Contents

8.17 Rotating Caliper 8.18 Rotating SweepLine 8.18 Rotating SweepLine 8.19 Half Plane Intersection 8.20 Minkowski Sum 1.1 Shell Script 1 8.20 Loelaunay Triangulation 1.2 Default Code 1 8.22 Triangulation Vonoroi 1.3 Increase Stack Size 1 1.4 Debug Macro 1 9 Else 1.5 Stress Test Shell* 1 9.1 Bit Hack 1.6 Pragma / FastIO 2 9.2 Dynamic Programming Condition 1.5 Pragma / FastIO 2 9.2 Dynamic Programming Condition 1.5 Pragma / FastIO 2 9.2 Dynamic Programming Condition 1.5 Pragma / FastIO 2 9.2 Dynamic Programming Condition .	2
8.19Half Plane Intersection 1 Basic 1.1 Shell Script 1.2 Default Code 1.3 Increase Stack Size 1.4 Debug Macro 1.5 Stress Test Shell* 1.8 S.19Half Plane Intersection 1.8 S.20Minkowski Sum 1.8 S.21Delaunay Triangulation 1.9 Else 1.9 Else 1.9 Flse 9.1 Bit Hack	2
1 Basic 1 8.20Minkowski Sum 1.1 Shell Script 1 8.21Delaunay Triangulation 1.2 Default Code 1 8.22Triangulation Vonoroi 1.3 Increase Stack Size 1 1 1.4 Debug Macro 1 9 Else 1.5 Stress Test Shell* 1 9.1 Bit Hack	
1.2 Default Code 1 8.22Triangulation Vonoroi <th> 2</th>	2
1.3 Increase Stack Size	
1.4 Debug Macro	2
1.5 Stress Test Shell*	2
1.6 Pragma / FastIO	2
	2
1.7 Divide*	
9.2.2 Monge Condition (Concave/Convex) 2 Data Structure 2 9.2.3 Ontimal Split Point	2
2 Data Structure 2 9.2.3 Optimal Split Point	2
2.2 Splay Tree	
2.3 Link Cut Tree	
2.4 Treap	
2.5 Persistent Segment Tree*	
2.6 2D Segment Tree*	
2.8 Chtholly Tree*	
2.9 Incremental Min Sum*	
9.12Matroid Intersection	2
3 Flow / Matching 5 9.13Python Misc	2
3.2 Min Cost Max Flow	
3.3 Kuhn Munkres	
3.4 SW Min Cut	
3.5 Gomory Hu Tree	
3.6 Blossom	
3.7 Weighted Blossom	
3.0 Flow House	
4 Graph 9 g++ -std=c++17 -DABS -Wall -Wextra -Wshadow	\$1.cpp -o
4.1 Heavy-Light Decomposition	
4.2 Centroid Decomposition	
4.3 Edge BCC	
4.4 Block Cut Tree	
4.6 Negative Cycle*	
4.7 Virtual Tree	
4.8 Directed MST	
4.9 Dominator Tree	
#define nh nush hack	
5 String 5.1 Aho-Corasick Automaton	
5.2 KMP Algorithm	
5.3 Z Algorithm	
5.4 Manacher	
5.5 Suffix Array	
5.6 SAIS	
5.8 Minimum Rotation	
5.9 Palindrome Tree	
	(char*)rsp
5.10Main Lorentz	(/
· (cital) married (STEC) * STEC, STA	
6 Math	
6 Math	
6 Math ; ; ; asm_("movq %0, %%rsp\n"::"r"(p)); 6.2 Miller Rabin / Pollard Rho	
6 Math ; 6.1 Fraction*	
6 Math	
6 Math	
6 Math 6.1 Fraction*	
6 Math 6.1 Fraction*	(T i, U
6 Math 6.1 Fraction*	(T i, U
6 Math 6.1 Fraction*	(T i, U
6 Math 6.1 Fraction*	(T i, U
6 Math 6.1 Fraction* 6.2 Miller Rabin / Pollard Rho 6.3 Ext GCD 6.4 PiCount 6.5 Linear Function Mod Min 6.6 Determinant 6.6 Determinant 6.7 Floor Sum 6.8 Quadratic Residue 6.9 Simplex 6.10Berlekamp Massey 6.11Linear Programming Construction 6.12Euclidean 6.13Theorem 6.13Theorem 6.14 Index ("movq %0, %%rsp\n"::"r"(bak)); 14	
6 Math 6.1 Fraction* 6.2 Miller Rabin / Pollard Rho 6.3 Ext GCD 6.4 PiCount 6.5 Linear Function Mod Min 6.6 Determinant 6.7 Floor Sum 6.8 Quadratic Residue 6.9 Simplex 6.9 Simplex 6.10Berlekamp Massey 6.11Linear Programming Construction 6.12Euclidean 6.13Theorem 6.14Estimation 13 asm("movq %0, %%rsp\n"::"r"(p)); // main	
6 Math 6.1 Fraction* 6.2 Miller Rabin / Pollard Rho 6.3 Ext GCD 6.4 PiCount 6.5 Linear Function Mod Min 6.6 Determinant 6.7 Floor Sum 6.8 Quadratic Residue 6.9 Simplex 6.10Berlekamp Massey 6.11Linear Programming Construction 6.12Euclidean 6.13Theorem 6.14Estimation 6.15General Purpose Numbers 6.17 General Purpose Numbers 6.18 Math 13 3 3 3 3 3 3 3 3 3 3 3 3 3 3 3 3 3 3	
6 Math 6.1 Fraction* 6.2 Miller Rabin / Pollard Rho 6.3 Ext GCD 6.4 PiCount 6.5 Linear Function Mod Min 6.6 Determinant 6.7 Floor Sum 6.8 Quadratic Residue 6.9 Simplex 6.9 Simplex 6.10Berlekamp Massey 6.11Linear Programming Construction 6.12Euclidean 6.13Theorem 6.14Estimation 6.15General Purpose Numbers 15 6.16 Fraction* 13 7 8	
6 Math 6.1 Fraction*	
6 Math 6.1 Fraction*	
6 Math 6.1 Fraction* 6.2 Miller Rabin / Pollard Rho 6.3 Ext GCD 6.4 PiCount 6.5 Linear Function Mod Min 6.6 Determinant 6.7 Floor Sum 6.8 Quadratic Residue 6.9 Simplex 6.10Berlekamp Massey 6.11Linear Programming Construction 6.12Euclidean 6.13Theorem 6.14Estimation 6.14Estimation 6.15General Purpose Numbers 6.16Tips for Generating Funtion 7 Polynomial 7.1 Number Theoretic Transform 7.2 Primes 6.1 Fraction* 13 3	
6 Math	
Section Sect	
6 Math	x)
Second	x)
Section Sect	x)
Section Sect	x)
Section 13	x)
Section	x)
6 Math 13	x)
Second	x)
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```
exit 1
done
echo OK!
./stress.sh main good gen
```

1.6 Pragma / FastIO

```
#pragma GCC optimize("Ofast,inline,unroll-loops")
#pragma GCC target("bmi,bmi2,lzcnt,popcnt,avx2")
#include<unistd.h>
char OB[65536]; int OP;
inline char RC() {
  static char buf[65536], *p = buf, *q = buf;
  return p == q && (q = (p = buf) + read(0, buf, 65536)
       ) == buf ? -1 : *p++;
inline int R() {
  static char c;
  while((c = RC()) < '0'); int a = c ^ '0';
while((c = RC()) >= '0') a *= 10, a += c ^ '0';
inline void W(int n) {
  static char buf[12], p;
  if (n == 0) OB[OP++]='0'; p = 0;
while (n) buf[p++] = '0' + (n % 10), n /= 10;
  for (--p; p >= 0; --p) OB[OP++] = buf[p];
  if (OP > 65520) write(1, OB, OP), OP = 0;
```

1.7 Divide*

```
11 divdown(ll a, ll b) {
  return a / b - (a < 0 && a % b);
11 divup(ll a, ll b) {
  return a / b + (a > 0 && a % b);
a / b < x \rightarrow divdown(a, b) + 1 <= x
a / b \le x \rightarrow divup(a, b) \le x
x < a / b \rightarrow x <= divup(a, b) - 1
x \leftarrow a / b \rightarrow x \leftarrow divdown(a, b)
```

Data Structure

2.1 Leftist Tree

```
struct node {
        ll rk, data, sz, sum;
         node *1, *r;
         node(11 \ k) : rk(0), data(k), sz(1), l(0), r(0), sum(k)
                           ) {}
11 sz(node *p) { return p ? p->sz : 0; }
11 rk(node *p) { return p ? p->rk : -1; }
11 sum(node *p) { return p ? p->sum : 0; }
node *merge(node *a, node *b) {
       if (!a || !b) return a ? a : b;
       if (a->data < b->data) swap(a, b);
       a->r = merge(a->r, b);
       if (rk(a->r) > rk(a->l)) swap(a->r, a->l);
        a - rk = rk(a - r) + 1, a - rk = sz(a - r) + sz(a 
                         1;
       a \rightarrow sum = sum(a \rightarrow 1) + sum(a \rightarrow r) + a \rightarrow data;
       return a;
void pop(node *&o) {
      node *tmp = o;
        o = merge(o->1, o->r);
         delete tmp;
```

2.2 Splay Tree

```
struct Splay {
  int pa[N], ch[N][2], sz[N], rt, _id;
  11 v[N];
  Splay() {}
  void init() {
    rt = 0, pa[0] = ch[0][0] = ch[0][1] = -1;
```

```
sz[0] = 1, v[0] = inf;
  int newnode(int p, int x) {
    int id = _id++;
    v[id] = x, pa[id] = p;
    ch[id][0] = ch[id][1] = -1, sz[id] = 1;
    return id;
  void rotate(int i) {
    int p = pa[i], x = ch[p][1] == i, gp = pa[p], c =
        ch[i][!x];
    sz[p] -= sz[i], sz[i] += sz[p];
    if (~c) sz[p] += sz[c], pa[c] = p;
    ch[p][x] = c, pa[p] = i;
    pa[i] = gp, ch[i][!x] = p;
    if (~gp) ch[gp][ch[gp][1] == p] = i;
  void splay(int i) {
    while (~pa[i]) {
      int p = pa[i];
      if (~pa[p]) rotate(ch[pa[p]][1] == p ^ ch[p][1]
          == i ? i : p);
      rotate(i);
    rt = i:
  int lower_bound(int x) {
    int i = rt, last = -1;
    while (true) {
      if (v[i] == x) return splay(i), i;
      if (v[i] > x) {
        last = i;
        if (ch[i][0] == -1) break;
        i = ch[i][0];
      else {
        if (ch[i][1] == -1) break;
        i = ch[i][1];
      }
    }
    splav(i):
    return last; // -1 if not found
  void insert(int x) {
    int i = lower_bound(x);
    if (i == -1) {
      // assert(ch[rt][1] == -1);
      int id = newnode(rt, x);
      ch[rt][1] = id, ++sz[rt];
      splay(id);
    else if (v[i] != x) {
      splay(i);
      int id = newnode(rt, x), c = ch[rt][0];
      ch[rt][0] = id;
      ch[id][0] = c;
      if (~c) pa[c] = id, sz[id] += sz[c];
      ++sz[rt];
      splay(id);
  }
};
```

2.3 Link Cut Tree

```
// weighted subtree size, weighted path max
struct LCT {
  int ch[N][2], pa[N], v[N], sz[N], sz2[N], w[N], mx[N
  ], _id;
// sz := sum of v in splay, sz2 := sum of v in
      virtual subtree
  // mx := max w in splay
  bool rev[N];
  LCT() : _id(1) {}
  int newnode(int _v, int _w) {
    int x = _id++;
    ch[x][0] = ch[x][1] = pa[x] = 0;
    v[x] = sz[x] = _v;
    sz2[x] = 0;
    w[x] = mx[x] = w;
    rev[x] = false;
    return x;
```

```
void pull(int i) {
    sz[i] = v[i] + sz2[i];
    mx[i] = w[i];
    if (ch[i][0])
      sz[i] += sz[ch[i][0]], mx[i] = max(mx[i], mx[ch[i])
          ][0]]);
    if (ch[i][1])
      sz[i] += sz[ch[i][1]], mx[i] = max(mx[i], mx[ch[i])
          ][1]]);
  void push(int i) {
    if (rev[i]) reverse(ch[i][0]), reverse(ch[i][1]),
        rev[i] = false;
  void reverse(int i) {
    if (!i) return;
    swap(ch[i][0], ch[i][1]);
    rev[i] ^= true;
  bool isrt(int i) {// rt of splay
    if (!pa[i]) return true;
    return ch[pa[i]][0] != i && ch[pa[i]][1] != i;
  void rotate(int i) {
    int p = pa[i], x = ch[p][1] == i, c = ch[i][!x], gp
         = pa[p];
    if (ch[gp][0] == p) ch[gp][0] = i;
    else if (ch[gp][1] == p) ch[gp][1] = i;
    pa[i] = gp, ch[i][!x] = p, pa[p] = i;
    ch[p][x] = c, pa[c] = p;
    pull(p), pull(i);
  void splay(int i) {
    vector<int> anc;
    anc.push_back(i);
    while (!isrt(anc.back())) anc.push_back(pa[anc.back
        ()]);
    while (!anc.empty()) push(anc.back()), anc.pop_back
        ();
    while (!isrt(i)) {
      int p = pa[i];
      if (!isrt(p)) rotate(ch[p][1] == i ^ ch[pa[p]][1]
          == p ? i : p);
      rotate(i);
   }
  }
  void access(int i) {
    int last = 0;
    while (i) {
      splay(i);
      if (ch[i][1])
        sz2[i] += sz[ch[i][1]];
      sz2[i] -= sz[last];
      ch[i][1] = last;
      pull(i), last = i, i = pa[i];
   }
  void makert(int i) {
    access(i), splay(i), reverse(i);
  void link(int i, int j) {
    // assert(findrt(i) != findrt(j));
    makert(i);
    makert(j);
    pa[i] = j;
    sz2[j] += sz[i];
    pull(j);
  void cut(int i, int j) {
    makert(i), access(j), splay(i);
    // assert(sz[i] == 2 && ch[i][1] == j);
    ch[i][1] = pa[j] = 0, pull(i);
  int findrt(int i) {
    access(i), splay(i);
    while (ch[i][0]) push(i), i = ch[i][0];
    splay(i);
    return i;
};
```

2.4 Treap

```
struct node {
  int data, sz;
  node *1, *r;
  node(int k) : data(k), sz(1), l(0), r(0) {}
  void up() {
    sz = 1;
    if (1) sz += 1->sz;
    if (r) sz += r->sz;
  void down() {}
};
int sz(node *a) { return a ? a->sz : 0; }
node *merge(node *a, node *b) {
  if (!a || !b) return a ? a : b;
  if (rand() % (sz(a) + sz(b)) < sz(a))
    return a \rightarrow down(), a \rightarrow r = merge(a \rightarrow r, b), a \rightarrow up(), a
  return b \rightarrow down(), b \rightarrow 1 = merge(a, b \rightarrow 1), b \rightarrow up(), b;
void split(node *o, node *&a, node *&b, int k) {
  if (!o) return a = b = 0, void();
  o->down();
  if (o->data <= k)
    a = o, split(o->r, a->r, b, k), a->up();
  else b = o, split(o->1, a, b->1, k), b->up();
void split2(node *o, node *&a, node *&b, int k) {
  if (sz(o) <= k) return a = o, b = 0, void();
  o->down():
  if (sz(o->1) + 1 <= k)
    a = o, split2(o->r, a->r, b, k - <math>sz(o->l) - 1);
  else b = o, split2(o->1, a, b->1, k);
  o->up();
node *kth(node *o, int k) {
  if (k <= sz(o->1)) return kth(o->1, k);
  if (k == sz(o->1) + 1) return o;
  return kth(o\rightarrow r, k - sz(o\rightarrow l) - 1);
int Rank(node *o, int key) {
  if (!o) return 0;
  if (o->data < key)</pre>
    return sz(o->1) + 1 + Rank(o->r, key);
  else return Rank(o->1, key);
bool erase(node *&o, int k) {
  if (!o) return 0;
  if (o->data == k) {
    node *t = o;
    o->down(), o = merge(o->1, o->r);
    delete t:
    return 1;
  node *&t = k < o->data ? o->l : o->r;
  return erase(t, k) ? o->up(), 1 : 0;
void insert(node *&o, int k) {
  node *a, *b;
  o->down(), split(o, a, b, k)
  o = merge(a, merge(new node(k), b));
  o->up();
void interval(node *&o, int 1, int r) {
 node *a, *b, *c; // [l, r)
  o->down();
  split2(o, a, b, 1), split2(b, b, c, r - 1);
  // operate
  o = merge(a, merge(b, c)), o->up();
```

2.5 Persistent Segment Tree*

```
struct Seg {
    // Persistent Segment Tree, single point modify,
        range query sum
    // 0-indexed, [l, r)
    static Seg mem[M], *pt;
    int l, r, m, val;
    Seg* ch[2];
    Seg () = default;
```

```
Seg (int _l, int _r) : l(_l), r(_r), m(l + r >> 1),
      val(0) {
    if (r - 1 > 1) {
      ch[0] = new (pt++) Seg(1, m);
      ch[1] = new (pt++) Seg(m, r);
  void pull() {val = ch[0]->val + ch[1]->val;}
  Seg* modify(int p, int v) {
    Seg *now = new (pt++) Seg(*this);
    if (r - 1 == 1) {
      now->val = v;
    } else {
      now \rightarrow ch[p >= m] = ch[p >= m] \rightarrow modify(p, v);
      now->pull();
    return now:
  int query(int a, int b) {
    if (a <= 1 && r <= b) return val;</pre>
    int ans = 0;
    if (a < m) ans += ch[0]->query(a, b);
    if (m < b) ans += ch[1]->query(a, b);
    return ans;
} Seg::mem[M], *Seg::pt = mem;
// Init Tree
Seg *root = new (Seg::pt++) Seg(0, n);
```

2.6 2D Segment Tree*

