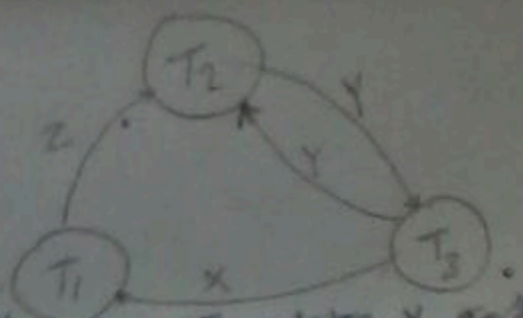


(a) Draw the serializability (precedence) graphs for  $S_1$  and  $S_2$  and state whether each schedule is serializable or not. If a schedule is serializable, write down the equivalent serial schedule(s).

$S_1$

not serializable.



Non Recoverable because  $T_2$  read  $y$  after  $T_3$  write  $y$  and  $T_2$  is Committed First

$S_2$

Serializable

and serial schedule



$T_3 T_1 T_2$

and any cascadeless - if strict schedule is also recoverable also cascadeless

(b) Determine whether  $S_2$  is strict, cascadeless, recoverable or nonrecoverable and justify your answer.

recovery :-  $T_1$  و  $T_2$  (  $T_2$  قبل  $T_1$  read  $x$  )

کتابخانه  $T_1$  و  $T_2$  Commit قبل  $T_2$

non recoverable

ام  $T_2$  قبل  $T_1$  Commit

STRICT  $T_1$  قرأت  $x$  میفهمند حد قرا لها و بکتابخانه بها می  
ما  $T_1$  قبل Commit

cascadeless

انها مثلا "قراوت"  $x$  في  $T_1$  و بعديه كتيبت  $T_1$  بيقى كده قبل  
قراوت  $x$  في  $T_2$

... schedule is serializable or not. If a schedule is serializable, write down the equivalent serial schedule(s).

$S_1$  :

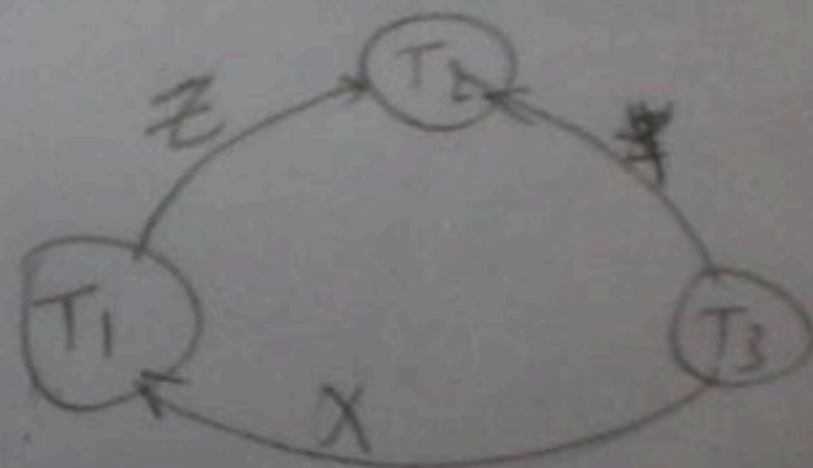
not serializable.



Non recoverable because  $T_2$  read  $y$  after  $T_3$  write  $y$  and  $T_2$  is committed first.

$S_2$  :

Serializable



and serial schedule

$\boxed{T_3 T_1 T_2}$

and any cascadeless schedule is also recoverable

Notes

if strict schedule is also cascadeless

(b) Determine whether  $S_2$  is strict, cascadeless, recoverable or nonrecoverable.



and any cascadeless  
Schedule is also recoverable

Notes

If Strict schedule  
also cascadeless

(b) Determine whether  $S_2$  is strict, cascadeless, recoverable or nonrecoverable and justify your answer.

re recoverable :-  $T_1$  و  $T_2$  (  $T_2$  read  $X$  )  $T_1$  کندی  
کتابها خراب و  $T_1$  Commit قبل  $T_2$

non recoverable

م  $T_2$  Commit Commit اول

Strict  $T_1$  قرأت  $X$  متوقف شد تا  $T_2$  تمام کارها را  
ما  $T_1$  رقیل Commit

cascadeless

انها مثلا "قرآوت"  $X$  فی  $T_1$  و بعد  $T_1$  یعنی کرده قبل  
قرأت  $X$  فی  $T_2$  و بعد به کتبیت  $T_1$

بسیار  
ی داده

2. Discuss the system of propagation of privileges and the revocation thereof slide 17

Propagation: Giving Grant of Privilege to other accounts without knowledge of owner account  
ex GRANT OPTION

Revocation: is desirable to grant a Privilege to a user temporarily

ex Owner grant Privilege to user for specific task and wants to revoke it after task end  
he use Revoke Command



3. Consider the relation shown below using mandatory access control.

- Fill in the proper values for TC (tuple classification)
- How would it appear to a user with classification U?
- Suppose a classification U user tries to update the salary of "Ahmed" to 50,000; what would be the result of this action? And why is this necessary?

National ID	Name	Salary	Job Performance	TC
123456789	Ahmed	40,000 C	Fair S	S
234567891	Hassan	60,000 C	Good S	S

(b) When User U read table: he will see

ID	Name	Salary	Job	TC
12---	Ahmed	NULL	NULL	<del>S</del>
234--	Hassan	NULL	NULL	<del>S</del>

(c)

ID	Name	Salary	Job	TC
12---	Ahmed	40,000 C	Fair S	S
234--	Hassan	50,000	NULL	U
	Hassan	NULL	NULL	U

invisible & true value

User U will see that tables only

this is called Polyinstantiation

This an option organization gives to users to update secret values as they need because there is security on salary