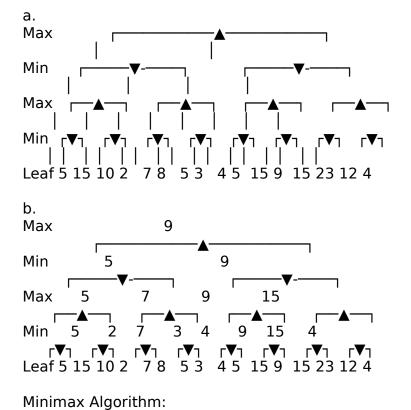
Adversarial Search

Chapter 5

A tree with branching factor 2 and depth 4, has the following values for its leaves: 5, 15, 10, 2, 7, 8, 5, 3, 4, 5, 15, 9, 15, 23, 12, 4

- a. Draw the described tree.
- b. Find the MinMax value of the tree.
- c. Apply Alpha-Beta pruning showing the nodes that will not be evaluated.



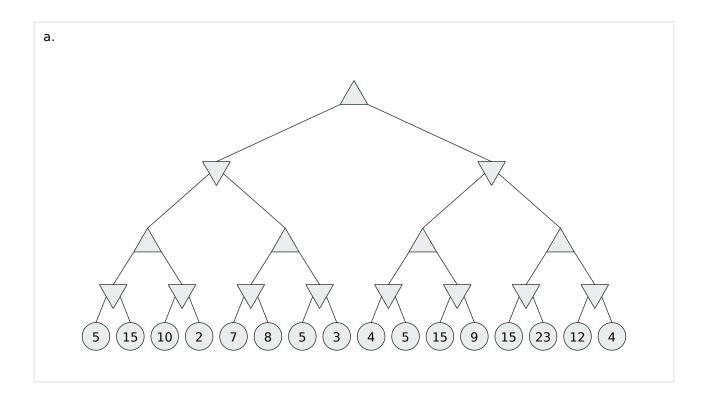
If is terminal(state): return value of(state)

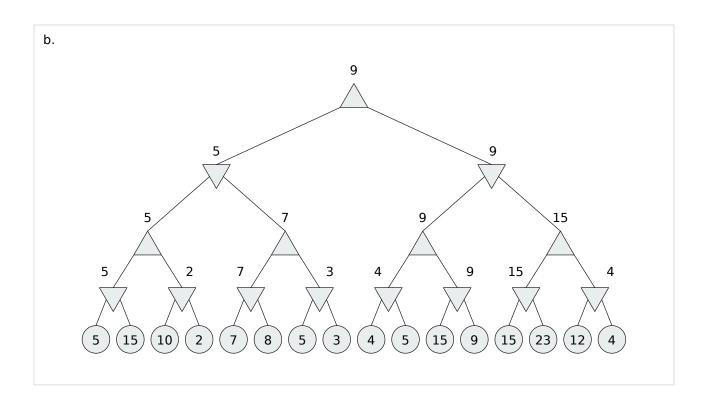
For action in get actions(state, max player):

child = get successor(state, max player, action)

Function max_search(state):

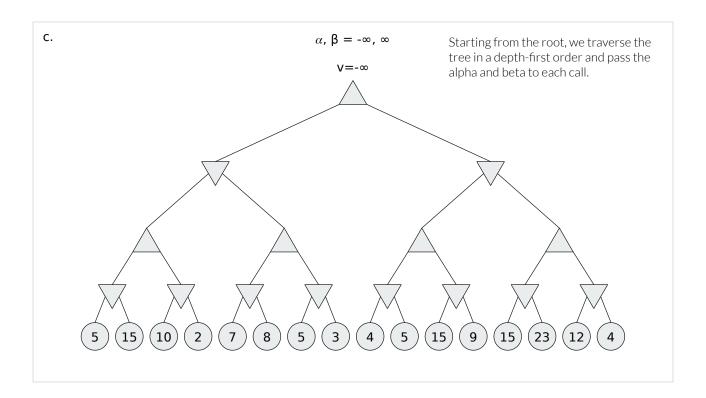
value = $-\infty$

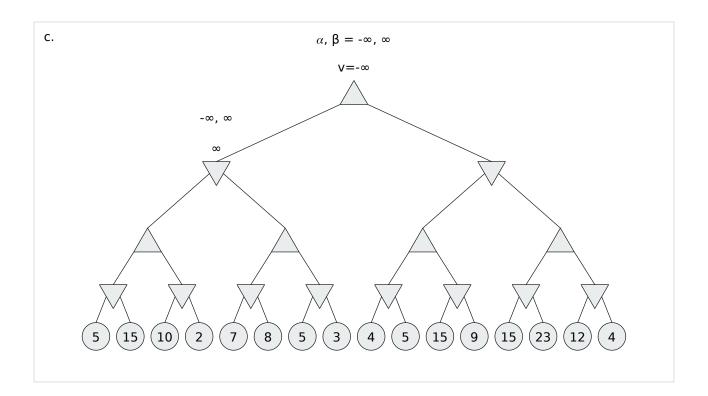


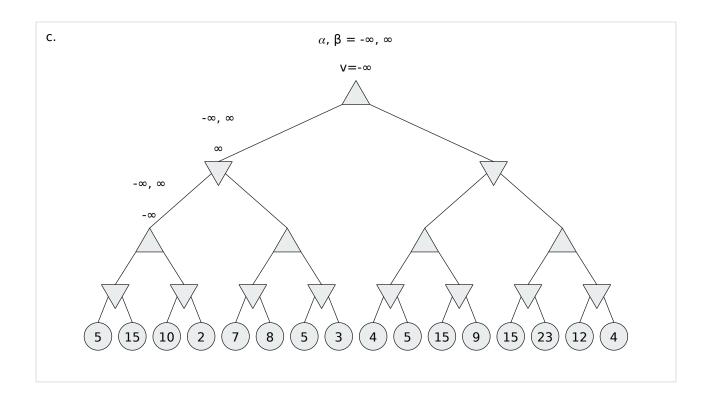


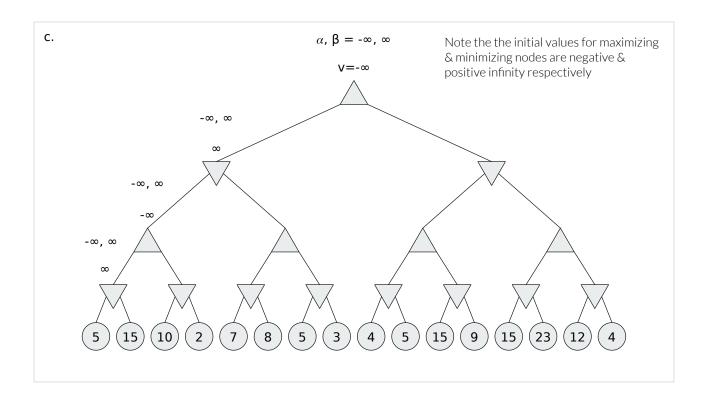
```
function Alpha-Beta-Search(state) returns an action
   v \leftarrow \mathsf{Max}\text{-}\mathsf{Value}(state, -\infty, +\infty)
   return the action in ACTIONS(state) with value v
function Max-Value(state, \alpha, \beta) returns a utility value
   if TERMINAL-TEST(state) then return UTILITY(state)
   for each a in ACTIONS(state) do
      v \leftarrow \mathsf{MAX}(v, \mathsf{MIN}\text{-}\mathsf{VALUE}(\mathsf{RESULT}(s, a), \alpha, \beta))
      if v \geq \beta then return v
      \alpha \leftarrow \overline{\mathsf{MAX}}(\alpha, v)
   return v
function Min-Value(state, \alpha, \beta) returns a utility value
   if Terminal-Test(state) then return Utility(state)
   v \leftarrow +\infty
   for each a in ACTIONS(state) do
      v \leftarrow \text{MIN}(v, \text{MAX-VALUE}(\text{RESULT}(s, a), \alpha, \beta))
      if v \leq \alpha then return v
      \beta \leftarrow \text{MIN}(\beta, v)
   return v
```

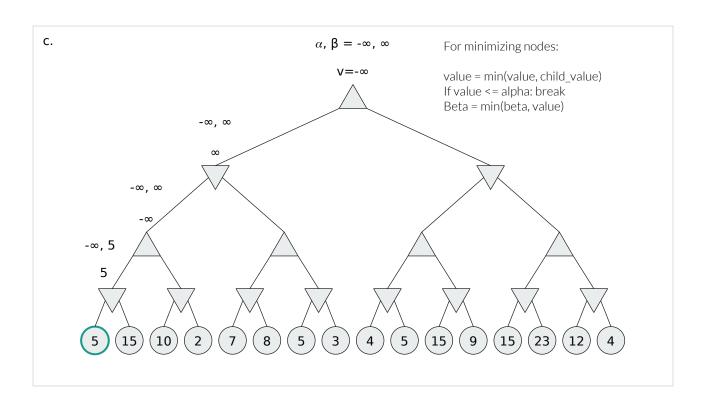
Figure 5.7 The alpha-beta search algorithm. Notice that these routines are the same as the MINIMAX functions in Figure 5.3, except for the two lines in each of MIN-VALUE and MAX-VALUE that maintain α and β (and the bookkeeping to pass these parameters along).

