \leq n$  .

• Fulkerson-Chen-Anstee theorem

A sequence  $(a_1,b_1),\ldots,(a_n,b_n)$  of nonnegative integer pairs with A sequence  $(a_1, a_1), \dots, (a_n, a_n)$   $a_1 \ge \dots \ge a_n$  is digraphic if and only if  $\sum_{i=1}^n a_i = \sum_{i=1}^n b_i$  and

$$\sum_{i=1}^k a_i \leq \sum_{i=1}^k \min(b_i,k-1) + \sum_{i=k+1}^n \min(b_i,k) \text{ holds for every } 1 \leq k \leq n.$$

• Möbius inversion formula

- 
$$f(n) = \sum_{d|n} g(d) \Leftrightarrow g(n) = \sum_{d|n} \mu(d) f(\frac{n}{d})$$
  
-  $f(n) = \sum_{n|d} g(d) \Leftrightarrow g(n) = \sum_{n|d} \mu(\frac{d}{n}) f(d)$ 

- Spherical cap
  - A portion of a sphere cut off by a plane. r: sphere radius, a: radius of the base of the cap, h: height of the cap,  $\theta$ :  $\arcsin(a/r)$ .  $1/2 = \pi h^2(3r-h)/3 = \pi h(3a^2+h^2)/6 = \pi r^3(2+\cos\theta)(1-h^2)/6$

  - $\cos\theta)^2/3. \\ \text{ Area} = 2\pi r h = \pi (a^2 + h^2) = 2\pi r^2 (1 \cos\theta).$
- Chinese Remainder Theorem
  - $x \equiv a_i \pmod{m_i}$
  - $M = \prod m_i, M_i = M/m_i$
  - $t_i M_i \equiv 1 \pmod{m_i}$
  - $x = \sum a_i t_i M_i \pmod{M}$

### 6.14 Estimation

- The number of divisors of n is at most around  $100\ {\rm for}\ n<5e4$  ,
- 500 for n<1e7, 2000 for n<1e10, 200000 for n<1e19. The number of ways of writing n as a sum of positive integers, disregarding the order of the summands. 1,1,2,3,5,7,11,15,22,30for  $n=0\sim 9$ , 627 for n=20,  $\sim 2e5$  for n=50,  $\sim 2e8$  for
- Total number of partitions of n distinct elements: B(n)=1,1,2,5,15,52,203,877,4140,21147,115975,678570,4213597,27644437, 190899322, . . . .

## 6.15 General Purpose Numbers

Bernoulli numbers

$$\begin{split} B_0 - 1, B_1^{\pm} &= \pm \tfrac{1}{2}, B_2 = \tfrac{1}{6}, B_3 = 0 \\ \sum_{j=0}^m \binom{m+1}{j} B_j &= 0 \text{, EGF is } B(x) = \tfrac{x}{e^x-1} = \sum_{n=0}^\infty B_n \frac{x^n}{n!} \,. \end{split}$$

$$S_m(n) = \sum_{k=1}^n k^m = \frac{1}{m+1} \sum_{k=0}^m {m+1 \choose k} B_k^+ n^{m+1-k}$$

ullet Stirling numbers of the second kind Partitions of n distinct elements into exactly k groups.

$$\begin{split} S(n,k) &= S(n-1,k-1) + kS(n-1,k), S(n,1) = S(n,n) = 1 \\ S(n,k) &= \frac{1}{k!} \sum_{i=0}^k (-1)^{k-i} {k \choose i} i^n \\ x^n &= \sum_{i=0}^n S(n,i)(x)_i \end{split}$$

• Pentagonal number theorem

$$\prod_{n=1}^{\infty} (1-x^n) = 1 + \sum_{k=1}^{\infty} (-1)^k \left( x^{k(3k+1)/2} + x^{k(3k-1)/2} \right)$$

• Catalan numbers

$$C_n^{(k)} = \frac{1}{(k-1)n+1} {kn \choose n}$$
 
$$C^{(k)}(x) = 1 + x[C^{(k)}(x)]^k$$

• Eulerian numbers

Number of permutations  $\pi \in S_n$  in which exactly k elements are greater than the previous element. k j:s s.t.  $\pi(j) > \pi(j+1)$ , k+1 j:s s.t.  $\pi(j) \geq j$ , k j:s s.t.  $\pi(j) > j$ .

$$E(n,k) = (n-k)E(n-1,k-1) + (k+1)E(n-1,k)$$

$$E(n,0) = E(n,n-1) = 1$$

$$E(n,k) = \sum_{j=0}^{k} (-1)^{j} {n+1 \choose j} (k+1-j)^{n}$$

## 6.16 Tips for Generating Funtion

• Ordinary Generating Function  $A(x) = \sum_{i \geq 0} a_i x^i$ 

```
- A(rx) \Rightarrow r^n a_n
- A(x) + B(x) \Rightarrow a_n + b_n

- A(x)B(x) \Rightarrow \sum_{i=0}^{n} a_i b_{n-i}
- A(x)^k \Rightarrow \sum_{i_1+i_2+\dots+i_k=n} a_{i_1} a_{i_2} \dots a_{i_k}
- xA(x)' \Rightarrow na_n
-\frac{A(x)}{1-x} \Rightarrow \sum_{i=0}^{n} a_i
```

• Exponential Generating Function  $A(x) = \sum_{i \ge 0} \frac{a_i}{i!} x_i$ 

```
- A(x) + B(x) \Rightarrow a_n + b_n

- A^{(k)}(x) \Rightarrow a_{n+k}

- A(x)B(x) \Rightarrow \sum_{i=0}^{n} \binom{n}{i} a_i b_{n-i}
 - A(x)^k \Rightarrow \sum_{i_1+i_2+\dots+i_k=n} \binom{n}{i_1,i_2,\dots,i_k} a_{i_1} a_{i_2} \dots a_{i_k}
```

Special Generating Function

```
-(1+x)^n = \sum_{i>0} \binom{n}{i} x^i
-\frac{1}{(1-x)^n} = \sum_{i>0} {i \choose n-1} x
```