```
// 2D range add, range sum in Log^2
struct seg {
  int 1, r;
  11 sum, 1z;
  seg *ch[2]{};
  seg(int _l, int _r) : 1(_l), r(_r), sum(0), lz(0) {}
  void push() {
    if (lz) ch[0]->add(l, r, lz), ch[1]->modify(l, r,
         1z), 1z = 0;
  void pull() \{sum = ch[0] -> sum + ch[1] -> sum;\}
  void add(int _1, int _r, 11 d) {
  if (_1 <= 1 && r <= _r) {</pre>
       sum += d * (r - 1);
      1z += d;
      return;
    if (!ch[0]) ch[0] = new seg(1, 1 + r >> 1), ch[1] =
          new seg(l + r >> 1, r);
    push();
    if (_1 < 1 + r >> 1) ch[0]->add(_1, _r, d);
    if (l + r >> 1 < _r) ch[1]->add(_l, _r, d);
    pull();
  ill qsum(int _1, int _r) {
   if (_1 <= 1 && r <= _r) return sum;</pre>
    if (!ch[0]) return lz * (min(r, _r) - max(l, _l));
    push();
    11 \text{ res} = 0:
    if (_1 < 1 + r >> 1) res += ch[0]->qsum(_1, _r);
    if (1 + r >> 1 < _r) res += ch[1]->qsum(_1, _r);
    return res:
  }
};
struct seg2 {
 int 1, r;
  seg v, lz;
seg2 *ch[2]{};
  seg2(int _1, int _r) : 1(_1), r(_r), v(0, N), lz(0, N)
    if (1 < r - 1) ch[0] = new seg2(1, 1 + r >> 1), ch
         [1] = new seg2(1 + r >> 1, r);
  void add(int _1, int _r, int _12, int _r2, 11 d) {
  v.add(_12, _r2, d * (min(r, _r) - max(1, _1)));
}
    if (_1 <= 1 && r <= _r) {</pre>
      lz.add(_12, _r2, d);
      return:
    if (_l < l + r >> 1) ch[0]->add(_l, _r, _l2, _r2, d
```

2.7 Zkw*

```
ll mx[N << 1], sum[N << 1], lz[N << 1];
void add(int 1, int r, 11 d) { // [l, r), 0-based
  int len = 1, cntl = 0, cntr = 0;
  for (1 += N, r += N + 1; 1 ^ r ^ 1; 1 >>= 1, r >>= 1,
       len <<= 1) {
     sum[1] += cntl * d, sum[r] += cnt[r] * d;
    if (len > 1) {
      mx[1] = max(mx[1 << 1], mx[1 << 1 | 1]) + lz[1];
      mx[r] = max(mx[r << 1], mx[r << 1 | 1]) + lz[r];
    if (~1 & 1)
      sum[1 ^ 1] += d * len, mx[1 ^ 1] += d, lz[1 ^ 1]
          += d, cntl += len;
    if (r & 1)
      sum[r ^ 1] += d * len, mx[r ^ 1] += d, lz[r ^ 1]
          += d, cntr += len;
  sum[1] += cntl * d, sum[r] += cntr * d;
  if (len > 1) {
    mx[1] = max(mx[1 << 1], mx[1 << 1 | 1]) + lz[1];
    mx[r] = max(mx[r << 1], mx[r << 1 | 1]) + lz[r];
  cntl += cntr;
  for (1 >>= 1; 1; 1 >>= 1) {
    sum[1] += cntl * d;
    mx[1] = max(mx[1 << 1], mx[1 << 1 | 1]) + lz[1];
11 qsum(int 1, int r) {
  ll res = 0, len = 1, cntl = 0, cntr = 0;
  for (1 += N, r += N + 1; 1 ^ r ^ 1; 1 >>= 1, r >>= 1,
       len <<= 1) {
    res += cntl * lz[1] + cntr * lz[r];
    if (~l & 1) res += sum[l ^ 1], cntl += len;
    if (r & 1) res += sum[r ^ 1], cntr += len;
  res += cntl * lz[1] + cntr * lz[r];
  cntl += cntr;
  for (1 >>= 1; 1; 1 >>= 1) res += cnt1 * lz[1];
  return res;
11 qmax(int 1, int r) {
  11 max1 = -INF, maxr = -INF;
  for (1 += N, r += N + 1; 1 ^ r ^ 1; 1 >>= 1, r >>= 1)
    \max 1 += 1z[1], \max[r] += 1z[r];
    if (~1 & 1) maxl = max(maxl, mx[1 ^ 1]);
    if (r & 1) maxr = max(maxr, mx[r ^ 1]);
  \max l = \max(\max l + lz[l], \max r + lz[r]);
  for (1 >>= 1; 1; 1 >>= 1) max1 += lz[1];
  return max1:
}
```

2.8 Chtholly Tree*

```
bool operator () (const interval &a, const interval
        & b) const {
      return a.l < b.l;
    }
  set <interval, cmp> s;
  vector <interval> split(int 1, int r) {
    // split into [0, l), [l, r), [r, n) and return [l, r]
    vector <interval> del, ans, re;
    auto it = s.lower_bound(interval(l, -1, 0));
    if (it != s.begin() && (it == s.end() || 1 < it->1)
        ) {
      --it;
      del.pb(*it);
      if (r < it->r) {
        re.pb(interval(it->l, l, it->v));
        ans.pb(interval(l, r, it->v));
        re.pb(interval(r, it->r, it->v));
      } else {
        re.pb(interval(it->l, l, it->v));
        ans.pb(interval(1, it->r, it->v));
      }
      ++it;
    for (; it != s.end() && it->r <= r; ++it) {</pre>
      ans.pb(*it);
      del.pb(*it);
    if (it != s.end() && it->l < r) {</pre>
      del.pb(*it);
      ans.pb(interval(it->l, r, it->v));
      re.pb(interval(r, it->r, it->v));
    for (interval &i : del)
      s.erase(i);
    for (interval &i : re)
      s.insert(i);
    return ans;
  void merge(vector <interval> a) {
    for (interval &i : a)
      s.insert(i);
  }
};
```

2.9 Incremental Min Sum*

```
struct IncrementalMinSum {
 multiset <int, greater <int>> in;
 multiset <int> out;
 11 sum; int cap;
 DS () : sum(0), cap(0) {}
 void enlarge() {
   if (!out.empty()) {
      int mx = *out.begin();
      sum += mx, in.insert(mx), out.erase(out.begin());
    }
    cap++;
  void insert(int x) {
   if (!cap) {
      out.insert(x);
      return;
    if (in.size() < cap) {</pre>
      in.insert(x), sum += x;
      return;
    int mx = *in.begin();
    if (x < mx) {
      sum -= mx, out.insert(mx), in.erase(in.begin());
      sum += x, in.insert(x);
    } else {
      out.insert(x);
  void erase(int x) {
   if (out.find(x) != out.end()) {
      out.erase(out.lower_bound(x));
    } else {
      in.erase(in.lower_bound(x)), sum -= x;
```

```
if (!out.empty()) {
    int mx = *out.begin();
    sum += mx, out.erase(out.begin()), in.insert(mx
    );
    }
}
```

3 Flow / Matching

3.1 Dinic

```
struct Dinic { // 0-base
  struct edge {
    int to, cap, flow, rev;
  };
  vector<edge> adj[N];
  int s, t, dis[N], cur[N], n;
  int dfs(int u, int cap) {
    if (u == t || !cap) return cap;
     for (int &i = cur[u]; i < (int)adj[u].size(); ++i)</pre>
       edge &e = adj[u][i];
       if (dis[e.to] == dis[u] + 1 && e.flow != e.cap) {
         int df = dfs(e.to, min(e.cap - e.flow, cap));
         if (df) {
           e.flow += df;
           adj[e.to][e.rev].flow -= df;
           return df;
      }
    dis[u] = -1;
    return 0;
  bool bfs() {
     fill_n(dis, n, -1);
     queue<int> q;
     q.push(s), dis[s] = 0;
     while (!q.empty()) {
       int tmp = q.front();
       q.pop();
       for (auto &u : adj[tmp])
         if (!~dis[u.to] && u.flow != u.cap) {
           q.push(u.to);
           dis[u.to] = dis[tmp] + 1;
    return dis[t] != -1;
  int maxflow(int _s, int _t) {
         _s, t = _t;
     int flow = 0, df;
     while (bfs()) {
      fill_n(cur, n, 0);
       while ((df = dfs(s, INF))) flow += df;
     }
     return flow;
  void init(int _n) {
     for (int i = 0; i < n; ++i) adj[i].clear();</pre>
  void reset() {
     for (int i = 0; i < n; ++i)</pre>
       for (auto &j : adj[i]) j.flow = 0;
  void add_edge(int u, int v, int cap) {
    adj[u].pb(edge{v, cap, 0, (int)adj[v].size()});
     adj[v].pb(edge{u, 0, 0, (int)adj[u].size() - 1});
};
```

3.2 Min Cost Max Flow

```
template <typename T>
struct MCMF {
  const T INF = 111 << 60;
  struct edge {
   int v;</pre>
```

```
T f, c:
    edge (int _v, T _f, T _c) : v(_v), f(_f), c(_c) {}
  vector <edge> E;
  vector <vector <int>> adj;
  vector <T> dis, pot;
  vector <int> rt;
  int n, s, t;
  MCMF (int _n, int _s, int _t) : n(_n), s(_s), t(_t) {
    adj.resize(n);
  void add_edge(int u, int v, T f, T c) {
    adj[u].pb(E.size()), E.pb(edge(v, f, c));
    adj[v].pb(E.size()), E.pb(edge(u, 0, -c));
  bool SPFA() {
    rt.assign(n, -1), dis.assign(n, INF);
vector <bool> vis(n, false);
    queue <int> q;
    q.push(s), dis[s] = 0, vis[s] = true;
    while (!q.empty()) {
      int v = q.front(); q.pop();
      vis[v] = false;
      for (int id : adj[v]) if (E[id].f > 0 && dis[E[id
          ].v] > dis[v] + E[id].c + pot[v] - pot[E[id].
          v]) {
          dis[E[id].v] = dis[v] + E[id].c + pot[v] -
               pot[E[id].v], rt[E[id].v] = id;
          if (!vis[E[id].v]) vis[E[id].v] = true, q.
               push(E[id].v);
        }
    return dis[t] != INF;
  bool dijkstra() {
    rt.assign(n, -1), dis.assign(n, INF);
    priority_queue <pair <T, int>, vector <pair <T, int</pre>
        >>, greater <pair <T, int>>> pq;
    dis[s] = 0, pq.emplace(dis[s], s);
    while (!pq.empty()) {
      int d, v; tie(d, v) = pq.top(); pq.pop();
      if (dis[v] < d) continue;</pre>
      for (int id : adj[v]) if (E[id].f > 0 && dis[E[id
          ].v] > dis[v] + E[id].c + pot[v] - pot[E[id].
          v]) {
          dis[E[id].v] = dis[v] + E[id].c + pot[v] -
               pot[E[id].v], rt[E[id].v] = id;
          pq.emplace(dis[E[id].v], E[id].v);
    return dis[t] != INF;
  }
  pair <T, T> solve() {
    pot.assign(n, 0);
    T cost = 0, flow = 0;
    bool fr = true;
    while ((fr ? SPFA() : dijkstra())) {
      for (int i = 0; i < n; i++) {</pre>
        dis[i] += pot[i] - pot[s];
      T add = INF;
      for (int i = t; i != s; i = E[rt[i] ^ 1].v) {
        add = min(add, E[rt[i]].f);
      for (int i = t; i != s; i = E[rt[i] ^ 1].v) {
        E[rt[i]].f -= add, E[rt[i] ^ 1].f += add;
      flow += add, cost += add * dis[t];
      fr = false;
      swap(dis, pot);
    return make_pair(flow, cost);
  }
};
```

Kuhn Munkres

```
template <typename T>
struct KM { // 0-based
  T w[N][N], h1[N], hr[N], slk[N];
  T fl[N], fr[N], pre[N]; int n;
  bool vl[N], vr[N];
```

```
const T INF = 1e9:
   queue <int> q;
   KM (int _n) : n(_n) {
     for (int i = 0; i < n; ++i) for (int j = 0; j < n;
         ++j)
         w[i][j] = -INF;
  void add_edge(int a, int b, int wei) {
    w[a][b] = wei;
  bool check(int x) {
     if (vl[x] = 1, ~fl[x]) return q.push(fl[x]), vr[fl[
         x]] = 1;
     while (\sim x) swap(x, fr[fl[x] = pre[x]]);
    return 0;
  void bfs(int s) {
     fill(slk, slk + n, INF), fill(vl, vl + n, 0), fill(
         vr, vr + n, 0);
     q.push(s), vr[s] = 1;
     while (1) {
      T d:
       while (!q.empty()) {
         int y = q.front(); q.pop();
         for (int x = 0; x < n; ++x)
           if (!v1[x] \&\& s1k[x] >= (d = h1[x] + hr[y] -
               w[x][y])
             if (pre[x] = y, d) slk[x] = d;
             else if (!check(x)) return;
       d = INF;
       for (int x = 0; x < n; ++x)
        if (!v1[x] \&\& d > s1k[x]) d = s1k[x];
       for (int x = 0; x < n; ++x) {
         if (v1[x]) h1[x] += d;
         else slk[x] -= d;
         if (vr[x]) hr[x] -= d;
       for (int x = 0; x < n; ++x) if (!v1[x] && !s1k[x]
            && !check(x)) return;
    }
  T solve() {
     fill(fl, fl + n, -1), fill(fr, fr + n, -1), fill(hr
         , hr + n, 0);
     for (int i = 0; i < n; ++i) hl[i] = *max_element(w[</pre>
         i], w[i] + n);
     for (int i = 0; i < n; ++i) bfs(i);</pre>
     T res = 0;
     for (int i = 0; i < n; ++i) res += w[i][fl[i]];</pre>
     return res;
  }
};
```

3.4 SW Min Cut

```
template <typename T>
struct SW { // 0-based
  T g[N][N], sum[N]; int n;
  bool vis[N], dead[N];
  void init(int _n) {
    n = _n;
for (int i = 0; i < n; ++i) fill(g[i], g[i] + n, 0)</pre>
    fill(dead, dead + n, false);
  void add_edge(int u, int v, T w) {
    g[u][v] += w, g[v][u] += w;
  T solve() {
    T ans = 1 << 30;
    for (int round = 0; round + 1 < n; ++round) {</pre>
      fill(vis, vis + n, false), fill(sum, sum + n, 0);
      int num = 0, s = -1, t = -1;
      while (num < n - round) {</pre>
        int now = -1;
         for (int i = 0; i < n; ++i) if (!vis[i] && !</pre>
             dead[i]) {
             if (now == -1 || sum[now] < sum[i]) now = i</pre>
         s = t, t = now;
```

3.5 Gomory Hu Tree

```
vector <array <int, 3>> GomoryHu(vector <vector <pii>>>
    adj, int n) {
// Tree edge min -> mincut (0-based)
 Dinic flow(n);
  for (int i = 0; i < n; ++i) for (auto [j, w] : adj[i</pre>
      1)
      flow.add_edge(i, j, w);
  flow.record();
 vector <array <int, 3>> ans;
  vector <int> rt(n);
  for (int i = 0; i < n; ++i) rt[i] = 0;
for (int i = 1; i < n; ++i) {</pre>
    int t = rt[i];
    flow.reset(); // clear flows on all edge
    ans.push_back({i, t, flow.solve(i, t)});
    flow.runbfs(i);
    for (int j = i + 1; j < n; ++j) if (rt[j] == t &&
         flow.vis[j]) {
        rt[j] = i;
  return ans:
```

3.6 Blossom

```
struct Matching { // 0-based
 int fa[N], pre[N], match[N], s[N], v[N], n, tk;
  vector <int> g[N];
  queue <int> q;
 Matching (int _n) : n(_n), tk(0) {
    for (int i = 0; i <= n; ++i) match[i] = pre[i] = n;</pre>
    for (int i = 0; i < n; ++i) g[i].clear();</pre>
 void add_edge(int u, int v) {
   g[u].push_back(v), g[v].push_back(u);
 int Find(int u) {
    return u == fa[u] ? u : fa[u] = Find(fa[u]);
 int lca(int x, int y) {
    x = Find(x), y = Find(y);
    for (; ; swap(x, y)) {
      if (x != n) {
        if (v[x] == tk) return x;
        v[x] = tk;
        x = Find(pre[match[x]]);
  void blossom(int x, int y, int 1) {
   while (Find(x) != 1) {
      pre[x] = y, y = match[x];
      if (s[y] == 1) q.push(y), s[y] = 0;
      if (fa[x] == x) fa[x] = 1;
      if (fa[y] == y) fa[y] = 1;
      x = pre[y];
   }
 }
  bool bfs(int r) {
   for (int i = 0; i <= n; ++i) fa[i] = i, s[i] = -1;</pre>
    while (!q.empty()) q.pop();
```

```
q.push(r);
    s[r] = 0;
     while (!q.empty()) {
      int x = q.front(); q.pop();
       for (int u : g[x]) {
         if (s[u] == -1) {
           pre[u] = x, s[u] = 1;
           if (match[u] == n) {
             for (int a = u, b = x, last; b != n; a =
                 last, b = pre[a])
               last = match[b], match[b] = a, match[a] =
                    b;
             return true;
           q.push(match[u]);
           s[match[u]] = 0;
         } else if (!s[u] && Find(u) != Find(x)) {
  int l = lca(u, x);
           blossom(x, u, 1);
           blossom(u, x, 1);
      }
    return false;
  int solve() {
    int res = 0;
    for (int x = 0; x < n; ++x) {
      if (match[x] == n) res += bfs(x);
    return res;
  }
};
```

3.7 Weighted Blossom