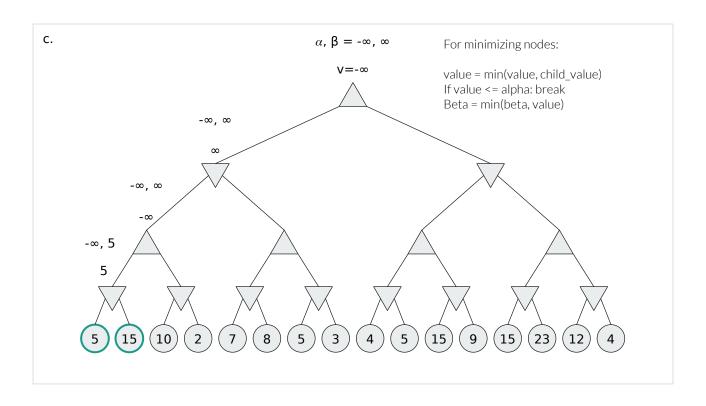


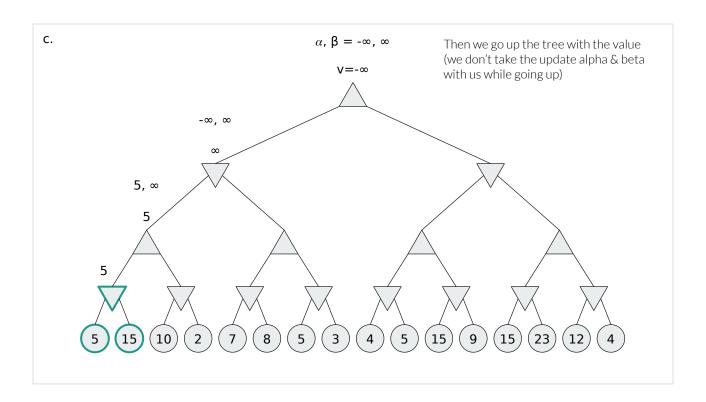


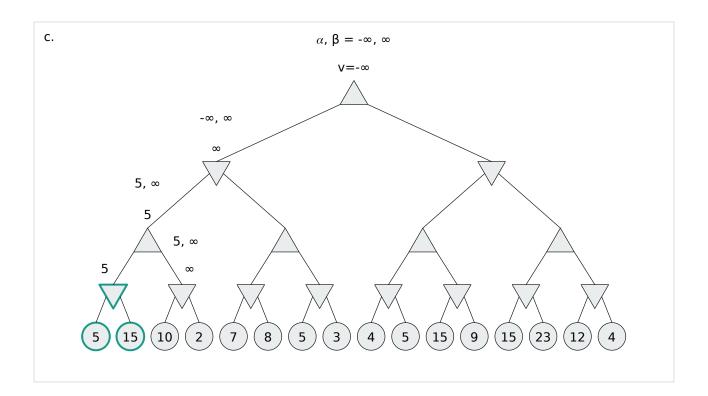


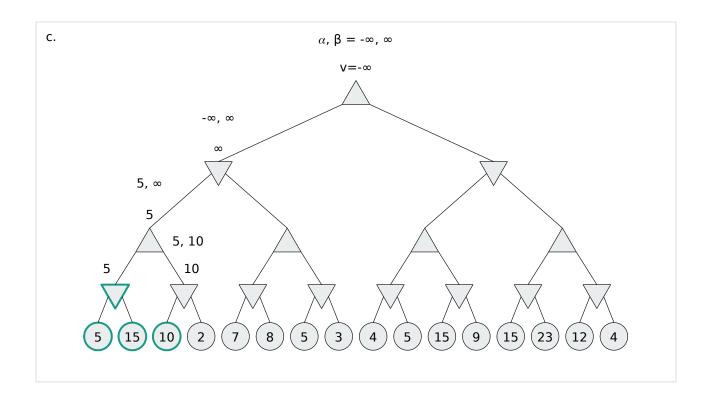


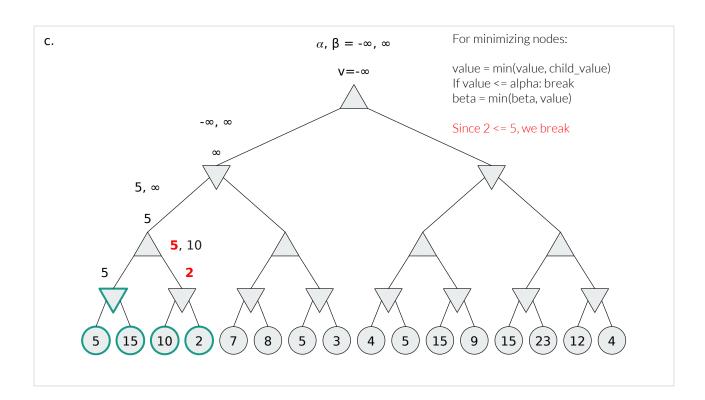


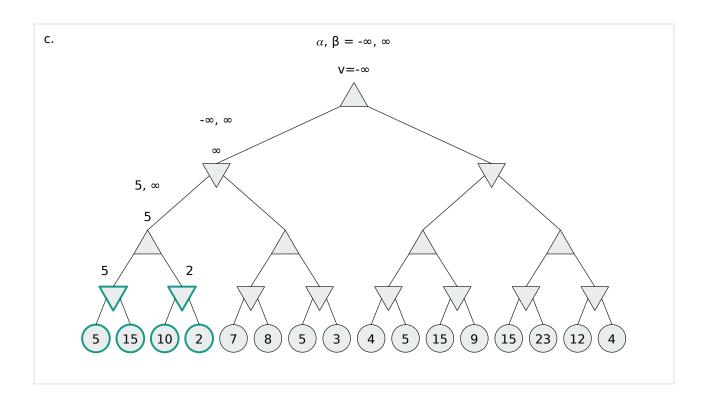


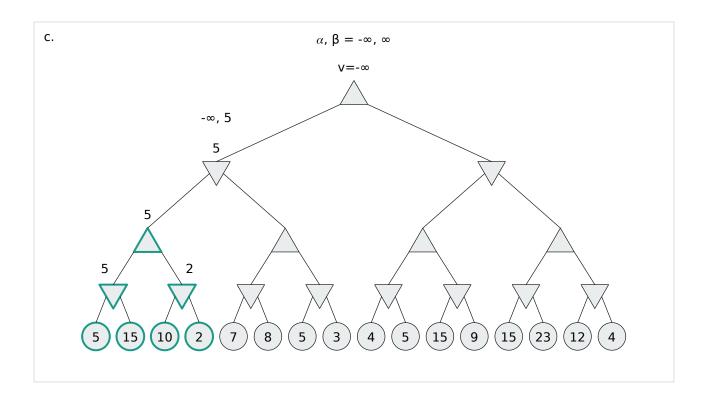


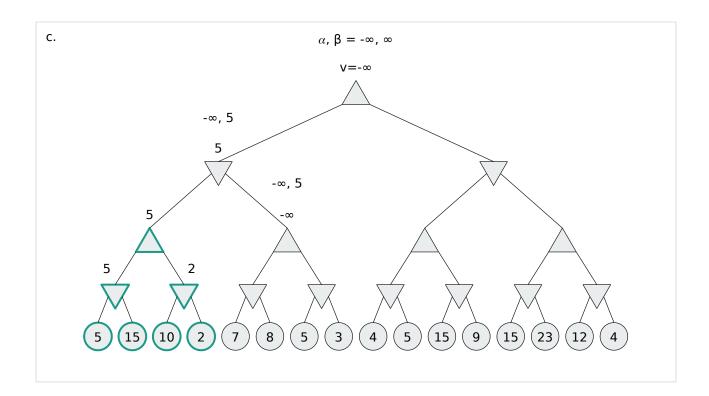


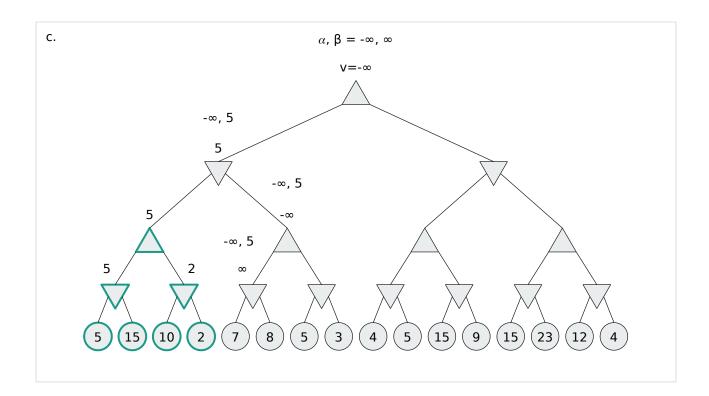


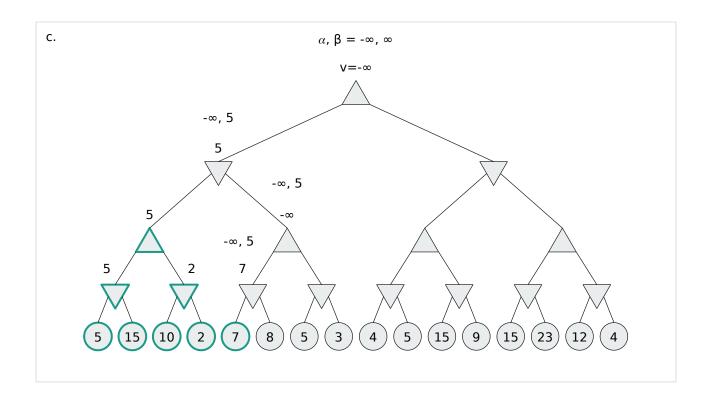


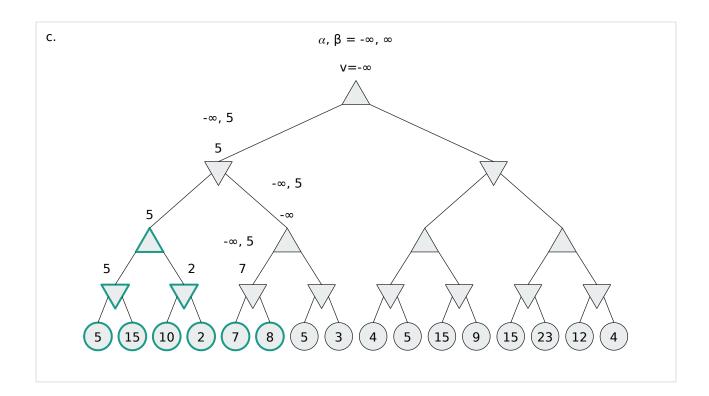


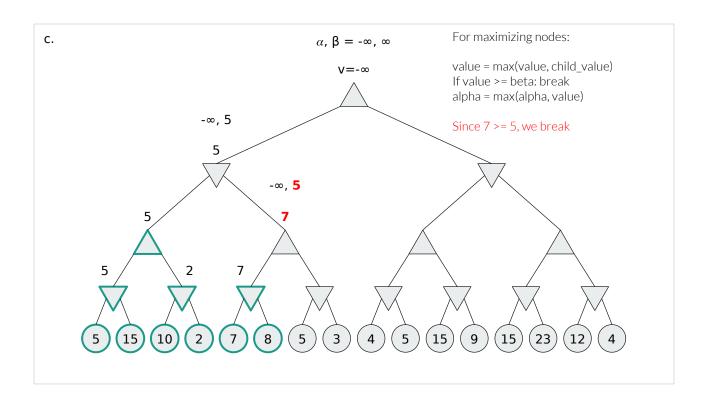


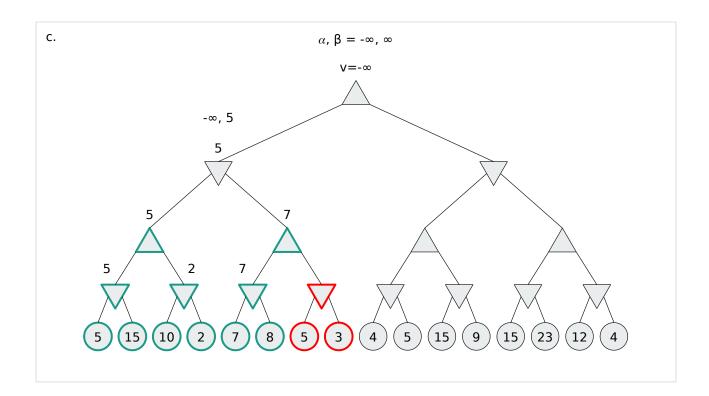


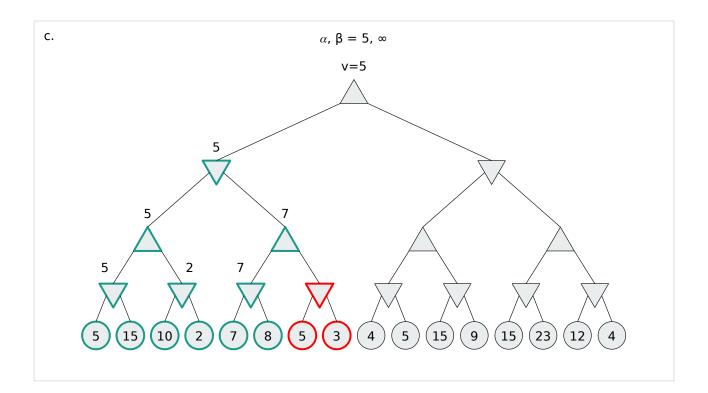


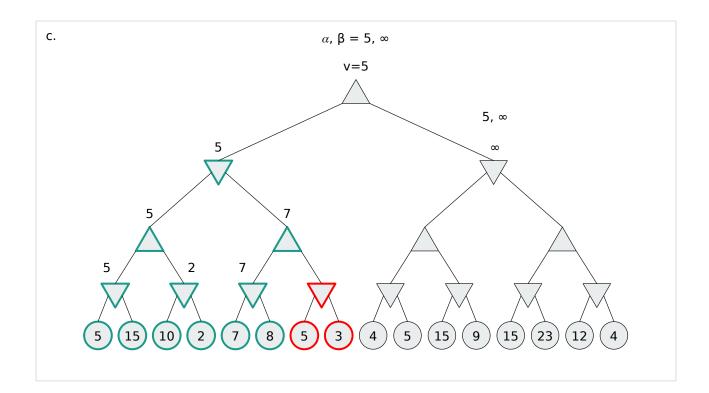


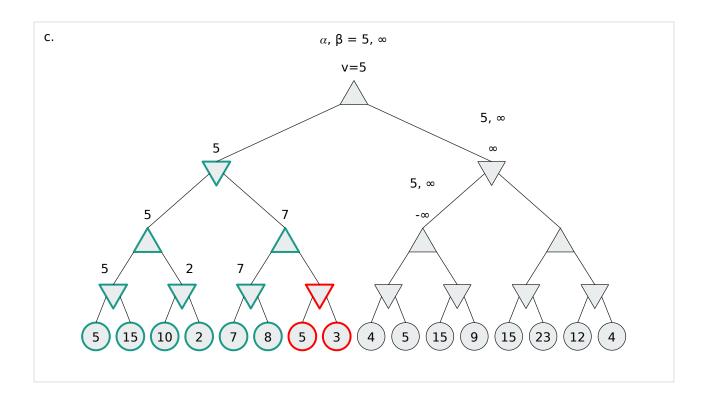


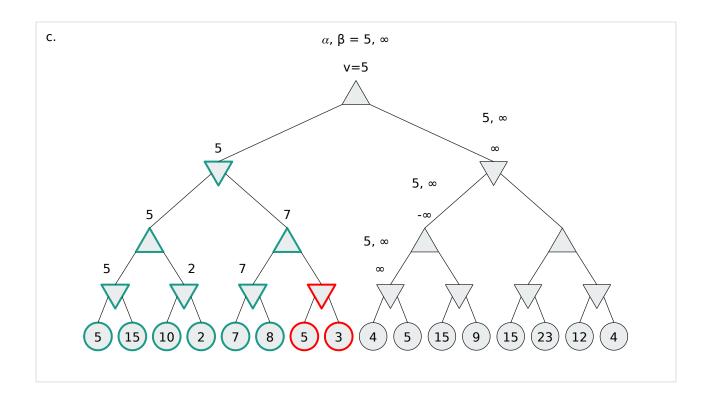


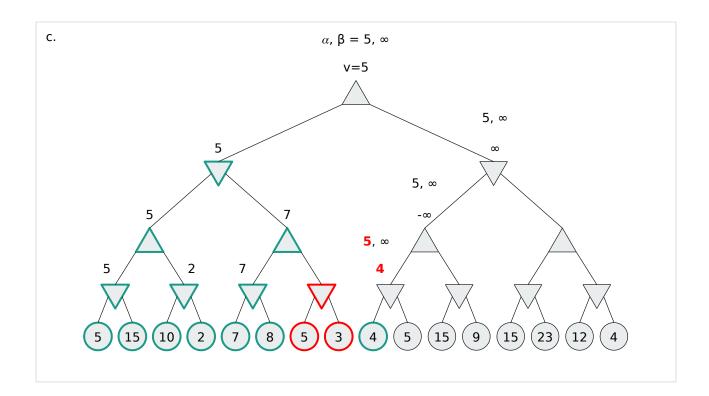


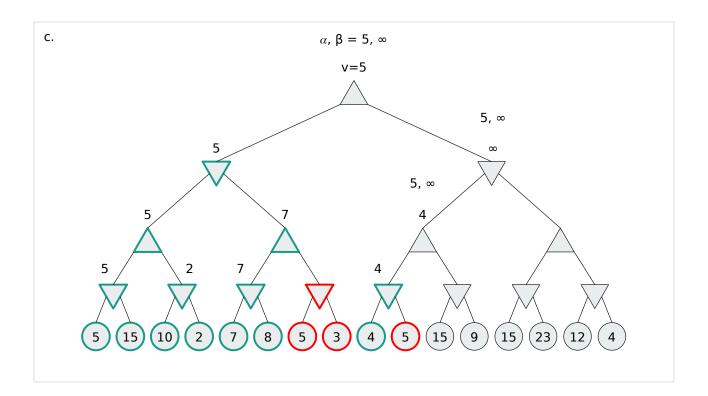


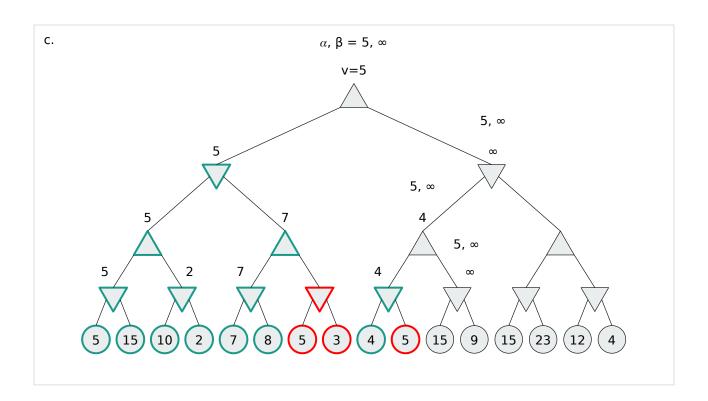


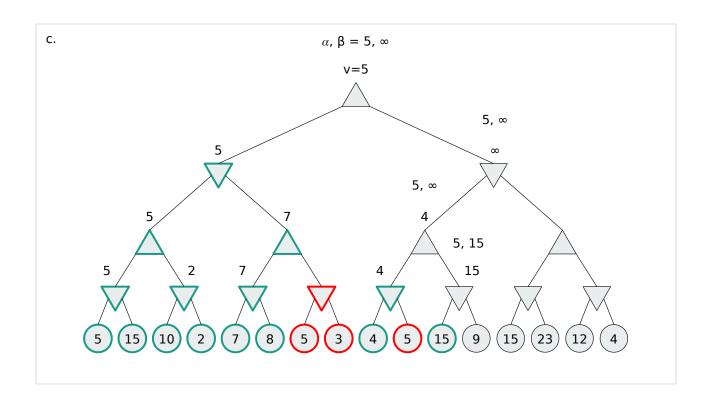


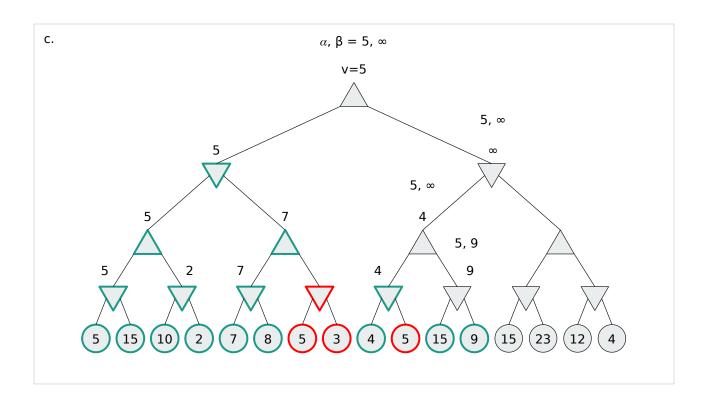


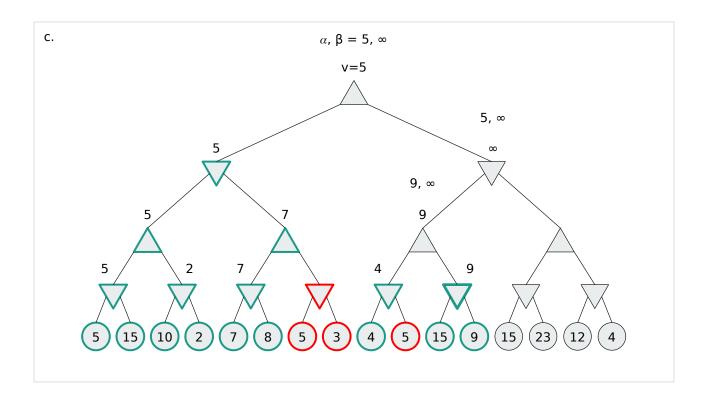