# **Polynomial**

## Number Theoretic Transform

```
// mul, add, sub, mpow
// ll -> int if too slow
struct NTT {
  11 w[N];
  NTT() {
     ll dw = mpow(G, (mod - 1) / N);
     for (int i = 1; i < N; ++i) w[i] = mul(w[i - 1], dw</pre>
  void operator()(vector<ll>& a, bool inv = false) { //
       0 \leftarrow a[i] \leftarrow P
     int x = 0, n = a.size();
     for (int j = 1; j < n - 1; ++j) {
  for (int k = n >> 1; (x ^= k) < k; k >>= 1);
       if (j < x) swap(a[x], a[j]);</pre>
     for (int L = 2; L <= n; L <<= 1) {
       int dx = N / L, dl = L >> 1;
for (int i = 0; i < n; i += L) {</pre>
          for (int j = i, x = 0; j < i + dl; ++j, x += dx
            ll tmp = mul(a[j + dl], w[x]);
            a[j + dl] = sub(a[j], tmp);
            a[j] = add(a[j], tmp);
       }
     if (inv) {
       reverse(a.begin() + 1, a.end());
       11 invn = mpow(n, mod - 2);
for (int i = 0; i < n; ++i) a[i] = mul(a[i], invn</pre>
  }
} ntt;
```

# 7.2 Primes

Prime	Root	Prime	Root
7681	17	167772161	3
12289	11	104857601	3
40961	3	985661441	3
65537	3	998244353	3
786433	10	1107296257	10
5767169	3	2013265921	31
7340033	3	2810183681	11
23068673	3	2885681153	3
469762049	3	605028353	3
2061584302081	2748779069441	3	
1945555039024054273	5	9223372036737335297	3

## 7.3 Polynomial Operations

```
vector <ll> Mul(vector <ll> a, vector <ll> b, int bound
     = N) {
  int m = a.size() + b.size() - 1, n = 1;
  while (n < m) n <<= 1;</pre>
  a.resize(n), b.resize(n);
 ntt(a), ntt(b);
 vector <11> out(n);
  for (int i = 0; i < n; ++i) out[i] = mul(a[i], b[i]);</pre>
 ntt(out, true), out.resize(min(m, bound));
 return out;
vector <ll> Inverse(vector <ll> a) {
 // O(NlogN), a[0] != 0
 int n = a.size();
  vector <11> res(1, mpow(a[0], mod - 2));
  for (int m = 1; m < n; m <<= 1) {</pre>
    if (n < m * 2) a.resize(m * 2);</pre>
    vector <ll> v1(a.begin(), a.begin() + m * 2), v2 =
        res;
    v1.resize(m * 4), v2.resize(m * 4);
    ntt(v1), ntt(v2);
    for (int i = 0; i < m * 4; ++i) v1[i] = mul(mul(v1[</pre>
         i], v2[i]), v2[i]);
    ntt(v1, true);
    res.resize(m * 2);
    for (int i = 0; i < m; ++i) res[i] = add(res[i],</pre>
         res[i]);
    for (int i = 0; i < m * 2; ++i) res[i] = sub(res[i</pre>
         ], v1[i]);
 res.resize(n):
 return res;
pair <vector <ll>, vector <ll>> Divide(vector <ll> a,
    vector <ll> b) {
  // a = bQ + R, O(NLogN), b.back() != 0
 int n = a.size(), m = b.size(), k = n - m + 1;
  if (n < m) return {{0}, a};</pre>
 vector <11> ra = a, rb = b;
 reverse(all(ra)), ra.resize(k);
 reverse(all(rb)), rb.resize(k);
 vector <11> Q = Mul(ra, Inverse(rb), k);
 reverse(all(Q));
 vector <1l> res = Mul(b, Q), R(m - 1);
for (int i = 0; i < m - 1; ++i) R[i] = sub(a[i], res[</pre>
      i]);
 return {Q, R};
vector <1l> SqrtImpl(vector <1l> a) {
  if (a.empty()) return {0};
  int z = QuadraticResidue(a[0], mod), n = a.size();
  if (z == -1) return {-1};
  vector \langle 11 \rangle q(1, z);
  const int inv2 = (mod + 1) / 2;
  for (int m = 1; m < n; m <<= 1) {</pre>
    if (n < m * 2) a.resize(m * 2);</pre>
    q.resize(m * 2);
    vector <11> f2 = Mul(q, q, m * 2);
for (int i = 0; i < m * 2; ++i) f2[i] = sub(f2[i],</pre>
        a[i]);
    f2 = Mul(f2, Inverse(q), m * 2);
    for (int i = 0; i < m * 2; ++i) q[i] = sub(q[i],</pre>
        mul(f2[i], inv2));
  q.resize(n);
  return q;
vector <11> Sqrt(vector <11> a) {
  // O(NlogN), return {-1} if not exists
  int n = a.size(), m = 0;
 while (m < n && a[m] == 0) m++;</pre>
  if (m == n) return vector <11>(n);
  if (m & 1) return {-1};
  vector <ll> s = SqrtImpl(vector <ll>(a.begin() + m, a
      .end()));
  if (s[0] == -1) return {-1};
  vector <11> res(n);
  for (int i = 0; i < s.size(); ++i) res[i + m / 2] = s</pre>
      [i];
  return res;
```

```
vector <11> Derivative(vector <11> a) {
  int n = a.size();
  vector <ll> res(n - 1);
  for (int i = 0; i < n - 1; ++i) res[i] = mul(a[i +</pre>
      1], i + 1);
  return res;
vector <ll> Integral(vector <ll> a) {
  int n = a.size();
  vector \langle 11 \rangle res(n + 1);
  for (int i = 0; i < n; ++i) {
    res[i + 1] = mul(a[i], mpow(i + 1, mod - 2));
  return res;
}
vector <ll> Ln(vector <ll> a) {
  // O(NlogN), a[0] = 1
  int n = a.size();
  if (n == 1) return {0};
  vector <1l> d = Derivative(a);
  a.pop_back();
  return Integral(Mul(d, Inverse(a), n - 1));
vector <11> Exp(vector <11> a) {
  // O(NlogN), a[0] = 0
  int n = a.size();
  vector \langle 11 \rangle q(1, 1);
  a[0] = add(a[0], 1);
  for (int m = 1; m < n; m <<= 1) {
   if (n < m * 2) a.resize(m * 2);</pre>
    vector <ll> g(a.begin(), a.begin() + m * 2), h(all(
         a));
    h.resize(m * 2), h = Ln(h);
    for (int i = 0; i < m * 2; ++i) {</pre>
      g[i] = sub(g[i], h[i]);
    q = Mul(g, q, m * 2);
  }
  q.resize(n);
  return a:
vector <ll> Pow(vector <ll> a, ll k) {
  int n = a.size(), m = 0;
  vector <11> ans(n, 0);
  while (m < n \&\& a[m] == 0) m++;
  if (k \&\& m \&\& (k >= n | | k * m >= n)) return ans;
  if (m == n) return ans[0] = 1, ans;
  ll lead = m * k;
  vector <1l> b(a.begin() + m, a.end());
  11 base = mpow(b[0], k), inv = mpow(b[0], mod - 2);
  for (int i = 0; i < n - m; ++i) b[i] = mul(b[i], inv)</pre>
  b = Ln(b);
  for (int i = 0; i < n - m; ++i) b[i] = mul(b[i], k %</pre>
      mod);
  b = Exp(b);
  for (int i = lead; i < n; ++i) ans[i] = mul(b[i -</pre>
      lead], base);
  return ans;
vector <ll> Evaluate(vector <ll> a, vector <ll> x) {
  if (x.empty()) return {};
  int n = x.size();
  vector <vector <11>> up(n * 2);
  for (int i = 0; i < n; ++i) up[i + n] = {sub(0, x[i])}
      , 1};
  for (int i = n - 1; i > 0; --i) up[i] = Mul(up[i *
      2], up[i * 2 + 1]);
  vector <vector <11>> down(n * 2);
  down[1] = Divide(a, up[1]).second;
  for (int i = 2; i < n * 2; ++i) down[i] = Divide(down</pre>
      [i >> 1], up[i]).second;
  vector <11> y(n);
  for (int i = 0; i < n; ++i) y[i] = down[i + n][0];</pre>
  return v;
vector <1l> Interpolate(vector <1l> x, vector <1l> y) {
  int n = x.size();
  vector <vector <11>> up(n * 2);
  for (int i = 0; i < n; ++i) up[i + n] = {sub(0, x[i])}
      , 1};
```

