Contents

1	Basi	ic										
	1.1	Compiler Shell										
	1.2	Create File										
	1.3	Default Code										
		Testing Todo List										
	1.5	Debug Macro										
	1.6	Stress Test Shell										
	1.7	Pragma										
	1.8	Fast IO										
	1.9	Divide			•	•		•			•	
_												
2		Structure										
		Leftist Tree										
		Splay Tree										
		Link Cut Tree										
		Treap										
		Persistent Segment Tree 2D Segment Tree										
		Zkw										
		Chtholly Tree										
		Incremental Min Sum										
	2.,,	The emerical right sum	•	• •	•	•	•	•	•	•	•	 •
3	Flov	/ Matching										
	3.1	Dinic										
	3.2	Min Cost Max Flow										
	3.3	Kuhn Munkres										
		SW Min Cut										
	3.5	Gomory Hu Tree										
		Blossom										
		Weighted Blossom										
	3.8	Flow Model	•		•	٠	•	•	•	•	•	 ٠
,	Cnar	s h										
4	Grap	Heavy-Light Decomposition										
		Centroid Decomposition										
		Edge BCC										
		Block Cut Tree										
		SCC / 2SAT										
		Negative Cycle										
		Virtual Tree										
		Directed MST										
	4.9	Dominator Tree										
5	Stri											
		Aho-Corasick Automaton										
		KMP Algorithm										
		Z Algorithm										
		Manacher			:						•	 •
		SAIS			:						•	 •
		JA13										•
										•	•	•
		Suffix Automaton										
	5.8	Suffix Automaton										
	5.8 5.9	Suffix Automaton	:						:			
	5.8 5.9	Suffix Automaton	:						:			
6	5.8 5.9 5.10 Math	Suffix Automaton			:	:	:	•			•	 •
6	5.8 5.9 5.10 Math 6.1	Suffix Automaton									•	
6	5.8 5.9 5.16 Math 6.1 6.2	Suffix Automaton										
6	5.8 5.9 5.16 Math 6.1 6.2 6.3	Suffix Automaton										 :
6	5.8 5.9 5.16 Math 6.1 6.2 6.3 6.4	Suffix Automaton										
6	5.8 5.9 5.16 Math 6.1 6.2 6.3 6.4 6.5	Suffix Automaton										
6	5.8 5.9 5.16 Math 6.1 6.2 6.3 6.4 6.5 6.6	Suffix Automaton Minimum Rotation Palindrome Tree										
6	5.8 5.9 5.16 Math 6.1 6.2 6.3 6.4 6.5 6.6 6.7	Suffix Automaton Minimum Rotation Palindrome Tree										
6	5.8 5.9 5.16 Math 6.1 6.2 6.3 6.4 6.5 6.6 6.7	Suffix Automaton Minimum Rotation Palindrome Tree										
6	5.8 5.9 5.16 6.1 6.2 6.3 6.4 6.5 6.6 6.7 6.8 6.9	Suffix Automaton Minimum Rotation Palindrome Tree Main Lorentz Fraction Miller Rabin / Pollard Rho Ext GCD PiCount Linear Function Mod Min Floor Sum Quadratic Residue Simplex Linear Programming Construction										
6	5.8 5.9 5.16 6.1 6.2 6.3 6.4 6.5 6.6 6.7 6.8 6.9 6.16	Suffix Automaton Minimum Rotation Palindrome Tree Main Lorentz Fraction Miller Rabin / Pollard Rho Ext GCD PiCount Linear Function Mod Min Floor Sum Quadratic Residue Simplex Linear Programming Construction Euclidean										
6	5.8 5.9 5.16 6.1 6.2 6.3 6.4 6.5 6.6 6.7 6.8 6.9 6.16	Suffix Automaton Minimum Rotation Palindrome Tree Main Lorentz Fraction Miller Rabin / Pollard Rho Ext GCD PiCount Linear Function Mod Min Floor Sum Quadratic Residue Simplex Linear Programming Construction										
6	5.8 5.9 5.16 6.1 6.2 6.3 6.4 6.5 6.6 6.7 6.8 6.9 6.10 6.11	Suffix Automaton Minimum Rotation Palindrome Tree Main Lorentz Fraction Miller Rabin / Pollard Rho Ext GCD PiCount Linear Function Mod Min Floor Sum Quadratic Residue Simplex Linear Programming Construction Euclidean										
	5.8 5.9 5.16 6.1 6.2 6.3 6.4 6.5 6.6 6.7 6.8 6.9 6.10 6.11	Suffix Automaton Minimum Rotation Palindrome Tree Main Lorentz Fraction Miller Rabin / Pollard Rho Ext GCD PiCount Linear Function Mod Min Floor Sum Quadratic Residue Simplex Linear Programming Construction Euclidean Theorem										
	5.8 5.9 5.16 6.1 6.2 6.3 6.4 6.5 6.6 6.7 6.8 6.9 6.10 6.11 Geom	Suffix Automaton Minimum Rotation Palindrome Tree Main Lorentz Fraction Miller Rabin / Pollard Rho Ext GCD PiCount Linear Function Mod Min Floor Sum Quadratic Residue Simplex Linear Programming Construction Euclidean LTheorem										
	5.8 5.9 5.16 6.1 6.2 6.3 6.4 6.5 6.6 6.7 6.8 6.9 6.10 6.11 7.1	Suffix Automaton Minimum Rotation Palindrome Tree Main Lorentz Main Lorentz Fraction Miller Rabin / Pollard Rho Ext GCD PiCount Linear Function Mod Min Floor Sum Quadratic Residue Simplex Linear Programming Construction Buclidean LTheorem Metry Basic										
	5.8 5.9 5.10 6.1 6.2 6.3 6.4 6.5 6.6 6.7 6.8 6.9 6.10 6.11 7.1 7.2 7.3 7.4	Suffix Automaton Minimum Rotation Palindrome Tree Main Lorentz Fraction Fraction Miller Rabin / Pollard Rho Ext GCD PiCount Linear Function Mod Min Floor Sum Quadratic Residue Simplex Linear Programming Construction Euclidean Unterprogramming Extery Basic Heart External Bisector Intersection of Segments										
	5.8 5.9 5.10 6.1 6.2 6.3 6.4 6.5 6.6 6.7 6.8 6.9 6.10 7.1 7.2 7.3 7.4	Suffix Automaton Minimum Rotation Palindrome Tree Main Lorentz Fraction Fraction Miller Rabin / Pollard Rho Ext GCD PiCount Linear Function Mod Min Floor Sum Quadratic Residue Simplex Linear Programming Construction Euclidean Theorem Metry Basic Heart External Bisector Intersection of Segments Intersection of Circle and Line										
	5.8 5.9 5.10 6.1 6.2 6.3 6.4 6.5 6.6 6.7 6.8 6.9 6.10 6.11 Geom 7.1 7.2 7.3 7.4 7.5 7.6	Suffix Automaton Minimum Rotation Palindrome Tree Main Lorentz Fraction Miller Rabin / Pollard Rho Ext GCD PiCount Linear Function Mod Min Floor Sum Quadratic Residue Simplex Linear Programming Construction Deuclidean Theorem Metry Basic Heart External Bisector Intersection of Segments Intersection of Circle and Line Intersection of Circles Minimum Rotation										
	5.8 5.9 5.10 6.1 6.2 6.3 6.4 6.5 6.6 6.7 7.1 7.2 7.3 7.4 7.5 7.6 7.7	Suffix Automaton Minimum Rotation Palindrome Tree Main Lorentz Fraction Miller Rabin / Pollard Rho Ext GCD PiCount Linear Function Mod Min Floor Sum Quadratic Residue Simplex Linear Programming Construction Buclidean Theorem Buch Simplex Linear Function Mod Min Linear Programming Construction Floor Sum Construction Floor Sum Linear Programming Construction Floor Simplex Linear Programmi	· · · · · · · · · · · · · · · · e									
	5.8 5.9 5.16 Math 6.2 6.3 6.4 6.5 6.6 6.7 6.8 6.9 6.10 7.1 7.2 7.3 7.4 7.7 7.6 7.7	Suffix Automaton Minimum Rotation Palindrome Tree Main Lorentz Fraction Miller Rabin / Pollard Rho Ext GCD PiCount Linear Function Mod Min Floor Sum Quadratic Residue Simplex Linear Programming Construction Buclidean Theorem Buclidean Theorem Buch Construction Floor Sum Con	· · · · · · · · · · · · · · · · · e									
	5.8 5.9 5.16 Math 6.1 6.2 6.3 6.4 6.5 6.6 6.7 6.8 6.9 6.10 7.1 7.2 7.3 7.4 7.5 7.7 7.8 7.9	Suffix Automaton Minimum Rotation Palindrome Tree Main Lorentz Fraction Fraction Miller Rabin / Pollard Rho Ext GCD PiCount Linear Function Mod Min Floor Sum Quadratic Residue Simplex Linear Programming Construction Deuclidean Theorem Extery Basic Heart External Bisector Intersection of Circle and Line Intersection of Polygon and Circl Tangent Lines of Circle and Point Tangent Lines of Circles Intersection Tangent Lines of Circles Intergent Lines Intergent Intergent Lines Intergent In										
	5.8 5.9 5.16 6.1 6.2 6.3 6.4 6.5 6.6 6.7 7.1 7.2 7.3 7.4 7.5 7.6 7.7 7.7 7.7 7.7	Suffix Automaton Minimum Rotation Palindrome Tree Main Lorentz Fraction Fraction Miller Rabin / Pollard Rho Ext GCD PiCount Linear Function Mod Min Floor Sum Quadratic Residue Simplex Linear Programming Construction Deuclidean Theorem Universe Simple External Bisector Intersection of Segments Intersection of Circle and Line Intersection of Polygon and Circl Tangent Lines of Circle and Point Tangent Lines of Circles Opoint In Convex	· · · · · · · · · · · · · · · · · · ·									
	5.8 5.9 5.16 6.1 6.2 6.3 6.4 6.6 6.7 6.8 6.9 6.10 6.11 7.2 7.4 7.5 7.6 7.7 7.8 7.1 7.1	Suffix Automaton Minimum Rotation Palindrome Tree Main Lorentz Fraction Fraction Miller Rabin / Pollard Rho Ext GCD PiCount Linear Function Mod Min Floor Sum Quadratic Residue Simplex Linear Programming Construction Duclidean Theorem Basic Heart External Bisector Intersection of Segments Intersection of Circle and Line Intersection of Polygon and Circl Tangent Lines of Circle and Point Tangent Lines of Circles Point In Convex Point Segment Distance										
	5.8 5.9 5.16 6.1 6.2 6.3 6.4 6.5 6.6 6.7 6.8 6.9 6.10 7.1 7.2 7.3 7.4 7.7 7.8 7.9 7.1 7.1 7.1	Suffix Automaton Minimum Rotation Palindrome Tree Main Lorentz Fraction Fraction Miller Rabin / Pollard Rho Ext GCD PiCount Linear Function Mod Min Floor Sum Quadratic Residue Simplex Linear Programming Construction Duclidean Theorem External Bisector Intersection of Segments Intersection of Circle and Line Intersection of Polygon and Circl Tangent Lines of Circles Piont In Convex Piont Segment Distance										
	5.8 5.9 5.16 6.1 6.2 6.3 6.6 6.7 6.8 6.6 6.7 7.1 7.2 7.3 7.4 7.5 7.7 7.8 7.9 7.11 7.12	Suffix Automaton Minimum Rotation Palindrome Tree Main Lorentz Fraction Miller Rabin / Pollard Rho Ext GCD PiCount Linear Function Mod Min Floor Sum Quadratic Residue Simplex Linear Programming Construction Buclidean Theorem Buch dean Theorem Intersection of Segments Intersection of Circle and Line Intersection of Polygon and Circl Tangent Lines of Circles Proint In Convex Proint Segment Distance Convex Hull Convex Hull Distance										
	5.8 5.9 5.16 6.1 6.2 6.3 6.4 6.6 6.7 6.8 6.9 6.10 7.1 7.2 7.3 7.4 7.5 7.6 7.7 7.1 7.1 7.1 7.1	Suffix Automaton Minimum Rotation Palindrome Tree Main Lorentz Fraction Miller Rabin / Pollard Rho Ext GCD PiCount Linear Function Mod Min Floor Sum Quadratic Residue Simplex Linear Programming Construction Deuclidean Theorem External Bisector Intersection of Circle and Line Intersection of Circle and Line Intersection of Circle and Circl Tangent Lines of Circle and Point Tangent Lines of Circle and Point Tangent Lines of Circles Point Segment Distance Convex Hull Sconvex Hull Distance Minimum Enclosing Circle Minimum Enclosing Circle Minimum Enclosing Circle										
	5.8 5.9 5.16 6.1 6.2 6.3 6.6 6.7 6.8 6.10 6.11 7.2 7.3 7.4 7.7 7.8 7.9 7.11 7.12 7.13 7.14 7.15 7.16 7.17	Suffix Automaton Minimum Rotation Palindrome Tree Main Lorentz Fraction Miller Rabin / Pollard Rho Ext GCD PiCount Linear Function Mod Min Floor Sum Quadratic Residue Simplex Linear Programming Construction Puclidean Theorem Basic Theorem External Bisector Intersection of Segments Intersection of Circle and Line Intersection of Circle and Point Tangent Lines of Circle and Point Tangent Lines of Circle and Point Tangent Lines of Circles Point Segment Distance Convex Hull Convex Hull Convex Hull Convex Hull Distance Minimum Enclosing Circle Gunion of Circles	· · · · · · · · · · · · · · · · · · ·									
	5.8 5.9 5.16 6.1 6.2 6.3 6.6 6.7 6.8 6.10 6.11 7.2 7.3 7.4 7.7 7.8 7.9 7.11 7.12 7.13 7.14 7.15 7.16 7.17	Suffix Automaton Minimum Rotation Palindrome Tree Main Lorentz Fraction Miller Rabin / Pollard Rho Ext GCD PiCount Linear Function Mod Min Floor Sum Quadratic Residue Simplex Linear Programming Construction Puclidean Theorem Basic Theorem External Bisector Intersection of Segments Intersection of Circle and Line Intersection of Circle and Point Tangent Lines of Circle and Point Tangent Lines of Circle and Point Tangent Lines of Circles Point Segment Distance Convex Hull Convex Hull Convex Hull Convex Hull Distance Minimum Enclosing Circle Gunion of Circles	· · · · · · · · · · · · · · · · · · ·									
	5.8 5.9 5.16 6.1 6.2 6.3 6.4 6.6 6.7 6.8 6.9 6.10 6.11 7.2 7.1 7.7 7.8 7.1 7.1 7.1 7.1 7.1 7.1 7.1 7.1 7.1 7.1	Suffix Automaton Minimum Rotation Palindrome Tree Main Lorentz Fraction Fraction Miller Rabin / Pollard Rho Ext GCD PiCount Linear Function Mod Min Floor Sum Quadratic Residue Simplex Linear Programming Construction Deuclidean Theorem External Bisector Intersection of Circle and Line Intersection of Circle and Line Intersection of Circle and Point Tangent Lines of Circle and Point Tangent Lines of Circle and Point Tangent Lines of Circles Point In Convex Point Segment Distance Convex Hull Sconvex Hull Distance Minimum Enclosing Circle Sunion of Circles Sunion of Circles Sunion of Circles Sunion of Circles	· · · · · · · · · · · · · · · · · · ·									

```
8 Polvnomial
                 19
19
    19
19
19
9 Else
                 20
9.1 Bit Hack
                 20
20
 9.2.2 Monge Condition (Concave/Convex) . . . .
                 20
 9.2.3 Optimal Split Point . . . . . . . . . . . . . . . .
                 20
20
21
21
9.9 Random
     9.10Smawk Algorithm
                 21
9.11Two Dimension Add Sum . . . . . . . . . . . . .
                 22
```

1 Basic

1

1

1

1

1

2

3

6

7

9

9

9

9

9

10 10

10 10

10 11

11 12

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1.1 Compiler Shell

```
if [ $# -ne 2 ] ; then
 g++ -std=c++17 -DABS -Wall -Wextra -Wshadow $1.cpp -o
else
 g++ -std=c++17 -DABS -Wall -Wextra -Wshadow $1.cpp -o
       $1 -fsanitize=address
fi
./$1
chmod +x ./run.sh
./run.sh main [1]
```

1.2 Create File

```
for i in {A..M}
 cp tem.cpp "$i".cpp
done
```

1.3 Default Code

```
#include <bits/stdc++.h>
using namespace std;
#define ll long long
#define pb push_back
#define all(x) x.begin(), x.end()
#define pii pair<int, int>
#define pi pii
```