```
struct WeightGraph { // 1-based
  static const int inf = INT_MAX;
  static const int maxn = 514;
  struct edge {
    int u, v, w;
    edge(){}
    edge(int u, int v, int w): u(u), v(v), w(w) {}
  int n, n_x;
  edge g[maxn * 2][maxn * 2];
  int lab[maxn * 2];
  int match[maxn * 2], slack[maxn * 2], st[maxn * 2],
      pa[maxn * 2];
  int flo_from[maxn * 2][maxn + 1], S[maxn * 2], vis[
      maxn * 2];
  vector<int> flo[maxn * 2];
  queue<int> q;
  int e_delta(const edge &e) { return lab[e.u] + lab[e.
      v] - g[e.u][e.v].w * 2; }
  void update_slack(int u, int x) { if (!slack[x] ||
      e_delta(g[u][x]) < e_delta(g[slack[x]][x])) slack</pre>
      [x] = u;
  void set_slack(int x) {
    slack[x] = 0;
    for (int u = 1; u <= n; ++u)</pre>
      if (g[u][x].w > 0 \&\& st[u] != x \&\& S[st[u]] == 0)
        update_slack(u, x);
  void q_push(int x) {
    if (x \le n) q.push(x);
    else for (size_t i = 0; i < flo[x].size(); i++)</pre>
        q_push(flo[x][i]);
  void set_st(int x, int b) {
    st[x] = b;
    if (x > n) for (size_t i = 0; i < flo[x].size(); ++</pre>
        i) set_st(flo[x][i], b);
  int get_pr(int b, int xr) {
    int pr = find(flo[b].begin(), flo[b].end(), xr) -
        flo[b].begin();
    if (pr % 2 == 1) {
      reverse(flo[b].begin() + 1, flo[b].end());
      return (int)flo[b].size() - pr;
    return pr;
```

```
void set_match(int u, int v) {
  match[u] = g[u][v].v;
  if (u <= n) return;</pre>
  edge e = g[u][v];
  int xr = flo_from[u][e.u], pr = get_pr(u, xr);
  for (int i = 0; i < pr; ++i) set_match(flo[u][i],</pre>
      flo[u][i ^ 1]);
  set match(xr, v):
  rotate(flo[u].begin(), flo[u].begin() + pr, flo[u].
void augment(int u, int v) {
  for (; ; ) {
    int xnv = st[match[u]];
    set_match(u, v);
    if (!xnv) return;
    set_match(xnv, st[pa[xnv]]);
    u = st[pa[xnv]], v = xnv;
 }
int get_lca(int u, int v) {
  static int t = 0;
  for (++t; u || v; swap(u, v)) {
    if (u == 0) continue;
    if (vis[u] == t) return u;
    vis[u] = t;
    u = st[match[u]];
    if (u) u = st[pa[u]];
  return 0;
void add_blossom(int u, int lca, int v) {
  int b = n + 1;
  while (b <= n_x && st[b]) ++b;</pre>
  if (b > n_x) ++n_x;
  lab[b] = 0, S[b] = 0;
  match[b] = match[lca];
  flo[b].clear();
  flo[b].push_back(lca);
  for (int x = u, y; x != lca; x = st[pa[y]])
    flo[b].push_back(x), flo[b].push_back(y = st[
        match[x]]), q_push(y);
  reverse(flo[b].begin() + 1, flo[b].end());
  for (int x = v, y; x != lca; x = st[pa[y]])
    flo[b].push_back(x), flo[b].push_back(y = st[
        match[x]]), q_push(y);
  set_st(b, b);
  for (int x = 1; x <= n_x; ++x) g[b][x].w = g[x][b].
  for (int x = 1; x \leftarrow n; ++x) flo_from[b][x] = 0;
  for (size_t i = 0; i < flo[b].size(); ++i) {</pre>
    int xs = flo[b][i];
    for (int x = 1; x <= n_x; ++x)</pre>
      if (g[b][x].w == 0 | e_delta(g[xs][x]) <</pre>
          e_delta(g[b][x]))
        g[b][x] = g[xs][x], g[x][b] = g[x][xs];
    for (int x = 1; x <= n; ++x)
      if (flo_from[xs][x]) flo_from[b][x] = xs;
  set_slack(b);
void expand_blossom(int b) {
  for (size_t i = 0; i < flo[b].size(); ++i)</pre>
    set_st(flo[b][i], flo[b][i]);
  int xr = flo_from[b][g[b][pa[b]].u], pr = get_pr(b,
  for (int i = 0; i < pr; i += 2) {</pre>
    int xs = flo[b][i], xns = flo[b][i + 1];
    pa[xs] = g[xns][xs].u;
    S[xs] = 1, S[xns] = 0;
slack[xs] = 0, set_slack(xns);
    q_push(xns);
  S[xr] = 1, pa[xr] = pa[b];
  for (size_t i = pr + 1; i < flo[b].size(); ++i) {</pre>
    int xs = flo[b][i];
    S[xs] = -1, set_slack(xs);
  }
  st[b] = 0;
bool on_found_edge(const edge &e) {
```

```
int u = st[e.u], v = st[e.v];
  if (S[v] == -1) {
    pa[v] = e.u, S[v] = 1;
    int nu = st[match[v]];
    slack[v] = slack[nu] = 0;
    S[nu] = 0, q_push(nu);
  } else if (S[v] == 0) {
    int lca = get_lca(u, v);
    if (!lca) return augment(u,v), augment(v,u), true
    else add_blossom(u, lca, v);
  }
  return false;
bool matching() {
  memset(S + 1, -1, sizeof(int) * n_x);
  memset(slack + 1, 0, sizeof(int) * n_x);
  q = queue<int>();
  for (int x = 1; x <= n_x; ++x)</pre>
    if (st[x] == x && !match[x]) pa[x] = 0, S[x] = 0,
          q_push(x);
  if (q.empty()) return false;
  for (; ; ) {
    while (q.size()) {
      int u = q.front(); q.pop();
      if (S[st[u]] == 1) continue;
      for (int v = 1; v <= n; ++v)
        if (g[u][v].w > 0 && st[u] != st[v]) {
          if (e_delta(g[u][v]) == 0) {
             if (on_found_edge(g[u][v])) return true;
          } else update_slack(u, st[v]);
    int d = inf;
    for (int b = n + 1; b <= n_x; ++b)
      if (st[b] == b && S[b] == 1) d = min(d, lab[b]
    for (int x = 1; x <= n_x; ++x)</pre>
      if (st[x] == x && slack[x]) {
        if (S[x] == -1) d = min(d, e_delta(g[slack[x
             11[x1)):
        else if (S[x] == 0) d = min(d, e_delta(g[
             slack[x]][x]) / 2);
    for (int u = 1; u <= n; ++u) {</pre>
      if (S[st[u]] == 0) {
   if (lab[u] <= d) return 0;</pre>
        lab[u] -= d;
      } else if (S[st[u]] == 1) lab[u] += d;
    for (int b = n + 1; b <= n_x; ++b)</pre>
      if (st[b] == b) {
        if (S[st[b]] == 0) lab[b] += d * 2;
        else if (S[st[b]] == 1) lab[b] -= d * 2;
    q = queue<int>();
    for (int x = 1; x <= n_x; ++x)
      if (st[x] == x && slack[x] && st[slack[x]] != x
            && e_delta(g[slack[x]][x]) == 0)
         if (on_found_edge(g[slack[x]][x])) return
             true;
    for (int b = n + 1; b <= n_x; ++b)
      if (st[b] == b & & S[b] == 1 & & lab[b] == 0)
           expand_blossom(b);
  return false;
pair<long long, int> solve() {
  memset(match + 1, 0, sizeof(int) * n);
  n_x = n;
  int n_matches = 0;
  long long tot_weight = 0;
  for (int u = 0; u <= n; ++u) st[u] = u, flo[u].
      clear();
  int w_max = 0;
  for (int u = 1; u <= n; ++u)</pre>
    for (int v = 1; v <= n; ++v) {</pre>
      flo_from[u][v] = (u == v ? u : 0);
      w_{max} = max(w_{max}, g[u][v].w);
  for (int u = 1; u <= n; ++u) lab[u] = w_max;</pre>
  while (matching()) ++n_matches;
```

```
for (int u = 1; u <= n; ++u)
    if (match[u] && match[u] < u)
        tot_weight += g[u][match[u]].w;
    return make_pair(tot_weight, n_matches);
}
void add_edge(int ui, int vi, int wi) { g[ui][vi].w =
        g[vi][ui].w = wi; }
void init(int _n) {
    n = _n;
    for (int u = 1; u <= n; ++u)
        for (int v = 1; v <= n; ++v)
        g[u][v] = edge(u, v, 0);
}
};</pre>
```

3.8 Flow Model

- Maximum/Minimum flow with lower bound / Circulation problem
 - 1. Construct super source S and sink T.
 - 2. For each edge (x,y,l,u), connect x o y with capacity u-l.
 - 3. For each vertex v, denote by in(v) the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.
 - 4. If in(v)>0, connect $S\to v$ with capacity in(v), otherwise, connect $v\to T$ with capacity -in(v).
 - To maximize, connect $t \to s$ with capacity ∞ (skip this in circulation problem), and let f be the maximum flow from S to T. If $f \neq \sum_{v \in V, in(v) > 0} in(v)$, there's no solution. Otherwise, the maximum flow from s to t is
 - the answer. - To minimize, let f be the maximum flow from S to T. Connect $t \to s$ with capacity ∞ and let the flow from S to T be f'. If $f+f' \neq \sum_{v \in V, in(v)>0} in(v)$, there's no solution. Otherwise, f' is the answer.
 - 5. The solution of each edge e is l_e+f_e , where f_e corresponds to the flow of edge e on the graph.
- Construct minimum vertex cover from maximum matching ${\cal M}$ on bipartite graph (X,Y)
 - 1. Redirect every edge: $y \to x$ if $(x,y) \in M$, $x \to y$ otherwise.
 - 2. DFS from unmatched vertices in X.
 - 3. $x \in X$ is chosen iff x is unvisited.
 - 4. $y \in Y$ is chosen iff y is visited.
- Maximum density induced subgraph
 - 1. Binary search on answer, suppose we're checking answer ${\cal T}$
 - 2. Construct a max flow model, let K be the sum of all weights
 - 3. Connect source $s \to v$, $v \in G$ with capacity K
 - 4. For each edge (u,v,w) in G, connect $u \to v$ and $v \to u$ with capacity w
 - 5. For $v\in G$, connect it with sink $v\to t$ with capacity $K+2T-(\sum_{e\in E(v)}w(e))-2w(v)$
 - 6. T is a valid answer if the maximum flow $f < K \vert V \vert$
- Minimum weight edge cover
 - 1. For each $v \in V$ create a copy v' , and connect $u' \to v'$ with weight w(u,v) .
 - 2. Connect $v \to v'$ with weight $2\mu(v)$, where $\mu(v)$ is the cost of the cheapest edge incident to v.
 - 3. Find the minimum weight perfect matching on G^\prime .
- Project selection problem
 - 1. If $p_v>0$, create edge (s,v) with capacity p_v ; otherwise, create edge (v,t) with capacity $-p_v$.
 - 2. Create edge (u,v) with capacity w with w being the cost of choosing u without choosing v.
 - 3. The mincut is equivalent to the maximum profit of a subset of projects.
- 0/1 quadratic programming

$$\sum_{x} c_{x} x + \sum_{y} c_{y} \bar{y} + \sum_{xy} c_{xy} x \bar{y} + \sum_{xyx'y'} c_{xyx'y'} (x \bar{y} + x' \bar{y'})$$

can be minimized by the mincut of the following graph:

- 1. Create edge (x,t) with capacity c_x and create edge (s,y) with capacity c_y .
- 2. Create edge (x,y) with capacity c_{xy} .
- 3. Create edge (x,y) and edge (x',y') with capacity $c_{xyx'y'}$.

4 Graph

4.1 Heavy-Light Decomposition

```
vector<int> dep, pa, sz, ch, hd, id;
int _id;
void dfs(int i, int p) {
  dep[i] = \sim p ? dep[p] + 1 : 0;
  pa[i] = p, sz[i] = 1, ch[i] = -1;
  for (int j : g[i])
    if (j != p) {
      dfs(j, i);
      if (ch[i] == -1 || sz[ch[i]] < sz[j]) ch[i] = j;</pre>
      sz[i] += sz[j];
void hld(int i, int p, int h) {
  hd[i] = h;
  id[i] = _id++;
  if (~ch[i]) hld(ch[i], i, h);
  for (int j : g[i]) if (j != p && j != ch[i])
    hld(j, i, j);
void query(int i, int j) {
  while (hd[i] != hd[j]) {
    if (dep[hd[i]] < dep[hd[j]]) swap(i, j);</pre>
    query2(id[hd[i]], id[i] + 1), i = pa[hd[i]];
  if (dep[i] < dep[j]) swap(i, j);</pre>
  query2(id[j], id[i] + 1);
```

4.2 Centroid Decomposition

```
vector<vector<int>> dis; // dis[n][logn]
vector<int> pa, sz, dep;
vector<bool> vis;
void dfs_sz(int i, int p) {
  sz[i] = 1;
  for (int j : g[i]) if (j != p && !vis[j])
    dfs_sz(j, i), sz[i] += sz[j];
int cen(int i, int p, int _n) {
  for (int j : g[i]) if (j != p && !vis[j] && sz[j] >
       n / 2)
    return cen(j, i, _n);
  return i;
void dfs_dis(int i, int p, int d) { // from i to
    ancestor with depth d
  dis[i][d] = \sim p ? dis[p][d] + 1 : 0;
  for (int j : g[i]) if (j != p && !vis[j])
    dfs_dis(j, i, d);
void cd(int i, int p, int d) {
  dfs_sz(i, -1), i = cen(i, -1, sz[i]);
  vis[i] = true, pa[i] = p, dep[i] = d;
  dfs_dis(i, -1, d);
  for (int j : g[i]) if (!vis[j])
    cd(j, i, d + 1);
}
```

4.3 Edge BCC

```
vector<int> low, dep, bcc_id, stk;
vector<bool> vis;
int id;
void dfs(int i, int p) {
 low[i] = dep[i] = \sim p ? dep[p] + 1 : 0;
 stk.push_back(i);
 vis[i] = true;
 for (int j : g[i])
    if (j != p)
      if (!vis[j])
        dfs(j, i), low[i] = min(low[i], low[j]);
       low[i] = min(low[i], dep[j]);
 if (low[i] == dep[i]) {
    int id =
             id++
    while (stk.back() != i) {
      int x = stk.back();
      stk.pop_back();
      bcc_id[x] = id;
    stk.pop_back();
```

```
National Taiwan University std_abs
     bcc_id[i] = id;
                                                                     else
                                                                       is[i + n] = true;
| }
                                                                   return true:
       Block Cut Tree
                                                              };
vector<vector<int>> g, _g;
vector<int> dep, low, stk;
void dfs(int i, int p) {
   dep[i] = low[i] = \sim p ? dep[p] + 1 : 0;
   stk.push_back(i);
                                                               template <typename
   for (int j : g[i]) if (j != p) {
     if (dep[j] == -1) {
                                                                 vector <T> dis;
       dfs(j, i), low[i] = min(low[i], low[j]);
                                                                 vector <int> rt;
                                                                 int n; T INF;
       if (low[j] >= dep[i]) {
         int id = _g.size();
         _g.emplace_back();
         while (stk.back() != j) {
           int x = stk.back();
                                                                     >::max()) {
           stk.pop_back();
           _g[x].push_back(id), _g[id].push_back(x);
                                                                   int relax = -1;
         stk.pop_back();
                                                                     relax = -1;
         _g[j].push_back(id), _g[id].push_back(j); _g[i].push_back(id), _g[id].push_back(i);
                                                                            + w) {
      else low[i] = min(low[i], dep[j]);
  }
| }
                                                                       }
                                                                     }
4.5 SCC / 2SAT
struct SAT {
  vector<vector<int>> g;
   vector<int> dep, low, scc_id;
  vector<bool> is;
   vector<int> stk;
  int n, _id, _t;
                                                                       s = rt[s];
  SAT() {}
  void init(int _n) {
  n = _n, _id = _t = 0;
     g.assign(2 * n, vector<int>());
                                                                 }
     dep.assign(2 * n, -1), low.assign(2 * n, -1);
                                                              };
     scc_id.assign(2 * n, -1), is.assign(2 * n, false);
                                                               4.7
     stk.clear();
   void add_edge(int x, int y) {g[x].push_back(y);}
   int rev(int i) {return i < n ? i + n : i - n;}</pre>
   void add_ifthen(int x, int y) {add_clause(rev(x), y)
```

void add_clause(int x, int y) {

add_edge(rev(x), y);

add_edge(rev(y), x);

 $dep[i] = low[i] = _t++;$

if (scc_id[j] == -1) {

if (dep[j] == -1)

if (low[i] == dep[i]) {

stk.pop_back(); scc_id[x] = id;

while (stk.back() != i) {
 int x = stk.back();

for (int i = 0; i < 2 * n; ++i)</pre>

for (int i = 0; i < n; ++i) {</pre>

if (scc_id[i] < scc_id[i + n])</pre>

if (scc_id[i] == scc_id[i + n]) return false;

low[i] = min(low[i], low[j]);

void dfs(int i) {

stk.push_back(i);
for (int j : g[i])

dfs(j);

int id = _id++;

stk.pop_back();

scc_id[i] = id;

if (dep[i] == -1)

is[i] = true;

dfs(i);

bool solve() {

}