#### 7.4 Fast Linear Recursion

```
11 FastLinearRecursion(vector <11> a, vector <11> c, 11
      k) {
  // a_n = sigma c_j * a_{n - j - 1}, 0-based
// O(NLogNLogK), |a| = |c|
  int n = a.size();
  if (k < n) return a[k];</pre>
  vector <ll> base(n + 1, 1);
  for (int i = 0; i < n; ++i) base[i] = sub(0, c[n - i</pre>
       - 1]);
  vector <11> poly(n);
  (n == 1 ? poly[0] = c[n - 1] : poly[1] = 1);
  auto calc = [&](vector <ll> p1, vector <ll> p2) {
    // O(n^2) bruteforce or O(nlogn) NTT
    return Divide(Mul(p1, p2), base).second;
  };
  vector \langle 11 \rangle res(n, 0); res[0] = 1;
  for (; k; k >>= 1, poly = calc(poly, poly)) {
    if (k & 1) res = calc(res, poly);
  11 \text{ ans} = 0;
  for (int i = 0; i < n; ++i) {</pre>
    (ans += res[i] * a[i]) %= mod;
  return ans;
```

#### 7.5 Fast Walsh Transform

```
void fwt(vector <int> &a) {
 // \ and : a[j] += x;
  //
         : a[j] -= x;
        : a[j ^ (1 << i)] += y;
         : a[j ^ (1 << i)] -= y;
 //
  // xor : a[j] = x - y, a[j ^ (1 << i)] = x + y;
          : a[j] = (x - y) / 2, a[j ^ (1 << i)] = (x + y)
      ) / 2;
 int n = __lg(a.size());
  for (int i = 0; i < n; ++i) {</pre>
    for (int j = 0; j < 1 << n; ++j) if (j >> i & 1) {
  int x = a[j ^ (1 << i)], y = a[j];</pre>
      // do something
 }
vector<int> subs_conv(vector<int> a, vector<int> b) {
 // c_i = sum_{j \& k = 0, j | k = i} a_j * b_k
 int n = __lg(a.size());
  vector<vector<int>> ha(n + 1, vector<int>(1 << n));</pre>
  vector<vector<int>> hb(n + 1, vector<int>(1 << n));</pre>
  vector<vector<int>> c(n + 1, vector<int>(1 << n));</pre>
  for (int i = 0; i < 1 << n; ++i) {</pre>
    ha[__builtin_popcount(i)][i] = a[i];
    hb[__builtin_popcount(i)][i] = b[i];
  for (int i = 0; i <= n; ++i) fwt(ha[i]),fwt(hb[i]);</pre>
  for (int i = 0; i <= n; ++i)</pre>
    for (int j = 0; i + j <= n; ++j)</pre>
      for (int k = 0; k < 1 << n; ++k)
        // mind overflow
        c[i + j][k] += ha[i][k] * hb[j][k];
  for (int i = 0; i <= n; ++i)</pre>
```

```
fwt(c[i], true);
vector <int> ans(1 << n);
for (int i = 0; i < 1 << n; ++i)
    ans[i] = c[__builtin_popcount(i)][i];
return ans;
}</pre>
```

# 8 Geometry

#### 8.1 Basic

```
const double eps = 1e-8, pi = acos(-1);
int sign(double x) \{return abs(x) \leftarrow eps ? 0 : (x > 0 ?
     1 : -1);}
struct Pt {
  double x, y;
  Pt (double _x, double _y) : x(_x), y(_y) {}
  Pt operator + (Pt o) {return Pt(x + o.x, y + o.y);}
  Pt operator - (Pt o) {return Pt(x - o.x, y - o.y);}
  Pt operator * (double k) {return Pt(x * k, y * k);}
  Pt operator / (double k) {return Pt (x / k, y / k);}
double operator * (Pt o) {return x * o.x + y * o.y;}
  double operator ^ (Pt o) {return x * o.y - y * o.x;}
};
struct Line {
  Pt a, b;
struct Cir {
  Pt o; double r;
double abs2(Pt o) {return o.x * o.x + o.y * o.y;}
double abs(Pt o) {return sqrt(abs2(o));}
int ori(Pt o, Pt a, Pt b) {return sign((o - a) ^ (o - b
    ));}
bool btw(Pt a, Pt b, Pt c) { // c on segment ab?
  return ori(a, b, c) == 0 && sign((c - a) * (c - b))
double area(Pt a, Pt b, Pt c) {return abs((a - b) ^ (a
    - c)) / 2;}
Pt unit(Pt o) {return o / abs(o);}
Pt rot(Pt a, double o) { // CCW
  double c = cos(o), s = sin(o);
  return Pt(c * a.x - s * a.y, s * a.x + c * a.y);
Pt proj_vector(Pt a, Pt b, Pt c) { // vector ac proj to
  return (b - a) * ((c - a) * (b - a)) / ((b - a) * (b
      - a));
Pt proj_pt(Pt a, Pt b, Pt c) { // point c proj to ab
  return proj_vector(a, b, c) + a;
```

