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# 1 Basic

#### 1.1 Shell Script

#### 1.2 Debug Macro [0c06cd]

```
void db() { cout << endl; }
template <typename T, typename ...U>
void db(T i, U ...j) { cout << i << ' ', db(j...); }
#ifdef ABS
#define bug(x...) db("[" + string(#x) + "]", x)
#else
#define bug(x...) void(0)
#endif</pre>
```

### 1.3 Pragma / FastIO

```
#pragma GCC optimize("Ofast, no-stack-protector")
#pragma GCC optimize("no-math-errno,unroll-loops")
#pragma GCC target("sse,sse2,sse3,ssse3,sse4")
#pragma GCC target("popcnt,abm,mmx,avx,arch=skylake")
 _builtin_ia32_ldmxcsr(__builtin_ia32_stmxcsr()|0x8040)
#include<unistd.h>
char OB[65536]; int OP;
inline char RC() {
  static char buf[65536], *p = buf, *q = buf;
  return p == q \& (q = (p = buf) + read(0, buf, 65536)
      ) == buf ? -1 : *p++;
inline int R() {
  static char c;
  while((c = RC()) < '0'); int a = c ^ '0';
  while((c = RC()) >= '0') a *= 10, a += c ^ '0';
  return a;
inline void W(int n) {
  static char buf[12], p;
  if (n == 0) OB[OP++]='0'; p = 0;
while (n) buf[p++] = '0' + (n % 10), n /= 10;
  for (--p; p >= 0; --p) OB[OP++] = buf[p];
  if (OP > 65520) write(1, OB, OP), OP = 0;
```

# 1.4 Divide

```
11 floor(ll a, ll b) {return a / b - (a < 0 && a % b);}
11 ceil(ll a, ll b) {return a / b + (a > 0 && a % b);}
a / b < x -> floor(a, b) + 1 <= x
a / b <= x -> ceil(a, b) <= x
x < a / b -> x <= ceil(a, b) - 1
x <= a / b -> x <= floor(a, b)</pre>
```

#### 2 Data Structure

#### 2.1 Leftist Tree [414ab9]

```
struct node {
  11 rk, data, sz, sum;
  node *1, *r;
  node(11 k) : rk(0), data(k), sz(1), l(0), r(0), sum(k)
11 sz(node *p) { return p ? p->sz : 0; }
11 rk(node *p) { return p ? p->rk : -1;
11 sum(node *p) { return p ? p->sum : 0; }
node *merge(node *a, node *b) {
  if (!a || !b) return a ? a : b;
  if (a->data < b->data) swap(a, b);
  a->r = merge(a->r, b);
  if (rk(a->r) > rk(a->l)) swap(a->r, a->l);
  a - rk = rk(a - r) + 1:
  a->sz = sz(a->1) + sz(a->r) + 1;
  a\rightarrow sum = sum(a\rightarrow 1) + sum(a\rightarrow r) + a\rightarrow data;
  return a:
void pop(node *&o) {
  node *tmp = o;
  o = merge(o->1, o->r);
  delete tmp;
}
```

# 2.2 Splay Tree [21142b]

```
struct Splay {
   int pa[N], ch[N][2], sz[N], rt, _id;
   ll v[N];
   Splay() {}
   void init() {
      rt = 0, pa[0] = ch[0][0] = ch[0][1] = -1;
      sz[0] = 1, v[0] = inf;
   }
   int newnode(int p, int x) {
      int id = _id++;
      v[id] = x, pa[id] = p;
      ch[id][0] = ch[id][1] = -1, sz[id] = 1;
      return id;
   }
   void rotate(int i) {
```

```
int p = pa[i], x = ch[p][1] == i;
    int gp = pa[p], c = ch[i][!x];
    sz[p] -= sz[i], sz[i] += sz[p];
    if (~c) sz[p] += sz[c], pa[c] = p;
    ch[p][x] = c, pa[p] = i;
    pa[i] = gp, ch[i][!x] = p;
    if (~gp) ch[gp][ch[gp][1] == p] = i;
  void splay(int i) {
    while (~pa[i]) {
      int p = pa[i];
      if (~pa[p]) rotate(ch[pa[p]][1] == p ^ ch[p][1]
          == i ? i : p);
      rotate(i);
    }
    rt = i;
  int lower_bound(int x) {
    int i = rt, last = -1;
    while (true) {
      if (v[i] == x) return splay(i), i;
      if (v[i] > x) {
        last = i;
        if (ch[i][0] == -1) break;
        i = ch[i][0];
      else {
        if (ch[i][1] == -1) break;
        i = ch[i][1];
    splay(i);
    return last; // -1 if not found
  void insert(int x) {
    int i = lower_bound(x);
    if (i == -1) {
      // assert(ch[rt][1] == -1);
      int id = newnode(rt, x);
      ch[rt][1] = id, ++sz[rt];
      splay(id);
    else if (v[i] != x) {
      splay(i);
      int id = newnode(rt, x), c = ch[rt][0];
      ch[rt][0] = id;
      ch[id][0] = c;
      if (~c) pa[c] = id, sz[id] += sz[c];
      ++sz[rt];
      splay(id);
  }
};
      Link Cut Tree [bca367]
```

```
// weighted subtree size, weighted path max
struct LCT {
 int ch[N][2], pa[N], v[N], sz[N];
 int sz2[N], w[N], mx[N], _id;
 // sz := sum \ of \ v \ in \ splay, \ sz2 := sum \ of \ v \ in
      virtual subtree
  // mx := max w in splav
  bool rev[N];
 LCT() : _id(1) {}
 int newnode(int _v, int _w) {
   int x = _id++;
    ch[x][0] = ch[x][1] = pa[x] = 0;
    v[x] = sz[x] = v;
    sz2[x] = 0;
   w[x] = mx[x] = w;
    rev[x] = false;
    return x;
  void pull(int i) {
   sz[i] = v[i] + sz2[i];
    mx[i] = w[i];
    if (ch[i][0]) {
      sz[i] += sz[ch[i][0]];
      mx[i] = max(mx[i], mx[ch[i][0]]);
    if (ch[i][1]) {
```

```
sz[i] += sz[ch[i][1]];
       mx[i] = max(mx[i], mx[ch[i][1]]);
  }
  void push(int i) {
    if (rev[i]) reverse(ch[i][0]), reverse(ch[i][1]),
         rev[i] = false;
  void reverse(int i) {
    if (!i) return;
     swap(ch[i][0], ch[i][1]);
     rev[i] ^= true;
  bool isrt(int i) {// rt of splay
    if (!pa[i]) return true;
     return ch[pa[i]][0] != i && ch[pa[i]][1] != i;
  void rotate(int i) {
     int p = pa[i], x = ch[p][1] == i;
     int c = ch[i][!x], gp = pa[p];
     if (ch[gp][0] == p) ch[gp][0] = i;
     else if (ch[gp][1] == p) ch[gp][1] = i;
    pa[i] = gp, ch[i][!x] = p, pa[p] = i;
     ch[p][x] = c, pa[c] = p;
    pull(p), pull(i);
  void splay(int i) {
    vector<int> anc;
     anc.push_back(i);
    while (!isrt(anc.back()))
       anc.push_back(pa[anc.back()]);
     while (!anc.empty())
      push(anc.back()), anc.pop_back();
     while (!isrt(i)) {
       int p = pa[i];
       if (!isrt(p)) rotate(ch[p][1] == i ^ ch[pa[p]][1]
            == p ? i : p);
       rotate(i);
    }
  void access(int i) {
    int last = 0;
     while (i) {
       splay(i);
       if (ch[i][1])
       sz2[i] += sz[ch[i][1]];
sz2[i] -= sz[last];
       ch[i][1] = last;
       pull(i), last = i, i = pa[i];
  void makert(int i) {
    access(i), splay(i), reverse(i);
  void link(int i, int j) {
     // assert(findrt(i) != findrt(j));
    makert(i);
    makert(j);
    pa[i] = j;
sz2[j] += sz[i];
    pull(j);
  void cut(int i, int j) {
    makert(i), access(j), splay(i);
     // assert(sz[i] == 2 && ch[i][1] == j);
    ch[i][1] = pa[j] = 0, pull(i);
  int findrt(int i) {
    access(i), splay(i);
     while (ch[i][0]) push(i), i = ch[i][0];
     splay(i);
     return i;
};
2.4 Treap [9d5c2a]
struct node {
  int data, sz;
  node *1, *r;
  node(int k) : data(k), sz(1), l(0), r(0) {}
```

void up() {

```
sz = 1;
    if (1) sz += 1->sz;
    if (r) sz += r->sz;
  void down() {}
// delete default code sz
int sz(node *a) { return a ? a->sz : 0; }
node *merge(node *a, node *b) {
  if (!a || !b) return a ? a : b;
  if (rand() % (sz(a) + sz(b)) < sz(a))
    return a->down(), a->r = merge(a->r, b), a->up(),a;
  return b->down(), b->l = merge(a, b->l), b->up(), b;
void split(node *o, node *&a, node *&b, int k) {
  if (!o) return a = b = 0, void();
  o->down();
  if (o->data <= k)
    a = o, split(o->r, a->r, b, k), <math>a->up();
  else b = o, split(o->1, a, b->1, k), b->up();
void split2(node *o, node *&a, node *&b, int k) {
  if (sz(o) <= k) return a = o, b = 0, void();
  o->down();
  if (sz(o->1) + 1 <= k)
    a = o, split2(o->r, a->r, b, k - <math>sz(o->l) - 1);
  else b = o, split2(o->1, a, b->1, k);
  o->up();
node *kth(node *o, int k) {
  if (k <= sz(o->1)) return kth(o->1, k);
  if (k == sz(o\rightarrow 1) + 1) return o;
  return kth(o\rightarrow r, k - sz(o\rightarrow 1) - 1);
int Rank(node *o, int key) {
  if (!o) return 0;
  if (o->data < key)</pre>
    return sz(o\rightarrow 1) + 1 + Rank(o\rightarrow r, key);
  else return Rank(o->1, key);
bool erase(node *&o, int k) {
  if (!o) return 0;
  if (o->data == k) {
    node *t = o;
    o->down(), o = merge(o->1, o->r);
    delete t;
    return 1;
  node *&t = k < o \rightarrow data ? o \rightarrow 1 : o \rightarrow r;
  return erase(t, k) ? o->up(), 1 : 0;
void insert(node *&o, int k) {
 node *a, *b;
  o->down(), split(o, a, b, k),
  o = merge(a, merge(new node(k), b));
  o->up();
void interval(node *&o, int 1, int r) {
 node *a, *b, *c; // [l, r)
  o->down();
  split2(o, a, b, 1), split2(b, b, c, r - 1);
  // operate
  o = merge(a, merge(b, c)), o->up();
}
```

# 2.5 2D Segment Tree [23be4b]

```
if (!ch[0]) ch[0] = new seg(1, 1 + r >> 1), ch[1] =
           new seg(l + r >> 1, r);
     push();
     if (_1 < 1 + r >> 1) ch[0]->add(_1, _r, d);
     if (1 + r >> 1 < _r) ch[1]->add(_1, _r, d);
     pull();
  push();
     11 \text{ res} = 0;
     if (_1 < 1 + r >> 1) res += ch[0]->qsum(_1, _r);
     if (1 + r >> 1 < _r) res += ch[1]->qsum(_1, _r);
     return res;
   }
};
struct seg2 {
  int 1, r;
   seg v, lz;
   seg2 *ch[2]{};
   seg2(int _1, int _r) : l(_1), r(_r), v(0, N), lz(0, N
     if (1 < r - 1) ch[0] = new seg2(1, 1 + r >> 1), ch
          [1] = new seg2(1 + r >> 1, r);
   void add(int _1, int _r, int _12, int _r2, 11 d) {
  v.add(_12, _r2, d * (min(r, _r) - max(1, _1)));
     if (_1 <= 1 && r <= _r)
       return lz.add(_12, _r2, d), void(0);
     if (_l < l + r >> 1)
         ch[0]->add(_1, _r, _12, _r2, d);
     if (1 + r >> 1 < _r)
          ch[1]->add(_l, _r, _l2, _r2, d);
  il qsum(int _1, int _r, int _12, int _r2) {
   if (_1 <= 1 && r <= _r) return v.qsum(_12, _r2);
   ll d = min(r, _r) - max(1, _1);
   ll res = lz.qsum(_12, _r2) * d;
}</pre>
     if (_1 < 1 + r >> 1)
          res += ch[0]->qsum(_1, _r, _12, _r2);
     if (1 + r >> 1 < _r)
         res += ch[1]->qsum(_1, _r, _12, _r2);
     return res;
  }
};
```

# 2.6 vEB Tree [087d11]

```
using u64=uint64_t;
constexpr int lsb(u64 x)
{ return x?__builtin_ctzll(x):1<<30; }
constexpr int msb(u64 x)
{ return x?63-__builtin_clzll(x):-1; }
template<int N, class T=void>
struct veb{
  static const int M=N>>1;
  veb<M> ch[1<<N-M];</pre>
  veb<N-M> aux;
  int mn, mx;
  veb():mn(1<<30),mx(-1){}
  constexpr int mask(int x){return x&((1<<M)-1);}</pre>
  bool empty(){return mx==-1;}
  int min(){return mn;}
  int max(){return mx;}
  bool have(int x){
    return x==mn?true:ch[x>>M].have(mask(x));
  void insert_in(int x){
    if(empty()) return mn=mx=x,void();
    if(x<mn) swap(x,mn);</pre>
    if(x>mx) mx=x;
    if(ch[x>>M].empty()) aux.insert_in(x>>M);
    ch[x>>M].insert_in(mask(x));
  void erase_in(int x){
    if(mn==mx) return mn=1<<30, mx=-1, void();</pre>
    if(x==mn) mn=x=(aux.min()<<M)^ch[aux.min()].min();</pre>
    ch[x>>M].erase_in(mask(x));
    if(ch[x>>M].empty()) aux.erase_in(x>>M);
    if(x==mx){
```

```
if(aux.emptv()) mx=mn;
      else mx=(aux.max()<<M)^ch[aux.max()].max();</pre>
    }
  }
  void insert(int x){
    if(!have(x)) insert_in(x);
  void erase(int x){
    if(have(x)) erase_in(x);
  int next(int x){//} >= x
    if(x>mx) return 1<<30;
    if(x<=mn) return mn;</pre>
    if(mask(x)<=ch[x>>M].max())
      return ((x>>M)<<M)^ch[x>>M].next(mask(x));
    int y=aux.next((x>>M)+1);
    return (y<<M)^ch[y].min();</pre>
  int prev(int x){// <x</pre>
    if(x<=mn) return -1;</pre>
    if(x>mx) return mx;
    if(x<=(aux.min()<<M)+ch[aux.min()].min())</pre>
      return mn;
    if(mask(x)>ch[x>>M].min())
      return ((x>>M)<<M)^ch[x>>M].prev(mask(x));
    int y=aux.prev(x>>M);
    return (y<<M)^ch[y].max();</pre>
 }
template<int N>
struct veb<N,typename enable_if<N<=6>::type>{
  u64 a;
  veb():a(0){}
  void insert_in(int x){a|=1ull<<x;}</pre>
  void insert(int x){a|=1ull<<x;}</pre>
 void erase_in(int x){a&=~(1ull<<x);}</pre>
  void erase(int x){a&=~(1ull<<x);}</pre>
  bool have(int x){return a>>x&1;}
  bool empty(){return a==0;}
  int min(){return lsb(a);}
 int max(){return msb(a);}
 int next(int x){return lsb(a&~((1ull<<x)-1));}</pre>
  int prev(int x){return msb(a&((1ull<<x)-1));}</pre>
```

#### 2.7 Range Set [a1cb88]

```
struct RangeSet { // [l, r)
 set <pii>> S:
  void cut(int x) {
    auto it = S.lower_bound(\{x + 1, -1\});
    if (it == S.begin()) return;
    auto [1, r] = *prev(it);
    if (1 >= x \mid | x >= r) return;
   S.erase(prev(it));
    S.insert({1, x});
    S.insert({x, r});
  vector <pii> split(int l, int r) {
   // remove and return ranges in [l, r)
    cut(1), cut(r);
    vector <pii> res;
    while (true) {
      auto it = S.lower_bound({1, -1});
      if (it == S.end() || r <= it->first) break;
      res.pb(*it), S.erase(it);
   }
    return res;
  void insert(int 1, int r) {
   // add a range [l, r), [l, r) not in S
    auto it = S.lower_bound({1, r});
    if (it != S.begin() && prev(it)->second == 1)
      1 = prev(it)->first, S.erase(prev(it));
    if (it != S.end() && r == it->first)
     r = it->second, S.erase(it);
    S.insert({1, r});
  bool count(int x) {
    auto it = S.lower_bound(\{x + 1, -1\});
    return it != S.begin() && prev(it)->first <= x</pre>
            && x < prev(it)->second;
```

```
};
```

Flow / Matching

# 3.1 Dinic [8898fb]

}

3

```
template <typename T>
struct Dinic { // 0-based
  const T INF = numeric_limits<T>::max() / 2;
  struct edge { int to, rev; T cap, flow; };
  int n, s, t;
  vector <vector <edge>> g;
  vector <int> dis, cur;
  T dfs(int u, T cap) {
     if (u == t || !cap) return cap;
     for (int &i = cur[u]; i < (int)g[u].size(); ++i) {</pre>
       edge &e = g[u][i];
       if (dis[e.to] == dis[u] + 1 && e.flow != e.cap) {
         T df = dfs(e.to, min(e.cap - e.flow, cap));
         if (df) {
           e.flow += df:
           g[e.to][e.rev].flow -= df;
           return df;
      }
    dis[u] = -1;
    return 0;
  bool bfs() {
     fill(all(dis), -1);
     queue<int> q;
     q.push(s), dis[s] = 0;
     while (!q.empty()) {
       int v = q.front(); q.pop();
       for (auto &u : g[v])
         if (!~dis[u.to] && u.flow != u.cap) {
           q.push(u.to);
           dis[u.to] = dis[v] + 1;
         }
     return dis[t] != -1;
  T solve(int _s, int _t) {
    s = _s, t = _t;
T flow = 0, df;
     while (bfs()) {
       fill(all(cur), 0);
       while ((df = dfs(s, INF))) flow += df;
    return flow;
  }
  void reset() {
     for (int i = 0; i < n; ++i)</pre>
       for (auto &j : g[i]) j.flow = 0;
  void add_edge(int u, int v, T cap) {
   g[u].pb(edge{v, (int)g[v].size(), cap, 0});
   g[v].pb(edge{u, (int)g[u].size() - 1, 0, 0});
  Dinic (int _n) : n(_n), g(n), dis(n), cur(n) {}
};
```