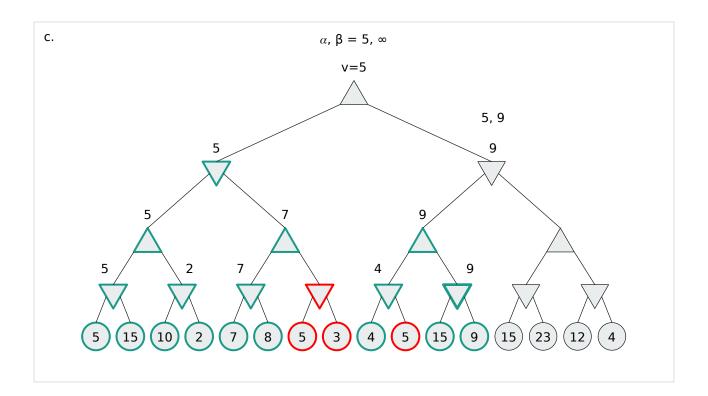


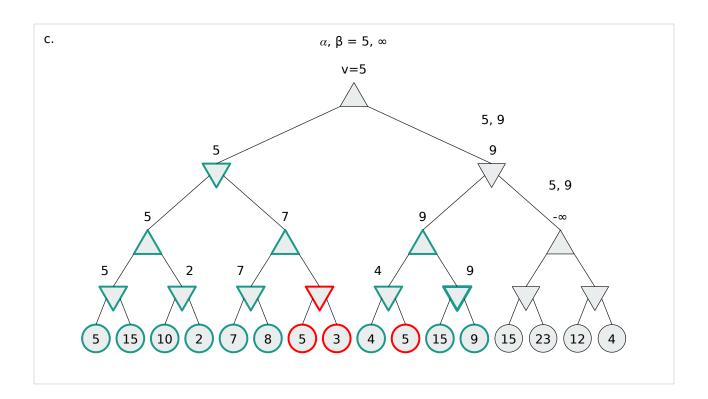


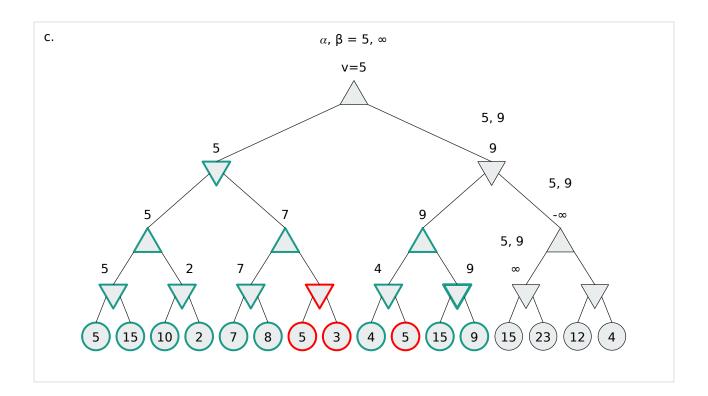


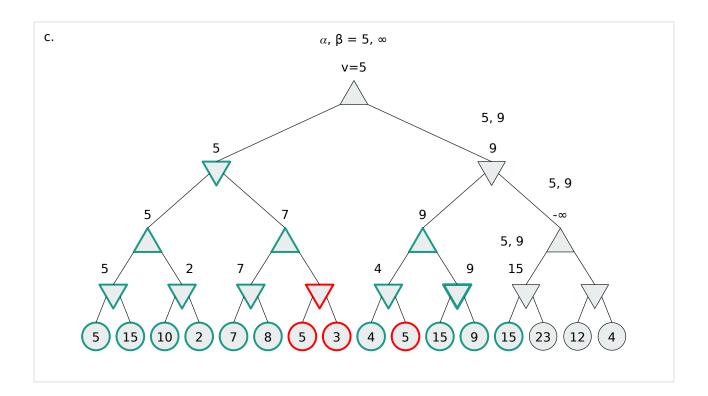


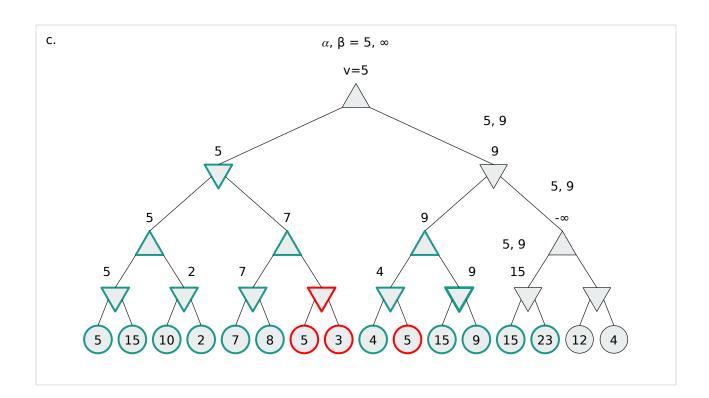


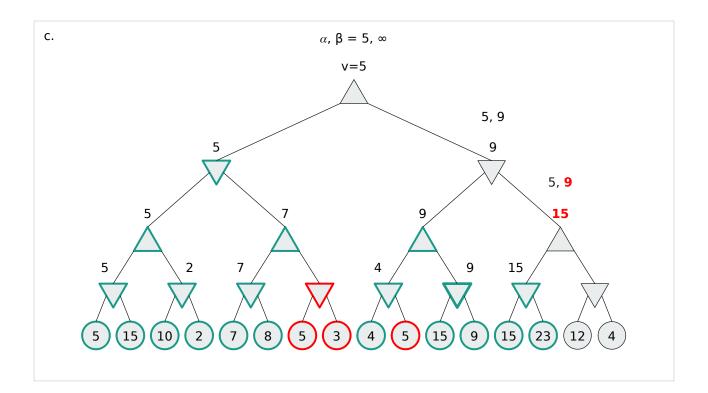


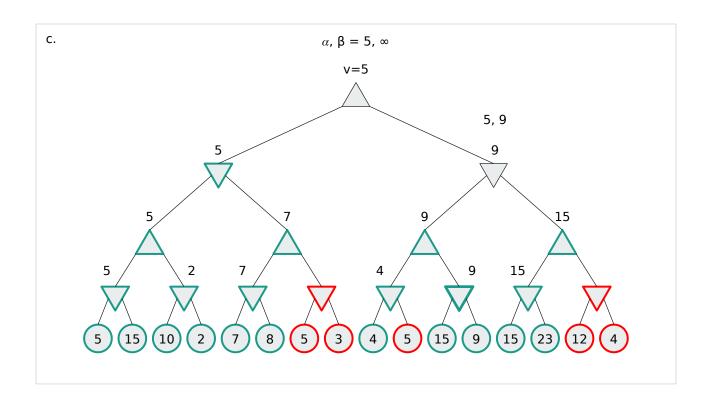


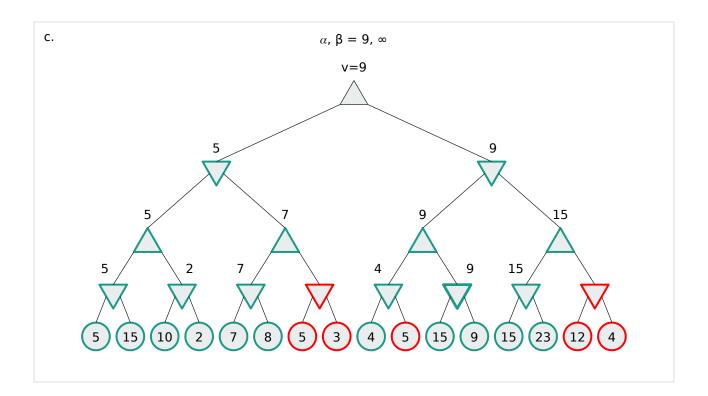


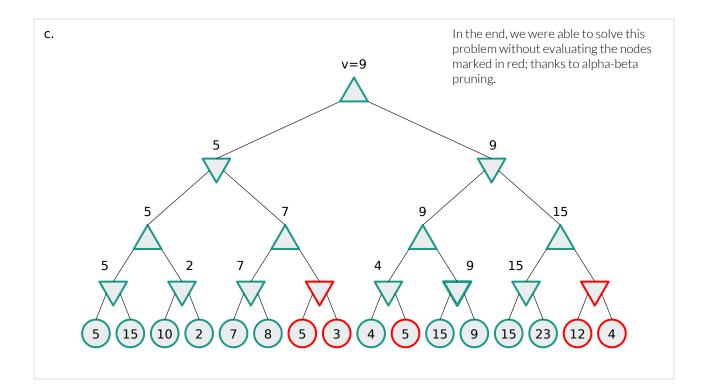


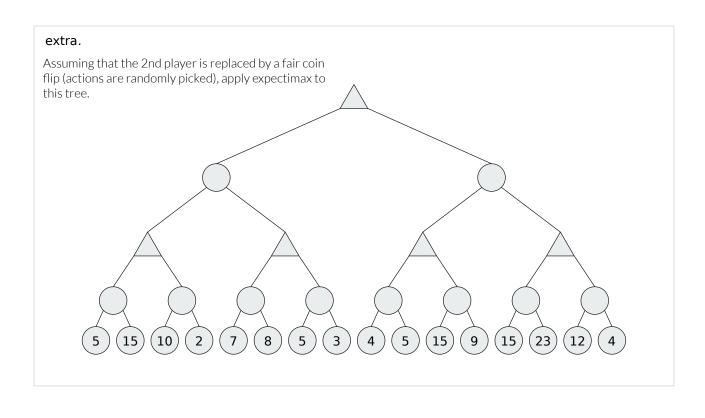


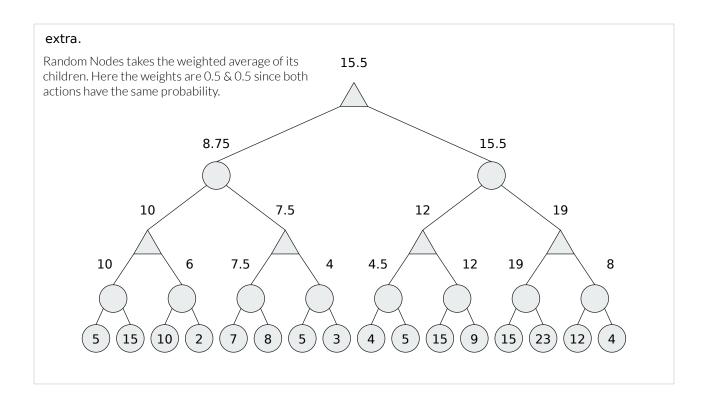


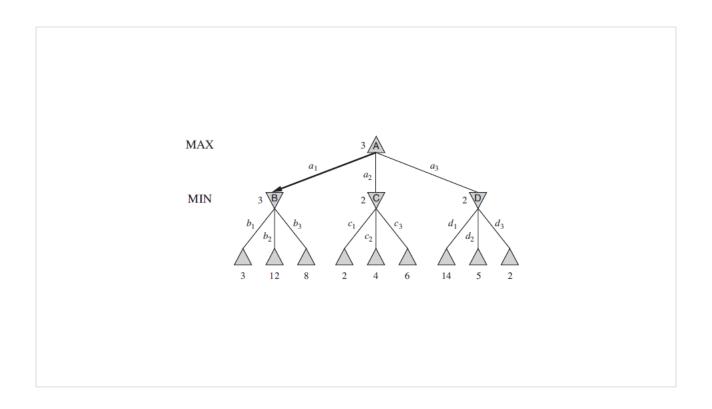












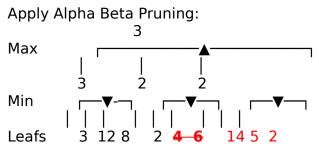




Figure 5.17 The starting position of a simple game. Player A moves first. The two players take turns moving, and each player must move his token to an open adjacent space in either direction. If the opponent occupies an adjacent space, then a player may jump over the opponent to the next open space if any. (For example, if A is on 3 and B is on 2, then A may move back to 1.) The game ends when