## 8.2 Heart

```
Pt circenter(Pt p0, Pt p1, Pt p2) { // radius = abs(
   center)
  p1 = p1 - p0, p2 = p2 - p0;
  double x1 = p1.x, y1 = p1.y, x2 = p2.x, y2 = p2.y;
  double m = 2. * (x1 * y2 - y1 * x2);
  Pt center(0, 0);
  center.x = (x1 * x1 * y2 - x2 * x2 * y1 + y1 * y2 * (
      y1 - y2)) / m;
  center.y = (x1 * x2 * (x2 - x1) - y1 * y1 * x2 + x1 *
       y2 * y2) / m;
  return center + p0;
Pt incenter(Pt p1, Pt p2, Pt p3) { // radius = area / s
  double a = abs(p2 - p3), b = abs(p1 - p3), c = abs(p1
       - p2);
  double s = a + b + c;
  return (p1 * a + p2 * b + p3 * c) / s;
Pt masscenter(Pt p1, Pt p2, Pt p3)
{ return (p1 + p2 + p3) / 3; }
Pt orthocenter(Pt p1, Pt p2, Pt p3)
{ return masscenter(p1, p2, p3) * 3 - circenter(p1, p2,
     p3) * 2; }
```

#### 8.3 External Bisector

```
Pt external_bisector(Pt p1, Pt p2, Pt p3) { //213
Pt L1 = p2 - p1, L2 = p3 - p1;
L2 = L2 * abs(L1) / abs(L2);
return L1 + L2;
}
```

## 8.4 Intersection of Segments

```
Pt LinesInter(Line a, Line b) {
    double abc = (a.b - a.a) ^ (b.a - a.a);
    double abd = (a.b - a.a) ^ (b.b - a.a);
    if (sign(abc - abd) == 0) return b.b;// no inter
    return (b.b * abc - b.a * abd) / (abc - abd);
}

vector<Pt> SegsInter(Line a, Line b) {
    if (btw(a.a, a.b, b.a)) return {b.a};
    if (btw(a.a, a.b, b.b)) return {b.b};
    if (btw(b.a, b.b, a.a)) return {a.a};
    if (btw(b.a, b.b, a.b)) return {a.b};
    if (ori(a.a, a.b, b.a) * ori(a.a, a.b, b.b) == -1 &&
        ori(b.a, b.b, a.a) * ori(b.a, b.b, a.b) == -1)
    return {LinesInter(a, b)};
    return {};
}
```

#### 8.5 Intersection of Circle and Line

#### 8.6 Intersection of Circles

## 8.7 Intersection of Polygon and Circle

```
double _area(Pt pa, Pt pb, double r){
  if(abs(pa) < abs(pb)) swap(pa, pb);</pre>
  if(abs(pb) < eps) return 0;</pre>
  double S, h, theta;
  double a = abs(pb), b = abs(pa), c = abs(pb - pa);
  double cosB = pb * (pb - pa) / a / c, B = acos(cosB);
double cosC = (pa * pb) / a / b, C = acos(cosC);
  if (a > r) {
    S = (C / 2) * r * r;

h = a * b * sin(C) / c;
    if (h < r && B < pi / 2) S -= (acos(h / r) * r * r</pre>
         - h * sqrt(r * r - h * h));
  } else if (b > r) {
    theta = pi - B - asin(sin(B) / r * a);
    S = .5 * a * r * sin(theta) + (C - theta) / 2 * r *
  } else
    S = .5 * sin(C) * a * b;
  return S;
double area_poly_circle(vector<Pt> poly, Pt 0, double r
    ) {
  double S = 0; int n = poly.size();
```

```
for(int i = 0; i < n; ++i)
   S += _area(poly[i] - 0, poly[(i + 1) % n] - 0, r) *
        ori(0, poly[i], poly[(i + 1) % n]);
return fabs(S);
}</pre>
```

## 8.8 Tangent Lines of Circle and Point

## 8.9 Tangent Lines of Circles

```
vector<Line> tangent(Cir a, Cir b) {
#define Pij \
  Pt i = unit(b.o - a.o) * a.r, j = Pt(i.y, -i.x);\
  z.push_back({a.o + i, a.o + i + j});
#define deo(I,J) \
  double d = abs(a.o - b.o), e = a.r I b.r, o = acos(e
      / d);\
  Pt i = unit(b.o - a.o), j = rot(i, o), k = rot(i, -o)
  z.push_back({a.o + j * a.r, b.o J j * b.r});\
z.push_back({a.o + k * a.r, b.o J k * b.r});
  if (a.r < b.r) swap(a, b);</pre>
  vector<Line> z;
  if (abs(a.o - b.o) + b.r < a.r) return z;</pre>
  else if (sign(abs(a.o - b.o) + b.r - a.r) == 0) { Pij
      ; }
  else {
    deo(-,+); // inter
    // outer
    if (sign(d - a.r - b.r) == 0) { Pij; }
    else if (d > a.r + b.r) { deo(+,-); }
  return z;
```

#### 8.10 Point In Convex

## 8.11 Point Segment Distance

```
double PointSegDist(Pt q0, Pt q1, Pt p) {
   if (sign(abs(q0 - q1)) == 0) return abs(q0 - p);
   if (sign((q1 - q0) * (p - q0)) >= 0 && sign((q0 - q1)
        * (p - q1)) >= 0)
      return fabs(((q1 - q0) ^ (p - q0)) / abs(q0 - q1));
   return min(abs(p - q0), abs(p - q1));
}
```

## 8.12 Convex Hull

## 8.13 Convex Hull Distance

# 8.14 Minimum Enclosing Circle

```
Cir min_enclosing(vector<Pt> &p) {
  random_shuffle(p.begin(), p.end());
  double r = 0.0;
  Pt cent = p[0];
  for (int i = 1; i < p.size(); ++i) {</pre>
    if (abs2(cent - p[i]) <= r) continue;</pre>
    cent = p[i];
    r = 0.0;
    for (int j = 0; j < i; ++j) {</pre>
      if (abs2(cent - p[j]) <= r) continue;</pre>
      cent = (p[i] + p[j]) / 2;
      r = abs2(p[j] - cent);
      for (int k = 0; k < j; ++k) {
        if (abs2(cent - p[k]) <= r) continue;</pre>
        cent = circenter(p[i], p[j], p[k]);
        r = abs2(p[k] - cent);
   }
  return {cent, sqrt(r)};
```

