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```

Basic

1.1 Shell Script

```
++ -std=c++17 -DABS -Wall -Wextra -Wshadow $1.cpp -o
   $1 && ./$1
or i in {A..J}; do cp tem.cpp $i.cpp; done;
```

1.2 Default Code

```
include <bits/stdc++.h>
sing namespace std;
ypedef long long ll;
define pb push_back
define pii pair<int, int>
define all(a) a.begin(), a.end()
define sz(a) ((int)a.size())
```

L.3 Increase Stack Size

```
onst int size = 256 << 20;
egister long rsp asm("rsp");
har *p = (char*)malloc(size) + size, *bak = (char*)rsp
     _("movq %0, %%rsp\n"::"r"(p));
asm
/ main
_asm__("movq %0, %%rsp\n"::"r"(bak));
```

1.4 Debug Macro

exit 1

```
oid db() {cout << endl;}</pre>
emplate <typename T, typename ...U> void db(T i, U ...
   i) {
 cout << i << ' ', db(j...);</pre>
define test(x...) db("[" + string(x) + "]", x)
```

1.5 Stress Test Shell*

```
!/usr/bin/env bash
++ $1.cpp -o $1
g++ $2.cpp -o $2
++ $3.cpp -o $3
or i in {1..100} ; do
 ./$3 > input.txt
 # st=$(date +%s%N)
 ./$1 < input.txt > output1.txt
 # echo "$((($(date +%s%N) - $st)/1000000))ms"
 ./$2 < input.txt > output2.txt
 if cmp --silent -- "output1.txt" "output2.txt"; then
   continue
 fi
 echo Input:
 cat input.txt
 echo Your Output:
 cat output1.txt
 echo Correct Output:
 cat output2.txt
```

```
done
echo OK!
./stress.sh main good gen
```

1.6 Pragma / FastIO

```
#pragma GCC optimize("Ofast,inline,unroll-loops")
#pragma GCC target("bmi,bmi2,lzcnt,popcnt,avx2")
#include<unistd.h>
char OB[65536]; int OP;
inline char RC() {
  static char buf[65536], *p = buf, *q = buf;
  return p == q && (q = (p = buf) + read(0, buf, 65536)
      ) == buf ? -1 : *p++;
inline int R() {
 static char c:
  while((c = RC()) < '0'); int a = c ^ '0';</pre>
 while((c = RC()) >= '0') a *= 10, a += c ^ '0';
 return a;
inline void W(int n) {
 static char buf[12], p;
 if (n == 0) OB[OP++] = '0'; p = 0;
 while (n) buf[p++] = '0' + (n % 10), n /= 10;
 for (--p; p >= 0; --p) OB[OP++] = buf[p];
 if (OP > 65520) write(1, OB, OP), OP = 0;
```

1.7 Divide*

```
ll divdown(ll a, ll b) {
  return a / b - (a < 0 && a % b);
}
ll divup(ll a, ll b) {
  return a / b + (a > 0 && a % b);
}
a / b < x -> divdown(a, b) + 1 <= x
a / b <= x -> divup(a, b) <= x
x < a / b -> x <= divup(a, b) - 1
x <= a / b -> x <= divdown(a, b)</pre>
```

2 Data Structure

2.1 Leftist Tree

```
struct node {
  ll rk, data, sz, sum;
  node *1, *r;
  node(ll \ k) : rk(0), data(k), sz(1), l(0), r(0), sum(k)
11 sz(node *p) { return p ? p->sz : 0; }
11 rk(node *p) { return p ? p->rk : -1; }
11 sum(node *p) { return p ? p->sum : 0; }
node *merge(node *a, node *b) {
 if (!a || !b) return a ? a : b;
 if (a->data < b->data) swap(a, b);
  a->r = merge(a->r, b);
  if (rk(a->r) > rk(a->l)) swap(a->r, a->l);
  a->rk = rk(a->r) + 1, a->sz = sz(a->l) + sz(a->r) +
      1:
  a \rightarrow sum = sum(a \rightarrow 1) + sum(a \rightarrow r) + a \rightarrow data;
  return a;
void pop(node *&o) {
  node *tmp = o;
  o = merge(o \rightarrow 1, o \rightarrow r):
  delete tmp;
```

2.2 Splay Tree

```
struct Splay {
  int pa[N], ch[N][2], sz[N], rt, _id;
  ll v[N];
  Splay() {}
  void init() {
    rt = 0, pa[0] = ch[0][0] = ch[0][1] = -1;
    sz[0] = 1, v[0] = inf;
```

```
int newnode(int p, int x) {
    int id = _id++;
    v[id] = x, pa[id] = p;
    ch[id][0] = ch[id][1] = -1, sz[id] = 1;
  void rotate(int i) {
    int p = pa[i], x = ch[p][1] == i, gp = pa[p], c =
        ch[i][!x];
    sz[p] -= sz[i], sz[i] += sz[p];
    if (~c) sz[p] += sz[c], pa[c] = p;
    ch[p][x] = c, pa[p] = i;
    pa[i] = gp, ch[i][!x] = p;
    if (~gp) ch[gp][ch[gp][1] == p] = i;
  void splay(int i) {
    while (~pa[i])
      int p = pa[i];
      if (~pa[p]) rotate(ch[pa[p]][1] == p ^ ch[p][1]
          == i ? i : p);
      rotate(i);
    rt = i;
  int lower_bound(int x) {
    int i = rt, last = -1;
    while (true) {
      if (v[i] == x) return splay(i), i;
      if (v[i] > x) {
        last = i:
        if (ch[i][0] == -1) break;
        i = ch[i][0];
      else {
        if (ch[i][1] == -1) break;
        i = ch[i][1];
     }
    splay(i);
    return last; // -1 if not found
  void insert(int x) {
    int i = lower_bound(x);
    if (i == -1) {
      // assert(ch[rt][1] == -1);
      int id = newnode(rt, x);
      ch[rt][1] = id, ++sz[rt];
      splav(id):
    else if (v[i] != x) {
      splay(i);
      int id = newnode(rt, x), c = ch[rt][0];
      ch[rt][0] = id;
      ch[id][0] = c;
      if (~c) pa[c] = id, sz[id] += sz[c];
      ++sz[rt];
      splay(id);
 }
};
```

2.3 Link Cut Tree

```
// weighted subtree size, weighted path max
struct LCT {
 int ch[N][2], pa[N], v[N], sz[N], sz2[N], w[N], mx[N
      ], _id;
  // sz := sum \ of \ v \ in \ splay, \ sz2 := sum \ of \ v \ in
      virtual subtree
  // mx := max w in splay
  bool rev[N];
  LCT() : _id(1) {}
  int newnode(int _v, int _w) {
    int x = _id++;
ch[x][0] = ch[x][1] = pa[x] = 0;
    v[x] = sz[x] = v;
    sz2[x] = 0;
    w[x] = mx[x] = w;
    rev[x] = false;
    return x;
```

```
void pull(int i) {
    sz[i] = v[i] + sz2[i];
    mx[i] = w[i];
    if (ch[i][0])
      sz[i] += sz[ch[i][0]], mx[i] = max(mx[i], mx[ch[i])
           ][0]]);
    if (ch[i][1])
       sz[i] += sz[ch[i][1]], mx[i] = max(mx[i], mx[ch[i])
          ][1]]);
  void push(int i) {
    rev[i] = false;
  void reverse(int i) {
    if (!i) return;
    swap(ch[i][0], ch[i][1]);
    rev[i] ^= true;
  bool isrt(int i) {// rt of splay
    if (!pa[i]) return true;
    return ch[pa[i]][0] != i && ch[pa[i]][1] != i;
  void rotate(int i) {
    int p = pa[i], x = ch[p][1] == i, c = ch[i][!x], gp
         = pa[p];
    if (ch[gp][0] == p) ch[gp][0] = i;
    else if (ch[gp][1] == p) ch[gp][1] = i;
    pa[i] = gp, ch[i][!x] = p, pa[p] = i;
    ch[p][x] = c, pa[c] = p;
    pull(p), pull(i);
  void splay(int i) {
    vector<int> anc;
    anc.push_back(i);
    while (!isrt(anc.back())) anc.push_back(pa[anc.back
         ()]);
    while (!anc.empty()) push(anc.back()), anc.pop_back
        ();
    while (!isrt(i)) {
      int p = pa[i];
      if (!isrt(p)) rotate(ch[p][1] == i ^ ch[pa[p]][1]
            == p ? i : p);
      rotate(i);
    }
  }
  void access(int i) {
    int last = 0;
    while (i) {
      splay(i);
      if (ch[i][1])
        sz2[i] += sz[ch[i][1]];
       sz2[i] -= sz[last];
      ch[i][1] = last;
      pull(i), last = i, i = pa[i];
  }
  void makert(int i) {
    access(i), splay(i), reverse(i);
  void link(int i, int j) {
    // assert(findrt(i) != findrt(j));
    makert(i);
    makert(j);
    pa[i] = j;
    sz2[j] += sz[i];
    pull(j);
  void cut(int i, int j) {
    makert(i), access(j), splay(i);
// assert(sz[i] == 2 && ch[i][1] == j);
    ch[i][1] = pa[j] = 0, pull(i);
  int findrt(int i) {
    access(i), splay(i);
    while (ch[i][0]) push(i), i = ch[i][0];
    splay(i);
    return i:
  }
|};
```

2.4 Treap

```
struct node {
  int data, sz;
  node *1, *r;
  node(int k) : data(k), sz(1), l(0), r(0) {}
  void up() {
    sz = 1;
    if (1) sz += 1->sz;
    if (r) sz += r->sz;
  void down() {}
int sz(node *a) { return a ? a->sz : 0; }
node *merge(node *a, node *b) {
  if (!a || !b) return a ? a : b;
  if (rand() % (sz(a) + sz(b)) < sz(a))
    return a \rightarrow down(), a \rightarrow r = merge(a \rightarrow r, b), a \rightarrow up(), a
  return b \rightarrow down(), b \rightarrow 1 = merge(a, b \rightarrow 1), b \rightarrow up(), b;
void split(node *o, node *&a, node *&b, int k) {
  if (!o) return a = b = 0, void();
  o->down();
  if (o->data <= k)
    a = o, split(o->r, a->r, b, k), <math>a->up();
  else b = o, split(o->1, a, b->1, k), b->up();
void split2(node *o, node *&a, node *&b, int k) {
  if (sz(o) <= k) return a = o, b = 0, void();
  o->down():
  if (sz(o->1) + 1 <= k)
    a = o, split2(o->r, a->r, b, k - <math>sz(o->l) - 1);
  else b = o, split2(o->1, a, b->1, k);
  o->up();
node *kth(node *o, int k) {
  if (k <= sz(o->1)) return kth(o->1, k);
  if (k == sz(o->1) + 1) return o;
  return kth(o\rightarrow r, k - sz(o\rightarrow l) - 1);
int Rank(node *o, int key) {
  if (!o) return 0;
  if (o->data < key)</pre>
    return sz(o->1) + 1 + Rank(o->r, key);
  else return Rank(o->1, key);
bool erase(node *&o, int k) {
  if (!o) return 0;
  if (o->data == k) {
    node *t = o;
    o->down(), o = merge(o->1, o->r);
    delete t:
    return 1;
  node *&t = k < o->data ? o->l : o->r;
  return erase(t, k) ? o->up(), 1 : 0;
void insert(node *&o, int k) {
 node *a, *b;
  o->down(), split(o, a, b, k)
  o = merge(a, merge(new node(k), b));
  o->up();
void interval(node *&o, int 1, int r) {
 node *a, *b, *c; // [l, r)
  o->down();
  split2(o, a, b, 1), split2(b, b, c, r - 1);
  // operate
  o = merge(a, merge(b, c)), o->up();
```

2.5 Persistent Segment Tree*

```
struct Seg {
    // Persistent Segment Tree, single point modify,
        range query sum
    // 0-indexed, [l, r)
    static Seg mem[M], *pt;
    int l, r, m, val;
    Seg* ch[2];
    Seg () = default;
```

```
Seg (int _l, int _r) : l(_l), r(_r), m(l + r >> 1),
      val(0) {
    if (r - 1 > 1) {
      ch[0] = new (pt++) Seg(1, m);
      ch[1] = new (pt++) Seg(m, r);
  void pull() {val = ch[0]->val + ch[1]->val;}
  Seg* modify(int p, int v) {
    Seg *now = new (pt++) Seg(*this);
    if (r - 1 == 1) {
      now->val = v;
    } else {
      now \rightarrow ch[p >= m] = ch[p >= m] \rightarrow modify(p, v);
      now->pull();
    return now:
  int query(int a, int b) {
    if (a <= 1 && r <= b) return val;</pre>
    int ans = 0;
    if (a < m) ans += ch[0]->query(a, b);
    if (m < b) ans += ch[1]->query(a, b);
    return ans;
} Seg::mem[M], *Seg::pt = mem;
// Init Tree
Seg *root = new (Seg::pt++) Seg(0, n);
```

2.6 2D Segment Tree*

```
// 2D range add, range sum in Log^2
struct seg {
  int 1, r;
  11 sum, 1z;
  seg *ch[2]{};
  seg(int _l, int _r) : 1(_l), r(_r), sum(0), lz(0) {}
  void push() {
    if (lz) ch[0]->add(l, r, lz), ch[1]->modify(l, r, lz)
         1z), 1z = 0;
  void pull() \{sum = ch[0] -> sum + ch[1] -> sum;\}
  void add(int _1, int _r, 11 d) {
  if (_1 <= 1 && r <= _r) {</pre>
       sum += d * (r - 1);
      1z += d;
      return;
    if (!ch[0]) ch[0] = new seg(1, 1 + r >> 1), ch[1] =
          new seg(l + r >> 1, r);
    push();
    if (_1 < 1 + r >> 1) ch[0]->add(_1, _r, d);
    if (l + r >> 1 < _r) ch[1]->add(_l, _r, d);
    pull();
  ill qsum(int _1, int _r) {
   if (_1 <= 1 && r <= _r) return sum;</pre>
    if (!ch[0]) return lz * (min(r, _r) - max(l, _l));
    push();
    11 \text{ res} = 0:
    if (_1 < 1 + r >> 1) res += ch[0]->qsum(_1, _r);
    if (1 + r >> 1 < _r) res += ch[1]->qsum(_1, _r);
    return res:
  }
};
struct seg2 {
 int 1, r;
  seg v, lz;
seg2 *ch[2]{};
  seg2(int _1, int _r) : 1(_1), r(_r), v(0, N), lz(0, N)
    if (1 < r - 1) ch[0] = new seg2(1, 1 + r >> 1), ch
         [1] = new seg2(1 + r >> 1, r);
  void add(int _1, int _r, int _12, int _r2, 11 d) {
  v.add(_12, _r2, d * (min(r, _r) - max(1, _1)));
}
    if (_1 <= 1 && r <= _r) {</pre>
      lz.add(_12, _r2, d);
      return:
    if (_l < l + r >> 1) ch[0]->add(_l, _r, _l2, _r2, d
```

2.7 Zkw*

```
ll mx[N << 1], sum[N << 1], lz[N << 1];
void add(int 1, int r, 11 d) { // [l, r), 0-based
  int len = 1, cntl = 0, cntr = 0;
  for (1 += N, r += N + 1; 1 ^ r ^ 1; 1 >>= 1, r >>= 1,
       len <<= 1) {
     sum[1] += cntl * d, sum[r] += cnt[r] * d;
    if (len > 1) {
      mx[1] = max(mx[1 << 1], mx[1 << 1 | 1]) + lz[1];
      mx[r] = max(mx[r << 1], mx[r << 1 | 1]) + lz[r];
    if (~1 & 1)
      sum[1 ^ 1] += d * len, mx[1 ^ 1] += d, lz[1 ^ 1]
          += d, cntl += len;
    if (r & 1)
      sum[r ^ 1] += d * len, mx[r ^ 1] += d, lz[r ^ 1]
          += d, cntr += len;
  sum[1] += cntl * d, sum[r] += cntr * d;
  if (len > 1) {
    mx[1] = max(mx[1 << 1], mx[1 << 1 | 1]) + lz[1];
    mx[r] = max(mx[r << 1], mx[r << 1 | 1]) + lz[r];
  cntl += cntr;
  for (1 >>= 1; 1; 1 >>= 1) {
    sum[1] += cntl * d;
    mx[1] = max(mx[1 << 1], mx[1 << 1 | 1]) + lz[1];
11 qsum(int 1, int r) {
  ll res = 0, len = 1, cntl = 0, cntr = 0;
  for (1 += N, r += N + 1; 1 ^ r ^ 1; 1 >>= 1, r >>= 1,
       len <<= 1) {
    res += cntl * lz[1] + cntr * lz[r];
    if (~l & 1) res += sum[l ^ 1], cntl += len;
    if (r & 1) res += sum[r ^ 1], cntr += len;
  res += cntl * lz[1] + cntr * lz[r];
  cntl += cntr;
  for (1 >>= 1; 1; 1 >>= 1) res += cnt1 * lz[1];
  return res;
11 qmax(int 1, int r) {
  11 max1 = -INF, maxr = -INF;
  for (1 += N, r += N + 1; 1 ^ r ^ 1; 1 >>= 1, r >>= 1)
    \max 1 += 1z[1], \max[r] += 1z[r];
    if (~1 & 1) maxl = max(maxl, mx[1 ^ 1]);
    if (r & 1) maxr = max(maxr, mx[r ^ 1]);
  \max l = \max(\max l + lz[l], \max r + lz[r]);
  for (1 >>= 1; 1; 1 >>= 1) max1 += lz[1];
  return max1:
}
```

