

Dynamic Programming

Top down
Bottom up

- * TOP down \Rightarrow Function + Memo
Recursive
- * BOTTOM up \Rightarrow For Loop + array

Dnc Vs DP
kompleksitas Rungkat
Memori Sacor

Memori Rungkat
kompleksitas Sacor

Kita punya lantai berukuran 1×10 , mau diisi dengan ubin berukuran 1×1 dan 1×2 , pemasangan tidak boleh tumpang tindih. Berapa banyak cara pemasangan ubin pada lantai jika pemasangan ubin dapat di rotasi atau refleksi sedemikian rupa

- variable dinamis Perm
DP
- ① initialisasi $dp(n) = \dots ?$
 $f(n) = \dots ?$
 - ② Model Recursif /
Transisi DP \rightarrow Abstraksi
 - ③ Base case
- ① optimasi Min - Max
 - ② Kombinasi Kas
 - ③ Traverse

1

Lantai berukuran $1 \times N$

2

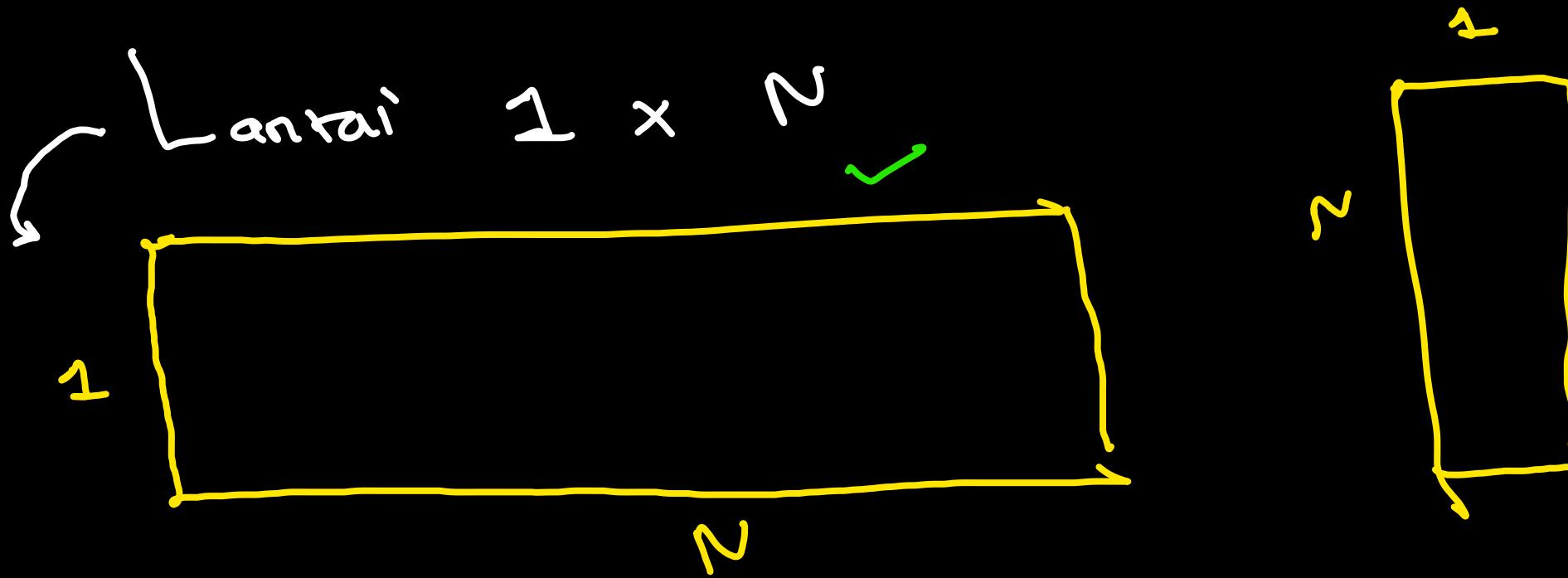
$dp(n)$ = menyatakan banyak cara memasang ubin pada lantai berukuran $1 \times N$.

