Problem Set 1

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Problem 1

In the prisoner's dilemma

 $fink, quiet >_1 fink, fink >_1 quiet, quiet >_1 quiet, fink$

With a punishment to finking worse than any amount of time in prison

$$quiet, quiet >_1 quiet, fink >_1 fink, quiet >_1 fink, fink$$

Symmetrically for player 2.

The game is

$$\begin{array}{ccc} & \text{quiet} & \text{fink} \\ \text{quiet} & 3,3 & 2,1 \\ \text{fink} & 1,2 & 0,0 \end{array}$$

Alternatively, start with the prisoner's dilemma and punish finking until it becomes strictly dominated:

$$\begin{array}{ccc} & \text{quiet} & \text{fink} \\ \text{quiet} & 2,2 & 0,3\text{-x} \\ \text{fink} & 3\text{-x},0 & 1\text{-x},1\text{-x} \end{array}$$

for some x > 3

Nash: quiet, quiet

— common mistake is to call (3,3) the Nash equilibrium—this is the *outcome*. obviously that would be bad terminology as in general the same outcome might be induced by a strategy profile that is not a Nash equilibrium

Problem 2

Part a

 $sit, stand >_1 sit, sit >_1 stand, sit >_1 stand, stand$

$$\begin{array}{ccc} & \text{sit} & \text{stand} \\ \text{sit} & 2,2 & 3,1 \\ \text{stand} & 1,3 & 0,0 \end{array}$$

Nash: sit,sit

Part b

$$stand, sit>_1 sit, sit>_1 sit, stand>_1 stand, stand$$

	$_{ m sit}$	stand
sit	2,2	1,3
stand	3,1	0,0

Nash: stand, stand

Both players now less comfortable

— prisoner's dilemma shows how selfish behavior can be Pareto inefficient; here, (in (a)) it is efficient

Problem 3

 $aggressive, passive >_1 passive, passive$

 $>_1 passive, aggressive >_1 aggressive, aggressive$

	aggressive	passive
aggressive	0,0	3,1
passive	1,3	2,2

Nash: aggressive, passive, and passive, aggressive

Problem 4

Straightforward.

- Point is to gain familiarity with best response correspondence and their characterization of Nash equilibrium
- I saw two methods used for computing Nash: drawing arrows to indicate that a player has a profitable deviation, and drawing stars under a player's payoff when they are best responding
- "star" method corresponds to best response description of Nash: all stars, all best responding

Problem 5

Straightforward.

- There are no strictly dominated strategies here, but when presented with a large game try to eliminate strictly dominated strategies: obviously, they are never part of a Nash equilibrium
- This problem illustrates that weakly dominated strategies may be played in a Nash equilibrium, so don't flippantly eliminate those when looking for Nash

Problem 6

Observe that it is always a Nash equilibrium for both players to put in zero effort. On the other hand, if both players exert effort 1, each gets $\alpha/2 - 1 > 0$, so there always exists (for any $\alpha > 2$) an inefficient Nash equilibrium.

For remainder, try to solve for all Nash:

$$u_i(x_1, x_2) = \frac{\alpha x_1 x_2}{2} - x_i^2$$

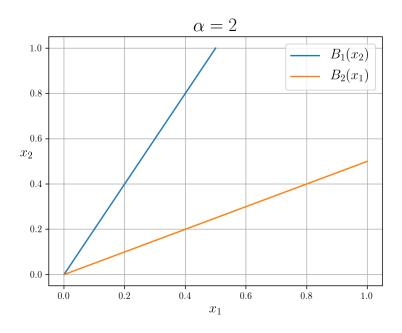
$$\frac{\partial u_1(x_1, x_2)}{\partial x_1} = 0 \implies x_1 = \frac{\alpha}{4} x_2$$

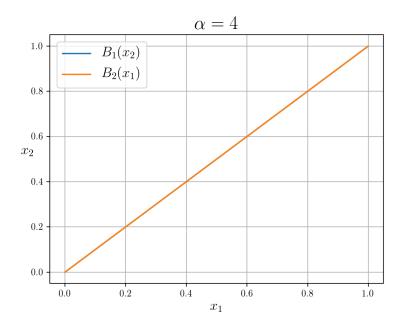
$$B_1(x_2) = \min\{\frac{\alpha}{4} x_2, 1\}$$