2.8 Chtholly Tree*

```
bool operator () (const interval &a, const interval
        & b) const {
      return a.1 < b.1;</pre>
    }
  set <interval, cmp> s;
  vector <interval> split(int 1, int r) {
    // split into [0, l), [l, r), [r, n) and return [l, r]
    vector <interval> del, ans, re;
    auto it = s.lower_bound(interval(l, -1, 0));
    if (it != s.begin() && (it == s.end() || 1 < it->1)
        ) {
      --it;
      del.pb(*it);
      if (r < it->r) {
        re.pb(interval(it->l, l, it->v));
        ans.pb(interval(l, r, it->v));
        re.pb(interval(r, it->r, it->v));
      } else {
        re.pb(interval(it->l, l, it->v));
        ans.pb(interval(1, it->r, it->v));
      }
      ++it;
    for (; it != s.end() && it->r <= r; ++it) {</pre>
      ans.pb(*it);
      del.pb(*it);
    if (it != s.end() && it->l < r) {</pre>
      del.pb(*it);
      ans.pb(interval(it->l, r, it->v));
      re.pb(interval(r, it->r, it->v));
    for (interval &i : del)
      s.erase(i);
    for (interval &i : re)
      s.insert(i);
    return ans;
  void merge(vector <interval> a) {
    for (interval &i : a)
      s.insert(i);
  }
};
```

3 Flow / Matching

3.1 Dinic

```
struct Dinic { // 0-base
 struct edge {
   int to, cap, flow, rev;
  vector<edge> adj[N];
 int s, t, dis[N], cur[N], n;
 int dfs(int u, int cap) {
    if (u == t || !cap) return cap;
    for (int &i = cur[u]; i < (int)adj[u].size(); ++i)</pre>
      edge &e = adj[u][i];
      if (dis[e.to] == dis[u] + 1 && e.flow != e.cap) {
        int df = dfs(e.to, min(e.cap - e.flow, cap));
        if (df) {
          e.flow += df;
          adj[e.to][e.rev].flow -= df;
          return df;
     }
    dis[u] = -1;
   return 0;
  bool bfs() {
   fill_n(dis, n, -1);
    queue<int> q;
    q.push(s), dis[s] = 0;
    while (!q.empty()) {
      int tmp = q.front();
      q.pop();
      for (auto &u : adj[tmp])
```

```
if (!~dis[u.to] && u.flow != u.cap) {
           q.push(u.to);
           dis[u.to] = dis[tmp] + 1;
     return dis[t] != -1;
  int maxflow(int _s, int _t) {
    s = _s, t = _t;
     int flow = 0, df;
     while (bfs()) {
      fill_n(cur, n, 0);
       while ((df = dfs(s, INF))) flow += df;
    return flow;
  void init(int _n) {
     for (int i = 0; i < n; ++i) adj[i].clear();</pre>
  void reset() {
     for (int i = 0; i < n; ++i)</pre>
      for (auto &j : adj[i]) j.flow = 0;
  void add_edge(int u, int v, int cap) {
    adj[u].pb(edge{v, cap, 0, (int)adj[v].size()});
     adj[v].pb(edge{u, 0, 0, (int)adj[u].size() - 1});
};
```

3.2 Min Cost Max Flow

```
template <typename T1, typename T2>
struct MCMF { // T1 -> flow, T2 -> cost, 0-based
  const T1 INF1 = 1 << 30;</pre>
  const T2 INF2 = 1 << 30;</pre>
  struct edge {
    int v; T1 f; T2 c;
  } E[M << 1];
  vector <int> adj[N];
  T2 dis[N], pot[N];
  int rt[N], vis[N], n, m, s, t;
  bool SPFA() {
    fill_n(rt, n, -1), fill_n(dis, n, INF2);
    fill_n(vis, n, false);
    queue <int> q;
    q.push(s), dis[s] = 0, vis[s] = true;
    while (!q.empty()) {
      int v = q.front(); q.pop();
      vis[v] = false;
      for (int id : adj[v]) if (E[id].f > 0 && dis[E[id
           ].v] > dis[v] + E[id].c + pot[v] - pot[E[id].
           v]) {
          dis[E[id].v] = dis[v] + E[id].c + pot[v] -
               pot[E[id].v], rt[E[id].v] = id;
           if (!vis[E[id].v]) vis[E[id].v] = true, q.
               push(E[id].v);
        }
    return dis[t] != INF2;
  bool dijkstra() {
    fill_n(rt, n, -1), fill_n(dis, n, INF2);
    priority_queue <pair <T2, int>, vector <pair <T2,</pre>
        int>>, greater <pair <T2, int>>> pq;
    dis[s] = 0, pq.emplace(dis[s], s);
    while (!pq.empty()) {
      auto [d, v] = pq.top(); pq.pop();
      if (dis[v] < d) continue;</pre>
      for (int id : adj[v]) if (E[id].f > 0 && dis[E[id
           ].v] > dis[v] + E[id].c + pot[v] - pot[E[id].
           v1) {
          dis[E[id].v] = dis[v] + E[id].c + pot[v] -
               pot[E[id].v], rt[E[id].v] = id;
          pq.emplace(dis[E[id].v], E[id].v);
    return dis[t] != INF2;
  pair <T1, T2> solve(int _s, int _t) {
    s = _s, t = _t, fill_n(pot, n, 0);
T1 flow = 0; T2 cost = 0;
```

```
bool fr = true;
    while ((fr ? SPFA() : dijkstra())) {
      for (int i = 0; i < n; i++) {</pre>
        dis[i] += pot[i] - pot[s];
      T1 add = INF1;
      for (int i = t; i != s; i = E[rt[i] ^ 1].v) {
        add = min(add, E[rt[i]].f);
      for (int i = t; i != s; i = E[rt[i] ^ 1].v) {
        E[rt[i]].f -= add, E[rt[i] ^ 1].f += add;
      flow += add, cost += add * dis[t];
      fr = false;
      for (int i = 0; i < n; ++i) swap(dis[i], pot[i]);</pre>
    return make_pair(flow, cost);
  void init(int _n) {
    n = n, m = 0;
    for (int i = 0; i < n; ++i) adj[i].clear();</pre>
  void reset() {
    for (int i = 0; i < m; ++i) E[i].f = 0;</pre>
  void add_edge(int u, int v, T1 f, T2 c) {
    adj[u].pb(m), E[m++] = \{v, f, c\};
    adj[v].pb(m), E[m++] = \{u, 0, -c\};
};
```

3.3 Kuhn Munkres

```
template <typename T>
struct KM { // 0-based
 T w[N][N], h1[N], hr[N], slk[N];
 T fl[N], fr[N], pre[N]; int n;
 bool vl[N], vr[N];
  const T INF = 1e9;
  queue <int> q;
  KM (int _n) : n(_n) {
    for (int i = 0; i < n; ++i) for (int j = 0; j < n;
        ++j)
        w[i][j] = -INF;
  void add_edge(int a, int b, int wei) {
   w[a][b] = wei;
 bool check(int x) {
   if (vl[x] = 1, ~fl[x]) return q.push(fl[x]), vr[fl[
       x]] = 1;
    while (\sim x) swap(x, fr[fl[x] = pre[x]]);
    return 0;
  void bfs(int s) {
   fill(slk, slk + n, INF), fill(vl, vl + n, 0), fill(
        vr, vr + n, 0);
    q.push(s), vr[s] = 1;
    while (1) {
     T d;
      while (!q.empty()) {
        int y = q.front(); q.pop();
        for (int x = 0; x < n; ++x)
          if (!vl[x] \&\& slk[x] >= (d = hl[x] + hr[y] -
              w[x][y])
            if (pre[x] = y, d) slk[x] = d;
            else if (!check(x)) return;
     d = INF:
     for (int x = 0; x < n; ++x)
       if (!v1[x] && d > s1k[x]) d = s1k[x];
      for (int x = 0; x < n; ++x) {
       if (v1[x]) h1[x] += d;
        else slk[x] -= d;
       if (vr[x]) hr[x] -= d;
      for (int x = 0; x < n; ++x) if (!v1[x] && !s1k[x]
           && !check(x)) return;
    }
  T solve() {
```

3.4 SW Min Cut

```
template <typename T>
struct SW { // 0-based
  T g[N][N], sum[N]; int n;
  bool vis[N], dead[N];
  void init(int _n) {
    n = _n;
    for (int i = 0; i < n; ++i) fill(g[i], g[i] + n, 0)
    fill(dead, dead + n, false);
  void add_edge(int u, int v, T w) {
    g[u][v] += w, g[v][u] += w;
  T solve() {
    T ans = 1 << 30;
    for (int round = 0; round + 1 < n; ++round) {</pre>
      fill(vis, vis + n, false), fill(sum, sum + n, 0);
      int num = 0, s = -1, t = -1;
      while (num < n - round) {</pre>
        int now = -1;
        for (int i = 0; i < n; ++i) if (!vis[i] && !</pre>
             dead[i]) {
             if (now == -1 \mid | sum[now] < sum[i]) now = i
          }
        s = t, t = now;
        vis[now] = true, num++;
        for (int i = 0; i < n; ++i) if (!vis[i] && !</pre>
             dead[i]) {
             sum[i] += g[now][i];
      ans = min(ans, sum[t]);
      for (int i = 0; i < n; ++i) {</pre>
        g[i][s] += g[i][t];
        g[s][i] += g[t][i];
      dead[t] = true;
    return ans;
  }
};
```

3.5 Gomory Hu Tree

```
vector <array <int, 3>> GomoryHu(vector <vector <pii>>
    adj, int n) {
// Tree edge min -> mincut (0-based)
  Dinic flow(n);
  for (int i = 0; i < n; ++i) for (auto [j, w] : adj[i</pre>
      1)
      flow.add_edge(i, j, w);
  flow.record();
  vector <array <int, 3>> ans;
  vector <int> rt(n);
  for (int i = 0; i < n; ++i) rt[i] = 0;</pre>
  for (int i = 1; i < n; ++i) {</pre>
    int t = rt[i];
    flow.reset(); // clear flows on all edge
    ans.push_back({i, t, flow.solve(i, t)});
    flow.runbfs(i);
    for (int j = i + 1; j < n; ++j) if (rt[j] == t &&</pre>
        flow.vis[j]) {
        rt[j] = i;
      }
  return ans;
```

3.6 Blossom

```
struct Matching { // 0-based
  int fa[N], pre[N], match[N], s[N], v[N], n, tk;
  vector <int> g[N];
  queue <int> q;
  Matching (int _n) : n(_n), tk(0) {
    for (int i = 0; i <= n; ++i) match[i] = pre[i] = n;</pre>
    for (int i = 0; i < n; ++i) g[i].clear();</pre>
  void add_edge(int u, int v) {
    g[u].push_back(v), g[v].push_back(u);
  int Find(int u) {
    return u == fa[u] ? u : fa[u] = Find(fa[u]);
  int lca(int x, int y) {
    tk++:
    x = Find(x), y = Find(y);
    for (; ; swap(x, y)) {
  if (x != n) {
        if (v[x] == tk) return x;
        v[x] = tk;
        x = Find(pre[match[x]]);
    }
  void blossom(int x, int y, int 1) {
    while (Find(x) != 1) {
      pre[x] = y, y = match[x];
      if (s[y] == 1) q.push(y), s[y] = 0;
      if (fa[x] == x) fa[x] = 1;
      if (fa[y] == y) fa[y] = 1;
      x = pre[y];
    }
  bool bfs(int r) {
    for (int i = 0; i <= n; ++i) fa[i] = i, s[i] = -1;</pre>
    while (!q.empty()) q.pop();
    q.push(r);
    s[r] = 0;
    while (!q.empty()) {
      int x = q.front(); q.pop();
      for (int u : g[x]) {
        if (s[u] == -1) {
          pre[u] = x, s[u] = 1;
           if (match[u] == n) {
             for (int a = u, b = x, last; b != n; a =
                 last, b = pre[a])
               last = match[b], match[b] = a, match[a] =
                    b;
             return true;
          }
          q.push(match[u]);
          s[match[u]] = 0;
        } else if (!s[u] && Find(u) != Find(x)) {
          int 1 = 1ca(u, x);
          blossom(x, u, 1);
          blossom(u, x, 1);
        }
      }
    return false;
  int solve() {
    int res = 0;
    for (int x = 0; x < n; ++x) {
      if (match[x] == n) res += bfs(x);
    return res;
  }
};
```

3.7 Weighted Blossom

```
struct WeightGraph { // 1-based
  static const int inf = INT_MAX;
  static const int maxn = 514;
  struct edge {
   int u, v, w;
   edge(){}
   edge(int u, int v, int w): u(u), v(v), w(w) {}
```

```
int n, n_x;
edge g[maxn * 2][maxn * 2];
int lab[maxn * 2];
int match[maxn * 2], slack[maxn * 2], st[maxn * 2],
    pa[maxn * 2];
int flo_from[maxn * 2][maxn + 1], S[maxn * 2], vis[
    maxn * 2];
vector<int> flo[maxn * 2];
queue<int> q;
int e_delta(const edge &e) { return lab[e.u] + lab[e.
    v] - g[e.u][e.v].w * 2; }
void update_slack(int u, int x) { if (!slack[x] ||
    e_{delta}(g[u][x]) < e_{delta}(g[slack[x]][x])) slack
    [x] = u;
void set_slack(int x) {
  slack[x] = 0;
  for (int u = 1; u <= n; ++u)</pre>
    if (g[u][x].w > 0 \&\& st[u] != x \&\& S[st[u]] == 0)
      update_slack(u, x);
void q_push(int x) {
  if (x \le n) q.push(x);
  else for (size_t i = 0; i < flo[x].size(); i++)</pre>
      q_push(flo[x][i]);
void set_st(int x, int b) {
  st[x] = b;
  if (x > n) for (size_t i = 0; i < flo[x].size(); ++</pre>
      i) set_st(flo[x][i], b);
int get_pr(int b, int xr) {
  int pr = find(flo[b].begin(), flo[b].end(), xr) -
      flo[b].begin();
  if (pr % 2 == 1) {
    reverse(flo[b].begin() + 1, flo[b].end());
    return (int)flo[b].size() - pr;
  return pr;
void set_match(int u, int v) {
  match[u] = g[u][v].v;
  if (u <= n) return;</pre>
  edge e = g[u][v];
  int xr = flo_from[u][e.u], pr = get_pr(u, xr);
  for (int i = 0; i < pr; ++i) set_match(flo[u][i],
    flo[u][i ^ 1]);</pre>
  set_match(xr, v);
  rotate(flo[u].begin(), flo[u].begin() + pr, flo[u].
      end());
void augment(int u, int v) {
  for (; ; ) {
   int xnv = st[match[u]];
    set_match(u, v);
    if (!xnv) return;
    set_match(xnv, st[pa[xnv]]);
    u = st[pa[xnv]], v = xnv;
int get_lca(int u, int v) {
  static int t = 0;
  for (++t; u || v; swap(u, v)) {
    if (u == 0) continue;
    if (vis[u] == t) return u;
    vis[u] = t;
    u = st[match[u]];
   if (u) u = st[pa[u]];
  return 0:
void add_blossom(int u, int lca, int v) {
  int b = n + 1;
  while (b <= n_x && st[b]) ++b;</pre>
  if (b > n_x) ++n_x;
  lab[b] = 0, S[b] = 0;
  match[b] = match[lca];
  flo[b].clear();
  flo[b].push_back(lca);
  for (int x = u, y; x != lca; x = st[pa[y]])
    flo[b].push_back(x), flo[b].push_back(y = st[
```