1.4 Testing Todo List

```
14 0. choose editor
    1. shell script
    2. _

    __int128, __lg, __builtin_pe
    judge speed v.s.local speed

                              _builtin_popcount
    3.1 bitset, +, ^, segment tree
    4. pragma CE?
16 | 5. CE penalty?
```

1.5 Debug Macro

```
void db() {cout << endl;}</pre>
template <typename T, typename ...U> void db(T i, U ...
    j) {
  cout << i << ' ', db(j...);
#define test(x...) db("[" + string(x) + "]", x)
```

1.6 Stress Test Shell

```
g++ $1.cpp -o $1
g++ $2.cpp -o $2
g++ $3.cpp -o $3
for i in {1..100} ; do
    ./$3 > input.txt
  # st=$(date +%s%N)
  ./$1 < input.txt > output1.txt
# echo "$((($(date +%s%N) - $st)/1000000))ms"
```

```
./$2 < input.txt > output2.txt
if cmp --silent -- "output1.txt" "output2.txt"; then
    continue
fi
    echo Input:
    cat input.txt
    echo Your Output:
    cat output1.txt
    echo Correct Output:
    cat output2.txt
    break
done
echo OK!
./stress.sh main good gen
```

1.7 Pragma

```
#pragma GCC optimize("Ofast,inline,unroll-loops")
#pragma GCC target("bmi,bmi2,lzcnt,popcnt,avx2")
```

1.8 Fast IO

```
#include<unistd.h>
char OB[65536]; int OP;
inline char RC() {
  static char buf[65536], *p = buf, *q = buf;
  return p == q \&\& (q = (p = buf) + read(0, buf, 65536)
      ) == buf ? -1 : *p++;
inline int R() {
  static char c;
 while((c = RC()) < '0'); int a = c ^ '0';
while((c = RC()) >= '0') a *= 10, a += c ^ '0';
 return a;
inline void W(int n) {
  static char buf[12], p;
  if (n == 0) OB[OP++] = '0'; p = 0;
 while (n) buf[p++] = '0' + (n \% 10), n /= 10;
  for (--p; p >= 0; --p) OB[OP++] = buf[p];
  if (OP > 65520) write(1, OB, OP), OP = 0;
```

1.9 Divide

```
ll divdown(ll a, ll b) {
  return a / b - (a < 0 && a % b);
}
ll divup(ll a, ll b) {
  return a / b + (a > 0 && a % b);
}
a / b < x -> divdown(a, b) + 1 <= x
a / b <= x -> divup(a, b) <= x
x < a / b -> x <= divup(a, b) - 1
x <= a / b -> x <= divdown(a, b)</pre>
```

2 Data Structure

2.1 Leftist Tree

```
struct node {
  ll rk, data, sz, sum;
  node *1, *r;
  node(11 \ k) : rk(0), data(k), sz(1), l(0), r(0), sum(k)
       ) {}
11 sz(node *p) { return p ? p->sz : 0; }
11 rk(node *p) { return p ? p->rk : -1;
11 sum(node *p) { return p ? p->sum : 0; }
node *merge(node *a, node *b) {
  if (!a || !b) return a ? a : b;
  if (a->data < b->data) swap(a, b);
  a->r = merge(a->r, b);
  if (rk(a->r) > rk(a->l)) swap(a->r, a->l);
  a->rk = rk(a->r) + 1, a->sz = sz(a->l) + sz(a->r) +
      1;
  a\rightarrow sum = sum(a\rightarrow 1) + sum(a\rightarrow r) + a\rightarrow data;
  return a;
void pop(node *&o) {
  node *tmp = o;
```

```
o = merge(o->1, o->r);
  delete tmp;
2.2 Splay Tree
struct Splay {
  int pa[N], ch[N][2], sz[N], rt, _id;
  11 v[N];
  Splay() {}
  void init() {
    rt = 0, pa[0] = ch[0][0] = ch[0][1] = -1;
    sz[0] = 1, v[0] = inf;
  int newnode(int p, int x) {
    int id = _id++;
    v[id] = x, pa[id] = p;
    ch[id][0] = ch[id][1] = -1, sz[id] = 1;
    return id:
  void rotate(int i) {
    int p = pa[i], x = ch[p][1] == i, gp = pa[p], c =
        ch[i][!x];
    sz[p] -= sz[i], sz[i] += sz[p];
    if (~c) sz[p] += sz[c], pa[c] = p;
    ch[p][x] = c, pa[p] = i;
    pa[i] = gp, ch[i][!x] = p;
    if (~gp) ch[gp][ch[gp][1] == p] = i;
  void splay(int i) {
    while (~pa[i]) {
      int p = pa[i];
      if (~pa[p]) rotate(ch[pa[p]][1] == p ^ ch[p][1]
           == i ? i : p);
      rotate(i);
    }
    rt = i;
  int lower_bound(int x) {
    int i = rt, last = -1;
    while (true) {
      if (v[i] == x) return splay(i), i;
      \quad \textbf{if} \ (v[\texttt{i}] \ \Rightarrow \ x) \ \{
        last = i;
        if (ch[i][0] == -1) break;
        i = ch[i][0];
      else {
        if (ch[i][1] == -1) break;
        i = ch[i][1];
      }
    splay(i);
    return last; // -1 if not found
  void insert(int x) {
    int i = lower_bound(x);
    if (i == -1) {
      // assert(ch[rt][1] == -1);
      int id = newnode(rt, x);
      ch[rt][1] = id, ++sz[rt];
      splay(id);
    else if (v[i] != x) {
      splay(i);
      int id = newnode(rt, x), c = ch[rt][0];
      ch[rt][0] = id;
      ch[id][0] = c;
      if (~c) pa[c] = id, sz[id] += sz[c];
      ++sz[rt]:
      splay(id);
```

2.3 Link Cut Tree

}

};

```
// assert(sz[i] == 2 && ch[i][1] == j);
// mx := max w in splay
bool rev[N];
                                                                ch[i][1] = pa[j] = 0, pull(i);
LCT() : _id(1) {}
                                                             int findrt(int i) {
int newnode(int _v, int _w) {
  int x = _id++;
ch[x][0] = ch[x][1] = pa[x] = 0;
                                                               access(i), splay(i);
                                                                while (ch[i][0]) push(i), i = ch[i][0];
  v[x] = sz[x] = _v;
                                                                splay(i);
  sz2[x] = 0;
                                                                return i;
  w[x] = mx[x] = w;
  rev[x] = false;
                                                          };
                                                           2.4 Treap
void pull(int i) {
                                                           struct node {
  sz[i] = v[i] + sz2[i];
                                                             int data, sz;
  mx[i] = w[i];
                                                             node *1, *r;
  if (ch[i][0])
                                                             node(int k) : data(k), sz(1), l(0), r(0) {}
    sz[i] += sz[ch[i][0]], mx[i] = max(mx[i], mx[ch[i])
                                                             void up() {
         ][0]]);
                                                                sz = 1;
  if (ch[i][1])
                                                               if (1) sz += 1->sz;
    sz[i] += sz[ch[i][1]], mx[i] = max(mx[i], mx[ch[i])
                                                               if (r) sz += r->sz;
        ][1]]);
}
                                                             void down() {}
void push(int i) {
  if (rev[i]) reverse(ch[i][0]), reverse(ch[i][1]),
                                                           int sz(node *a) { return a ? a->sz : 0; }
      rev[i] = false;
                                                           node *merge(node *a, node *b) {
                                                             if (!a || !b) return a ? a : b;
void reverse(int i) {
                                                             if (rand() % (sz(a) + sz(b)) < sz(a))
  if (!i) return;
                                                                return a \rightarrow down(), a \rightarrow r = merge(a \rightarrow r, b), a \rightarrow up(), a
  swap(ch[i][0], ch[i][1]);
  rev[i] ^= true;
                                                             return b->down(), b->l = merge(a, b->l), b->up(), b;
bool isrt(int i) {// rt of splay
                                                           void split(node *o, node *&a, node *&b, int k) {
  if (!pa[i]) return true;
                                                             if (!o) return a = b = 0, void();
  return ch[pa[i]][0] != i && ch[pa[i]][1] != i;
                                                             o->down();
                                                             if (o->data <= k)
void rotate(int i) {
                                                               a = o, split(o->r, a->r, b, k), a->up();
  int p = pa[i], x = ch[p][1] == i, c = ch[i][!x], gp
                                                             else b = o, split(o \rightarrow l, a, b \rightarrow l, k), b \rightarrow up();
       = pa[p];
  if (ch[gp][0] == p) ch[gp][0] = i;
                                                           void split2(node *o, node *&a, node *&b, int k) {
  else if (ch[gp][1] == p) ch[gp][1] = i;
                                                             if (sz(o) <= k) return a = o, b = 0, void();</pre>
  pa[i] = gp, ch[i][!x] = p, pa[p] = i;
                                                             o->down();
  ch[p][x] = c, pa[c] = p;
                                                             if (sz(o->1) + 1 <= k)
  pull(p), pull(i);
                                                               a = o, split2(o->r, a->r, b, k - <math>sz(o->l) - 1);
                                                             else b = o, split2(o->1, a, b->1, k);
void splay(int i) {
                                                             o->up();
  vector<int> anc;
  anc.push back(i):
                                                           node *kth(node *o, int k) {
  while (!isrt(anc.back())) anc.push_back(pa[anc.back
                                                             if (k <= sz(o->1)) return kth(o->1, k);
      ()1):
                                                             if (k == sz(o\rightarrow 1) + 1) return o;
  while (!anc.empty()) push(anc.back()), anc.pop_back
                                                             return kth(o\rightarrow r, k - sz(o\rightarrow l) - 1);
      ();
  while (!isrt(i)) {
                                                           int Rank(node *o, int key) {
    int p = pa[i];
                                                             if (!o) return 0;
    if (!isrt(p)) rotate(ch[p][1] == i ^ ch[pa[p]][1]
                                                             if (o->data < key)</pre>
         == p ? i : p);
                                                                return sz(o->1) + 1 + Rank(o->r, key);
    rotate(i);
                                                             else return Rank(o->1, key);
  }
                                                           bool erase(node *&o, int k) {
void access(int i) {
                                                             if (!o) return 0;
  int last = 0;
                                                             if (o->data == k) {
  while (i) {
                                                               node *t = o;
    splay(i);
                                                               o->down(), o = merge(o->1, o->r);
    if (ch[i][1])
                                                                delete t;
      sz2[i] += sz[ch[i][1]];
                                                               return 1;
    sz2[i] -= sz[last];
    ch[i][1] = last;
                                                             node *&t = k < o->data ? o->l : o->r;
    pull(i), last = i, i = pa[i];
                                                             return erase(t, k) ? o->up(), 1 : 0;
}
                                                           void insert(node *&o, int k) {
void makert(int i) {
                                                             node *a, *b;
  access(i), splay(i), reverse(i);
                                                             split(o, a, b, k),
                                                             o = merge(a, merge(new node(k), b));
void link(int i, int j) {
 // assert(findrt(i) != findrt(j));
                                                           void interval(node *&o, int 1, int r) {
  makert(i);
                                                             node *a, *b, *c;
  makert(j);
                                                             split2(o, a, b, l - 1), split2(b, b, c, r);
  pa[i] = j;
                                                             // operate
  sz2[j] += sz[i];
                                                             o = merge(a, merge(b, c));
 pull(j);
void cut(int i, int j) {
                                                           2.5 Persistent Segment Tree
  makert(i), access(j), splay(i);
```

```
struct Seg {
 // Persistent Segment Tree, single point modify,
      range query sum
  // 0-indexed, [l, r)
  static Seg mem[M], *pt;
  int 1, r, m, val;
  Seg* ch[2];
  Seg () = default;
  Seg (int _1, int _r) : l(_1), r(_r), m(1 + r >> 1),
      val(0) {
    if (r - 1 > 1) {
      ch[0] = new (pt++) Seg(1, m);
      ch[1] = new (pt++) Seg(m, r);
   }
 }
  void pull() {val = ch[0]->val + ch[1]->val;}
  Seg* modify(int p, int v) {
    Seg *now = new (pt++) Seg(*this);
    if (r - l == 1) {
      now \rightarrow val = v;
    } else {
      now \rightarrow ch[p >= m] = ch[p >= m] \rightarrow modify(p, v);
      now->pull();
    return now;
  int query(int a, int b) {
   if (a <= 1 && r <= b) return val;</pre>
    int ans = 0;
    if (a < m) ans += ch[0]->query(a, b);
    if (m < b) ans += ch[1]->query(a, b);
    return ans;
} Seg::mem[M], *Seg::pt = mem;
// Init Tree
Seg *root = new (Seg::pt++) Seg(0, n);
```

2.6 2D Segment Tree

```
// 2D range add, range sum in log^2
struct seg {
  int 1, r;
  11 sum, 1z;
  seg *ch[2]{};
  seg(int _1, int _r) : 1(_1), r(_r), sum(0), lz(0) {}
  void push() {
    if (lz) ch[0]->add(l, r, lz), ch[1]->modify(l, r,
        1z), 1z = 0;
  void pull() \{sum = ch[0] -> sum + ch[1] -> sum;\}
  void add(int _l, int _r, ll d) {
    if (_1 <= 1 && r <= .
      sum += d * (r - 1);
      1z += d;
      return;
    if (!ch[0]) ch[0] = new seg(1, 1 + r >> 1), ch[1] =
         new seg(1 + r \gg 1, r);
    push();
    if (_1 < 1 + r >> 1) ch[0]->add(_1, _r, d);
    if (l + r >> 1 < _r) ch[1]->add(_l, _r, d);
    pull();
  11 qsum(int _1, int _r) {
    if (_1 <= 1 && r <= _r) return sum;
if (!ch[0]) return lz * (min(r, _r) - max(1, _1));</pre>
    push();
    11 \text{ res} = 0;
    if (_1 < 1 + r >> 1) res += ch[0]->qsum(_1, _r);
    if (1 + r >> 1 < _r) res += ch[1]->qsum(_1, _r);
    return res:
 }
};
struct seg2 {
  int 1, r;
  seg v, lz;
  seg2 *ch[2]{};
  seg2(int _1, int _r) : l(_1), r(_r), v(0, N), lz(0, N
    if (1 < r - 1) ch[0] = new seg2(1, 1 + r >> 1), ch
        [1] = new seg2(1 + r >> 1, r);
  }
```

```
void add(int _1, int _r, int _12, int _r2, 11 d) {
  v.add(_12, _r2, d * (min(r, _r) - max(1, _1)));
  if (_1 <= 1 && r <= _r) {</pre>
        lz.add(_12, _r2, d);
        return;
     if (_1 < 1 + r >> 1) ch[0]->add(_1, _r, _12, _r2, d)
     if (1 + r >> 1 < _r) ch[1]->add(_1, _r, _12, _r2, d)
          );
   11 qsum(int _1, int _r, int _12, int _r2) {
     11 res = v.qsum(_12, _r2);
     if (_1 <= 1 && r <= _r) return res;</pre>
     res += lz.qsum(_12, _r2) * (min(r, _r) - max(1, _1)
     if (_1 < 1 + r >> 1) res += ch[0]->query(_1, _r,
           _12, _r2);
      if (1 + r >> 1 < _r) res += ch[1]->query(_1, _r,
           _12, _r2);
     return res;
};
```