one player reaches the opposite end of the board. If player A reaches space 4 first, then the value of the game to A is +1; if player B reaches space 1 first, then the value of the game to A is -1.

- 5.8 Consider the two-player game described in Figure 5.17.
 - a. Draw the complete game tree, using the following conventions:
 - Write each state as (s_A, s_B), where s_A and s_B denote the token locations.
 - Put each terminal state in a square box and write its game value in a circle.
 - Put loop states (states that already appear on the path to the root) in double square boxes. Since their value is unclear, annotate each with a "?" in a circle.
 - b. Now mark each node with its backed-up minimax value (also in a circle). Explain how you handled the "?" values and why.
 - c. Explain why the standard minimax algorithm would fail on this game tree and briefly sketch how you might fix it, drawing on your answer to (b). Does your modified algorithm give optimal decisions for all games with loops?
 - **d**. This 4-square game can be generalized to n squares for any n > 2. Prove that A wins if n is even and loses if n is odd.

b.

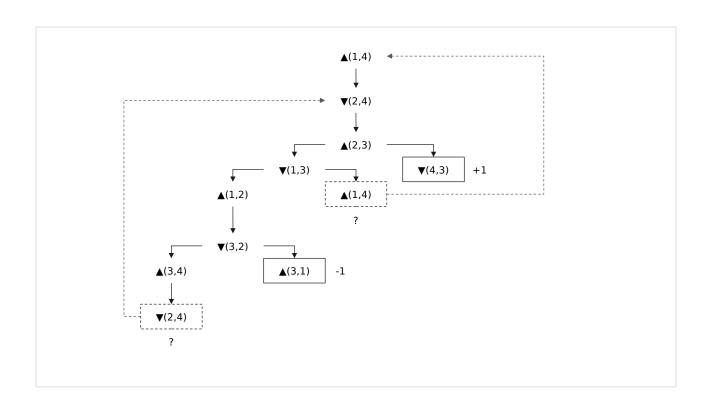
Since there are two possible values only {-1, 1}, we can do:

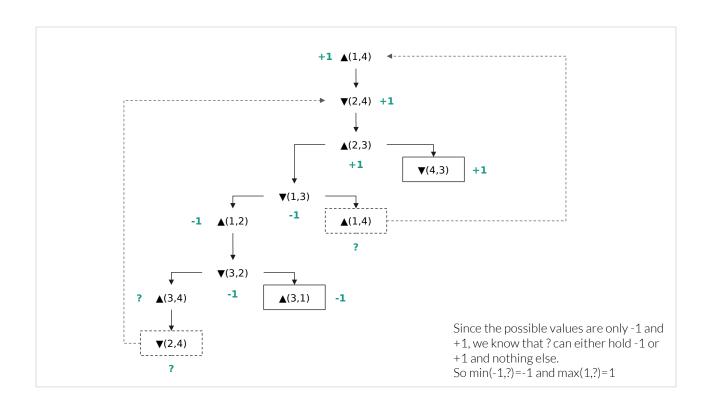
$$\max(?,?) = ?, \max(?,-1) = ?, \max(?,1) = 1$$
 $\max(-1,-1) = -1, \max(-1,1) = 1$
 $\max(1,1) = 1$
 $\min(?,?) = ?, \min(?,-1) = -1, \min(?,1) = ?$
 $\min(-1,-1) = -1, \min(-1,1) = -1$
 $\min(1,1) = 1$

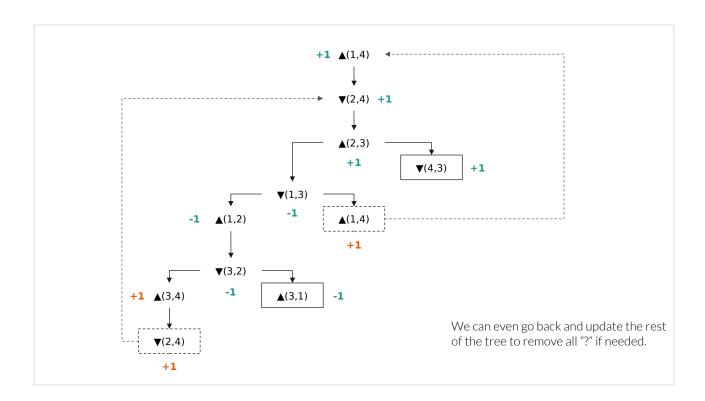
c. The standard minimax algorithm would get stuck in an infinite loop. One way to fix it is to store states along the path to the root and whenever we visit a loop node, we return with "?" and handle it as shown in "b". This will not handle all cases since we may end up propagating "?" to the root and fail to get an answer. Also we have no way to handle "?" if the known values are not binary (not just -1 & 1).

d.