```
4.6 Negative Cycle*
vector <pair <int, long long>> adj[N];
struct NegativeCycle {
  vector <int> cycle;
  NegativeCycle () = default;
  NegativeCycle (int _n) : n(_n), INF(numeric_limits<T</pre>
    dis.assign(n, 0), rt.assign(n, -1);
    for (int t = 0; t < n; ++t) {</pre>
      for (int i = 0; i < n; ++i) {</pre>
        for (auto [j, w] : adj[i]) if (dis[j] > dis[i]
          dis[j] = dis[i] + w, rt[j] = i;
          relax = j;
    if (relax != -1) {
      int s = relax:
      for (int i = 0; i < n; ++i) s = rt[s];</pre>
      vector <bool> vis(n, false);
      while (!vis[s]) {
        cycle.push_back(s), vis[s] = true;
      reverse(cycle.begin(), cycle.end());
      Virtual Tree
vector<vector<int>> g
vector<int> st, ed, stk;
void solve(vector<int> v) {
  sort(all(v), [&](int x, int y) {return st[x] < st[y</pre>
      ];});
  int sz = v.size();
  for (int i = 0; i < sz - 1; ++i)
    v.push_back(lca(v[i], v[i + 1]));
  sort(all(v), [&](int x, int y) {return st[x] < st[y</pre>
  v.resize(unique(all(v)) - v.begin());
  stk.clear(); stk.push_back(v[0]);
  for (int i = 1; i < v.size(); ++i) {</pre>
    int x = v[i]:
    while (ed[stk.back()] < ed[x]) stk.pop_back();</pre>
    _g[stk.back()].push_back(x), stk.push_back(x);
  // do something
  for (int i : v) _g[i].clear();
}
4.8 Directed MST
template <typename T> struct DMST { // 1-based
  T g[maxn][maxn], fw[maxn];
  int n, fr[maxn];
  bool vis[maxn], inc[maxn];
  void clear() {
    for (int i = 0; i < maxn; ++i) {</pre>
      for (int j = 0; j < maxn; ++j) g[i][j] = inf;</pre>
      vis[i] = inc[i] = false;
    }
  }
  void addedge(int u, int v, T w) {
    g[u][v] = min(g[u][v], w);
```

```
T query(int root, int _n) {
    if (dfs(root) != n) return -1;
    T ans = 0:
    while (true) {
      for (int i = 1; i <= n; ++i) fw[i] = inf, fr[i] =</pre>
            i;
      for (int i = 1; i <= n; ++i) if (!inc[i]) {</pre>
           for (int j = 1; j <= n; ++j) {
             if (!inc[j] \&\& i != j \&\& g[j][i] < fw[i]) {
               fw[i] = g[j][i];
               fr[i] = j;
            }
          }
        }
      int x = -1;
      for (int i = 1; i <= n; ++i) if (i != root && !</pre>
           inc[i]) {
           int j = i, c = 0;
           while (j != root && fr[j] != i && c <= n) ++c</pre>
               , j = fr[j];
           if (j == root || c > n) continue;
           else { x = i; break; }
      if (!~x) {
        for (int i = 1; i <= n; ++i) if (i != root && !</pre>
             inc[i]) ans += fw[i];
        return ans;
      int y = x;
      for (int i = 1; i <= n; ++i) vis[i] = false;</pre>
      do { ans += fw[y]; y = fr[y]; vis[y] = inc[y] =
           true; } while (y != x);
      inc[x] = false;
      for (int k = 1; k <= n; ++k) if (vis[k]) {</pre>
           for (int j = 1; j <= n; ++j) if (!vis[j]) {</pre>
               if (g[x][j] > g[k][j]) g[x][j] = g[k][j];
               if (g[j][k] < inf \&\& g[j][k] - fw[k] < g[
                    j][x]) g[j][x] = g[j][k] - fw[k];
        }
    return ans;
  int dfs(int now) {
    int r = 1;
    vis[now] = true;
    for (int i = 1; i <= n; ++i) if (g[now][i] < inf &&</pre>
          !vis[i]) r += dfs(i);
    return r;
  }
};
```

4.9 Dominator Tree

```
struct Dominator_tree {
 int n, id;
  vector <vector <int>> adj, radj, bucket;
  vector <int> sdom, dom, vis, rev, par, rt, mn;
 Dominator_tree (int _n) : n(_n), id(0) {
    adj.resize(n), radj.resize(n), bucket.resize(n);
    sdom.resize(n), dom.resize(n, -1), vis.resize(n,
       -1);
    rev.resize(n), rt.resize(n), mn.resize(n), par.
       resize(n);
 }
 void add_edge(int u, int v) {adj[u].pb(v);}
 int query(int v, bool x) {
    if (rt[v] == v) return x ? -1 : v;
    int p = query(rt[v], true);
    if (p == -1) return x ? rt[v] : mn[v];
    if (sdom[mn[v]] > sdom[mn[rt[v]]]) mn[v] = mn[rt[v
        ]];
    rt[v] = p;
    return x ? p : mn[v];
  void dfs(int v) {
    vis[v] = id, rev[id] = v;
    rt[id] = mn[id] = sdom[id] = id, id++;
    for (int u : adj[v]) {
      if (vis[u] == -1) dfs(u), par[vis[u]] = vis[v];
      radj[vis[u]].pb(vis[v]);
```

```
}
  void build(int s) {
    dfs(s);
    for (int i = id - 1; ~i; --i) {
      for (int u : radj[i]) {
        sdom[i] = min(sdom[i], sdom[query(u, false)]);
      if (i) bucket[sdom[i]].pb(i);
      for (int u : bucket[i]) {
        int p = query(u, false);
        dom[u] = sdom[p] == i ? i : p;
      if (i) rt[i] = par[i];
    vector <int> res(n, -1);
    for (int i = 1; i < id; ++i) {</pre>
      if (dom[i] != sdom[i]) dom[i] = dom[dom[i]];
    for (int i = 1; i < id; ++i) res[rev[i]] = rev[dom[</pre>
        i]];
    res[s] = s;
    dom = res;
};
```

5 String

5.1 Aho-Corasick Automaton

```
struct AC {
   int ch[N][26], to[N][26], fail[N], sz;
   vector <int> g[N];
   int cnt[N];
   AC () \{sz = 0, extend();\}
   void extend() {fill(ch[sz], ch[sz] + 26, 0), sz++;}
   int nxt(int u, int v) {
     if (!ch[u][v]) ch[u][v] = sz, extend();
     return ch[u][v];
   int insert(string s) {
     int now = 0;
     for (char c : s) now = nxt(now, c - 'a');
     cnt[now]++;
     return now;
   void build_fail() {
     queue <int> q;
     for (int i = 0; i < 26; ++i) if (ch[0][i]) {</pre>
         q.push(ch[0][i]);
         g[0].push_back(ch[0][i]);
     while (!q.empty()) {
       int v = q.front(); q.pop();
       for (int j = 0; j < 26; ++j) {
  to[v][j] = ch[v][j] ? v : to[fail[v]][j];</pre>
       for (int i = 0; i < 26; ++i) if (ch[v][i]) {</pre>
            int u = ch[v][i], k = fail[v];
            while (k && !ch[k][i]) k = fail[k];
            if (ch[k][i]) k = ch[k][i];
            fail[u] = k;
            cnt[u] += cnt[k], g[k].push_back(u);
            q.push(u);
     }
   int match(string &s) {
     int now = 0, ans = 0;
     for (char c : s) {
       now = to[now][c - 'a'];
if (ch[now][c - 'a']) now = ch[now][c - 'a'];
       ans += cnt[now];
     return ans;
};
```

5.2 KMP Algorithm

```
| vector <int> build_fail(string s) {
```

```
vector <int> f(s.length() + 1, 0);
int k = 0;
for (int i = 1; i < s.length(); ++i) {
    while (k && s[k] != s[i]) k = f[k];
    if (s[k] == s[i]) k++;
    f[i + 1] = k;
}
return f;
}
int match(string s, string t) {
    vector <int> f = build_fail(t);
    int k = 0, ans = 0;
    for (int i = 0; i < s.length(); ++i) {
        while (k && s[i] != t[k]) k = f[k];
        if (s[i] == t[k]) k++;
        if (k == t.length()) ans++, k = f[k];
}
return ans;
}</pre>
```

5.3 Z Algorithm

```
vector <int> build(string s) {
  int n = s.length();
  vector <int> Z(n);
  int l = 0, r = 0;
  for (int i = 0; i < n; ++i) {
    Z[i] = max(min(Z[i - 1], r - i), 0);
    while (i + Z[i] < s.size() && s[Z[i]] == s[i + Z[i] ]) {
        l = i, r = i + Z[i], Z[i]++;
    }
  }
  return Z;
}</pre>
```

5.4 Manacher

```
vector <int> manacher(string &s) {
   string t = "^#";
   for (char c : s) t += c, t += '#';
   t += '&';
   int n = t.length();
   vector <int> r(n, 0);
   int C = 0, R = 0;
   for (int i = 1; i < n - 1; ++i) {
      int mirror = 2 * C - i;
      r[i] = (i < R ? min(r[mirror], R - i) : 0);
      while (t[i - 1 - r[i]] == t[i + 1 + r[i]]) r[i]++;
      if (i + r[i] > R) R = i + r[i], C = i;
   }
   return r;
}
```

5.5 Suffix Array

```
int sa[N], tmp[2][N], c[N], rk[N], lcp[N];
void buildSA(string s) {
 int *x = tmp[0], *y = tmp[1], m = 256, n = s.length()
  for (int i = 0; i < m; ++i) c[i] = 0;</pre>
  for (int i = 0; i < n; ++i) c[x[i] = s[i]]++;
for (int i = 1; i < m; ++i) c[i] += c[i - 1];</pre>
  for (int i = n - 1; ~i; --i) sa[--c[x[i]]] = i;
  for (int k = 1; k < n; k <<= 1) {</pre>
    for (int i = 0; i < m; ++i) c[i] = 0;</pre>
    for (int i = 0; i < n; ++i) c[x[i]]++;</pre>
    for (int i = 1; i < m; ++i) c[i] += c[i - 1];
    int p = 0;
    for (int i = n - k; i < n; ++i) y[p++] = i;
    for (int i = 0; i < n; ++i) if (sa[i] >= k) y[p++]
         = sa[i] - k;
    for (int i = n - 1; ~i; --i) sa[--c[x[y[i]]]] = y[i
    y[sa[0]] = p = 0;
    for (int i = 1; i < n; ++i) {</pre>
      int a = sa[i], b = sa[i - 1];
      if (!(x[a] == x[b] \&\& a + k < n \&\& b + k < n \&\& x
           [a + k] == x[b + k])) p++;
      y[sa[i]] = p;
    if (n == p + 1) break;
```

```
swap(x, y), m = p + 1;
  }
}
void buildLCP(string s) {
  // lcp[i] = LCP(sa[i - 1], sa[i])
  // lcp(i, j) = min(lcp[rk[i] + 1], lcp[rk[i] + 2],
       ..., lcp[rk[j]])
  int n = s.length(), val = 0;
  for (int i = 0; i < n; ++i) rk[sa[i]] = i;</pre>
  for (int i = 0; i < n; ++i) {</pre>
    if (!rk[i]) lcp[rk[i]] = 0;
    else {
      if (val) val--;
      int p = sa[rk[i] - 1];
      while (val + i < n && val + p < n && s[val + i]</pre>
           == s[val + p]) val++;
      lcp[rk[i]] = val;
}
```

5.6 SAIS

```
namespace sfx {
bool _t[N * 2];
int SA[N * 2], H[N], RA[N];
int _s[N * 2], _c[N * 2], x[N], _p[N], _q[N * 2];
void pre(int *sa, int *c, int n, int z) {
  fill_n(sa, n, 0), copy_n(c, z, x);
void induce(int *sa, int *c, int *s, bool *t, int n,
     int z) {
   copy_n(c, z - 1, x + 1);
for (int i = 0; i < n; ++i) if (sa[i] && !t[sa[i] -</pre>
       1]) sa[x[s[sa[i] - 1]]++] = sa[i] - 1;
   copy_n(c, z, x);
   for (int i = n - 1; i >= 0; --i) if (sa[i] && t[sa[i]
        - 1]) sa[--x[s[sa[i] - 1]]] = sa[i] - 1;
void sais(int *s, int *sa, int *p, int *q, bool *t, int
      *c, int n, int z) {
   bool uniq = t[n - 1] = true;
   int nn = 0, nmxz = -1, *nsa = sa + n, *ns = s + n,
       last = -1;
   fill_n(c, z, 0);
   for (int i = 0; i < n; ++i) uniq &= ++c[s[i]] < 2;</pre>
   partial_sum(c, c + z, c);
   if (uniq) {
     for (int i = 0; i < n; ++i) sa[--c[s[i]]] = i;</pre>
     return;
   for (int i = n - 2; i >= 0; --i)
     t[i] = (s[i] == s[i + 1] ? t[i + 1] : s[i] < s[i +
         1]);
   pre(sa, c, n, z);
   for (int i = 1; i <= n - 1; ++i)</pre>
     if (t[i] && !t[i - 1])
       sa[--x[s[i]]] = p[q[i] = nn++] = i;
   induce(sa, c, s, t, n, z);
   for (int i = 0; i < n; ++i)</pre>
     if (sa[i] && t[sa[i]] && !t[sa[i] - 1]) {
       bool neq = last < 0 || !equal(s + sa[i], s + p[q[</pre>
            sa[i]] + 1], s + last);
       ns[q[last = sa[i]]] = nmxz += neq;
   sais(ns, nsa, p + nn, q + n, t + n, c + z, nn, nmxz +
        1);
   pre(sa, c, n, z);
   for (int i = nn - 1; i >= 0; --i)
     sa[--x[s[p[nsa[i]]]]] = p[nsa[i]];
   induce(sa, c, s, t, n, z);
vector<int> build(int *s, int n) {
   copy_n(s, n, _s), _s[n] = 0;
   sais(_s, SA, _p, _q, _t, _c, n + 1, 256);
   vector <int> sa(n);
   for (int i = 0; i < n; ++i)</pre>
     sa[i] = SA[i + 1];
   return sa;
}
```

5.7 Suffix Automaton

```
struct SAM {
  int ch[N][26], len[N], link[N], cnt[N], sz;
  // link -> suffix endpos
  SAM () \{len[0] = 0, link[0] = -1, sz = 1;\}
  void build(string s) {
    int last = 0;
    for (char c : s) {
       int cur = sz++;
       len[cur] = len[last] + 1;
       int p = last;
      while (p \& \ !ch[p][c - 'a']) ch[p][c - 'a'] =
           cur, p = link[p];
       if (p == -1) {
         link[cur] = 0;
       } else {
         int q = ch[p][c - 'a'];
         if (len[p] + 1 == len[q]) {
           link[cur] = q;
         } else {
           int nxt = sz++;
           len[nxt] = len[p] + 1, link[nxt] = link[q];
           for (int j = 0; j < 26; ++j) ch[nxt][j] = ch[</pre>
               q][j];
           while (\sim p \&\& ch[p][c - 'a'] == q) ch[p][c - '
               a'] = nxt, p = link[p];
           link[q] = link[cur] = nxt;
        }
      cnt[cur]++;
      last = cur;
    vector <int> p(sz);
    iota(all(p), 0);
    sort(all(p), [&](int i, int j) {return len[i] > len
         [j];});
    for (int i = 0; i < sz; ++i) cnt[link[p[i]]] += cnt</pre>
         [p[i]];
  }
};
```

5.8 Minimum Rotation

```
string rotate(const string &s) {
  int n = s.length();
  string t = s + s;
  int i = 0, j = 1;
  while (i < n && j < n) {
   int k = 0;
  while (k < n && t[i + k] == t[j + k]) ++k;
   if (t[i + k] <= t[j + k]) j += k + 1;
   else i += k + 1;
   if (i == j) ++j;
  }
  int pos = (i < n ? i : j);
  return t.substr(pos, n);
}</pre>
```

5.9 Palindrome Tree

```
struct PAM {
 int ch[N][26], cnt[N], fail[N], len[N], sz;
  // 0 -> even root, 1 -> odd root
 PAM (string _s) : s(_s) {
   sz = 0;
    extend(), extend();
    len[0] = 0, fail[0] = 1, len[1] = -1;
    int lst = 1;
    for (int i = 0; i < s.length(); ++i) {</pre>
      while (s[i - len[lst] - 1] != s[i]) lst = fail[
          1stl;
      if (!ch[lst][s[i] - 'a']) {
        int idx = extend();
        len[idx] = len[lst] + 2;
        int now = fail[lst];
        while (s[i - len[now] - 1] != s[i]) now = fail[
            now];
        fail[idx] = ch[now][s[i] - 'a'];
        ch[lst][s[i] - 'a'] = idx;
```

```
lst = ch[lst][s[i] - 'a'], cnt[lst]++;
}

void build_count() {
  for (int i = sz - 1; i > 1; --i)
     cnt[fail[i]] += cnt[i];
}
int extend() {
  fill(ch[sz], ch[sz] + 26, 0), sz++;
  return sz - 1;
}
};
```

5.10 Main Lorentz