match[x]]), q_push(y);

```
reverse(flo[b].begin() + 1, flo[b].end());
for (int x = v, y; x != lca; x = st[pa[y]])
    flo[b].push_back(x), flo[b].push_back(y = st[
        match[x]]), q_push(y);
  set_st(b, b);
  for (int x = 1; x <= n_x; ++x) g[b][x].w = g[x][b].
      w = 0;
  for (int x = 1; x <= n; ++x) flo_from[b][x] = 0;
  for (size_t i = 0; i < flo[b].size(); ++i) {</pre>
    int xs = flo[b][i];
    for (int x = 1; x <= n_x; ++x)</pre>
      if (g[b][x].w == 0 | e_delta(g[xs][x]) <</pre>
           e_delta(g[b][x]))
        g[b][x] = g[xs][x], g[x][b] = g[x][xs];
    for (int x = 1; x <= n; ++x)
      if (flo_from[xs][x]) flo_from[b][x] = xs;
  set_slack(b);
void expand_blossom(int b) {
  for (size_t i = 0; i < flo[b].size(); ++i)</pre>
    set_st(flo[b][i], flo[b][i]);
  int xr = flo_from[b][g[b][pa[b]].u], pr = get_pr(b,
       xr);
  for (int i = 0; i < pr; i += 2) {</pre>
    int xs = flo[b][i], xns = flo[b][i + 1];
    pa[xs] = g[xns][xs].u;
    S[xs] = 1, S[xns] = 0;
    slack[xs] = 0, set_slack(xns);
    q_push(xns);
  S[xr] = 1, pa[xr] = pa[b];
  for (size_t i = pr + 1; i < flo[b].size(); ++i) {</pre>
    int xs = flo[b][i];
    S[xs] = -1, set_slack(xs);
  st[b] = 0;
bool on_found_edge(const edge &e) {
  int u = st[e.u], v = st[e.v];
  if (S[v] == -1) {
    pa[v] = e.u, S[v] = 1;
    int nu = st[match[v]];
    slack[v] = slack[nu] = 0;
    S[nu] = 0, q_push(nu);
  } else if (S[v] == 0) {
  int lca = get_lca(u, v);
    if (!lca) return augment(u,v), augment(v,u), true
    else add_blossom(u, lca, v);
  return false;
bool matching() {
 memset(S + 1, -1, sizeof(int) * n_x);
memset(slack + 1, 0, sizeof(int) * n_x);
  q = queue<int>();
  for (int x = 1; x <= n_x; ++x)
    if (st[x] == x \&\& !match[x]) pa[x] = 0, S[x] = 0,
         q_push(x);
  if (q.empty()) return false;
  for (;;) {
    while (q.size()) {
      int u = q.front(); q.pop();
      if (S[st[u]] == 1) continue;
      for (int v = 1; v <= n; ++v)</pre>
        if (g[u][v].w > 0 && st[u] != st[v]) {
           if (e_delta(g[u][v]) == 0) {
             \quad \text{if } (on\_found\_edge(g[u][v])) \ \textit{return true}; \\
           } else update_slack(u, st[v]);
        }
    int d = inf;
    for (int b = n + 1; b <= n_x; ++b)</pre>
      if (st[b] == b && S[b] == 1) d = min(d, lab[b]
           / 2);
    for (int x = 1; x <= n_x; ++x)
      if (st[x] == x && slack[x]) {
        if (S[x] == -1) d = min(d, e_delta(g[slack[x
             ]][x]));
        else if (S[x] == 0) d = min(d, e_delta(g[
             slack[x]][x]) / 2);
```

```
for (int u = 1; u <= n; ++u) {</pre>
         if (S[st[u]] == 0) {
           if (lab[u] <= d) return 0;</pre>
           lab[u] -= d;
         } else if (S[st[u]] == 1) lab[u] += d;
       for (int b = n + 1; b <= n_x; ++b)
         if (st[b] == b) {
           if (S[st[b]] == 0) lab[b] += d * 2;
           else if (S[st[b]] == 1) lab[b] -= d * 2;
       q = queue<int>();
       for (int x = 1; x <= n_x; ++x)
        if (st[x] == x && slack[x] && st[slack[x]] != x
              && e_delta(g[slack[x]][x]) == 0)
           if (on_found_edge(g[slack[x]][x])) return
               true;
       for (int b = n + 1; b <= n_x; ++b)
         if (st[b] == b && S[b] == 1 && lab[b] == 0)
             expand_blossom(b);
    }
    return false:
  pair<long long, int> solve() {
    memset(match + 1, 0, sizeof(int) * n);
     n_x = n;
     int n_matches = 0;
     long long tot_weight = 0;
     for (int u = 0; u <= n; ++u) st[u] = u, flo[u].
         clear();
     int w_max = 0;
     for (int u = 1; u <= n; ++u)</pre>
       for (int v = 1; v <= n; ++v) {</pre>
         flo_from[u][v] = (u == v ? u : 0);
         w_max = max(w_max, g[u][v].w);
     for (int u = 1; u <= n; ++u) lab[u] = w_max;</pre>
     while (matching()) ++n_matches;
     for (int u = 1; u <= n; ++u)</pre>
       if (match[u] && match[u] < u)</pre>
         tot_weight += g[u][match[u]].w;
    return make_pair(tot_weight, n_matches);
  void add_edge(int ui, int vi, int wi) { g[ui][vi].w =
        g[vi][ui].w = wi; }
  void init(int _n) {
    n = _n;
     for (int u = 1; u <= n; ++u)</pre>
       for (int v = 1; v <= n; ++v)
         g[u][v] = edge(u, v, 0);
  }
};
```

3.8 Flow Model

- Maximum/Minimum flow with lower bound / Circulation problem
 - 1. Construct super source S and sink T.
 - 2. For each edge (x,y,l,u), connect x o y with capacity u-l. 3. For each vertex v, denote by in(v) the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.
 - 4. If in(v)>0, connect $S\to v$ with capacity in(v), otherwise, connect $v\to T$ with capacity -in(v).
 - To maximize, connect t o s with capacity ∞ (skip this in circulation problem), and let f be the maximum flow from S to T. If $f \neq \sum_{v \in V, in(v) > 0} in(v)$, there's no solution. Otherwise, the maximum flow from s to t is the answer.
 - To minimize, let f be the maximum flow from S to TConnect $t \to s$ with capacity ∞ and let the flow from Sto T be f' . If $f+f'
 eq \sum_{v \in V, in(v)>0} in(v)$, there's no solution. Otherwise, f^\prime is the answer.
 - 5. The solution of each edge e is l_e+f_e , where f_e corresponds to the flow of edge e on the graph.
- ullet Construct minimum vertex cover from maximum matching M on bipartite graph (X, Y)
 - 1. Redirect every edge: $y \to x$ if $(x,y) \in M$, $x \to y$ otherwise.

 - 2. DFS from unmatched vertices in X. 3. $x \in X$ is chosen iff x is unvisited. 4. $y \in Y$ is chosen iff y is visited.
- Maximum density induced subgraph
 - 1. Binary search on answer, suppose we're checking answer ${\it T}$ 2. Construct a max flow model, let K be the sum of all weights
 - 3. Connect source $s \to v$, $v \in G$ with capacity K

```
4. For each edge (u,v,w) in G, connect u \to v and v \to u with
   capacity w
5. For v \in G, connect it with sink v 	o t with capacity K + t
   2T - \left(\sum_{e \in E(v)} w(e)\right) - 2w(v)
```

6. T is a valid answer if the maximum flow f < K|V|

• Minimum weight edge cover

- 1. For each $v \in V$ create a copy v', and connect $u' \to v'$ with
- weight w(u,v) . 2. Connect $v\to v'$ with weight $2\mu(v)$, where $\mu(v)$ is the cost of the cheapest edge incident to v.
- 3. Find the minimum weight perfect matching on G^{\prime} .
- Project selection problem
 - 1. If $p_v>0$, create edge (s,v) with capacity p_v ; otherwise, create edge $\left(v,t\right)$ with capacity $-p_{v}$
 - 2. Create edge (u,v) with capacity w with w being the cost of
 - choosing u without choosing v. 3. The mincut is equivalent to the maximum profit of a subset of projects.
- 0/1 quadratic programming

$$\sum_{x} c_{x} x + \sum_{y} c_{y} \bar{y} + \sum_{xy} c_{xy} x \bar{y} + \sum_{xyx'y'} c_{xyx'y'} (x \bar{y} + x' \bar{y'})$$

can be minimized by the mincut of the following graph:

- 1. Create edge (x,t) with capacity c_x and create edge (s,y) with
- 2. Create edge (x,y) with capacity c_{xy}
- 3. Create edge (x,y) and edge (x',y') with capacity $c_{xyx'y'}$.

Graph

4.1 Heavy-Light Decomposition

```
vector<int> dep, pa, sz, ch, hd, id;
int _id;
void dfs(int i, int p) {
  dep[i] = \sim p ? dep[p] + 1 : 0;
  pa[i] = p, sz[i] = 1, ch[i] = -1;
  for (int j : g[i])
    if (j != p) {
      dfs(j, i);
      if (ch[i] == -1 || sz[ch[i]] < sz[j]) ch[i] = j;</pre>
      sz[i] += sz[j];
void hld(int i, int p, int h) {
  hd[i] = h;
  id[i] = _id++;
  if (~ch[i]) hld(ch[i], i, h);
  for (int j : g[i]) if (j != p && j != ch[i])
    hld(j, i, j);
void query(int i, int j) {
  while (hd[i] != hd[j]) {
    if (dep[hd[i]] < dep[hd[j]]) swap(i, j);</pre>
    query2(id[hd[i]], id[i] + 1), i = pa[hd[i]];
  if (dep[i] < dep[j]) swap(i, j);</pre>
  query2(id[j], id[i] + 1);
```

4.2 Centroid Decomposition

```
vector<vector<int>> dis; // dis[n][logn]
vector<int> pa, sz, dep;
vector<bool> vis;
void dfs_sz(int i, int p) {
  sz[i] = 1;
  for (int j : g[i]) if (j != p && !vis[j])
    dfs_sz(j, i), sz[i] += sz[j];
int cen(int i, int p, int _n) {
 for (int j : g[i]) if (j != p && !vis[j] && sz[j] >
      _n / 2)
    return cen(j, i, _n);
 return i;
void dfs_dis(int i, int p, int d) { // from i to
    ancestor with depth d
  dis[i][d] = \sim p ? dis[p][d] + 1 : 0;
 for (int j : g[i]) if (j != p && !vis[j])
```

```
dfs_dis(j, i, d);
void cd(int i, int p, int d) {
 dfs_sz(i, -1), i = cen(i, -1, sz[i]);
  vis[i] = true, pa[i] = p, dep[i] = d;
  dfs_dis(i, -1, d);
  for (int j : g[i]) if (!vis[j])
    cd(j, i, d + 1);
```

4.3 Edge BCC

```
vector<int> low, dep, bcc_id, stk;
vector<bool> vis;
int _id;
void dfs(int i, int p) {
  low[i] = dep[i] = \sim p ? dep[p] + 1 : 0;
  stk.push_back(i);
  vis[i] = true;
  for (int j : g[i])
    if (j != p)
      if (!vis[j])
        dfs(j, i), low[i] = min(low[i], low[j]);
        low[i] = min(low[i], dep[j]);
  if (low[i] == dep[i]) {
    int id =
              _id++;
    while (stk.back() != i) {
      int x = stk.back();
      stk.pop_back();
      bcc_id[x] = id;
    stk.pop_back();
    bcc_id[i] = id;
}
```

4.4 Block Cut Tree

```
vector<vector<int>> g, _g;
vector<int> dep, low, stk;
void dfs(int i, int p) {
  dep[i] = low[i] = \sim p ? dep[p] + 1 : 0;
  stk.push_back(i);
  for (int j : g[i]) if (j != p) {
    if (dep[j] == -1) {
      dfs(j, i), low[i] = min(low[i], low[j]);
      if (low[j] >= dep[i]) {
        int id = _g.size();
         _g.emplace_back();
        while (stk.back() != j) {
          int x = stk.back();
           stk.pop_back();
          _g[x].push_back(id), _g[id].push_back(x);
        stk.pop_back();
         _g[j].push_back(id), _g[id].push_back(j);
         _g[i].push_back(id), _g[id].push_back(i);
    } else low[i] = min(low[i], dep[j]);
}
```

4.5 SCC / 2SAT

```
struct SAT {
  vector<vector<int>> g;
  vector<int> dep, low, scc_id;
  vector<bool> is;
  vector<int> stk;
  int n, _id, _t;
  SAT() {}
  void init(int _n) {
    n = _n, _id = _t = 0;
    g.assign(2 * n, vector<int>());
    dep.assign(2 * n, -1), low.assign(2 * n, -1);
scc_id.assign(2 * n, -1), is.assign(2 * n, false);
    stk.clear();
  void add_edge(int x, int y) {g[x].push_back(y);}
  int rev(int i) {return i < n ? i + n : i - n;}</pre>
```

```
void add_ifthen(int x, int y) {add_clause(rev(x), y)
  void add_clause(int x, int y) {
    add_edge(rev(x), y);
    add_edge(rev(y), x);
  void dfs(int i) {
    dep[i] = low[i] = _t++;
    stk.push_back(i);
    for (int j : g[i])
      if (scc_id[j] == -1) {
        if (dep[j] == -1)
          dfs(j);
        low[i] = min(low[i], low[j]);
    if (low[i] == dep[i]) {
      int id = _id++;
      while (stk.back() != i) {
        int x = stk.back();
        stk.pop_back();
        scc_id[x] = id;
      }
      stk.pop_back();
      scc_id[i] = id;
   }
  bool solve() {
    for (int i = 0; i < 2 * n; ++i)
      if (dep[i] == -1)
        dfs(i);
    for (int i = 0; i < n; ++i) {</pre>
      if (scc_id[i] == scc_id[i + n]) return false;
      if (scc_id[i] < scc_id[i + n])</pre>
        is[i] = true;
      else
        is[i + n] = true;
    return true;
 }
};
```

4.6 Negative Cycle*

```
vector <pair <int, long long>> adj[N];
template <typename T>
struct NegativeCycle {
  vector <T> dis:
  vector <int> rt;
  int n; T INF;
  vector <int> cycle;
  NegativeCycle () = default;
  NegativeCycle \ ( \mbox{int } \mbox{\_n} ) \ : \ n(\mbox{\_n}) \ , \ INF(numeric\_limits < T
       >::max()) {
     dis.assign(n, 0), rt.assign(n, -1);
    int relax = -1;
    for (int t = 0; t < n; ++t) {</pre>
       relax = -1;
       for (int i = 0; i < n; ++i) {</pre>
         for (auto [j, w] : adj[i]) if (dis[j] > dis[i]
             + w) {
           dis[j] = dis[i] + w, rt[j] = i;
           relax = j;
         }
       }
    if (relax != -1) {
       int s = relax;
       for (int i = 0; i < n; ++i) s = rt[s];</pre>
       vector <bool> vis(n, false);
       while (!vis[s]) {
         cycle.push_back(s), vis[s] = true;
         s = rt[s];
       reverse(cycle.begin(), cycle.end());
  }
};
```

4.7 Virtual Tree

```
vector<vector<int>> _g;
vector<int> st, ed, stk;
```

```
void solve(vector<int> v) {
   sort(all(v), [&](int x, int y) {return st[x] < st[y</pre>
       ];});
   int sz = v.size();
   for (int i = 0; i < sz - 1; ++i)</pre>
     v.push_back(lca(v[i], v[i + 1]));
   sort(all(v), [\&](int x, int y) \{return st[x] < st[y]\}
       ];});
   v.resize(unique(all(v)) - v.begin());
   stk.clear(); stk.push_back(v[0]);
   for (int i = 1; i < v.size(); ++i)</pre>
     int x = v[i];
     while (ed[stk.back()] < ed[x]) stk.pop_back();</pre>
     _g[stk.back()].push_back(x), stk.push_back(x);
   // do something
   for (int i : v) _g[i].clear();
```

4.8 Directed MST

vis[now] = true;

```
template <typename T> struct DMST { // 1-based
  T g[maxn][maxn], fw[maxn];
  int n, fr[maxn];
  bool vis[maxn], inc[maxn];
  void clear() {
    for (int i = 0; i < maxn; ++i) {</pre>
      for (int j = 0; j < maxn; ++j) g[i][j] = inf;</pre>
      vis[i] = inc[i] = false;
    }
  void addedge(int u, int v, T w) {
    g[u][v] = min(g[u][v], w);
  T query(int root, int _n) {
    if (dfs(root) != n) return -1;
    T ans = 0;
    while (true) {
      for (int i = 1; i <= n; ++i) fw[i] = inf, fr[i] =</pre>
      for (int i = 1; i <= n; ++i) if (!inc[i]) {</pre>
        for (int j = 1; j <= n; ++j) {</pre>
          if (!inc[j] && i != j && g[j][i] < fw[i]) {</pre>
            fw[i] = g[j][i];
            fr[i] = j;
          }
        }
      int x = -1;
      for (int i = 1; i <= n; ++i) if (i != root && !</pre>
          inc[i]) {
        int j = i, c = 0;
        while (j != root && fr[j] != i && c <= n) ++c,</pre>
            j = fr[j];
        if (j == root || c > n) continue;
        else { x = i; break; }
      if (!~x) {
        for (int i = 1; i <= n; ++i) if (i != root && !</pre>
            inc[i]) ans += fw[i];
        return ans;
      int y = x;
      for (int i = 1; i <= n; ++i) vis[i] = false;</pre>
      do { ans += fw[y]; y = fr[y]; vis[y] = inc[y] =
           true; } while (y != x);
      inc[x] = false;
      for (int k = 1; k <= n; ++k) if (vis[k]) {</pre>
        for (int j = 1; j <= n; ++j) if (!vis[j]) {</pre>
            if (g[x][j] > g[k][j]) g[x][j] = g[k][j];
            ][x]) g[j][x] = g[j][k] - fw[k];
          }
      }
    return ans;
  int dfs(int now) {
    int r = 1;
```

```
for (int i = 1; i <= n; ++i) if (g[now][i] < inf &&
    !vis[i]) r += dfs(i);
return r;
}
};</pre>
```