# 3.2 Min Cost Max Flow [8083d7]

```
template <typename T1, typename T2>
struct MCMF { // T1 -> flow, T2 -> cost, 0-based
   const T1 INF1 = numeric_limits<T1>::max() / 2;
   const T2 INF2 = numeric_limits<T2>::max() / 2;
   struct edge {
     int v; T1 f; T2 c;
   };
   int n, s, t;
   vector <vector <int>> g;
   vector <vedge> e;
   vector <T2> dis, pot;
   vector <int> rt, vis;
   // bool DAG()...
   bool SPFA() {
     fill(all(rt), -1), fill(all(dis), INF2);
     fill(all(vis), false);
```

```
queue <int> q;
  q.push(s), dis[s] = 0, vis[s] = true;
  while (!q.empty()) {
    int v = q.front(); q.pop();
    vis[v] = false;
    for (int id : g[v]) {
      auto [u, f, c] = e[id];
      T2 ndis = dis[v] + c + pot[v] - pot[u];
      if (f > 0 && dis[u] > ndis) {
        dis[u] = ndis, rt[u] = id;
        if (!vis[u]) vis[u] = true, q.push(u);
      }
    }
  }
  return dis[t] != INF2;
bool dijkstra() {
  fill(all(rt), -1), fill(all(dis), INF2);
  priority_queue <pair <T2, int>, vector <pair <T2,</pre>
      int>>, greater <pair <T2, int>>> pq;
  dis[s] = 0, pq.emplace(dis[s], s);
  while (!pq.empty()) {
    auto [d, v] = pq.top(); pq.pop();
    if (dis[v] < d) continue;</pre>
    for (int id : g[v]) {
      auto [u, f, c] = e[id];
      T2 ndis = dis[v] + c + pot[v] - pot[u];
      if (f > 0 \&\& dis[u] > ndis) {
        dis[u] = ndis, rt[u] = id;
        pq.emplace(ndis, u);
    }
  }
  return dis[t] != INF2;
vector <pair <T1, T2>> solve(int _s, int _t) {
   s = _s, t = _t, fill(all(pot), 0);
  vector <pair <T1, T2>> ans; bool fr = true;
  while ((fr ? SPFA() : dijkstra())) {
    for (int i = 0; i < n; i++)</pre>
      dis[i] += pot[i] - pot[s];
    T1 add = INF1;
    for (int i = t; i != s; i = e[rt[i] ^ 1].v)
      add = min(add, e[rt[i]].f);
    for (int i = t; i != s; i = e[rt[i] ^ 1].v)
      e[rt[i]].f -= add, e[rt[i] ^ 1].f += add;
    ans.emplace_back(add, dis[t]), fr = false;
    for (int i = 0; i < n; ++i) swap(dis[i], pot[i]);</pre>
  return ans;
void reset() {
  for (int i = 0; i < (int)e.size(); ++i) e[i].f = 0;</pre>
void add_edge(int u, int v, T1 f, T2 c) {
  g[u].pb((int)e.size()), e.pb({v, f, c});
  g[v].pb((int)e.size()), e.pb({u, 0, -c});
MCMF (int _n): n(_n), g(n), e(), dis(n), pot(n),
  rt(n), vis(n) {}
```

# 3.3 Kuhn Munkres [b880ad]

```
template <typename T>
struct KM { // 0-based, remember to init
 const T INF = numeric_limits<T>::max() / 2;
 int n; vector <vector <T>> w;
 vector <T> hl, hr, slk;
 vector <int> fl, fr, vl, vr, pre;
 queue <int> q;
 bool check(int x) {
   if (vl[x] = 1, \sim fl[x])
      return q.push(fl[x]), vr[fl[x]] = 1;
    while (\sim x) swap(x, fr[fl[x] = pre[x]]);
    return 0;
 void bfs(int s) {
   fill(all(slk), INF), fill(all(vl), 0);
    fill(all(vr), 0);
    while (!q.empty()) q.pop();
    q.push(s), vr[s] = 1;
```

```
while (true) {
    T d;
    while (!q.empty()) {
      int y = q.front(); q.pop();
      for (int x = 0; x < n; ++x) {
        d = h1[x] + hr[y] - w[x][y];
        if (!v1[x] \&\& s1k[x] >= d) {
          if (pre[x] = y, d) slk[x] = d;
          else if (!check(x)) return;
        }
     }
    }
    d = INF;
    for (int x = 0; x < n; ++x)
     if (!vl[x] && d > slk[x]) d = slk[x];
    for (int x = 0; x < n; ++x) {
      if (v1[x]) h1[x] += d;
      else slk[x] -= d;
      if (vr[x]) hr[x] -= d;
    for (int x = 0; x < n; ++x)
      if (!v1[x] && !s1k[x] && !check(x)) return;
T solve() {
  fill(all(fl), -1), fill(all(fr), -1);
  fill(all(hr), 0);
  for (int i = 0; i < n; ++i)</pre>
    hl[i] = *max_element(all(w[i]));
  for (int i = 0; i < n; ++i) bfs(i);</pre>
  T res = 0:
  for (int i = 0; i < n; ++i) res += w[i][fl[i]];</pre>
 return res:
void add_edge(int a, int b, T wei) { w[a][b] = wei; }
KM (int _n) : n(_n), w(n, vector < T > (n, -INF)), hl(n),
 hr(n), slk(n), fl(n), fr(n), vl(n), vr(n), pre(n){}
```

```
3.4 Hopcroft Karp [33c68d]
struct HopcroftKarp { // 0-based
  const int INF = 1 << 30;</pre>
  int n, m;
  vector <int>> g;
  vector <int> match, dis, matched, vis;
  bool dfs(int x) {
    vis[x] = true;
    for (int y : g[x])
      if (match[y] == -1 || (dis[match[y]] == dis[x] +
          1 && !vis[match[y]] && dfs(match[y]))) {
        match[y] = x, matched[x] = true;
        return true;
    return false:
  bool bfs() {
    fill(all(dis), -1);
    queue <int> q;
    for (int x = 0; x < n; ++x) if (!matched[x])
      dis[x] = 0, q.push(x);
    int mx = INF;
    while (!q.empty()) {
      int x = q.front(); q.pop();
      for (int y : g[x]) {
        if (match[y] == -1) {
          mx = dis[x];
          break;
        } else if (dis[match[y]] == -1)
          dis[match[y]] = dis[x] + 1, q.push(match[y]);
     }
    }
    return mx < INF;</pre>
  int solve() {
    int res = 0;
    fill(all(match), -1);
    fill(all(matched), 0);
    while (bfs()) {
      fill(all(vis), 0);
```

for (int x = 0; x < n; ++x) if (!matched[x])

res += dfs(x);

}

```
void blossom(int x, int y, int 1) {
    return res;
                                                               while (Find(x) != 1) {
  void add_edge(int x, int y) { g[x].pb(y); }
                                                                 pre[x] = y, y = match[x];
 HopcroftKarp (int _n, int _m) : n(_n), m(_m), g(n),
                                                                 if (s[y] == 1) q.push(y), s[y] = 0;
    match(m), dis(n), matched(n), vis(n) {}
                                                                 if (fa[x] == x) fa[x] = 1;
                                                                 if (fa[y] == y) fa[y] = 1;
                                                                 x = pre[y];
3.5 SW Min Cut [b9af94]
                                                               }
template <typename T>
                                                             bool bfs(int r) {
struct SW { // 0-based
                                                               iota(all(fa), 0), fill(all(s), -1);
  const T INF = numeric_limits<T>::max() / 2;
                                                               while (!q.empty()) q.pop();
  vector <vector <T>> g;
                                                               q.push(r);
  vector <T> sum;
                                                               s[r] = 0;
  vector <bool> vis, dead;
                                                               while (!q.empty()) {
 int n:
                                                                 int x = q.front(); q.pop();
 T solve() -
                                                                 for (int u : g[x]) {
    T ans = INF;
                                                                   if (s[u] == -1) {
    for (int r = 0; r + 1 < n; ++r) {
                                                                     pre[u] = x, s[u] = 1;
      fill(all(vis), 0), fill(all(sum), 0);
                                                                     if (match[u] == n) {
      int num = 0, s = -1, t = -1;
                                                                       for (int a = u, b = x, last; b != n; a =
      while (num < n - r) {
                                                                            last, b = pre[a])
        int now = -1;
                                                                         last = match[b], match[b] = a, match[a] =
        for (int i = 0; i < n; ++i)</pre>
                                                                              b:
          if (!vis[i] && !dead[i] &&
                                                                       return true;
            (now == -1 \mid \mid sum[now] > sum[i])) now = i;
        s = t, t = now;
                                                                     q.push(match[u]);
        vis[now] = true, num++;
                                                                     s[match[u]] = 0;
        for (int i = 0; i < n; ++i)</pre>
                                                                   } else if (!s[u] && Find(u) != Find(x)) {
          if (!vis[i] && !dead[i]) sum[i] += g[now][i];
                                                                     int 1 = lca(u, x);
                                                                     blossom(x, u, 1);
      ans = min(ans, sum[t]);
                                                                     blossom(u, x, 1);
      for (int i = 0; i < n; ++i)</pre>
        g[i][s] += g[i][t], g[s][i] += g[t][i];
                                                                 }
      dead[t] = true;
                                                               }
                                                               return false;
    return ans;
                                                             int solve() {
  void add_edge(int u, int v, T w) {
                                                               int res = 0;
    g[u][v] += w, g[v][u] += w; }
                                                               for (int x = 0; x < n; ++x) {
  SW (int _n) : n(_n), g(n, vector <T>(n)), vis(n),
                                                                 if (match[x] == n) res += bfs(x);
    sum(n), dead(n) {}
                                                               return res;
3.6 Gomory Hu Tree [90ead2]
                                                             void add_edge(int u, int v) {
                                                               g[u].push_back(v), g[v].push_back(u);
vector <array <int, 3>> GomoryHu(Dinic <int> flow) {
  // Tree edge min = mincut (0-based)
                                                             Matching (int _n): n(_n), tk(0), g(n), fa(n + 1),
  int n = flow.n;
                                                               pre(n + 1, n), match(n + 1, n), s(n + 1), t(n) {}
  vector <array <int, 3>> ans;
  vector <int> rt(n);
  for (int i = 1; i < n; ++i) {</pre>
                                                           3.8 Min Cost Circulation [bd1e15]
    int t = rt[i];
    flow.reset();
                                                           struct MinCostCirculation { // 0-base
    ans.pb({i, t, flow.solve(i, t)});
                                                             struct Edge {
    flow.bfs();
                                                               ll from, to, cap, fcap, flow, cost, rev;
    for (int j = i + 1; j < n; ++j)
                                                             } *past[N];
      if (rt[j] == t && flow.dis[j] != -1) rt[j] = i;
                                                             vector<Edge> G[N];
                                                             11 dis[N], inq[N], n;
  return ans;
                                                             void BellmanFord(int s) {
                                                               fill_n(dis, n, INF), fill_n(inq, n, 0);
3.7 Blossom [6092d8]
                                                               queue<int> q;
                                                               auto relax = [&](int u, ll d, Edge *e) {
struct Matching { // O-based
                                                                 if (dis[u] > d) {
                                                                   dis[u] = d, past[u] = e;
 int n, tk;
  vector <vector <int>> g;
                                                                   if (!inq[u]) inq[u] = 1, q.push(u);
  vector <int> fa, pre, match, s, t;
                                                                 }
  queue <int> q;
  int Find(int u) {
                                                               relax(s, 0, 0);
    return u == fa[u] ? u : fa[u] = Find(fa[u]);
                                                               while (!q.empty()) {
                                                                 int u = q.front();
 int lca(int x, int y) {
                                                                 q.pop(), inq[u] = 0;
    tk++;
                                                                 for (auto &e : G[u])
    x = Find(x), y = Find(y);
                                                                   if (e.cap > e.flow)
    for (; ; swap(x, y)) {
                                                                     relax(e.to, dis[u] + e.cost, &e);
      if (x != n) {
        if (t[x] == tk) return x;
                                                             void try_edge(Edge &cur) {
        t[x] = tk:
        x = Find(pre[match[x]]);
                                                               if (cur.cap > cur.flow) return ++cur.cap, void();
```

BellmanFord(cur.to);

if (dis[cur.from] + cur.cost < 0) {</pre>

```
++cur.flow, --G[cur.to][cur.rev].flow;
       for (int i = cur.from; past[i]; i = past[i]->from
           ) {
         auto &e = *past[i];
         ++e.flow, --G[e.to][e.rev].flow;
    }
    ++cur.cap;
  }
  void solve(int mxlg) {
    for (int b = mxlg; b >= 0; --b) {
      for (int i = 0; i < n; ++i)
         for (auto &e : G[i])
      e.cap *= 2, e.flow *= 2;
for (int i = 0; i < n; ++i)
         for (auto &e : G[i])
           if (e.fcap >> b & 1)
             try_edge(e);
    }
  }
  void init(int _n) { n = _n;
    for (int i = 0; i < n; ++i) G[i].clear();</pre>
  void add_edge(ll a, ll b, ll cap, ll cost) {
    G[a].pb(Edge{a, b, 0, cap, 0, cost, sz(G[b]) + (a)}
         == b));
    G[b].pb(Edge{b, a, 0, 0, 0, -cost, sz(G[a]) - 1});
  }
|} mcmf; // O(VE * ELogC)
```

#### 3.9 Flow Model

- Maximum/Minimum flow with lower bound / Circulation problem
  - 1. Construct super source  ${\cal S}$  and sink  ${\cal T}$  .
  - 2. For each edge (x,y,l,u), connect x o y with capacity u-l.
  - 3. For each vertex  $\boldsymbol{v}$ , denote by  $in(\boldsymbol{v})$  the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.
  - 4. If in(v)>0, connect  $S\to v$  with capacity in(v), otherwise, connect  $v\to T$  with capacity -in(v).
    - To maximize, connect t o s with capacity  $\infty$  (skip this in circulation problem), and let f be the maximum flow from S to T. If  $f \neq \sum_{v \in V, in(v) > 0} in(v)$ , there's no solution. Otherwise, the maximum flow from s to t is
    - the answer. – To minimize, let f be the maximum flow from S to T . Connect  $t \to s$  with capacity  $\infty$  and let the flow from Sto T be f' . If  $f+f' 
      eq \sum_{v \in V, in(v)>0} in(v)$  , there's no solution. Otherwise, f' is the answer.
  - 5. The solution of each edge e is  $l_e+f_e$ , where  $f_e$  corresponds to the flow of edge e on the graph.
- ullet Construct minimum vertex cover from maximum matching M on bipartite graph (X,Y)
  - 1. Redirect every edge:  $y \to x$  if  $(x,y) \in M$ ,  $x \to y$  otherwise.

  - 2. DFS from unmatched vertices in X. 3.  $x \in X$  is chosen iff x is unvisited.
  - 4.  $y \in Y$  is chosen iff y is visited.
- Minimum cost cyclic flow
  - 1. Consruct super source S and sink T
  - 2. For each edge (x, y, c), connect  $x \to y$  with (cost, cap) = (c, 1)if c>0, otherwise connect  $y\to x$  with (cost, cap)=(-c,1)
  - 3. For each edge with c < 0, sum these cost as K, then increase
  - d(y) by 1, decrease d(x) by 1 4. For each vertex v with d(v)>0, connect  $S\to v$  with (cost, cap)=(0,d(v))
  - 5. For each vertex v with d(v) < 0, connect v o T with
  - (cost, cap) = (0, -d(v)) 6. Flow from S to T , the answer is the cost of the flow C+K
- Maximum density induced subgraph

  - 1. Binary search on answer, suppose we're checking answer T 2. Construct a max flow model, let K be the sum of all weights 3. Connect source  $s \to v$ ,  $v \in G$  with capacity K 4. For each edge (u,v,w) in G, connect  $u \to v$  and  $v \to u$  with capacity  $\boldsymbol{w}$
  - 5. For  $v\in G$ , connect it with sink  $v\to t$  with capacity  $K+2T-(\sum_{e\in E(v)}w(e))-2w(v)$
  - 6. T is a valid answer if the maximum flow f < K |V|
- Minimum weight edge cover
  - 1. Change the weight of each edge to  $\mu(u)+\mu(v)-w(u,v)$ , where  $\mu(v)$  is the cost of the cheapest edge incident to v.

    2. Let the maximum weight matching of the graph be x, the answer
  - will be  $\sum \mu(v) x$ .
- Project selection problem
  - 1. If  $p_v>0$ , create edge (s,v) with capacity  $p_v$ ; otherwise, create edge (v,t) with capacity  $-p_v$ .
  - 2. Create edge (u,v) with capacity w with w being the cost of choosing u without choosing v.

- 3. The mincut is equivalent to the maximum profit of a subset of projects.
- 0/1 quadratic programming

$$\sum_{x} c_{x}x + \sum_{y} c_{y}\bar{y} + \sum_{xy} c_{xy}x\bar{y} + \sum_{xyx'y'} c_{xyx'y'}(x\bar{y} + x'\bar{y'})$$

can be minimized by the mincut of the following graph:

- 1. Create edge (x,t) with capacity  $c_x$  and create edge (s,y) with capacity  $c_y$  .
- 2. Create edge (x,y) with capacity  $c_{xy}$
- 3. Create edge (x,y) and edge (x',y') with capacity  $c_{xyx'y'}$ .