2.7 Zkw

```
ll mx[N << 1], sum[N << 1], lz[N << 1];
void add(int 1, int r, 11 d) { // [l, r), 0-based
  int len = 1, cntl = 0, cntr = 0;
  for (1 += N, r += N + 1; 1 ^ r ^ 1; 1 >>= 1, r >>= 1,
       len <<= 1) {
    sum[1] += cnt1 * d, sum[r] += cnt[r] * d;
    if (len > 1) {
      mx[1] = max(mx[1 << 1], mx[1 << 1 | 1]) + lz[1];
      mx[r] = max(mx[r << 1], mx[r << 1 | 1]) + lz[r];
    if (~1 & 1)
      sum[1 ^ 1] += d * len, mx[l ^ 1] += d, lz[l ^ 1]
          += d, cntl += len;
    if (r & 1)
      sum[r ^ 1] += d * len, mx[r ^ 1] += d, lz[r ^ 1]
          += d, cntr += len;
  sum[1] += cntl * d, sum[r] += cntr * d;
  if (len > 1) {
    mx[1] = max(mx[1 << 1], mx[1 << 1 | 1]) + lz[1];
    mx[r] = max(mx[r << 1], mx[r << 1 | 1]) + lz[r];
  cntl += cntr;
  for (1 >>= 1; 1; 1 >>= 1) {
    sum[1] += cntl * d;
    mx[1] = max(mx[1 << 1], mx[1 << 1 | 1]) + lz[1];
  }
11 qsum(int 1, int r) {
  ll res = 0, len = 1, cntl = 0, cntr = 0;
  for (1 += N, r += N + 1; 1 ^ r ^ 1; 1 >>= 1, r >>= 1,
       len <<= 1) {
    res += cntl * lz[1] + cntr * lz[r];
    if (~l & 1) res += sum[l ^ 1], cntl += len;
    if (r & 1) res += sum[r ^ 1], cntr += len;
  res += cntl * lz[1] + cntr * lz[r];
  cntl += cntr;
  for (1 >>= 1; 1; 1 >>= 1) res += cntl * lz[1];
  return res;
11 qmax(int 1, int r) {
  11 max1 = -INF, maxr = -INF;
  for (1 += N, r += N + 1; 1 ^ r ^ 1; 1 >>= 1, r >>= 1)
    \max 1 += 1z[1], \max[r] += 1z[r];
    if (~1 & 1) maxl = max(maxl, mx[l ^ 1]);
    if (r & 1) maxr = max(maxr, mx[r ^ 1]);
  \max l = \max(\max l + lz[l], \max r + lz[r]);
  for (1 >>= 1; 1; 1 >>= 1) max1 += lz[1];
  return max1;
```

2.8 Chtholly Tree

```
struct ChthollyTree {
  struct interval {
    int 1, r;
    11 v:
    interval (int _l, int _r, ll _v) : l(_l), r(_r), v(
  struct cmp {
    bool operator () (const interval &a, const interval
        & b) const {
      return a.l < b.l:
    }
  set <interval, cmp> s;
  vector <interval> split(int 1, int r) {
    // split into [0, l), [l, r), [r, n) and return [l, r]
         r)
    vector <interval> del, ans, re;
    auto it = s.lower_bound(interval(l, -1, 0));
    if (it != s.begin() && (it == s.end() || 1 < it->1)
        ) {
      --it;
      del.pb(*it);
      if (r < it->r) {
        re.pb(interval(it->1, 1, it->v));
        ans.pb(interval(l, r, it->v));
        re.pb(interval(r, it->r, it->v));
      } else {
        re.pb(interval(it->1, 1, it->v));
        ans.pb(interval(l, it->r, it->v));
      ++it;
    for (; it != s.end() && it->r <= r; ++it) {</pre>
      ans.pb(*it);
      del.pb(*it);
    if (it != s.end() && it->l < r) {</pre>
      del.pb(*it);
      ans.pb(interval(it->l, r, it->v));
      re.pb(interval(r, it->r, it->v));
    for (interval &i : del)
      s.erase(i);
    for (interval &i : re)
      s.insert(i);
    return ans;
  void merge(vector <interval> a) {
    for (interval &i : a)
      s.insert(i);
  }
};
```

2.9 Incremental Min Sum

```
struct IncrementalMinSum {
 multiset <int, greater <int>> in;
  multiset <int> out;
 11 sum; int cap;
 DS () : sum(0), cap(0) {}
 void enlarge() {
    if (!out.empty()) {
      int mx = *out.begin();
      sum += mx, in.insert(mx), out.erase(out.begin());
    }
    cap++;
  void insert(int x) {
    if (!cap) {
      out.insert(x);
      return;
    if (in.size() < cap) {</pre>
      in.insert(x), sum += x;
      return;
    int mx = *in.begin();
    if (x < mx) {
      sum -= mx, out.insert(mx), in.erase(in.begin());
      sum += x, in.insert(x);
    } else {
```

```
out.insert(x);
}

void erase(int x) {
  if (out.find(x) != out.end()) {
    out.erase(out.lower_bound(x));
} else {
    in.erase(in.lower_bound(x)), sum -= x;
    if (!out.empty()) {
        int mx = *out.begin();
        sum += mx, out.erase(out.begin()), in.insert(mx );
    }
}

}
}
```

3 Flow / Matching

3.1 Dinic

```
struct Dinic { // 0-base
  struct edge {
    int to, cap, flow, rev;
  vector<edge> adj[N];
  int s, t, dis[N], cur[N], n;
  int dfs(int u, int cap) {
    if (u == t || !cap) return cap;
    for (int &i = cur[u]; i < (int)adj[u].size(); ++i)</pre>
       edge &e = adj[u][i];
       if (dis[e.to] == dis[u] + 1 && e.flow != e.cap) {
         int df = dfs(e.to, min(e.cap - e.flow, cap));
         if (df) {
           e.flow += df;
           adj[e.to][e.rev].flow -= df;
           return df;
         }
      }
    dis[u] = -1;
    return 0;
  bool bfs() {
    fill_n(dis, n, -1);
    queue<int> q;
    q.push(s), dis[s] = 0;
    while (!q.empty()) {
      int tmp = q.front();
       q.pop();
       for (auto &u : adj[tmp])
         if (!~dis[u.to] && u.flow != u.cap) {
           q.push(u.to);
           dis[u.to] = dis[tmp] + 1;
         }
    return dis[t] != -1;
  int maxflow(int _s, int _t) {
    s = _s, t = _t;
int flow = 0, df;
    while (bfs()) {
      fill_n(cur, n, 0);
while ((df = dfs(s, INF))) flow += df;
    return flow;
  void init(int _n) {
    n = _n;
    for (int i = 0; i < n; ++i) adj[i].clear();</pre>
  void reset() {
    for (int i = 0; i < n; ++i)</pre>
       for (auto &j : adj[i]) j.flow = 0;
  void add_edge(int u, int v, int cap) {
  adj[u].pb(edge{v, cap, 0, (int)adj[v].size()});
     adj[v].pb(edge{u, 0, 0, (int)adj[u].size() - 1});
};
```

3.2 Min Cost Max Flow

```
template <typename T>
struct MCMF {
  const T INF = 111 << 60;</pre>
  struct edge {
    int v;
    T f, c;
    edge (int _v, T _f, T _c) : v(_v), f(_f), c(_c) {}
  vector <edge> E;
  vector <vector <int>> adj;
  vector <T> dis, pot;
  vector <int> rt;
  int n, s, t;
  \label{eq:mcmf} \mbox{MCMF (int _n, int _s, int _t) : n(_n), s(_s), t(_t) {} \{ \mbox{}
    adj.resize(n);
  void add_edge(int u, int v, T f, T c) {
    adj[u].pb(E.size()), E.pb(edge(v, f, c));
    adj[v].pb(E.size()), E.pb(edge(u, 0, -c));
  bool SPFA() {
    rt.assign(n, -1), dis.assign(n, INF);
vector <bool> vis(n, false);
    queue <int> q;
    q.push(s), dis[s] = 0, vis[s] = true;
    while (!q.empty()) {
       int v = q.front(); q.pop();
      vis[v] = false;
for (int id : adj[v]) if (E[id].f > 0 && dis[E[id].f]
           ].v] > dis[v] + E[id].c + pot[v] - pot[E[id].
           v1) {
           dis[E[id].v] = dis[v] + E[id].c + pot[v] -
               pot[E[id].v], rt[E[id].v] = id;
           if (!vis[E[id].v]) vis[E[id].v] = true, q.
                push(E[id].v);
         }
    return dis[t] != INF;
  bool dijkstra() {
    rt.assign(n, -1), dis.assign(n, INF);
    priority_queue <pair <T, int>, vector <pair <T, int</pre>
         >>, greater <pair <T, int>>> pq;
    dis[s] = 0, pq.emplace(dis[s], s);
    while (!pq.empty()) {
       int d, v; tie(d, v) = pq.top(); pq.pop();
       if (dis[v] < d) continue;
for (int id : adj[v]) if (E[id].f > 0 && dis[E[id
           ].v] > dis[v] + E[id].c + pot[v] - pot[E[id].
           v]) {
           dis[E[id].v] = dis[v] + E[id].c + pot[v] -
               pot[E[id].v], rt[E[id].v] = id;
           pq.emplace(dis[E[id].v], E[id].v);
    return dis[t] != INF;
  pair <T, T> solve() {
    pot.assign(n, 0);
     T cost = 0, flow = 0;
    bool fr = true;
    while ((fr ? SPFA() : dijkstra())) {
       for (int i = 0; i < n; i++) {</pre>
         dis[i] += pot[i] - pot[s];
       T add = INF;
       for (int i = t; i != s; i = E[rt[i] ^ 1].v) {
         add = min(add, E[rt[i]].f);
       for (int i = t; i != s; i = E[rt[i] ^ 1].v) {
         E[rt[i]].f -= add, E[rt[i] ^ 1].f += add;
       flow += add, cost += add * dis[t];
       fr = false;
       swap(dis, pot);
     return make_pair(flow, cost);
  }
|};
```

3.3 Kuhn Munkres

```
template <typename T>
struct KM { // 0-based
  T w[N][N], h1[N], hr[N], slk[N];
  T fl[N], fr[N], pre[N]; int n;
  bool vl[N], vr[N];
  const T INF = 1e9;
  queue <int> q;
  KM (int _n) : n(_n) {
    for (int i = 0; i < n; ++i) for (int j = 0; j < n;
         ++j)
         w[i][j] = -INF;
  void add_edge(int a, int b, int wei) {
    w[a][b] = wei;
  bool check(int x) {
     if (vl[x] = 1, ~fl[x]) return q.push(fl[x]), vr[fl[
         x]] = 1;
     while (\sim x) swap(x, fr[fl[x] = pre[x]]);
     return 0;
  void bfs(int s) {
     fill(slk, slk + n, INF), fill(vl, vl + n, 0), fill(
         vr, vr + n, 0);
     q.push(s), vr[s] = 1;
     while (1) {
       T d;
       while (!q.empty()) {
         int y = q.front(); q.pop();
         for (int x = 0; x < n; ++x)
           if (!vl[x] \&\& slk[x] >= (d = hl[x] + hr[y] -
               w[x][y])
             if (pre[x] = y, d) slk[x] = d;
             else if (!check(x)) return;
       d = INF;
       for (int x = 0; x < n; ++x)
        if (!v1[x] && d > s1k[x]) d = s1k[x];
       for (int x = 0; x < n; ++x) {
         if (v1[x]) h1[x] += d;
         else slk[x] -= d;
         if (vr[x]) hr[x] -= d;
       for (int x = 0; x < n; ++x) if (!v1[x] && !s1k[x]
            && !check(x)) return;
    }
  T solve() {
     fill(fl, fl + n, -1), fill(fr, fr + n, -1), fill(hr
          hr + n, 0);
     for (int i = 0; i < n; ++i) hl[i] = *max_element(w[</pre>
         i], w[i] + n);
     for (int i = 0; i < n; ++i) bfs(i);</pre>
    T res = 0:
     for (int i = 0; i < n; ++i) res += w[i][fl[i]];</pre>
     return res;
| };
```

3.4 SW Min Cut

```
int now = -1;
        for (int i = 0; i < n; ++i) if (!vis[i] && !</pre>
             dead[i]) {
             if (now == -1 \mid | sum[now] < sum[i]) now = i
        s = t, t = now;
        vis[now] = true, num++;
        for (int i = 0; i < n; ++i) if (!vis[i] && !</pre>
             dead[i]) {
             sum[i] += g[now][i];
      ans = min(ans, sum[t]);
      for (int i = 0; i < n; ++i) {</pre>
        g[i][s] += g[i][t];
        g[s][i] += g[t][i];
      dead[t] = true;
    return ans;
  }
};
```

3.5 Gomory Hu Tree

```
vector <array <int, 3>> GomoryHu(vector <vector <pii>>>
    adj, int n) {
// Tree edge min -> mincut (0-based)
 Dinic flow(n);
  for (int i = 0; i < n; ++i) for (auto [j, w] : adj[i
      1)
      flow.add_edge(i, j, w);
 flow.record();
 vector <array <int, 3>> ans;
  vector <int> rt(n);
  for (int i = 0; i < n; ++i) rt[i] = 0;</pre>
  for (int i = 1; i < n; ++i) {</pre>
    int t = rt[i];
    flow.reset(); // clear flows on all edge
    ans.push_back({i, t, flow.solve(i, t)});
    flow.runbfs(i);
    for (int j = i + 1; j < n; ++j) if (rt[j] == t &&</pre>
        flow.vis[j]) {
        rt[j] = i;
 }
  return ans;
```

3.6 Blossom

```
struct Matching { // 0-based
 int fa[N], pre[N], match[N], s[N], v[N], n, tk;
  vector <int> g[N];
  queue <int> q;
  Matching (int _n) : n(_n), tk(0) {
    for (int i = 0; i <= n; ++i) match[i] = pre[i] = n;</pre>
    for (int i = 0; i < n; ++i) g[i].clear();</pre>
  void add_edge(int u, int v) {
    g[u].push_back(v), g[v].push_back(u);
  int Find(int u) {
   return u == fa[u] ? u : fa[u] = Find(fa[u]);
  int lca(int x, int y) {
    x = Find(x), y = Find(y);
    for (; ; swap(x, y)) {
  if (x != n) {
        if (v[x] == tk) return x;
        v[x] = tk;
        x = Find(pre[match[x]]);
   }
  void blossom(int x, int y, int 1) {
    while (Find(x) != 1) {
      pre[x] = y, y = match[x];
if (s[y] == 1) q.push(y), s[y] = 0;
      if (fa[x] == x) fa[x] = 1;
```

```
if (fa[y] == y) fa[y] = 1;
      x = pre[y];
    }
  bool bfs(int r) {
    for (int i = 0; i <= n; ++i) fa[i] = i, s[i] = -1;
    while (!q.empty()) q.pop();
    q.push(r);
    s[r] = 0:
    while (!q.empty()) {
      int x = q.front(); q.pop();
      for (int u : g[x]) {
        if (s[u] == -1) {
          pre[u] = x, s[u] = 1;
          if (match[u] == n) {
            for (int a = u, b = x, last; b != n; a =
                 last, b = pre[a])
              last = match[b], match[b] = a, match[a] =
                   b;
            return true;
          q.push(match[u]);
          s[match[u]] = 0;
        } else if (!s[u] && Find(u) != Find(x)) {
          int 1 = lca(u, x);
          blossom(x, u, 1);
          blossom(u, x, 1);
        }
      }
    return false:
  int solve() {
    int res = 0;
    for (int x = 0; x < n; ++x) {
      if (match[x] == n) res += bfs(x);
    return res;
  }
};
```