4.9 Dominator Tree

```
struct Dominator_tree {
  int n, id;
  vector <vector <int>> adj, radj, bucket;
  vector <int> sdom, dom, vis, rev, par, rt, mn;
  Dominator_tree (int _n) : n(_n), id(0) {
    adj.resize(n), radj.resize(n), bucket.resize(n);
    sdom.resize(n), dom.resize(n, -1), vis.resize(n,
         -1);
    rev.resize(n), rt.resize(n), mn.resize(n), par.
        resize(n);
  void add_edge(int u, int v) {adj[u].pb(v);}
  int query(int v, bool x) {
    if (rt[v] == v) return x ? -1 : v;
    int p = query(rt[v], true);
    if (p == -1) return x ? rt[v] : mn[v];
    if (sdom[mn[v]] > sdom[mn[rt[v]]]) mn[v] = mn[rt[v
        ]];
    rt[v] = p;
    return x ? p : mn[v];
  void dfs(int v) {
    vis[v] = id, rev[id] = v;
    rt[id] = mn[id] = sdom[id] = id, id++;
    for (int u : adj[v]) {
      if (vis[u] == -1) dfs(u), par[vis[u]] = vis[v];
      radj[vis[u]].pb(vis[v]);
    }
  void build(int s) {
    dfs(s);
    for (int i = id - 1; ~i; --i) {
      for (int u : radj[i]) {
        sdom[i] = min(sdom[i], sdom[query(u, false)]);
      if (i) bucket[sdom[i]].pb(i);
      for (int u : bucket[i]) {
        int p = query(u, false);
        dom[u] = sdom[p] == i ? i : p;
      if (i) rt[i] = par[i];
    }
    vector <int> res(n, -1);
    for (int i = 1; i < id; ++i) {</pre>
      if (dom[i] != sdom[i]) dom[i] = dom[dom[i]];
    for (int i = 1; i < id; ++i) res[rev[i]] = rev[dom[</pre>
        i]];
    res[s] = s;
    dom = res;
|};
```

5 String

5.1 Aho-Corasick Automaton

```
struct AC {
   int ch[N][26], to[N][26], fail[N], sz;
   vector <int> g[N];
   int cnt[N];
   AC () {sz = 0, extend();}
   void extend() {fill(ch[sz], ch[sz] + 26, 0), sz++;}
   int nxt(int u, int v) {
     if (!ch[u][v]) ch[u][v] = sz, extend();
     return ch[u][v];
   }
   int insert(string s) {
     int now = 0;
     for (char c : s) now = nxt(now, c - 'a');
     cnt[now]++;
     return now;
   }
```

```
void build_fail() {
    queue <int> q;
    for (int i = 0; i < 26; ++i) if (ch[0][i]) {
      q.push(ch[0][i]);
      g[0].push_back(ch[0][i]);
    while (!q.empty()) {
      int v = q.front(); q.pop();
      for (int j = 0; j < 26; ++j) {</pre>
        to[v][j] = ch[v][j] ? v : to[fail[v]][j];
      for (int i = 0; i < 26; ++i) if (ch[v][i]) {</pre>
         int u = ch[v][i], k = fail[v];
        while (k && !ch[k][i]) k = fail[k];
        if (ch[k][i]) k = ch[k][i];
         fail[u] = k;
        cnt[u] += cnt[k], g[k].push_back(u);
        q.push(u);
      }
    }
  int match(string &s) {
    int now = 0, ans = 0;
for (char c : s) {
      now = to[now][c - 'a'];
      if (ch[now][c - 'a']) now = ch[now][c - 'a'];
      ans += cnt[now];
    }
    return ans;
};
```

5.2 KMP Algorithm

```
vector <int> build_fail(string s) {
  vector <int> f(s.length() + 1, 0);
  int k = 0;
  for (int i = 1; i < s.length(); ++i) {</pre>
    while (k && s[k] != s[i]) k = f[k];
    if (s[k] == s[i]) k++;
    f[i + 1] = k;
  return f;
int match(string s, string t) {
  vector <int> f = build_fail(t);
  int k = 0, ans = 0;
  for (int i = 0; i < s.length(); ++i) {</pre>
    while (k && s[i] != t[k]) k = f[k];
    if (s[i] == t[k]) k++;
    if (k == t.length()) ans++, k = f[k];
  return ans;
```

5.3 Z Algorithm

5.4 Manacher

```
vector <int> manacher(string &s) {
   string t = "^#";
   for (char c : s) t += c, t += '#';
   t += '&';
   int n = t.length();
   vector <int> r(n, 0);
   int C = 0, R = 0;
   for (int i = 1; i < n - 1; ++i) {
     int mirror = 2 * C - i;
}</pre>
```

```
r[i] = (i < R ? min(r[mirror], R - i) : 0);
while (t[i - 1 - r[i]] == t[i + 1 + r[i]]) r[i]++;
if (i + r[i] > R) R = i + r[i], C = i;
}
return r;
}
```

5.5 Suffix Array

```
int sa[N], tmp[2][N], c[N], rk[N], lcp[N];
void buildSA(string s) {
  int *x = tmp[0], *y = tmp[1], m = 256, n = s.length()
  for (int i = 0; i < m; ++i) c[i] = 0;</pre>
  for (int i = 0; i < n; ++i) c[x[i] = s[i]]++;</pre>
  for (int i = 1; i < m; ++i) c[i] += c[i - 1];</pre>
  for (int i = n - 1; ~i; --i) sa[--c[x[i]]] = i;
  for (int k = 1; k < n; k <<= 1) {</pre>
    for (int i = 0; i < m; ++i) c[i] = 0;</pre>
    for (int i = 0; i < n; ++i) c[x[i]]++;</pre>
    for (int i = 1; i < m; ++i) c[i] += c[i - 1];</pre>
    int p = 0;
    for (int i = n - k; i < n; ++i) y[p++] = i;
    for (int i = 0; i < n; ++i) if (sa[i] >= k) y[p++]
         = sa[i] - k;
    for (int i = n - 1; \sim i; --i) sa[--c[x[y[i]]]] = y[i]
    y[sa[0]] = p = 0;
    for (int i = 1; i < n; ++i) {</pre>
      int a = sa[i], b = sa[i - 1];
      if (!(x[a] == x[b] \&\& a + k < n \&\& b + k < n \&\& x)
           [a + k] == x[b + k])) p++;
      y[sa[i]] = p;
    if (n == p + 1) break;
    swap(x, y), m = p + 1;
void buildLCP(string s) {
  // lcp[i] = LCP(sa[i - 1], sa[i])
  // lcp(i, j) = min(lcp[rk[i] + 1], lcp[rk[i] + 2],
..., lcp[rk[j]])
  int n = s.length(), val = 0;
  for (int i = 0; i < n; ++i) rk[sa[i]] = i;</pre>
  for (int i = 0; i < n; ++i) {</pre>
    if (!rk[i]) lcp[rk[i]] = 0;
    else {
       if (val) val--;
       int p = sa[rk[i] - 1];
       while (val + i < n && val + p < n && s[val + i]</pre>
           == s[val + p]) val++;
      lcp[rk[i]] = val;
  }
}
```

5.6 SAIS

```
namespace sfx {
bool _t[N * 2];
int SA[N * 2], H[N], RA[N];
int _s[N * 2], _c[N * 2], x[N], _p[N], _q[N * 2];
void pre(int *sa, int *c, int n, int z) {
  fill_n(sa, n, 0), copy_n(c, z, x);
void induce(int *sa, int *c, int *s, bool *t, int n,
    int z) {
  copy_n(c, z - 1, x + 1);
for (int i = 0; i < n; ++i) if (sa[i] && !t[sa[i] -</pre>
       1]) sa[x[s[sa[i] - 1]]++] = sa[i] - 1;
  copy_n(c, z, x);
  for (int i = n - 1; i >= 0; --i) if (sa[i] && t[sa[i]
         - 1]) sa[--x[s[sa[i] - 1]]] = sa[i] - 1;
void sais(int *s, int *sa, int *p, int *q, bool *t, int
      *c, int n, int z) {
  bool uniq = t[n - 1] = true;
  int nn = 0, nmxz = -1, *nsa = sa + n, *ns = s + n,
       last = -1:
  fill_n(c, z, 0);
  for (int i = 0; i < n; ++i) uniq &= ++c[s[i]] < 2;</pre>
  partial_sum(c, c + z, c);
```

```
if (uniq) {
     for (int i = 0; i < n; ++i) sa[--c[s[i]]] = i;</pre>
   for (int i = n - 2; i >= 0; --i)
     t[i] = (s[i] == s[i + 1] ? t[i + 1] : s[i] < s[i +
         1]);
   pre(sa, c, n, z);
   for (int i = 1; i <= n - 1; ++i)</pre>
     if (t[i] && !t[i - 1])
       sa[--x[s[i]]] = p[q[i] = nn++] = i;
   induce(sa, c, s, t, n, z);
   for (int i = 0; i < n; ++i)</pre>
     if (sa[i] && t[sa[i]] && !t[sa[i] - 1]) {
       bool neq = last < 0 || !equal(s + sa[i], s + p[q[</pre>
           sa[i]] + 1], s + last);
       ns[q[last = sa[i]]] = nmxz += neq;
   sais(ns, nsa, p + nn, q + n, t + n, c + z, nn, nmxz +
        1);
   pre(sa, c, n, z);
   for (int i = nn - 1; i >= 0; --i)
     sa[--x[s[p[nsa[i]]]] = p[nsa[i]];
   induce(sa, c, s, t, n, z);
vector<int> build(int *s, int n) {
   copy_n(s, n, _s), _s[n] = 0;
   sais(_s, SA, _p, _q, _t, _c, n + 1, 256);
   vector <int> sa(n);
   for (int i = 0; i < n; ++i)</pre>
     sa[i] = SA[i + 1];
   return sa;
}
}
```

5.7 Suffix Automaton

```
struct SAM {
  int ch[N][26], len[N], link[N], cnt[N], sz;
  // link -> suffix endpos
  SAM () \{len[0] = 0, link[0] = -1, sz = 1;\}
  void build(string s) {
    int last = 0;
    for (char c : s) {
      int cur = sz++;
      len[cur] = len[last] + 1;
      int p = last;
      while (\sim p \&\& !ch[p][c - 'a']) ch[p][c - 'a'] =
           cur, p = link[p];
      if (p == -1) {
        link[cur] = 0;
      } else {
        int q = ch[p][c - 'a'];
        if (len[p] + 1 == len[q]) {
           link[cur] = q;
           int nxt = sz++;
          len[nxt] = len[p] + 1, link[nxt] = link[q];
           for (int j = 0; j < 26; ++j) ch[nxt][j] = ch[</pre>
               q][j];
           while (\sim p \&\& ch[p][c - 'a'] == q) ch[p][c - '
               a'] = nxt, p = link[p];
          link[q] = link[cur] = nxt;
        }
      }
      cnt[cur]++;
      last = cur;
    vector <int> p(sz);
    iota(all(p), 0);
    sort(all(p), [&](int i, int j) {return len[i] > len
        [j];});
    for (int i = 0; i < sz; ++i) cnt[link[p[i]]] += cnt</pre>
        [p[i]];
};
```

5.8 Minimum Rotation

```
string rotate(const string &s) {
  int n = s.length();
  string t = s + s;
```

```
int i = 0, j = 1;
while (i < n && j < n) {
  int k = 0;
while (k < n && t[i + k] == t[j + k]) ++k;
  if (t[i + k] <= t[j + k]) j += k + 1;
  else i += k + 1;
  if (i == j) ++j;
}
int pos = (i < n ? i : j);
return t.substr(pos, n);
}</pre>
```

5.9 Palindrome Tree

```
struct PAM {
  int ch[N][26], cnt[N], fail[N], len[N], sz;
  // 0 -> even root, 1 -> odd root
  PAM (string _s) : s(_s) {
    sz = 0;
    extend(), extend();
    len[0] = 0, fail[0] = 1, len[1] = -1;
    int lst = 1;
    for (int i = 0; i < s.length(); ++i) {</pre>
      while (s[i - len[lst] - 1] != s[i]) lst = fail[
          lst];
      if (!ch[lst][s[i] - 'a']) {
        int idx = extend();
        len[idx] = len[lst] + 2;
        int now = fail[lst];
        while (s[i - len[now] - 1] != s[i]) now = fail[
            now];
        fail[idx] = ch[now][s[i] - 'a'];
        ch[lst][s[i] - 'a'] = idx;
      lst = ch[lst][s[i] - 'a'], cnt[lst]++;
    }
  }
  void build_count() {
    for (int i = sz - 1; i > 1; --i)
      cnt[fail[i]] += cnt[i];
  int extend() {
    fill(ch[sz], ch[sz] + 26, 0), sz++;
    return sz - 1;
 }
};
```

5.10 Main Lorentz

```
int to_left[N], to_right[N];
vector <array <int, 3>> rep; // l, r, len.
// substr(l ~ r, len * 2) are tandem
void findRep(string &s, int 1, int r) {
 if (r - 1 == 1) return;
  int m = 1 + r >> 1;
 findRep(s, 1, m), findRep(s, m, r);
  string sl = s.substr(1, m - 1), sr = s.substr(m, r - 1)
      m);
  vector <int> Z = buildZ(sr + "#" + sl);
  for (int i = 1; i < m; ++i) to_right[i] = Z[r - m + 1</pre>
       + i - 1];
  reverse(all(sl)):
  Z = buildZ(sl);
  for (int i = 1; i < m; ++i) to_left[i] = Z[m - i -</pre>
      1];
  reverse(all(sl));
  for (int i = 1; i + 1 < m; ++i) {</pre>
    int k1 = to_left[i], k2 = to_right[i + 1], len = m
         - i - 1;
    if (k1 < 1 || k2 < 1 || len < 2) continue;
    int tl = max(1, len - k2), tr = min(len - 1, k1);
    if (tl <= tr) rep.pb({i + 1 - tr, i + 1 - tl, len})</pre>
  Z = buildZ(sr);
  for (int i = m; i < r; ++i) to_right[i] = Z[i - m];</pre>
 reverse(all(sl)), reverse(all(sr));
Z = buildZ(sl + "#" + sr);
  for (int i = m; i < r; ++i) to_left[i] = Z[m - l + 1</pre>
      + r - i - 1];
  reverse(all(sl)), reverse(all(sr));
```

6 Math

6.1 Fraction*

```
struct fraction {
  11 n, d;
  fraction(const ll _n=0, const ll _d=1): n(_n), d(_d)
     11 t = gcd(n, d);
    n /= t, d /= t;
    if (d < 0) n = -n, d = -d;
  fraction operator-() const
  { return fraction(-n, d);
  fraction operator+(const fraction &b) const
  { return fraction(n * b.d + b.n * d, d * b.d); }
  fraction operator-(const fraction &b) const
  { return fraction(n * b.d - b.n * d, d * b.d); }
  fraction operator*(const fraction &b) const
  { return fraction(n * b.n, d * b.d); }
  fraction operator/(const fraction &b) const
  { return fraction(n * b.d, d * b.n); }
  void print() {
    cout << n;
     if (d != 1) cout << "/" << d;</pre>
  }
};
```

6.2 Miller Rabin / Pollard Rho

```
11 mul(11 x, 11 y, 11 p) {return (x * y - (11)((long
    double)x / p * y) * p + p) % p;}
vector<ll> chk = {2, 325, 9375, 28178, 450775, 9780504,
     1795265022};
11 Pow(11 a, 11 b, 11 n) {11 res = 1; for (; b; b >>=
    1, a = mul(a, a, n)) if (b \& 1) res = mul(res, a, n)
    ); return res;}
bool check(ll a, ll d, int s, ll n) {
  a = Pow(a, d, n);
  if (a <= 1) return 1;</pre>
  for (int i = 0; i < s; ++i, a = mul(a, a, n)) {</pre>
    if (a == 1) return 0;
    if (a == n - 1) return 1;
  return 0;
bool IsPrime(ll n) {
  if (n < 2) return 0;
  if (n % 2 == 0) return n == 2;
  11 d = n - 1, s = 0;
  while (d % 2 == 0) d >>= 1, ++s;
  for (ll i : chk) if (!check(i, d, s, n)) return 0;
  return 1;
const vector<ll> small = {2, 3, 5, 7, 11, 13, 17, 19};
11 FindFactor(11 n) {
  if (IsPrime(n)) return 1;
  for (11 p : small) if (n % p == 0) return p;
  11 x, y = 2, d, t = 1;
  auto f = [&](11 a) {return (mul(a, a, n) + t) % n;};
  for (int 1 = 2; ; 1 <<= 1) {
   x = y;
    int m = min(1, 32);
    for (int i = 0; i < 1; i += m) {</pre>
      d = 1;
```

```
for (int j = 0; j < m; ++j) {</pre>
        y = f(y), d = mul(d, abs(x - y), n);
      ll g = \_gcd(d, n);
      if (g == n) {
        1 = 1, y = 2, ++t;
        break:
      if (g != 1) return g;
   }
 }
}
map <ll, int> res;
void PollardRho(ll n) {
 if (n == 1) return;
 if (IsPrime(n)) return ++res[n], void(0);
 11 d = FindFactor(n);
 PollardRho(n / d), PollardRho(d);
```

6.3 Ext GCD

```
//a * p.first + b * p.second = gcd(a, b)
pair<ll, ll> extgcd(ll a, ll b) {
  pair<ll, ll> res;
  if (a < 0) {
    res = extgcd(-a, b);
    res.first *= -1;
    return res;
  }
  if (b < 0) {
    res = extgcd(a, -b);
    res.second *= -1;
    return res;
  }
  if (b == 0) return {1, 0};
  res = extgcd(b, a % b);
  return {res.second, res.first - res.second * (a / b)
    };
}</pre>
```