# 4 Graph

# 4.1 Heavy-Light Decomposition [9ec77f]

```
struct HLD { // 0-based, remember to build
 int n, _id;
vector <vector <int>> g;
  vector <int> dep, pa, tsz, ch, hd, id;
  void dfs(int v, int p) {
    dep[v] = \sim p ? dep[p] + 1 : 0;
    pa[v] = p, tsz[v] = 1, ch[v] = -1;
    for (int u : g[v]) if (u != p) {
      dfs(u, v);
      if (ch[v] == -1 || tsz[ch[v]] < tsz[u])</pre>
        ch[v] = u;
      tsz[v] += tsz[u];
    }
  void hld(int v, int p, int h) {
    hd[v] = h, id[v] = _id++;
    if (~ch[v]) hld(ch[v], v, h);
    for (int u : g[v]) if (u != p && u != ch[v])
      hld(u, v, u);
  vector <pii> query(int u, int v) {
    vector <pii> ans;
    while (hd[u] != hd[v]) {
      if (dep[hd[u]] > dep[hd[v]]) swap(u, v);
      ans.emplace_back(id[hd[v]], id[v] + 1);
      v = pa[hd[v]];
    if (dep[u] > dep[v]) swap(u, v);
    ans.emplace_back(id[u], id[v] + 1);
    return ans;
  void build() {
    for (int i = 0; i < n; ++i) if (id[i] == -1)</pre>
      dfs(i, -1), hld(i, -1, i);
  void add_edge(int u, int v) {
    g[u].pb(v), g[v].pb(u); }
  HLD (int _n) : n(_n), _id(0), g(n), dep(n), pa(n),
    tsz(n), ch(n), hd(n), id(n, -1) {}
```

# 4.2 Centroid Decomposition [28b80a]

```
struct CD { // 0-based, remember to build
  int n, lg; // pa, dep are centroid tree attributes
  vector <vector <int>> g, dis;
  vector <int> pa, tsz, dep, vis;
  void dfs1(int v, int p) {
    tsz[v] = 1;
    for (int u : g[v]) if (u != p && !vis[u])
      dfs1(u, v), tsz[v] += tsz[u];
  int dfs2(int v, int p, int _n) {
    for (int u : g[v])
      if (u != p && !vis[u] && tsz[u] > _n / 2)
        return dfs2(u, v, _n);
    return v;
  void dfs3(int v, int p, int d) {
    dis[v][d] = \sim p ? dis[p][d] + 1 : 0;
    for (int u : g[v]) if (u != p && !vis[u])
      dfs3(u, v, d);
  void cd(int v, int p, int d) {
    dfs1(v, -1), v = dfs2(v, -1, tsz[v]);
vis[v] = true, pa[v] = p, dep[v] = d;
```

```
dfs3(v, -1, d);
  for (int u : g[v]) if (!vis[u])
    cd(u, v, d + 1);
}
void build() { cd(0, -1, 0); }
void add_edge(int u, int v) {
    g[u].pb(v), g[v].pb(u); }
CD (int _n) : n(_n), lg(__lg(n) + 1), g(n),
    dis(n, vector <int>(lg)), pa(n), tsz(n),
    dep(n), vis(n) {}
};
```

#### 4.3 Edge BCC [cf5e55]

```
struct EBCC { // 0-based, remember to build
 int n, m, nbcc;
 vector <vector <pii>>> g;
  vector <int> pa, low, dep, bcc_id, stk, is_bridge;
 void dfs(int v, int p, int f)
   low[v] = dep[v] = \sim p ? dep[p] + 1 : 0;
    stk.pb(v), pa[v] = p;
    for (auto [u, e] : g[v]) {
      if (low[u] == -1)
      dfs(u, v, e), low[v] = min(low[v], low[u]);
else if (e != f)
        low[v] = min(low[v], dep[u]);
    if (low[v] == dep[v]) {
      if (~f) is_bridge[f] = true;
      int id = nbcc++, x;
      do {
        x = stk.back(), stk.pop_back();
        bcc_id[x] = id;
      } while (x != v);
   }
 }
  void build() {
    is_bridge.assign(m, 0);
    for (int i = 0; i < n; ++i) if (low[i] == -1)</pre>
      dfs(i, -1, -1);
 void add_edge(int u, int v) {
   g[u].emplace_back(v, m), g[v].emplace_back(u, m++);
 EBCC (int _n): n(_n), m(0), nbcc(0), g(n), pa(n),
    low(n, -1), dep(n), bcc_id(n), stk() {}
```

# 4.4 Vertex BCC / Round Square Tree [3818e9]

```
struct BCC { // 0-based, remember to build
 int n, nbcc; // note for isolated point
  vector <vector <int>> g, _g; // id >= n: bcc
  vector <int> pa, dep, low, stk, pa2, dep2;
 void dfs(int v, int p) {
    dep[v] = low[v] = \sim p ? dep[p] + 1 : 0;
    stk.pb(v), pa[v] = p;
    for (int u : g[v]) if (u != p) {
      if (low[u] == -1) {
        dfs(u, v), low[v] = min(low[v], low[u]);
        if (low[u] >= dep[v]) {
          int id = nbcc++, x;
          do {
            x = stk.back(), stk.pop_back();
             g[id + n].pb(x), g[x].pb(id + n);
          } while (x != u);
          _g[id + n].pb(v), _g[v].pb(id + n);
      } else low[v] = min(low[v], dep[u]);
   }
  bool is_cut(int x) { return (int)_g[x].size() != 1; }
 vector <int> bcc(int id) { return _g[id + n]; }
 int bcc_id(int u, int v) {
    return pa2[dep2[u] < dep2[v] ? v : u] - n; }</pre>
  void dfs2(int v, int p) {
    dep2[v] = \sim p ? dep2[p] + 1 : 0, pa2[v] = p;
    for (int u : _g[v]) if (u != p) dfs2(u, v);
  void build() {
    low.assign(n, -1);
    for (int i = 0; i < n; ++i) if (low[i] == -1)</pre>
```

```
dfs(i, -1), dfs2(i, -1);
}
void add_edge(int u, int v) {
   g[u].pb(v), g[v].pb(u); }
BCC (int _n) : n(_n), nbcc(0), g(n), _g(2 * n),
   pa(n), dep(n), low(n), stk(), pa2(n * 2),
   dep2(n * 2) {}
};
```

# 4.5 SCC [9bee8c]

```
struct SCC {
  int n, nscc, _id;
  vector <vector <int>> g;
  vector <int> dep, low, scc_id, stk;
  void dfs(int v) {
    dep[v] = low[v] = _id++, stk.pb(v);
for (int u : g[v]) if (scc_id[u] == -1) {
   if (low[u] == -1) dfs(u);
       low[v] = min(low[v], low[u]);
     if (low[v] == dep[v]) {
       int id = nscc++, x;
         x = stk.back(), stk.pop_back(), scc_id[x] = id;
       } while (x != v);
    }
  }
  void build() {
    for (int i = 0; i < n; ++i) if (low[i] == -1)</pre>
       dfs(i);
  void add_edge(int u, int v) { g[u].pb(v); }
  SCC (int _n) : n(_n), nscc(0), _id(0), g(n), dep(n),
     low(n, -1), scc_id(n, -1), stk() {}
};
```

#### **4.6 2SAT** [938072]

```
struct SAT { // 0-based, need SCC
  int n; vector <pii> edge; vector <int> is;
  int rev(int x) { return x < n ? x + n : x - n; }</pre>
  void add_ifthen(int x, int y) {
    add_clause(rev(x), y); }
  void add_clause(int x, int y) {
    edge.emplace_back(rev(x), y);
    edge.emplace_back(rev(y), x); }
  bool solve() {
    // is[i] = true -> i, is[i] = false -> -i
    SCC scc(2 * n);
    for (auto [u, v] : edge) scc.add_edge(u, v);
    scc.build();
    for (int i = 0; i < n; ++i) {</pre>
      if (scc.scc_id[i] == scc.scc_id[i + n])
        return false;
      is[i] = scc.scc_id[i] < scc.scc_id[i + n];</pre>
    }
    return true;
  SAT (int _n) : n(_n), edge(), is(n) {}
};
```

#### 4.7 Virtual Tree [9e4a93]

```
// need lca, in, out
vector <pii> virtual_tree(vector <int> &v) {
  auto cmp = [&](int x, int y) {return in[x] < in[y];};</pre>
  sort(all(v), cmp);
  int sz = (int)v.size();
  for (int i = 0; i + 1 < sz; ++i)
    v.pb(lca(v[i], v[i + 1]));
  sort(all(v), cmp);
  v.resize(unique(all(v)) - v.begin());
  vector <int> stk(1, v[0]);
  vector <pii> res;
  for (int i = 1; i < (int)v.size(); ++i) {</pre>
    int x = v[i];
    while (out[stk.back()] < out[x]) stk.pop_back();</pre>
    res.emplace_back(stk.back(), x), stk.pb(x);
  return res;
```

#### 4.8 Directed MST [d6cf86]

```
using D = int;
struct edge { int u, v; D w; };
// 0-based, return index of edges
vector<int> dmst(vector<edge> &e, int n, int root) {
 using T = pair <D, int>;
  using PQ = pair <pri>priority_queue <T, vector <T>,
      greater <T>>, D>;
  auto push = [](PQ &pq, T v) {
   pq.first.emplace(v.first - pq.second, v.second);
  auto top = [](const PQ &pq) -> T {
    auto r = pq.first.top();
    return {r.first + pq.second, r.second};
  auto join = [&push, &top](PQ &a, PQ &b) {
   if (a.first.size() < b.first.size()) swap(a, b);</pre>
    while (!b.first.empty())
      push(a, top(b)), b.first.pop();
  vector<PQ> h(n * 2);
  for (int i = 0; i < e.size(); ++i)</pre>
  push(h[e[i].v], {e[i].w, i});
vector<int> a(n * 2), v(n * 2, -1), pa(n * 2, -1), r(
      n * 2);
  iota(all(a), 0);
  auto o = [\&](int x) \{ int y;
    for (y = x; a[y] != y; y = a[y]);
    for (int ox = x; x != y; ox = x)
      x = a[x], a[ox] = y;
    return y;
  };
  v[root] = n + 1;
  int pc = n;
  for (int i = 0; i < n; ++i) if (v[i] == -1) {</pre>
    for (int p = i; v[p] == -1 || v[p] == i; p = o(e[r[
        p]].u)) {
      if (v[p] == i) {
        int q = p; p = pc++;
        do {
          h[q].second = -h[q].first.top().first;
          join(h[pa[q] = a[q] = p], h[q]);
        } while ((q = o(e[r[q]].u)) != p);
      v[p] = i;
      while (!h[p].first.empty() && o(e[top(h[p]).
          second].u) == p)
        h[p].first.pop();
      r[p] = top(h[p]).second;
   }
 vector<int> ans;
for (int i = pc - 1; i >= 0; i--)
    if (i != root && v[i] != n) {
      for (int f = e[r[i]].v; f != -1 && v[f] != n; f =
    pa[f]) v[f] = n;
      ans.pb(r[i]);
    }
  return ans;
```

#### 4.9 Dominator Tree [49fc16]

```
struct DominatorTree {
 int n, id;
  vector <vector <int>> g, rg, bucket;
  vector <int> sdom, dom, vis, rev, pa, rt, mn, res;
  // dom[s] = s, dom[v] = -1 if s \rightarrow v not exists
 int query(int v, int x) {
    if (rt[v] == v) return x ? -1 : v;
    int p = query(rt[v], 1);
    if (p == -1) return x ? rt[v] : mn[v];
    if (sdom[mn[v]] > sdom[mn[rt[v]]])
      mn[v] = mn[rt[v]];
    rt[v] = p;
    return x ? p : mn[v];
 void dfs(int v) {
   vis[v] = id, rev[id] = v;
    rt[id] = mn[id] = sdom[id] = id, id++;
    for (int u : g[v]) {
      if (vis[u] == -1) dfs(u), pa[vis[u]] = vis[v];
```

```
rg[vis[u]].pb(vis[v]);
  }
  void build(int s) {
     dfs(s);
     for (int i = id - 1; ~i; --i) {
       for (int u : rg[i]) {
         sdom[i] = min(sdom[i], sdom[query(u, 0)]);
       if (i) bucket[sdom[i]].pb(i);
       for (int u : bucket[i]) {
         int p = query(u, 0);
         dom[u] = sdom[p] == i ? i : p;
       if (i) rt[i] = pa[i];
     fill(all(res), -1);
for (int i = 1; i < id; ++i) {
       if (dom[i] != sdom[i]) dom[i] = dom[dom[i]];
     for (int i = 1; i < id; ++i)</pre>
         res[rev[i]] = rev[dom[i]];
     res[s] = s;
     for (int i = 0; i < n; ++i) dom[i] = res[i];</pre>
  void add_edge(int u, int v) { g[u].pb(v); }
Dominator_tree (int _n) : n(_n), id(0), g(n), rg(n),
     bucket(n), sdom(n), dom(n, -1), vis(n, -1),
     rev(n), pa(n), rt(n), mn(n), res(n) {}
};
```

#### 4.10 Bipartite Edge Coloring [a22d96]

```
struct BipartiteEdgeColoring { // 1-based
  // returns edge coloring in adjacent matrix G
  int n, m;
  vector <vector <int>> col, G;
  int find_col(int x) {
    int c = 1;
    while (col[x][c]) c++;
    return c;
  void dfs(int v, int c1, int c2) {
    if (!col[v][c1]) return col[v][c2] = 0, void(0);
    int u = col[v][c1];
    dfs(u, c2, c1);
    col[v][c1] = 0, col[v][c2] = u, col[u][c2] = v;
  void solve() {
    for (int i = 1; i <= n + m; ++i)</pre>
      for (int j = 1; j <= max(n, m); ++j)</pre>
        if (col[i][j])
          G[i][col[i][j]] = G[col[i][j]][i] = j;
  } // u = left index, v = right index
void add_edge(int u, int v) {
    int c1 = find_col(u), c2 = find_col(v + n);
    dfs(u, c2, c1);
    col[u][c2] = v + n, col[v + n][c2] = u;
  BipartiteEdgeColoring (int _n, int _m) : n(_n),
    m(_m), col(n + m + 1, vector < int > (max(n, m) + 1)),
    G(n + m + 1, vector < int > (n + m + 1)) {}
```

# 4.11 Edge Coloring [60e200]

```
struct Vizing { // 1-based
  // returns edge coloring in adjacent matrix G
  int n;
  vector <vector <int>> C, G;
  vector <int> X, vst;
  vector <pii> E;
  void solve() {
    auto update = [&](int u)
    { for (X[u] = 1; C[u][X[u]]; ++X[u]); };
    auto color = [&](int u, int v, int c) {
      int p = G[u][v];
      G[u][v] = G[v][u] = c;
      C[u][c] = v, C[v][c] = u;
      C[u][p] = C[v][p] = 0;
      if (p) X[u] = X[v] = p;
      else update(u), update(v);
```

```
return p:
  auto flip = [&](int u, int c1, int c2) {
    int p = C[u][c1];
    swap(C[u][c1], C[u][c2]);
    if (p) G[u][p] = G[p][u] = c2;
    if (!C[u][c1]) X[u] = c1;
    if (!C[u][c2]) X[u] = c2;
    return p;
  fill(1 + all(X), 1);
  for (int t = 0; t < (int)E.size(); ++t) {</pre>
    auto [u, v0] = E[t];
    int v = v0, c0 = X[u], c = c0, d;
    vector<pii> L;
    fill(1 + all(vst), 0);
    while (!G[u][v0]) {
      L.emplace_back(v, d = X[v]);
      if (!C[v][c]) {
        for (int a = sz(L) - 1; a >= 0; --a)
          c = color(u, L[a].first, c);
      } else if (!C[u][d]) {
        for (int a = sz(L) - 1; a >= 0; --a)
          color(u, L[a].first, L[a].second);
      } else if (vst[d]) break;
      else vst[d] = 1, v = C[u][d];
    if (!G[u][v0]) {
      for (; v; v = flip(v, c, d), swap(c, d));
      if (int a; C[u][c0]) {
        for (a = sz(L) - 2;
          a >= 0 && L[a].second != c; --a);
        for (; a >= 0; --a)
          color(u, L[a].first, L[a].second);
      else --t;
 }
void add_edge(int u, int v) { E.emplace_back(u, v); }
Vizing(int _n) : n(_n), C(n + 1, vector < int > (n + 1)),
G(n + 1, vector < int > (n + 1)), X(n + 1), vst(n + 1) {}
```

#### 4.12 Maximum Clique [f99a13]

```
struct MaxClique { // Maximum Clique
 bitset<N> a[N], cs[N];
  int ans, sol[N], q, cur[N], d[N], n;
  void init(int _n) {
    n = _n;
    for (int i = 0; i < n; i++) a[i].reset();</pre>
  void add_edge(int u, int v) { a[u][v] = a[v][u] = 1;
      }
  void csort(vector<int> &r, vector<int> &c) {
    int mx = 1, km = max(ans - q + 1, 1), t = 0;
    int m = r.size();
    cs[1].reset(), cs[2].reset();
    for (int i = 0; i < m; i++) {</pre>
      int p = r[i], k = 1;
      while ((cs[k] & a[p]).count()) k++;
      if (k > mx) mx++, cs[mx + 1].reset();
      cs[k][p] = 1;
      if (k < km) r[t++] = p;
    c.resize(m);
   if (t) c[t - 1] = 0;
    for (int k = km; k <= mx; k++)</pre>
      for (int p = cs[k]._Find_first(); p < N;</pre>
              p = cs[k]._Find_next(p))
        r[t] = p, c[t] = k, t++;
  void dfs(vector<int> &r, vector<int> &c, int 1,
    bitset<N> mask) {
   while (!r.empty()) {
      int p = r.back();
      r.pop_back(), mask[p] = 0;
      if (q + c.back() <= ans) return;</pre>
      cur[q++] = p;
      vector<int> nr, nc;
      bitset<N> nmask = mask & a[p];
```

```
for (int i : r)
        if (a[p][i]) nr.push_back(i);
      if (!nr.empty()) {
        if (1 < 4) {
           for (int i : nr)
            d[i] = (a[i] \& nmask).count();
          sort(nr.begin(), nr.end(),
            [&](int x, int y) { return d[x] > d[y]; });
        csort(nr, nc), dfs(nr, nc, l + 1, nmask);
      } else if (q > ans) ans = q, copy_n(cur, q, sol);
      c.pop_back(), q--;
  int solve(bitset<N> mask = bitset<N>(
               string(N, '1'))) { // vertex mask
    vector<int> r, c;
    ans = q = 0;
    for (int i = 0; i < n; i++)</pre>
      if (mask[i]) r.push_back(i);
    for (int i = 0; i < n; i++)</pre>
      d[i] = (a[i] & mask).count();
    sort(r.begin(), r.end(),
      [&](int i, int j) { return d[i] > d[j]; });
    csort(r, c), dfs(r, c, 1, mask);
    return ans; // sol[0 ~ ans-1]
};
```

# 5 String

#### 5.1 Aho-Corasick Automaton [99a712]

```
struct AC {
  int ch[N][26], to[N][26], fail[N], sz;
  vector <int> g[N];
  int cnt[N];
  AC () {sz = 0, extend();}
  void extend() {fill(ch[sz], ch[sz] + 26, 0), sz++;}
  int nxt(int u, int v) {
    if (!ch[u][v]) ch[u][v] = sz, extend();
     return ch[u][v];
  int insert(string s) {
     int now = 0;
     for (char c : s) now = nxt(now, c - 'a');
     cnt[now]++;
     return now:
  void build_fail() {
     queue <int> q;
     for (int i = 0; i < 26; ++i) if (ch[0][i]) {</pre>
       q.push(ch[0][i]);
       g[0].push_back(ch[0][i]);
       to[0][i] = ch[0][i];
     while (!q.empty()) {
       int v = q.front(); q.pop();
       for (int j = 0; j < 26; ++j) {</pre>
         to[v][j] = ch[v][j] ? ch[v][j] : to[fail[v]][j
       for (int i = 0; i < 26; ++i) if (ch[v][i]) {</pre>
         int u = ch[v][i], k = fail[v];
         while (k && !ch[k][i]) k = fail[k];
         if (ch[k][i]) k = ch[k][i];
         fail[u] = k;
         cnt[u] += cnt[k], g[k].push_back(u);
         q.push(u);
      }
    }
  int match(string &s) {
    int now = 0, ans = 0;
     for (char c : s) {
      now = to[now][c - 'a'];
       ans += cnt[now];
     return ans;
};
```

#### 5.2 KMP Algorithm [f379fc]

```
vector <int> build_fail(string s) {
 vector <int> f(s.size() + 1, 0);
  int k = 0;
  for (int i = 1; i < (int)s.size(); ++i) {</pre>
   while (k \&\& s[k] != s[i]) k = f[k];
    if (s[k] == s[i]) k++;
   f[i + 1] = k;
  return f;
int match(string s, string t) {
 vector <int> f = build_fail(t);
  int k = 0, ans = 0;
  for (int i = 0; i < (int)s.size(); ++i) {</pre>
   while (k && s[i] != t[k]) k = f[k];
    if (s[i] == t[k]) k++;
   if (k == (int)t.size()) ans++, k = f[k];
 }
  return ans;
}
```