# 8.15 Union of Circles

```
vector<pair<double, double>> CoverSegment(Cir a, Cir b)
  double d = abs(a.o - b.o);
  vector<pair<double, double>> res;
  if (sign(a.r + b.r - d) == 0);
  else if (d <= abs(a.r - b.r) + eps) {</pre>
    if (a.r < b.r) res.emplace_back(0, 2 * pi);</pre>
  } else if (d < abs(a.r + b.r) - eps) {</pre>
    double o = acos((sqrt(a.r) + sqrt(d) - sqrt(b.r)) /
         (2 * a.r * d)), z = atan2((b.o - a.o).y, (b.o
        - a.o).x);
   if (z < 0) z += 2 * pi;
    double l = z - o, r = z + o;
    if (1 < 0) 1 += 2 * pi;</pre>
   if (r > 2 * pi) r -= 2 * pi;
    if (1 > r) res.emplace_back(1, 2 * pi), res.
        emplace_back(0, r);
    else res.emplace_back(l, r);
 return res:
double CircleUnionArea(vector<Cir> c) { // circle
    should be identical
```

```
int n = c.size();
double a = 0, w;
for (int i = 0; w = 0, i < n; ++i) {</pre>
  vector<pair<double, double>> s = {{2 * pi, 9}}, z;
  for (int j = 0; j < n; ++j) if (i != j) {</pre>
   z = CoverSegment(c[i], c[j]);
    for (auto &e : z) s.push_back(e);
  sort(s.begin(), s.end());
  auto F = [&] (double t) { return c[i].r * (c[i].r *
       t + c[i].o.x * sin(t) - c[i].o.y * cos(t)); };
  for (auto &e : s) {
    if (e.first > w) a += F(e.first) - F(w);
    w = max(w, e.second);
  }
}
return a * 0.5;
```

## 8.16 Polar Angle Sort

# 8.17 Rotating Caliper

## 8.18 Rotating SweepLine

```
void RotatingSweepLine(vector <Pt> &pt) {
  int n = pt.size();
  vector <int> id(n), pos(n);
  vector <pair <int, int>> line;
  sort(line.begin(), line.end(), [&](pair <int, int> i,
      pair <int, int> j) {
    Pt a = pt[i.second] - pt[i.first], b = pt[j.second]
         - pt[j.first];
    return (a.pos() == b.pos() ? sign(a ^ b) > 0 : a.
        pos() < b.pos());
  });
  iota(id.begin(), id.end(), 0);
  sort(id.begin(), id.end(), [&](int i, int j) {
    return (sign(pt[i].y - pt[j].y) == 0 ? pt[i].x < pt</pre>
        [j].x : pt[i].y < pt[j].y);
  for (int i = 0; i < n; ++i)</pre>
   pos[id[i]] = i;
  for (auto [i, j] : line) {
   // point sort by the distance to line(i, j)  
    // do something
    tie(pos[i], pos[j], id[pos[i]], id[pos[j]]) =
        make_tuple(pos[j], pos[i], j, i);
}
```

## 8.19 Half Plane Intersection

```
// first ----> second
auto pos = [&](Pt a) {return sign(a.y) == 0 ? sign(a
    .x) < 0 : sign(a.y) > 0;};
sort(all(vec), [&](pair <Pt, Pt> a, pair <Pt, Pt> b)
  Pt A = a.second - a.first, B = b.second - b.first;
  if (pos(A) == pos(B)) {
    if (sign(A ^ B) == 0) return sign((b.first - a.
        first) * (b.second - a.first)) > 0;
    return sign(A ^ B) > 0;
  return pos(A) < pos(B);</pre>
});
deque <Pt> inter;
deque <pair <Pt, Pt>> seg;
int n = vec.size();
auto get = [&](pair <Pt, Pt> a, pair <Pt, Pt> b) {
    return intersect(a.first, a.second, b.first, b.
    second);};
for (int i = 0; i < n; ++i) if (!i || vec[i] != vec[i</pre>
      - 1]) {
  while (seg.size() >= 2 && sign((vec[i].second 
      inter.back()) ^ (vec[i].first - inter.back()))
      == 1) seg.pop_back(), inter.pop_back();
  while (seg.size() >= 2 && sign((vec[i].second -
      inter.front()) ^ (vec[i].first - inter.front())
      ) == 1) seg.pop_front(), inter.pop_front();
  seg.push_back(vec[i]);
  if (seg.size() >= 2) inter.pb(get(seg[seg.size() -
      2], seg.back()));
while (seg.size() >= 2 && sign((seg.front().second -
    inter.back()) ^ (seg.front().first - inter.back()
    )) == 1) seg.pop_back(), inter.pop_back();
inter.push_back(get(seg.front(), seg.back()));
return vector <Pt>(all(inter));
```

## 8.20 Minkowski Sum

```
void reorder(vector <Pt> &P) {
 rotate(P.begin(), min_element(all(P), [&](Pt a, Pt b)
       { return make_pair(a.y, a.x) < make_pair(b.y, b.
      x); }), P.end());
vector <Pt> Minkowski(vector <Pt> P, vector <Pt> Q) {
 // P, Q: convex polygon
  reorder(P), reorder(Q);
  int n = P.size(), m = Q.size();
 P.pb(P[0]), P.pb(P[1]), Q.pb(Q[0]), Q.pb(Q[1]);
 vector <Pt> ans;
 for (int i = 0, j = 0; i < n || j < m; ) {
   ans.pb(P[i] + Q[j]);
    auto val = (P[i + 1] - P[i]) ^ (Q[j + 1] - Q[j]);
   if (val >= 0) i++;
    if (val <= 0) j++;</pre>
  return ans;
```