3.7 Weighted Blossom

```
struct WeightGraph { // 1-based
  static const int inf = INT_MAX;
  static const int maxn = 514;
  struct edge {
    int u, v, w;
    edge(){}
    edge(int u, int v, int w): u(u), v(v), w(w) {}
  };
  int n, n_x;
  edge g[maxn * 2][maxn * 2];
  int lab[maxn * 2];
  int match[maxn * 2], slack[maxn * 2], st[maxn * 2],
      pa[maxn * 2];
  int flo_from[maxn * 2][maxn + 1], S[maxn * 2], vis[
      maxn * 2];
  vector<int> flo[maxn * 2];
  queue<int> q;
  int e_delta(const edge &e) { return lab[e.u] + lab[e.
  v] - g[e.u][e.v].w * 2; }
void update_slack(int u, int x) { if (!slack[x] ||
      e_delta(g[u][x]) < e_delta(g[slack[x]][x])) slack</pre>
      [x] = u;
  void set_slack(int x) {
    slack[x] = 0;
    for (int u = 1; u <= n; ++u)</pre>
      if (g[u][x].w > 0 \&\& st[u] != x \&\& S[st[u]] == 0)
        update_slack(u, x);
  void q_push(int x) {
    if (x <= n) q.push(x);</pre>
    else for (size_t i = 0; i < flo[x].size(); i++)</pre>
        q_push(flo[x][i]);
  void set_st(int x, int b) {
    st[x] = b;
    if (x > n) for (size_t i = 0; i < flo[x].size(); ++
        i) set_st(flo[x][i], b);
  int get_pr(int b, int xr) {
```

```
int pr = find(flo[b].begin(), flo[b].end(), xr) -
      flo[b].begin();
  if (pr % 2 == 1) {
    reverse(flo[b].begin() + 1, flo[b].end());
    return (int)flo[b].size() - pr;
  return pr;
void set_match(int u, int v) {
  match[u] = g[u][v].v;
  if (u <= n) return;</pre>
  edge e = g[u][v];
  int xr = flo_from[u][e.u], pr = get_pr(u, xr);
  for (int i = 0; i < pr; ++i) set_match(flo[u][i],</pre>
      flo[u][i ^ 1]);
  set match(xr, v);
  rotate(flo[u].begin(), flo[u].begin() + pr, flo[u].
      end());
void augment(int u, int v) {
  for (; ; ) {
    int xnv = st[match[u]];
    set_match(u, v);
    if (!xnv) return;
    set_match(xnv, st[pa[xnv]]);
    u = st[pa[xnv]], v = xnv;
 }
}
int get_lca(int u, int v) {
  static int t = 0;
  for (++t; u || v; swap(u, v)) {
    if (u == 0) continue;
    if (vis[u] == t) return u;
    vis[u] = t;
    u = st[match[u]];
    if (u) u = st[pa[u]];
  return 0;
void add_blossom(int u, int lca, int v) {
 int b = n + 1;
  while (b <= n_x && st[b]) ++b;</pre>
  if (b > n_x) ++n_x;
  lab[b] = 0, S[b] = 0;
  match[b] = match[lca];
  flo[b].clear();
  flo[b].push_back(lca);
  for (int x = u, y; x != lca; x = st[pa[y]])
    flo[b].push_back(x), flo[b].push_back(y = st[
        match[x]]), q_push(y);
  reverse(flo[b].begin() + 1, flo[b].end());
  for (int x = v, y; x != lca; x = st[pa[y]])
    flo[b].push_back(x), flo[b].push_back(y = st[
        match[x]]), q_push(y);
  set_st(b, b);
  for (int x = 1; x <= n_x; ++x) g[b][x].w = g[x][b].
  for (int x = 1; x \le n; ++x) flo_from[b][x] = 0;
  for (size_t i = 0; i < flo[b].size(); ++i) {</pre>
    int xs = flo[b][i];
    for (int x = 1; x <= n_x; ++x)
      if (g[b][x].w == 0 || e_delta(g[xs][x]) <</pre>
          e_delta(g[b][x]))
        g[b][x] = g[xs][x], g[x][b] = g[x][xs];
    for (int x = 1; x <= n; ++x)
      if (flo_from[xs][x]) flo_from[b][x] = xs;
  set_slack(b);
void expand_blossom(int b) {
  for (size_t i = 0; i < flo[b].size(); ++i)</pre>
    set_st(flo[b][i], flo[b][i]);
  int xr = flo_from[b][g[b][pa[b]].u], pr = get_pr(b,
       xr);
  for (int i = 0; i < pr; i += 2) {</pre>
    int xs = flo[b][i], xns = flo[b][i + 1];
    pa[xs] = g[xns][xs].u;
   S[xs] = 1, S[xns] = 0;
slack[xs] = 0, set_slack(xns);
    q_push(xns);
  S[xr] = 1, pa[xr] = pa[b];
```

```
for (size_t i = pr + 1; i < flo[b].size(); ++i) {</pre>
    int xs = flo[b][i];
    S[xs] = -1, set_slack(xs);
  st[b] = 0;
bool on_found_edge(const edge &e) {
  int u = st[e.u], v = st[e.v];
  if (S[v] == -1) {
    pa[v] = e.u, S[v] = 1;
    int nu = st[match[v]];
    slack[v] = slack[nu] = 0;
    S[nu] = 0, q_push(nu);
  } else if (S[v] == 0) {
    int lca = get_lca(u, v);
    if (!lca) return augment(u,v), augment(v,u), true
    else add_blossom(u, lca, v);
  return false;
bool matching() {
  memset(S + 1, -1, sizeof(int) * n_x);
  memset(slack + 1, 0, sizeof(int) * n_x);
  q = queue<int>();
  for (int x = 1; x <= n_x; ++x)
    if (st[x] == x && !match[x]) pa[x] = 0, S[x] = 0,
         q_push(x);
  if (q.empty()) return false;
  for (;;) {
    while (q.size()) {
      int u = q.front(); q.pop();
      if (S[st[u]] == 1) continue;
      for (int v = 1; v <= n; ++v)
        if (g[u][v].w > 0 && st[u] != st[v]) {
          if (e_delta(g[u][v]) == 0) {
            if (on_found_edge(g[u][v])) return true;
          } else update_slack(u, st[v]);
        }
    int d = inf;
    for (int b = n + 1; b \le n_x; ++b)
      if (st[b] == b && S[b] == 1) d = min(d, lab[b]
          / 2);
    for (int x = 1; x <= n_x; ++x)
      if (st[x] == x && slack[x]) {
        if (S[x] == -1) d = min(d, e_delta(g[slack[x
            ]][x]));
        else if (S[x] == 0) d = min(d, e delta(g[
            slack[x]][x]) / 2);
    for (int u = 1; u <= n; ++u) {</pre>
      if (S[st[u]] == 0) {
        if (lab[u] <= d) return 0;</pre>
        lab[u] -= d;
      } else if (S[st[u]] == 1) lab[u] += d;
    for (int b = n + 1; b \le n_x; ++b)
      if (st[b] == b) {
        if (S[st[b]] == 0) lab[b] += d * 2;
        else if (S[st[b]] == 1) lab[b] -= d * 2;
    q = queue<int>();
    for (int x = 1; x <= n_x; ++x)
      if (st[x] == x && slack[x] && st[slack[x]] != x
           && e_delta(g[slack[x]][x]) == 0)
        if (on_found_edge(g[slack[x]][x])) return
            true:
    for (int b = n + 1; b <= n_x; ++b)
      if (st[b] == b && S[b] == 1 && lab[b] == 0)
          expand_blossom(b);
  return false;
pair<long long, int> solve() {
  memset(match + 1, 0, sizeof(int) * n);
  n_x = n;
  int n_matches = 0;
  long long tot_weight = 0;
  for (int u = 0; u <= n; ++u) st[u] = u, flo[u].
      clear();
  int w_max = 0;
```

```
for (int u = 1; u <= n; ++u)</pre>
       for (int v = 1; v <= n; ++v) {</pre>
         flo_from[u][v] = (u == v ? u : 0);
         w_max = max(w_max, g[u][v].w);
    for (int u = 1; u <= n; ++u) lab[u] = w_max;</pre>
    while (matching()) ++n_matches;
    for (int u = 1; u <= n; ++u)</pre>
      if (match[u] && match[u] < u)</pre>
         tot_weight += g[u][match[u]].w;
    return make_pair(tot_weight, n_matches);
  }
  void add_edge(int ui, int vi, int wi) { g[ui][vi].w =
        g[vi][ui].w = wi; }
  void init(int _n) {
    n = _n;
    for (int u = 1; u <= n; ++u)</pre>
      for (int v = 1; v <= n; ++v)</pre>
         g[u][v] = edge(u, v, 0);
  }
};
```

3.8 Flow Model

- Maximum/Minimum flow with lower bound / Circulation problem
 - 1. Construct super source S and sink T.
 - 2. For each edge (x,y,l,u), connect $x \to y$ with capacity u-l.
 - 3. For each vertex v, denote by in(v) the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.
 - 4. If in(v)>0, connect $S\to v$ with capacity in(v), otherwise, connect $v \to T$ with capacity -in(v).
 - To maximize, connect t o s with capacity ∞ (skip this in circulation problem), and let f be the maximum flow from S to T. If $f \neq \sum_{v \in V, in(v) > 0} in(v)$, there's no solution. Otherwise, the maximum flow from s to t is the answer.
 - To minimize, let f be the maximum flow from S to T. Connect $t \to s$ with capacity ∞ and let the flow from S to T be f'. If $f+f' \neq \sum_{v \in V, in(v)>0} in(v)$, there's no solution. Otherwise, f' is the answer.
 - 5. The solution of each edge e is $l_e + f_e$, where f_e corresponds to the flow of edge e on the graph.
- ullet Construct minimum vertex cover from maximum matching M on bipar- ${\rm tite\ graph\ }(X,Y)$
 - 1. Redirect every edge: $y \to x$ if $(x,y) \in M$, $x \to y$ otherwise.
 - 2. DFS from unmatched vertices in \boldsymbol{X} .
 - 3. $x \in X$ is chosen iff x is unvisited.
 - 4. $y \in Y$ is chosen iff y is visited.
- Maximum density induced subgraph
 - 1. Binary search on answer, suppose we're checking answer \boldsymbol{T}
 - 2. Construct a max flow model, let ${\cal K}$ be the sum of all weights
 - 3. Connect source $s \to v$, $v \in G$ with capacity K
 - 4. For each edge (u,v,w) in G, connect $u \to v$ and $v \to u$ with capacity \boldsymbol{w}
 - 5. For $v\in G$, connect it with sink $v\to t$ with capacity $K+2T-(\sum_{e\in E(v)}w(e))-2w(v)$
 - 6. ${\cal T}$ is a valid answer if the maximum flow $f < K |{\cal V}|$
- Minimum weight edge cover
 - 1. For each $v \in V$ create a copy v', and connect u' o v' with
 - weight w(u,v) . 2. Connect $v \to v'$ with weight $2\mu(v)$, where $\mu(v)$ is the cost of the cheapest edge incident to \widetilde{v} .
 - 3. Find the minimum weight perfect matching on G^\prime .
- Project selection problem
 - 1. If $p_v>0$, create edge (s,v) with capacity p_v ; otherwise, create edge (v,t) with capacity $-p_v$
 - 2. Create edge (u,v) with capacity w with w being the cost of choosing u without choosing v.
 - 3. The mincut is equivalent to the maximum profit of a subset of projects.

$$\sum_{x} c_{x}x + \sum_{y} c_{y}\bar{y} + \sum_{xy} c_{xy}x\bar{y} + \sum_{xyx'y'} c_{xyx'y'}(x\bar{y} + x'\bar{y'})$$

can be minimized by the mincut of the following graph:

- 1. Create edge (x,t) with capacity c_x and create edge (s,y) with capacity $c_y\,.$
- 2. Create edge (x,y) with capacity c_{xy}
- 3. Create edge (x,y) and edge (x^\prime,y^\prime) with capacity $c_{xyx^\prime y^\prime}$.

Graph

4.1 Heavy-Light Decomposition

```
vector<int> dep, pa, sz, ch, hd, id;
int _id;
void dfs(int i, int p) {
  dep[i] = \sim p ? dep[p] + 1 : 0;
  pa[i] = p, sz[i] = 1, ch[i] = -1;
  for (int j : g[i])
    if (j != p) {
      dfs(j, i);
      if (ch[i] == -1 || sz[ch[i]] < sz[j]) ch[i] = j;</pre>
      sz[i] += sz[j];
void hld(int i, int p, int h) {
  hd[i] = h;
  id[i] = _id++;
  if (~ch[i]) hld(ch[i], i, h);
  for (int j : g[i]) if (j != p && j != ch[i])
    hld(j, i, j);
void query(int i, int j) {
  while (hd[i] != hd[j]) {
    if (dep[hd[i]] < dep[hd[j]]) swap(i, j);</pre>
    query2(id[hd[i]], id[i] + 1), i = pa[hd[i]];
  if (dep[i] < dep[j]) swap(i, j);</pre>
  query2(id[j], id[i] + 1);
```

4.2 Centroid Decomposition

```
vector<vector<int>> dis;
vector<int> pa, sz;
vector<bool> vis;
void dfs_sz(int i, int p) {
  sz[i] = 1;
  for (int j : g[i]) if (j != p && !vis[j])
    dfs_sz(j, i), sz[i] += sz[j];
void cen(int i, int p, int _n) {
  for (int j : g[i]) if (j != p && !vis[j] && sz[j] >
      _n / 2)
    return cen(j, i, _n);
  return i;
void dfs_dis(int i, int p, int d) { // from i to
    ancestor with depth d
  dis[i][d] = \sim p ? dis[p][d] + 1 : 0;
  for (int j : g[i]) if (j != p && !vis[j])
    dfs_dis(j, i, d);
void cd(int i, int p, int d) {
  dfs sz(i), i = cen(i, -1, sz[i]);
  vis[i] = true, pa[i] = p;
  dfs_dis(i, -1, d);
  for (int j : g[i]) if (!vis[j])
    cd(j, i, d + 1);
```

4.3 Edge BCC

```
vector<int> low, dep, bcc_id, stk;
vector<bool> vis;
int _id;
void dfs(int i, int p) {
  low[i] = dep[i] = \sim p ? dep[p] + 1 : 0;
  stk.push_back(i);
  vis[i] = true;
  for (int j : g[i])
    if (j != p) {
      if (!vis[j])
        dfs(j, i), low[i] = min(low[i], low[j]);
        low[i] = min(low[i], dep[j]);
  if (low[i] == dep[i]) {
    int id = _id++;
    while (stk.back() != i) {
```

```
int x = stk.back();
    stk.pop_back();
    bcc_id[x] = id;
}
stk.pop_back();
bcc_id[i] = id;
}
}
```

4.4 Block Cut Tree

```
vector<vector<int>> g,
vector<int> dep, low, stk;
void dfs(int i, int p) {
  dep[i] = low[i] = \sim p ? dep[p] + 1 : 0;
  stk.push_back(i);
  for (int j : g[i]) if (j != p) {
  if (dep[j] == -1) {
      dfs(j, i), low[i] = min(low[i], low[j]);
       if (low[j] >= dep[i]) {
        int id = _g.size();
         _g.emplace_back();
        while (stk.back() != j) {
           int x = stk.back();
           stk.pop_back();
           _g[x].push_back(id), _g[id].push_back(x);
        stk.pop_back();
        _g[j].push_back(id), _g[id].push_back(j);
        _g[i].push_back(id), _g[id].push_back(i);
      else low[i] = min(low[i], dep[j]);
  }
}
```

4.5 SCC / 2SAT

```
struct SAT {
 vector<vector<int>> g;
  vector<int> dep, low, scc_id;
 vector<bool> is:
 vector<int> stk;
  int n, _id, _t;
 SAT() {}
  void init(int _n) {
   n = _n, _id = _t = 0;
    g.assign(2 * n, vector<int>());
    dep.assign(2 * n, -1), low.assign(2 * n, -1);
    scc_id.assign(2 * n, -1), is.assign(2 * n, false);
    stk.clear();
  void add_edge(int x, int y) {g[x].push_back(y);}
  int rev(int i) {return i < n ? i + n : i - n;}</pre>
  void add_ifthen(int x, int y) {add_clause(rev(x), y)
      ;}
  void add_clause(int x, int y) {
    add_edge(rev(x), y);
    add_edge(rev(y), x);
  void dfs(int i) {
    dep[i] = low[i] = _t++;
    stk.push_back(i);
    for (int j : g[i])
      if (scc_id[j] == -1) {
        if (dep[j] == -1)
          dfs(j);
        low[i] = min(low[i], low[j]);
    if (low[i] == dep[i]) {
      int id = _id++;
      while (stk.back() != i) {
        int x = stk.back();
        stk.pop_back();
        scc_id[x] = id;
      stk.pop_back();
      scc_id[i] = id;
   }
 }
  bool solve() {
    for (int i = 0; i < 2 * n; ++i)</pre>
```

if (dep[i] == -1)

```
dfs(i);
for (int i = 0; i < n; ++i) {
    if (scc_id[i] == scc_id[i + n]) return false;
    if (scc_id[i] < scc_id[i + n])
        is[i] = true;
    else
        is[i + n] = true;
}
return true;
}
</pre>
```

4.6 Negative Cycle

```
vector <pair <int, long long>> adj[N];
template <typename T>
struct NegativeCycle {
  vector <T> dis;
  vector <int> rt;
  int n; T INF;
  vector <int> cycle;
  NegativeCycle () = default;
  NegativeCycle (int _n) : n(_n), INF(numeric_limits<T</pre>
       >::max()) {
     dis.assign(n, 0), rt.assign(n, -1);
     int relax = -1;
     for (int t = 0; t < n; ++t) {
      relax = -1;
       for (int i = 0; i < n; ++i) {</pre>
         for (auto [j, w] : adj[i]) if (dis[j] > dis[i]
             + w) {
           dis[j] = dis[i] + w, rt[j] = i;
           relax = j;
         }
      }
     if (relax != -1) {
       int s = relax;
       for (int i = 0; i < n; ++i) s = rt[s];</pre>
       vector <bool> vis(n, false);
       while (!vis[s]) {
         cycle.push_back(s), vis[s] = true;
         s = rt[s];
       reverse(cycle.begin(), cycle.end());
  }
};
```