# 5.3 Z Algorithm [7d5c7c]

```
vector <int> buildZ(string s) {
  int n = (int)s.size(), l = 0, r = 0;
  vector <int> Z(n);
  for (int i = 0; i < n; ++i) {
    Z[i] = max(min(Z[i - 1], r - i), 0);
    while (i + Z[i] < n && s[Z[i]] == s[i + Z[i]]) {
        l = i, r = i + Z[i], Z[i]++;
    }
  }
  return Z;
}</pre>
```

### 5.4 Manacher [c18d8b]

```
// return value only consider string tmp, not s
vector <int> manacher(string tmp) {
   string s = "&";
   for (char c : tmp) s.pb(c), s.pb('%');
   int l = 0, r = 0, n = (int)s.size();
   vector <int> Z(n);
   for (int i = 0; i < n; ++i) {
        Z[i] = r > i ? min(Z[2 * l - i], r - i) : 1;
        while (s[i + Z[i]] == s[i - Z[i]]) Z[i]++;
        if (Z[i] + i > r) l = i, r = Z[i] + i;
   }
   for (int i = 0; i < n; ++i) {
        Z[i] = (Z[i] - (i & 1)) / 2 * 2 + (i & 1);
   }
   return Z;
}</pre>
```

#### 5.5 Suffix Array [ba4998]

```
int sa[N], tmp[2][N], c[N], rk[N], lcp[N];
void buildSA(string s) {
 int *x = tmp[0], *y = tmp[1], m = 256, n = s.size();
  for (int i = 0; i < m; ++i) c[i] = 0;</pre>
 for (int i = 0; i < n; ++i) c[x[i] = s[i]]++;
for (int i = 1; i < m; ++i) c[i] += c[i - 1];</pre>
  for (int i = n - 1; ~i; --i) sa[--c[x[i]]] = i;
  for (int k = 1; k < n; k <<= 1) {</pre>
    for (int i = 0; i < m; ++i) c[i] = 0;
    for (int i = 0; i < n; ++i) c[x[i]]++;</pre>
    for (int i = 1; i < m; ++i) c[i] += c[i - 1];
    int p = 0;
    for (int i = n - k; i < n; ++i) y[p++] = i;</pre>
    for (int i = 0; i < n; ++i) if (sa[i] >= k)
      y[p++] = sa[i] - k;
    for (int i = n - 1; ~i; --i)
      sa[--c[x[y[i]]]] = y[i];
    y[sa[0]] = p = 0;
    for (int i = 1; i < n; ++i) {
      int a = sa[i], b = sa[i - 1];
      if (!(x[a] == x[b] \&\& a + k < n \&\& b + k < n \&\& x
           [a + k] == x[b + k])) p++;
      y[sa[i]] = p;
    if (n == p + 1) break;
```

```
swap(x, y), m = p + 1;
 }
}
void buildLCP(string s) {
 // lcp[i] = LCP(sa[i - 1], sa[i])
  // lcp(i, j) = query_lcp_min[rk[i] + 1, rk[j] + 1)
  int n = s.length(), val = 0;
  for (int i = 0; i < n; ++i) rk[sa[i]] = i;</pre>
  for (int i = 0; i < n; ++i) {
    if (!rk[i]) lcp[rk[i]] = 0;
    else {
      if (val) val--;
      int p = sa[rk[i] - 1];
      while (val + i < n && val + p < n && s[val + i]
          == s[val + p]) val++;
      lcp[rk[i]] = val;
    }
  }
```

#### 5.6 SAIS [fbc167]

```
int sa[N << 1], rk[N], lcp[N];</pre>
// string ASCII value need > 0
namespace sfx {
bool _t[N << 1];</pre>
int _s[N << 1], _c[N << 1], x[N], _p[N], _q[N << 1];
void pre(int *sa, int *c, int n, int z) {</pre>
  fill_n(sa, n, 0), copy_n(c, z, x);
void induce(int *sa, int *c, int *s, bool *t, int n,
    int z) {
  copy_n(c, z - 1, x + 1);
for (int i = 0; i < n; ++i)
    if (sa[i] && !t[sa[i] - 1])
      sa[x[s[sa[i] - 1]]++] = sa[i] - 1;
  copy_n(c, z, x);
  for (int i = n - 1; i >= 0; --i)
    if (sa[i] && t[sa[i] - 1])
      sa[--x[s[sa[i] - 1]]] = sa[i] - 1;
void sais(int *s, int *sa, int *p, int *q, bool *t, int
      *c, int n, int z) {
  bool uniq = t[n - 1] = true;
  int nn = 0, nmxz = -1, *nsa = sa + n, *ns = s + n,
      last = -1;
  fill_n(c, z, 0);
  for (int i = 0; i < n; ++i) uniq &= ++c[s[i]] < 2;</pre>
  partial_sum(c, c + z, c);
  if (uniq) {
    for (int i = 0; i < n; ++i) sa[--c[s[i]]] = i;</pre>
    return;
  for (int i = n - 2; i >= 0; --i)
    if (s[i] == s[i + 1]) t[i] = t[i + 1];
    else t[i] = s[i] < s[i + 1];</pre>
  pre(sa, c, n, z);
  for (int i = 1; i <= n - 1; ++i)</pre>
    if (t[i] && !t[i - 1])
      sa[--x[s[i]]] = p[q[i] = nn++] = i;
  induce(sa, c, s, t, n, z);
for (int i = 0; i < n; ++i)
    if (sa[i] && t[sa[i]] && !t[sa[i] - 1]) {
       bool neq = last < 0 || !equal(s + sa[i], s + p[q[</pre>
           sa[i]] + 1], s + last);
      ns[q[last = sa[i]]] = nmxz += neq;
  sais(ns, nsa, p + nn, q + n, t + n, c + z, nn, nmxz +
       1);
  pre(sa, c, n, z);
  for (int i = nn - 1; i >= 0; --i)
    sa[--x[s[p[nsa[i]]]]] = p[nsa[i]];
  induce(sa, c, s, t, n, z);
void buildSA(string s) {
  int n = s.length();
  for (int i = 0; i < n; ++i) _s[i] = s[i];</pre>
   s[n] = 0;
  sais(_s, sa, _p, _q, _t, _c, n + 1, 256);
  for (int i = 1; i <= n; ++i) sa[i - 1] = sa[i];</pre>
} // buildLCP()...
```

#### 5.7 Suffix Automaton [24c521]

```
struct SAM {
  int ch[N][26], len[N], link[N], pos[N], cnt[N], sz;
  // node -> strings with the same endpos set
  // length in range [len(link) + 1, len]
  // node's endpos set -> pos in the subtree of node
  // link -> longest suffix with different endpos set
  // len -> longest suffix
// pos -> end position
  // cnt -> size of endpos set
  SAM () \{len[0] = 0, link[0] = -1, pos[0] = 0, cnt[0]
       = 0, sz = 1;
  void build(string s) {
    int last = 0;
    for (int i = 0; i < s.length(); ++i) {</pre>
       char c = s[i];
      int cur = sz++;
       len[cur] = len[last] + 1, pos[cur] = i + 1;
      int p = last;
      while (~p && !ch[p][c - 'a'])
  ch[p][c - 'a'] = cur, p = link[p];
       if (p == -1) link[cur] = 0;
       else {
         int q = ch[p][c - 'a'];
         if (len[p] + 1 == len[q]) {
           link[cur] = q;
         } else {
           int nxt = sz++;
           len[nxt] = len[p] + 1, link[nxt] = link[q];
           pos[nxt] = 0;
           for (int j = 0; j < 26; ++j)
           ch[nxt][j] = ch[q][j];

while (~p && ch[p][c - 'a'] == q)

ch[p][c - 'a'] = nxt, p = link[p];
           link[q] = link[cur] = nxt;
        }
      cnt[cur]++;
      last = cur;
    vector <int> p(sz);
    iota(all(p), 0);
    sort(all(p),
      [&](int i, int j) {return len[i] > len[j];});
     for (int i = 0; i < sz; ++i)</pre>
       cnt[link[p[i]]] += cnt[p[i]];
  }
} sam;
```

# 5.8 Minimum Rotation [aa3a61]

```
string rotate(const string &s) {
  int n = (int)s.size(), i = 0, j = 1;
  string t = s + s;
  while (i < n && j < n) {
    int k = 0;
    while (k < n && t[i + k] == t[j + k]) ++k;
    if (t[i + k] <= t[j + k]) j += k + 1;
    else i += k + 1;
    if (i == j) ++j;
  }
  int pos = (i < n ? i : j);
  return t.substr(pos, n);
}</pre>
```

#### 5.9 Palindrome Tree [0518a5]

```
struct PAM {
   int ch[N][26], cnt[N], fail[N], len[N], sz;
   string s;
   // 0 -> even root, 1 -> odd root
   PAM () {}
   void init(string s) {
      sz = 0, extend(), extend();
      len[0] = 0, fail[0] = 1, len[1] = -1;
      int lst = 1;
      for (int i = 0; i < s.length(); ++i) {
       while (s[i - len[lst] - 1] != s[i])
        lst = fail[lst];
      if (!ch[lst][s[i] - 'a']) {
        int idx = extend();
        len[idx] = len[lst] + 2;</pre>
```

```
int now = fail[lst];
        while (s[i - len[now] - 1] != s[i])
          now = fail[now];
        fail[idx] = ch[now][s[i] - 'a'];
        ch[lst][s[i] - 'a'] = idx;
      lst = ch[lst][s[i] - 'a'], cnt[lst]++;
   }
  }
  void build_count() {
    for (int i = sz - 1; i > 1; --i)
      cnt[fail[i]] += cnt[i];
  int extend() {
    fill(ch[sz], ch[sz] + 26, 0), sz++;
    return sz - 1;
  }
};
```

# 5.10 Lyndon Factorization [a9eeb0]

```
// partition s = w[0] + w[1] + ... + w[k-1],
// w[0] >= w[1] >= ... >= w[k-1]
// each w[i] strictly smaller than all its suffix
vector <string> duval(const string &s) {
  vector <string> ans;
  for (int n = (int)s.size(), i = 0, j, k; i < n; ) {
    for (j = i + 1, k = i; j < n && s[k] <= s[j]; j++)
        k = (s[k] < s[j] ? i : k + 1);
    for (; i <= k; i += j - k)
        ans.pb(s.substr(i, j - k)); // s.substr(l, len)
  }
  return ans;
}</pre>
```

#### 5.11 Main Lorentz [f3da14]

```
int to_left[N], to_right[N];
vector <array <int, 3>> rep; // L, r, Len.
// substr([l, r], Len * 2) are tandem
void findRep(string &s, int 1, int r) {
   if (r - l == 1) return;
   int m = 1 + r >> 1;
   findRep(s, 1, m), findRep(s, m, r);
   string sl = s.substr(1, m - 1);
   string sr = s.substr(m, r - m);
   vector <int> Z = buildZ(sr + "#" + sl);
   for (int i = 1; i < m; ++i)</pre>
     to_right[i] = Z[r - m + 1 + i - 1];
   reverse(all(sl));
   Z = buildZ(s1);
   for (int i = 1; i < m; ++i)</pre>
     to_left[i] = Z[m - i - 1];
   reverse(all(sl));
   for (int i = 1; i + 1 < m; ++i) {</pre>
     int k1 = to_left[i], k2 = to_right[i + 1];
     int len = m - i - 1;
     if (k1 < 1 || k2 < 1 || len < 2) continue;</pre>
     int t1 = max(1, len - k2), tr = min(len - 1, k1);
     if (tl <= tr) rep.pb({i + 1 - tr, i + 1 - tl,len});</pre>
   Z = buildZ(sr);
for (int i = m; i < r; ++i) to_right[i] = Z[i - m];</pre>
   reverse(all(s1)), reverse(all(sr));
   Z = buildZ(sl + "#" + sr);
   for (int i = m; i < r; ++i)</pre>
     to_left[i] = Z[m - l + 1 + r - i - 1];
   reverse(all(s1)), reverse(all(sr));
   for (int i = m; i + 1 < r; ++i) {</pre>
     int k1 = to_left[i], k2 = to_right[i + 1];
     int len = i - m + 1;
     if (k1 < 1 || k2 < 1 || len < 2) continue;</pre>
     int tl = max(len - k2, 1), tr = min(len - 1, k1);
     if (tl <= tr)
       rep.pb(\{i + 1 - len - tr, i + 1 - len - tl, len\});
   Z = buildZ(sr + "#" + sl);
  for (int i = 1; i < m; ++i)
  if (Z[r - m + 1 + i - 1] >= m - i)
       rep.pb({i, i, m - i});
}
```

#### 6 Math

#### 6.1 Miller Rabin / Pollard Rho [6c9c33]

```
11 mul(11 x, 11 y, 11 p) {return (x * y - (11)((long
double)x / p * y) * p + p) % p;} // __int128
vector<11> chk = {2, 325, 9375, 28178, 450775, 9780504,
     1795265022}:
ll Pow(ll a, ll b, ll n) {
  11 \text{ res} = 1;
  for (; b; b >>= 1, a = mul(a, a, n))
    if (b & 1) res = mul(res, a, n);
  return res;
bool check(ll a, ll d, int s, ll n) {
  a = Pow(a, d, n);
  if (a <= 1) return 1;</pre>
  for (int i = 0; i < s; ++i, a = mul(a, a, n)) {</pre>
   if (a == 1) return 0;
    if (a == n - 1) return 1;
  }
  return 0;
bool IsPrime(ll n) {
  if (n < 2) return 0;
  if (n % 2 == 0) return n == 2;
  ll d = n - 1, s = 0;
  while (d % 2 == 0) d >>= 1, ++s;
  for (ll i : chk) if (!check(i, d, s, n)) return 0;
const vector<ll> small = {2, 3, 5, 7, 11, 13, 17, 19};
11 FindFactor(ll n) {
  if (IsPrime(n)) return 1;
  for (ll p : small) if (n % p == 0) return p;
  11 x, y = 2, d, t = 1;
  auto f = [&](ll a) {return (mul(a, a, n) + t) % n;};
  for (int 1 = 2; ; 1 <<= 1) {
    x = y;
    int m = min(1, 32);
    for (int i = 0; i < 1; i += m) {</pre>
      d = 1;
      for (int j = 0; j < m; ++j) {
        y = f(y), d = mul(d, abs(x - y), n);
      11 g = \_gcd(d, n);
      if (g == n) {
       1 = 1, y = 2, ++t;
        break;
      if (g != 1) return g;
    }
 }
map <11, int> res;
void PollardRho(ll n) {
  if (n == 1) return;
  if (IsPrime(n)) return ++res[n], void(0);
  11 d = FindFactor(n);
  PollardRho(n / d), PollardRho(d);
6.2 Ext GCD [a4b22d]
//a * p.first + b * p.second = gcd(a, b)
pair<ll, ll> extgcd(ll a, ll b) {
```

```
//a * p.first + b * p.second = gcd(a, b)
pair<ll, ll> extgcd(ll a, ll b) {
   if (b == 0) return {1, 0};
   auto [y, x] = extgcd(b, a % b);
   return pair<ll, ll>(x, y - (a / b) * x);
}
```

#### 6.3 Chinese Remainder Theorem [90d2ce]

```
pair<11, 11> CRT(11 x1, 11 m1, 11 x2, 11 m2) {
    11 g = gcd(m1, m2);
    if ((x2 - x1) % g) return make_pair(-1, -1);// no sol
    m1 /= g, m2 /= g;
    pair <11, 11> p = extgcd(m1, m2);
    11 lcm = m1 * m2 * g;
    11 res = p.first * (x2 - x1) * m1 + x1;
    // be careful with overflow
    return make_pair((res % lcm + lcm) % lcm, lcm);
}
```

#### 6.4 PiCount [1db46f]

```
const int V = 10000000, N = 100, M = 100000;
vector<int> primes;
bool isp[V];
int small_pi[V], dp[N][M];
void sieve(int x){
  for(int i = 2; i < x; ++i) isp[i] = true;</pre>
  isp[0] = isp[1] = false;
  for(int i = 2; i * i < x; ++i) if(isp[i])</pre>
  for(int j = i * i; j < x; j += i) isp[j] = false;
for(int i = 2; i < x; ++i) if(isp[i]) primes.pb(i);</pre>
void init(){
  sieve(V)
  small_pi[0] = 0;
  for(int i = 1; i < V; ++i)</pre>
     small_pi[i] = small_pi[i - 1] + isp[i];
  for(int i = 0; i < M; ++i) dp[0][i] = i;</pre>
  for(int i = 1; i < N; ++i) for(int j = 0; j < M; ++j)
     dp[i][j] = dp[i - 1][j] - dp[i - 1][j / primes[i -
11 phi(ll n, int a){
  if(!a) return n;
  if(n < M && a < N) return dp[a][n];</pre>
  if(primes[a - 1] > n) return 1;
  if(111 * primes[a - 1] * primes[a - 1] >= n && n < V)
     return small_pi[n] - a + 1;
   return phi(n, a - 1) - phi(n / primes[a - 1], a - 1);
11 PiCount(11 n){
  if(n < V) return small_pi[n];</pre>
  int s = sqrt(n + 0.5), y = cbrt(n + 0.5), a =
       small_pi[y];
  ll res = phi(n, a) + a - 1;
  for(; primes[a] <= s; ++a) res -= max(PiCount(n /</pre>
       primes[a]) - PiCount(primes[a]) + 1, 011);
  return res:
}
```

# 6.5 Linear Function Mod Min [5552e3]

```
11 topos(ll x, ll m)
{ x \%= m; if (x < 0) x += m; return x; }
 //min value of ax + b \pmod{m} for x \in [0, n - 1]. O(
     Log m)
ll min_rem(ll n, ll m, ll a, ll b) {
   a = topos(a, m), b = topos(b, m);
   for (ll g = \_gcd(a, m); g > 1;) return g * min_rem(n)
       , m / g, a / g, b / g) + (b % g);
   for (11 nn, nm, na, nb; a; n = nn, m = nm, a = na, b
       = nb) {
     if (a <= m - a) {
       nn = (a * (n - 1) + b) / m;
       if (!nn) break;
       nn += (b < a);
       nm = a, na = topos(-m, a);
       nb = b < a ? b : topos(b - m, a);
     } else {
       ll lst = b - (n - 1) * (m - a);
       if (lst >= 0) {b = lst; break;}
       nn = -(1st / m) + (1st % m < -a) + 1;
       nm = m - a, na = m % (m - a), nb = b % (m - a);
    }
  }
  return b;
}
//min value of ax + b \pmod{m} for x \in [0, n - 1],
     also return min \times to get the value. O(log m)
//{value, x}
pair<ll, ll> min_rem_pos(ll n, ll m, ll a, ll b) {
   a = topos(a, m), b = topos(b, m);
   11 mn = min_rem(n, m, a, b), g = __gcd(a, m);
   //ax = (mn - b) (mod m)
   11 x = (extgcd(a, m).first + m) * ((mn - b + m) / g)
       % (m / g);
   return {mn, x};
}
```

