

CNF Algorithm

Unit
Production
 $\square \rightarrow \square$

→ We have 4 main steps

1. Eliminate empty production:

ای case تعلق فیہ ال ع بتیلا و بروج یکد
Possible cases

ie: $A \rightarrow B$ & $\{A \& B \text{ may be nullable}\}$

∴ $A \rightarrow B \mid \epsilon \mid A \mid \epsilon \rightarrow$ the 4 possible cases.

2. Eliminate Variable Unit Production

لو فی variable موفی غیر قیلا واحد قیلا
عوج بی.

3. Replace Long Productions by short ones:

ای case فیہ اکثر مد حرفیہ، قیلا.

ie: $A \rightarrow BCDEF$

↓
 $A \rightarrow KL$

$K \rightarrow BC$ و $L \rightarrow DE$

وخذ بالک ان عندنا only allowed cases

$X \rightarrow XX$ or $X \rightarrow x$

Capital یا حری Capital یا حری
small یا حری Capital یا حری

4. Move Terminal to Unit Production

$A \rightarrow bC \Rightarrow A \rightarrow BC$

$B \rightarrow b$

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$$X \rightarrow 0X \mid 1X \mid \epsilon$$

Steps For Removal of Unit production:-

→ We have 3 main steps:-

1. To remove $A \rightarrow B$, add production $A \rightarrow X$ to the grammar rule whenever $B \rightarrow X$ occurs in the grammar where $[X \rightarrow \text{Terminal}]$

2. Delete $A \rightarrow B$ from the grammar

3. Repeat from step 1 until all unit productions are removed.

Example:-

$S \rightarrow XY, X \rightarrow a, Y \rightarrow Z \mid b, Z \rightarrow M, M \rightarrow N, N \rightarrow a$

Solution:

add this to Unit Productions:

1) Since $N \rightarrow a \therefore M \rightarrow a$ our grammar
& remove $M \rightarrow N$

$Y \rightarrow Z, Z \rightarrow M$
 $M \rightarrow N$

$S \rightarrow XY, X \rightarrow a, Y \rightarrow Z \mid b, Z \rightarrow M$
 $M \rightarrow N, M \rightarrow a, N \rightarrow a$

1 variable \rightarrow Z is used
another 1 variable \rightarrow M is used

2) Repeat:

$\because Z \rightarrow M \& M \rightarrow a \therefore Z \rightarrow a$

$\because Y \rightarrow Z \& Z \rightarrow a \therefore Y \rightarrow a$

$S \rightarrow XY, X \rightarrow a, Y \rightarrow a \mid b, Z \rightarrow a, M \rightarrow a$

Final Grammar But, Now Z, M, N can not be reached from S , so we need to apply modification.

→ Remove unreachable variables ✓

Final →

$S \rightarrow XY, X \rightarrow a, Y \rightarrow a \mid b$



Example:-

$$\begin{aligned} S &\rightarrow aXbX \\ X &\rightarrow aY|bY|\epsilon \\ Y &\rightarrow X|c \end{aligned}$$

Important Note

لو عندك $B \rightarrow A$ و $A \rightarrow X| \epsilon$ فبت
واضح ان A ممكن تبقي Null
لكن متأكد ان B ح مكان ممكن
تبقي Null فلازم تعد نفس الـ Null
على

Step 1: Replace all Nullables

$$\begin{aligned} S &\rightarrow aXbX | abX | ~~abX~~ | aXb | ab \\ X &\rightarrow aY | bY | a | b \rightarrow \text{In case } bY = X \text{ \& } X \text{ is Null} \\ Y &\rightarrow X | c \end{aligned}$$

Step 2: Eliminate variable unit production

الـ Y ح و a و b ح Variable في الـ X فبت ح

$$\begin{aligned} S &\rightarrow aXbX | abX | aXb | ab \\ X &\rightarrow aY | bY | a | b \\ Y &\rightarrow aY | bY | a | b | c \end{aligned}$$

Step 3: Replace all long terms with short terms.

$$\begin{aligned} EF &\rightarrow AXBX \rightarrow aXbX \\ E &\rightarrow AX \\ F &\rightarrow BX \\ A &\rightarrow a \\ B &\rightarrow b \end{aligned} \quad \left. \begin{aligned} \therefore S &\rightarrow EF | AF | EB | AB \\ X &\rightarrow AY | BY | a | b \\ Y &\rightarrow AY | BY | a | b | c \\ E &\rightarrow AX & F &\rightarrow BX \\ A &\rightarrow a & B &\rightarrow b \end{aligned} \right\}$$

Done ✓✓