4.7 Virtual Tree

```
vector<vector<int>> _g;
 vector<int> st, ed, stk;
 void solve(vector<int> v) {
   sort(all(v), [&](int x, int y) {return st[x] < st[y</pre>
       ];});
   int sz = v.size();
   for (int i = 0; i < sz - 1; ++i)
     v.push_back(lca(v[i], v[i + 1]));
   sort(all(v), [&](int x, int y) {return st[x] < st[y</pre>
       ];});
   v.resize(unique(all(v)) - v.begin());
   stk.clear(); stk.push_back(v[0]);
   for (int i = 1; i < v.size(); ++i) {</pre>
     int x = v[i];
     while (ed[stk.back()] < ed[x]) stk.pop_back();</pre>
     g[stk.back()].push_back(x), stk.push_back(x);
   // do something
   for (int i : v) _g[i].clear();
}
```

4.8 Directed MST

```
template <typename T> struct DMST { // 1-based
  T g[maxn][maxn], fw[maxn];
  int n, fr[maxn];
  bool vis[maxn], inc[maxn];
  void clear() {
    for (int i = 0; i < maxn; ++i) {
       for (int j = 0; j < maxn; ++j) g[i][j] = inf;
       vis[i] = inc[i] = false;</pre>
```

```
}
  void addedge(int u, int v, T w) {
    g[u][v] = min(g[u][v], w);
  T query(int root, int _n) {
    n = _n;
    if (dfs(root) != n) return -1;
    T ans = 0:
    while (true) {
      for (int i = 1; i <= n; ++i) fw[i] = inf, fr[i] =</pre>
            i;
      for (int i = 1; i <= n; ++i) if (!inc[i]) {</pre>
           for (int j = 1; j <= n; ++j) {</pre>
             if (!inc[j] && i != j && g[j][i] < fw[i]) {</pre>
               fw[i] = g[j][i];
               fr[i] = j;
           }
        }
      int x = -1;
      for (int i = 1; i <= n; ++i) if (i != root && !</pre>
           inc[i]) {
           int j = i, c = 0;
           while (j != root && fr[j] != i && c <= n) ++c</pre>
                , j = fr[j];
           if (j == root || c > n) continue;
           else { x = i; break; }
        }
      if (!~x) {
         for (int i = 1; i <= n; ++i) if (i != root && !</pre>
             inc[i]) ans += fw[i];
        return ans:
      int y = x;
      for (int i = 1; i <= n; ++i) vis[i] = false;</pre>
      do { ans += fw[y]; y = fr[y]; vis[y] = inc[y] =
           true; } while (y != x);
      inc[x] = false;
      for (int k = 1; k <= n; ++k) if (vis[k]) {</pre>
           for (int j = 1; j <= n; ++j) if (!vis[j]) {</pre>
               if (g[x][j] > g[k][j]) g[x][j] = g[k][j];
               if (g[j][k] < inf && g[j][k] - fw[k] < g[
                    j][x]) g[j][x] = g[j][k] - fw[k];
             }
        }
    }
    return ans;
  int dfs(int now) {
    int r = 1;
    vis[now] = true;
    for (int i = 1; i <= n; ++i) if (g[now][i] < inf &&</pre>
          !vis[i]) r += dfs(i);
    return r;
  }
};
```

4.9 Dominator Tree

```
struct Dominator_tree {
  int n, id;
  vector <vector <int>> adj, radj, bucket;
  vector <int> sdom, dom, vis, rev, par, rt, mn;
 Dominator_tree (int _n) : n(_n), id(0) {
  adj.resize(n), radj.resize(n), bucket.resize(n);
    sdom.resize(n), dom.resize(n, -1), vis.resize(n,
        -1);
    rev.resize(n), rt.resize(n), mn.resize(n), par.
        resize(n);
  }
  void add_edge(int u, int v) {adj[u].pb(v);}
  int query(int v, bool x) {
    if (rt[v] == v) return x ? -1 : v;
    int p = query(rt[v], true);
    if (p == -1) return x ? rt[v] : mn[v];
    if (sdom[mn[v]] > sdom[mn[rt[v]]]) mn[v] = mn[rt[v
        ]];
    rt[v] = p;
    return x ? p : mn[v];
  void dfs(int v) {
```

```
vis[v] = id, rev[id] = v;
    rt[id] = mn[id] = sdom[id] = id, id++;
    for (int u : adj[v]) {
      if (vis[u] == -1) dfs(u), par[vis[u]] = vis[v];
      radj[vis[u]].pb(vis[v]);
  void build(int s) {
    dfs(s);
    for (int i = id - 1; ~i; --i) {
      for (int u : radj[i]) {
        sdom[i] = min(sdom[i], sdom[query(u, false)]);
      if (i) bucket[sdom[i]].pb(i);
      for (int u : bucket[i]) {
        int p = query(u, false);
        dom[u] = sdom[p] == i ? i : p;
      if (i) rt[i] = par[i];
    }
    vector <int> res(n, -1);
    for (int i = 1; i < id; ++i) {</pre>
      if (dom[i] != sdom[i]) dom[i] = dom[dom[i]];
    for (int i = 1; i < id; ++i) res[rev[i]] = rev[dom[</pre>
        i]];
    res[s] = s;
    dom = res;
};
```

5 String

5.1 Aho-Corasick Automaton

```
struct AC {
  int ch[N][26], to[N][26], fail[N], sz;
  vector <int> g[N];
  int cnt[N];
  AC () \{sz = 0, extend();\}
  void extend() {fill(ch[sz], ch[sz] + 26, 0), sz++;}
  int nxt(int u, int v) {
    if (!ch[u][v]) ch[u][v] = sz, extend();
    return ch[u][v];
  int insert(string s) {
    int now = 0;
    for (char c : s) now = nxt(now, c - 'a');
    cnt[now]++;
    return now;
  void build_fail() {
    queue <int> q;
    for (int i = 0; i < 26; ++i) if (ch[0][i]) {</pre>
        q.push(ch[0][i]);
        g[0].push_back(ch[0][i]);
    while (!q.empty()) {
      int v = q.front(); q.pop();
      for (int j = 0; j < 26; ++j) {
        to[v][j] = ch[v][j] ? v : to[fail[v]][j];
      for (int i = 0; i < 26; ++i) if (ch[v][i]) {</pre>
          int u = ch[v][i], k = fail[v];
          while (k \&\& !ch[k][i]) k = fail[k];
          if (ch[k][i]) k = ch[k][i];
          fail[u] = k;
          cnt[u] += cnt[k], g[k].push_back(u);
          q.push(u);
        }
   }
  int match(string &s) {
    int now = 0, ans = 0;
    for (char c : s) {
      now = to[now][c - 'a'];
if (ch[now][c - 'a']) now = ch[now][c - 'a'];
      ans += cnt[now];
    return ans;
```

5.2 KMP Algorithm

|};

```
vector <int> build_fail(string s) {
 vector <int> f(s.length() + 1, 0);
 int k = 0;
  for (int i = 1; i < s.length(); ++i) {</pre>
   while (k \&\& s[k] != s[i]) k = f[k];
    if (s[k] == s[i]) k++;
    f[i + 1] = k;
 }
 return f;
int match(string s, string t) {
 vector <int> f = build_fail(t);
  int k = 0, ans = 0;
  for (int i = 0; i < s.length(); ++i) {</pre>
   while (k \&\& s[i] != t[k]) k = f[k];
   if (s[i] == t[k]) k++;
   if (k == t.length()) ans++, k = f[k];
 return ans;
```

5.3 Z Algorithm

5.4 Manacher

```
vector <int> manacher(string &s) {
   string t = "^#";
   for (char c : s) t += c, t += '#';
   t += '&';
   int n = t.length();
   vector <int> r(n, 0);
   int C = 0, R = 0;
   for (int i = 1; i < n - 1; ++i) {
      int mirror = 2 * C - i;
      r[i] = (i < R ? min(r[mirror], R - i) : 0);
      while (t[i - 1 - r[i]] == t[i + 1 + r[i]]) r[i]++;
      if (i + r[i] > R) R = i + r[i], C = i;
   }
   return r;
}
```

5.5 Suffix Array

```
int sa[N], tmp[2][N], c[N], rk[N], lcp[N];
void buildSA(string s) {
  int *x = tmp[0], *y = tmp[1], m = 256, n = s.length()
  for (int i = 0; i < m; ++i) c[i] = 0;</pre>
  for (int i = 0; i < n; ++i) c[x[i] = s[i]]++;
  for (int i = 1; i < m; ++i) c[i] += c[i - 1];</pre>
  for (int i = n - 1; ~i; --i) sa[--c[x[i]]] = i;
  for (int k = 1; k < n; k <<= 1) {
    for (int i = 0; i < m; ++i) c[i] = 0;</pre>
    for (int i = 0; i < n; ++i) c[x[i]]++;</pre>
    for (int i = 1; i < m; ++i) c[i] += c[i - 1];</pre>
    int p = 0;
    for (int i = n - k; i < n; ++i) y[p++] = i;
    for (int i = 0; i < n; ++i) if (sa[i] >= k) y[p++]
        = sa[i] - k;
    for (int i = n - 1; \sim i; --i) sa[--c[x[y[i]]]] = y[i]
    y[sa[0]] = p = 0;
    for (int i = 1; i < n; ++i) {
      int a = sa[i], b = sa[i - 1];
```

```
if (!(x[a] == x[b] \&\& a + k < n \&\& b + k < n \&\& x)
           [a + k] == x[b + k])) p++;
      y[sa[i]] = p;
    }
    if (n == p + 1) break;
    swap(x, y), m = p + 1;
  }
}
void buildLCP(string s) {
  // lcp[i] = LCP(sa[i - 1], sa[i])
  // lcp(i, j) = min(lcp[rk[i] + 1], lcp[rk[i] + 2],
       ..., lcp[rk[j]])
  int n = s.length(), val = 0;
  for (int i = 0; i < n; ++i) rk[sa[i]] = i;
for (int i = 0; i < n; ++i) {</pre>
    if (!rk[i]) lcp[rk[i]] = 0;
    else {
       if (val) val--;
       int p = sa[rk[i] - 1];
       while (val + i < n && val + p < n && s[val + i]</pre>
           == s[val + p]) val++;
       lcp[rk[i]] = val;
}
```

5.6 SAIS

```
namespace sfx {
bool _t[N * 2];
int SA[N * 2], H[N], RA[N];
int _s[N * 2], _c[N * 2], _x[N], _p[N], _q[N * 2];
void pre(int *sa, int *c, int n, int z) {
 fill_n(sa, n, 0), copy_n(c, z, x);
void induce(int *sa, int *c, int *s, bool *t, int n,
    int z) {
  copy_n(c, z - 1, x + 1);
for (int i = 0; i < n; ++i) if (sa[i] && !t[sa[i] -</pre>
      1]) sa[x[s[sa[i] - 1]]++] = sa[i] - 1;
  copy_n(c, z, x);
  for (int i = n - 1; i >= 0; --i) if (sa[i] && t[sa[i]
        - 1]) sa[--x[s[sa[i] - 1]]] = sa[i] - 1;
void sais(int *s, int *sa, int *p, int *q, bool *t, int
     *c, int n, int z) {
  bool uniq = t[n - 1] = true;
  int nn = 0, nmxz = -1, *nsa = sa + n, *ns = s + n,
      last = -1;
  fill_n(c, z, 0);
  for (int i = 0; i < n; ++i) uniq &= ++c[s[i]] < 2;</pre>
  partial_sum(c, c + z, c);
  if (uniq) {
    for (int i = 0; i < n; ++i) sa[--c[s[i]]] = i;</pre>
    return;
  for (int i = n - 2; i >= 0; --i)
    t[i] = (s[i] == s[i + 1] ? t[i + 1] : s[i] < s[i +
        1]);
  pre(sa, c, n, z);
  for (int i = 1; i <= n - 1; ++i)
    if (t[i] && !t[i - 1])
      sa[--x[s[i]]] = p[q[i] = nn++] = i;
  induce(sa, c, s, t, n, z);
for (int i = 0; i < n; ++i)
    if (sa[i] && t[sa[i]] && !t[sa[i] - 1]) {
      bool neq = last < 0 \mid | !equal(s + sa[i], s + p[q[
           sa[i]] + 1], s + last);
      ns[q[last = sa[i]]] = nmxz += neq;
  sais(ns, nsa, p + nn, q + n, t + n, c + z, nn, nmxz +
       1);
  pre(sa, c, n, z);
  for (int i = nn - 1; i >= 0; --i)
    sa[--x[s[p[nsa[i]]]]] = p[nsa[i]];
  induce(sa, c, s, t, n, z);
vector<int> build(int *s, int n) {
  copy_n(s, n, _s), _s[n] = 0;
  sais(_s, SA, _p, _q, _t, _c, n + 1, 256);
  vector <int> sa(n);
  for (int i = 0; i < n; ++i)</pre>
```

```
sa[i] = SA[i + 1];
return sa;
}
}
```

5.7 Suffix Automaton

```
struct SAM +
  int ch[N][26], len[N], link[N], cnt[N], sz;
  // link -> suffix endpos
  SAM () \{len[0] = 0, link[0] = -1, sz = 1;\}
  void build(string s) {
    int last = 0;
    for (char c : s) {
      int cur = sz++;
      len[cur] = len[last] + 1;
      int p = last;
      while (\sim p \&\& !ch[p][c - 'a']) ch[p][c - 'a'] =
          cur, p = link[p];
      if (p == -1) {
        link[cur] = 0;
      } else {
        int q = ch[p][c - 'a'];
        if (len[p] + 1 == len[q]) {
          link[cur] = q;
        } else {
           int nxt = sz++;
          len[nxt] = len[p] + 1, link[nxt] = link[q];
           for (int j = 0; j < 26; ++j) ch[nxt][j] = ch[</pre>
               al[i];
          while (\sim p && ch[p][c - 'a'] == q) ch[p][c - '
               a'] = nxt, p = link[p];
          link[q] = link[cur] = nxt;
        }
      }
      cnt[cur]++;
      last = cur;
    vector <int> p(sz);
    iota(all(p), 0);
    sort(all(p), [\&](int i, int j) \{return len[i] > len
         [j];});
    for (int i = 0; i < sz; ++i) cnt[link[p[i]]] += cnt</pre>
         [p[i]];
  }
};
```

5.8 Minimum Rotation

```
string rotate(const string &s) {
   int n = s.length();
   string t = s + s;
   int i = 0, j = 1;
   while (i < n && j < n) {
    int k = 0;
   while (k < n && t[i + k] == t[j + k]) ++k;
      if (t[i + k] <= t[j + k]) j += k + 1;
      else i += k + 1;
      if (i == j) ++j;
   }
   int pos = (i < n ? i : j);
   return t.substr(pos, n);
}</pre>
```

5.9 Palindrome Tree

```
struct PAM {
   int ch[N][26], cnt[N], fail[N], len[N], sz;
   string s;
   // 0 -> even root, 1 -> odd root
   PAM (string _s) : s(_s) {
      sz = 0;
      extend(), extend();
   len[0] = 0, fail[0] = 1, len[1] = -1;
   int lst = 1;
   for (int i = 0; i < s.length(); ++i) {
      while (s[i - len[lst] - 1] != s[i]) lst = fail[
            lst];
   if (!ch[lst][s[i] - 'a']) {
      int idx = extend();
      len[idx] = len[lst] + 2;
      int now = fail[lst];</pre>
```

5.10 Main Lorentz

```
int to_left[N], to_right[N];
vector <array <int, 3>> rep; // l, r, len.
// substr(l ~ r, len * 2) are tandem
void findRep(string &s, int 1, int r) {
  if (r - l == 1) return;
  int m = 1 + r >> 1;
  findRep(s, 1, m), findRep(s, m, r);
  string sl = s.substr(1, m - 1), sr = s.substr(m, r - 1)
       m);
  vector <int> Z = buildZ(sr + "#" + sl);
  for (int i = 1; i < m; ++i) to_right[i] = Z[r - m + 1</pre>
        + i - 1];
  reverse(all(sl));
  Z = buildZ(s1);
  for (int i = 1; i < m; ++i) to_left[i] = Z[m - i -</pre>
       1];
  reverse(all(sl));
  for (int i = 1; i + 1 < m; ++i) {</pre>
    int k1 = to_left[i], k2 = to_right[i + 1], len = m
          - i - 1;
     if (k1 < 1 || k2 < 1 || len < 2) continue;</pre>
     int tl = max(1, len - k2), tr = min(len - 1, k1);
     if (tl <= tr) rep.pb({i + 1 - tr, i + 1 - tl, len})</pre>
  Z = buildZ(sr);
  for (int i = m; i < r; ++i) to right[i] = Z[i - m];
  reverse(all(sl)), reverse(all(sr));
Z = buildZ(sl + "#" + sr);
  for (int i = m; i < r; ++i) to_left[i] = Z[m - 1 + 1</pre>
       + r - i - 1];
  reverse(all(sl)), reverse(all(sr));
for (int i = m; i + 1 < r; ++i) {</pre>
     int k1 = to_left[i], k2 = to_right[i + 1], len = i
          - m + 1;
     if (k1 < 1 || k2 < 1 || len < 2) continue;</pre>
     int tl = max(len - k2, 1), tr = min(len - 1, k1);
     if (tl \leftarrow tr) rep.pb({i + 1 - len - tr, i + 1 - len}
           - tl, len});
  Z = buildZ(sr + "#" + sl);
  for (int i = 1; i < m; ++i) {</pre>
    if (Z[r - m + 1 + i - 1] >= m - i) {
       rep.pb({i, i, m - i});
  }
}
```