6.4 PiCount

```
const int V = 10000000, N = 100, M = 100000;
vector<int> primes;
bool isp[V];
int small_pi[V], dp[N][M];
void sieve(int x){
  for(int i = 2; i < x; ++i) isp[i] = true;</pre>
 isp[0] = isp[1] = false;
  for(int i = 2; i * i < x; ++i) if(isp[i]) for(int j =</pre>
       i * i; j < x; j += i) isp[j] = false;</pre>
  for(int i = 2; i < x; ++i) if(isp[i]) primes.</pre>
      push_back(i);
void init(){
 sieve(V);
  small_pi[0] = 0;
  for(int i = 1; i < V; ++i) small_pi[i] = small_pi[i -</pre>
       1] + isp[i];
 for(int i = 0; i < M; ++i) dp[0][i] = i;</pre>
 - 1]];
11 phi(ll n, int a){
 if(!a) return n;
  if(n < M && a < N) return dp[a][n];</pre>
  if(primes[a - 1] > n) return 1;
  if(((ll)primes[a - 1]) * primes[a - 1] >= n && n < V)</pre>
       return small_pi[n] - a + 1;
  11 de = phi(n, a - 1) - phi(n / primes[a - 1], a - 1)
  return de;
11 PiCount(11 n){
  if(n < V) return small_pi[n];</pre>
  int s = sqrt(n + 0.5), y = cbrt(n + 0.5), a =
      small_pi[y];
  ll res = phi(n, a) + a - 1;
  for(; primes[a] <= s; ++a) res -= max(PiCount(n /</pre>
```

primes[a]) - PiCount(primes[a]) + 1, 0ll);

```
return res;
```

```
6.5 Linear Function Mod Min
ll topos(ll x, ll m) {x %= m; if (x < 0) x += m; return
     x;}
//min value of ax + b \pmod{m} for x \in [0, n - 1]. O(
    Log m)
ll min_rem(ll n, ll m, ll a, ll b) {
  a = topos(a, m), b = topos(b, m);
  for (ll g = \_gcd(a, m); g > 1;) return g * min_rem(n)
      , m / g, a / g, b / g) + (b % g);
  for (11 nn, nm, na, nb; a; n = nn, m = nm, a = na, b
      = nb) {
    if (a <= m - a) {
      nn = (a * (n - 1) + b) / m;
      if (!nn) break;
      nn += (b < a);
      nm = a, na = topos(-m, a);
      nb = b < a ? b : topos(b - m, a);
    } else {
      11 lst = b - (n - 1) * (m - a);
      if (lst >= 0) {b = lst; break;}
nn = -(lst / m) + (lst % m < -a) + 1;</pre>
      nm = m - a, na = m % (m - a), nb = b % (m - a);
 }
  return b;
}
//min value of ax + b \pmod{m} for x \in [0, n - 1],
    also return min x to get the value. O(\log m)
//{value, x}
pair<11, 11> min_rem_pos(11 n, 11 m, 11 a, 11 b) {
  a = topos(a, m), b = topos(b, m);
  ll mn = min_rem(n, m, a, b), g = __gcd(a, m);
  //ax = (mn - b) \pmod{m}
  11 x = (extgcd(a, m).first + m) * ((mn - b + m) / g)
```

6.6 Determinant

}

% (m / g);
return {mn, x};

```
11 Det(vector <vector <11>>> a) {
  int n = a.size();
  ll det = 1;
  for (int i = 0; i < n; ++i) {</pre>
    if (!a[i][i]) {
       det = -det;
       if (det < 0) det += mod;</pre>
       for (int j = i + 1; j < n; ++j) if (a[j][i]) {</pre>
         swap(a[j], a[i]);
         break:
       if (!a[i][i]) return 0;
    det = det * a[i][i] % mod;
     11 mul = mpow(a[i][i], mod - 2);
    for (int j = 0; j < n; ++j) a[i][j] = a[i][j] * mul</pre>
          % mod;
    for (int j = 0; j < n; ++j) if (i ^ j) {
    ll mul = a[j][i];</pre>
       for (int k = 0; k < n; ++k) {
         a[j][k] -= a[i][k] * mul % mod;
         if (a[j][k] < 0) a[j][k] += mod;</pre>
      }
    }
  return det;
```

6.7 Floor Sum

```
ans += (n - (x_max + a - 1) / a) * y_max;
ans += floor_sum(y_max, a, m, (a - x_max % a) % a);
return ans;
}
```

6.8 Quadratic Residue

```
int Jacobi(int a, int m) {
 int s = 1;
  for (; m > 1; ) {
   a %= m;
   if (a == 0) return 0;
    const int r = __builtin_ctz(a);
    if ((r \& 1) \&\& ((m + 2) \& 4)) s = -s;
    a >>= r;
    if (a \& m \& 2) s = -s;
   swap(a, m);
 return s;
int QuadraticResidue(int a, int p) {
 if (p == 2) return a & 1;
  const int jc = Jacobi(a, p);
  if (jc == 0) return 0;
 if (jc == -1) return -1;
 int b, d;
  for (; ; )
   b = rand() % p;
d = (1LL * b * b + p - a) % p;
   if (Jacobi(d, p) == -1) break;
  int f0 = b, f1 = 1, g0 = 1, g1 = 0, tmp;
 for (int e = (1LL + p) >> 1; e; e >>= 1) {
   if (e & 1) {
      tmp = (1LL * g0 * f0 + 1LL * d * (1LL * g1 * f1 %
          p)) % p;
      g1 = (1LL * g0 * f1 + 1LL * g1 * f0) % p;
      g0 = tmp;
    tmp = (1LL * f0 * f0 + 1LL * d * (1LL * f1 * f1 % p)
       )) % p;
    f1 = (2LL * f0 * f1) % p;
    f0 = tmp;
 }
  return g0;
```

6.9 Simplex

```
struct Simplex { // 0-based
  using T = long double;
  static const int N = 410, M = 30010;
  const T eps = 1e-7;
  int n, m;
  int Left[M], Down[N];
  // Ax <= b, max c^T x
  // result : v, xi = sol[i]
  T a[M][N], b[M], c[N], v, sol[N];
  bool eq(T a, T b) {return fabs(a - b) < eps;}
bool ls(T a, T b) {return a < b && !eq(a, b);}</pre>
  void init(int _n, int _m) {
    n = _n, m = _m, v = 0;
    for (int i = 0; i < m; ++i) for (int j = 0; j < n;</pre>
         ++j) a[i][j] = 0;
    for (int i = 0; i < m; ++i) b[i] = 0;
for (int i = 0; i < n; ++i) c[i] = sol[i] = 0;</pre>
  void pivot(int x, int y) {
    swap(Left[x], Down[y]);
    T k = a[x][y]; a[x][y] = 1;
    vector <int> nz;
    for (int i = 0; i < n; ++i) {</pre>
       a[x][i] /= k;
       if (!eq(a[x][i], 0)) nz.push_back(i);
    b[x] /= k;
     for (int i = 0; i < m; ++i) {</pre>
       if (i == x || eq(a[i][y], 0)) continue;
       k = a[i][y], a[i][y] = 0;
       b[i] -= k * b[x];
       for (int j : nz) a[i][j] -= k * a[x][j];
```

```
if (eq(c[y], 0)) return;
  k = c[y], c[y] = 0, v += k * b[x];
  for (int i : nz) c[i] -= k * a[x][i];
// 0: found solution, 1: no feasible solution, 2:
    unbounded
int solve() {
  for (int i = 0; i < n; ++i) Down[i] = i;</pre>
  for (int i = 0; i < m; ++i) Left[i] = n + i;</pre>
  while (1) {
    int x = -1, y = -1;
    for (int i = 0; i < m; ++i) if (ls(b[i], 0) && (x</pre>
         == -1 \mid \mid b[i] < b[x])) x = i;
    if (x == -1) break;
    for (int i = 0; i < n; ++i) if (ls(a[x][i], 0) &&</pre>
         (y == -1 \mid | a[x][i] < a[x][y])) y = i;
    if (y == -1) return 1;
    pivot(x, y);
  while (1) {
    int x = -1, y = -1;
    for (int i = 0; i < n; ++i) if (ls(0, c[i]) && (y
         == -1 \mid \mid c[i] > c[y])) y = i;
    if (y == -1) break;
    for (int i = 0; i < m; ++i) if (ls(0, a[i][y]) &&</pre>
         (x == -1 \mid | b[i] / a[i][y] < b[x] / a[x][y])
         ) x = i;
    if (x == -1) return 2;
   pivot(x, y);
  for (int i = 0; i < m; ++i) if (Left[i] < n) sol[</pre>
      Left[i]] = b[i];
  return 0:
```

6.10 Berlekamp Massey

```
vector <11> BerlekampMassey(vector <11> a) {
  // find min |c| such that a_n = sum c_j * a_{n - j - 1}
      1}, 0-based
  // O(N^2), if |c| = k, |a| >= 2k sure correct
  auto f = [&](vector<ll> v, ll c) {
    for (11 &x : v) x = mul(x, c);
    return v;
 };
  vector <11> c, best;
  int pos = 0, n = a.size();
  for (int i = 0; i < n; ++i) {</pre>
    ll error = a[i];
    for (int j = 0; j < c.size(); ++j) error = sub(</pre>
        error, mul(c[j], a[i - 1 - j]));
    if (error == 0) continue;
    11 inv = mpow(error, mod - 2);
    if (c.empty()) {
      c.resize(i + 1);
      pos = i;
      best.pb(inv);
    } else {
      vector <1l> fix = f(best, error);
      fix.insert(fix.begin(), i - pos - 1, 0);
      if (fix.size() >= c.size()) {
        best = f(c, sub(0, inv));
        best.insert(best.begin(), inv);
        pos = i;
        c.resize(fix.size());
      for (int j = 0; j < fix.size(); ++j) c[j] = add(c</pre>
          [j], fix[j]);
   }
  }
  return c;
```

6.11 Linear Programming Construction

Standard form: maximize $\mathbf{c}^T\mathbf{x}$ subject to $A\mathbf{x} \leq \mathbf{b}$ and $\mathbf{x} \geq 0$. Dual LP: minimize $\mathbf{b}^T\mathbf{y}$ subject to $A^T\mathbf{y} \geq \mathbf{c}$ and $\mathbf{y} \geq 0$. $\bar{\mathbf{x}}$ and $\bar{\mathbf{y}}$ are optimal if and only if for all $i \in [1,n]$, either $\bar{x}_i = 0$ or $\sum_{j=1}^m A_{ji}\bar{y}_j = c_i$ holds and for all $i \in [1,m]$ either $\bar{y}_i = 0$ or $\sum_{j=1}^n A_{ij}\bar{x}_j = b_j$ holds.

1. In case of minimization, let $c_i^\prime = -c_i$

- 2. $\sum_{1 \leq i \leq n} A_{ji} x_i \geq b_j \to \sum_{1 \leq i \leq n} -A_{ji} x_i \leq -b_j$ 3. $\sum_{1 \leq i \leq n} A_{ji} x_i = b_j$
- - $\begin{array}{ll} \bullet & \sum_{1 \leq i \leq n} A_{ji} x_i \leq b_j \\ \bullet & \sum_{1 \leq i \leq n} A_{ji} x_i \geq b_j \end{array}$
- 4. If x_i has no lower bound, replace x_i with $x_i x_i'$

6.12 Euclidean

- $m = \lfloor \frac{an+b}{c} \rfloor$
- Time complexity: $O(\log n)$

$$\begin{split} f(a,b,c,n) &= \sum_{i=0}^n \lfloor \frac{ai+b}{c} \rfloor \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)}{2} + \lfloor \frac{b}{c} \rfloor \cdot (n+1) \\ + f(a \bmod c, b \bmod c, c, n), & a \geq c \lor b \geq c \\ 0, & n < 0 \lor a = 0 \\ nm - f(c, c - b - 1, a, m - 1), & \text{otherwise} \end{cases} \end{split}$$

$$\begin{split} g(a,b,c,n) &= \sum_{i=0}^n i \lfloor \frac{ai+b}{c} \rfloor \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor \cdot \frac{n(n+1)}{2} \\ +g(a \mod c, b \mod c, c, n), & a \geq c \vee b \geq c \\ 0, & n < 0 \vee a = 0 \end{cases} \\ &= \begin{pmatrix} \frac{1}{2} \cdot (n(n+1)m - f(c, c-b-1, a, m-1)) \\ -h(c, c-b-1, a, m-1)), & \text{otherwise} \end{cases} \end{split}$$

$$\begin{split} h(a,b,c,n) &= \sum_{i=0}^n \lfloor \frac{ai+b}{c} \rfloor^2 \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor^2 \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor^2 \cdot (n+1) \\ + \lfloor \frac{a}{c} \rfloor \cdot \lfloor \frac{b}{c} \rfloor \cdot n(n+1) \\ + h(a \bmod c, b \bmod c, c, n) \\ + 2 \lfloor \frac{a}{c} \rfloor \cdot g(a \bmod c, b \bmod c, c, n) \\ + 2 \lfloor \frac{b}{c} \rfloor \cdot f(a \bmod c, b \bmod c, c, n), & a \geq c \lor b \geq c \\ 0, & n < 0 \lor a = 0 \\ nm(m+1) - 2g(c, c-b-1, a, m-1) \\ -2f(c, c-b-1, a, m-1) - f(a, b, c, n), & \text{otherwise} \end{cases} \end{split}$$

6.13 Theorem

• Kirchhoff's Theorem

Denote L be a $n \times n$ matrix as the Laplacian matrix of graph G, where $L_{ii}=d(i)$, $L_{ij}=-c$ where c is the number of edge (i,j) in

- The number of undirected spanning in G is $|\det(\tilde{L}_{11})|$.
- The number of directed spanning tree rooted at r in G is $|\det(\tilde{L}_{rr})|$.
- Tutte's Matrix

Let D be a n imes n matrix, where $d_{ij} = x_{ij}$ (x_{ij} is chosen uniformly at random) if i < j and $(i, j) \in E$, otherwise $d_{ij} = -d_{ji}$. $\frac{rank(D)}{2}$ is the maximum matching on \widehat{G} .

- Cayley's Formula
 - Given a degree sequence d_1, d_2, \dots, d_n for each $\emph{labeled}$ vertices, there are

$$\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$$

spanning trees.

- Let $T_{n,k}$ be the number of labeled forests on n vertices with k components, such that vertex $1,2,\dots,k$ belong to different components. Then $T_{n,k}=kn^{n-k-1}$.
- Erdős-Gallai Theorem

A sequence of non-negative integers $d_1 \geq d_2 \geq \ldots \geq d_n$ can be represented as the degree sequence of a finite simple graph on n vertices if and only if $d_1+d_2+\ldots+d_n$ is even and

$$\sum_{i=1}^{k} d_i \le k(k-1) + \sum_{i=k+1}^{n} \min(d_i, k)$$

holds for all $1 \le k \le n$.

• Burnside's Lemma

Let X be a set and G be a group that acts on X . For $g\in G$, denote by X^g the elements fixed by g:

$$X^g = \{x \in X \mid gx \in X\}$$

Then

$$|X/G| = \frac{1}{|G|} \sum_{g \in G} |X^g|.$$

• Gale-Ryser theorem

A pair of sequences of nonnegative integers $a_1 \geq \cdots \geq a_n$ and b_1,\dots,b_n is bigraphic if and only if $\sum_{i=1}^n a_i = \sum_{i=1}^n b_i$ and $\sum_{i=1}^k a_i \leq a_i$ $\sum_{i=1}^{n} \min(b_i, k)$ holds for every $1 \leq k \leq n$.