#### 8.21 Delaunay Triangulation

```
/* Delaunay Triangulation:
Given a sets of points on 2D plane, find a
triangulation such that no points will strictly
inside circumcircle of any triangle.
find : return a triangle contain given point
add_point : add a point into triangulation
A Triangle is in triangulation iff. its has_chd is 0.
Region of triangle u: iterate each u.edge[i].tri,
each points are u.p[(i+1)\%3], u.p[(i+2)\%3]
Voronoi diagram: for each triangle in triangulation,
the bisector of all its edges will split the region.
nearest point will belong to the triangle containing it
const ll inf = MAXC * MAXC * 100; // Lower_bound
    unknown
struct Tri;
struct Edge {
  Tri* tri; int side;
  Edge(): tri(0), side(0){}
  Edge(Tri* _tri, int _side): tri(_tri), side(_side){}
```

```
struct Tri {
  pll p[3];
 Edge edge[3];
Tri* chd[3];
  Tri() {}
  Tri(const pll& p0, const pll& p1, const pll& p2) {
    p[0] = p0; p[1] = p1; p[2] = p2;
    chd[0] = chd[1] = chd[2] = 0;
  bool has_chd() const { return chd[0] != 0; }
  int num_chd() const {
    return !!chd[0] + !!chd[1] + !!chd[2];
  bool contains(pll const& q) const {
    for (int i = 0; i < 3; ++i)
      if (ori(p[i], p[(i + 1) % 3], q) < 0)</pre>
       return 0;
} pool[N * 10], *tris;
void edge(Edge a, Edge b) {
 if(a.tri) a.tri->edge[a.side] = b;
  if(b.tri) b.tri->edge[b.side] = a;
struct Trig { // Triangulation
 Trig() {
    the_root = // Tri should at least contain all
        points
      new(tris++) Tri(pll(-inf, -inf), pll(inf + inf, -
          inf), pll(-inf, inf + inf));
  Tri* find(pll p) { return find(the_root, p); }
  void add_point(const pll &p) { add_point(find(
      the_root, p), p); }
  Tri* the root;
  static Tri* find(Tri* root, const pll &p) {
    while (1) {
      if (!root->has_chd())
        return root;
      for (int i = 0; i < 3 && root->chd[i]; ++i)
        if (root->chd[i]->contains(p)) {
          root = root->chd[i];
          break:
        }
    assert(0); // "point not found"
  void add point(Tri* root, pll const& p) {
    Tri* t[3];
    /* split it into three triangles */
    for (int i = 0; i < 3; ++i)</pre>
      t[i] = new(tris++) Tri(root->p[i], root->p[(i +
          1) % 3], p);
    for (int i = 0; i < 3; ++i)</pre>
      edge(Edge(t[i], 0), Edge(t[(i + 1) % 3], 1));
    for (int i = 0; i < 3; ++i)
      edge(Edge(t[i], 2), root->edge[(i + 2) % 3]);
    for (int i = 0; i < 3; ++i)</pre>
     root->chd[i] = t[i];
    for (int i = 0; i < 3; ++i)
      flip(t[i], 2);
  void flip(Tri* tri, int pi) {
    Tri* trj = tri->edge[pi].tri;
    int pj = tri->edge[pi].side;
    if (!trj) return;
    if (!in_cc(tri->p[0], tri->p[1], tri->p[2], trj->p[
        pj])) return;
    /* flip edge between tri,trj */
    Tri* trk = new(tris++) Tri(tri->p[(pi + 1) % 3],
        trj->p[pj], tri->p[pi]);
    Tri* trl = new(tris++) Tri(trj->p[(pj + 1) % 3],
        tri->p[pi], trj->p[pj]);
    edge(Edge(trk, 0), Edge(trl, 0));
    edge(Edge(trk, 1), tri->edge[(pi + 2) % 3]);
    edge(Edge(trk, 2), trj->edge[(pj + 1) % 3]);
    edge(Edge(trl, 1), trj->edge[(pj + 2) % 3]);
    edge(Edge(trl, 2), tri->edge[(pi + 1) % 3]);
    tri->chd[0] = trk; tri->chd[1] = trl; tri->chd[2] =
    trj->chd[0] = trk; trj->chd[1] = trl; trj->chd[2] =
```

```
0:
    flip(trk, 1); flip(trk, 2);
    flip(trl, 1); flip(trl, 2);
vector<Tri*> triang; // vector of all triangle
set<Tri*> vst;
void go(Tri* now) { // store all tri into triang
 if (vst.find(now) != vst.end())
    return;
  vst.insert(now);
 if (!now->has_chd())
    return triang.pb(now);
 for (int i = 0; i < now->num_chd(); ++i)
   go(now->chd[i]);
void build(int n, pll* ps) { // build triangulation
 tris = pool; triang.clear(); vst.clear();
 random_shuffle(ps, ps + n);
 Trig tri; // the triangulation structure
 for (int i = 0; i < n; ++i)</pre>
   tri.add_point(ps[i]);
 go(tri.the_root);
```

## 8.22 Triangulation Vonoroi

```
vector<Line> ls[N];
pll arr[N];
Line make_line(pdd p, Line 1) {
  pdd d = 1.Y - 1.X; d = perp(d);
  pdd m = (1.X + 1.Y) / 2;
  l = Line(m, m + d);
 if (ori(1.X, 1.Y, p) < 0)</pre>
   l = Line(m + d, m);
 return 1;
double calc_area(int id) {
 // use to calculate the area of point "strictly in
      the convex hull"
  vector<Line> hpi = halfPlaneInter(ls[id]);
 vector<pdd> ps;
  for (int i = 0; i < SZ(hpi); ++i)</pre>
    ps.pb(intersect(hpi[i].X, hpi[i].Y, hpi[(i + 1) %
        SZ(hpi)].X, hpi[(i + 1) % SZ(hpi)].Y));
  double rt = 0;
  for (int i = 0; i < SZ(ps); ++i)</pre>
    rt += cross(ps[i], ps[(i + 1) % SZ(ps)]);
  return fabs(rt) / 2;
void solve(int n, pii *oarr) {
 map<pll, int> mp;
for (int i = 0; i < n; ++i)</pre>
    arr[i] = pll(oarr[i].X, oarr[i].Y), mp[arr[i]] = i;
 build(n, arr); // Triangulation
for (auto *t : triang) {
    vector<int> p;
    for (int i = 0; i < 3; ++i)</pre>
      if (mp.find(t->p[i]) != mp.end())
        p.pb(mp[t->p[i]]);
    for (int i = 0; i < SZ(p); ++i)</pre>
      for (int j = i + 1; j < SZ(p); ++j) {</pre>
        Line l(oarr[p[i]], oarr[p[j]]);
        ls[p[i]].pb(make_line(oarr[p[i]], 1));
        ls[p[j]].pb(make_line(oarr[p[j]], 1));
```

## 9 Else

## 9.1 Bit Hack

```
9.2 Dynamic Programming Condition9.2.1 Totally Monotone (Concave/Convex)
```

```
\begin{array}{l} \forall i < i', j < j' \text{, } B[i][j] \leq B[i'][j] \implies B[i][j'] \leq B[i'][j'] \\ \forall i < i', j < j' \text{, } B[i][j] \geq B[i'][j] \implies B[i][j'] \geq B[i'][j'] \end{array}
```

# 9.2.2 Monge Condition (Concave/Convex)

```
\begin{array}{l} \forall i < i', j < j' \text{, } B[i][j] + B[i'][j'] \geq B[i][j'] + B[i'][j] \\ \forall i < i', j < j' \text{, } B[i][j] + B[i'][j'] \leq B[i][j'] + B[i'][j] \end{array}
```

#### 9.2.3 Optimal Split Point

```
If B[i][j] + B[i+1][j+1] \geq B[i][j+1] + B[i+1][j] then H_{i,j-1} \leq H_{i,j} \leq H_{i+1,j}
```