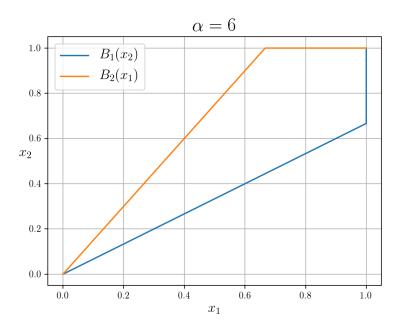
Three cases: (1) $\alpha < 4$ only (0,0) Nash. (2) $\alpha = 4$ any $(x,x), x \in [0,1]$ Nash, $\alpha > 4$ only (1,1) Nash

To wrap up: all Nash equilibria are Pareto dominated by both player's exerting maximal effort, except obviously when this is itself Nash, which it is once $\alpha \geq 4$. That is, all Nash equilibria are inefficient if and only if $2 < \alpha < 4$.

— Always good to graph best response correspondences. here are plots for $\alpha=2,\,4,\,$ and 6 (for $\alpha=4,\,$ the graphs overlap exactly)





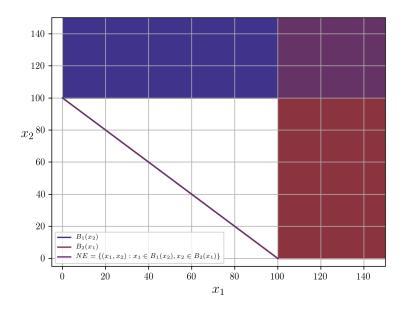


Problem 7

Part a There are no strictly dominated strategies: if $x_2 \ge 100$, player 1 is indifferent between all strategies

Part b Bids of 0, and greater than 100 are weakly dominated—they pay zero no matter what the other player's bid; any other bid x is the uniquely optimal response to the opponent bidding 100 - x, so can't be weakly dominated

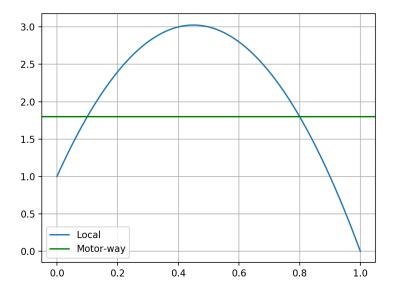
Part c As some of you pointed out, graphing the best responses, and thus solving for the Nash equilibria, is quick and painless:



Nash equilibria at (x, 100 - x) for $0 \le x \le 100$ and any two bids (x_1, x_2) both at least equal to 100

Problem 8

Part a See graph below



Part b Equilibrium traffic patterns at x = 0, 0.1, and 0.8. By considering small perturbation to equilibrium, easy to show x = 0, 0.8 stable, x = 0.1 unstable.

— Note though that if a single agent controlled a positive fraction of traffic, then x = 0.1 would not be an equilibrium to begin with.

Part c We maximize total benefit to the population $x(1+9x-10x^2)+(1-x)1.8$. the maximum cannot occur at x=0 or x=1, since there are obviously better options. (in fact the shape of the graphs immediatly imply that the solution will be where the blue curve dominates the green, and to the right of the blue's maximizer).

Solving for the critical points of the benefit function and choosing the one that yields most benefits gives $x^* = 0.5516$.

One possibility for implementing x^* is capping motorway usage at $1 - x^*$ and road usage at x^* .

This could be implemented as an equilibrium by subsidizing those on the motor-way, or taxing those on the local road so that the green and blue curves intersect at x^* .

It is also desirable to eliminate the two other equilibria that remain after such a subsidy/tax.

The caps achieve this.

With the subsidy/tax, inducing congestion on the motor-way by closing lanes, say, when x is less than some $\hat{x} < x^*$ could do this. Better policing on

local roads (paid for from the revenue of the tax) could work, too.