• Fulkerson-Chen-Anstee theorem

A sequence $(a_1,b_1),\ldots,(a_n,b_n)$ of nonnegative integer pairs with $a_1 \geq \cdots \geq a_n$ is digraphic if and only if $\sum_{i=1}^n a_i = \sum_{i=1}^n b_i$ and $\sum_{i=1}^k a_i \leq \sum_{i=1}^k \min(b_i,k-1) + \sum_{i=k+1}^n \min(b_i,k) \text{ holds for every } 1 \leq k \leq n.$

- Möbius inversion formula
 - $f(n) = \sum_{d|n} g(d) \Leftrightarrow g(n) = \sum_{d|n} \mu(d) f(\frac{n}{d})$ $f(n) = \sum_{n|d} g(d) \Leftrightarrow g(n) = \sum_{n|d} \mu(\frac{d}{n}) f(d)$
- Spherical cap

- A portion of a sphere cut off by a plane. r: sphere radius, a: radius of the base of the cap, h: height of the cap, θ : $\arcsin(a/r)$. Volume = $\pi h^2(3r-h)/3 = \pi h(3a^2+h^2)/6 = \pi r^3(2+\cos\theta)(1-\cos\theta)^2/3$. Area = $2\pi rh = \pi(a^2+h^2) = 2\pi r^2(1-\cos\theta)$.
- Chinese Remainder Theorem
 - $x \equiv a_i \pmod{m_i}$
 - $M = \prod m_i, M_i = M/m_i$
 - $t_i M_i \equiv 1 \pmod{m_i}$
 - $x = \sum a_i t_i M_i \pmod{M}$

6.14 Estimation

- The number of divisors of n is at most around $100\ {\rm for}\ n<5e4$, 500 for n < 1e7, 2000 for n < 1e10, 200000 for n < 1e19.
- The number of ways of writing n as a sum of positive integers, disregarding the order of the summands. $1,1,2,3,5,7,11,15,22,30\,$ for $n=0\sim 9$, 627 for n=20, $\sim 2e5$ for n=50, $\sim 2e8$ for n = 100.
- Total number of partitions of n distinct elements: B(n)=1,1,2,5,15,52,203,877,4140,21147,115975,678570,4213597,27644437, 190899322,

6.15 General Purpose Numbers

• Bernoulli numbers

$$\begin{split} B_0 - 1, B_1^{\pm} &= \pm \frac{1}{2}, B_2 = \frac{1}{6}, B_3 = 0 \\ &\sum_{j=0}^m {m+1 \choose j} B_j = 0 \text{, EGF is } B(x) = \frac{x}{e^x - 1} = \sum_{n=0}^\infty B_n \frac{x^n}{n!} \text{.} \\ S_m(n) &= \sum_{k=1}^n k^m = \frac{1}{m+1} \sum_{k=0}^m {m+1 \choose k} B_k^+ n^{m+1-k} \end{split}$$

- Stirling numbers of the second kind Partitions of \boldsymbol{n} distinct elements into exactly \boldsymbol{k} groups.

$$S(n,k) = S(n-1,k-1) + kS(n-1,k), S(n,1) = S(n,n) = 1$$

$$S(n,k) = \frac{1}{k!} \sum_{i=0}^{k} (-1)^{k-i} {k \choose i} i^{n}$$

$$x^{n} = \sum_{i=0}^{n} S(n,i)(x)_{i}$$

• Pentagonal number theorem

$$\prod_{n=1}^{\infty} (1-x^n) = 1 + \sum_{k=1}^{\infty} (-1)^k \left(x^{k(3k+1)/2} + x^{k(3k-1)/2} \right)$$

• Catalan numbers

$$C_n^{(k)} = \frac{1}{(k-1)n+1} {kn \choose n}$$
$$C^{(k)}(x) = 1 + x[C^{(k)}(x)]^k$$

• Eulerian numbers

Number of permutations $\pi \in S_n$ in which exactly k elements are greater than the previous element. k j:s s.t. $\pi(j) > \pi(j+1)$, k+1 j:s s.t. $\pi(j) \geq j$, k j:s s.t. $\pi(j) > j$.

$$E(n,k) = (n-k)E(n-1,k-1) + (k+1)E(n-1,k)$$

$$E(n,0) = E(n, n-1) = 1$$

$$E(n,k) = \sum_{j=0}^{k} (-1)^{j} {n+1 \choose j} (k+1-j)^{n}$$

6.16 Tips for Generating Funtion

```
• Ordinary Generating Function A(x)=\sum_{i\geq 0}a_ix^i
\begin{array}{l} -A(rx)\Rightarrow r^na_n\\ -A(x)+B(x)\Rightarrow \sum_{i=0}^na_ib_{n-i}\\ -A(x)B(x)\Rightarrow \sum_{i=0}^na_ib_{n-i}\\ -A(x)^k\Rightarrow \sum_{i_1+i_2+\dots+i_k=n}a_{i_1}a_{i_2}\dots a_{i_k}\\ -xA(x)'\Rightarrow na_n\\ -\frac{A(x)}{1-x}\Rightarrow \sum_{i=0}^na_i \end{array}
• Exponential Generating Function A(x)=\sum_{i\geq 0}\frac{a_i}{i!}x_i
\begin{array}{l} -A(x)+B(x)\Rightarrow a_n+b_n\\ -A(x)B(x)\Rightarrow \sum_{i=0}^n(i)a_ib_{n-i}\\ -A(x)B(x)\Rightarrow \sum_{i=0}^n(i)a_ib_{n-i}\\ -A(x)B(x)\Rightarrow \sum_{i=0}^n(i)a_ib_{n-i}\\ -A(x)B(x)\Rightarrow na_n + i(i_1,i_2,\dots,i_k)a_{i_1}a_{i_2}\dots a_{i_k} \end{array}
• Special Generating Function
\begin{array}{l} -(1+x)^n=\sum_{i\geq 0}\binom{n}{i}x^i\\ -\frac{1}{(1-x)^n}=\sum_{i\geq 0}\binom{n}{i}x^i\\ -\frac{1}{(1-x)^n}=\sum_{i\geq 0}\binom{n}{i}x^i\\ -\frac{1}{(1-x)^n}=\sum_{i\geq 0}\binom{n}{i}x^i\\ -\frac{1}{(1-x)^n}=\sum_{i\geq 0}\binom{n}{i}x^i\\ -\frac{1}{(1-x)^n}=\sum_{i\geq 0}\binom{n}{i}x^i\\ -\frac{1}{(1-x)^n}=\sum_{i\geq 0}\binom{n}{i}x^i\\ \end{array}
```

7 Polynomial

7.1 Number Theoretic Transform

```
// mul, add, sub, mpow
// ll -> int if too slow
struct NTT {
  11 w[N];
  NTT() {
    ll dw = mpow(G, (mod - 1) / N);
    w[0] = 1;
    for (int i = 1; i < N; ++i) w[i] = w[i - 1] * dw %</pre>
  void operator()(vector<ll>& a, bool inv = false) { //
       \theta \leftarrow a[i] \leftarrow P
    int x = 0, n = a.size();
    for (int j = 1; j < n - 1; ++j) {
       for (int k = n >> 1; (x ^= k) < k; k >>= 1);
       if (j < x) swap(a[x], a[j]);</pre>
    for (int L = 2; L <= n; L <<= 1) {
       int dx = N / L, dl = L >> 1;
for (int i = 0; i < n; i += L) {</pre>
         for (int j = i, x = 0; j < i + dl; ++j, x += dx
              ) {
           ll tmp = mul(a[j + dl], w[x]);
           a[j + dl] = sub(a[j], tmp);
           a[j] = add(a[j], tmp);
         }
      }
    if (inv) {
       reverse(a.begin() + 1, a.end());
       11 invn = mpow(n, mod - 2);
       for (int i = 0; i < n; ++i) a[i] = mul(a[i], invn</pre>
  }
} ntt;
```

7.2 Primes

```
Prime
            Root
                    Prime
                                  Root
                    167772161
7681
            17
12289
                    104857601
            11
40961
                    985661441
65537
                    998244353
786433
            10
                    1107296257
                                  10
5767169
            3
                    2013265921
                                  31
7340033
                    2810183681
                                  11
23068673
                    2885681153
469762049
                    605028353
```

7.3 Polynomial Operations

}

```
vector <11> out(n);
  for (int i = 0; i < n; ++i) out[i] = mul(a[i], b[i]);</pre>
  ntt(out, true), out.resize(min(m, bound));
  return out:
vector <ll> Inverse(vector <ll> a) {
  // O(NlogN), a[0] != 0
  int n = a.size();
  vector <1l> res(1, mpow(a[0], mod - 2));
  for (int m = 1; m < n; m <<= 1) {</pre>
    if (n < m * 2) a.resize(m * 2);</pre>
    vector <11> v1(a.begin(), a.begin() + m * 2), v2 =
         res;
    v1.resize(m * 4), v2.resize(m * 4);
    ntt(v1), ntt(v2);
    for (int i = 0; i < m * 4; ++i) v1[i] = mul(mul(v1[</pre>
    i], v2[i]), v2[i]);
ntt(v1, true);
res.resize(m * 2);
    for (int i = 0; i < m; ++i) res[i] = add(res[i],</pre>
         res[i]);
    for (int i = 0; i < m * 2; ++i) res[i] = sub(res[i])</pre>
         ], v1[i]);
  res.resize(n):
  return res;
pair <vector <ll>, vector <ll>> Divide(vector <ll> a,
    vector <ll> b) {
  // a = bQ + R, O(NlogN), b.back() != 0
  int n = a.size(), m = b.size(), k = n - m + 1;
  if (n < m) return {{0}, a};</pre>
  vector <1l> ra = a, rb = b;
  reverse(all(ra)), ra.resize(k);
  reverse(all(rb)), rb.resize(k);
  vector <11> Q = Mul(ra, Inverse(rb), k);
  reverse(all(Q));
  vector \langle 11 \rangle res = Mul(b, Q), R(m - 1);
  for (int i = 0; i < m - 1; ++i) R[i] = sub(a[i], res[
      i]);
  return {Q, R};
vector <11> SqrtImpl(vector <11> a) {
  if (a.empty()) return {0};
  int z = QuadraticResidue(a[0], mod), n = a.size();
  if (z == -1) return {-1};
  vector <ll> q(1, z);
  const int inv2 = (mod + 1) / 2;
  for (int m = 1; m < n; m <<= 1) {</pre>
    if (n < m * 2) a.resize(m * 2);</pre>
    q.resize(m * 2);
    vector i + f2 = Mul(q, q, m * 2);
for (int i = 0; i < m * 2; ++i) f2[i] = sub(f2[i],</pre>
         a[i]);
    f2 = Mul(f2, Inverse(q), m * 2);
for (int i = 0; i < m * 2; ++i) q[i] = sub(q[i],</pre>
         mul(f2[i], inv2));
  q.resize(n);
  return q;
vector <11> Sqrt(vector <11> a) {
  // O(NlogN), return {-1} if not exists
  int n = a.size(), m = 0;
  while (m < n && a[m] == 0) m++;</pre>
  if (m == n) return vector <11>(n);
  if (m & 1) return {-1};
  vector <11> s = SqrtImpl(vector <11>(a.begin() + m, a
       .end()));
  if (s[0] == -1) return {-1};
  vector <11> res(n);
  for (int i = 0; i < s.size(); ++i) res[i + m / 2] = s</pre>
      [i];
  return res;
vector <ll> Derivative(vector <ll> a) {
  int n = a.size();
  vector <ll> res(n - 1);
  for (int i = 0; i < n - 1; ++i) res[i] = mul(a[i +</pre>
      1], i + 1);
  return res;
```

```
vector <ll> Integral(vector <ll> a) {
 int n = a.size();
  vector <ll> res(n + 1);
 for (int i = 0; i < n; ++i) {</pre>
   res[i + 1] = mul(a[i], mpow(i + 1, mod - 2));
 return res;
vector <11> Ln(vector <11> a) {
  // O(NlogN), a[0] = 1
  int n = a.size();
 if (n == 1) return {0};
 vector <1l> d = Derivative(a);
 a.pop_back();
 return Integral(Mul(d, Inverse(a), n - 1));
vector <11> Exp(vector <11> a) {
  // O(NlogN), a[0] = 0
 int n = a.size();
 vector <ll> q(1, 1);
  a[0] = add(a[0], 1);
  for (int m = 1; m < n; m <<= 1) {</pre>
    if (n < m * 2) a.resize(m * 2);</pre>
    vector <ll> g(a.begin(), a.begin() + m * 2), h(all(
        a));
    h.resize(m * 2), h = Ln(h);
    for (int i = 0; i < m * 2; ++i) {</pre>
      g[i] = sub(g[i], h[i]);
   q = Mul(g, q, m * 2);
  q.resize(n);
 return q;
vector <ll> Pow(vector <ll> a, ll k) {
 int n = a.size(), m = 0;
  vector <11> ans(n, 0);
 while (m < n && a[m] == 0) m++;</pre>
 if (k \&\& m \&\& (k >= n || k * m >= n)) return ans;
 if (m == n) return ans[0] = 1, ans;
 11 lead = m * k;
 vector <1l> b(a.begin() + m, a.end());
 11 base = mpow(b[0], k), inv = mpow(b[0], mod - 2);
 for (int i = 0; i < n - m; ++i) b[i] = mul(b[i], inv)</pre>
 b = Ln(b);
 for (int i = 0; i < n - m; ++i) b[i] = mul(b[i], k %</pre>
      mod);
 b = Exp(b);
  for (int i = lead; i < n; ++i) ans[i] = mul(b[i -</pre>
      lead], base);
 return ans;
vector <ll> Evaluate(vector <ll> a, vector <ll> x) {
 if (x.empty()) return {};
 int n = x.size();
 vector <vector <11>> up(n * 2);
 for (int i = 0; i < n; ++i) up[i + n] = {sub(0, x[i])}
       1};
 for (int i = n - 1; i > 0; --i) up[i] = Mul(up[i *
      2], up[i * 2 + 1]);
  vector <vector <11>> down(n * 2);
  down[1] = Divide(a, up[1]).second;
  for (int i = 2; i < n * 2; ++i) down[i] = Divide(down</pre>
      [i >> 1], up[i]).second;
  vector <ll> y(n);
 for (int i = 0; i < n; ++i) y[i] = down[i + n][0];</pre>
 return v;
vector <ll> Interpolate(vector <ll> x, vector <ll> y) {
 int n = x.size();
  vector <vector <11>> up(n * 2);
  for (int i = 0; i < n; ++i) up[i + n] = {sub(0, x[i])}
       1};
  for (int i = n - 1; i > 0; --i) up[i] = Mul(up[i *
      2], up[i * 2 + 1]);
  vector <ll> a = Evaluate(Derivative(up[1]), x);
  for (int i = 0; i < n; ++i) {</pre>
   a[i] = mul(y[i], mpow(a[i], mod - 2));
  vector <vector <11>> down(n * 2);
 for (int i = 0; i < n; ++i) down[i + n] = {a[i]};</pre>
```

```
for (int i = n - 1; i > 0; --i) {
   vector <ll> lhs = Mul(down[i * 2], up[i * 2 + 1]);
   vector <ll> rhs = Mul(down[i * 2 + 1], up[i * 2]);
   down[i].resize(lhs.size());
   for (int j = 0; j < lhs.size(); ++j) {
      down[i][j] = add(lhs[j], rhs[j]);
   }
}
return down[1];
}</pre>
```

7.4 Fast Linear Recursion

```
ll FastLinearRecursion(vector <ll> a, vector <ll> c, ll
      k) {
  // a_n = sigma c_j * a_{n - j - 1}, 0-based
  // O(NlogNlogK), |a| = |c|
  int n = a.size();
  if (k < n) return a[k];</pre>
  vector <ll> base(n + 1, 1);
  for (int i = 0; i < n; ++i) base[i] = sub(0, c[n - i</pre>
       - 1]);
  vector <1l> poly(n);
  (n == 1 ? poly[0] = c[n - 1] : poly[1] = 1);
  auto calc = [&](vector <ll> p1, vector <ll> p2) {
    // O(n^2) bruteforce or O(nlogn) NTT
    return Divide(Mul(p1, p2), base).second;
  vector \langle 11 \rangle res(n, 0); res[0] = 1;
  for (; k; k >>= 1, poly = calc(poly, poly)) {
    if (k & 1) res = calc(res, poly);
  11 ans = 0;
  for (int i = 0; i < n; ++i) {</pre>
    (ans += res[i] * a[i]) %= mod;
  return ans:
}
```

7.5 Fast Walsh Transform

```
void fwt(vector <int> &a) {
 // and : a[j] += x;
         : a[j] -= x;
  //
         : a[j ^ (1 << i)] += y;
  // or
          : a[j ^ (1 << i)] -= y;
 //
  // xor : a[j] = x - y, a[j ^ (1 << i)] = x + y;
          : a[j] = (x - y) / 2, a[j ^ (1 << i)] = (x + y)
      ) / 2;
  int n = __lg(a.size());
  for (int i = 0; i < n; ++i) {</pre>
    for (int j = 0; j < 1 << n; ++j) if (j >> i & 1) {
  int x = a[j ^ (1 << i)], y = a[j];</pre>
      // do something
  }
vector<int> subs_conv(vector<int> a, vector<int> b) {
  // c_i = sum_{j} & k = 0, j | k = i a_j * b_k
  int n = __lg(a.size());
  vector<vector<int>> ha(n + 1, vector<int>(1 << n));</pre>
  vector<vector<int>> hb(n + 1, vector<int>(1 << n));</pre>
  vector<vector<int>> c(n + 1, vector<int>(1 << n));</pre>
  for (int i = 0; i < 1 << n; ++i) {</pre>
    ha[__builtin_popcount(i)][i] = a[i];
    hb[__builtin_popcount(i)][i] = b[i];
  for (int i = 0; i <= n; ++i) fwt(ha[i]),fwt(hb[i]);</pre>
  for (int i = 0; i <= n; ++i)</pre>
    for (int j = 0; i + j <= n; ++j)
      for (int k = 0; k < 1 << n; ++k)
        // mind overflow
         c[i + j][k] += ha[i][k] * hb[j][k];
  for (int i = 0; i <= n; ++i)</pre>
    fwt(c[i], true);
  vector <int> ans(1 << n);</pre>
  for (int i = 0; i < 1 << n; ++i)</pre>
    ans[i] = c[__builtin_popcount(i)][i];
  return ans;
```

8 Geometry

8.1 Basic

```
const double eps = 1e-8, pi = acos(-1);
int sign(double x) \{return abs(x) \leftarrow eps ? 0 : (x > 0 ?
    1 : -1);}
struct Pt {
 double x, y;
 Pt (double _x, double _y) : x(_x), y(_y) {}
 Pt operator + (Pt o) {return Pt(x + o.x, y + o.y);}
 Pt operator - (Pt o) {return Pt(x - o.x, y - o.y);}
 Pt operator * (double k) {return Pt(x * k, y * k);}
 Pt operator / (double k) {return Pt (x / k, y / k);}
 double operator * (Pt o) {return x * o.x + y * o.y;}
 double operator ^ (Pt o) {return x * o.y - y * o.x;}
struct Line {
 Pt a, b;
struct Cir {
 Pt o; double r;
double abs2(Pt o) {return o.x * o.x + o.y * o.y;}
double abs(Pt o) {return sqrt(abs2(o));}
int ori(Pt o, Pt a, Pt b) {return sign((o - a) ^ (o - b
    ));}
bool btw(Pt a, Pt b, Pt c) { // c on segment ab?
 return ori(a, b, c) == 0 && sign((c - a) * (c - b))
      <= 0:
double area(Pt a, Pt b, Pt c) {return abs((a - b) ^ (a
    - c)) / 2;}
Pt unit(Pt o) {return o / abs(o);}
Pt rot(Pt a, double o) { // CCW
 double c = cos(o), s = sin(o);
  return Pt(c * a.x - s * a.y, s * a.x + c * a.y);
Pt proj_vector(Pt a, Pt b, Pt c) { // vector ac proj to
 return (b - a) * ((c - a) * (b - a)) / ((b - a) * (b
      - a));
Pt proj_pt(Pt a, Pt b, Pt c) { // point c proj to ab
 return proj_vector(a, b, c) + a;
```

