

Thanks [zzq](#) for his improvement on my original observation.

Given a [multiplicative function](#)  $f$ , we'd like to calculate its prefix sum  $F(n) = \sum_{i=1}^n f(i)$ .

For any multiplicative function  $g$ , there exists another multiplicative function  $h$  satisfying that  $f = g * h$  where  $*$  denotes the [Dirichlet convolution](#). Since they are multiplicative, it is enough to expand and check their values at prime powers:  $f(p^e) = \sum_{i=0}^e g(p^i)h(p^{e-i})$  where  $p$  is prime. Specifically,  $f(p) = g(1)h(p) + g(p)h(1) = g(p) + h(p)$ . Denote the prefix sum of  $g$  as  $G$ , i.e.,  $G(n) = \sum_{i=1}^n g(i)$ . We have  $F(n) = \sum_{i=1}^n f(i) = \sum_{i=1}^n \sum_{j|i} h(j)g(\frac{i}{j}) = \sum_{j=1}^n \sum_{k=1}^{\lfloor \frac{n}{j} \rfloor} h(j)g(k) = \sum_{j=1}^n h(j)G(\lfloor \frac{n}{j} \rfloor)$ . If there is a multiplicative function  $g$  such that the corresponding  $h$  satisfies that  $h(p) = 0$  at primes (or, equivalently,  $f(p) = g(p)$ ), to calculate  $F$  we only need to consider those  $j$  which is a [powerful number](#). (Explanation: if  $j$  is not powerful,  $h(j)$  is zero, thus they do not contribute to the sum and we can safely ignore them). Since there are  $O(\sqrt{n})$  powerful numbers, if  $G(n)$  can be calculated in time complexity  $O(n^\alpha)$ ,  $F(n)$  can be calculated in time complexity  $O(\max(\sqrt{n}, n^\alpha \cdot \frac{\zeta(2\alpha)\zeta(3\alpha)}{\zeta(6\alpha)}))$  where  $\zeta$  is the Riemann zeta function. In practice, it runs very fast.

A few examples:

1.  $f(p^e) = p$ : we can let  $g(x) = x$ . The final time complexity is  $O(\sqrt{n})$ .
2.  $f(p^e) = e + 1$ : we can let  $g(x) = \sigma_0(x)$ , the divisor counting function. The final time complexity is  $O(\sqrt{n})$ .
3.  $f(p^e) = p^e - 1$ : we can let  $g(x) = \varphi(x)$ , the Euler's totient function. The final time complexity is  $O(n^{\frac{2}{3}})$ .

Note that if we let the multiplicative function  $g(p^e) = f(p)^e$ , the above method converts the problem "summing a multiplicative function" into a problem "summing a completely multiplicative function". We can just focus on the latter to solve the class of these problems.