* Abstraction

Visual — Tertulis

Rumah →

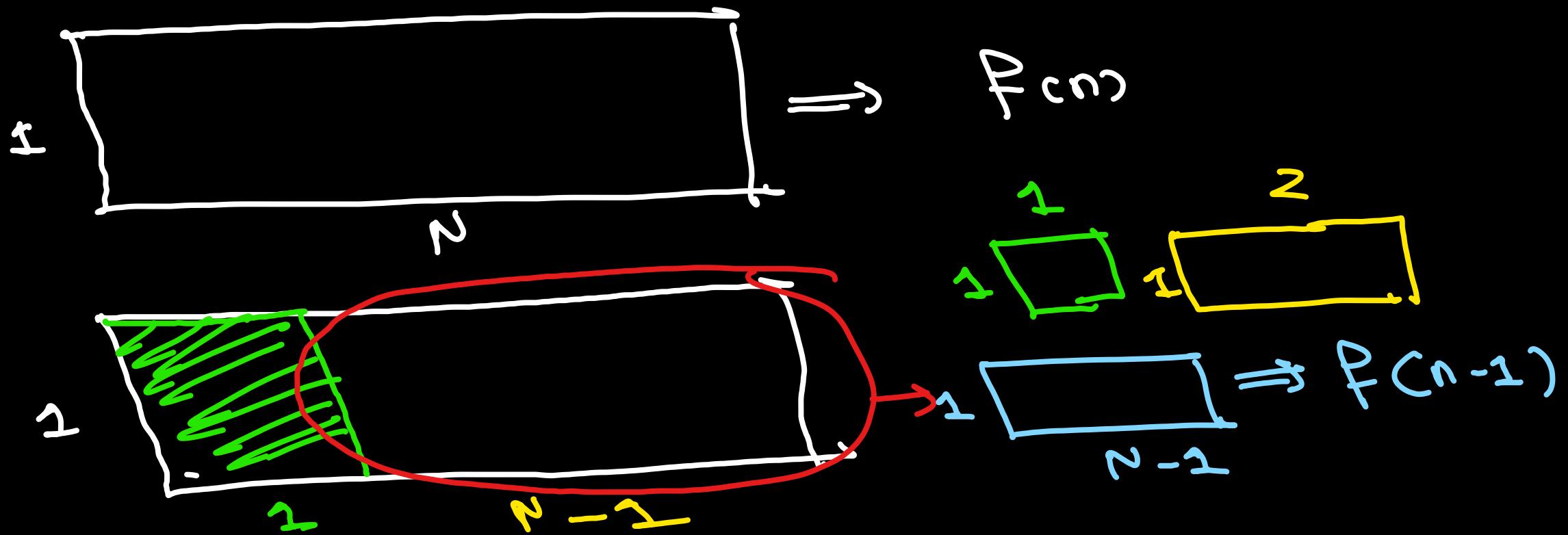


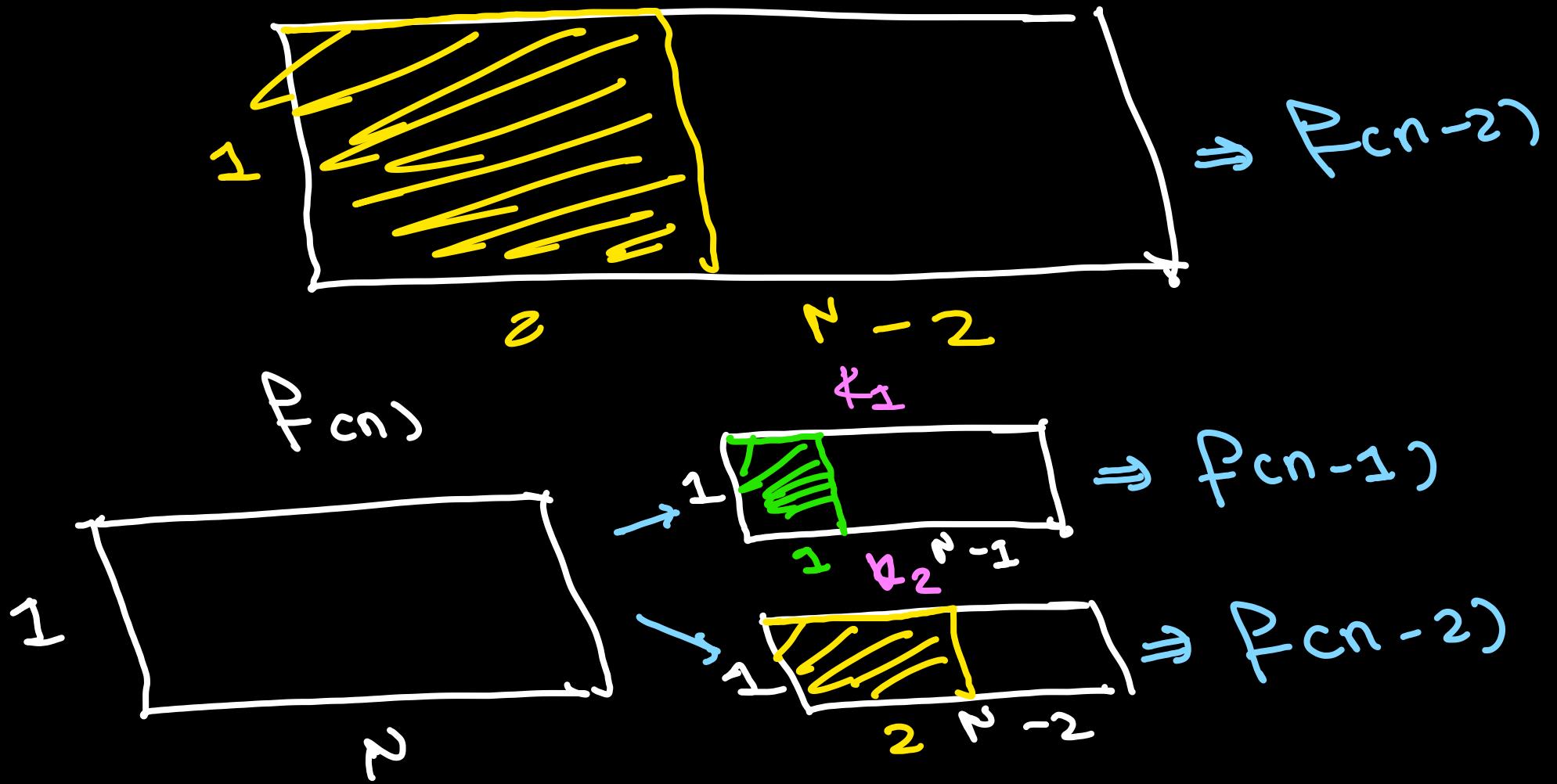


- ↖ Bottom - up Transisi: $f_{ci}) \rightarrow f_{c(i+1)} \rightarrow f_{c(i+2)} \rightarrow \dots \rightarrow$
 - ↖ Top down Transisi: $\text{answer} \rightarrow f_{c(i-1)} \rightarrow \dots \rightarrow \text{answer}$
- $f_{c(1)} \rightarrow f_{c(0)}$
 $f_{c(n)} / 2^p c_n)$

Literal base : Teori, Panduan → ikutin aja

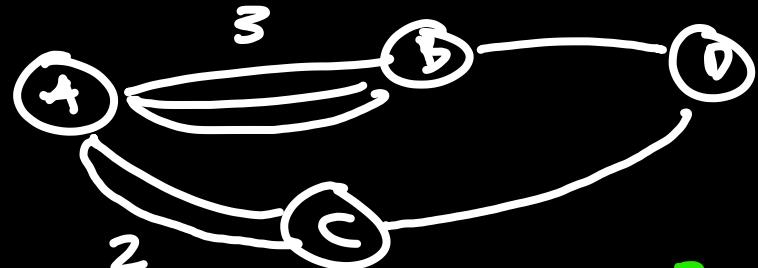
Experimental : Intuitif, berubah bentuk
↳ Kasus Real sgt banyak
Life





$$P(n) = P(n-1) + P(n-2)$$

Kaidah Pencacahan



2.

$K_1 : A - B - D$ $\left. \begin{matrix} 3 \text{ cara} \\ \end{matrix} \right\} K_1 \text{ & } K_2$ *Sesing Lepas*

$K_2 : A - C - D$ $\left. \begin{matrix} 2 \text{ cara} \\ \end{matrix} \right\}$

$3 + 2 = 5$ cara $+ \quad \text{dan} \quad A \neq D$

(3)

$$F(n) = F(n-1) + F(n-2)$$

$$F(1) = 1 \begin{array}{|c|} \hline \text{---} \\ \hline \end{array} = 1 \text{ cara}$$

$$F(2) = \begin{array}{|c|c|} \hline 1 & \\ \hline 1 & 1 \\ \hline \end{array} + \begin{array}{|c|} \hline \text{---} \\ \hline \end{array} = 2 \text{ cara}$$

$$\begin{aligned} F(10) &= F(7) \\ F(9) &= F(6) \\ F(8) &= F(5) \\ \vdots & \quad F(10) = F(9) + F(8) \end{aligned}$$

$$\begin{aligned} P(10) &= P(9) + P(8) \\ P(9) &= P(8) + P(7) \\ \vdots & \\ P(3) &= P(2) + P(1) = 3 \end{aligned}$$