6 Math

6.1 Fraction

```
struct fraction {
    ll n, d;
    fraction(const ll _n=0, const ll _d=1): n(_n), d(_d)
        {
        ll t = gcd(n, d);
        n /= t, d /= t;
        if (d < 0) n = -n, d = -d;
    }
    fraction operator-() const</pre>
```

```
{ return fraction(-n, d); }
fraction operator+(const fraction &b) const
{ return fraction(n * b.d + b.n * d, d * b.d); }
fraction operator-(const fraction &b) const
{ return fraction(n * b.d - b.n * d, d * b.d); }
fraction operator*(const fraction &b) const
{ return fraction(n * b.n, d * b.d); }
fraction operator/(const fraction &b) const
{ return fraction(n * b.d, d * b.n); }
void print() {
   cout << n;
   if (d != 1) cout << "/" << d;
}
};</pre>
```

6.2 Miller Rabin / Pollard Rho

```
11 mul(11 x, 11 y, 11 p) \{return (x * y - (11))((long x + y - (1
double)x / p * y) * p + p) % p;}
vector<ll> chk = {2, 325, 9375, 28178, 450775, 9780504,
              1795265022};
ll Pow(ll a, ll b, ll n) {ll res = 1; for (; b; b >>=
           1, a = mul(a, a, n)) if (b \& 1) res = mul(res, a, n)
            ); return res;}
bool check(ll a, ll d, int s, ll n) {
      a = Pow(a, d, n);
      if (a <= 1) return 1;</pre>
     for (int i = 0; i < s; ++i, a = mul(a, a, n)) {</pre>
          if (a == 1) return 0;
           if (a == n - 1) return 1;
      return 0;
bool IsPrime(ll n) {
     if (n < 2) return 0;
     if (n % 2 == 0) return n == 2;
11 d = n - 1, s = 0;
      while (d % 2 == 0) d >>= 1, ++s;
      for (ll i : chk) if (!check(i, d, s, n)) return 0;
const vector<ll> small = {2, 3, 5, 7, 11, 13, 17, 19};
11 FindFactor(ll n) {
      if (IsPrime(n)) return 1;
      for (11 p : small) if (n % p == 0) return p;
      11 x, y = 2, d, t = 1;
      auto f = [&](11 a) {return (mul(a, a, n) + t) % n;};
      for (int 1 = 2; ; 1 <<= 1) {
           x = y;
           int m = min(1, 32);
            for (int i = 0; i < 1; i += m) {</pre>
                d = 1;
                 for (int j = 0; j < m; ++j) {</pre>
                      y = f(y), d = mul(d, abs(x - y), n);
                ll g = \_gcd(d, n);
                if (g == n) {
                      1 = 1, y = 2, ++t;
                      break;
                 if (g != 1) return g;
          }
     }
map <ll, int> res;
void PollardRho(ll n) {
     if (n == 1) return;
      if (IsPrime(n)) return ++res[n], void(0);
     11 d = FindFactor(n);
      PollardRho(n / d), PollardRho(d);
}
```

6.3 Ext GCD

```
//a * p.first + b * p.second = gcd(a, b)
pair<11, 11> extgcd(11 a, 11 b) {
   pair<11, 11> res;
   if (a < 0) {
      res = extgcd(-a, b);
      res.first *= -1;
      return res;
   }</pre>
```

```
if (b < 0) {
    res = extgcd(a, -b);
    res.second *= -1;
    return res;
}
if (b == 0) return {1, 0};
res = extgcd(b, a % b);
return {res.second, res.first - res.second * (a / b)
    };
}</pre>
```

6.4 PiCount

```
const int V = 10000000, N = 100, M = 100000;
vector<int> primes;
bool isp[V];
int small_pi[V], dp[N][M];
void sieve(int x){
  for(int i = 2; i < x; ++i) isp[i] = true;</pre>
  isp[0] = isp[1] = false;
  for(int i = 2; i * i < x; ++i) if(isp[i]) for(int j =</pre>
       i * i; j < x; j += i) isp[j] = false;
  for(int i = 2; i < x; ++i) if(isp[i]) primes.
      push back(i);
void init(){
  sieve(V);
  small_pi[0] = 0;
  for(int i = 1; i < V; ++i) small_pi[i] = small_pi[i -</pre>
       1] + isp[i];
  for(int i = 0; i < M; ++i) dp[0][i] = i;</pre>
  for(int i = 1; i < N; ++i) for(int j = 0; j < M; ++j)
       dp[i][j] = dp[i - 1][j] - dp[i - 1][j / primes[i]
11 phi(ll n, int a){
  if(!a) return n;
  if(n < M && a < N) return dp[a][n];</pre>
  if(primes[a - 1] > n) return 1;
  if(((ll)primes[a - 1]) * primes[a - 1] >= n && n < V)</pre>
       return small_pi[n] - a + 1;
  11 de = phi(n, a - 1) - phi(n / primes[a - 1], a - 1)
  return de;
11 PiCount(11 n){
  if(n < V) return small_pi[n];</pre>
  int s = sqrt(n + 0.5), y = cbrt(n + 0.5), a =
      small_pi[y];
  ll res = phi(n, a) + a - 1;
  for(; primes[a] <= s; ++a) res -= max(PiCount(n /</pre>
      primes[a]) - PiCount(primes[a]) + 1, 0ll);
  return res;
```

6.5 Linear Function Mod Min

```
ll topos(ll x, ll m) {x %= m; if (x < 0) x += m; return
     x;}
//min value of ax + b \pmod{m} for x \in [0, n - 1]. O(
    Log m)
11 min_rem(ll n, ll m, ll a, ll b) {
  a = topos(a, m), b = topos(b, m);
  for (ll g = __gcd(a, m); g > 1;) return g * min_rem(n
  , m / g, a / g, b / g) + (b % g); for (11 nn, nm, na, nb; a; n = nn, m = nm, a = na, b
       = nb) {
    if (a <= m - a) {
      nn = (a * (n' - 1) + b) / m;
      if (!nn) break;
      nn += (b < a);
      nm = a, na = topos(-m, a);
      nb = b < a ? b : topos(b - m, a);
    } else {
      11 lst = b - (n - 1) * (m - a);
      if (lst >= 0) {b = lst; break;}
      nn = -(1st / m) + (1st % m < -a) + 1;
      nm = m - a, na = m % (m - a), nb = b % (m - a);
    }
  return b;
```

```
//min value of ax + b \pmod{m} for x \in [0, n - 1],
    also return min x to get the value. O(\log m)
//{value, x}
pair<ll, 11> min_rem_pos(11 n, 11 m, 11 a, 11 b) {
  a = topos(a, m), b = topos(b, m);
 11 mn = min_rem(n, m, a, b), g = __gcd(a, m);
  //ax = (mn - b) \pmod{m}
  11 x = (extgcd(a, m).first + m) * ((mn - b + m) / g)
     % (m / g);
  return {mn, x};
```

6.6 Floor Sum

```
// sum^{n-1}_0 floor((a * i + b) / m) in log(n + m + a
11 floor_sum(ll n, ll m, ll a, ll b) {
  11 \text{ ans} = 0;
  if (a >= m) ans += (n - 1) * n * (a / m) / 2, a %= m;
  if (b >= m) ans += n * (b / m), b %= m;
11 y_max = (a * n + b) / m, x_max = (y_max * m - b);
  if (y_max == 0) return ans;
  ans += (n - (x_max + a - 1) / a) * y_max;
  ans += floor_sum(y_max, a, m, (a - x_max % a) % a);
  return ans:
```

6.7 Quadratic Residue

```
int Jacobi(int a, int m) {
 int s = 1;
  for (; m > 1; ) {
   a %= m;
    if (a == 0) return 0;
    const int r = __builtin_ctz(a);
   if ((r \& 1) \& \& ((m + 2) \& 4)) s = -s;
    a >>= r;
   if (a \& m \& 2) s = -s;
    swap(a, m);
 return s:
int QuadraticResidue(int a, int p) {
 if (p == 2) return a & 1;
  const int jc = Jacobi(a, p);
  if (jc == 0) return 0;
 if (jc == -1) return -1;
  int b, d;
 for (; ; ) {
   b = rand() \% p;
    d = (1LL * b * b + p - a) \% p;
   if (Jacobi(d, p) == -1) break;
 int f0 = b, f1 = 1, g0 = 1, g1 = 0, tmp;
  for (int e = (1LL + p) >> 1; e; e >>= 1) {
    if (e & 1) {
      tmp = (1LL * g0 * f0 + 1LL * d * (1LL * g1 * f1 %
          p)) % p;
      g1 = (1LL * g0 * f1 + 1LL * g1 * f0) % p;
      g0 = tmp;
   tmp = (1LL * f0 * f0 + 1LL * d * (1LL * f1 * f1 % p)
   )) % p;
f1 = (2LL * f0 * f1) % p;
    f0 = tmp;
 }
 return g0;
```

6.8 Simplex

```
struct Simplex { // 0-based
 using T = long double;
  static const int N = 410, M = 30010;
  const T eps = 1e-7;
  int n, m;
  int Left[M], Down[N];
  // Ax <= b, max c^T x
// result : v, xi = sol[i]
  T a[M][N], b[M], c[N], v, sol[N];
  bool eq(T a, T b) {return fabs(a - b) < eps;}
bool ls(T a, T b) {return a < b && !eq(a, b);}</pre>
```

```
void init(int _n, int _m) {
      n = _n, m = _m, v = \overline{0};
for (int i = 0; i < m; ++i) for (int j = 0; j < n;
           ++j) a[i][j] = 0;
      for (int i = 0; i < m; ++i) b[i] = 0;</pre>
      for (int i = 0; i < n; ++i) c[i] = sol[i] = 0;</pre>
   void pivot(int x, int y) {
      swap(Left[x], Down[y]);
      T k = a[x][y]; a[x][y] = 1;
      vector <int> nz;
      for (int i = 0; i < n; ++i) {
        a[x][i] /= k;
        if (!eq(a[x][i], 0)) nz.push_back(i);
      b[x] /= k;
      for (int i = 0; i < m; ++i) {
  if (i == x || eq(a[i][y], 0)) continue;</pre>
        k = a[i][y], a[i][y] = 0;
b[i] -= k * b[x];
        for (int j : nz) a[i][j] -= k * a[x][j];
      if (eq(c[y], 0)) return;
      k = c[y], c[y] = 0, v += k * b[x];
      for (int i : nz) c[i] -= k * a[x][i];
   // 0: found solution, 1: no feasible solution, 2:
        unbounded
   int solve() {
      for (int i = 0; i < n; ++i) Down[i] = i;
for (int i = 0; i < m; ++i) Left[i] = n + i;</pre>
      while (1) {
        int x = -1, y = -1;
for (int i = 0; i < m; ++i) if (ls(b[i], 0) && (x
              == -1 \mid \mid b[i] < b[x]) x = i;
        if (x == -1) break;
        for (int i = 0; i < n; ++i) if (ls(a[x][i], 0) &&</pre>
              (y == -1 \mid | a[x][i] < a[x][y])) y = i;
        if (y == -1) return 1;
        pivot(x, y);
      while (1) {
        int x = -1, y = -1;
        for (int i = 0; i < n; ++i) if (ls(0, c[i]) && (y
              == -1 \mid \mid c[i] > c[y])) y = i;
        if (y == -1) break;
        for (int i = 0; i < m; ++i) if (ls(0, a[i][y]) &&</pre>
              (x == -1 \mid | b[i] / a[i][y] < b[x] / a[x][y])
             ) x = i;
        if (x == -1) return 2;
        pivot(x, y);
      for (int i = 0; i < m; ++i) if (Left[i] < n) sol[</pre>
           Left[i]] = b[i];
      return 0;
};
```

Linear Programming Construction

Standard form: maximize $\mathbf{c}^T\mathbf{x}$ subject to $A\mathbf{x} \leq \mathbf{b}$ and $\mathbf{x} \geq 0$. Dual LP: minimize $\mathbf{b}^T\mathbf{y}$ subject to $A^T\mathbf{y} \geq \mathbf{c}$ and $\mathbf{y} \geq \mathbf{0}$. $\bar{\mathbf{x}}$ and $\bar{\mathbf{y}}$ are optimal if and only if for all $i \in [1,n]$, either $\bar{x}_i = 0$ or $\sum_{j=1}^m A_{ji}\bar{y}_j = c_i$ holds and for all $i \in [1,m]$ either $\bar{y}_i = 0$ or $\sum_{j=1}^{n} A_{ij} \bar{x}_j = b_j$ holds.

- 1. In case of minimization, let $c_i'=-c_i$ 2. $\sum_{1\leq i\leq n}A_{ji}x_i\geq b_j\to \sum_{1\leq i\leq n}-A_{ji}x_i\leq -b_j$ $3. \sum_{1 \le i \le n} A_{ji} x_i = b_j$
 - $\begin{array}{ll} \bullet & \sum_{1 \leq i \leq n} A_{ji} x_i \leq b_j \\ \bullet & \sum_{1 \leq i \leq n} A_{ji} x_i \geq b_j \end{array}$
- 4. If x_i has no lower bound, replace x_i with $x_i x_i'$

6.10 Euclidean

• $m = \lfloor \frac{an+b}{c} \rfloor$ • Time complexity: $O(\log n)$

$$\begin{split} f(a,b,c,n) &= \sum_{i=0}^n \lfloor \frac{ai+b}{c} \rfloor \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)}{2} + \lfloor \frac{b}{c} \rfloor \cdot (n+1) \\ + f(a \mod c, b \mod c, c, n), & a \geq c \vee b \geq c \\ 0, & n < 0 \vee a = 0 \\ nm - f(c, c - b - 1, a, m - 1), & \text{otherwise} \end{cases} \end{split}$$

$$\begin{split} g(a,b,c,n) &= \sum_{i=0}^n i \lfloor \frac{ai+b}{c} \rfloor \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor \cdot \frac{n(n+1)}{2} \\ +g(a \bmod c, b \bmod c, c, n), & a \geq c \lor b \geq c \\ 0, & n < 0 \lor a = 0 \\ \frac{1}{2} \cdot (n(n+1)m - f(c, c-b-1, a, m-1) \\ -h(c, c-b-1, a, m-1)), & \text{otherwise} \end{cases} \end{split}$$

$$\begin{split} h(a,b,c,n) &= \sum_{i=0}^n \lfloor \frac{ai+b}{c} \rfloor^2 \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor^2 \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor^2 \cdot (n+1) \\ + \lfloor \frac{a}{c} \rfloor \cdot \lfloor \frac{b}{c} \rfloor \cdot n(n+1) \\ + h(a \bmod c, b \bmod c, c, n) \\ + 2 \lfloor \frac{a}{c} \rfloor \cdot g(a \bmod c, b \bmod c, c, n) \\ + 2 \lfloor \frac{b}{c} \rfloor \cdot f(a \bmod c, b \bmod c, c, n), & a \geq c \lor b \geq c \\ 0, & n < 0 \lor a = 0 \\ nm(m+1) - 2g(c, c-b-1, a, m-1) \\ - 2f(c, c-b-1, a, m-1) - f(a, b, c, n), & \text{otherwise} \end{cases} \end{split}$$

6.11 Theorem

• Kirchhoff's Theorem

Denote L be a $n\times n$ matrix as the Laplacian matrix of graph G, where $L_{ii}=d(i)$, $L_{ij}=-c$ where c is the number of edge (i,j) in G.

- The number of undirected spanning in G is $|\det(\tilde{L}_{11})|$.
- The number of directed spanning tree rooted at r in G is $|{\rm det}(\tilde{L}_{rr})|$.
- Tutte's Matrix

Let D be a $n \times n$ matrix, where $d_{ij} = x_{ij}$ (x_{ij} is chosen uniformly at random) if i < j and $(i,j) \in E$, otherwise $d_{ij} = -d_{ji}$. $\frac{rank(D)}{2}$ is the maximum matching on G.

