

# Sheet 1

1)

P<sub>1</sub>

P<sub>2</sub>

P<sub>3</sub>

#Inst

1.5

1.0

2.2

CPI

$3 \times 10^9$

$2.5 \times 10^9$

$4 \times 10^9$

C<sub>r</sub>

a)  $IPS = \frac{C_r}{CPI}$

P<sub>2</sub> has the highest instructions Per Second ( $2.5 \times 10^9 \text{ s}^{-1}$ )

b)  $\#Inst = \frac{E_T C_r}{CPI} = E_T IPS$

P<sub>1</sub>

P<sub>2</sub>

P<sub>3</sub>

$2 \times 10^{10}$

$2.5 \times 10^{10}$

$1.82 \times 10^{10}$

$\#Cycles = \#Inst. CPI = E_T C_r$

$3 \times 10^{10}$

$2.5 \times 10^{10}$

$4 \times 10^{10}$

C) P1 P2 P3

1.5 1.0 2.2 CPI old

$3 \times 10^9$   $2.5 \times 10^9$   $4 \times 10^9$  Cr old

Reduce Execution time by 30%.

•  $\frac{E_{old} - E_{new}}{E_{old}} = 1 - \frac{E_{new}}{E_{old}} = 0.3$

•  $E_{new} = 0.7 E_{old}$

$(\# \text{Inst. CPI} \cdot C_T)_{new} = 0.7 (\text{CPI} \# \text{Inst.} C_T)_{old}$

→ #Inst. is the Same

→  $\frac{\text{CPI}_{new}}{\text{CPI}_{old}} = 1.2$

•  $C_{T_{new}} = (0.7/1.2) C_{T_{old}}$

•  $C_{r_{new}} = \frac{1.2}{0.7} C_{r_{old}} = 1.71 C_{r_{old}}$

P1

P2

P3

5.13GHz

4.275GHz

6.84GHz

$C_{r_{new}}$

$$MIPS = \frac{\#Instructions}{E_T (10^6)} = \frac{C_r}{CPI(10^6)}$$

Changes with Program.  
• Million instructions per second.

$$E_T = \#Cycles \times C_T \text{ with } \#Cycles = CPI \times \#Inst.$$

→ dep. on Cores

$$\text{and } C_T = \frac{1}{C}$$

dep. on Compiler, language, algorithm

dep. on tech, Cores \* All dep. on ISA

$$E_T|_{\text{after}} = E_T|_{\text{Affected}} + E_T|_{\text{Unaffected}}$$

Improvement

→  $\frac{E_T|_{\text{before}}}{n}$   
n ← amount of improvement.

Percentage difference of A w.r.t B

$$\frac{Q_A - Q_B}{Q_B} \times 100$$

Q { Performance (Posix)  
time (Slowx)

Speed up/Slow down of A w.r.t B

$$\frac{Q_A}{Q_B}$$

FOELW Stages

- load All 5
- Store No Wr
- branch No Mem, Wr
- ALU No Mem.

Processor

Pipelined

$$T_{\text{Cyc}} = T_{\text{longest stage}}$$

$$T_{\text{Tot.}} = K T_c + (n-1) T_c$$

n ← no. of instr.

Non Pipelined

Single Cycle

$$T_{\text{Inst.}} = T_{\text{Cyc}} = \sum T_{\text{stages}}$$

$$T_{\text{Tot.}} = n T_{\text{Inst.}}$$

Multi Cycle

$$T_c = T_{\text{longest stage}}$$

$$T_{\text{Inst.}} = K T_c$$

$$T_{\text{Tot.}} = \sum K_i T_c$$

K: no. of needed stages.

2) P1

$10^6$

$$0.1(1) + 0.2(2) + 0.5(3) + 0.2(3) = 2.6$$

$$2.5 \times 10^9$$

$$2.6 \times 10^6$$

$$\frac{2.6}{2.5} \times 10^6 = 1.04 \times 10^6$$

→ P1 is Slower.