8.2 Heart

```
Pt circenter(Pt p0, Pt p1, Pt p2) { // radius = abs(
   center)
  p1 = p1 - p0, p2 = p2 - p0;
 double x1 = p1.x, y1 = p1.y, x2 = p2.x, y2 = p2.y;
 double m = 2. * (x1 * y2 - y1 * x2);
 Pt center(0, 0);
 center.x = (x1 * x1 * y2 - x2 * x2 * y1 + y1 * y2 * (
     y1 - y2)) / m;
 center.y = (x1 * x2 * (x2 - x1) - y1 * y1 * x2 + x1 *
      y2 * y2) / m;
 return center + p0;
Pt incenter(Pt p1, Pt p2, Pt p3) { // radius = area / s
     * 2
  double a = abs(p2 - p3), b = abs(p1 - p3), c = abs(p1
       - p2);
  double s = a + b + c;
 return (p1 * a + p2 * b + p3 * c) / s;
Pt masscenter(Pt p1, Pt p2, Pt p3)
{ return (p1 + p2 + p3) / 3; }
Pt orthocenter(Pt p1, Pt p2, Pt p3)
{ return masscenter(p1, p2, p3) * 3 - circenter(p1, p2,
     p3) * 2; }
```

8.3 External Bisector

```
Pt external_bisector(Pt p1, Pt p2, Pt p3) { //213
Pt L1 = p2 - p1, L2 = p3 - p1;
L2 = L2 * abs(L1) / abs(L2);
return L1 + L2;
```

8.4 Intersection of Segments

```
Pt LinesInter(Line a, Line b) {
    double abc = (a.b - a.a) ^ (b.a - a.a);
    double abd = (a.b - a.a) ^ (b.b - a.a);
    if (sign(abc - abd) == 0) return b.b;// no inter
    return (b.b * abc - b.a * abd) / (abc - abd);
}

vector<Pt> SegsInter(Line a, Line b) {
    if (btw(a.a, a.b, b.a)) return {b.a};
    if (btw(a.a, a.b, b.b)) return {b.b};
    if (btw(b.a, b.b, a.a)) return {a.a};
    if (btw(b.a, b.b, a.b)) return {a.b};
    if (ori(a.a, a.b, b.a) * ori(a.a, a.b, b.b) == -1 &&
        ori(b.a, b.b, a.a) * ori(b.a, b.b, a.b) == -1)
    return {LinesInter(a, b)};
    return {};
}
```

8.5 Intersection of Circle and Line

8.6 Intersection of Circles

8.7 Intersection of Polygon and Circle

```
double _area(Pt pa, Pt pb, double r){
  if(abs(pa) < abs(pb)) swap(pa, pb);</pre>
  if(abs(pb) < eps) return 0;</pre>
  double S, h, theta;
  double a = abs(pb), b = abs(pa), c = abs(pb - pa);
  double cosB = pb * (pb - pa) / a / c, B = acos(cosB);
  double cosC = (pa * pb) / a / b, C = acos(cosC);
  if (a > r) {
    S = (C / 2) * r * r;
    h = a * b * sin(C) / c;
    if (h < r && B < pi / 2) S -= (acos(h / r) * r * r
        - h * sqrt(r * r - h * h));
  } else if (b > r) {
  theta = pi - B - asin(sin(B) / r * a);
    S = .5 * a * r * sin(theta) + (C - theta) / 2 * r *
         r;
  } else
    S = .5 * sin(C) * a * b;
  return S:
double area_poly_circle(vector<Pt> poly, Pt 0, double r
    ) {
  double S = 0; int n = poly.size();
  for(int i = 0; i < n; ++i)</pre>
    S += _area(poly[i] - 0, poly[(i + 1) % n] - 0, r) *
         ori(0, poly[i], poly[(i + 1) % n]);
  return fabs(S);
```

8.8 Tangent Lines of Circle and Point

8.9 Tangent Lines of Circles

```
vector<Line> tangent(Cir a, Cir b) {
#define Pij \
 Pt i = unit(b.o - a.o) * a.r, j = Pt(i.y, -i.x);\
  z.push_back({a.o + i, a.o + i + j});
#define deo(I,J) \
  double d = abs(a.o - b.o), e = a.r I b.r, o = acos(e)
      / d);\
 Pt i = unit(b.o - a.o), j = rot(i, o), k = rot(i, -o)
 z.push\_back({a.o + j * a.r, b.o J j * b.r});\
 z.push_back({a.o + k * a.r, b.o J k * b.r});
 if (a.r < b.r) swap(a, b);
 vector<Line> z;
  if (abs(a.o - b.o) + b.r < a.r) return z;</pre>
 else if (sign(abs(a.o - b.o) + b.r - a.r) == 0) { Pij
 else {
   deo(-,+); // inter
    // outer
    if (sign(d - a.r - b.r) == 0) { Pij; }
    else if (d > a.r + b.r) { deo(+,-); }
 return z;
```

8.10 Point In Convex

8.11 Point Segment Distance

8.12 Convex Hull

```
ans.pop_back();
ans.push_back(pt[i]);
}
reverse(all(pt));
}
ans.pop_back();
return ans;
}
```

8.13 Convex Hull Distance

8.14 Minimum Enclosing Circle

```
Cir min_enclosing(vector<Pt> &p)
  random_shuffle(p.begin(), p.end());
  double r = 0.0;
  Pt cent = p[0];
  for (int i = 1; i < p.size(); ++i) {</pre>
    if (abs2(cent - p[i]) <= r) continue;</pre>
    cent = p[i];
    r = 0.0;
    for (int j = 0; j < i; ++j) {</pre>
      if (abs2(cent - p[j]) <= r) continue;</pre>
      cent = (p[i] + p[j]) / 2;
      r = abs2(p[j] - cent);
      for (int k = 0; k < j; ++k) {
        if (abs2(cent - p[k]) <= r) continue;</pre>
        cent = circenter(p[i], p[j], p[k]);
        r = abs2(p[k] - cent);
      }
    }
  return {cent, sqrt(r)};
```

8.15 Union of Circles

```
vector<pair<double, double>> CoverSegment(Cir a, Cir b)
  double d = abs(a.o - b.o);
  vector<pair<double, double>> res;
  if (sign(a.r + b.r - d) == 0);
  else if (d <= abs(a.r - b.r) + eps) {
    if (a.r < b.r) res.emplace_back(0, 2 * pi);</pre>
  } else if (d < abs(a.r + b.r) - eps) {</pre>
    double o = acos((sqrt(a.r) + sqrt(d) - sqrt(b.r)) /
         (2 * a.r * d)), z = atan2((b.o - a.o).y, (b.o
        - a.o).x);
    if (z < 0) z += 2 * pi;
    double 1 = z - o, r = z + o;
    if (1 < 0) 1 += 2 * pi;
    if (r > 2 * pi) r -= 2 * pi;
    if (1 > r) res.emplace_back(1, 2 * pi), res.
        emplace_back(0, r);
    else res.emplace_back(l, r);
  return res;
double CircleUnionArea(vector<Cir> c) { // circle
    should be identical
  int n = c.size();
  double a = 0, w;
  for (int i = 0; w = 0, i < n; ++i) {
    vector<pair<double, double>> s = {{2 * pi, 9}}, z;
    for (int j = 0; j < n; ++j) if (i != j) {</pre>
      z = CoverSegment(c[i], c[j]);
      for (auto &e : z) s.push_back(e);
    sort(s.begin(), s.end());
```

8.16 Polar Angle Sort

8.17 Rotating Caliper

8.18 Rotating SweepLine

```
void RotatingSweepLine(vector <Pt> &pt) {
  int n = pt.size();
   vector <int> id(n), pos(n);
  vector <pair <int, int>> line;
  for (int i = 0; i < n; ++i) for (int j = 0; j < n; ++
       j) if (i ^ j) line.emplace_back(i, j);
   sort(line.begin(), line.end(), [&](pair <int, int> i,
        pair <int, int> j) {
     Pt a = pt[i.second] - pt[i.first], b = pt[j.second]
          - pt[j.first];
     return (a.pos() == b.pos() ? sign(a ^ b) > 0 : a.
         pos() < b.pos());</pre>
  iota(id.begin(), id.end(), 0);
   sort(id.begin(), id.end(), [&](int i, int j) {
     return (sign(pt[i].y - pt[j].y) == 0 ? pt[i].x < pt</pre>
         [j].x : pt[i].y < pt[j].y);
  for (int i = 0; i < n; ++i)</pre>
    pos[id[i]] = i;
  for (auto [i, j] : line) {
    // point sort by the distance to line(i, j) // do something.
     tie(pos[i], pos[j], id[pos[i]], id[pos[j]]) =
         make_tuple(pos[j], pos[i], j, i);
  }
| }
```

8.19 Half Plane Intersection

```
return sign(A ^ B) > 0;
 return pos(A) < pos(B);</pre>
});
deque <Pt> inter;
deque <pair <Pt, Pt>> seg;
int n = vec.size();
auto get = [&](pair <Pt, Pt> a, pair <Pt, Pt> b) {
    return intersect(a.first, a.second, b.first, b.
    second);};
for (int i = 0; i < n; ++i) if (!i || vec[i] != vec[i</pre>
     - 1]) {
  while (seg.size() >= 2 && sign((vec[i].second -
      inter.back()) ^ (vec[i].first - inter.back()))
      == 1) seg.pop_back(), inter.pop_back();
  while (seg.size() >= 2 && sign((vec[i].second -
      inter.front()) ^ (vec[i].first - inter.front())
      ) == 1) seg.pop_front(), inter.pop_front();
  seg.push_back(vec[i]);
  if (seg.size() >= 2) inter.pb(get(seg[seg.size() -
      2], seg.back()));
while (seg.size() >= 2 && sign((seg.front().second -
    inter.back()) ^ (seg.front().first - inter.back()
    )) == 1) seg.pop_back(), inter.pop_back();
inter.push_back(get(seg.front(), seg.back()));
return vector <Pt>(all(inter));
```

8.20 Minkowski Sum

```
void reorder(vector <Pt> &P) {
  rotate(P.begin(), min_element(all(P), [&](Pt a, Pt b)
       { return make_pair(a.y, a.x) < make_pair(b.y, b.
      x); }), P.end());
vector <Pt> Minkowski(vector <Pt> P, vector <Pt> Q) {
  // P, Q: convex polygon
  reorder(P), reorder(Q);
  int n = P.size(), m = Q.size();
  P.pb(P[0]), P.pb(P[1]), Q.pb(Q[0]), Q.pb(Q[1]);
  vector <Pt> ans;
  for (int i = 0, j = 0; i < n || j < m; ) {</pre>
    ans.pb(P[i] + Q[j]);
    auto val = (P[i + 1] - P[i]) ^ (Q[j + 1] - Q[j]);
    if (val >= 0) i++;
    if (val <= 0) j++;</pre>
  return ans;
```

8.21 Delaunay Triangulation

```
/* Delaunay Triangulation:
Given a sets of points on 2D plane, find a
triangulation such that no points will strictly
inside circumcircle of any triangle.
find : return a triangle contain given point
add_point : add a point into triangulation
A Triangle is in triangulation iff. its has_chd is 0.
Region of triangle u: iterate each u.edge[i].tri,
each points are u.p[(i+1)\%3], u.p[(i+2)\%3]
Voronoi diagram: for each triangle in triangulation,
the bisector of all its edges will split the region.
nearest point will belong to the triangle containing it
const ll inf = MAXC * MAXC * 100; // Lower_bound
    unknown
struct Tri;
struct Edge {
 Tri* tri; int side;
  Edge(): tri(0), side(0){}
  Edge(Tri* _tri, int _side): tri(_tri), side(_side){}
struct Tri {
  pll p[3];
  Edge edge[3];
  Tri* chd[3];
  Tri() {}
  Tri(const pll& p0, const pll& p1, const pll& p2) {
    p[0] = p0; p[1] = p1; p[2] = p2;
    chd[0] = chd[1] = chd[2] = 0;
```

```
bool has_chd() const { return chd[0] != 0; }
  int num_chd() const {
    return !!chd[0] + !!chd[1] + !!chd[2];
  bool contains(pll const& q) const {
    for (int i = 0; i < 3; ++i)
      if (ori(p[i], p[(i + 1) % 3], q) < 0)</pre>
        return 0:
    return 1;
} pool[N * 10], *tris;
void edge(Edge a, Edge b) {
  if(a.tri) a.tri->edge[a.side] = b;
  if(b.tri) b.tri->edge[b.side] = a;
struct Trig { // Triangulation
  Trig() {
    the_root = // Tri should at least contain all
        points
      new(tris++) Tri(pll(-inf, -inf), pll(inf + inf, -
          inf), pll(-inf, inf + inf));
  Tri* find(pll p) { return find(the_root, p); }
  void add_point(const pll &p) { add_point(find(
      the_root, p), p); }
  Tri* the_root;
  static Tri* find(Tri* root, const pll &p) {
    while (1) {
      if (!root->has_chd())
        return root;
      for (int i = 0; i < 3 && root->chd[i]; ++i)
        if (root->chd[i]->contains(p)) {
          root = root->chd[i];
          break;
        }
    assert(0); // "point not found"
  void add_point(Tri* root, pll const& p) {
   Tri* t[3];
    /* split it into three triangles */
    for (int i = 0; i < 3; ++i)
      t[i] = new(tris++) Tri(root->p[i], root->p[(i +
          1) % 3], p);
    for (int i = 0; i < 3; ++i)
      edge(Edge(t[i], 0), Edge(t[(i + 1) % 3], 1));
    for (int i = 0; i < 3; ++i)
      edge(Edge(t[i], 2), root->edge[(i + 2) % 3]);
    for (int i = 0; i < 3; ++i)</pre>
      root->chd[i] = t[i];
    for (int i = 0; i < 3; ++i)
      flip(t[i], 2);
  void flip(Tri* tri, int pi) {
    Tri* trj = tri->edge[pi].tri;
    int pj = tri->edge[pi].side;
    if (!trj) return;
    if (!in_cc(tri->p[0], tri->p[1], tri->p[2], trj->p[
        pj])) return;
    /* flip edge between tri,trj */
    Tri* trk = new(tris++) Tri(tri->p[(pi + 1) % 3],
        trj->p[pj], tri->p[pi]);
    Tri* trl = new(tris++) Tri(trj->p[(pj + 1) % 3],
        tri->p[pi], trj->p[pj]);
    edge(Edge(trk, 0), Edge(trl, 0));
    edge(Edge(trk, 1), tri->edge[(pi + 2) % 3]);
    edge(Edge(trk, 2), trj->edge[(pj + 1) % 3]);
edge(Edge(trl, 1), trj->edge[(pj + 2) % 3]);
    edge(Edge(trl, 2), tri->edge[(pi + 1) % 3]);
    tri->chd[0] = trk; tri->chd[1] = trl; tri->chd[2] =
    trj->chd[0] = trk; trj->chd[1] = trl; trj->chd[2] =
    flip(trk, 1); flip(trk, 2);
    flip(trl, 1); flip(trl, 2);
 }
};
vector<Tri*> triang; // vector of all triangle
set<Tri*> vst;
void go(Tri* now) { // store all tri into triang
 if (vst.find(now) != vst.end())
```

```
return;
vst.insert(now);
if (!now->has_chd())
    return triang.pb(now);
for (int i = 0; i < now->num_chd(); ++i)
    go(now->chd[i]);
}
void build(int n, pll* ps) { // build triangulation
    tris = pool; triang.clear(); vst.clear();
    random_shuffle(ps, ps + n);
Trig tri; // the triangulation structure
    for (int i = 0; i < n; ++i)
        tri.add_point(ps[i]);
    go(tri.the_root);
}</pre>
```

8.22 Triangulation Vonoroi

```
vector<Line> ls[N];
pll arr[N];
Line make_line(pdd p, Line 1) {
  pdd d = 1.Y - 1.X; d = perp(d);
  pdd m = (1.X + 1.Y) / 2;
  l = Line(m, m + d);
  if (ori(1.X, 1.Y, p) < 0)</pre>
    l = Line(m + d, m);
  return 1;
double calc_area(int id) {
  // use to calculate the area of point "strictly in
      the convex hull"
  vector<Line> hpi = halfPlaneInter(ls[id]);
  vector<pdd> ps;
  for (int i = 0; i < SZ(hpi); ++i)</pre>
    ps.pb(intersect(hpi[i].X, \, hpi[i].Y, \, hpi[(i\, +\, 1)\, \, \%
        SZ(hpi)].X, hpi[(i + 1) % SZ(hpi)].Y));
  double rt = 0:
  for (int i = 0; i < SZ(ps); ++i)</pre>
    rt += cross(ps[i], ps[(i + 1) % SZ(ps)]);
  return fabs(rt) / 2;
void solve(int n, pii *oarr) {
  map<pll, int> mp;
  for (int i = 0; i < n; ++i)</pre>
    arr[i] = pll(oarr[i].X, oarr[i].Y), mp[arr[i]] = i;
  build(n, arr); // Triangulation
  for (auto *t : triang) {
    vector<int> p;
    for (int i = 0; i < 3; ++i)</pre>
      if (mp.find(t->p[i]) != mp.end())
        p.pb(mp[t->p[i]]);
    for (int i = 0; i < SZ(p); ++i)</pre>
      for (int j = i + 1; j < SZ(p); ++j) {
        Line l(oarr[p[i]], oarr[p[j]]);
        ls[p[i]].pb(make_line(oarr[p[i]], 1));
        ls[p[j]].pb(make_line(oarr[p[j]], 1));
```