# **6.6 Determinant** [fb958e]

```
int Det(vector <vector <int>> a) {
```

```
int n = (int)a.size(), det = 1;
for (int i = 0; i < n; ++i) {</pre>
  if (!a[i][i]) {
    det = sub(0, det);
    for (int j = i + 1; j < n; ++j) if (a[j][i]) {
      swap(a[j], a[i]);
      break:
    if (!a[i][i]) return 0;
  det = mul(det, a[i][i]);
  int tmp = Pow(a[i][i], mod - 2);
  for (int j = 0; j < n; ++j)</pre>
    a[i][j] = mul(a[i][j], tmp);
  for (int j = 0; j < n; ++j) if (i ^ j) {</pre>
    tmp = a[j][i];
    for (int k = 0; k < n; ++k) {
      a[j][k] = sub(a[j][k], mul(a[i][k], tmp));
 }
return det;
```

#### **6.7 Floor Sum** [49de67]

# 6.8 Quadratic Residue [51ec55]

```
int Jacobi(int a, int m) {
 int s = 1;
  for (; m > 1; ) {
    a %= m;
    if (a == 0) return 0;
    const int r = __builtin_ctz(a);
    if ((r \& 1) \&\& ((m + 2) \& 4)) s = -s;
    a >>= r;
    if (a \& m \& 2) s = -s;
    swap(a, m);
  }
  return s;
int QuadraticResidue(int a, int p) {
 if (p == 2) return a & 1;
  const int jc = Jacobi(a, p);
  if (jc == 0) return 0;
  if (jc == -1) return -1;
  int b, d;
  for (;;) {
   b = rand() % p;
    d = (111 * b * b + p - a) % p;
    if (Jacobi(d, p) == -1) break;
  11 f0 = b, f1 = 1, g0 = 1, g1 = 0, tmp;
  for (int e = (p + 1) >> 1; e; e >>= 1) {
    if (e & 1) {
      tmp = (g0 * f0 + d * (g1 * f1 % p)) % p;
      g1 = (g0 * f1 + g1 * f0) % p;
      g0 = tmp;
    tmp = (f0 * f0 + d * (f1 * f1 % p)) % p;
    f1 = (2 * f0 * f1) % p;
    f0 = tmp;
  }
  return g0;
```

# 6.9 Discrete Log [8f7f93]

```
11 DiscreteLog(11 a, 11 b, 11 m) { // a^x = b (mod m)
  const int B = 35000;
```

```
ll k = 1 % m, ans = 0, g;
while ((g = gcd(a, m)) > 1) {
   if (b == k) return ans;
   if (b % g) return -1;
   b /= g, m /= g, ans++, k = (k * a / g) % m;
}
if (b == k) return ans;
unordered_map <1l, int> m1;
ll tot = 1;
for (int i = 0; i < B; ++i)
   m1[tot * b % m] = i, tot = tot * a % m;
ll cur = k * tot % m;
for (int i = 1; i <= B; ++i, cur = cur * tot % m)
   if (m1.count(cur)) return i * B - m1[cur] + ans;
return -1;
}</pre>
```

# **6.10 Factorial without Prime Factor**[c324f3]

```
// O(p^k + Log^2 n), pk = p^k
ll prod[MAXP];
ll fac_no_p(ll n, ll p, ll pk) {
  prod[0] = 1;
  for (int i = 1; i <= pk; ++i)
    if (i % p) prod[i] = prod[i - 1] * i % pk;
    else prod[i] = prod[i - 1];
ll rt = 1;
for (; n; n /= p) {
    rt = rt * mpow(prod[pk], n / pk, pk) % pk;
    rt = rt * prod[n % pk] % pk;
}
return rt;
} // (n! without factor p) % p^k</pre>
```

#### 6.11 Berlekamp Massey [f867ec]

```
// need add, sub, mul
vector <int> BerlekampMassey(vector <int> a) {
  // find min |c| such that a_n = sum c_j * a_{n - j - 1}
       1}, 0-based
  // O(N^2), if |c| = k, |a| >= 2k sure correct auto f = [&](vector<int> v, 11 c) {
    for (int &x : v) x = mul(x, c);
    return v;
  };
  vector <int> c, best;
  int pos = 0, n = (int)a.size();
  for (int i = 0; i < n; ++i) {</pre>
    int error = a[i];
    for (int j = 0; j < (int)c.size(); ++j)</pre>
       error = sub(error, mul(c[j], a[i - 1 - j]));
    if (error == 0) continue;
    int inv = Pow(error, mod - 2);
    if (c.empty()) {
      c.resize(i + 1), pos = i, best.pb(inv);
    } else {
      vector <int> fix = f(best, error);
fix.insert(fix.begin(), i - pos - 1, 0);
       if (fix.size() >= c.size()) {
         best = f(c, sub(0, inv));
         best.insert(best.begin(), inv);
         pos = i, c.resize(fix.size());
       for (int j = 0; j < (int)fix.size(); ++j)</pre>
         c[j] = add(c[j], fix[j]);
    }
  }
  return c;
```

## **6.12** Simplex [b68fb9]

```
struct Simplex { // 0-based
  using T = long double;
  static const int N = 410, M = 30010;
  const T eps = 1e-7;
  int n, m;
  int Left[M], Down[N];
  // Ax <= b, max c^T x
  // result : v, xi = sol[i]
  T a[M][N], b[M], c[N], v, sol[N];
  bool eq(T a, T b) {return fabs(a - b) < eps;}</pre>
```

```
bool ls(T a, T b) {return a < b && !eq(a, b);}
void init(int _n, int _m) {</pre>
    n = _n, m = _m, v = 0;
    for (int i = 0; i < m; ++i)</pre>
      for (int j = 0; j < n; ++j) a[i][j] = 0;</pre>
    for (int i = 0; i < m; ++i) b[i] = 0;</pre>
    for (int i = 0; i < n; ++i) c[i] = sol[i] = 0;</pre>
  void pivot(int x, int y) {
    swap(Left[x], Down[y]);
    T k = a[x][y]; a[x][y] = 1;
    vector <int> nz;
    for (int i = 0; i < n; ++i) {</pre>
      a[x][i] /= k;
      if (!eq(a[x][i], 0)) nz.push_back(i);
    b[x] /= k;
    for (int i = 0; i < m; ++i) {</pre>
      if (i == x || eq(a[i][y], 0)) continue;
      k = a[i][y], a[i][y] = 0;
      b[i] -= k * b[x];
      for (int j : nz) a[i][j] -= k * a[x][j];
    if (eq(c[y], 0)) return;
    k = c[y], c[y] = 0, v += k * b[x];
for (int i : nz) c[i] -= k * a[x][i];
  // 0: found solution, 1: no feasible solution, 2:
       unbounded
  int solve() {
    for (int i = 0; i < n; ++i) Down[i] = i;</pre>
    for (int i = 0; i < m; ++i) Left[i] = n + i;</pre>
    while (true) {
      int x = -1, y = -1;
       for (int i = 0; i < m; ++i) if (ls(b[i], 0) && (x
            == -1 \mid \mid b[i] < b[x])) x = i;
      if (x == -1) break;
      for (int i = 0; i < n; ++i) if (ls(a[x][i], 0) &&</pre>
            (y == -1 \mid | a[x][i] < a[x][y])) y = i;
      if (y == -1) return 1;
      pivot(x, y);
    while (true) {
      int x = -1, y = -1;
      for (int i = 0; i < n; ++i) if (ls(0, c[i]) && (y
            == -1 || c[i] > c[y])) y = i;
      if (y == -1) break;
      for (int i = 0; i < m; ++i)</pre>
        if (x == -1) return 2;
      pivot(x, y);
    for (int i = 0; i < m; ++i) if (Left[i] < n)</pre>
      sol[Left[i]] = b[i];
    return 0;
  }
};
```

# 6.13 Linear Programming Construction

Standard form: maximize  $\mathbf{c}^T\mathbf{x}$  subject to  $A\mathbf{x} \leq \mathbf{b}$  and  $\mathbf{x} \geq 0$ . Dual LP: minimize  $\mathbf{b}^T\mathbf{y}$  subject to  $A^T\mathbf{y} \geq \mathbf{c}$  and  $\mathbf{y} \geq 0$ .  $\bar{\mathbf{x}}$  and  $\bar{\mathbf{y}}$  are optimal if and only if for all  $i \in [1,n]$ , either  $\bar{x}_i = 0$  or  $\sum_{j=1}^m A_{ji}\bar{y}_j = c_i$  holds and for all  $i \in [1,m]$  either  $\bar{y}_i = 0$  or  $\sum_{j=1}^n A_{ij}\bar{x}_j = b_j$  holds.

- 1. In case of minimization, let  $c_i^\prime = -c_i$
- 2.  $\sum_{1 \leq i \leq n} A_{ji} x_i \geq b_j \rightarrow \sum_{1 \leq i \leq n} -A_{ji} x_i \leq -b_j$
- 3.  $\sum_{1 \leq i \leq n} A_{ji} x_i = b_j$ 
  - $\sum_{1 \le i \le n} A_{ji} x_i \le b_j$ •  $\sum_{1 \le i \le n} A_{ji} x_i \ge b_j$
- 4. If  $x_i$  has no lower bound, replace  $x_i$  with  $x_i x_i'$

# 6.14 Euclidean

- $m = \lfloor \frac{an+b}{c} \rfloor$
- Time complexity:  $O(\log n)$

$$\begin{split} g(a,b,c,n) &= \sum_{i=0}^n i \lfloor \frac{ai+b}{c} \rfloor \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor \cdot \frac{n(n+1)}{2} \\ +g(a \mod c, b \mod c, c, n), & a \geq c \vee b \geq c \\ 0, & n < 0 \vee a = 0 \\ \frac{1}{2} \cdot (n(n+1)m - f(c, c-b-1, a, m-1) \\ -h(c, c-b-1, a, m-1)), & \text{otherwise} \end{cases} \end{split}$$

$$\begin{split} h(a,b,c,n) &= \sum_{i=0}^n \lfloor \frac{ai+b}{c} \rfloor^2 \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor^2 \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor^2 \cdot (n+1) \\ + \lfloor \frac{a}{c} \rfloor \cdot \lfloor \frac{b}{c} \rfloor \cdot n(n+1) \\ + h(a \bmod c, b \bmod c, c, n) \\ + 2 \lfloor \frac{a}{c} \rfloor \cdot g(a \bmod c, b \bmod c, c, n) \\ + 2 \lfloor \frac{b}{c} \rfloor \cdot f(a \bmod c, b \bmod c, c, n), & a \geq c \lor b \geq c \\ 0, & n < 0 \lor a = 0 \\ nm(m+1) - 2g(c, c-b-1, a, m-1) \\ -2f(c, c-b-1, a, m-1) - f(a, b, c, n), & \text{otherwise} \end{cases} \end{split}$$

#### 6.15 Theorem

• Kirchhoff's Theorem

Denote L be a  $n\times n$  matrix as the Laplacian matrix of graph G , where  $L_{ii}=d(i)$  ,  $L_{ij}=-c$  where c is the number of edge (i,j) in G .

- The number of undirected spanning in G is  $|\det(\tilde{L}_{11})|$ .
- The number of directed spanning tree rooted at r in G is  $|{\rm det}(\tilde{L}_{rr})|$  .
- Tutte's Matrix

Let D be a  $n \times n$  matrix, where  $d_{ij} = x_{ij}$  ( $x_{ij}$  is chosen uniformly at random) if i < j and  $(i,j) \in E$ , otherwise  $d_{ij} = -d_{ji}$ .  $\frac{rank(D)}{2}$  is the maximum matching on G.

- Cayley's Formula
  - Given a degree sequence  $d_1,d_2,\ldots,d_n$  for each *Labeled* vertices, there are

$$\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$$

spanning trees.

- Let  $T_{n,k}$  be the number of *labeled* forests on n vertices with k components, such that vertex  $1,2,\ldots,k$  belong to different components. Then  $T_{n,k}=kn^{n-k-1}$ .
- Erdős-Gallai Theorem

A sequence of non-negative integers  $d_1 \geq d_2 \geq \ldots \geq d_n$  can be represented as the degree sequence of a finite simple graph on n vertices if and only if  $d_1+d_2+\ldots+d_n$  is even and

$$\sum_{i=1}^k d_i \leq k(k-1) + \sum_{i=k+1}^n \min(d_i,k)$$

holds for all  $1 \leq k \leq n$ .

• Burnside's Lemma

Let X be a set and G be a group that acts on X . For  $g\in G$ , denote by  $X^g$  the elements fixed by g:

$$X^g = \{ x \in X \mid gx \in X \}$$

Then

$$|X/G| = \frac{1}{|G|} \sum_{g \in G} |X^g|.$$

• Gale-Ryser theorem

A pair of sequences of nonnegative integers  $a_1 \geq \cdots \geq a_n$  and  $b_1,\ldots,b_n$  is bigraphic if and only if  $\sum_{i=1}^n a_i = \sum_{i=1}^n b_i$  and  $\sum_{i=1}^k a_i \leq a_i$ 

 $\sum_{i=1}^{n} \min(b_i,k)$  holds for every  $1 \leq k \leq n$ . Sequences a and b called bigraphic if there is a labeled simple bipartite graph such that a and b is the degree sequence of this bipartite graph.

• Fulkerson-Chen-Anstee theorem

A sequence  $(a_1,b_1),\ldots,(a_n,b_n)$  of nonnegative integer pairs with  $a_1 \geq \cdots \geq a_n$  is digraphic if and only if  $\sum_{i=1}^n a_i = \sum_{i=1}^n b_i$  and  $\sum_{i=1}^k a_i \leq \sum_{i=1}^k \min(b_i,k-1) + \sum_{i=k+1}^n \min(b_i,k) \text{ holds for every } 1 \leq k \leq n$ 

Sequences  $\boldsymbol{a}$  and  $\boldsymbol{b}$  called digraphic if there is a labeled simple directed graph such that each vertex  $\boldsymbol{v}_i$  has indegree  $\boldsymbol{a}_i$  and outdegree  $b_i$ .

• Pick's theorem

For simple polygon, when points are all integer, we have  ${\cal A}\,=\,$  $\#\{\text{lattice points in the interior}\} + \frac{\#\{\text{lattice points on the boundary}\}}{2} - 1$ 

• Möbius inversion formula

- 
$$f(n) = \sum_{d|n} g(d) \Leftrightarrow g(n) = \sum_{d|n} \mu(d) f(\frac{n}{d})$$
  
-  $f(n) = \sum_{n|d} g(d) \Leftrightarrow g(n) = \sum_{n|d} \mu(\frac{d}{n}) f(d)$ 

- Spherical cap
  - A portion of a sphere cut off by a plane.

  - r: sphere radius, a: radius of the base of the cap, h: height of the cap,  $\theta$ :  $\arcsin(a/r)$ .
     Volume  $=\pi h^2(3r-h)/3=\pi h(3a^2+h^2)/6=\pi r^3(2+\cos\theta)(1-\sin\theta)$
  - Area  $= 2\pi rh = \pi(a^2 + h^2) = 2\pi r^2(1 \cos\theta)$ .

#### 6.16 Estimation

- The number of divisors of n is at most around  $100\ {\rm for}\ n<5e4$  , 500 for n < 1e7, 2000 for n < 1e10, 2000000 for n < 1e19.
- ullet The number of ways of writing n as a sum of positive integers, disregarding the order of the summands. 1,1,2,3,5,7,11,15,22,30 for  $n=0\sim9$ , 627 for n=20,  $\sim2e5$  for n=50,  $\sim2e8$  for n = 100.
- ullet Total number of partitions of n distinct elements: B(n) = 11, 1, 2, 5, 15, 52, 203, 877, 4140, 21147, 115975, 678570, 4213597,27644437, 190899322, . . . .

#### 6.17 General Purpose Numbers

• Bernoulli numbers

$$B_0 = 1, B_1^{\pm} = \pm \frac{1}{2}, B_2 = \frac{1}{6}, B_3 = 0$$

$$\sum_{j=0}^m \binom{m+1}{j} B_j = 0 \text{, EGF is } B(x) = \frac{x}{e^x-1} = \sum_{n=0}^\infty B_n \frac{x^n}{n!} \,.$$

$$S_m(n) = \sum_{k=1}^n k^m = \frac{1}{m+1} \sum_{k=0}^m {m+1 \choose k} B_k^+ n^{m+1-k}$$

ullet Stirling numbers of the second kind Partitions of n distinct elements into exactly k groups.

$$\begin{split} S(n,k) &= S(n-1,k-1) + kS(n-1,k), S(n,1) = S(n,n) = 1 \\ S(n,k) &= \frac{1}{k!} \sum_{i=0}^k (-1)^{k-i} {k \choose i} i^n \\ x^n &= \sum_{i=0}^n S(n,i)(x)_i \end{split}$$

• Pentagonal number theorem

$$\prod_{n=1}^{\infty} (1 - x^n) = 1 + \sum_{k=1}^{\infty} (-1)^k \left( x^{k(3k+1)/2} + x^{k(3k-1)/2} \right)$$

• Catalan numbers

$$C_n^{(k)} = \frac{1}{(k-1)n+1} \binom{kn}{n}$$

$$C^{(k)}(x) = 1 + x[C^{(k)}(x)]^k$$

• Eulerian numbers

Number of permutations  $\pi \in S_n$  in which exactly k elements are greater than the previous element. k j:s s.t.  $\pi(j) > \pi(j+1)$ ,  $k+1 \ j\text{:s s.t.} \ \pi(j) \geq j\text{,} \ k \ j\text{:s s.t.} \ \pi(j) > j\text{.}$ 

$$E(n,k) = (n-k)E(n-1,k-1) + (k+1)E(n-1,k)$$

$$E(n,0) = E(n, n-1) = 1$$

$$E(n,k) = \sum_{j=0}^{k} (-1)^{j} {n+1 \choose j} (k+1-j)^{n}$$

# Polynomial

# 7.1 Number Theoretic Transform [536cc5]

```
// mul, add, sub, Pow
struct NTT {
  int w[N];
  NTT() {
     int dw = Pow(G, (mod - 1) / N);
     w[0] = 1;
     for (int i = 1; i < N; ++i)</pre>
       w[i] = mul(w[i - 1], dw);
  void operator()(vector<int>& a, bool inv = false) {
     //0 <= a[i] < P
int x = 0, n = a.size();
     for (int j = 1; j < n - 1; ++j) {
  for (int k = n >> 1; (x ^= k) < k; k >>= 1);
       if (j < x) swap(a[x], a[j]);</pre>
     for (int L = 2; L <= n; L <<= 1) {
       int dx = N / L, dl = L >> 1;
for (int i = 0; i < n; i += L) {
         for (int j = i, x = 0; j < i + dl; ++j, x += dx
            int tmp = mul(a[j + dl], w[x]);
            a[j + dl] = sub(a[j], tmp);
           a[j] = add(a[j], tmp);
      }
     if (inv) {
       reverse(a.begin() + 1, a.end());
       int invn = Pow(n, mod - 2);
       for (int i = 0; i < n; ++i)</pre>
         a[i] = mul(a[i], invn);
  }
} ntt;
```