- Cayley's Formula
 - Given a degree sequence d_1,d_2,\ldots,d_n for each *labeled* vertices, there are

$$\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$$

spanning trees.

- Let $T_{n,k}$ be the number of *labeled* forests on n vertices with k components, such that vertex $1,2,\ldots,k$ belong to different components. Then $T_{n,k}=kn^{n-k-1}$.
- Erdős-Gallai Theorem

A sequence of non-negative integers $d_1 \geq d_2 \geq \ldots \geq d_n$ can be represented as the degree sequence of a finite simple graph on n vertices if and only if $d_1+d_2+\ldots+d_n$ is even and

$$\sum_{i=1}^k d_i \leq k(k-1) + \sum_{i=k+1}^n \min(d_i,k)$$

holds for all $1 \leq k \leq n$.

Burnside's Lemma

Let X be a set and G be a group that acts on X . For $g\in G$, denote by X^g the elements fixed by g:

$$X^g = \{x \in X \mid gx \in X\}$$

Then

$$|X/G| = \frac{1}{|G|} \sum_{g \in G} |X^g|.$$

• Gale-Ryser theorem

A pair of sequences of nonnegative integers $a_1\geq\cdots\geq a_n$ and b_1,\ldots,b_n is bigraphic if and only if $\sum_{i=1}^n a_i=\sum_{i=1}^n b_i$ and $\sum_{i=1}^k a_i\leq\sum_{i=1}^n \min(b_i,k)$ holds for every $1\leq k\leq n$.

• Fulkerson-Chen-Anstee theorem

A sequence $(a_1,b_1),\ldots,(a_n,b_n)$ of nonnegative integer pairs with $a_1\geq\cdots\geq a_n$ is digraphic if and only if $\sum_{i=1}^n a_i=\sum_{i=1}^n b_i$ and $\sum_{i=1}^k a_i\leq\sum_{i=1}^k \min(b_i,k-1)+\sum_{i=k+1}^n \min(b_i,k)$ holds for every $1\leq k\leq n$.

• Möbius inversion formula

```
- f(n) = \sum_{d|n} g(d) \Leftrightarrow g(n) = \sum_{d|n} \mu(d) f(\frac{n}{d})

- f(n) = \sum_{n|d} g(d) \Leftrightarrow g(n) = \sum_{n|d} \mu(\frac{d}{n}) f(d)
```

• Spherical cap

```
– A portion of a sphere cut off by a plane. – r: sphere radius, a: radius of the base of the cap, h: height of the cap, \theta: \arcsin(a/r). – Volume = \pi h^2(3r-h)/3 = \pi h(3a^2+h^2)/6 = \pi r^3(2+\cos\theta)(1-\cos\theta)^2/3. – Area = 2\pi rh = \pi(a^2+h^2) = 2\pi r^2(1-\cos\theta).
```

• Chinese Remainder Theorem

```
- x\equiv a_i\pmod{m_i}

- M=\prod m_i, M_i=M/m_i

- t_iM_i\equiv 1\pmod{m_i}

- x=\sum a_it_iM_i\pmod{M}
```

7 Geometry

7.1 Basic

```
const double eps = 1e-8, pi = acos(-1);
int sign(double x) \{return abs(x) \leftarrow eps ? 0 : (x > 0 ?
     1 : -1);}
struct Pt {
  double x, y;
  Pt (double _x, double _y) : x(_x), y(_y) {}
  Pt operator + (Pt o) {return Pt(x + o.x, y + o.y);}
  Pt operator - (Pt o) {return Pt(x - o.x, y - o.y);}
  Pt operator * (double k) {return Pt(x * k, y * k);}
  Pt operator / (double k) {return Pt (x / k, y / k);}
  double operator * (Pt o) {return x * o.x + y * o.y;}
  double operator ^ (Pt o) {return x * o.y - y * o.x;}
struct Line {
 Pt a, b;
};
struct Cir {
  Pt o; double r;
double abs2(Pt o) {return o.x * o.x + o.y * o.y;}
double abs(Pt o) {return sqrt(abs2(o));}
int ori(Pt o, Pt a, Pt b) {return sign((o - a) ^ (o - b
    ));}
bool btw(Pt a, Pt b, Pt c) { // c on segment ab?
  return ori(a, b, c) == 0 && sign((c - a) * (c - b))
double area(Pt a, Pt b, Pt c) {return abs((a - b) ^ (a
    - c)) / 2;}
Pt unit(Pt o) {return o / abs(o);}
Pt rot(Pt a, double o) { // CCW
  double c = cos(o), s = sin(o);
  return Pt(c * a.x - s * a.y, s * a.x + c * a.y);
Pt proj_vector(Pt a, Pt b, Pt c) { // vector ac proj to
  return (b - a) * ((c - a) * (b - a)) / ((b - a) * (b
      - a));
Pt proj_pt(Pt a, Pt b, Pt c) { // point c proj to ab
 return proj_vector(a, b, c) + a;
```

7.2 Heart

7.3 External Bisector

```
Pt external_bisector(Pt p1, Pt p2, Pt p3) { //213
Pt L1 = p2 - p1, L2 = p3 - p1;
L2 = L2 * abs(L1) / abs(L2);
return L1 + L2;
}
```

7.4 Intersection of Segments

```
Pt LinesInter(Line a, Line b) {
    double abc = (a.b - a.a) ^ (b.a - a.a);
    double abd = (a.b - a.a) ^ (b.b - a.a);
    if (sign(abc - abd) == 0) return b.b;// no inter
    return (b.b * abc - b.a * abd) / (abc - abd);
}

vector<Pt> SegsInter(Line a, Line b) {
    if (btw(a.a, a.b, b.a)) return {b.a};
    if (btw(a.a, a.b, b.b)) return {b.b};
    if (btw(b.a, b.b, a.a)) return {a.a};
    if (btw(b.a, b.b, a.b)) return {a.b};
    if (ori(a.a, a.b, b.a) * ori(a.a, a.b, b.b) == -1 &&
        ori(b.a, b.b, a.a) * ori(b.a, b.b, a.b) == -1)
    return {LinesInter(a, b)};
    return {};
}
```

7.5 Intersection of Circle and Line

7.6 Intersection of Circles

7.7 Intersection of Polygon and Circle

```
double _area(Pt pa, Pt pb, double r){
  if(abs(pa) < abs(pb)) swap(pa, pb);
  if(abs(pb) < eps) return 0;
  double S, h, theta;
  double a = abs(pb), b = abs(pa), c = abs(pb - pa);
  double cosB = pb * (pb - pa) / a / c, B = acos(cosB);
  double cosC = (pa * pb) / a / b, C = acos(cosC);
  if (a > r) {
    S = (C / 2) * r * r;
    h = a * b * sin(C) / c;
```

7.8 Tangent Lines of Circle and Point

7.9 Tangent Lines of Circles

```
vector<Line> tangent(Cir a, Cir b) {
#define Pij \
 Pt i = unit(b.o - a.o) * a.r, j = Pt(i.y, -i.x);\
  z.push_back({a.o + i, a.o + i + j});
#define deo(I,J) \
  double d = abs(a.o - b.o), e = a.r I b.r, o = acos(e
      / d);\
  Pt i = unit(b.o - a.o), j = rot(i, o), k = rot(i, -o)
  z.push\_back({a.o + j * a.r, b.o J j * b.r});\
  z.push_back({a.o + k * a.r, b.o J k * b.r});
  if (a.r < b.r) swap(a, b);</pre>
  vector<Line> z;
  if (abs(a.o - b.o) + b.r < a.r) return z;</pre>
  else if (sign(abs(a.o - b.o) + b.r - a.r) == 0) { Pij
  else {
    deo(-,+); // inter
    // outer
    if (sign(d - a.r - b.r) == 0) { Pij; }
    else if (d > a.r + b.r) { deo(+,-); }
  return z;
```

7.10 Point In Convex

7.11 Point Segment Distance

```
double PointSegDist(Pt q0, Pt q1, Pt p) {
  if (sign(abs(q0 - q1)) == 0) return abs(q0 - p);
  if (sign((q1 - q0) * (p - q0)) >= 0 && sign((q0 - q1)
          * (p - q1)) >= 0)
    return fabs(((q1 - q0) ^ (p - q0)) / abs(q0 - q1));
  return min(abs(p - q0), abs(p - q1));
}
```

7.12 Convex Hull

7.13 Convex Hull Distance

7.14 Minimum Enclosing Circle

```
Cir min_enclosing(vector<Pt> &p) {
  random_shuffle(p.begin(), p.end());
  double r = 0.0;
  Pt cent = p[0];
  for (int i = 1; i < p.size(); ++i) {</pre>
   if (abs2(cent - p[i]) <= r) continue;</pre>
    cent = p[i];
    r = 0.0;
    for (int j = 0; j < i; ++j) {
      if (abs2(cent - p[j]) <= r) continue;</pre>
      cent = (p[i] + p[j]) / 2;
      r = abs2(p[j] - cent);
      for (int k = 0; k < j; ++k) {
        if (abs2(cent - p[k]) <= r) continue;</pre>
        cent = circenter(p[i], p[j], p[k]);
        r = abs2(p[k] - cent);
   }
  return {cent, sqrt(r)};
```

7.15 Union of Circles

```
if (z < 0) z += 2 * pi;
    double l = z - o, r = z + o;
    if (1 < 0) 1 += 2 * pi;</pre>
    if (r > 2 * pi) r -= 2 * pi;
    if (1 > r) res.emplace_back(1, 2 * pi), res.
        emplace_back(0, r);
    else res.emplace_back(l, r);
  }
  return res:
double CircleUnionArea(vector<Cir> c) { // circle
    should be identical
  int n = c.size();
  double a = 0, w;
  for (int i = 0; w = 0, i < n; ++i) {
    vector<pair<double, double>> s = {{2 * pi, 9}}, z;
    for (int j = 0; j < n; ++j) if (i != j) {</pre>
      z = CoverSegment(c[i], c[j]);
      for (auto &e : z) s.push_back(e);
    sort(s.begin(), s.end());
    auto F = [&] (double t) { return c[i].r * (c[i].r *
         t + c[i].o.x * sin(t) - c[i].o.y * cos(t)); };
    for (auto &e : s) {
      if (e.first > w) a += F(e.first) - F(w);
      w = max(w, e.second);
   }
 }
  return a * 0.5;
```

7.16 Polar Angle Sort

7.17 Rotating Caliper

7.18 Rotating SweepLine

```
void RotatingSweepLine(vector <Pt> &pt) {
  int n = pt.size();
  vector <int> id(n), pos(n);
  vector <pair <int, int>> line;
  for (int i = 0; i < n; ++i) for (int j = 0; j < n; ++
      j) if (i ^ j) line.emplace_back(i, j);
  sort(line.begin(), line.end(), [&](pair <int, int> i,
       pair <int, int> j) {
    Pt a = pt[i.second] - pt[i.first], b = pt[j.second]
         - pt[j.first];
    return (a.pos() == b.pos() ? sign(a ^ b) > 0 : a.
        pos() < b.pos());
  });
  iota(id.begin(), id.end(), 0);
  sort(id.begin(), id.end(), [&](int i, int j) {
    return (sign(pt[i].y - pt[j].y) == 0 ? pt[i].x < pt</pre>
        [j].x : pt[i].y < pt[j].y);
  for (int i = 0; i < n; ++i)</pre>
    pos[id[i]] = i;
  for (auto [i, j] : line) {
```

```
// point sort by the distance to line(i, j)
// do something.
tie(pos[i], pos[j], id[pos[i]], id[pos[j]]) =
          make_tuple(pos[j], pos[i], j, i);
}
}
```

7.19 Half Plane Intersection

```
vector <Pt> HalfPlaneInter(vector <pair <Pt, Pt>> vec)
  // first -----> second
  auto pos = [&](Pt a) {return sign(a.y) == 0 ? sign(a
      .x) < 0 : sign(a.y) > 0;};
  sort(all(vec), [&](pair <Pt, Pt> a, pair <Pt, Pt> b)
    Pt A = a.second - a.first, B = b.second - b.first;
    if (pos(A) == pos(B)) {
      if (sign(A ^ B) == 0) return sign((b.first - a.
          first) * (b.second - a.first)) > 0;
      return sign(A ^ B) > 0;
    }
    return pos(A) < pos(B);</pre>
  });
  deque <Pt> inter;
  deque <pair <Pt, Pt>> seg;
  int n = vec.size();
  auto get = [&](pair <Pt, Pt> a, pair <Pt, Pt> b) {
      return intersect(a.first, a.second, b.first, b.
      second);};
  for (int i = 0; i < n; ++i) if (!i || vec[i] != vec[i</pre>
        - 1]) {
    while (seg.size() >= 2 && sign((vec[i].second -
        inter.back()) ^ (vec[i].first - inter.back()))
    == 1) seg.pop_back(), inter.pop_back();
while (seg.size() >= 2 && sign((vec[i].second -
        inter.front()) ^ (vec[i].first - inter.front())
        ) == 1) seg.pop_front(), inter.pop_front();
    seg.push_back(vec[i]);
    if (seg.size() >= 2) inter.pb(get(seg[seg.size() -
        2], seg.back()));
 while (seg.size() >= 2 && sign((seg.front().second -
      inter.back()) ^ (seg.front().first - inter.back()
      )) == 1) seg.pop_back(), inter.pop_back();
  inter.push_back(get(seg.front(), seg.back()));
  return vector <Pt>(all(inter));
```

7.20 Minkowski Sum

8 Polynomial

8.1 Number Theoretic Transform

```
const int N = 1 << 18, mod = 998244353, G = 3;
struct NTT {
    11 w[N];
    NTT() {
        11 dw = mpow(G, (mod - 1) / N);
        w[0] = 1;
    for (int i = 1; i < N; ++i) w[i] = w[i - 1] * dw %
        mod;</pre>
```

```
void bitrev(vector<11>& a, int n) {
    int i = 0;
    for (int j = 1; j < n - 1; ++j) {</pre>
      for (int k = n >> 1; (i ^{-} k) < k; k >>= 1);
      if (j < i) swap(a[i], a[j]);</pre>
  }
  void operator()(vector<ll>& a, int n, bool inv =
      false) { //0 <= a[i] < P
    bitrev(a, n);
    for (int L = 2; L <= n; L <<= 1) {</pre>
      int dx = N / L, dl = L >> 1;
      for (int i = 0; i < n; i += L) {</pre>
        for (int j = i, x = 0; j < i + dl; ++j, x += dx
          ll tmp = mul(a[j + dl], w[x]);
          a[j + dl] = sub(a[j], tmp);
          a[j] = add(a[j], tmp);
        }
      }
    if (inv) {
      reverse(a.begin() + 1, a.end());
      ll invn = mpow(n, mod - 2);
      for (int i = 0; i < n; ++i) a[i] = mul(a[i], invn</pre>
 }
};
vector<ll> mul(vector<ll> a, vector<ll> b, int M = N /
    2){
  int m = a.size() + b.size() - 1, n = 1;
  while(n < m) n <<= 1;</pre>
  a.resize(n), b.resize(n);
  ntt(a, n), ntt(b, n);
  for(int i = 0; i < n; ++i) a[i] = mul(a[i], b[i]);</pre>
  ntt(a, n, 1);
  a.resize(min(m, M));
  return a;
```

8.2 Primes

```
Prime
             Root
                     Prime
                                   Root
                     167772161
7681
             17
                     104857601
12289
             11
40961
                     985661441
65537
                     998244353
786433
                     1107296257
             10
                                   10
5767169
             3
                     2013265921
                                   31
7340033
                     2810183681
                                   11
23068673
                     2885681153
             3
469762049
                     605028353
```

8.3 Inverse of Polynomial

```
vector<ll> inv(vector<ll> a){
  int m = a.size();
  vector<ll> res(1, modpow(a[0], mod - 2));
  for(int n = 2; n / 2 < m; n <<= 1){</pre>
    if(a.size() < n) a.resize(n);</pre>
    vector<ll> v1(a.begin(), a.begin() + n), v2 = res;
    v1.resize(n * 2), v2.resize(n * 2);
    ntt(v1, n * 2), ntt(v2, n * 2);
    for(int i = 0; i < n * 2; ++i) v1[i] = mul(mul(v1[i</pre>
         ], v2[i]), v2[i]);
    ntt(v1, n * 2, 1);
    vector<ll> nres(n);
    for(int i = 0; i < n / 2; ++i) nres[i] = add(res[i</pre>
         ], res[i]);
    for(int i = 0; i < n; ++i) nres[i] = sub(nres[i],</pre>
         v1[i]);
    res = nres;
  res.resize(m);
  return res;
```