9 Else

9.1 Bit Hack

9.2 Dynamic Programming Condition

9.2.1 Totally Monotone (Concave/Convex)

 $\begin{array}{l} \forall i < i', j < j' \text{, } B[i][j] \leq B[i'][j] \implies B[i][j'] \leq B[i'][j'] \\ \forall i < i', j < j' \text{, } B[i][j] \geq B[i'][j] \implies B[i][j'] \geq B[i'][j'] \end{array}$

9.2.2 Monge Condition (Concave/Convex)

```
\begin{array}{l} \forall i < i', j < j' \text{, } B[i][j] + B[i'][j'] \geq B[i][j'] + B[i'][j] \\ \forall i < i', j < j' \text{, } B[i][j] + B[i'][j'] \leq B[i][j'] + B[i'][j] \end{array}
```

9.2.3 Optimal Split Point

```
If B[i][j] + B[i+1][j+1] \geq B[i][j+1] + B[i+1][j] then H_{i,j-1} \leq H_{i,j} \leq H_{i+1,j}
```

9.3 Slope Trick

```
template<typename T>
struct slope_trick_convex {
 T minn = 0, ground_1 = 0, ground_r = 0;
 priority_queue<T, vector<T>, less<T>> left;
priority_queue<T, vector<T>, greater<T>> right;
 slope_trick_convex() {left.push(numeric_limits<T>::
      min() / 2), right.push(numeric_limits<T>::max() /
       2);}
 void push_left(T x) {left.push(x - ground_l);}
 void push_right(T x) {right.push(x - ground_r);}
 //add a line with slope 1 to the right starting from
 void add_right(T x) {
   T 1 = left.top() + ground_1;
   if (1 <= x) push_right(x);</pre>
    else push_left(x), push_right(l), left.pop(), minn
        += 1 - x;
 //add a line with slope -1 to the left starting from
  void add_left(T x) {
    T r = right.top() + ground_r;
   if (r >= x) push_left(x);
    else push_right(x), push_left(r), right.pop(), minn
         += x - r;
  //val[i]=min(val[j]) for all i-l<=j<=i+r
 void expand(T 1, T r) {ground_1 -= 1, ground_r += r;}
  void shift_up(T x) {minn += x;}
 T get_val(T x) {
    T l = left.top() + ground_l, r = right.top() +
        ground_r;
    if (x >= 1 && x <= r) return minn;
    if (x < 1) {
      vector<T> trash;
      T cur_val = minn, slope = 1, res;
      while (1) {
        trash.push_back(left.top());
        left.pop();
        if (left.top() + ground_l <= x) {</pre>
          res = cur_val + slope * (1 - x);
          break;
        cur_val += slope * (1 - (left.top() + ground_1)
        1 = left.top() + ground_l;
        slope += 1;
      for (auto i : trash) left.push(i);
      return res;
    if (x > r) {
      vector<T> trash;
      T cur_val = minn, slope = 1, res;
      while (1) {
        trash.push_back(right.top());
        right.pop();
        if (right.top() + ground_r >= x) {
          res = cur_val + slope * (x - r);
          break;
        }
        cur_val += slope * ((right.top() + ground_r) -
        r = right.top() + ground_r;
        slope += 1;
      for (auto i : trash) right.push(i);
      return res;
```

```
assert(0);
}
};
```

9.4 Manhattan MST

```
void solve(int n) {
  init();
  vector<int> v(n), ds;
  for (int i = 0; i < n; ++i) {</pre>
    v[i] = i;
    ds.push_back(x[i] - y[i]);
  sort(ds.begin(), ds.end());
  ds.resize(unique(ds.begin(), ds.end()) - ds.begin());
  sort(v.begin(), v.end(), [&](int i, int j) { return x
      [i] == x[j] ? y[i] > y[j] : x[i] > x[j]; });
  int i = 0;
  for (int i = 0; i < n; ++i) {
    int p = lower_bound(ds.begin(), ds.end(), x[v[i]] -
         y[v[i]]) - ds.begin() + 1;
    pair<int, int> q = query(p);
    // query return prefix minimum
    if (~q.second) add_edge(v[i], q.second);
    add(p, make_pair(x[v[i]] + y[v[i]], v[i]));
  }
void make_graph() {
  solve(n);
  for (int i = 0; i < n; ++i) swap(x[i], y[i]);</pre>
  solve(n);
  for (int i = 0; i < n; ++i) x[i] = -x[i];</pre>
  solve(n);
  for (int i = 0; i < n; ++i) swap(x[i], y[i]);</pre>
  solve(n);
```

9.5 Dynamic MST

```
int cnt[maxn], cost[maxn], st[maxn], ed[maxn];
pair<int, int> qr[maxn];
// qr[i].first = id of edge to be changed, qr[i].second
     = weight after operation
// cnt[i] = number of operation on edge i
// call solve(0, q - 1, v, 0), where v contains edges i
     such that cnt[i] == 0
void contract(int 1, int r, vector<int> v, vector<int>
    &x, vector<int> &y) {
  sort(v.begin(), v.end(), [&](int i, int j) {
      if (cost[i] == cost[j]) return i < j;</pre>
      return cost[i] < cost[j];</pre>
      });
  djs.save();
  for (int i = 1; i <= r; ++i) djs.merge(st[qr[i].first</pre>
      ], ed[qr[i].first]);
  for (int i = 0; i < (int)v.size(); ++i) {</pre>
    if (djs.find(st[v[i]]) != djs.find(ed[v[i]])) {
      x.push_back(v[i]);
      djs.merge(st[v[i]], ed[v[i]]);
 djs.undo();
  djs.save();
  for (int i = 0; i < (int)x.size(); ++i) djs.merge(st[</pre>
      x[i]], ed[x[i]]);
  for (int i = 0; i < (int)v.size(); ++i) {</pre>
    if (djs.find(st[v[i]]) != djs.find(ed[v[i]])) {
      y.push_back(v[i]);
      djs.merge(st[v[i]], ed[v[i]]);
    }
 djs.undo();
void solve(int 1, int r, vector<int> v, long long c) {
  if (1 == r) {
    cost[qr[1].first] = qr[1].second;
    if (st[qr[1].first] == ed[qr[1].first]) {
      printf("%lld\n", c);
      return;
```

```
int minv = qr[1].second;
  for (int i = 0; i < (int)v.size(); ++i) minv = min(</pre>
      minv, cost[v[i]]);
  printf("%lld\n", c + minv);
  return;
int m = (1 + r) >> 1;
vector<int> lv = v, rv = v;
vector<int> x, y;
for (int i = m + 1; i <= r; ++i) {
  cnt[qr[i].first]--;
  if (cnt[qr[i].first] == 0) lv.push_back(qr[i].first
}
contract(l, m, lv, x, y);
long long lc = c, rc = c;
djs.save();
for (int i = 0; i < (int)x.size(); ++i) {</pre>
  lc += cost[x[i]];
  djs.merge(st[x[i]], ed[x[i]]);
solve(1, m, y, lc);
djs.undo();
x.clear(), y.clear();
for (int i = m + 1; i <= r; ++i) cnt[qr[i].first]++;</pre>
for (int i = 1; i <= m; ++i) {</pre>
  cnt[qr[i].first]--;
  if (cnt[qr[i].first] == 0) rv.push_back(qr[i].first
contract(m + 1, r, rv, x, y);
djs.save();
for (int i = 0; i < (int)x.size(); ++i) {</pre>
 rc += cost[x[i]];
  djs.merge(st[x[i]], ed[x[i]]);
solve(m + 1, r, y, rc);
djs.undo();
for (int i = 1; i <= m; ++i) cnt[qr[i].first]++;</pre>
```

9.6 ALL LCS

```
void all_lcs(string s, string t) { // 0-base
  vector<int> h(t.size());
  iota(all(h), 0);
  for (int a = 0; a < s.size(); ++a) {
    int v = -1;
    for (int c = 0; c < t.size(); ++c)
        if (s[a] == t[c] || h[c] < v)
            swap(h[c], v);
        // LCS(s[0, a], t[b, c]) =
        // c - b + 1 - sum([h[i] >= b] | i <= c)
        // h[i] might become -1 !!
  }
}</pre>
```

9.7 Hilbert Curve

```
long long hilbertOrder(int x, int y, int pow, int
    rotate) {
  if (pow == 0) return 0;
  int hpow = 1 << (pow-1);</pre>
  int seg = (x < hpow) ? ((y < hpow) ? 0 : 3) : ((y < hpow)
      hpow) ? 1 : 2);
  seg = (seg + rotate) & 3;
  const int rotateDelta[4] = {3, 0, 0, 1};
  int nx = x & (x ^ hpow), ny = y & (y ^ hpow);
  int nrot = (rotate + rotateDelta[seg]) & 3;
  long long subSquareSize = 111 << (pow * 2 - 2);</pre>
  long long ans = seg * subSquareSize;
  long long add = hilbertOrder(nx, ny, pow - 1, nrot);
  ans += (seg == 1 || seg == 2) ? add : (subSquareSize
      - add - 1);
  return ans;
}
```

9.8 Pbds

```
#include <ext/pb_ds/priority_queue.hpp>
#include <ext/pb_ds/assoc_container.hpp>
using namespace __gnu_pbds;
```

```
#include <ext/rope>
using namespace __gnu_cxx;
int main () {
     _gnu_pbds::priority_queue <<mark>int</mark>> pq1, pq2;
  pq1.join(pq2); // pq1 += pq2, pq2 = {}
  cc_hash_table<int, int> m1;
  tree<int, null_type, less<int>, rb_tree_tag,
      tree_order_statistics_node_update> oset;
  oset.insert(2), oset.insert(4);
  cout << *oset.find_by_order(1) << ' ' << oset.</pre>
      order_of_key(1) << '\n'; // 4 0
  bitset <100> BS;
  BS.flip(3), BS.flip(5);
  cout << BS._Find_first() << ' ' << BS._Find_next(3)</pre>
      << '\n'; // 3 5
  rope <int> rp1, rp2;
  rp1.push_back(1), rp1.push_back(3);
  rp1.insert(0, 2); // pos, num
  rp1.erase(0, 2); // pos, len
  rp1.substr(0, 2); // pos, len
  rp2.push_back(4);
  rp1 += rp2, rp2 = rp1;
  cout << rp2[0] << ' ' << rp2[1] << '\n'; // 3 4
```

9.9 Random

```
struct custom_hash {
    static uint64_t splitmix64(uint64_t x) {
        x += 0x9e3779b97f4a7c15;
        x = (x ^ (x >> 30)) * 0xbf58476d1ce4e5b9;
        x = (x ^ (x >> 27)) * 0x94d049bb133111eb;
        return x ^ (x >> 31);
    }
    size_t operator()(uint64_t a) const {
        static const uint64_t FIXED_RANDOM = chrono::
            steady_clock::now().time_since_epoch().count();
        return splitmix64(i + FIXED_RANDOM);
    }
};
unordered_map <int, int, custom_hash> m1;
random_device rd; mt19937 rng(rd());
```

9.10 Smawk Algorithm

```
11 query(int 1, int r) {
 // ...
struct SMAWK {
  // Condition:
  // If M[1][0] < M[1][1] then M[0][0] < M[0][1]
  // If M[1][0] == M[1][1] then M[0][0] <= M[0][1]
  // For all i, find r_i s.t. M[i][r_i] is maximum ||
      minimum.
  int ans[N], tmp[N];
  void interpolate(vector <int> 1, vector <int> r) {
    int n = 1.size(), m = r.size();
    vector <int> nl;
    for (int i = 1; i < n; i += 2) {</pre>
      nl.push_back(l[i]);
    run(nl, r);
    for (int i = 1, j = 0; i < n; i += 2) {
      while (j < m && r[j] < ans[l[i]])</pre>
        j++;
      assert(j < m && ans[l[i]] == r[j]);
      tmp[l[i]] = j;
    for (int i = 0; i < n; i += 2) {</pre>
      int curl = 0, curr = m - 1;
      if (i)
        curl = tmp[l[i - 1]];
      if (i + 1 < n)
        curr = tmp[l[i + 1]];
      11 res = query(1[i], r[curl]);
      ans[l[i]] = r[curl];
      for (int j = curl + 1; j <= curr; ++j) {</pre>
        11 nxt = query(l[i], r[j]);
        if (res < nxt)</pre>
          res = nxt, ans[l[i]] = r[j];
```

```
void reduce(vector <int> 1, vector <int> r) {
    int n = l.size(), m = r.size();
    vector <int> nr;
    for (int j : r) {
      while (!nr.empty()) {
        int i = nr.size() - 1;
        if (query(l[i], nr.back()) <= query(l[i], j))</pre>
          nr.pop_back();
        else
          break:
      if (nr.size() < n)</pre>
        nr.push_back(j);
    }
    run(1, nr);
  void run(vector <int> 1, vector <int> r) {
    int n = 1.size(), m = r.size();
    if (max(n, m) <= 2) {
      for (int i : 1) {
        ans[i] = r[0];
        if (m > 1) +
          if (query(i, r[0]) < query(i, r[1]))</pre>
            ans[i] = r[1];
      }
    } else if (n >= m) {
      interpolate(l, r);
    } else {
      reduce(1, r);
  }
};
```

9.11 Two Dimension Add Sum*

```
struct TwoDimensionAddAndSum {
  // 0-index, [l, r)
  struct Seg {
    int 1, r, m;
    ll vala, valb, lza, lzb;
    Seg* ch[2];
    Seg (int _1, int _r) : l(_1), r(_r), m(1 + r >> 1),
         vala(0), valb(0), lza(0), lzb(0) {
      if (r - 1 > 1) {
       ch[0] = new Seg(1, m);
        ch[1] = new Seg(m, r);
      }
    void pull() {vala = ch[0]->vala + ch[1]->vala, valb
        = ch[0]->valb + ch[1]->valb;}
    void give(ll a, ll b) {
      lza += a, lzb += b;
      vala += a * (r - 1), valb += b * (r - 1);
    void push() {
      ch[0]->give(lza, lzb), ch[1]->give(lza, lzb), lza
           = 1zb = 0;
    void add(int a, int b, ll va, ll vb) {
      if (a <= 1 && r <= b)
        give(va, vb);
      else {
        push();
        if (a < m) ch[0]->add(a, b, va, vb);
        if (m < b) ch[1]->add(a, b, va, vb);
        pull();
    long long query(int a, int b, int v) {
      if (a <= 1 && r <= b) return vala * v + valb;</pre>
      push();
      long long ans = 0;
      if (a < m) ans += ch[0]->query(a, b, v);
      if (m < b) ans += ch[1]->query(a, b, v);
      return ans;
   }
 }:
 // note integer overflow.
 vector <array <int, 4>> E[N];
 vector <array <int, 4>> Q[N];
```

```
vector <ll> ans:
  void add_event(int x1, int y1, int x2, int y2, ll v)
    E[x1].pb({y1, y2, v, -v * x1});
    E[x2].pb({y1, y2, -v, v * x2});
  void add_query(int x1, int y1, int x2, int y2, int id
    Q[x1].pb({y1, y2, -1, id});
    Q[x2].pb({y1, y2, 1, id});
    ans.pb(0);
  void solve(int n) {
    Seg root(0, n);
    for (int i = 0; i <= n; ++i) {
      for (auto j : E[i]) root.add(j[0], j[1], j[2], j
       for (auto j : Q[i]) ans[j[3]] += j[2] * root.
           query(j[0], j[1], i);
    }
  }
};
```

9.12 Matroid Intersection

Start from $S=\emptyset$. In each iteration, let

```
• Y_1 = \{x \notin S \mid S \cup \{x\} \in I_1\}
• Y_2 = \{x \notin S \mid S \cup \{x\} \in I_2\}
```

If there exists $x\in Y_1\cap Y_2$, insert x into S. Otherwise for each $x\in S, y\not\in S$, create edges

```
\begin{array}{ll} \bullet & x \rightarrow y \text{ if } S - \{x\} \cup \{y\} \in I_1\text{.} \\ \bullet & y \rightarrow x \text{ if } S - \{x\} \cup \{y\} \in I_2\text{.} \end{array}
```

Find a shortest path (with BFS) starting from a vertex in Y_1 and ending at a vertex in Y_2 which doesn't pass through any other vertices in Y_2 , and alternate the path. The size of S will be incremented by 1 in each iteration. For the weighted case, assign weight w(x) to vertex x if $x \in S$ and -w(x) if $x \not\in S$. Find the path with the minimum number of edges among all minimum length paths and alternate it.

9.13 Python Misc