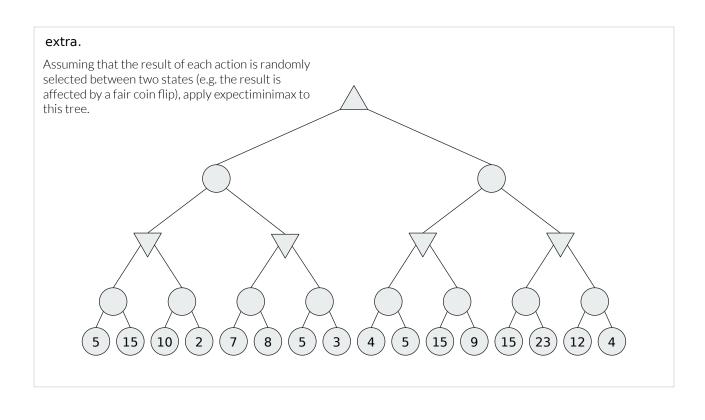
This question is a little tricky. One approach is a proof by induction on the size of the game. Clearly, the base case n=3 is a loss for A and the base

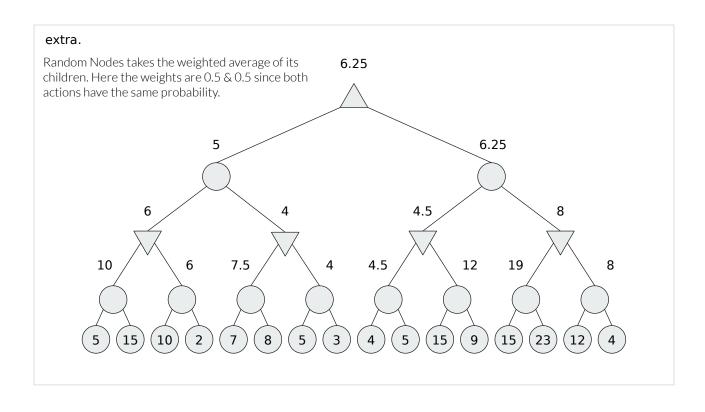


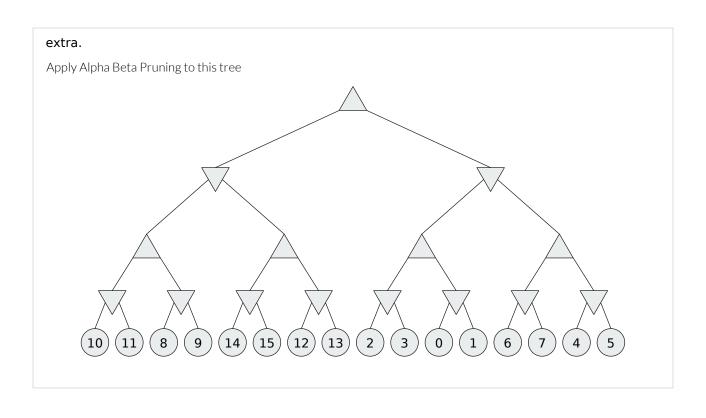


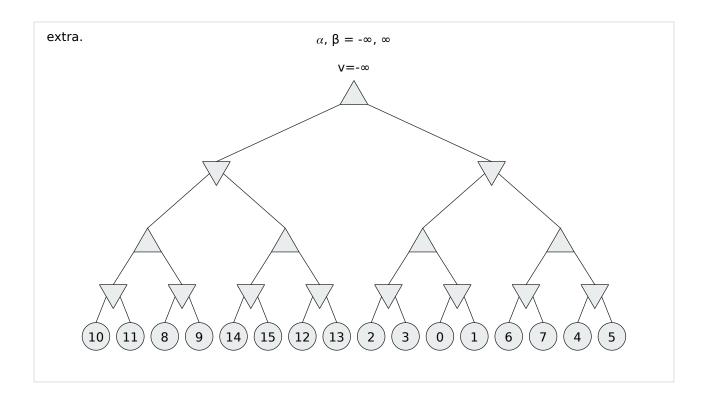


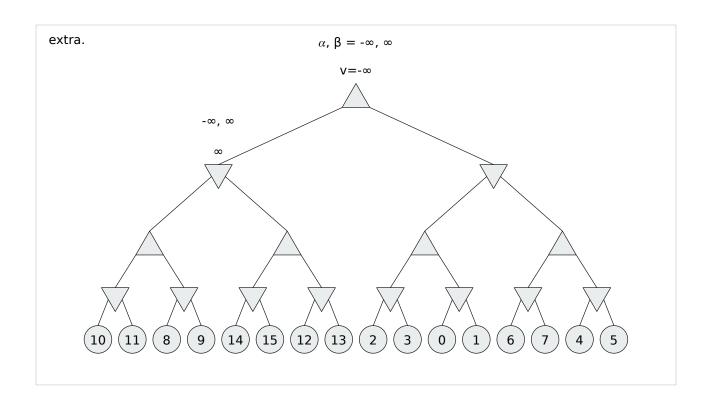
More Questions

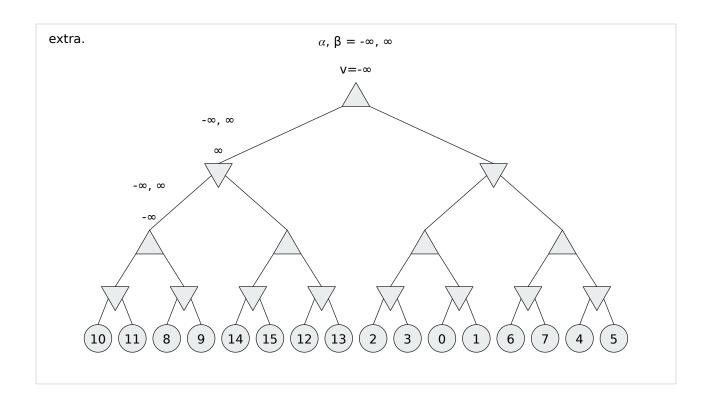


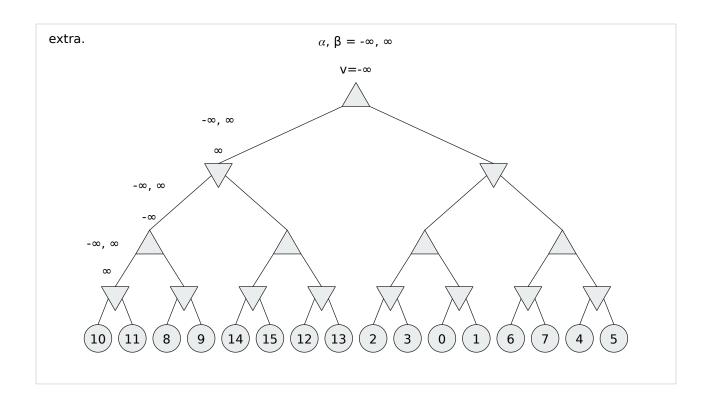


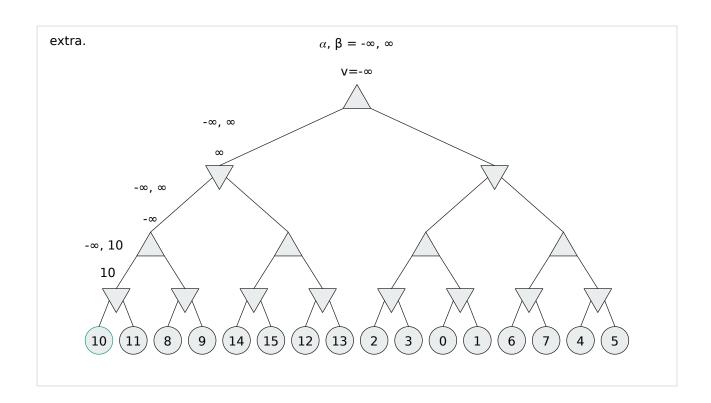