```
int to_left[N], to_right[N];
vector <array <int, 3>> rep; // l, r, len.
// substr(l ~ r, len * 2) are tandem
void findRep(string &s, int 1, int r) {
  if (r - l == 1) return;
  int m = 1 + r >> 1;
  findRep(s, 1, m), findRep(s, m, r);
  string sl = s.substr(1, m - 1), sr = s.substr(m, r - 1)
       m);
  vector <int> Z = buildZ(sr + "#" + sl);
  for (int i = 1; i < m; ++i) to_right[i] = Z[r - m + 1</pre>
       + i - 1];
  reverse(all(sl));
  Z = buildZ(s1);
  for (int i = 1; i < m; ++i) to_left[i] = Z[m - i -</pre>
      11:
  reverse(all(sl));
  for (int i = 1; i + 1 < m; ++i) {</pre>
    int k1 = to_left[i], k2 = to_right[i + 1], len = m
         - i - 1;
    if (k1 < 1 || k2 < 1 || len < 2) continue;</pre>
    int tl = max(1, len - k2), tr = min(len - 1, k1);
    if (tl <= tr) rep.pb({i + 1 - tr, i + 1 - tl, len})</pre>
  Z = buildZ(sr);
  for (int i = m; i < r; ++i) to_right[i] = Z[i - m];</pre>
  reverse(all(sl)), reverse(all(sr));
Z = buildZ(sl + "#" + sr);
  for (int i = m; i < r; ++i) to_left[i] = Z[m - l + 1</pre>
       + r - i - 1];
  reverse(all(sl)), reverse(all(sr));
  for (int i = m; i + 1 < r; ++i) {</pre>
    int k1 = to_left[i], k2 = to_right[i + 1], len = i
         - m + 1;
    if (k1 < 1 || k2 < 1 || len < 2) continue;</pre>
    int tl = max(len - k2, 1), tr = min(len - 1, k1);
    if (tl <= tr) rep.pb(\{i + 1 - len - tr, i + 1 - len
          - tl, len});
  Z = buildZ(sr + "#" + sl);
  for (int i = 1; i < m; ++i) {</pre>
    if (Z[r - m + 1 + i - 1] >= m - i) {
      rep.pb({i, i, m - i});
  }
}
```

6 Math

6.1 Fraction*

```
struct fraction {
    11 n, d;
    fraction(const l1 _n=0, const l1 _d=1): n(_n), d(_d)
        {
        l1 t = gcd(n, d);
        n /= t, d /= t;
        if (d < 0) n = -n, d = -d;
    }
    fraction operator-() const
    { return fraction(-n, d); }
    fraction operator+(const fraction &b) const
    { return fraction(n * b.d + b.n * d, d * b.d); }
    fraction operator-(const fraction &b) const
    { return fraction(n * b.d - b.n * d, d * b.d); }</pre>
```

```
fraction operator*(const fraction &b) const
{ return fraction(n * b.n, d * b.d); }
fraction operator/(const fraction &b) const
{ return fraction(n * b.d, d * b.n); }
void print() {
  cout << n;
  if (d != 1) cout << "/" << d;
}
};</pre>
```

6.2 Miller Rabin / Pollard Rho

```
11 mul(11 x, 11 y, 11 p) {return (x * y - (11)((long
    double)x / p * y) * p + p) % p;}
vector<11> chk = {2, 325, 9375, 28178, 450775, 9780504,
     1795265022};
ll Pow(ll a, ll b, ll n) {ll res = 1; for (; b; b >>=
    1, a = mul(a, a, n)) if (b \& 1) res = mul(res, a, n)
    ); return res;}
bool check(ll a, ll d, int s, ll n) {
  a = Pow(a, d, n);
  if (a <= 1) return 1;</pre>
  for (int i = 0; i < s; ++i, a = mul(a, a, n)) {</pre>
    if (a == 1) return 0;
    if (a == n - 1) return 1;
  return 0;
bool IsPrime(ll n) {
  if (n < 2) return 0;
  if (n % 2 == 0) return n == 2;
  11 d = n - 1, s = 0;
  while (d \% 2 == 0) d >>= 1, ++s;
  for (ll i : chk) if (!check(i, d, s, n)) return 0;
  return 1;
const vector<ll> small = {2, 3, 5, 7, 11, 13, 17, 19};
11 FindFactor(ll n) {
  if (IsPrime(n)) return 1;
  for (11 p : small) if (n % p == 0) return p;
  11 x, y = 2, d, t = 1;
  auto f = [&](ll a) {return (mul(a, a, n) + t) % n;};
  for (int 1 = 2; ; 1 <<= 1) {
    x = y;
    int m = min(1, 32);
    for (int i = 0; i < 1; i += m) {</pre>
      d = 1:
      for (int j = 0; j < m; ++j) {</pre>
        y = f(y), d = mul(d, abs(x - y), n);
      11 g = \_gcd(d, n);
      if (g == n) {
        1 = 1, y = 2, ++t;
        break:
      if (g != 1) return g;
    }
  }
map <ll, int> res;
void PollardRho(ll n) {
  if (n == 1) return;
  if (IsPrime(n)) return ++res[n], void(0);
  11 d = FindFactor(n);
  PollardRho(n / d), PollardRho(d);
}
```

6.3 Ext GCD

```
//a * p.first + b * p.second = gcd(a, b)
pair<ll, ll> extgcd(ll a, ll b) {
  pair<ll, ll> res;
  if (a < 0) {
    res = extgcd(-a, b);
    res.first *= -1;
    return res;
  }
  if (b < 0) {
    res = extgcd(a, -b);
    res.second *= -1;
    return res;
}</pre>
```

```
if (b == 0) return {1, 0};
res = extgcd(b, a % b);
return {res.second, res.first - res.second * (a / b)
     };
}
```

6.4 PiCount

```
const int V = 10000000, N = 100, M = 100000;
vector<int> primes;
 bool isp[V];
int small_pi[V], dp[N][M];
void sieve(int x){
  for(int i = 2; i < x; ++i) isp[i] = true;</pre>
   isp[0] = isp[1] = false;
   for(int i = 2; i * i < x; ++i) if(isp[i]) for(int j =</pre>
  i * i; j < x; j += i) isp[j] = false;
for(int i = 2; i < x; ++i) if(isp[i]) primes.</pre>
       push_back(i);
void init(){
  sieve(V);
   small_pi[0] = 0;
   for(int i = 1; i < V; ++i) small_pi[i] = small_pi[i -</pre>
        1] + isp[i];
   for(int i = 0; i < M; ++i) dp[0][i] = i;</pre>
   for(int i = 1; i < N; ++i) for(int j = 0; j < M; ++j)</pre>
        dp[i][j] = dp[i - 1][j] - dp[i - 1][j / primes[i]
11 phi(ll n, int a){
   if(!a) return n;
   if(n < M && a < N) return dp[a][n];</pre>
   if(primes[a - 1] > n) return 1;
   if(((ll)primes[a - 1]) * primes[a - 1] >= n && n < V)</pre>
        return small_pi[n] - a + 1;
   11 de = phi(n, a - 1) - phi(n / primes[a - 1], a - 1)
   return de:
11 PiCount(11 n){
   if(n < V) return small_pi[n];</pre>
   int s = sqrt(n + 0.5), y = cbrt(n + 0.5), a =
       small_pi[y];
   ll res = phi(n, a) + a - 1;
   for(; primes[a] <= s; ++a) res -= max(PiCount(n /</pre>
       primes[a]) - PiCount(primes[a]) + 1, 0ll);
}
```

6.5 Linear Function Mod Min

```
ll topos(ll x, ll m) {x %= m; if (x < 0) x += m; return
     x;}
//min value of ax + b \pmod{m} for x \in [0, n - 1]. O(
    Log m)
11 min_rem(11 n, 11 m, 11 a, 11 b) {
  a = topos(a, m), b = topos(b, m);
  for (ll g = \_gcd(a, m); g > 1;) return g * min_rem(n)
       m / g, a / g, b / g) + (b % g);
  for (11 nn, nm, na, nb; a; n = nn, m = nm, a = na, b
      = nb) {
    if (a <= m - a) {
      nn = (a * (n - 1) + b) / m;
      if (!nn) break;
      nn += (b < a);
      nm = a, na = topos(-m, a);
      nb = b < a ? b : topos(b - m, a);
    } else {
      ll lst = b - (n - 1) * (m - a);
      if (lst >= 0) {b = lst; break;}
      nn = -(1st / m) + (1st % m < -a) + 1;
      nm = m - a, na = m % (m - a), nb = b % (m - a);
    }
  }
  return b;
//min\ value\ of\ ax + b\ (mod\ m)\ for\ x\ in\ [0,\ n-1],
    also return min x to get the value. O(\log m)
//{value, x}
pair<ll, 11> min_rem_pos(ll n, ll m, ll a, ll b) {
 a = topos(a, m), b = topos(b, m);
```

6.6 Determinant

```
11 Det(vector <vector <11>>> a) {
  int n = a.size();
  ll det = 1;
  for (int i = 0; i < n; ++i) {</pre>
    if (!a[i][i]) {
      det = -det;
      if (det < 0) det += mod;
      for (int j = i + 1; j < n; ++j) if (a[j][i]) {</pre>
         swap(a[j], a[i]);
      if (!a[i][i]) return 0;
    det = det * a[i][i] % mod;
    ll mul = mpow(a[i][i], mod - 2);
    for (int j = 0; j < n; ++j) a[i][j] = a[i][j] * mul</pre>
          % mod;
    for (int j = 0; j < n; ++j) if (i ^ j) {</pre>
      11 mul = a[j][i];
       for (int k = 0; k < n; ++k) {</pre>
        a[j][k] -= a[i][k] * mul % mod;
        if (a[j][k] < 0) a[j][k] += mod;</pre>
    }
  }
  return det;
}
```

6.7 Floor Sum

6.8 Quadratic Residue

```
int Jacobi(int a, int m) {
  int s = 1;
  for (; m > 1; ) {
    a \% = m;
    if (a == 0) return 0;
    const int r = __builtin_ctz(a);
    if ((r \& 1) \&\& ((m + 2) \& 4)) s = -s;
    if (a \& m \& 2) s = -s;
    swap(a, m);
  }
  return s;
int QuadraticResidue(int a, int p) {
  if (p == 2) return a & 1;
  const int jc = Jacobi(a, p);
  if (jc == 0) return 0;
  if (jc == -1) return -1;
  int b, d;
  for (; ; ) {
   b = rand() % p;
d = (1LL * b * b + p - a) % p;
    if (Jacobi(d, p) == -1) break;
  int f0 = b, f1 = 1, g0 = 1, g1 = 0, tmp;
  for (int e = (1LL + p) >> 1; e; e >>= 1) {
    if (e & 1) {
      tmp = (1LL * g0 * f0 + 1LL * d * (1LL * g1 * f1 %
           p)) % p;
```

```
g1 = (1LL * g0 * f1 + 1LL * g1 * f0) % p;
g0 = tmp;
}
tmp = (1LL * f0 * f0 + 1LL * d * (1LL * f1 * f1 % p
                )) % p;
f1 = (2LL * f0 * f1) % p;
f0 = tmp;
}
return g0;
}
```

6.9 Simplex

```
struct Simplex { // O-based
  using T = long double;
  static const int N = 410, M = 30010;
  const T eps = 1e-7;
  int n, m;
  int Left[M], Down[N];
  // Ax <= b, max c^T x
  // result : v, xi = sol[i]
  T a[M][N], b[M], c[N], v, sol[N];
  bool eq(T a, T b) {return fabs(a - b) < eps;}</pre>
  bool ls(T a, T b) {return a < b && !eq(a, b);}
void init(int _n, int _m) {</pre>
    n = _n, m = _m, v = 0;
    for (int i = 0; i < m; ++i) for (int j = 0; j < n;
         ++j) a[i][j] = 0;
    for (int i = 0; i < m; ++i) b[i] = 0;</pre>
    for (int i = 0; i < n; ++i) c[i] = sol[i] = 0;</pre>
  void pivot(int x, int y) {
    swap(Left[x], Down[y]);
    T k = a[x][y]; a[x][y] = 1;
    vector <int> nz;
    for (int i = 0; i < n; ++i) {</pre>
      a[x][i] /= k;
      if (!eq(a[x][i], 0)) nz.push_back(i);
    b[x] /= k;
    for (int i = 0; i < m; ++i) {</pre>
      if (i == x || eq(a[i][y], 0)) continue;
      k = a[i][y], a[i][y] = 0;
b[i] -= k * b[x];
      for (int j : nz) a[i][j] -= k * a[x][j];
    if (eq(c[y], 0)) return;
    k = c[y], c[y] = 0, v += k * b[x];
    for (int i : nz) c[i] -= k * a[x][i];
  // 0: found solution, 1: no feasible solution, 2:
      unbounded
  int solve() {
    for (int i = 0; i < n; ++i) Down[i] = i;</pre>
    for (int i = 0; i < m; ++i) Left[i] = n + i;</pre>
    while (1) {
      int x = -1, y = -1;
      for (int i = 0; i < m; ++i) if (ls(b[i], 0) && (x</pre>
            == -1 \mid \mid b[i] < b[x]) x = i;
      if (x == -1) break;
      for (int i = 0; i < n; ++i) if (ls(a[x][i], 0) &&</pre>
            (y == -1 \mid | a[x][i] < a[x][y])) y = i;
      if (y == -1) return 1;
      pivot(x, y);
    while (1) {
      int x = -1, y = -1;
      for (int i = 0; i < n; ++i) if (ls(0, c[i]) && (y</pre>
            == -1 \mid \mid c[i] > c[y])) y = i;
      if (y == -1) break;
      for (int i = 0; i < m; ++i) if (ls(0, a[i][y]) &&</pre>
            (x == -1 \mid | b[i] / a[i][y] < b[x] / a[x][y])
           ) x = i;
      if (x == -1) return 2;
      pivot(x, y);
    for (int i = 0; i < m; ++i) if (Left[i] < n) sol[</pre>
         Left[i]] = b[i];
    return 0;
```

6.10 Berlekamp Massey