## 9.3 Slope Trick

}

```
template<typename T>
struct slope_trick_convex {
  T minn = 0, ground_1 = 0, ground_r = 0;
  priority_queue<T, vector<T>, less<T>> left;
  priority_queue<T, vector<T>, greater<T>> right;
  slope_trick_convex() {left.push(numeric_limits<T>::
      min() / 2), right.push(numeric_limits<T>::max() /
       2);}
  void push_left(T x) {left.push(x - ground_1);}
  void push_right(T x) {right.push(x - ground_r);}
  //add a line with slope 1 to the right starting from
  void add_right(T x) {
    T 1 = left.top() + ground_1;
    if (1 <= x) push_right(x);</pre>
    else push_left(x), push_right(l), left.pop(), minn
        += 1 - x;
  //add a line with slope -1 to the left starting from
  void add_left(T x) {
    T r = right.top() + ground_r;
    if (r >= x) push_left(x);
    else push_right(x), push_left(r), right.pop(), minn
  //val[i]=min(val[j]) for all i-l<=j<=i+r
  void expand(T 1, T r) {ground_1 -= 1, ground_r += r;}
  void shift_up(T x) {minn += x;}
  T get_val(T x) {
    T l = left.top() + ground_l, r = right.top() +
        ground_r;
    if (x >= 1 && x <= r) return minn;
    if (x < 1) {
      vector<T> trash;
      T cur_val = minn, slope = 1, res;
      while (1) {
        trash.push_back(left.top());
        left.pop();
        if (left.top() + ground_l <= x) {</pre>
          res = cur_val + slope * (1 - x);
          break:
        cur_val += slope * (1 - (left.top() + ground_1)
        1 = left.top() + ground_l;
        slope += 1;
      for (auto i : trash) left.push(i);
      return res;
    if(x > r) {
      vector<T> trash;
      T cur_val = minn, slope = 1, res;
      while (1) {
        trash.push_back(right.top());
        right.pop();
        if (right.top() + ground_r >= x) {
          res = cur_val + slope * (x - r);
          break;
        }
```

#### 9.4 Manhattan MST

```
void solve(int n) {
 init();
  vector<int> v(n), ds;
  for (int i = 0; i < n; ++i) {</pre>
    v[i] = i;
    ds.push_back(x[i] - y[i]);
  sort(ds.begin(), ds.end());
  ds.resize(unique(ds.begin(), ds.end()) - ds.begin());
  sort(v.begin(), v.end(), [&](int i, int j) { return x
      [i] == x[j] ? y[i] > y[j] : x[i] > x[j]; });
  int j = 0;
  for (int i = 0; i < n; ++i) {</pre>
   int p = lower_bound(ds.begin(), ds.end(), x[v[i]] -
         y[v[i]]) - ds.begin() + 1;
    pair<int, int> q = query(p);
    // query return prefix minimum
    if (~q.second) add_edge(v[i], q.second);
    add(p, make_pair(x[v[i]] + y[v[i]], v[i]));
 }
void make_graph() {
  solve(n);
  for (int i = 0; i < n; ++i) swap(x[i], y[i]);</pre>
  solve(n);
  for (int i = 0; i < n; ++i) x[i] = -x[i];
 solve(n):
 for (int i = 0; i < n; ++i) swap(x[i], y[i]);</pre>
  solve(n);
```

#### 9.5 Dynamic MST

```
int cnt[maxn], cost[maxn], st[maxn], ed[maxn];
pair<int, int> qr[maxn];
// qr[i].first = id of edge to be changed, qr[i].second
     = weight after operation
// cnt[i] = number of operation on edge i
// call solve(0, q - 1, v, 0), where v contains edges i
     such that cnt[i] == 0
void contract(int 1, int r, vector<int> v, vector<int>
    &x, vector<int> &y) {
  sort(v.begin(), v.end(), [&](int i, int j) {
      if (cost[i] == cost[j]) return i < j;</pre>
      return cost[i] < cost[j];</pre>
      });
 djs.save();
 for (int i = 1; i <= r; ++i) djs.merge(st[qr[i].first</pre>
      ], ed[qr[i].first]);
 for (int i = 0; i < (int)v.size(); ++i) {</pre>
    if (djs.find(st[v[i]]) != djs.find(ed[v[i]])) {
      x.push_back(v[i]);
      djs.merge(st[v[i]], ed[v[i]]);
 djs.undo();
  djs.save();
  for (int i = 0; i < (int)x.size(); ++i) djs.merge(st[</pre>
      x[i]], ed[x[i]]);
  for (int i = 0; i < (int)v.size(); ++i) {</pre>
    if (djs.find(st[v[i]]) != djs.find(ed[v[i]])) {
      y.push_back(v[i]);
      djs.merge(st[v[i]], ed[v[i]]);
   }
  djs.undo();
```

```
void solve(int 1, int r, vector<int> v, long long c) {
 if (1 == r) {
    cost[qr[1].first] = qr[1].second;
    if (st[qr[1].first] == ed[qr[1].first]) {
      printf("%lld\n", c);
      return;
    int minv = qr[1].second;
    for (int i = 0; i < (int)v.size(); ++i) minv = min(</pre>
        minv, cost[v[i]]);
    printf("%lld \ n", c + minv);
  int m = (1 + r) >> 1;
  vector<int> lv = v, rv = v;
  vector<int> x, y;
  for (int i = m + 1; i <= r; ++i) {</pre>
    cnt[qr[i].first]--
    if (cnt[qr[i].first] == 0) lv.push_back(qr[i].first
  }
  contract(l, m, lv, x, y);
  long long lc = c, rc = c;
  djs.save();
  for (int i = 0; i < (int)x.size(); ++i) {</pre>
    lc += cost[x[i]];
    djs.merge(st[x[i]], ed[x[i]]);
  solve(1, m, y, lc);
  djs.undo();
  x.clear(), y.clear();
  for (int i = m + 1; i <= r; ++i) cnt[qr[i].first]++;</pre>
  for (int i = 1; i <= m; ++i) {</pre>
    cnt[qr[i].first]--
    if (cnt[qr[i].first] == 0) rv.push_back(qr[i].first
  contract(m + 1, r, rv, x, y);
  djs.save();
  for (int i = 0; i < (int)x.size(); ++i) {</pre>
    rc += cost[x[i]];
    djs.merge(st[x[i]], ed[x[i]]);
  solve(m + 1, r, y, rc);
  djs.undo();
  for (int i = 1; i <= m; ++i) cnt[qr[i].first]++;</pre>
```