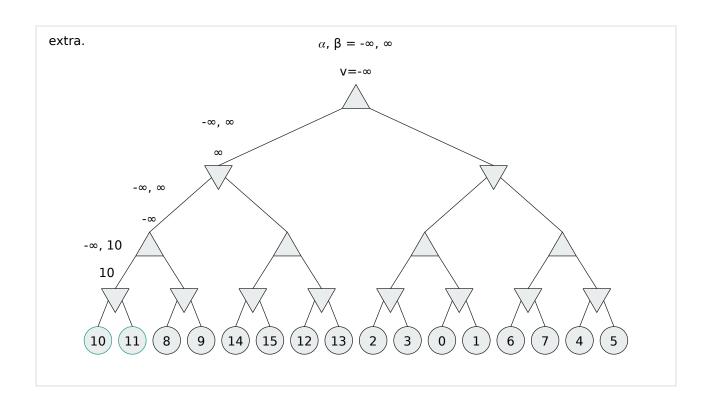


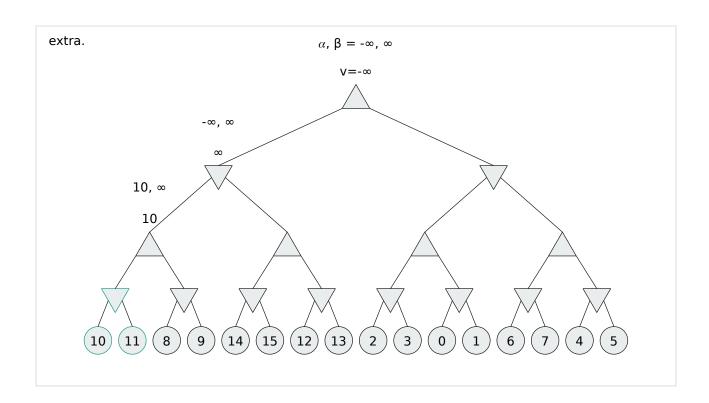


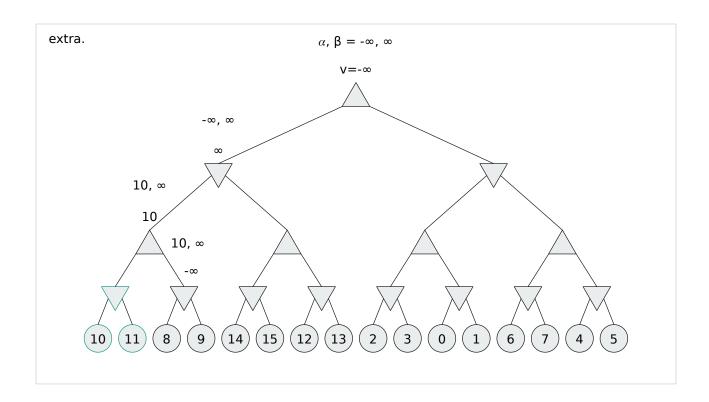


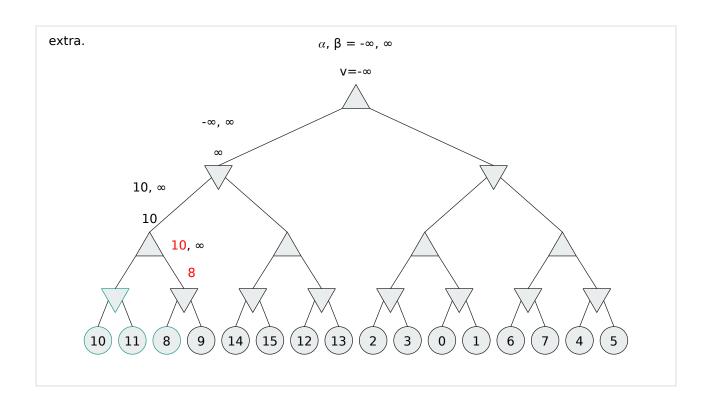


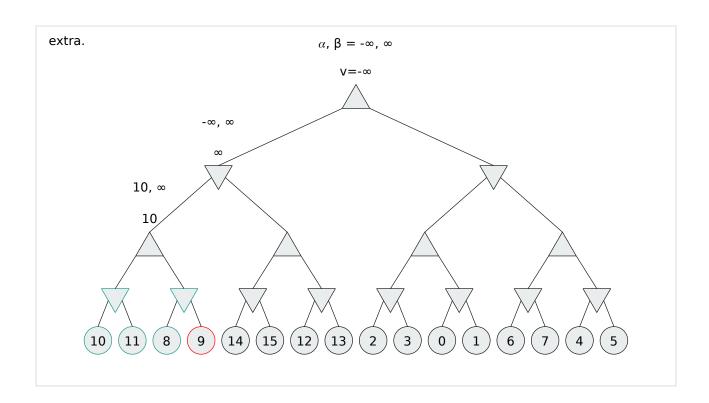


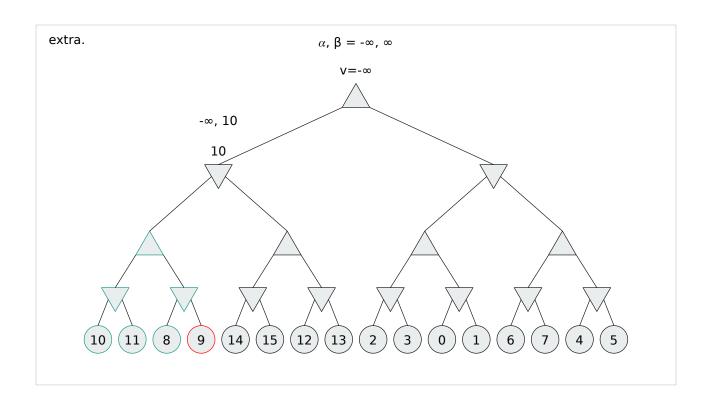


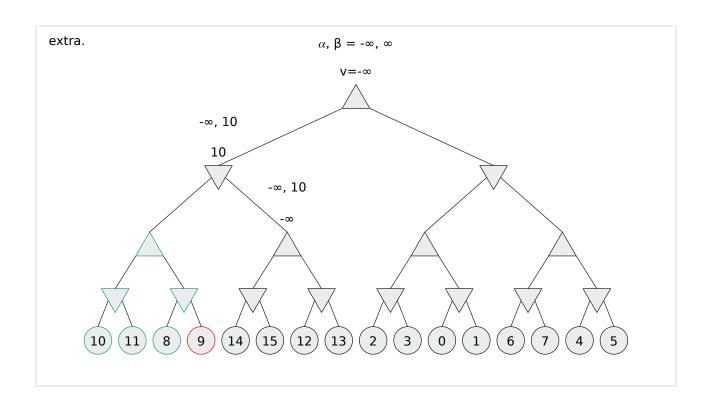


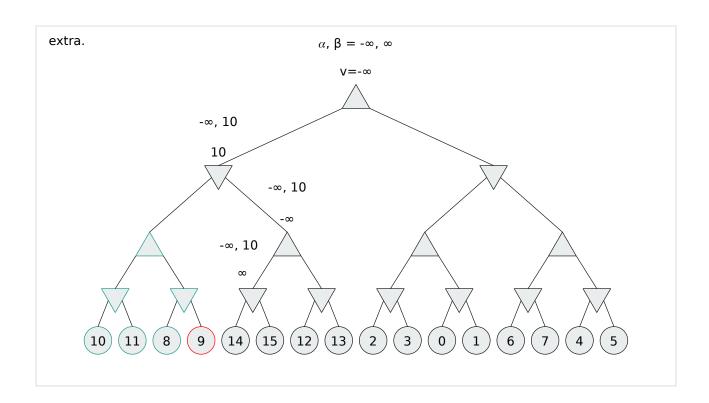


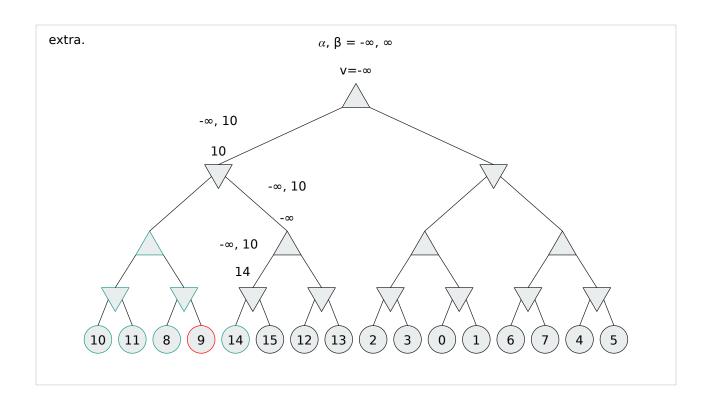


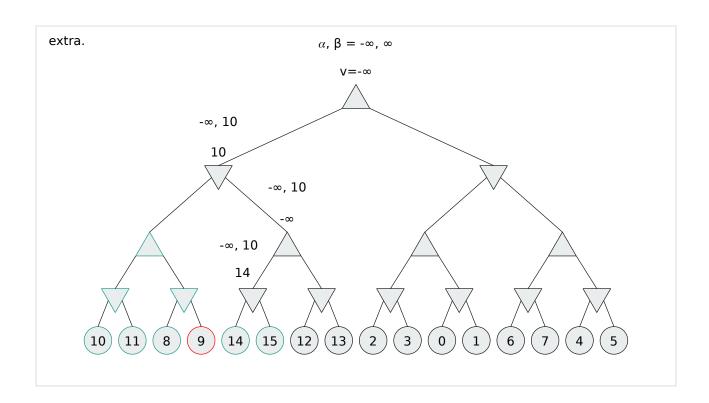


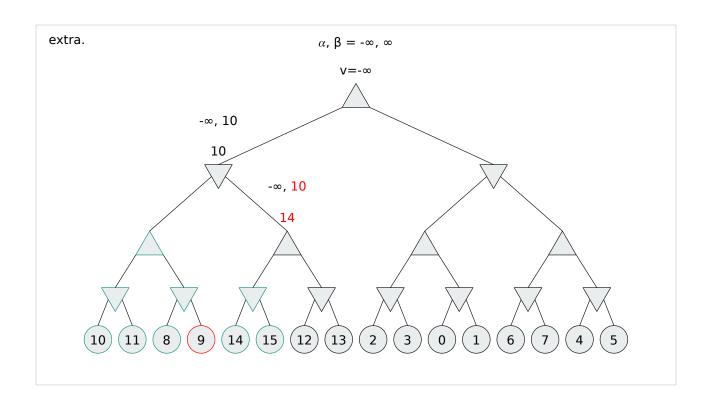


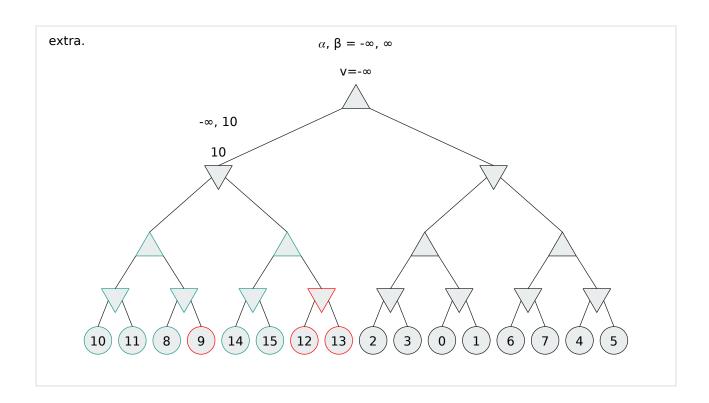


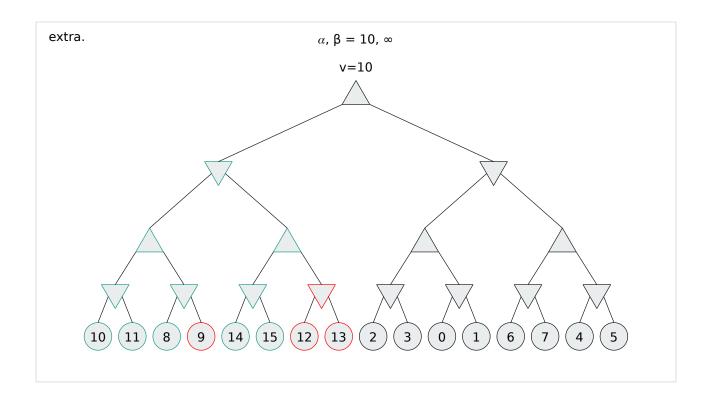