#### Fast Fourier Transform [21cf37]

```
using T = complex <double>;
const double PI = acos(-1);
struct NTT {
  T w[N];
  FFT() {
    T dw = \{\cos(2 * PI / N), \sin(2 * PI / N)\};
    w[0] = 1;
    for (int i = 1; i < N; ++i) w[i] = w[i - 1] * dw;</pre>
  void operator()(vector<T>& a, bool inv = false) {
    // see NTT, replace ll with T
    if (inv) {
      reverse(a.begin() + 1, a.end());
      T invn = 1.0 / n;
      for (int i = 0; i < n; ++i) a[i] = a[i] * invn;</pre>
  }
} ntt;
// after mul, round i.real()
```

#### 7.3 Primes

```
Prime
                 Root
                        Prime
                                                Root
7681
                        167772161
                 17
12289
                 11
                        104857601
40961
                        985661441
65537
                        998244353
                                                3
                        1107296257
786433
                 10
                                                10
5767169
                        2013265921
                                                31
7340033
                        2810183681
23068673
                        2885681153
469762049
                        605028353
2061584302081
                        1945555039024054273
2748779069441
                        9223372036737335297
```

#### Polynomial Operations [9be4e4]

```
typedef vector<int> Poly;
Poly Mul(Poly a, Poly b, int bound = N) {
  int m = a.size() + b.size() - 1, n = 1;
  while (n < m) n <<= 1;</pre>
  a.resize(n), b.resize(n);
  ntt(a), ntt(b);
  Poly out(n);
```

```
for (int i = 0; i < n; ++i) out[i] = mul(a[i], b[i]);</pre>
                                                                     res[i + 1] = mul(a[i], Pow(i + 1, mod - 2));
  ntt(out, true), out.resize(min(m, bound));
                                                                   return res;
  return out:
                                                                 Poly Ln(Poly a) {
                                                                   // O(NlogN), a[0] = 1
Poly Inverse(Poly a) {
  // O(NlogN), a[0] != 0
                                                                   int n = a.size();
                                                                   if (n == 1) return {0};
  int n = a.size();
  Poly res(1, Pow(a[0], mod - 2));
                                                                   Poly d = Derivative(a);
  for (int m = 1; m < n; m <<= 1) {</pre>
                                                                   a.pop_back();
    if (n < m * 2) a.resize(m * 2);</pre>
                                                                   return Integral(Mul(d, Inverse(a), n - 1));
    Poly v1(a.begin(), a.begin() + m * 2), v2 = res;
v1.resize(m * 4), v2.resize(m * 4);
                                                                 Poly Exp(Poly a) {
    ntt(v1), ntt(v2);
                                                                   // O(NlogN), a[0] = 0
    for (int i = 0; i < m * 4; ++i)</pre>
                                                                   int n = a.size();
      v1[i] = mul(mul(v1[i], v2[i]), v2[i]);
                                                                   Poly q(1, 1);
    ntt(v1, true);
                                                                   a[0] = add(a[0], 1);
    res.resize(m * 2);
                                                                   for (int m = 1; m < n; m <<= 1) {
  if (n < m * 2) a.resize(m * 2);</pre>
    for (int i = 0; i < m; ++i)</pre>
                                                                     Poly g(a.begin(), a.begin() + m * 2), h(all(q));
      res[i] = add(res[i], res[i]);
    for (int i = 0; i < m * 2; ++i)
                                                                     h.resize(m * 2), h = Ln(h);
                                                                     for (int i = 0; i < m * 2; ++i)</pre>
      res[i] = sub(res[i], v1[i]);
  }
                                                                       g[i] = sub(g[i], h[i]);
  res.resize(n);
                                                                     q = Mul(g, q, m * 2);
  return res;
                                                                   q.resize(n);
pair <Poly, Poly> Divide(Poly a, Poly b) {
                                                                   return q;
  // a = bQ + R, O(NlogN), b.back() != 0
  int n = a.size(), m = b.size(), k = n - m + 1;
                                                                 Poly PolyPow(Poly a, 11 k) {
  if (n < m) return {{0}, a};</pre>
                                                                   int n = a.size(), m = 0;
  Poly ra = a, rb = b;
                                                                   Poly ans(n, 0);
                                                                   while (m < n && a[m] == 0) m++;</pre>
 reverse(all(ra)), ra.resize(k);
                                                                   if (k \&\& m \&\& (k >= n || k * m >= n)) return ans;
  reverse(all(rb)), rb.resize(k);
 Poly Q = Mul(ra, Inverse(rb), k);
                                                                   if (m == n) return ans[0] = 1, ans;
                                                                   int lead = m * k;
  reverse(all(Q));
                                                                   Poly b(a.begin() + m, a.end());
  Poly res = Mul(b, Q), R(m - 1);
  for (int i = 0; i < m - 1; ++i)</pre>
                                                                   int base = Pow(b[0], k), inv = Pow(b[0], mod - 2);
    R[i] = sub(a[i], res[i]);
                                                                   for (int i = 0; i < n - m; ++i)</pre>
  return {Q, R};
                                                                     b[i] = mul(b[i], inv);
                                                                   b = Ln(b);
Poly SqrtImpl(Poly a) {
                                                                   for (int i = 0; i < n - m; ++i)</pre>
 if (a.empty()) return {0};
                                                                     b[i] = mul(b[i], k % mod);
  int z = QuadraticResidue(a[0], mod), n = a.size();
                                                                   b = Exp(b);
  if (z == -1) return {-1};
                                                                   for (int i = lead; i < n; ++i)</pre>
                                                                     ans[i] = mul(b[i - lead], base);
  Poly q(1, z);
  const int inv2 = (mod + 1) / 2;
                                                                   return ans;
  for (int m = 1; m < n; m <<= 1) {</pre>
    if (n < m * 2) a.resize(m * 2);</pre>
                                                                 vector <int> Evaluate(Poly a, vector <int> x) {
    q.resize(m * 2);
                                                                   if (x.empty()) return {};
    Poly f2 = Mul(q, q, m * 2);
for (int i = 0; i < m * 2; ++i)
                                                                   int n = x.size();
                                                                   vector <Poly> up(n * 2);
      f2[i] = sub(f2[i], a[i]);
                                                                   for (int i = 0; i < n; ++i)</pre>
    f2 = Mul(f2, Inverse(q), m * 2);
for (int i = 0; i < m * 2; ++i)
                                                                     up[i + n] = {sub(0, x[i]), 1};
                                                                   for (int i = n - 1; i > 0; --i)
  up[i] = Mul(up[i * 2], up[i * 2 + 1]);
      q[i] = sub(q[i], mul(f2[i], inv2));
                                                                   vector <Poly> down(n * 2);
  q.resize(n);
                                                                   down[1] = Divide(a, up[1]).second;
                                                                   for (int i = 2; i < n * 2; ++i)</pre>
  return q;
                                                                     down[i] = Divide(down[i >> 1], up[i]).second;
Poly Sqrt(Poly a) {
                                                                   Poly y(n);
 // O(NlogN), return {-1} if not exists
                                                                   for (int i = 0; i < n; ++i) y[i] = down[i + n][0];
  int n = a.size(), m = 0;
                                                                   return y;
  while (m < n && a[m] == 0) m++;</pre>
  if (m == n) return Poly(n);
                                                                 Poly Interpolate(vector <int> x, vector <int> y) {
  if (m & 1) return {-1};
                                                                   int n = x.size();
                                                                   vector <Poly> up(n * 2);
  Poly s = SqrtImpl(Poly(a.begin() + m, a.end()));
  if (s[0] == -1) return {-1};
                                                                   for (int i = 0; i < n; ++i)</pre>
  Poly res(n);
                                                                     up[i + n] = {sub(0, x[i]), 1};
                                                                   for (int i = n - 1; i > 0; --i)
  up[i] = Mul(up[i * 2], up[i * 2 + 1]);
  for (int i = 0; i < s.size(); ++i)</pre>
    res[i + m / 2] = s[i];
                                                                   Poly a = Evaluate(Derivative(up[1]), x);
  return res;
                                                                   for (int i = 0; i < n; ++i)
  a[i] = mul(y[i], Pow(a[i], mod - 2));</pre>
Poly Derivative(Poly a) {
 int n = a.size();
                                                                   vector <Poly> down(n * 2);
  Poly res(n - 1);
                                                                   for (int i = 0; i < n; ++i) down[i + n] = {a[i]};</pre>
                                                                   for (int i = n - 1; i > 0; --i) {
  Poly lhs = Mul(down[i * 2], up[i * 2 + 1]);
  for (int i = 0; i < n - 1; ++i)</pre>
   res[i] = mul(a[i + 1], i + 1);
                                                                     Poly rhs = Mul(down[i * 2 + 1], up[i * 2]);
  return res;
                                                                     down[i].resize(lhs.size());
Poly Integral(Poly a) {
                                                                     for (int j = 0; j < lhs.size(); ++j)</pre>
  int n = a.size();
                                                                        down[i][j] = add(lhs[j], rhs[j]);
  Poly res(n + 1);
 for (int i = 0; i < n; ++i)</pre>
                                                                   return down[1];
```

```
Poly TaylorShift(Poly a, int c) {
 // return sum a_i(x + c)^i;
  // fac[i] = i!, facp[i] = inv(i!)
  int n = a.size();
  for (int i = 0; i < n; ++i) a[i] = mul(a[i], fac[i]);</pre>
  reverse(all(a));
  Poly b(n);
  int w = 1:
  for (int i = 0; i < n; ++i)</pre>
   b[i] = mul(facp[i], w), w = mul(w, c);
  a = Mul(a, b, n), reverse(all(a));
  for (int i = 0; i < n; ++i) a[i] = mul(a[i],facp[i]);</pre>
  return a:
vector<int> SamplingShift(vector<int> a, int c, int m){
 // given f(0), f(1), ..., f(n-1)
// return f(c), f(c+1), ..., f(c+m-1)
  int n = a.size();
  for (int i = 0; i < n; ++i) a[i] = mul(a[i], facp[i]);</pre>
  Poly b(n);
  for (int i = 0; i < n; ++i) {</pre>
   b[i] = facp[i];
    if (i & 1) b[i] = sub(0, b[i]);
  a = Mul(a, b, n);
  for (int i = 0; i < n; ++i) a[i] = mul(a[i], fac[i]);</pre>
  reverse(all(a));
  int w = 1;
  for (int i = 0; i < n; ++i)</pre>
   b[i] = mul(facp[i], w), w = mul(w, sub(c, i));
  a = Mul(a, b, n);
 reverse(all(a));
 for (int i = 0; i < n; ++i) a[i] = mul(a[i], facp[i]);</pre>
  a.resize(m), b.resize(m);
 for (int i = 0; i < m; ++i) b[i] = facp[i];</pre>
  a = Mul(a, b, m);
  for (int i = 0; i < m; ++i) a[i] = mul(a[i], fac[i]);</pre>
  return a;
```

#### 7.5 Fast Linear Recursion [3f8e4e]

```
int FastLinearRecursion(vector <int> a, vector <int> c,
     11 k) {
  // a_n = sigma c_j * a_{n - j - 1}, 0-based
  // O(NlogNlogK), |a| = |c|
  int n = a.size();
  if (k < n) return a[k];</pre>
  vector <int> base(n + 1, 1);
  for (int i = 0; i < n; ++i)</pre>
    base[i] = sub(0, c[n - i - 1]);
  vector <int> poly(n);
  (n == 1 ? poly[0] = c[n - 1] : poly[1] = 1);
  auto calc = [&](vector <int> p1, vector <int> p2) {
    // O(n^2) bruteforce or O(nlogn) NTT
    return Divide(Mul(p1, p2), base).second;
  vector \langle int \rangle res(n, 0); res[0] = 1;
  for (; k; k \Rightarrow= 1, poly = calc(poly, poly)) {
    if (k & 1) res = calc(res, poly);
  int ans = 0:
  for (int i = 0; i < n; ++i)</pre>
    ans = add(ans, mul(res[i], a[i]));
  return ans;
```

# 7.6 Fast Walsh Transform

```
void fwt(vector <int> &a, bool inv = false) {
    // and : x += y * (1, -1)
    // or : y += x * (1, -1)
    // xor : x = (x + y) * (1, 1/2)
    // y = (x - y) * (1, 1/2)
    int n = __lg(a.size());
    for (int i = 0; i < n; ++i) {
        for (int j = 0; j < 1 << n; ++j) if (j >> i & 1) {
            int x = a[j ^ (1 << i)], y = a[j];
            // do something
        }
    }
}</pre>
```

```
vector<int> subs_conv(vector<int> a, vector<int> b) {
  // c_i = sum_{\{j \& k = 0, j \mid k = i\}} a_j * b_k
  int n = __lg(a.size());
  vector ha(n + 1, vector<int>(1 << n));</pre>
  vector hb(n + 1, vector < int > (1 << n));
  vector c(n + 1, vector<int>(1 << n));
for (int i = 0; i < 1 << n; ++i) {</pre>
    ha[__builtin_popcount(i)][i] = a[i];
    hb[__builtin_popcount(i)][i] = b[i];
  for (int i = 0; i <= n; ++i)
    or_fwt(ha[i]), or_fwt(hb[i]);
  for (int i = 0; i <= n; ++i)
  for (int j = 0; i + j <= n; ++j)</pre>
       for (int k = 0; k < 1 << n; ++k)
         c[i + j][k] = add(c[i + j][k],
           mul(ha[i][k], hb[j][k]));
  for (int i = 0; i <= n; ++i) or_fwt(c[i], true);</pre>
  vector <int> ans(1 << n);</pre>
  for (int i = 0; i < 1 << n; ++i)</pre>
    ans[i] = c[__builtin_popcount(i)][i];
  return ans;
```

# 8 Geometry

#### 8.1 Basic

```
const double eps = 1e-8, PI = acos(-1);
int sign(double x)
{ return fabs(x) \leftarrow eps ? 0 : (x > 0 ? 1 : -1); }
double normalize(double x) {
  while (x < -eps) x += PI * 2;</pre>
  while (x > PI * 2 + eps) x -= PI * 2;
  return x;
template <typename T>
struct P {
  T x, y;
P (T _x, T _y) : x(_x), y(_y) {}
    return P<T>(x + o.x, y + o.y);
  P<T> operator - (P<T> o) {
  return P<T>(x - 0.x, y - 0.y);}
P<T> operator * (T k) {return P<T>(x * k, y * k);}
  P<T> operator / (T k) {return P<T>(x / k, y / k);}
  T operator * (P<T> o) {return x * o.x + y * o.y;}
  T operator ^ (P<T> 0) {return x * 0.y - y * 0.x;}
using Pt = P<double>;
struct Line { Pt a, b; };
struct Cir { Pt o; double r; };
double abs2(Pt o) { return o * o; }
double abs(Pt o) { return sqrt(abs2(o)); }
int ori(Pt o, Pt a, Pt b)
{ return sign((o - a) ^ (o - b)); }
bool btw(Pt a, Pt b, Pt c) // c on segment ab?
{ return ori(a, b, c) == 0 &&
          sign((c - a) * (c - b)) <= 0; }
int pos(Pt a)
{ return sign(a.y) == 0 ? sign(a.x) < 0 : a.y < 0; }
int cmp(Pt a, Pt b)
{ return pos(a) == pos(b) ? sign(a ^ b) > 0 :
         pos(a) < pos(b); }
double area(Pt a, Pt b, Pt c)
{ return fabs((a - b) ^ (a - c)) / 2; }
double angle(Pt a, Pt b)
{ return normalize(atan2(b.y - a.y, b.x - a.x)); }
Pt unit(Pt o) { return o / abs(o); }
Pt rot(Pt a, double o) { // CCW
  double c = cos(o), s = sin(o);
  return Pt(c * a.x - s * a.y, s * a.x + c * a.y);
Pt perp(Pt a) {return Pt(-a.y, a.x);}
Pt proj_vec(Pt a, Pt b, Pt c) { // vector ac proj to ab return (b - a) * ((c - a) * (b - a)) / (abs2(b - a));
Pt proj_pt(Pt a, Pt b, Pt c) { // point c proj to ab
  return proj_vec(a, b, c) + a;
```

#### 8.2 Heart [043c0d]

```
Pt circenter(Pt p0, Pt p1, Pt p2) {
  // radius = abs(center)
  p1 = p1 - p0, p2 = p2 - p0;
  double x1 = p1.x, y1 = p1.y, x2 = p2.x, y2 = p2.y;
  double m = 2. * (x1 * y2 - y1 * x2);
 Pt center(0, 0);
  center.x = (x1 * x1 * y2 - x2 * x2 * y1 + y1 * y2 * (
      y1 - y2)) / m;
 center.y = (x1 * x2 * (x2 - x1) - y1 * y1 * x2 + x1 *
       y2 * y2) / m;
 return center + p0;
Pt incenter(Pt p1, Pt p2, Pt p3) {
 // radius = area / s * 2
double a = abs(p2 - p3), b = abs(p1 - p3), c = abs(p1
       - p2);
  double s = a + b + c;
  return (p1 * a + p2 * b + p3 * c) / s;
Pt masscenter(Pt p1, Pt p2, Pt p3)
{ return (p1 + p2 + p3) / 3; }
Pt orthocenter(Pt p1, Pt p2, Pt p3)
{ return masscenter(p1, p2, p3) * 3 - circenter(p1, p2,
     p3) * 2; }
```

#### 8.3 External Bisector [cafb92]

```
Pt external_bisector(Pt p1, Pt p2, Pt p3) { //213
Pt L1 = p2 - p1, L2 = p3 - p1;
L2 = L2 * abs(L1) / abs(L2);
return L1 + L2;
}
```

# 8.4 Intersection of Segments [e59919]

```
Pt LinesInter(Line a, Line b) {
    double abc = (a.b - a.a) ^ (b.a - a.a);
    double abd = (a.b - a.a) ^ (b.b - a.a);
    if (sign(abc - abd) == 0) return b.b;// no inter
    return (b.b * abc - b.a * abd) / (abc - abd);
}
vector<Pt> SegsInter(Line a, Line b) {
    if (btw(a.a, a.b, b.a)) return {b.a};
    if (btw(a.a, a.b, b.b)) return {b.b};
    if (btw(b.a, b.b, a.a)) return {a.a};
    if (btw(b.a, b.b, a.b)) return {a.b};
    if (ori(a.a, a.b, b.a) * ori(a.a, a.b, b.b) == -1 &&
        ori(b.a, b.b, a.a) * ori(b.a, b.b, a.b) == -1)
    return {LinesInter(a, b)};
    return {};
}
```