```
vector <1l> BerlekampMassey(vector <1l> a) {
  // find min |c| such that a_n = sum c_j * a_{n - j}
      1}, 0-based
  // O(N^2), if |c| = k, |a| >= 2k sure correct
  auto f = [&](vector<11> v, 11 c) {
    for (11 &x : v) x = mul(x, c);
  };
  vector <11> c, best;
  int pos = 0, n = a.size();
for (int i = 0; i < n; ++i) {</pre>
    ll error = a[i];
    for (int j = 0; j < c.size(); ++j) error = sub(</pre>
        error, mul(c[j], a[i - 1 - j]));
    if (error == 0) continue;
    11 inv = mpow(error, mod - 2);
    if (c.empty()) {
      c.resize(i + 1);
      pos = i;
      best.pb(inv);
    } else {
      vector <11> fix = f(best, error);
      fix.insert(fix.begin(), i - pos - 1, 0);
      if (fix.size() >= c.size()) {
        best = f(c, sub(0, inv));
        best.insert(best.begin(), inv);
        pos = i:
        c.resize(fix.size());
      for (int j = 0; j < fix.size(); ++j) c[j] = add(c</pre>
          [j], fix[j]);
    }
  }
  return c;
```

6.11 Linear Programming Construction

Standard form: maximize $\mathbf{c}^T \mathbf{x}$ subject to $A\mathbf{x} \leq \mathbf{b}$ and $\mathbf{x} \geq 0$. Dual LP: minimize $\mathbf{b}^T\mathbf{y}$ subject to $A^T\mathbf{y} \geq \mathbf{c}$ and $\mathbf{x} \geq 0$. $\bar{\mathbf{x}}$ and $\bar{\mathbf{y}}$ are optimal if and only if for all $i \in [1,n]$, either $\bar{x}_i = 0$ or $\sum_{j=1}^m A_{ji}\bar{y}_j = c_i$ holds and for all $i \in [1,m]$ either $\bar{y}_i = 0$ or $\sum_{j=1}^{n} A_{ij} \bar{x}_j = b_j$ holds.

- 1. In case of minimization, let $c_i'=-c_i$ 2. $\sum_{1\leq i\leq n}A_{ji}x_i\geq b_j\to \sum_{1\leq i\leq n}-A_{ji}x_i\leq -b_j$ 3. $\sum_{1\leq i\leq n}A_{ji}x_i=b_j$
 - $\begin{array}{ll} \bullet & \sum_{1 \leq i \leq n} A_{ji} x_i \leq b_j \\ \bullet & \sum_{1 \leq i \leq n} A_{ji} x_i \geq b_j \end{array}$
- 4. If x_i has no lower bound, replace x_i with $x_i x_i'$

6.12 Euclidean

- $m = \lfloor \frac{an+b}{c} \rfloor$ Time complexity: $O(\log n)$

$$\begin{split} f(a,b,c,n) &= \sum_{i=0}^n \lfloor \frac{ai+b}{c} \rfloor \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)}{2} + \lfloor \frac{b}{c} \rfloor \cdot (n+1) \\ +f(a \bmod c, b \bmod c, c, n), & a \geq c \vee b \geq c \\ 0, & n < 0 \vee a = 0 \\ nm - f(c, c-b-1, a, m-1), & \text{otherwise} \end{cases} \end{split}$$

$$\begin{split} g(a,b,c,n) &= \sum_{i=0}^n i \lfloor \frac{ai+b}{c} \rfloor \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor \cdot \frac{n(n+1)}{2} \\ + g(a \bmod c, b \bmod c, c, n), & a \geq c \vee b \geq c \\ 0, & n < 0 \vee a = 0 \\ \frac{1}{2} \cdot (n(n+1)m - f(c, c-b-1, a, m-1) \\ -h(c, c-b-1, a, m-1)), & \text{otherwise} \end{cases} \end{split}$$

$$\begin{split} h(a,b,c,n) &= \sum_{i=0}^n \lfloor \frac{ai+b}{c} \rfloor^2 \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor^2 \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor^2 \cdot (n+1) \\ + \lfloor \frac{a}{c} \rfloor \cdot \lfloor \frac{b}{c} \rfloor \cdot n(n+1) \\ + h(a \bmod c, b \bmod c, c, n) \\ + 2\lfloor \frac{a}{c} \rfloor \cdot g(a \bmod c, b \bmod c, c, n) \\ + 2\lfloor \frac{b}{c} \rfloor \cdot f(a \bmod c, b \bmod c, c, n), & a \geq c \vee b \geq c \\ 0, & n < 0 \vee a = 0 \\ nm(m+1) - 2g(c, c-b-1, a, m-1) \\ - 2f(c, c-b-1, a, m-1) - f(a, b, c, n), & \text{otherwise} \end{cases} \end{split}$$

6.13 Theorem

Kirchhoff's Theorem

Denote L be a $n \times n$ matrix as the Laplacian matrix of graph G, where $L_{ii}=d(i)$, $L_{ij}=-c$ where c is the number of edge (i,j) in

- The number of undirected spanning in G is $|\det(\tilde{L}_{11})|$.
- The number of directed spanning tree rooted at r in ${\it G}$ is
- Tutte's Matrix

Let D be a n imes n matrix, where $d_{ij} = x_{ij}$ (x_{ij} is chosen uniformly at random) if i < j and $(i,j) \in E$, otherwise $d_{ij} = -d_{ji}$. $\frac{rank(D)}{2}$ is the maximum matching on G.

- Cayley's Formula
 - Given a degree sequence d_1, d_2, \ldots, d_n for each $\emph{labeled}$ vertices, there are

$$\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$$

- Let $T_{n,k}$ be the number of labeled forests on n vertices with k components, such that vertex $1,2,\ldots,k$ belong to different components. Then $T_{n,k}=kn^{n-k-1}$.
- Erdős-Gallai Theorem

A sequence of non-negative integers $d_1 \geq d_2 \geq \ldots \geq d_n$ can be represented as the degree sequence of a finite simple graph on n vertices if and only if $d_1+d_2+\ldots+d_n$ is even and

$$\sum_{i=1}^k d_i \leq k(k-1) + \sum_{i=k+1}^n \min(d_i,k)$$

holds for all $1 \leq k \leq n$.

• Burnside's Lemma

Let X be a set and G be a group that acts on X . For $g\in G$, denote by X^g the elements fixed by g :

$$X^g = \{x \in X \mid gx \in X\}$$

Then

$$|X/G| = \frac{1}{|G|} \sum_{g \in G} |X^g|.$$

• Gale-Ryser theorem

A pair of sequences of nonnegative integers $a_1 \geq \cdots \geq a_n$ and b_1, \ldots, b_n is bigraphic if and only if $\sum_{i=1}^n a_i = \sum_{i=1}^n b_i$ and $\sum_{i=1}^k a_i \leq a_i$ $\sum \mathsf{min}(b_i,k)$ holds for every $1 \leq k \leq n$.

• Fulkerson-Chen-Anstee theorem

A sequence $(a_1,b_1),\ldots,(a_n,b_n)$ of nonnegative integer pairs with $a_1 \geq \cdots \geq a_n$ is digraphic if and only if $\sum_{i=1}^n a_i = \sum_{i=1}^n b_i$ and $\sum_{i=1}^k a_i \leq \sum_{i=1}^k \min(b_i,k-1) + \sum_{i=k+1}^n \min(b_i,k) \text{ holds for every } 1 \leq k \leq n.$

- Möbius inversion formula
 - $f(n) = \sum_{d|n} g(d) \Leftrightarrow g(n) = \sum_{d|n} \mu(d) f(\frac{n}{d})$ $f(n) = \sum_{n|d} g(d) \Leftrightarrow g(n) = \sum_{n|d} \mu(\frac{d}{n}) f(d)$
- Spherical cap
 - A portion of a sphere cut off by a plane. r: sphere radius, a: radius of the base of the cap, h: height of the cap, θ : $\arcsin(a/r)$. Volume = $\pi h^2(3r-h)/3 = \pi h(3a^2+h^2)/6 = \pi r^3(2+\cos\theta)(1-\cos\theta)^2/3$. Area = $2\pi rh = \pi(a^2+h^2) = 2\pi r^2(1-\cos\theta)$.
- Chinese Remainder Theorem
 - $x \equiv a_i \pmod{m_i}$
 - $M = \prod m_i, M_i = M/m_i$
 - $t_i M_i \equiv 1 \pmod{m_i}$
 - $x = \sum a_i t_i M_i \pmod{M}$

6.14 Estimation

- The number of divisors of n is at most around 100 for n<5e4, 500 for n<1e7, 2000 for n<1e10, 200000 for n<1e19.
- The number of ways of writing n as a sum of positive integers, disregarding the order of the summands. 1,1,2,3,5,7,11,15,22,30 for $n=0\sim 9$, 627 for n=20, $\sim 2e5$ for n=50, $\sim 2e8$ for n=100.
- Total number of partitions of n distinct elements: $B(n)=1,1,2,5,15,52,203,877,4140,21147,115975,678570,4213597, 27644437,190899322, <math display="inline">\ldots$

6.15 General Purpose Numbers

• Bernoulli numbers

$$\begin{split} &B_0-1, B_1^{\pm}=\pm\tfrac{1}{2}, B_2=\tfrac{1}{6}, B_3=0\\ &\sum_{j=0}^m \binom{m+1}{j} B_j=0\text{, EGF is } B(x)=\tfrac{x}{e^x-1}=\sum_{n=0}^\infty B_n \frac{x^n}{n!}\text{.}\\ &S_m(n)=\sum_{k=1}^n k^m=\frac{1}{m+1}\sum_{k=0}^m \binom{m+1}{k} B_k^+ n^{m+1-k} \end{split}$$

- Stirling numbers of the second kind Partitions of \boldsymbol{n} distinct elements into exactly \boldsymbol{k} groups.

$$\begin{split} S(n,k) &= S(n-1,k-1) + kS(n-1,k), S(n,1) = S(n,n) = 1 \\ S(n,k) &= \frac{1}{k!} \sum_{i=0}^k (-1)^{k-i} {k \choose i} i^n \\ x^n &= \sum_{i=0}^n S(n,i)(x)_i \end{split}$$

• Pentagonal number theorem

$$\prod_{n=1}^{\infty} (1 - x^n) = 1 + \sum_{k=1}^{\infty} (-1)^k \left(x^{k(3k+1)/2} + x^{k(3k-1)/2} \right)$$

Catalan numbers

$$C_n^{(k)} = \frac{1}{(k-1)n+1} {kn \choose n}$$
$$C^{(k)}(x) = 1 + x[C^{(k)}(x)]^k$$

Eulerian numbers

Number of permutations $\pi \in S_n$ in which exactly k elements are greater than the previous element. k j:s s.t. $\pi(j) > \pi(j+1)$, k+1 j:s s.t. $\pi(j) \geq j$, k j:s s.t. $\pi(j) > j$. E(n,k) = (n-k)E(n-1,k-1) + (k+1)E(n-1,k) E(n,0) = E(n,n-1) = 1 $E(n,k) = \sum_{j=0}^k (-1)^j \binom{n+1}{j} (k+1-j)^n$

6.16 Tips for Generating Funtion

• Ordinary Generating Function $A(x) = \sum_{i>0} a_i x^i$

```
 \begin{array}{l} - \ A(rx) \Rightarrow r^n a_n \\ - \ A(x) + B(x) \Rightarrow a_n + b_n \\ - \ A(x)B(x) \Rightarrow \sum_{i=0}^n a_i b_{n-i} \\ - \ A(x)^k \Rightarrow \sum_{i=1}_{i+1} a_{i} b_{n-i} \\ - \ A(x)^k \Rightarrow \sum_{i=0}^n a_i b_{n-i} \\ - \ xA(x)' \Rightarrow n a_n \\ - \ \frac{A(x)}{1-x} \Rightarrow \sum_{i=0}^n a_i \end{array}
```

• Exponential Generating Function $A(x) = \sum_{i>0} \frac{a_i}{i!} x_i$

```
- A(x) + B(x) \Rightarrow a_n + b_n

- A^{(k)}(x) \Rightarrow a_{n+k_n}

- A(x)B(x) \Rightarrow \sum_{i=0}^{k_n} {n \choose i} a_i b_{n-i}

- A(x)^k \Rightarrow \sum_{i_1+i_2+\dots+i_k=n}^{k_n} {n \choose i_1, i_2, \dots, i_k} a_{i_1} a_{i_2} \dots a_{i_k}

- xA(x) \Rightarrow na_n
```

• Special Generating Function

$$- (1+x)^n = \sum_{i\geq 0} \binom{n}{i} x^i$$

$$- \frac{1}{(1-x)^n} = \sum_{i\geq 0} \binom{n}{i-1} x^i$$

7 Polynomial

7.1 Number Theoretic Transform

```
void operator()(vector<ll>& a, bool inv = false) { //
       0 \leftarrow a[i] \leftarrow P
     int x = 0, n = a.size();
     for (int j = 1; j < n - 1; ++j) {
  for (int k = n >> 1; (x ^= k) < k; k >>= 1);
       if (j < x) swap(a[x], a[j]);</pre>
     for (int L = 2; L <= n; L <<= 1) {
       int dx = N / L, d1 = L >> 1;
       for (int i = 0; i < n; i += L) {</pre>
         for (int j = i, x = 0; j < i + dl; ++j, x += dx
            ll tmp = mul(a[j + dl], w[x]);
           a[j + dl] = sub(a[j], tmp);
           a[j] = add(a[j], tmp);
       }
     if (inv) {
       reverse(a.begin() + 1, a.end());
       ll invn = mpow(n, mod - 2);
       for (int i = 0; i < n; ++i) a[i] = mul(a[i], invn</pre>
  }
} ntt;
```

7.2 Primes

```
Prime
             Root
                    Prime
                                   Root
                    167772161
12289
                    104857601
             11
40961
                    985661441
65537
                    998244353
786433
             10
                    1107296257
                                  10
5767169
                    2013265921
             3
                                   31
7340033
                     2810183681
23068673
                    2885681153
469762049
                    605028353
```

7.3 Polynomial Operations

```
vector <ll> Mul(vector <ll> a, vector <ll> b, int bound
      = N) {
  int m = a.size() + b.size() - 1, n = 1;
  while (n < m) n <<= 1;</pre>
  a.resize(n), b.resize(n);
  ntt(a), ntt(b);
vector <ll> out(n);
  for (int i = 0; i < n; ++i) out[i] = mul(a[i], b[i]);</pre>
  ntt(out, true), out.resize(min(m, bound));
  return out;
vector <ll> Inverse(vector <ll> a) {
  // O(NlogN), a[0] != 0
  int n = a.size();
  vector <1l> res(1, mpow(a[0], mod - 2));
  for (int m = 1; m < n; m <<= 1) {</pre>
    if (n < m * 2) a.resize(m * 2);</pre>
    vector \langle 11 \rangle v1(a.begin(), a.begin() + m * 2), v2 =
        res;
    v1.resize(m * 4), v2.resize(m * 4);
    ntt(v1), ntt(v2);
    for (int i = 0; i < m * 4; ++i) v1[i] = mul(mul(v1[</pre>
         i], v2[i]), v2[i]);
    ntt(v1, true);
    res.resize(m * 2);
    for (int i = 0; i < m; ++i) res[i] = add(res[i],</pre>
         res[i]);
    for (int i = 0; i < m * 2; ++i) res[i] = sub(res[i</pre>
         ], v1[i]);
  res.resize(n);
  return res;
pair <vector <ll>, vector <ll>> Divide(vector <ll> a,
    vector <ll> b) {
  // a = bQ + R, O(NlogN), b.back() != 0
  int n = a.size(), m = b.size(), k = n - m + 1;
  if (n < m) return {{0}, a};</pre>
  vector <11> ra = a, rb = b;
  reverse(all(ra)), ra.resize(k);
  reverse(all(rb)), rb.resize(k);
  vector <11> Q = Mul(ra, Inverse(rb), k);
```

```
reverse(all(Q));
  vector \langle 11 \rangle res = Mul(b, Q), R(m - 1);
  for (int i = 0; i < m - 1; ++i) R[i] = sub(a[i], res[</pre>
      il);
  return {Q, R};
vector <1l> SqrtImpl(vector <1l> a) {
  if (a.empty()) return {0};
  int z = QuadraticResidue(a[0], mod), n = a.size();
  if (z == -1) return {-1};
  vector <ll> q(1, z);
  const int inv2 = (mod + 1) / 2;
  for (int m = 1; m < n; m <<= 1) {</pre>
    if (n < m * 2) a.resize(m * 2);</pre>
    q.resize(m * 2);
    vector <ll> f2 = Mul(q, q, m * 2);
for (int i = 0; i < m * 2; ++i) f2[i] = sub(f2[i],</pre>
         a[i]);
    f2 = Mul(f2, Inverse(q), m * 2);
    for (int i = 0; i < m * 2; ++i) q[i] = sub(q[i],</pre>
        mul(f2[i], inv2));
  }
  q.resize(n);
  return q;
vector <11> Sqrt(vector <11> a) {
  // O(NlogN), return {-1} if not exists
  int n = a.size(), m = 0;
  while (m < n && a[m] == 0) m++;</pre>
  if (m == n) return vector <11>(n);
  if (m & 1) return {-1};
  vector <ll> s = SqrtImpl(vector <ll>(a.begin() + m, a
      .end()));
  if (s[0] == -1) return {-1};
  vector <1l> res(n);
  for (int i = 0; i < s.size(); ++i) res[i + m / 2] = s</pre>
      [i];
  return res;
vector <1l> Derivative(vector <1l> a) {
 int n = a.size();
  vector <ll> res(n - 1);
  for (int i = 0; i < n - 1; ++i) res[i] = mul(a[i +</pre>
      1], i + 1);
  return res;
vector <ll> Integral(vector <ll> a) {
  int n = a.size();
  vector <ll> res(n + 1);
  for (int i = 0; i < n; ++i) {</pre>
    res[i + 1] = mul(a[i], mpow(i + 1, mod - 2));
  }
  return res;
vector <ll> Ln(vector <ll> a) {
  // O(NlogN), a[0] = 1
  int n = a.size();
  if (n == 1) return {0};
  vector <ll> d = Derivative(a);
  a.pop_back();
  return Integral(Mul(d, Inverse(a), n - 1));
vector <ll> Exp(vector <ll> a) {
  // O(NlogN), a[0] = 0
  int n = a.size();
  vector \langle 11 \rangle q(1, 1);
  a[0] = add(a[0], 1);
  for (int m = 1; m < n; m <<= 1) {</pre>
    if (n < m * 2) a.resize(m * 2);
    vector <ll> g(a.begin(), a.begin() + m * 2), h(all()
        a));
    h.resize(m * 2), h = Ln(h);
    for (int i = 0; i < m * 2; ++i) {</pre>
      g[i] = sub(g[i], h[i]);
    q = Mul(g, q, m * 2);
  }
  q.resize(n);
  return q;
vector <ll> Pow(vector <ll> a, ll k) {
 int n = a.size(), m = 0;
```

```
vector <11> ans(n, 0);
  while (m < n && a[m] == 0) m++;</pre>
  if (k && m && (k >= n || k * m >= n)) return ans;
  if (m == n) return ans[0] = 1, ans;
  ll lead = m * k;
  vector <11> b(a.begin() + m, a.end());
  ll base = mpow(b[0], k), inv = mpow(b[0], mod - 2);
  for (int i = 0; i < n - m; ++i) b[i] = mul(b[i], inv)</pre>
  b = Ln(b);
  for (int i = 0; i < n - m; ++i) b[i] = mul(b[i], k %</pre>
      mod);
  b = Exp(b);
  for (int i = lead; i < n; ++i) ans[i] = mul(b[i -</pre>
      lead], base);
vector <ll> Evaluate(vector <ll> a, vector <ll> x) {
  if (x.empty()) return {};
  int n = x.size();
  vector <vector <11>> up(n * 2);
  for (int i = 0; i < n; ++i) up[i + n] = {sub(0, x[i])
       , 1};
  for (int i = n - 1; i > 0; --i) up[i] = Mul(up[i *
      2], up[i * 2 + 1]);
  vector <vector <11>> down(n * 2);
  down[1] = Divide(a, up[1]).second;
  for (int i = 2; i < n * 2; ++i) down[i] = Divide(down</pre>
       [i >> 1], up[i]).second;
  vector <11> y(n);
  for (int i = 0; i < n; ++i) y[i] = down[i + n][0];</pre>
  return y;
vector <ll> Interpolate(vector <ll> x, vector <ll> y) {
  int n = x.size();
  vector <vector <11>> up(n * 2);
  for (int i = 0; i < n; ++i) up[i + n] = {sub(0, x[i])}
       , 1};
  for (int i = n - 1; i > 0; --i) up[i] = Mul(up[i *
      2], up[i * 2 + 1]);
  vector <ll> a = Evaluate(Derivative(up[1]), x);
  for (int i = 0; i < n; ++i) {</pre>
    a[i] = mul(y[i], mpow(a[i], mod - 2));
  vector <vector <11>> down(n * 2);
  for (int i = 0; i < n; ++i) down[i + n] = {a[i]};
for (int i = n - 1; i > 0; --i) {
  vector <ll> lhs = Mul(down[i * 2], up[i * 2 + 1]);
    vector <ll> rhs = Mul(down[i * 2 + 1], up[i * 2]);
    down[i].resize(lhs.size());
    for (int j = 0; j < lhs.size(); ++j) {</pre>
      down[i][j] = add(lhs[j], rhs[j]);
  return down[1];
```

7.4 Fast Linear Recursion

