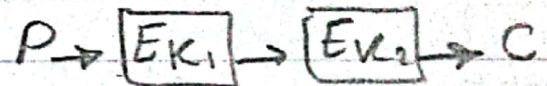


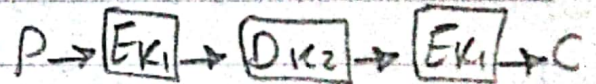
## Quiz 2.17

Q:1A

\* Double DES  $\rightarrow$  use 2 DES encryptions on each block  
 $\rightarrow C = E_{K_2}(E_{K_1}(P))$



\* Triple DES  $\rightarrow$  use 2 keys with E-D-E sequence  
 $\rightarrow C = E_{K_1}(D_{K_2}(E_{K_1}(P)))$



\* Triple DES is more secure, as double DES is vulnerable  
to the "meet-in-the-middle" attack, since  $X = E_{K_1}(P) = D_{K_2}(C)$ .

Moreover, it can be reduced to a single stage DES which is very vulnerable!

Q:1b 220 encry/sec, Key size = 40 bits

$\hookrightarrow$  Assume same rate for the decryption!  $\rightarrow 19.008 \times 10^6$  decry/day

$$\text{Time Rel} \Rightarrow 2^{40} / 19.008 \times 10^6 \approx 57'844.6 \text{ days} \\ \approx 158.47 \text{ years}$$

\* Brute force attack would be practical if we have much faster CPU OR the key size is really small, unless that happens, the brute force attack is not practical at all!

\* The attack will still be useless!



## 1-C 112-bit DES VS 56bit DES !

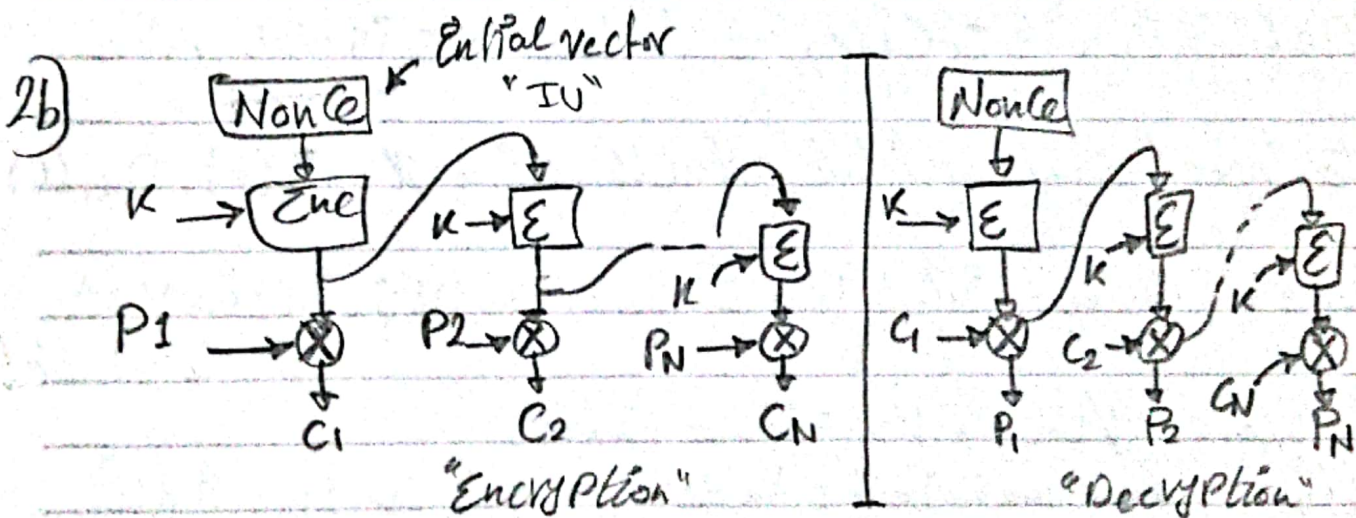
\* The Key-space sizes are  $2^{112}$  VS.  $2^{56}$ . Hence, the value of the time would be sth like:  $2^{112} / 2^{56} = 2^{56} \approx 7 \times 10^{16}$

## Q.2a

Stream

Block ciphers

- |                          |   |                                |
|--------------------------|---|--------------------------------|
| 1) Output Feedback (OFB) | } | 4) Electronic Codebook (ECB)   |
| 2) Cipher Feedback (CFB) |   | 5) Cipher block chaining (CBC) |
| 3) Counter (CTR)         |   |                                |



## 2C Adv → Bit errors don't propagate

- Disadv →
- Need an IV which is unique for each use
  - Sender & receiver must remain in sync
  - More vulnerable to message stream modification
  - only used with "Full Block Feedback"
  - e.g. CFB-64, CFB-128