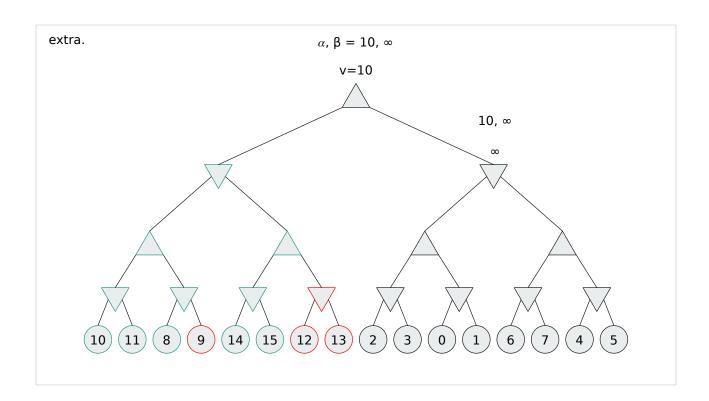


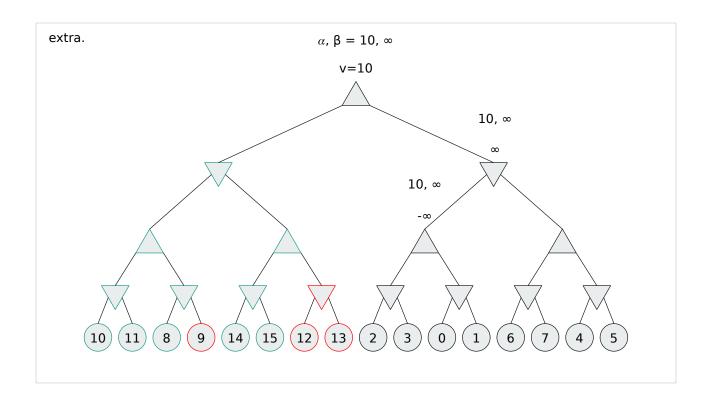


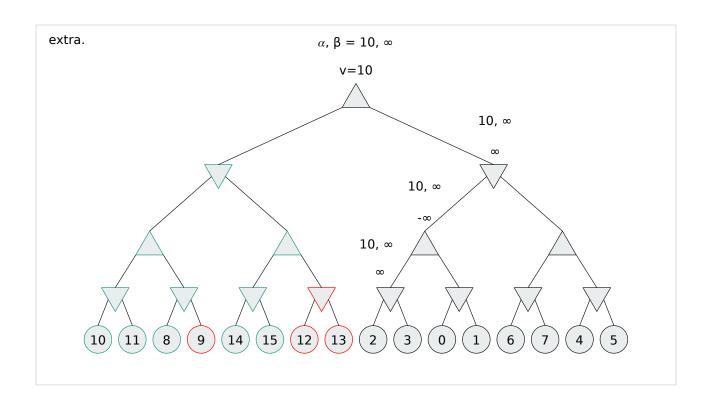


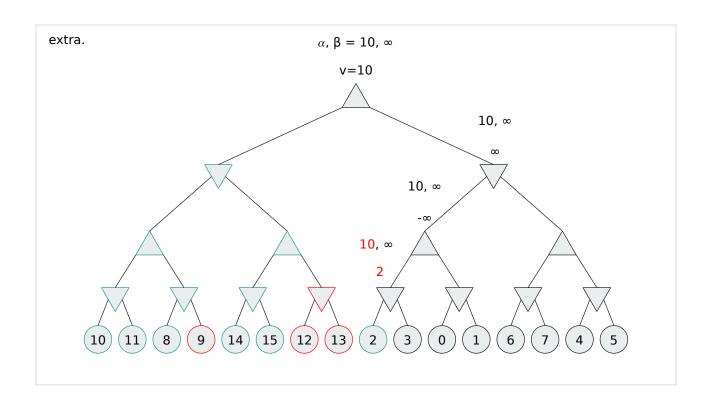


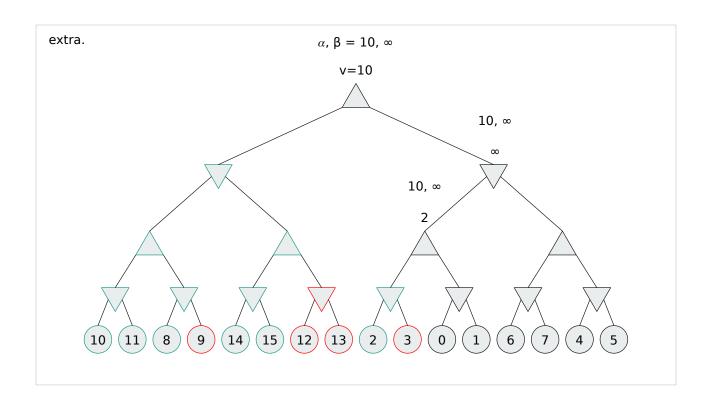


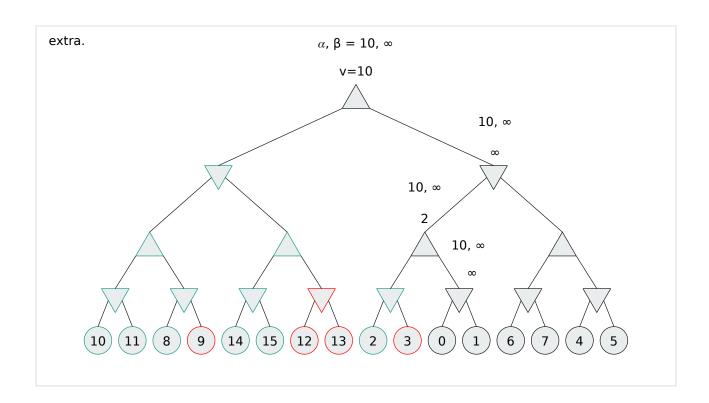


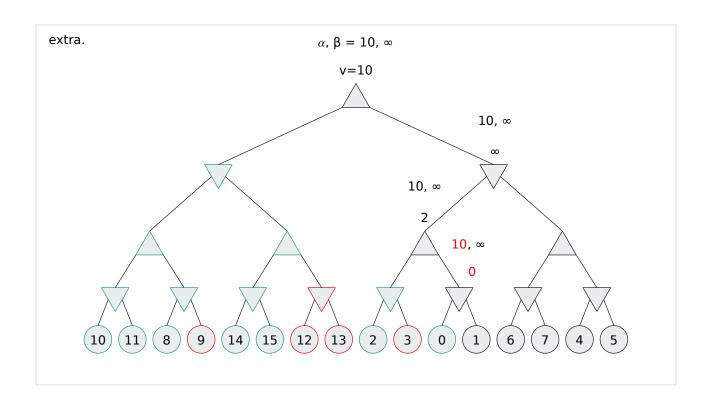


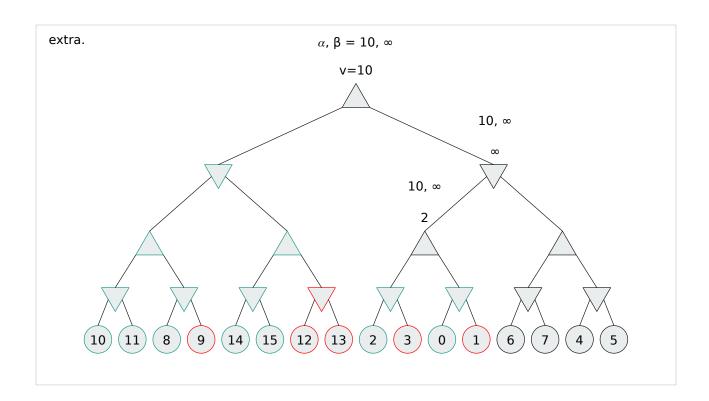


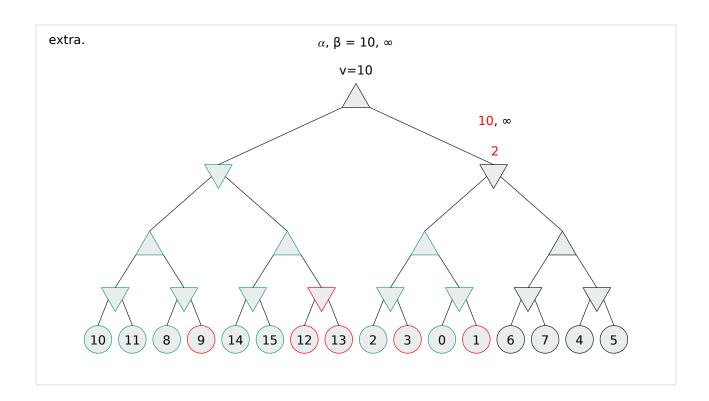


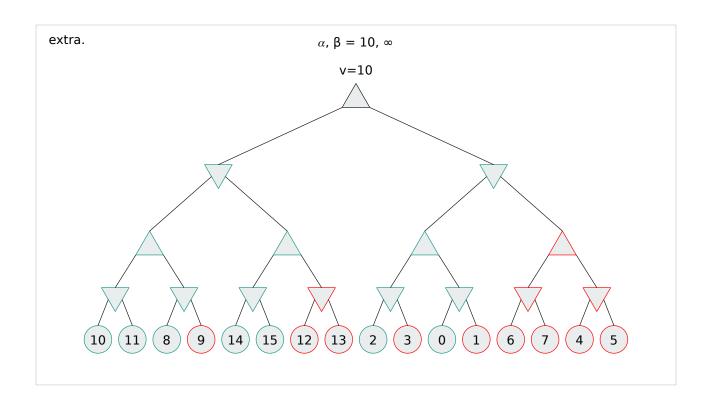


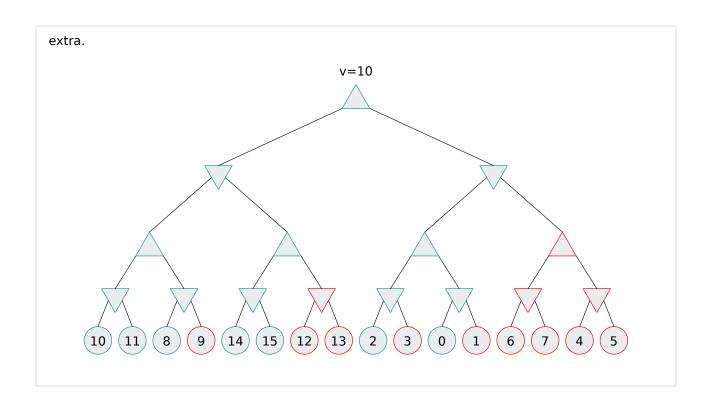












NEXT WEEK QUIZ on "GAMES"