# 8.5 Intersection of Circle and Line [75bb3e]

#### 8.6 Intersection of Circles [373889]

# 8.7 Intersection of Polygon and Circle [e005c9]

```
double _area(Pt pa, Pt pb, double r){
  if (abs(pa) < abs(pb)) swap(pa, pb);</pre>
  if (abs(pb) < eps) return 0;</pre>
  double S, h, theta;
  double a = abs(pb), b = abs(pa), c = abs(pb - pa);
  double cosB = pb * (pb - pa) / a / c, B = acos(cosB);
double cosC = (pa * pb) / a / b, C = acos(cosC);
  if (a > r) {
    S = (C / 2) * r * r;
    h = a * b * sin(C) / c;
    if (h < r && B < pi / 2) S -= (acos(h / r) * r * r</pre>
         - h * sqrt(r * r - h * h));
  } else if (b > r) {
    theta = pi - B - asin(sin(B) / r * a);
    S = 0.5 * a * r * sin(theta) + (C - theta) / 2 * r
         * r;
  } else S = 0.5 * sin(C) * a * b;
  return S;
double area_poly_circle(vector<Pt> poly, Pt 0, double r
  double S = 0; int n = poly.size();
  for (int i = 0; i < n; ++i)</pre>
    S += _area(poly[i] - 0, poly[(i + 1) % n] - 0, r) *
          ori(0, poly[i], poly[(i + 1) \% n]);\\
  return fabs(S);
```

# 8.8 Tangent Lines of Circle and Point

#### 8.9 Tangent Lines of Circles [4bf589]

```
vector <Line> tangent(Cir c1, Cir c2, int sign1) {
  // sign1 = 1 for outer tang, -1 for inter tang
  vector <Line> ret;
  double d_sq = abs2(c1.o - c2.o);
  if (sign(d_sq) == 0) return ret;
  double d = sqrt(d_sq);
  Pt v = (c2.0 - c1.0) / d;
  double c = (c1.r - sign1 * c2.r) / d;
  if (c * c > 1) return ret;
  double h = sqrt(max(0.0, 1.0 - c * c));
  Pt p1 = c1.o + n * c1.r;
    Pt p2 = c2.0 + n * (c2.r * sign1);
    if (sign(p1.x - p2.x) == 0 \& sign(p1.y - p2.y) ==
      p2 = p1 + perp(c2.o - c1.o);
    ret.pb({p1, p2});
  return ret;
}
```

#### 8.10 Point In Convex [771a90]

```
bool PointInConvex(const vector<Pt> &C, Pt p, bool
    strict = true) {
    // only works when no three points are collinear
    int a = 1, b = int(C.size()) - 1, r = !strict;
    if (C.size() == 0) return false;
    if (C.size() < 3) return r && btw(C[0], C.back(), p);</pre>
```

#### 8.11 Point In Circle [960672]

# 8.12 Point Segment Distance [651335]

```
double PointSegDist(Pt q0, Pt q1, Pt p) {
  if (sign(abs(q0 - q1)) == 0) return abs(q0 - p);
  if (sign((q1 - q0) * (p - q0)) >= 0 && sign((q0 - q1)
          * (p - q1)) >= 0)
    return fabs(((q1 - q0) ^ (p - q0)) / abs(q0 - q1));
  return min(abs(p - q0), abs(p - q1));
}
```

# 8.13 Convex Hull [eae9b2]

#### 8.14 Minimum Enclosing Circle [1f5028]

```
Cir min_enclosing(vector<Pt> &p) {
  random_shuffle(all(p));
  double r = 0.0;
  Pt cent = p[0];
  for (int i = 1; i < p.size(); ++i) {</pre>
    if (abs2(cent - p[i]) <= r) continue;</pre>
    cent = p[i], r = 0.0;
    for (int j = 0; j < i; ++j) {</pre>
      if (abs2(cent - p[j]) <= r) continue;</pre>
      cent = (p[i] + p[j]) / 2, r = abs2(p[j] - cent);
      for (int k = 0; k < j; ++k) {
        if (abs2(cent - p[k]) <= r) continue;</pre>
        cent = circenter(p[i], p[j], p[k]);
        r = abs2(p[k] - cent);
      }
   }
  return {cent, sqrt(r)};
```

#### 8.15 Union of Circles [53b8f9]

```
vector<pair<double, double>> CoverSegment(Cir a, Cir b)
  double d = abs(a.o - b.o);
  vector<pair<double, double>> res;
  if (sign(a.r + b.r - d) == 0);
  else if (d <= abs(a.r - b.r) + eps) {</pre>
    if (a.r < b.r) res.emplace_back(0, 2 * pi);</pre>
  } else if (d < abs(a.r + b.r) - eps) {</pre>
    double o = acos((a.r * a.r + d * d - b.r * b.r) /
        (2 * a.r * d));
    double z = norm(atan2((b.o - a.o).y, (b.o - a.o).x)
    double l = norm(z - o), r = norm(z + o);
    if (1 > r) res.emplace_back(1, 2 * pi), res.
        emplace_back(0, r);
    else res.emplace_back(1, r);
  }
  return res;
double CircleUnionArea(vector<Cir> c) { // circle
    should be identical
  int n = c.size();
  double a = 0, w;
  for (int i = 0; w = 0, i < n; ++i) {
    vector<pair<double, double>> s = \{\{2 * pi, 9\}\}, z;
    for (int j = 0; j < n; ++j) if (i != j) {</pre>
      z = CoverSegment(c[i], c[j]);
      for (auto &e : z) s.push_back(e);
    sort(s.begin(), s.end());
    auto F = [&] (double t) { return c[i].r * (c[i].r *
         t + c[i].o.x * sin(t) - c[i].o.y * cos(t)); };
    for (auto &e : s) {
      if (e.first > w) a += F(e.first) - F(w);
      w = max(w, e.second);
  return a * 0.5;
```

#### 8.16 Union of Polygons [1eca7c]

```
double polyUnion(vector <vector <Pt>> poly) {
  int n = poly.size();
  double ans = 0;
  auto solve = [&](Pt a, Pt b, int cid) {
    vector <pair <Pt, int>> event;
    for (int i = 0; i < n; ++i) {
      int st = 0, sz = poly[i].size();
      while (st < sz && ori(poly[i][st], a, b) != 1)</pre>
        st++;
      if (st == sz) continue;
      for (int j = 0; j < sz; ++j) {</pre>
        Pt c = poly[i][(j + st) % sz];
        Pt d = poly[i][(j + st + 1) % sz];
        if (sign((a - b) ^ (c - d)) != 0) {
           int ok1 = ori(c, a, b) == 1;
           int ok2 = ori(d, a, b) == 1;
           if (ok1 ^ ok2) event.emplace_back(LinesInter
               ({a, b}, {c, d}), ok1 ? 1 : -1);
        } else if (ori(c, a, b) == 0 && sign((a - b) *
    (c - d)) > 0 && i <= cid) {</pre>
           event.emplace_back(c, -1);
           event.emplace_back(d, 1);
      }
    sort(all(event), [&](pair <Pt, int> i, pair <Pt,</pre>
        int> j) {
      return ((a - i.first) * (a - b)) < ((a - j.first)</pre>
            * (a - b));
    });
    int now = 0;
    Pt lst = a;
    for (auto [x, y] : event) {
      if (btw(a, b, lst) && btw(a, b, x) && !now)
ans += lst ^ x;
      now += y, lst = x;
    }
  };
  for (int i = 0; i < n; ++i) {</pre>
```

```
int sz = poly[i].size();
  for (int j = 0; j < sz; ++j)
      solve(poly[i][j], poly[i][(j + 1) % sz], i);
}
return ans / 2;
}</pre>
```

#### 8.17 Rotating SweepLine [5e4c3d]

```
struct Event {
  Pt d; int u, v;
  bool operator < (const Event &b) const {</pre>
    return sign(d ^ b.d) > 0; }
Pt ref(Pt o) {return pos(o) == 1 ? Pt(-o.x, -o.y) : o;}
void RotatingSweepLine(vector <Pt> &pt) {
  int n = pt.size();
  vector <int> ord(n), pos(n);
  vector <Event> e;
  for (int i = 0; i < n; ++i)</pre>
    for (int j = i + 1; j < n; ++j) if (i ^ j)</pre>
      e.pb({ref(pt[i] - pt[j]), i, j});
  sort(all(e));
  iota(all(ord), 0);
  sort(all(ord), [&](int i, int j) {
    return (sign(pt[i].y - pt[j].y) == 0 ?
        pt[i].x < pt[j].x : pt[i].y < pt[j].y); });</pre>
  for (int i = 0; i < n; ++i) pos[ord[i]] = i;</pre>
  const auto makeReverse = [](auto &v) {
    sort(all(v)); v.resize(unique(all(v)) - v.begin());
    vector <pii> segs;
    for (int i = 0, j = 0; i < v.size(); i = j) {</pre>
      for (; j < v.size() && v[j] - v[i] <= j - i; ++j)</pre>
      segs.emplace_back(v[i], v[j-1]+1+1);
    }
    return segs;
  };
  for (int i = 0, j = 0; i < e.size(); i = j) {</pre>
    vector<int> tmp;
    for (; j < e.size() && !(e[i] < e[j]); j++)</pre>
      tmp.pb(min(pos[e[j].u], pos[e[j].v]));
    for (auto [1, r] : makeReverse(tmp)) {
      reverse(ord.begin() + 1, ord.begin() + r);
      for (int t = 1; t < r; ++t) pos[ord[t]] = t;</pre>
      // update value here
    }
  }
}
```

#### 8.18 Half Plane Intersection [58ae6c]

```
pair <11, 11> area_pair(Line a, Line b)
{ return {(a.b - a.a) ^ (b.a - a.a), (a.b - a.a) ^ (b.b
     - a.a)}; }
bool isin(Line 10, Line 11, Line 12) {
 // Check inter(l1, l2) strictly in l0
  auto [a02X, a02Y] = area_pair(10, 12);
  auto [a12X, a12Y] = area_pair(l1, l2);
 if (a12X - a12Y < 0) a12X *= -1, a12Y *= -1;
 return a02Y * a12X - a02X * a12Y > 0; // C^4
/* Having solution, check size > 2 */
/* --^-- Line.a --^-- Line.b --^-- */
vector<Line> HalfPlaneInter(vector<Line> arr) {
  sort(all(arr), [&](Line a, Line b) {
   Pt A = a.b - a.a, B = b.b - b.a;
    if (pos(A) != pos(B)) return pos(A) < pos(B);</pre>
    if (sign(A ^ B) != 0) return sign(A ^ B) > 0;
    return ori(a.a, a.b, b.b) < 0;</pre>
 });
  deque<Line> dq(1, arr[0]);
  auto same = [&](Pt a, Pt b)
  { return sign(a ^ b) == 0 && pos(a) == pos(b); };
  for (auto p : arr) {
    if (same(dq.back().b - dq.back().a, p.b - p.a))
      continue;
    while (sz(dq) \ge 2 \&\& !isin(p, dq[sz(dq) - 2], dq.
        back())) dq.pop_back();
    while (sz(dq) >= 2 \&\& !isin(p, dq[0], dq[1]))
      dq.pop_front();
    dq.pb(p);
```

#### 8.19 Minkowski Sum [6e64eb]

```
void reorder(vector <Pt> &P) {
    rotate(P.begin(), min_element(all(P), [&](Pt a, Pt b)
        { return make_pair(a.y, a.x) < make_pair(b.y, b.x);
    }), P.end());
}
vector <Pt> Minkowski(vector <Pt> P, vector <Pt> Q) {
    // P, Q: convex polygon, CCW order
    reorder(P), reorder(Q);
    int n = P.size(), m = Q.size();
    P.pb(P[0]), P.pb(P[1]), Q.pb(Q[0]), Q.pb(Q[1]);
    vector <Pt> ans;
    for (int i = 0, j = 0; i < n || j < m; ) {
        ans.pb(P[i] + Q[j]);
        auto val = (P[i + 1] - P[i]) ^ (Q[j + 1] - Q[j]);
        if (val >= 0) i++;
        if (val <= 0) j++;
    }
    return ans;
}</pre>
```

## 8.20 Vector In Polygon [6dac08]

```
// ori(a, b, c) >= 0, valid: "strict" angle from a-b to
    a-c
bool btwangle(Pt a, Pt b, Pt c, Pt p, int strict) {
    return ori(a, b, p) >= strict && ori(a, p, c) >=
        strict;
}

// whether vector{cur, p} in counter-clockwise order
    prv, cur, nxt
bool inside(Pt prv, Pt cur, Pt nxt, Pt p, int strict) {
    if (ori(cur, nxt, prv) >= 0)
        return btwangle(cur, nxt, prv, p, strict);
    return !btwangle(cur, prv, nxt, p, !strict);
}
```

### 8.21 Delaunay Triangulation [52180a]

```
/* Delaunay Triangulation:
Given a sets of points on 2D plane, find a
triangulation such that no points will strictly
inside circumcircle of any triangle.
find : return a triangle contain given point
add_point : add a point into triangulation
A Triangle is in triangulation iff. its has_chd is 0.
Region of triangle u: iterate each u.edge[i].tri,
each points are u.p[(i+1)\%3], u.p[(i+2)\%3]
Voronoi diagram: for each triangle in triangulation,
the bisector of all its edges will split the region.
nearest point will belong to the triangle containing it
const ll inf = MAXC * MAXC * 100;// Lower_bound unknown
struct Tri;
struct Edge {
  Tri* tri; int side;
  Edge(): tri(0), side(0){}
  Edge(Tri* _tri, int _side): tri(_tri), side(_side){}
};
struct Tri {
  Pt p[3];
  Edge edge[3];
  Tri* chd[3];
  Tri() {}
  Tri(const Pt &p0, const Pt &p1, const Pt &p2) {
    p[0] = p0; p[1] = p1; p[2] = p2;
    chd[0] = chd[1] = chd[2] = 0;
  bool has_chd() const { return chd[0] != 0; }
  int num_chd() const {
   return !!chd[0] + !!chd[1] + !!chd[2];
  bool contains(const Pt &q) const {
    for (int i = 0; i < 3; ++i)
```

```
if (ori(p[i], p[(i + 1) % 3], q) < 0)</pre>
        return 0;
    return 1;
} pool[N * 10], *tris;
void edge(Edge a, Edge b) {
  if(a.tri) a.tri->edge[a.side] = b;
  if(b.tri) b.tri->edge[b.side] = a;
struct Trig { // Triangulation
  Trig() {
    the_root = // Tri should at least contain all
        points
      new(tris++) Tri(Pt(-inf, -inf), Pt(inf + inf, -
          inf), Pt(-inf, inf + inf));
  Tri* find(Pt p) { return find(the_root, p); }
  void add_point(const Pt &p) { add_point(find(the_root
       , p), p); }
  Tri* the_root;
  static Tri* find(Tri* root, const Pt &p) {
    while (1) {
      if (!root->has_chd())
        return root;
      for (int i = 0; i < 3 && root->chd[i]; ++i)
        if (root->chd[i]->contains(p)) {
          root = root->chd[i];
          break;
        }
    assert(0); // "point not found"
  void add_point(Tri* root, Pt const& p) {
    Tri* t[3];
    /* split it into three triangles */
    for (int i = 0; i < 3; ++i)</pre>
      t[i] = new(tris++) Tri(root->p[i], root->p[(i +
          1) % 3], p);
    for (int i = 0; i < 3; ++i)</pre>
      edge(Edge(t[i], 0), Edge(t[(i + 1) % 3], 1));
    for (int i = 0; i < 3; ++i)</pre>
      edge(Edge(t[i], 2), root->edge[(i + 2) % 3]);
    for (int i = 0; i < 3; ++i)</pre>
      root->chd[i] = t[i];
    for (int i = 0; i < 3; ++i)
      flip(t[i], 2);
  void flip(Tri* tri, int pi) {
    Tri* trj = tri->edge[pi].tri;
    int pj = tri->edge[pi].side;
    if (!trj) return;
    if (!in_cc(tri->p[0], tri->p[1], tri->p[2], trj->p[
        pj])) return;
    /* flip edge between tri,trj */
    Tri* trk = new(tris++) Tri(tri->p[(pi + 1) % 3],
        trj->p[pj], tri->p[pi]);
    Tri* trl = new(tris++) Tri(trj->p[(pj + 1) % 3],
        tri->p[pi], trj->p[pj]);
    edge(Edge(trk, 0), Edge(trl, 0));
edge(Edge(trk, 1), tri->edge[(pi + 2) % 3]);
    edge(Edge(trk, 2), trj->edge[(pj + 1) % 3]);
    edge(Edge(trl, 1), trj->edge[(pj + 2) % 3]);
edge(Edge(trl, 2), tri->edge[(pi + 1) % 3]);
    tri->chd[0] = trk; tri->chd[1] = trl; tri->chd[2] =
    trj->chd[0] = trk; trj->chd[1] = trl; trj->chd[2] =
    flip(trk, 1); flip(trk, 2);
flip(trl, 1); flip(trl, 2);
 }
};
vector<Tri*> triang; // vector of all triangle
set<Tri*> vst;
void go(Tri* now) { // store all tri into triang
  if (vst.find(now) != vst.end())
    return;
  vst.insert(now);
  if (!now->has_chd())
    return triang.pb(now);
  for (int i = 0; i < now->num_chd(); ++i)
    go(now->chd[i]);
```

```
void build(vector <Pt> &arr) { // build triangulation
  int n = arr.size();
  tris = pool; triang.clear(); vst.clear();
  random_shuffle(all(arr));
  Trig tri; // the triangulation structure
  for (int i = 0; i < n; ++i)
     tri.add_point(arr[i]);
  go(tri.the_root);
}</pre>
```

# 8.22 Triangulation Vonoroi [5c6634]

```
vector<Line> ls[N];
Line make_line(Pt p, Line 1) {
 Pt d = 1.b - 1.a; d = perp(d);
  Pt m = (1.a + 1.b) / 2; // remember to *2
  1 = \{m, m + d\};
  if (ori(l.a, l.b, p) < 0) swap(l.a, l.b);</pre>
  return 1:
void solve(vector <Pt> &oarr) {
 int n = oarr.size();
  map<pair <11, 11>, int> mp;
  vector <Pt> arr = oarr;
  for (int i = 0; i < n; ++i)</pre>
    mp[{arr[i].x, arr[i].y}] = i;
  build(arr); // Triangulation
  for (auto *t : triang) {
    vector<int> p;
    for (int i = 0; i < 3; ++i) {
      pair <11, 11> tmp = {t->p[i].x, t->p[i].y};
      if (mp.count(tmp)) p.pb(mp[tmp]);
    for (int i = 0; i < sz(p); ++i)</pre>
      for (int j = i + 1; j < sz(p); ++j) {
        Line l = \{oarr[p[i]], oarr[p[j]]\};
        ls[p[i]].pb(make_line(oarr[p[i]], 1));
        ls[p[j]].pb(make_line(oarr[p[j]], 1));
      }
  for (int i = 0; i < n; ++i)
    ls[i] = HalfPlaneInter(ls[i]);
```