P2

$10^6$  #Inst

$$0.1(2) + 0.2(2) + 0.5(2) + 0.2(2) = 2.0$$

$$3 \times 10^9$$

$$2 \times 10^6$$

$$\frac{2}{3} \times 10^6 = 0.67 \times 10^6$$

CPI

$C_r$

# Cycles

$$E_T (ns) = \frac{\#Cycles}{C_r}$$



3)

A

B

$$10^9$$

$$1.2 \times 10^9 \quad \#I$$

$$1.1$$

$$1.5$$

$$E_T \text{ (sec.)}$$

a)  $1 \times 10^{-9}$

$$1 \times 10^{-9}$$

$$C$$

$$\frac{1.1}{10^9 \times 10^{-9}} = 1.1$$

$$\frac{1.5}{1.2 \times 10^9} = 1.25$$

$$\begin{aligned} \text{CPI} \\ &= \frac{E_T}{\#I \cdot \text{CPI}} \end{aligned}$$

b)  $E_{Tx} = E_{Ty}$   
 // Processor x running A's Code  
 // ... y running B's Code  
 (Compiled)

$$\frac{C_{rx}}{C_{ry}} = ?$$

The CPI is intact as only the clock has changed. (Assume two processors have same ISA, cores; only clock has changed)

$$\#I \text{ CPI } C_{rx} = \#I \text{ CPI } C_{ry}$$

$$\frac{C_{rx}}{C_{ry}} = \frac{C_{ry}}{C_{rx}} = \frac{\text{CPI}_x}{\text{CPI}_y} \cdot \frac{I_x}{I_y} = \frac{1.1}{1.25} \times \frac{10^9}{1.2 \times 10^9} = 0.733$$

It's slower!



C) New Compiler

$$\# I = 6 \times 10^8$$

$$CPI = 1.1$$

→ Cr Const. across all.

. It's explicitly the same Processor as in Part A  
So we can use the CPI values

$$SpeedUP|_C = \frac{Perf|_C}{Perf|_A} = \frac{10^9 \times 1.1}{6 \times 10^8 \times 1.1} = 1.67$$

(Perf =  $\frac{1}{\# \text{ Cycles at Given Cr}}$ )

$$SpeedUP|_C = \frac{P_C}{P_B} = \frac{1.2 \times 10^9 \times 1.25}{6 \times 10^8 \times 1.1} = 2.27$$

Performance ✓

$$4) T_{\text{before}} = 250s = 180 + 70$$

$$= (85 + 40 + 55) + 70$$

load/store    branch    Int    Float

$$a) T_{\text{After}} = T_{\text{Affected}} + T_{\text{Unaffected}} = 0.8 \times 70 + 85 + 40 + 55$$
$$= 236s$$
$$\therefore \frac{250 - 236}{250} \times 100 = \text{by } 5.6\%$$

b)  $T = 195 + t_{int}$

$$T_{After} = 250 - 250 \times 0.2 = 200$$

$$t_{int_{new}} = 5$$

$$t_{int_{old}} = 55$$

Should be improved by  $\frac{55-5}{55} \approx 91\%$

c)  $T = 210 + t_{branch}$

$T_{After} = 200$ , impossible.  $t_{branch}$  is at least zero.

5)

	P <sub>float</sub>	P <sub>int</sub>	LIS	branch
#Inst. (10 <sup>8</sup> )	50	110	80	16
CPI	1	1	4	2

$C_r = 2 \times 10^9$  // useless. Perf  $\propto \frac{1}{\#InstCPI}$  at const.  $C_r$

$$\#Cycles = 50 \times 1 + 110 \times 1 + 80 \times 4 + 16 \times 2 = 512 \times 10^6$$

a) Should require half the no. of Cycles

$$256 = 462 + 50 CPI_{new}$$

↙ unaffected  
time

↘ +ve float.

// IMPOSSIBLE

Time<sub>After</sub> > Time<sub>unaffected</sub> else it's impossible  
↳ or #Cycles

b)  $256 = (512 - 80 \times 4) + 80 CPI_{new}$

↘ unaffected

= total

- to be affected

\*

$$CPI_{new} = 0.8$$

// Improve by

reducing it by  $\frac{4 - 0.8}{4}$

= 80%.

c)  $T_{After} = 512 - ((110 + 50) \times 0.4 + (80 \times 4 + 16 \times 2) \times 0.3)$   
= 342.4

reduce by ↙

(33.125% Improvement)

// reduction