$P(9)$

Banyak yang valid > invalid

Jawaban = Semua kemungkinan - invalid
Punah - invalid

1. $f(n)$ = Banyak cara menjumlahkan sehingga hasilnya adalah n

$$\dots + \dots + \dots + \dots + \dots + \dots = n$$

$\brace{f(n)}$

$$1 + 2 + 3 = 6 \quad \brace{3} \geq f(6)$$
$$3 + 5 = 6$$

~~$$1 + \dots + \dots + \dots + \dots + \dots = n$$~~
$$\dots + \dots + \dots + \dots + \dots = n - 1$$

$\brace{f(n-1)}$

$$2 + \dots + \dots + \dots = n \Rightarrow P(n-2)$$

$$3 + \dots + \dots + \dots = n \Rightarrow P(n-3)$$

$$i + \dots + \dots + \dots + \dots = n \Rightarrow P(n-i)$$

$$\cancel{n = g} \quad \cancel{i = 10} \rightarrow \cancel{10 + \dots + \dots + \dots = g}$$

$$\cancel{n = g} \quad \cancel{i = g} \rightarrow \cancel{g + \dots + \dots + \dots = g}$$

$$i < n \quad P(n) = \sum_{i=1}^{n-1} P(n-i)$$

$$F(n) = F(n-1) + F(n-2) \quad F(3) + F(2) + \dots$$

$$F(g) = F(g-1) + F(g-2) + F(g-3) + \dots$$

$+ F(g-8)$

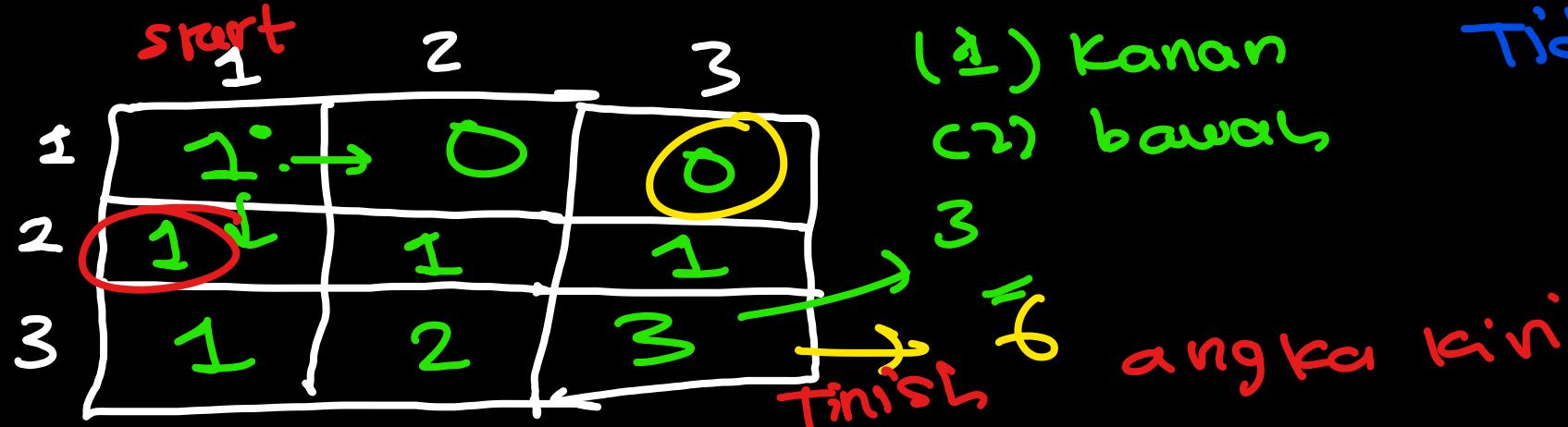
$\boxed{F(0)}$

- 1

9 → 7 → 5 → 3 → 1 → -1 → -3 → ...

$$2^0 = 1, \quad 0! = 1$$

$$F(0) = 1$$



Tidak boleh
lewat
(1,2)
angka atas

1

$$f(x,y) = f(x-1, y) + f(x, y-1)$$

$$f(3,3) = f(2,2) + f(2,1)$$

| Item ke - | <u>Wi</u> | <u>Vi</u> | <u>1/0</u> | <u>Max = n</u> |
|-----------|-----------|-----------|------------|------------------------|
| 1 | | | | $dp(n-w_i, i+1) + v_i$ |
| 2 | | | | $dp(n, i+1)$ |
| 3 | | | | |
| q | | | | <u>Max</u> |

$dp(n, i) = \begin{cases} \text{nilai Kepuasan max} \\ \geq v_i \text{ max} \end{cases}$