#### 9.6 ALL LCS

```
void all_lcs(string s, string t) { // 0-base
  vector<int> h(t.size());
  iota(all(h), 0);
  for (int a = 0; a < s.size(); ++a) {
    int v = -1;
    for (int c = 0; c < t.size(); ++c)
        if (s[a] == t[c] || h[c] < v)
            swap(h[c], v);
        // LCS(s[0, a], t[b, c]) =
        // c - b + 1 - sum([h[i] >= b] | i <= c)
        // h[i] might become -1 !!
    }
}</pre>
```

### 9.7 Hilbert Curve

```
long long hilbertOrder(int x, int y, int pow, int
    rotate) {
    if (pow == 0) return 0;
    int hpow = 1 << (pow-1);
    int seg = (x < hpow) ? ((y < hpow) ? 0 : 3) : ((y <
        hpow) ? 1 : 2);
    seg = (seg + rotate) & 3;
    const int rotateDelta[4] = {3, 0, 0, 1};
    int nx = x & (x ^ hpow), ny = y & (y ^ hpow);
    int nrot = (rotate + rotateDelta[seg]) & 3;
    long long subSquareSize = 1ll << (pow * 2 - 2);
    long long ans = seg * subSquareSize;
    long long add = hilbertOrder(nx, ny, pow - 1, nrot);
    ans += (seg == 1 | seg == 2) ? add : (subSquareSize - add - 1);</pre>
```

```
9.8 Pbds
```

return ans:

```
#include <ext/pb_ds/priority_queue.hpp>
#include <ext/pb_ds/assoc_container.hpp>
using namespace __gnu_pbds;
#include <ext/rope>
using namespace __gnu_cxx;
int main () {
     _gnu_pbds::priority_queue <<mark>int</mark>> pq1, pq2;
  pq1.join(pq2); // pq1 += pq2, pq2 = {}
  cc_hash_table<int, int> m1;
  tree<int, null_type, less<int>, rb_tree_tag,
      tree_order_statistics_node_update> oset;
  oset.insert(2), oset.insert(4);
  cout << *oset.find_by_order(1) << ' ' << oset.</pre>
      order_of_key(1) << '\n'; // 4 0
  bitset <100> BS:
  BS.flip(3), BS.flip(5);
 cout << BS._Find_first() << ' ' << BS._Find_next(3)</pre>
      << '\n'; // 3 5
  rope <int> rp1, rp2;
 rp1.push_back(1), rp1.push_back(3);
 rp1.insert(0, 2); // pos, num
 rp1.erase(0, 2); // pos, len
 rp1.substr(0, 2); // pos, len
 rp2.push_back(4);
 rp1 += rp2, rp2 = rp1;
cout << rp2[0] << ' ' << rp2[1] << '\n'; // 3 4
```

# 9.9 Random

```
struct custom_hash {
    static uint64_t splitmix64(uint64_t x) {
        x += 0x9e3779b97f4a7c15;
        x = (x ^ (x >> 30)) * 0xbf58476d1ce4e5b9;
        x = (x ^ (x >> 27)) * 0x94d049bb133111eb;
        return x ^ (x >> 31);
    }
    size_t operator()(uint64_t a) const {
        static const uint64_t FIXED_RANDOM = chrono::
            steady_clock::now().time_since_epoch().count();
        return splitmix64(i + FIXED_RANDOM);
    }
};
unordered_map <int, int, custom_hash> m1;
random_device rd; mt19937 rng(rd());
```

## 9.10 Smawk Algorithm

```
11 query(int 1, int r) {
 // ...
struct SMAWK {
  // Condition:
  // If M[1][0] < M[1][1] then M[0][0] < M[0][1]
  // If M[1][0] == M[1][1] then M[0][0] \leftarrow= M[0][1]
  // For all i, find r_i s.t. M[i][r_i] is maximum ||
      minimum.
  int ans[N], tmp[N];
  void interpolate(vector <int> 1, vector <int> r) {
    int n = 1.size(), m = r.size();
    vector <int> nl;
    for (int i = 1; i < n; i += 2) {</pre>
      nl.push_back(l[i]);
    run(nl, r);
    for (int i = 1, j = 0; i < n; i += 2) {
      while (j < m && r[j] < ans[l[i]])</pre>
        j++;
      assert(j < m && ans[l[i]] == r[j]);
      tmp[l[i]] = j;
    for (int i = 0; i < n; i += 2) {</pre>
      int curl = 0, curr = m - 1;
      if (i)
        curl = tmp[l[i - 1]];
      if (i + 1 < n)
        curr = tmp[l[i + 1]];
```

```
11 res = query(l[i], r[curl]);
      ans[l[i]] = r[curl];
      for (int j = curl + 1; j <= curr; ++j) {</pre>
        11 nxt = query(l[i], r[j]);
         if (res < nxt)</pre>
           res = nxt, ans[l[i]] = r[j];
      }
    }
  }
  void reduce(vector <int> 1, vector <int> r) {
    int n = l.size(), m = r.size();
    vector <int> nr;
    for (int j : r) {
      while (!nr.empty()) {
        int i = nr.size() - 1;
         if (query(1[i], nr.back()) <= query(1[i], j))</pre>
          nr.pop_back();
         else
          break:
      if (nr.size() < n)</pre>
        nr.push_back(j);
    run(1, nr);
  void run(vector <int> 1, vector <int> r) {
    int n = l.size(), m = r.size();
    if (max(n, m) <= 2) {
      for (int i : 1) {
        ans[i] = r[0];
        if (m > 1)
           if (query(i, r[0]) < query(i, r[1]))</pre>
             ans[i] = r[1];
        }
    } else if (n >= m) {
      interpolate(1, r);
    } else {
      reduce(1, r);
  }
};
```

#### 9.11 Matroid Intersection

Start from  $S=\emptyset$ . In each iteration, let

```
• Y_1 = \{x \notin S \mid S \cup \{x\} \in I_1\}
• Y_2 = \{x \notin S \mid S \cup \{x\} \in I_2\}
```

If there exists  $x\in Y_1\cap Y_2$ , insert x into S. Otherwise for each  $x\in S, y\not\in S$ , create edges

```
• x \to y if S - \{x\} \cup \{y\} \in I_1.
• y \to x if S - \{x\} \cup \{y\} \in I_2.
```

Find a shortest path (with BFS) starting from a vertex in  $Y_1$  and ending at a vertex in  $Y_2$  which doesn't pass through any other vertices in  $Y_2$ , and alternate the path. The size of S will be incremented by 1 in each iteration. For the weighted case, assign weight w(x) to vertex x if  $x \in S$  and -w(x) if  $x \not\in S$ . Find the path with the minimum number of edges among all minimum length paths and alternate it.

#### 9.12 Python Misc