8.4 Fast Walsh Transform

```
void fwt(vector <int> &a) {
   // and : a[j] += x;
   // : a[j] -= x;
```

9 Else

9.1 Bit Hack

9.2 Dynamic Programming Condition

9.2.1 Totally Monotone (Concave/Convex)

```
\begin{array}{l} \forall i < i', j < j' \text{, } B[i][j] \leq B[i'][j] \implies B[i][j'] \leq B[i'][j'] \\ \forall i < i', j < j' \text{, } B[i][j] \geq B[i'][j] \implies B[i][j'] \geq B[i'][j'] \end{array}
```

9.2.2 Monge Condition (Concave/Convex)

```
\begin{array}{l} \forall i < i', j < j' \text{, } B[i][j] + B[i'][j'] \geq B[i][j'] + B[i'][j] \\ \forall i < i', j < j' \text{, } B[i][j] + B[i'][j'] \leq B[i][j'] + B[i'][j] \end{array}
```

9.2.3 Optimal Split Point

```
If B[i][j] + B[i+1][j+1] \geq B[i][j+1] + B[i+1][j] then H_{i,j-1} \leq H_{i,j} \leq H_{i+1,j}
```

9.3 Slope Trick

```
template<typename T>
struct slope_trick_convex {
  T minn = 0, ground_1 = 0, ground_r = 0;
  priority_queue<T, vector<T>, less<T>> left;
priority_queue<T, vector<T>, greater<T>> right;
  slope_trick_convex() {left.push(numeric_limits<T>::
      min() / 2), right.push(numeric_limits<T>::max() /
  void push_left(T x) {left.push(x - ground_1);}
  void push_right(T x) {right.push(x - ground_r);}
  //add a line with slope 1 to the right starting from
  void add_right(T x) {
    T l = left.top() + ground_l;
    if (1 <= x) push_right(x);</pre>
    else push_left(x), push_right(1), left.pop(), minn
        += 1 - x;
  //add a line with slope -1 to the left starting from
  void add_left(T x) {
    T r = right.top() + ground_r;
    if (r >= x) push_left(x);
    else push_right(x), push_left(r), right.pop(), minn
  //val[i]=min(val[j]) for all i-l<=j<=i+r
  void expand(T 1, T r) {ground_1 -= 1, ground_r += r;}
  void shift_up(T x) {minn += x;}
  T get_val(T x) {
    T l = left.top() + ground_l, r = right.top() +
        ground_r;
    if (x >= 1 & x <= r) return minn;
    if (x < 1) {
```

```
vector<T> trash;
       T cur_val = minn, slope = 1, res;
       while (1) {
         trash.push_back(left.top());
         left.pop();
         if (left.top() + ground_l <= x) {</pre>
           res = cur_val + slope * (1 - x);
           break;
        }
         cur_val += slope * (1 - (left.top() + ground_1)
         1 = left.top() + ground_l;
         slope += 1;
       for (auto i : trash) left.push(i);
       return res;
     if (x > r) {
       vector<T> trash;
       T cur_val = minn, slope = 1, res;
       while (1) {
        trash.push_back(right.top());
         right.pop();
         if (right.top() + ground_r >= x) {
           res = cur_val + slope * (x - r);
           break:
         cur_val += slope * ((right.top() + ground_r) -
         r = right.top() + ground_r;
        slope += 1;
       for (auto i : trash) right.push(i);
       return res;
    assert(0);
};
```

9.4 Manhattan MST

```
void solve(int n) {
  init();
  vector<int> v(n), ds;
  for (int i = 0; i < n; ++i) {</pre>
    v[i] = i;
    ds.push_back(x[i] - y[i]);
  sort(ds.begin(), ds.end());
  ds.resize(unique(ds.begin(), ds.end()) - ds.begin());
  sort(v.begin(), v.end(), [\&](int i, int j) { return x}
      [i] == x[j] ? y[i] > y[j] : x[i] > x[j]; });
  int j = 0;
  for (int i = 0; i < n; ++i) {</pre>
    int p = lower_bound(ds.begin(), ds.end(), x[v[i]] -
         y[v[i]]) - ds.begin() + 1;
    pair<int, int> q = query(p);
    // query return prefix minimum
    if (~q.second) add_edge(v[i], q.second);
    add(p,\ make\_pair(x[v[i]]\ +\ y[v[i]],\ v[i]));
void make_graph() {
  solve(n);
  for (int i = 0; i < n; ++i) swap(x[i], y[i]);</pre>
  solve(n);
  for (int i = 0; i < n; ++i) x[i] = -x[i];
  solve(n);
  for (int i = 0; i < n; ++i) swap(x[i], y[i]);</pre>
  solve(n);
```

9.5 Dynamic MST

```
void contract(int 1, int r, vector<int> v, vector<int>
    &x, vector<int> &y) {
  sort(v.begin(), v.end(), [&](int i, int j) {
      if (cost[i] == cost[j]) return i < j;</pre>
      return cost[i] < cost[j];</pre>
      });
  djs.save();
  for (int i = 1; i <= r; ++i) djs.merge(st[qr[i].first</pre>
      ], ed[qr[i].first]);
  for (int i = 0; i < (int)v.size(); ++i) {</pre>
    if (djs.find(st[v[i]]) != djs.find(ed[v[i]])) {
      x.push back(v[i]);
      djs.merge(st[v[i]], ed[v[i]]);
   }
  djs.undo();
  djs.save();
  for (int i = 0; i < (int)x.size(); ++i) djs.merge(st[</pre>
      x[i]], ed[x[i]]);
  for (int i = 0; i < (int)v.size(); ++i) {</pre>
    if (djs.find(st[v[i]]) != djs.find(ed[v[i]])) {
      y.push_back(v[i]);
      djs.merge(st[v[i]], ed[v[i]]);
  djs.undo();
}
void solve(int 1, int r, vector<int> v, long long c) {
  if (1 == r) {
    cost[qr[l].first] = qr[l].second;
    if (st[qr[1].first] == ed[qr[1].first]) {
      printf("%lld\n", c);
      return;
    int minv = qr[1].second;
    for (int i = 0; i < (int)v.size(); ++i) minv = min(</pre>
        minv, cost[v[i]]);
    printf("%lld\n", c + minv);
    return;
  int m = (1 + r) >> 1;
  vector<int> lv = v, rv = v;
  vector<int> x, y;
  for (int i = m + 1; i <= r; ++i) {</pre>
    cnt[qr[i].first]--;
    if (cnt[qr[i].first] == 0) lv.push_back(qr[i].first
  contract(l, m, lv, x, y);
  long long lc = c, rc = c;
  djs.save();
  for (int i = 0; i < (int)x.size(); ++i) {</pre>
    lc += cost[x[i]];
    djs.merge(st[x[i]], ed[x[i]]);
  solve(1, m, y, 1c);
  djs.undo();
  x.clear(), y.clear();
  for (int i = m + 1; i <= r; ++i) cnt[qr[i].first]++;</pre>
  for (int i = 1; i <= m; ++i) {
    cnt[qr[i].first]--;
    if (cnt[qr[i].first] == 0) rv.push_back(qr[i].first
        );
  contract(m + 1, r, rv, x, y);
  djs.save();
  for (int i = 0; i < (int)x.size(); ++i) {</pre>
   rc += cost[x[i]];
    djs.merge(st[x[i]], ed[x[i]]);
  solve(m + 1, r, y, rc);
  djs.undo();
  for (int i = 1; i <= m; ++i) cnt[qr[i].first]++;</pre>
```

9.6 ALL LCS

```
void all_lcs(string s, string t) { // 0-base
vector<int> h(t.size());
iota(all(h), 0);
for (int a = 0; a < s.size(); ++a) {</pre>
```

```
int v = -1;
for (int c = 0; c < t.size(); ++c)
   if (s[a] == t[c] || h[c] < v)
       swap(h[c], v);
// LCS(s[0, a], t[b, c]) =
   // c - b + 1 - sum([h[i] >= b] | i <= c)
   // h[i] might become -1 !!
}</pre>
```

9.7 Hilbert Curve

```
long long hilbertOrder(int x, int y, int pow, int
    rotate) {
  if (pow == 0) return 0;
  int hpow = 1 << (pow-1);</pre>
  int seg = (x < hpow) ? ((y < hpow) ? 0 : 3) : ((y < hpow)
      hpow) ? 1 : 2);
  seg = (seg + rotate) & 3;
  const int rotateDelta[4] = {3, 0, 0, 1};
  int nx = x & (x ^ hpow), ny = y & (y ^ hpow);
  int nrot = (rotate + rotateDelta[seg]) & 3;
  long long subSquareSize = 111 << (pow * 2 - 2);</pre>
  long long ans = seg * subSquareSize;
  long long add = hilbertOrder(nx, ny, pow - 1, nrot);
  ans += (seg == 1 \mid \mid seg == 2) ? add : (subSquareSize
      - add - 1);
  return ans;
```

9.8 Pbds

```
#include <ext/pb_ds/priority_queue.hpp>
#include <ext/pb_ds/assoc_container.hpp>
using namespace __gnu_pbds;
#include <ext/rope>
using namespace __gnu_cxx;
int main () {
    _gnu_pbds::priority_queue <int> pq1, pq2;
  pq1.join(pq2); // pq1 += pq2, pq2 = {}
  cc_hash_table<int, int> m1;
  tree<int, null_type, less<int>, rb_tree_tag,
      tree_order_statistics_node_update> oset;
  oset.insert(2), oset.insert(4);
  cout << *oset.find_by_order(1) << ' ' << oset.</pre>
     order_of_key(1) << '\n'; // 4 0
  bitset <100> BS:
  BS.flip(3), BS.flip(5);
  rope <int> rp1, rp2;
  rp1.push_back(1), rp1.push_back(3);
  rp1.insert(0, 2); // pos, num
  rp1.erase(0, 2); // pos, len
  rp1.substr(0, 2); // pos, len
  rp2.push_back(4);
  rp1 += rp2, rp2 = rp1;
  cout << rp2[0] << '\' << rp2[1] << '\n'; // 3 4
```

9.9 Random

```
struct custom_hash {
    static uint64_t splitmix64(uint64_t x) {
        x += 0x9e3779b97f4a7c15;
        x = (x ^ (x >> 30)) * 0xbf58476d1ce4e5b9;
        x = (x ^ (x >> 27)) * 0x94d049bb133111eb;
        return x ^ (x >> 31);
    }
    size_t operator()(uint64_t a) const {
        static const uint64_t FIXED_RANDOM = chrono::
            steady_clock::now().time_since_epoch().count();
        return splitmix64(i + FIXED_RANDOM);
    }
};
unordered_map <int, int, custom_hash> m1;
random_device rd; mt19937 rng(rd());
```

9.10 Smawk Algorithm

```
11 query(int 1, int r) {
  // ...
```

```
struct SMAWK {
 // Condition:
  // If M[1][0] < M[1][1] then M[0][0] < M[0][1]
  // If M[1][0] == M[1][1] then M[0][0] <= M[0][1]
  // For all i, find r_i s.t. M[i][r_i] is maximum ||
       minimum.
  int ans[N], tmp[N];
  void interpolate(vector <int> 1, vector <int> r) {
    int n = 1.size(), m = r.size();
    vector <int> nl;
    for (int i = 1; i < n; i += 2) {
      nl.push_back(l[i]);
    run(nl, r);
    for (int i = 1, j = 0; i < n; i += 2) {
      \label{eq:while} \textbf{while} \ (\texttt{j} < \texttt{m} \ \&\& \ \texttt{r[j]} < \texttt{ans[l[i]])}
       assert(j < m && ans[l[i]] == r[j]);
      tmp[l[i]] = j;
    for (int i = 0; i < n; i += 2) {</pre>
      int curl = 0, curr = m - 1;
      if (i)
         curl = tmp[l[i - 1]];
      if (i + 1 < n)
         curr = tmp[l[i + 1]];
      11 res = query(l[i], r[curl]);
       ans[l[i]] = r[curl];
      for (int j = curl + 1; j <= curr; ++j) {
    ll nxt = query(l[i], r[j]);</pre>
         if (res < nxt)</pre>
           res = nxt, ans[l[i]] = r[j];
      }
    }
  }
  void reduce(vector <int> 1, vector <int> r) {
    int n = 1.size(), m = r.size();
    vector <int> nr;
    for (int j : r) {
      while (!nr.empty()) {
         int i = nr.size() - 1;
         if (query(l[i], nr.back()) <= query(l[i], j))</pre>
           nr.pop_back();
         else
           break:
       if (nr.size() < n)</pre>
         nr.push_back(j);
    run(1, nr);
  }
  void run(vector <int> 1, vector <int> r) {
    int n = 1.size(), m = r.size();
    if (max(n, m) <= 2) {
  for (int i : 1) {</pre>
         ans[i] = r[0];
         if (m > 1)
           if (query(i, r[0]) < query(i, r[1]))</pre>
             ans[i] = r[1];
         }
    } else if (n >= m) {
      interpolate(1, r);
    } else {
      reduce(1, r);
 }
};
```

9.11 Two Dimension Add Sum

```
struct TwoDimensionAddAndSum {
    // 0-index, [L, r)
    struct Seg {
        int l, r, m;
        ll vala, valb, lza, lzb;
        Seg* ch[2];
        Seg (int _l, int _r) : l(_l), r(_r), m(l + r >> 1),
            vala(0), valb(0), lza(0), lzb(0) {
        if (r - l > 1) {
            ch[0] = new Seg(l, m);
        }
}
```

```
ch[1] = new Seg(m, r);
      }
     void pull() {vala = ch[0]->vala + ch[1]->vala, valb
          = ch[0]->valb + ch[1]->valb;}
     void give(ll a, ll b) {
       lza += a, lzb += b;
       vala += a * (r - 1), valb += b * (r - 1);
     void push() {
       ch[0]->give(lza, lzb), ch[1]->give(lza, lzb), lza
            = 1zb = 0;
     void add(int a, int b, ll va, ll vb) {
       if (a <= 1 && r <= b)
         give(va, vb);
       else {
         push();
         if (a < m) ch[0]->add(a, b, va, vb);
         if (m < b) ch[1]->add(a, b, va, vb);
         pull();
       }
     long long query(int a, int b, int v) {
      if (a <= 1 && r <= b) return vala * v + valb;</pre>
       push();
       long long ans = 0;
       if (a < m) ans += ch[0]->query(a, b, v);
       if (m < b) ans += ch[1]->query(a, b, v);
       return ans;
    }
   };
   // note integer overflow.
   vector <array <int, 4>> E[N];
   vector <array <int, 4>> Q[N];
   vector <11> ans;
   void add_event(int x1, int y1, int x2, int y2, ll v)
     E[x1].pb({y1, y2, v, -v * x1});
E[x2].pb({y1, y2, -v, v * x2});
   void add_query(int x1, int y1, int x2, int y2, int id
     Q[x1].pb({y1, y2, -1, id});
     Q[x2].pb({y1, y2, 1, id});
     ans.pb(0);
   void solve(int n) {
     Seg root(0, n);
     for (int i = 0; i <= n; ++i) {</pre>
       for (auto j : E[i]) root.add(j[0], j[1], j[2], j
           [3]);
       for (auto j : Q[i]) ans[j[3]] += j[2] * root.
           query(j[0], j[1], i);
     }
  }
};
```

9.12 Matroid Intersection

Start from $S=\emptyset$. In each iteration, let

```
• Y_1 = \{x \notin S \mid S \cup \{x\} \in I_1\}
• Y_2 = \{x \notin S \mid S \cup \{x\} \in I_2\}
```

If there exists $x \in Y_1 \cap Y_2$, insert x into S. Otherwise for each $x \in S, y \not\in S$, create edges

```
 \begin{array}{l} \bullet \ x \to y \ \text{if} \ S - \{x\} \cup \{y\} \in I_1 \text{.} \\ \bullet \ y \to x \ \text{if} \ S - \{x\} \cup \{y\} \in I_2 \text{.} \\ \end{array}
```

Find a shortest path (with BFS) starting from a vertex in Y_1 and ending at a vertex in Y_2 which doesn't pass through any other vertices in Y_2 , and alternate the path. The size of S will be incremented by 1 in each iteration. For the weighted case, assign weight w(x) to vertex x if $x \in S$ and -w(x) if $x \not\in S$. Find the path with the minimum number of edges among all minimum length paths and alternate it.