```
11 FastLinearRecursion(vector <11> a, vector <11> c, 11
  // a_n = sigma c_j * a_{n - j - 1}, 0-based
// O(NLogNLogK), |a| = |c|
  int n = a.size();
  if (k < n) return a[k];</pre>
  vector \langle 11 \rangle base(n + 1, 1);
  for (int i = 0; i < n; ++i) base[i] = sub(0, c[n - i
       - 1]);
  vector <1l> poly(n);
  (n == 1 ? poly[0] = c[n - 1] : poly[1] = 1);
  auto calc = [&](vector <ll> p1, vector <ll> p2) {
    // O(n^2) bruteforce or O(nlogn) NTT
    return Divide(Mul(p1, p2), base).second;
  };
  vector \langle 11 \rangle res(n, 0); res[0] = 1;
  for (; k; k >>= 1, poly = calc(poly, poly)) {
    if (k & 1) res = calc(res, poly);
  ll ans = 0;
  for (int i = 0; i < n; ++i) {
    (ans += res[i] * a[i]) %= mod;
```

```
7.5 Fast Walsh Transform
```

8 Geometry

8.1 Basic

return ans:

```
const double eps = 1e-8, pi = acos(-1);
int sign(double x) \{return abs(x) \leftarrow eps ? 0 : (x > 0 ?
    1 : -1);}
struct Pt {
  double x, y;
 Pt (double _x, double _y) : x(_x), y(_y) {}
 Pt operator + (Pt o) {return Pt(x + o.x, y + o.y);}
 Pt operator - (Pt o) {return Pt(x - o.x, y - o.y);}
 Pt operator * (double k) {return Pt(x * k, y * k);}
 Pt operator / (double k) {return Pt (x / k, y / k);}
 double operator * (Pt o) {return x * o.x + y * o.y;}
 double operator ^ (Pt o) {return x * o.y - y * o.x;}
struct Line {
 Pt a, b;
struct Cir {
 Pt o; double r;
double abs2(Pt o) {return o.x * o.x + o.y * o.y;}
double abs(Pt o) {return sqrt(abs2(o));}
int ori(Pt o, Pt a, Pt b) {return sign((o - a) ^ (o - b
    ));}
bool btw(Pt a, Pt b, Pt c) { // c on segment ab?
  return ori(a, b, c) == 0 && sign((c - a) * (c - b))
double area(Pt a, Pt b, Pt c) {return abs((a - b) ^ (a
    - c)) / 2;}
Pt unit(Pt o) {return o / abs(o);}
Pt rot(Pt a, double o) { // CCW
 double c = cos(o), s = sin(o);
  return Pt(c * a.x - s * a.y, s * a.x + c * a.y);
Pt proj_vector(Pt a, Pt b, Pt c) { // vector ac proj to
    ab
  return (b - a) * ((c - a) * (b - a)) / ((b - a) * (b
      - a));
Pt proj_pt(Pt a, Pt b, Pt c) { // point c proj to ab
 return proj_vector(a, b, c) + a;
```

8.2 Heart

8.3 External Bisector

```
Pt external_bisector(Pt p1, Pt p2, Pt p3) { //213
Pt L1 = p2 - p1, L2 = p3 - p1;
L2 = L2 * abs(L1) / abs(L2);
return L1 + L2;
}
```

8.4 Intersection of Segments

```
Pt LinesInter(Line a, Line b) {
    double abc = (a.b - a.a) ^ (b.a - a.a);
    double abd = (a.b - a.a) ^ (b.b - a.a);
    if (sign(abc - abd) == 0) return b.b;// no inter
    return (b.b * abc - b.a * abd) / (abc - abd);
}

vector<Pt> SegsInter(Line a, Line b) {
    if (btw(a.a, a.b, b.a)) return {b.a};
    if (btw(a.a, a.b, b.b)) return {b.b};
    if (btw(b.a, b.b, a.a)) return {a.a};
    if (btw(b.a, b.b, a.b)) return {a.b};
    if (ori(a.a, a.b, b.a) * ori(a.a, a.b, b.b) == -1 &&
        ori(b.a, b.b, a.a) * ori(b.a, b.b, a.b) == -1)
    return {LinesInter(a, b)};
    return {};
}
```

8.5 Intersection of Circle and Line

```
vector<Pt> CircleLineInter(Cir c, Line 1) {
  Pt p = l.a + (l.b - l.a) * ((c.o - l.a) * (l.b - l.a)
        ) / abs2(l.b - l.a);
  double s = (l.b - l.a) ^ (c.o - l.a), h2 = c.r * c.r
        - s * s / abs2(l.b - l.a);
  if (sign(h2) == -1) return {};
  if (sign(h2) == 0) return {p};
  Pt h = (l.b - l.a) / abs(l.b - l.a) * sqrt(h2);
  return {p - h, p + h};
}
```

8.6 Intersection of Circles

8.7 Intersection of Polygon and Circle

```
double _area(Pt pa, Pt pb, double r){
   if(abs(pa) < abs(pb)) swap(pa, pb);
   if(abs(pb) < eps) return 0;
   double S, h, theta;
   double a = abs(pb), b = abs(pa), c = abs(pb - pa);
   double cosB = pb * (pb - pa) / a / c, B = acos(cosB);
   double cosC = (pa * pb) / a / b, C = acos(cosC);
   if (a > r) {
```

```
S = (C / 2) * r * r;
    h = a * b * sin(C) / c;
    if (h < r && B < pi / 2) S -= (acos(h / r) * r * r</pre>
        - h * sqrt(r * r - h * h));
  } else if (b > r) {
    theta = pi - B - asin(sin(B) / r * a);
    S = .5 * a * r * sin(theta) + (C - theta) / 2 * r *
  } else
    S = .5 * sin(C) * a * b;
  return S;
double area_poly_circle(vector<Pt> poly, Pt 0, double r
    ) {
  double S = 0; int n = poly.size();
  for(int i = 0; i < n; ++i)</pre>
    S += _area(poly[i] - 0, poly[(i + 1) % n] - 0, r) *
         ori(0, poly[i], poly[(i + 1) % n]);
  return fabs(S);
}
```

8.8 Tangent Lines of Circle and Point

8.9 Tangent Lines of Circles

```
vector<Line> tangent(Cir a, Cir b) {
#define Pij \
 Pt i = unit(b.o - a.o) * a.r, j = Pt(i.y, -i.x);\
  z.push_back({a.o + i, a.o + i + j});
#define deo(I,J) \
  double d = abs(a.o - b.o), e = a.r I b.r, o = acos(e
      / d);\
  Pt i = unit(b.o - a.o), j = rot(i, o), k = rot(i, -o)
 z.push_back({a.o + j * a.r, b.o J j * b.r});\
 z.push_back({a.o + k * a.r, b.o J k * b.r});
 if (a.r < b.r) swap(a, b);
  vector<Line> z;
  if (abs(a.o - b.o) + b.r < a.r) return z;</pre>
  else if (sign(abs(a.o - b.o) + b.r - a.r) == 0) { Pij
      ; }
  else {
    deo(-,+); // inter
    // outer
    if (sign(d - a.r - b.r) == 0) { Pij; }
    else if (d > a.r + b.r) { deo(+,-); }
  return z;
}
```

8.10 Point In Convex

8.11 Point Segment Distance

```
double PointSegDist(Pt q0, Pt q1, Pt p) {
  if (sign(abs(q0 - q1)) == 0) return abs(q0 - p);
  if (sign((q1 - q0) * (p - q0)) >= 0 && sign((q0 - q1)
      * (p - q1)) >= 0)
    return fabs(((q1 - q0) ^ (p - q0)) / abs(q0 - q1));
  return min(abs(p - q0), abs(p - q1));
}
```

8.12 Convex Hull

8.13 Convex Hull Distance

8.14 Minimum Enclosing Circle

```
Cir min_enclosing(vector<Pt> &p) {
  random_shuffle(p.begin(), p.end());
  double r = 0.0;
  Pt cent = p[0];
  for (int i = 1; i < p.size(); ++i) {</pre>
    if (abs2(cent - p[i]) <= r) continue;</pre>
    cent = p[i];
    r = 0.0;
    for (int j = 0; j < i; ++j) {
      if (abs2(cent - p[j]) <= r) continue;</pre>
      cent = (p[i] + p[j]) / 2;
      r = abs2(p[j] - cent);
      for (int k = 0; k < j; ++k) {
        if (abs2(cent - p[k]) <= r) continue;</pre>
        cent = circenter(p[i], p[j], p[k]);
        r = abs2(p[k] - cent);
      }
    }
  return {cent, sqrt(r)};
```

8.15 Union of Circles

```
if (z < 0) z += 2 * pi;
    double 1 = z - o, r = z + o;
    if (1 < 0) 1 += 2 * pi;
    if (r > 2 * pi) r -= 2 * pi;
    if (1 > r) res.emplace_back(1, 2 * pi), res.
       emplace_back(0, r);
    else res.emplace_back(l, r);
 return res:
double CircleUnionArea(vector<Cir> c) { // circle
    should be identical
  int n = c.size();
  double a = 0, w;
  for (int i = 0; w = 0, i < n; ++i) {
    vector<pair<double, double>> s = {{2 * pi, 9}}, z;
    for (int j = 0; j < n; ++j) if (i != j) {</pre>
     z = CoverSegment(c[i], c[j]);
      for (auto &e : z) s.push_back(e);
    sort(s.begin(), s.end());
   auto F = [&] (double t) { return c[i].r * (c[i].r *
         t + c[i].o.x * sin(t) - c[i].o.y * cos(t)); };
    for (auto &e : s) {
     if (e.first > w) a += F(e.first) - F(w);
     w = max(w, e.second);
   }
 }
  return a * 0.5;
```

8.16 Polar Angle Sort

8.17 Rotating Caliper

8.18 Rotating SweepLine

```
void RotatingSweepLine(vector <Pt> &pt) {
 int n = pt.size();
 vector <int> id(n), pos(n);
 sort(line.begin(), line.end(), [&](pair <int, int> i,
      pair <int, int> j) {
   Pt a = pt[i.second] - pt[i.first], b = pt[j.second]
        - pt[j.first];
   return (a.pos() == b.pos() ? sign(a ^ b) > 0 : a.
       pos() < b.pos());
 });
 iota(id.begin(), id.end(), 0);
 sort(id.begin(), id.end(), [&](int i, int j) {
   return (sign(pt[i].y - pt[j].y) == 0 ? pt[i].x < pt</pre>
       [j].x : pt[i].y < pt[j].y);
 for (int i = 0; i < n; ++i)</pre>
   pos[id[i]] = i;
 for (auto [i, j] : line) {
```

```
// point sort by the distance to line(i, j)
// do something.
tie(pos[i], pos[j], id[pos[i]], id[pos[j]]) =
          make_tuple(pos[j], pos[i], j, i);
}
```

8.19 Half Plane Intersection

```
vector <Pt> HalfPlaneInter(vector <pair <Pt, Pt>> vec)
   1
 11
 // first -----> second
 auto pos = [&](Pt a) {return sign(a.y) == 0 ? sign(a
      .x) < 0 : sign(a.y) > 0;};
  sort(all(vec), [&](pair <Pt, Pt> a, pair <Pt, Pt> b)
    Pt A = a.second - a.first, B = b.second - b.first;
    if (pos(A) == pos(B)) {
     if (sign(A ^ B) == 0) return sign((b.first - a.
          first) * (b.second - a.first)) > 0;
     return sign(A ^ B) > 0;
    return pos(A) < pos(B);</pre>
 }):
 deque <Pt> inter;
 deque <pair <Pt, Pt>> seg;
 int n = vec.size();
 auto get = [&](pair <Pt, Pt> a, pair <Pt, Pt> b) {
      return intersect(a.first, a.second, b.first, b.
      second);};
 for (int i = 0; i < n; ++i) if (!i || vec[i] != vec[i</pre>
       - 1]) {
    while (seg.size() >= 2 && sign((vec[i].second -
        inter.back()) ^ (vec[i].first - inter.back()))
        == 1) seg.pop_back(), inter.pop_back();
    while (seg.size() >= 2 && sign((vec[i].second -
        inter.front()) ^ (vec[i].first - inter.front())
        ) == 1) seg.pop_front(), inter.pop_front();
    seg.push_back(vec[i]);
    if (seg.size() >= 2) inter.pb(get(seg[seg.size() -
        2], seg.back()));
 while (seg.size() >= 2 && sign((seg.front().second -
      inter.back()) ^ (seg.front().first - inter.back()
      )) == 1) seg.pop_back(), inter.pop_back();
 inter.push_back(get(seg.front(), seg.back()));
  return vector <Pt>(all(inter));
```

8.20 Minkowski Sum

8.21 Delaunay Triangulation

```
/* Delaunay Triangulation:
Given a sets of points on 2D plane, find a
triangulation such that no points will strictly
inside circumcircle of any triangle.
find: return a triangle contain given point
add_point: add a point into triangulation
A Triangle is in triangulation iff. its has_chd is 0.
Region of triangle u: iterate each u.edge[i].tri,
each points are u.p[(i+1)%3], u.p[(i+2)%3]
Voronoi diagram: for each triangle in triangulation,
the bisector of all its edges will split the region.
```

```
nearest point will belong to the triangle containing it
                                                                    tri->p[pi], trj->p[pj]);
                                                                edge(Edge(trk, 0), Edge(trl, 0));
const ll inf = MAXC * MAXC * 100; // Lower_bound
                                                                edge(Edge(trk, 1), tri->edge[(pi + 2) % 3]);
    unknown
                                                                edge(Edge(trk, 2), trj->edge[(pj + 1) \% 3]);
struct Tri;
                                                                edge(Edge(trl, 1), trj->edge[(pj + 2) % 3]);
struct Edge {
                                                                edge(Edge(trl, 2), tri->edge[(pi + 1) % 3]);
 Tri* tri; int side;
                                                                tri->chd[0] = trk; tri->chd[1] = trl; tri->chd[2] =
  Edge(): tri(0), side(0){}
 Edge(Tri* _tri, int _side): tri(_tri), side(_side){}
                                                                trj->chd[0] = trk; trj->chd[1] = trl; trj->chd[2] =
                                                                flip(trk, 1); flip(trk, 2);
struct Tri {
 pll p[3];
                                                                flip(trl, 1); flip(trl, 2);
  Edge edge[3];
                                                              }
  Tri* chd[3];
                                                            };
                                                            vector<Tri*> triang; // vector of all triangle
 Tri() {}
                                                            set<Tri*> vst;
 Tri(const pll& p0, const pll& p1, const pll& p2) {
                                                            void go(Tri* now) { // store all tri into triang
   p[0] = p0; p[1] = p1; p[2] = p2;
chd[0] = chd[1] = chd[2] = 0;
                                                              if (vst.find(now) != vst.end())
                                                                return:
 bool has_chd() const { return chd[0] != 0; }
                                                              vst.insert(now);
 int num_chd() const {
                                                              if (!now->has_chd())
                                                                return triang.pb(now);
   return !!chd[0] + !!chd[1] + !!chd[2];
                                                              for (int i = 0; i < now->num_chd(); ++i)
 bool contains(pll const& q) const {
                                                                go(now->chd[i]);
    for (int i = 0; i < 3; ++i)</pre>
                                                            void build(int n, pll* ps) { // build triangulation
      if (ori(p[i], p[(i + 1) % 3], q) < 0)</pre>
        return 0;
                                                              tris = pool; triang.clear(); vst.clear();
                                                              random_shuffle(ps, ps + n);
    return 1;
                                                              Trig tri; // the triangulation structure
                                                              for (int i = 0; i < n; ++i)</pre>
} pool[N * 10], *tris;
void edge(Edge a, Edge b) {
                                                                tri.add_point(ps[i]);
 if(a.tri) a.tri->edge[a.side] = b;
                                                              go(tri.the_root);
  if(b.tri) b.tri->edge[b.side] = a;
                                                            8.22 Triangulation Vonoroi
struct Trig { // Triangulation
 Trig() {
                                                            vector<Line> ls[N];
    the_root = // Tri should at least contain all
                                                            pll arr[N];
                                                            Line make_line(pdd p, Line 1) {
      new(tris++) Tri(pll(-inf, -inf), pll(inf + inf, -
                                                              pdd d = 1.Y - 1.X; d = perp(d);
          inf), pll(-inf, inf + inf));
                                                              pdd m = (1.X + 1.Y) / 2;
                                                              l = Line(m, m + d);
 Tri* find(pll p) { return find(the_root, p); }
                                                              if (ori(1.X, 1.Y, p) < 0)
  void add_point(const pll &p) { add_point(find(
                                                                l = Line(m + d, m);
      the_root, p), p); }
                                                              return 1:
 Tri* the_root;
  static Tri* find(Tri* root, const pll &p) {
                                                            double calc_area(int id) {
    while (1) {
                                                              // use to calculate the area of point "strictly in
      if (!root->has_chd())
                                                                  the convex hull"
        return root:
                                                              vector<Line> hpi = halfPlaneInter(ls[id]);
      for (int i = 0; i < 3 && root->chd[i]; ++i)
                                                              vector<pdd> ps;
        if (root->chd[i]->contains(p)) {
                                                              for (int i = 0; i < SZ(hpi); ++i)</pre>
          root = root->chd[i];
                                                                ps.pb(intersect(hpi[i].X, \, hpi[i].Y, \, hpi[(i\, +\, 1)\, \, \%
          break;
                                                                    SZ(hpi)].X, hpi[(i + 1) % SZ(hpi)].Y));
                                                              double rt = 0;
                                                              for (int i = 0; i < SZ(ps); ++i)</pre>
    assert(0); // "point not found"
                                                                rt += cross(ps[i], ps[(i + 1) % SZ(ps)]);
                                                              return fabs(rt) / 2;
  void add_point(Tri* root, pll const& p) {
   Tri* t[3];
                                                            void solve(int n, pii *oarr) {
     '* split it into three triangles */
                                                              map<pll, int> mp;
    for (int i = 0; i < 3; ++i)
                                                              for (int i = 0; i < n; ++i)</pre>
      t[i] = new(tris++) Tri(root->p[i], root->p[(i +
                                                                arr[i] = pll(oarr[i].X, oarr[i].Y), mp[arr[i]] = i;
          1) % 3], p);
                                                              build(n, arr); // Triangulation
    for (int i = 0; i < 3; ++i)
                                                              for (auto *t : triang) {
      edge(Edge(t[i], 0), Edge(t[(i + 1) % 3], 1));
                                                                vector<int> p;
    for (int i = 0; i < 3; ++i)
                                                                for (int i = 0; i < 3; ++i)</pre>
      edge(Edge(t[i], 2), root->edge[(i + 2) % 3]);
                                                                  if (mp.find(t->p[i]) != mp.end())
    for (int i = 0; i < 3; ++i)</pre>
                                                                    p.pb(mp[t->p[i]]);
      root->chd[i] = t[i];
                                                                for (int i = 0; i < SZ(p); ++i)</pre>
    for (int i = 0; i < 3; ++i)
                                                                  for (int j = i + 1; j < SZ(p); ++j) {
      flip(t[i], 2);
                                                                    Line l(oarr[p[i]], oarr[p[j]]);
                                                                    ls[p[i]].pb(make_line(oarr[p[i]], 1));
  void flip(Tri* tri, int pi) {
                                                                    ls[p[j]].pb(make_line(oarr[p[j]], 1));
   Tri* trj = tri->edge[pi].tri;
                                                                  }
    int pj = tri->edge[pi].side;
    if (!trj) return;
    if (!in_cc(tri->p[0], tri->p[1], tri->p[2], trj->p[
        pj])) return;
                                                                 Else
       flip edge between tri,trj */
    Tri* trk = new(tris++) Tri(tri->p[(pi + 1) % 3],
        trj->p[pj], tri->p[pi]);
                                                            9.1 Bit Hack
    Tri* trl = new(tris++) Tri(trj->p[(pj + 1) % 3],
```

9.2 Dynamic Programming Condition

9.2.1 Totally Monotone (Concave/Convex)