#### 8.23 3D Point

```
struct Pt {
  double x, y, z;
  Pt(double _x = 0, double _y = 0, double _z = 0): x(_x)
      ), y(_y), z(_z)\{\}
  Pt operator + (const Pt &o) const
  { return Pt(x + o.x, y + o.y, z + o.z); }
  Pt operator - (const Pt &o) const
  { return Pt(x - o.x, y - o.y, z - o.z); }
  Pt operator * (const double &k) const
  { return Pt(x * k, y * k, z * k); }
  Pt operator / (const double &k) const
  { return Pt(x / k, y / k, z / k); }
  double operator * (const Pt &o) const
  { return x * o.x + y * o.y + z * o.z; }
  Pt operator ^ (const Pt &o) const
  { return {Pt(y`* o.z - z *´o.y, z * o.x - x * o.z, x * o.y - y * o.x)}; }
double abs2(Pt o) { return o * o; }
double abs(Pt o) { return sqrt(abs2(o)); }
Pt cross3(Pt a, Pt b, Pt c)
{ return (b - a) ^ (c - a); }
double area(Pt a, Pt b, Pt c)
{ return abs(cross3(a, b, c)); }
double volume(Pt a, Pt b, Pt c, Pt d)
{ return cross3(a, b, c) * (d - a);
bool coplaner(Pt a, Pt b, Pt c, Pt d)
{ return sign(volume(a, b, c, d)) == 0; }
Pt proj(Pt o, Pt a, Pt b, Pt c) // o proj to plane abc
{ Pt n = cross3(a, b, c);
return o - n * ((o - a) * (n / abs2(n)));}
Pt LinePlaneInter(Pt u, Pt v, Pt a, Pt b, Pt c) {
  // intersection of line uv and plane abc
  Pt n = cross3(a, b, c);
  double s = n * (u - v);
  if (sign(s) == 0) return {-1, -1, -1}; // not found
```

```
return v + (u - v) * ((n * (a - v)) / s);
8.24
        3D Convex Hull [2c9f0d]
struct CH3D {
  struct face{int a, b, c; bool ok;} F[8 * N];
  double dblcmp(Pt &p,face &f)
  {return cross3(P[f.a], P[f.b], P[f.c]) * (p - P[f.a])
  int g[N][N], num, n;
  Pt P[N];
  void deal(int p,int a,int b) {
    int f = g[a][b];
    face add;
    if (F[f].ok) {
      if (dblcmp(P[p],F[f]) > eps) dfs(p,f);
      else
        add.a = b, add.b = a, add.c = p, add.ok = 1, g[
            p][b] = g[a][p] = g[b][a] = num, F[num++]=
   }
  void dfs(int p, int now) {
    F[now].ok = 0;
    deal(p, F[now].b, F[now].a), deal(p, F[now].c, F[
        now].b), deal(p, F[now].a, F[now].c);
  bool same(int s,int t){
    Pt &a = P[F[s].a];
    Pt &b = P[F[s].b];
    Pt &c = P[F[s].c];
    return fabs(volume(a, b, c, P[F[t].a])) < eps &&</pre>
        fabs(volume(a, b, c, P[F[t].b])) < eps && fabs(</pre>
        volume(a, b, c, P[F[t].c])) < eps;</pre>
  void init(int _n){n = _n, num = 0;}
  void solve() {
    face add;
    num = 0;
    if(n < 4) return;</pre>
    if([&](){
        for (int i = 1; i < n; ++i)</pre>
        if (abs(P[0] - P[i]) > eps)
        return swap(P[1], P[i]), 0;
        return 1;
        }() || [&](){
        for (int i = 2; i < n; ++i)</pre>
        if (abs(cross3(P[i], P[0], P[1])) > eps)
        return swap(P[2], P[i]), 0;
        return 1;
        }() || [&](){
        for (int i = 3; i < n; ++i)</pre>
        if (fabs(((P[0] - P[1]) ^ (P[1] - P[2])) * (P
             [0] - P[i])) > eps)
                                                               }
                                                             };
        return swap(P[3], P[i]), 0;
        return 1;
        }())return;
                                                             9
    for (int i = 0; i < 4; ++i) {</pre>
      add.a = (i + 1) \% 4, add.b = (i + 2) \% 4, add.c =
            (i + 3) % 4, add.ok = true;
      if (dblcmp(P[i],add) > 0) swap(add.b, add.c);
      g[add.a][add.b] = g[add.b][add.c] = g[add.c][add.
          a] = num;
      F[num++] = add;
    for (int i = 4; i < n; ++i)</pre>
      for (int j = 0; j < num; ++j)</pre>
        if (F[j].ok && dblcmp(P[i],F[j]) > eps) {
          dfs(i, j);
          break:
    for (int tmp = num, i = (num = 0); i < tmp; ++i)</pre>
      if (F[i].ok) F[num++] = F[i];
  double get_area() {
    double res = 0.0;
    if (n == 3)
      return abs(cross3(P[0], P[1], P[2])) / 2.0;
    for (int i = 0; i < num; ++i)</pre>
      res += area(P[F[i].a], P[F[i].b], P[F[i].c]);
```

return res / 2.0;

```
double get_volume() {
  double res = 0.0;
  for (int i = 0; i < num; ++i)</pre>
    res += volume(Pt(0, 0, 0), P[F[i].a], P[F[i].b],
       P[F[i].c]);
 return fabs(res / 6.0);
int triangle() {return num;}
int polygon() {
  int res = 0:
  for (int i = 0, flag = 1; i < num; ++i, res += flag</pre>
      , flag = 1)
    for (int j = 0; j < i && flag; ++j)</pre>
     flag &= !same(i,j);
  return res;
Pt getcent(){
  Pt ans(0, 0, 0), temp = P[F[0].a];
  double v = 0.0, t2;
  for (int i = 0; i < num; ++i)</pre>
    if (F[i].ok == true) {
      Pt p1 = P[F[i].a], p2 = P[F[i].b], p3 = P[F[i].
          c];
      t2 = volume(temp, p1, p2, p3) / 6.0;
      if (t2>0)
        ans.x += (p1.x + p2.x + p3.x + temp.x) * t2,
            ans.y += (p1.y + p2.y + p3.y + temp.y) *
            t2, ans.z += (p1.z + p2.z + p3.z + temp.z
            ) * t2, v += t2;
  ans.x /= (4 * v), ans.y /= (4 * v), ans.z /= (4 * v)
      );
  return ans;
double pointmindis(Pt p) {
  double rt = 99999999;
  for(int i = 0; i < num; ++i)</pre>
    if(F[i].ok == true) {
      Pt p1 = P[F[i].a], p2 = P[F[i].b], p3 = P[F[i].
          c];
      double a = (p2.y - p1.y) * (p3.z - p1.z) - (p2.
          z - p1.z) * (p3.y - p1.y);
      double b = (p2.z - p1.z) * (p3.x - p1.x) - (p2.
          x - p1.x) * (p3.z - p1.z);
      double c = (p2.x - p1.x) * (p3.y - p1.y) - (p2.
          y - p1.y) * (p3.x - p1.x);
      double d = 0 - (a * p1.x + b * p1.y + c * p1.z)
      double temp = fabs(a * p.x + b * p.y + c * p.z
          + d) / sqrt(a * a + b * b + c * c);
      rt = min(rt, temp);
 return rt:
```

#### **Else**

#### 9.1 Pbds

```
#include <ext/pb_ds/priority_queue.hpp>
#include <ext/pb_ds/assoc_container.hpp>
using namespace __gnu_pbds;
#include <ext/rope>
using namespace _
                  _gnu_cxx;
 __gnu_pbds::priority_queue <<mark>int</mark>> pq1, pq2;
pq1.join(pq2); // pq1 += pq2, pq2 = {}
cc_hash_table<int, int> m1;
tree<int, null_type, less<int>, rb_tree_tag,
    tree_order_statistics_node_update> oset;
oset.insert(2), oset.insert(4);
*oset.find_by_order(1), oset.order_of_key(1);// 4 0
bitset <100> BS;
BS.flip(3), BS.flip(5);
BS._Find_first(), BS._Find_next(3); // 3 5
rope <int> rp1, rp2;
rp1.push_back(1), rp1.push_back(3);
rp1.insert(0, 2); // pos, num
rp1.erase(0, 2); // pos, len
rp1.substr(0, 2); // pos, len
rp2.push_back(4);
```

```
National Taiwan University std_abs
rp1 += rp2, rp2 = rp1;
rp2[0], rp2[1]; // 3 4
9.2 Bit Hack
ll next_perm(ll v) { ll t = v | (v - 1);}
   return (t + 1) |
     (((~t & -~t) - 1) >> (__builtin_ctz(v) + 1)); }
9.3 Dynamic Programming Condition
9.3.1 Totally Monotone (Concave/Convex)
\begin{array}{l} \forall i < i', j < j' \text{, } B[i][j] \leq B[i'][j] \implies B[i][j'] \leq B[i'][j'] \\ \forall i < i', j < j' \text{, } B[i][j] \geq B[i'][j] \implies B[i][j'] \geq B[i'][j'] \end{array}
```

# 9.3.2 Monge Condition (Concave/Convex)

```
\begin{array}{l} \forall i < i', j < j' \text{, } B[i][j] + B[i'][j'] \geq B[i][j'] + B[i'][j] \\ \forall i < i', j < j' \text{, } B[i][j] + B[i'][j'] \leq B[i][j'] + B[i'][j] \end{array}
```

#### 9.3.3 Optimal Split Point

```
B[i][j] + B[i+1][j+1] \ge B[i][j+1] + B[i+1][j]
then
                            H_{i,j-1} \le H_{i,j} \le H_{i+1,j}
```

#### 9.4 Smawk Algorithm [5a33b4]

```
11 f(int 1, int r) { }
bool select(int r, int u, int v) {
  // if f(r, v) is better than f(r, v), return true
 return f(r, u) < f(r, v);
// For all 2x2 submatrix:
// If M[1][0] < M[1][1], M[0][0] < M[0][1]
// If M[1][0] == M[1][1], M[0][0] <= M[0][1]
// M[i][ans_i] is the best value in the i-th row
vector<int> solve(vector<int> &r, vector<int> &c) {
 const int n = r.size();
 if (n == 0) return {};
  vector <int> c2;
 for (const int &i : c) {
   while (!c2.empty() && select(r[c2.size() - 1], c2.
        back(), i)) c2.pop_back();
   if (c2.size() < n) c2.pb(i);</pre>
 }
  vector <int> r2;
 for (int i = 1; i < n; i += 2) r2.pb(r[i]);</pre>
 const auto a2 = solve(r2, c2);
  vector <int> ans(n);
 for (int i = 0; i < a2.size(); i++)</pre>
   ans[i * 2 + 1] = a2[i];
  int j = 0;
 for (int i = 0; i < n; i += 2) {
    ans[i] = c2[j];
    const int end = i + 1 == n ? c2.back() : ans[i +
        1];
    while (c2[j] != end) {
      j++;
      if (select(r[i], ans[i], c2[j])) ans[i] = c2[j];
   }
 }
 return ans;
vector<int> smawk(int n, int m) {
  vector<int> row(n), col(m);
 iota(all(row), 0), iota(all(col), 0);
  return solve(row, col);
```

#### 9.5 Slope Trick [d51078]

```
template<typename T>
struct slope_trick_convex {
  T minn = 0, ground_1 = 0, ground_r = 0;
 priority_queue<T, vector<T>, less<T>> left;
 priority_queue<T, vector<T>, greater<T>> right;
  slope_trick_convex() {left.push(numeric_limits<T>::
      min() / 2), right.push(numeric_limits<T>::max() /
       2);}
  void push_left(T x) {left.push(x - ground_l);}
 void push_right(T x) {right.push(x - ground_r);}
```

```
//add a line with slope 1 to the right starting from
  void add_right(T x) {
    T l = left.top() + ground_l;
    if (1 <= x) push_right(x);</pre>
    else push_left(x), push_right(l), left.pop(), minn
  //add a line with slope -1 to the left starting from
  void add_left(T x) {
    T r = right.top() + ground_r;
    if (r >= x) push_left(x);
    else push_right(x), push_left(r), right.pop(), minn
  //val[i]=min(val[j]) for all i-l<=j<=i+r
  void expand(T 1, T r) {ground_1 -= 1, ground_r += r;}
  void shift_up(T x) {minn += x;}
  T get_val(T x) {
    T l = left.top() + ground_l, r = right.top() +
        ground r
    if (x >= 1 \&\& x <= r) return minn;
    if (x < 1) {
      vector<T> trash;
      T cur_val = minn, slope = 1, res;
      while (1) {
        trash.push_back(left.top());
        left.pop();
        if (left.top() + ground_l <= x) {</pre>
          res = cur_val + slope * (1 - x);
        }
        cur_val += slope * (1 - (left.top() + ground_1)
        1 = left.top() + ground_l;
        slope += 1;
      for (auto i : trash) left.push(i);
      return res;
    if(x > r) {
      vector<T> trash;
      T cur_val = minn, slope = 1, res;
      while (1) {
        trash.push_back(right.top());
        right.pop();
        if (right.top() + ground_r >= x) {
          res = cur_val + slope * (x - r);
          break:
        cur_val += slope * ((right.top() + ground_r) -
        r = right.top() + ground_r;
        slope += 1;
      for (auto i : trash) right.push(i);
      return res;
    assert(0);
};
9.6 ALL LCS [5ff948]
```

```
void all_lcs(string s, string t) { // 0-base
  vector<int> h(t.size());
  iota(all(h), 0);
  for (int a = 0; a < s.size(); ++a) {</pre>
    int v = -1;
    for (int c = 0; c < t.size(); ++c)</pre>
      if (s[a] == t[c] || h[c] < v)</pre>
        swap(h[c], v);
    // LCS(s[0, a], t[b, c]) =
    // c - b + 1 - sum([h[i] >= b] | i <= c)
    // h[i] might become -1 !!
```

#### 9.7 Hilbert Curve [1274a3]

```
| 11 hilbert(int n, int x, int y) {
```

```
1l res = 0;
for (int s = n / 2; s; s >>= 1) {
   int rx = (x & s) > 0;
   int ry = (y & s) > 0;
   res += s * 111 * s * ((3 * rx) ^ ry);
   if (ry == 0) {
      if (rx == 1) x = s - 1 - x, y = s - 1 - y;
      swap(x, y);
   }
}
return res;
} // n = 2^k
```

#### 9.8 Random

```
struct custom_hash {
    static uint64_t splitmix64(uint64_t x) {
        x += 0x9e3779b97f4a7c15;
        x = (x ^ (x >> 30)) * 0xbf58476d1ce4e5b9;
        x = (x ^ (x >> 27)) * 0x94d049bb133111eb;
        return x ^ (x >> 31);
    }
    size_t operator()(uint64_t a) const {
        static const uint64_t FIXED_RANDOM = chrono::
            steady_clock::now().time_since_epoch().count();
        return splitmix64(i + FIXED_RANDOM);
    }
};
unordered_map <int, int, custom_hash> m1;
random_device rd; mt19937 rng(rd());
```

#### 9.9 Line Container [673ffd]

```
// only works for integer coordinates!! maintain max
struct Line {
  mutable 11 a, b, p;
  bool operator<(const Line &rhs) const { return a <</pre>
      rhs.a; }
  bool operator<(11 x) const { return p < x; }</pre>
struct DynamicHull : multiset<Line, less<>>> {
  static const ll kInf = 1e18;
  11 Div(11 a, 11 b) { return a / b - ((a ^ b) < 0 && a</pre>
        % b); }
  bool isect(iterator x, iterator y) {
    if (y == end()) { x->p = kInf; return 0; }
    if (x->a == y->a) x->p = x->b > y->b? kInf : -kInf
    else x \rightarrow p = Div(y \rightarrow b - x \rightarrow b, x \rightarrow a - y \rightarrow a);
    return x->p >= y->p;
  void addline(ll a, ll b) { // ax + b
    auto z = insert({a, b, 0}), y = z++, x = y;
    while (isect(y, z)) z = erase(z);
    if (x != begin() \&\& isect(--x, y)) isect(x, y =
         erase(y))
    while ((y = x) != begin() && (--x)->p >= y->p)
        isect(x, erase(y));
  11 query(ll x) {
    auto 1 = *lower_bound(x);
    return 1.a * x + 1.b;
  }
};
```

#### 9.10 Min Plus Convolution [b34de3]

## 9.11 Matroid Intersection

```
Start from S=\emptyset. In each iteration, let • Y_1=\{x\not\in S\mid S\cup\{x\}\in I_1\} • Y_2=\{x\not\in S\mid S\cup\{x\}\in I_2\}
```

If there exists  $x\in Y_1\cap Y_2$  , insert x into S . Otherwise for each  $x\in S, y\not\in S$  , create edges

 $\begin{array}{l} \bullet \quad x \rightarrow y \text{ if } S - \{x\} \cup \{y\} \in I_1\text{.} \\ \bullet \quad y \rightarrow x \text{ if } S - \{x\} \cup \{y\} \in I_2\text{.} \end{array}$ 

Find a shortest path (with BFS) starting from a vertex in  $Y_1$  and ending at a vertex in  $Y_2$  which doesn't pass through any other vertices in  $Y_2$ , and alternate the path. The size of S will be incremented by 1 in each iteration. For the weighted case, assign weight w(x) to vertex x if  $x \in S$  and -w(x) if  $x \not \in S$ . Find the path with the minimum number of edges among all minimum length paths and alternate it.

# 9.12 Simulated Annealing

#### 9.13 Bitset LCS

```
cin >> n >> m;
for (int i = 1, x; i <= n; ++i)
  cin >> x, p[x].set(i);
for (int i = 1, x; i <= m; i++) {
  cin >> x, (g = f) |= p[x];
  f.shiftLeftByOne(), f.set(0);
  ((f = g - f) ^= g) &= g;
}
cout << f.count() << '\n';</pre>
```

# 9.14 Binary Search On Fraction

```
struct 0 {
  11 p, q;
  Q go(Q b, 11 d) { return {p + b.p*d, q + b.q*d}; }
};
bool pred(Q);
// returns smallest p/q in [lo, hi] such that
// pred(p/q) is true, and 0 <= p,q <= N
Q frac_bs(ll N) {
  Q lo{0, 1}, hi{1, 0};
  if (pred(lo)) return lo;
  assert(pred(hi));
  bool dir = 1, L = 1, H = 1;
  for (; L || H; dir = !dir) {
    ll len = 0, step = 1;
    for (int t = 0; t < 2 && (t ? step/=2 : step*=2);)</pre>
      if (Q mid = hi.go(lo, len + step);
          mid.p > N || mid.q > N || dir ^ pred(mid))
        t++;
      else len += step;
    swap(lo, hi = hi.go(lo, len));
    (dir ? L : H) = !!len;
  return dir ? hi : lo;
}
```

# 9.15 Python Misc

```
from [decimal, fractions, math, random] import *
arr = list(map(int, input().split())) # input
setcontext(Context(prec=10, Emax=MAX_EMAX, rounding=
    ROUND FLOOR))
Decimal('1.1') / Decimal('0.2')
Fraction(3, 7)
Fraction(Decimal('1.14'))
Fraction('1.2').limit_denominator(4).numerator
Fraction(cos(pi / 3)).limit_denominator()
S = set(), S.add((a, b)), S.remove((a, b)) # set
if not (a, b) in S:
D = dict(), D[(a, b)] = 1, del D[(a, b)] # dict
for (a, b) in D.items():
arr = [randint(1, C) for i in range(N)]
choice([8, 6, 4, 1]) # random pick one
shuffle(arr)
print(*arr, sep=' ')
```