```
\begin{array}{l} \forall i < i', j < j' \text{, } B[i][j] \leq B[i'][j] \implies B[i][j'] \leq B[i'][j'] \\ \forall i < i', j < j' \text{, } B[i][j] \geq B[i'][j] \implies B[i][j'] \geq B[i'][j'] \end{array}
```

9.2.2 Monge Condition (Concave/Convex)

```
\begin{array}{l} \forall i < i', j < j' \text{, } B[i][j] + B[i'][j'] \geq B[i][j'] + B[i'][j] \\ \forall i < i', j < j' \text{, } B[i][j] + B[i'][j'] \leq B[i][j'] + B[i'][j] \end{array}
```

9.2.3 Optimal Split Point

```
If B[i][j] + B[i+1][j+1] \geq B[i][j+1] + B[i+1][j] then H_{i,j-1} \leq H_{i,j} \leq H_{i+1,j}
```

9.3 Slope Trick

```
template<typename T>
struct slope_trick_convex {
 T minn = 0, ground_1 = 0, ground_r = 0;
 priority_queue<T, vector<T>, less<T>> left;
priority_queue<T, vector<T>, greater<T>> right;
  slope_trick_convex() {left.push(numeric_limits<T>::
      min() / 2), right.push(numeric_limits<T>::max() /
  void push_left(T x) {left.push(x - ground_l);}
  void push_right(T x) {right.push(x - ground_r);}
  //add a line with slope 1 to the right starting from
  void add_right(T x) {
    T l = left.top() + ground_l;
   if (1 <= x) push_right(x);</pre>
    else push_left(x), push_right(1), left.pop(), minn
        += 1 - x;
 //add a line with slope -1 to the left starting from
  void add_left(T x) {
    T r = right.top() + ground_r;
    if (r >= x) push_left(x);
    else push_right(x), push_left(r), right.pop(), minn
  //val[i]=min(val[j]) for all i-l<=j<=i+r
  void expand(T 1, T r) {ground_1 -= 1, ground_r += r;}
  void shift_up(T x) {minn += x;}
  T get_val(T x) {
    T l = left.top() + ground_l, r = right.top() +
        ground_r;
    if (x >= 1 \&\& x <= r) return minn;
    if (x < 1) {
      vector<T> trash;
      T cur_val = minn, slope = 1, res;
      while (1) {
        trash.push_back(left.top());
        left.pop();
        if (left.top() + ground_l <= x) {</pre>
          res = cur_val + slope * (1 - x);
          break:
        cur_val += slope * (1 - (left.top() + ground_1)
        1 = left.top() + ground_l;
        slope += 1;
      for (auto i : trash) left.push(i);
      return res;
    if(x > r) {
      vector<T> trash;
```

```
T cur_val = minn, slope = 1, res;
while (1) {
    trash.push_back(right.top());
    right.pop();
    if (right.top() + ground_r >= x) {
        res = cur_val + slope * (x - r);
        break;
    }
    cur_val += slope * ((right.top() + ground_r) -
        r);
    r = right.top() + ground_r;
    slope += 1;
    }
    for (auto i : trash) right.push(i);
    return res;
    }
    assert(0);
}
```

9.4 Manhattan MST

```
void solve(int n) {
  init();
  vector<int> v(n), ds;
  for (int i = 0; i < n; ++i) {</pre>
    v[i] = i;
    ds.push_back(x[i] - y[i]);
  sort(ds.begin(), ds.end());
  ds.resize(unique(ds.begin(), ds.end()) - ds.begin());
  sort(v.begin(), v.end(), [&](int i, int j) { return x
      [i] == x[j] ? y[i] > y[j] : x[i] > x[j]; });
  int j = 0;
  for (int i = 0; i < n; ++i) {</pre>
    int p = lower_bound(ds.begin(), ds.end(), x[v[i]] -
         y[v[i]]) - ds.begin() + 1;
    pair<int, int> q = query(p);
    // query return prefix minimum
    if (~q.second) add_edge(v[i], q.second);
    add(p, make_pair(x[v[i]] + y[v[i]], v[i]));
void make_graph() {
  solve(n);
  for (int i = 0; i < n; ++i) swap(x[i], y[i]);</pre>
  solve(n);
  for (int i = 0; i < n; ++i) x[i] = -x[i];
  solve(n);
  for (int i = 0; i < n; ++i) swap(x[i], y[i]);</pre>
  solve(n);
```

9.5 Dynamic MST

```
int cnt[maxn], cost[maxn], st[maxn], ed[maxn];
pair<int, int> qr[maxn];
// qr[i].first = id of edge to be changed, qr[i].second
     = weight after operation
// cnt[i] = number of operation on edge i
// call solve(0, q - 1, v, 0), where v contains edges i
     such that cnt[i] == 0
void contract(int 1, int r, vector<int> v, vector<int>
    &x, vector<int> &y) {
  sort(v.begin(), v.end(), [&](int i, int j) {
      if (cost[i] == cost[j]) return i < j;</pre>
      return cost[i] < cost[j];</pre>
      });
  djs.save();
  for (int i = 1; i <= r; ++i) djs.merge(st[qr[i].first</pre>
      ], ed[qr[i].first]);
  for (int i = 0; i < (int)v.size(); ++i) {</pre>
    if (djs.find(st[v[i]]) != djs.find(ed[v[i]])) {
      x.push_back(v[i]);
      djs.merge(st[v[i]], ed[v[i]]);
    }
 dis.undo():
  djs.save();
  for (int i = 0; i < (int)x.size(); ++i) djs.merge(st[</pre>
      x[i]], ed[x[i]]);
```

```
for (int i = 0; i < (int)v.size(); ++i) {</pre>
    if (djs.find(st[v[i]]) != djs.find(ed[v[i]])) {
      y.push_back(v[i]);
      djs.merge(st[v[i]], ed[v[i]]);
  djs.undo();
void solve(int 1, int r, vector<int> v, long long c) {
  if (1 == r) {
    cost[qr[1].first] = qr[1].second;
    if (st[qr[1].first] == ed[qr[1].first]) {
      printf("%lld\n", c);
      return;
    int minv = qr[1].second;
    for (int i = 0; i < (int)v.size(); ++i) minv = min(</pre>
        minv, cost[v[i]]);
    printf("%lld\n", c + minv);
    return;
  int m = (1 + r) >> 1;
  vector<int> lv = v, rv = v;
  vector<int> x, y;
  for (int i = m + 1; i <= r; ++i) {</pre>
    cnt[qr[i].first]--;
    if (cnt[qr[i].first] == 0) lv.push_back(qr[i].first
  contract(1, m, lv, x, y);
  long long lc = c, rc = c;
  djs.save();
  for (int i = 0; i < (int)x.size(); ++i) {</pre>
   lc += cost[x[i]];
    \label{eq:djs.merge} \mbox{djs.merge(st[x[i]], ed[x[i]]);}
  solve(1, m, y, 1c);
  djs.undo();
  x.clear(), y.clear();
  for (int i = m + 1; i <= r; ++i) cnt[qr[i].first]++;</pre>
  for (int i = 1; i <= m; ++i) {</pre>
    cnt[qr[i].first]--;
    if (cnt[qr[i].first] == 0) rv.push_back(qr[i].first
  contract(m + 1, r, rv, x, y);
  djs.save();
  for (int i = 0; i < (int)x.size(); ++i) {</pre>
   rc += cost[x[i]];
    djs.merge(st[x[i]], ed[x[i]]);
  solve(m + 1, r, y, rc);
  djs.undo();
  for (int i = 1; i <= m; ++i) cnt[qr[i].first]++;</pre>
```

9.6 ALL LCS

```
void all_lcs(string s, string t) { // 0-base
  vector<int> h(t.size());
  iota(all(h), 0);
  for (int a = 0; a < s.size(); ++a) {
    int v = -1;
    for (int c = 0; c < t.size(); ++c)
        if (s[a] == t[c] || h[c] < v)
            swap(h[c], v);
        // LCS(s[0, a], t[b, c]) =
        // c - b + 1 - sum([h[i] >= b] | i <= c)
        // h[i] might become -1 !!
  }
}</pre>
```

9.7 Hilbert Curve

```
long long hilbertOrder(int x, int y, int pow, int
    rotate) {
    if (pow == 0) return 0;
    int hpow = 1 << (pow-1);
    int seg = (x < hpow) ? ((y < hpow) ? 0 : 3) : ((y < hpow) ? 1 : 2);
    seg = (seg + rotate) & 3;</pre>
```

9.8 Pbds

```
#include <ext/pb_ds/priority_queue.hpp>
#include <ext/pb_ds/assoc_container.hpp>
using namespace __gnu_pbds;
#include <ext/rope>
using namespace __gnu_cxx;
int main () {
     _gnu_pbds::priority_queue <int> pq1, pq2;
  pq1.join(pq2); // pq1 += pq2, pq2 = {}
  cc_hash_table<int, int> m1;
  tree<int, null_type, less<int>, rb_tree_tag,
      tree_order_statistics_node_update> oset;
  oset.insert(2), oset.insert(4);
  cout << *oset.find_by_order(1) << ' ' << oset.</pre>
      order_of_key(1) << '\n'; // 4 0
  bitset <100> BS;
  BS.flip(3), BS.flip(5);
  cout << BS._Find_first() << ' ' << BS._Find_next(3)</pre>
      << '\n'; // 3 5
  rope <int> rp1, rp2;
  rp1.push_back(1), rp1.push_back(3);
  rp1.insert(0, 2); // pos, num
  rp1.erase(0, 2); // pos, len
  rp1.substr(0, 2); // pos, Len
  rp2.push_back(4);
  rp1 += rp2, rp2 = rp1;
  cout << rp2[0] << ' ' << rp2[1] << '\n'; // 3 4
```

9.9 Random

```
struct custom_hash {
    static uint64_t splitmix64(uint64_t x) {
        x += 0x9e3779b97f4a7c15;
        x = (x ^ (x >> 30)) * 0xbf58476d1ce4e5b9;
        x = (x ^ (x >> 27)) * 0x94d049bb133111eb;
        return x ^ (x >> 31);
    }
    size_t operator()(uint64_t a) const {
        static const uint64_t FIXED_RANDOM = chrono::
            steady_clock::now().time_since_epoch().count();
        return splitmix64(i + FIXED_RANDOM);
    }
};
unordered_map <int, int, custom_hash> m1;
random_device rd; mt19937 rng(rd());
```

9.10 Smawk Algorithm

```
11 query(int 1, int r) {
struct SMAWK {
  // Condition:
  // If M[1][0] < M[1][1] then M[0][0] < M[0][1]
  // If M[1][0] == M[1][1] then M[0][0] <= M[0][1]
  // For all i, find r_i s.t. M[i][r_i] is maximum ||
      minimum.
  int ans[N], tmp[N];
  void interpolate(vector <int> 1, vector <int> r) {
    int n = 1.size(), m = r.size();
    vector <int> nl;
    for (int i = 1; i < n; i += 2) {</pre>
      nl.push_back(l[i]);
    run(nl, r);
    for (int i = 1, j = 0; i < n; i += 2) {
      while (j < m && r[j] < ans[l[i]])</pre>
        j++;
      assert(j < m \&\& ans[l[i]] == r[j]);
```

```
tmp[l[i]] = j;
    for (int i = 0; i < n; i += 2) {
      int curl = 0, curr = m - 1;
      if (i)
        curl = tmp[l[i - 1]];
      if (i + 1 < n)
        curr = tmp[l[i + 1]];
      11 res = query(l[i], r[curl]);
       ans[l[i]] = r[curl];
      for (int j = curl + 1; j <= curr; ++j) {
    ll nxt = query(l[i], r[j]);</pre>
         if (res < nxt)</pre>
           res = nxt, ans[l[i]] = r[j];
      }
    }
  }
  void reduce(vector <int> 1, vector <int> r) {
    int n = 1.size(), m = r.size();
    vector <int> nr;
    for (int j : r) {
      while (!nr.empty()) {
         int i = nr.size() - 1;
         if (query(l[i], nr.back()) <= query(l[i], j))</pre>
          nr.pop_back();
         else
           break;
       if (nr.size() < n)</pre>
         nr.push_back(j);
    run(1, nr);
  void run(vector <int> 1, vector <int> r) {
    int n = l.size(), m = r.size();
    if (max(n, m) <= 2) {</pre>
       for (int i : 1) {
        ans[i] = r[0];
         if (m > 1)
           if (query(i, r[0]) < query(i, r[1]))</pre>
             ans[i] = r[1];
    } else if (n >= m) {
      interpolate(1, r);
    } else {
      reduce(1, r);
  }
};
```

9.11 Two Dimension Add Sum*

```
struct TwoDimensionAddAndSum {
 // 0-index, [l, r)
  struct Seg {
    int 1, r, m;
    ll vala, valb, lza, lzb;
    Seg* ch[2];
    Seg (int _l, int _r) : l(_l), r(_r), m(l + r >> 1),
         vala(0), valb(0), lza(0), lzb(0) {
      if (r - 1 > 1) {
        ch[0] = new Seg(1, m);
        ch[1] = new Seg(m, r);
     }
    void pull() {vala = ch[0]->vala + ch[1]->vala, valb
         = ch[0]->valb + ch[1]->valb;}
    void give(ll a, ll b) {
     lza += a, lzb += b;
     vala += a * (r - 1), valb += b * (r - 1);
    void push() {
     ch[0]->give(lza, lzb), ch[1]->give(lza, lzb), lza
           = 1zb = 0;
    void add(int a, int b, ll va, ll vb) {
     if (a <= 1 && r <= b)
        give(va, vb);
      else {
        push();
        if (a < m) ch[0]->add(a, b, va, vb);
```

```
if (m < b) ch[1]->add(a, b, va, vb);
        pull();
      }
     long long query(int a, int b, int v) {
      if (a <= 1 && r <= b) return vala * v + valb;</pre>
       push();
       long long ans = 0;
       if (a < m) ans += ch[0]->query(a, b, v);
       if (m < b) ans += ch[1]->query(a, b, v);
       return ans;
    }
  };
  // note integer overflow.
  vector <array <int, 4>> E[N];
  vector <array <int, 4>> Q[N];
  vector <11> ans;
  void add_event(int x1, int y1, int x2, int y2, l1 v)
     E[x1].pb({y1, y2, v, -v * x1});
    E[x2].pb({y1, y2, -v, v * x2});
  void add_query(int x1, int y1, int x2, int y2, int id
     Q[x1].pb({y1, y2, -1, id});
     Q[x2].pb({y1, y2, 1, id});
     ans.pb(0);
  void solve(int n) {
    Seg root(0, n);
     for (int i = 0; i <= n; ++i) {</pre>
       for (auto j : E[i]) root.add(j[0], j[1], j[2], j
           [3]);
       for (auto j : Q[i]) ans[j[3]] += j[2] * root.
           query(j[0], j[1], i);
    }
  }
};
```

9.12 Matroid Intersection

Start from $S=\emptyset$. In each iteration, let

```
• Y_1 = \{x \notin S \mid S \cup \{x\} \in I_1\}
• Y_2 = \{x \notin S \mid S \cup \{x\} \in I_2\}
```

If there exists $x\in Y_1\cap Y_2$, insert x into S . Otherwise for each $x\in S, y\not\in S$, create edges

```
• x \to y if S - \{x\} \cup \{y\} \in I_1.
• y \to x if S - \{x\} \cup \{y\} \in I_2.
```

Find a shortest path (with BFS) starting from a vertex in Y_1 and ending at a vertex in Y_2 which doesn't pass through any other vertices in Y_2 , and alternate the path. The size of S will be incremented by 1 in each iteration. For the weighted case, assign weight w(x) to vertex x if $x \in S$ and -w(x) if $x \notin S$. Find the path with the minimum number of edges among all minimum length paths and alternate it.

9.13 Python Misc