

Machine learning II

Learning theory

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Break point

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Proof that $m_{\mathcal{H}}(N)$ can replace M

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Main result

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No break point $\implies m_{\mathcal{H}}(N) = 2^N$

Any break point $\implies m_{\mathcal{H}}(N)$ is **polynomial** in N

Break point

k is a break point if $m_{\mathcal{H}}(k) < 2^k$

	1	2	3	N 4	5	...
2-D perceptron	2	4	8	14	...	
1-D pos. ray	2	3	4	5	...	
2-D pos. rectangles	2	4	8	16	$< 2^5$...

Quiz I

For every set of k^* points $\mathbf{x}_1, \dots, \mathbf{x}_{k^*}$, \mathcal{H} implements $< 2^{k^*}$ dichotomies. Which statement(s) is (are) true?

1. k^* is a break point
2. k^* is not a break point
3. all $k \geq k^*$ are break points
4. all $k < k^*$ are break points
5. this has nothing to do with break points!

Quiz I

For every set of k^* points $\mathbf{x}_1, \dots, \mathbf{x}_{k^*}$, \mathcal{H} implements $< 2^{k^*}$ dichotomies:

1. k^* is a break point.
2. k^* is not a break point.
3. all $k \geq k^*$ are break points.
4. all $k < k^*$ are break points.
5. this has nothing to do with break points!

Quiz II

To show that k is **not** a break point for \mathcal{H} :

1. Show a set of k points $\mathbf{x}_1, \dots, \mathbf{x}_k$ which \mathcal{H} can shatter.
2. Show \mathcal{H} can shatter any set of k points.
3. Show a set of k points $\mathbf{x}_1, \dots, \mathbf{x}_k$ which \mathcal{H} cannot shatter.
4. Show \mathcal{H} cannot shatter any set of k points.
5. Show $m_{\mathcal{H}}(k) = 2^k$.

Quiz II

To show that k **is not** a break point for \mathcal{H} :

1. Show a set of k points $\mathbf{x}_1, \dots, \mathbf{x}_k$ which \mathcal{H} can shatter.
2. Show \mathcal{H} can shatter any set of k points. (Overkill)
3. Show a set of k points $\mathbf{x}_1, \dots, \mathbf{x}_k$ which \mathcal{H} cannot shatter.
4. Show \mathcal{H} cannot shatter any set of k points.
5. Show $m_{\mathcal{H}}(k) = 2^k$.

Quiz III

To show that k is a break point for \mathcal{H} :

1. Show a set of k points $\mathbf{x}_1, \dots, \mathbf{x}_k$ which \mathcal{H} can shatter.
2. Show \mathcal{H} can shatter any set of k points.
3. Show a set of k points $\mathbf{x}_1, \dots, \mathbf{x}_k$ which \mathcal{H} cannot shatter.
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Quiz III

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4. Show \mathcal{H} cannot shatter any set of k points.
5. Show $m_{\mathcal{H}}(k) = 2^k$.

Back to our combinatorial puzzle

X_1	X_2	X_3	X_4
○	○	○	○
○	○	○	●
○	○	●	○
○	●	○	○
●	○	○	○

Try to add a 6th dichotomy.

We cannot add a 6th dichotomy

x_1	x_2	x_3	x_4
○	○	○	○
○	○	○	●
○	○	●	○
○	●	○	○
●	○	○	○
○	●	●	○

If $k = 2$ is a break point, the maximum number of dichotomies on $N = 4$ points is 5.

Intuition: any break point implies a huge combinatorial restriction, i.e. an enormous constraint on the number of dichotomies. The number of dichotomies, which is equal to 2^N (without a break point) will reduce to a polynomial (with a break point).

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Bounding $m_{\mathcal{H}}(N)$

- To show $m_{\mathcal{H}}(N)$ is polynomial, we will show that

$$m_{\mathcal{H}}(N) \leq \dots \leq \dots \leq \text{a polynomial}$$

- The key quantity is $B(N, k)$ which gives the maximum number of dichotomies on N points with break point k .
- How many dichotomies can you list on 4 points with break point 2?

x_1	x_2	x_3
○	○	○
○	○	●
○	●	○
●	○	○

$$B(3, 2) = 4 (< 8)$$

x_1	x_2	x_3	x_4
○	○	○	○
○	○	○	●
○	○	●	○
○	●	○	○
●	○	○	○

$$B(4, 2) = 5 (< 16)$$

$m_{\mathcal{H}}(N)$ is bounded by $B(N, k)$

Suppose that \mathcal{H} has a break point at k . Then

$$m_{\mathcal{H}}(N) \leq B(N, k).$$

\mathbf{x}_1	\mathbf{x}_2	\mathbf{x}_3	\mathbf{x}_4	\dots	\mathbf{x}_N
○	○	○	○	\dots	●
○	○	○	●	\dots	○
○	○	●	○	\dots	○
○	●	○	○	\dots	○
●	○	○	○	\dots	●
○	○	●	●	\dots	●
○	●	○	●	\dots	○
\vdots	\vdots	\vdots	\vdots	\dots	\vdots

- Consider any k points. They cannot be shattered (otherwise k would not be a break point)
- $B(N, k)$ is largest such list

$m_{\mathcal{H}}(N)$ is bounded by $B(N, k)$

Suppose that \mathcal{H} has a break point at k . Then

$$m_{\mathcal{H}}(N) \leq B(N, k).$$

x_1	x_2	x_3	x_4	...	x_N
○	○	○	○	...	●
○	○	○	●	...	○
○	○	●	○	...	○
○	●	○	○	...	○
●	○	○	○	...	●
○	○	●	●	...	●
○	●	○	●	...	○
⋮	⋮	⋮	⋮	...	⋮

- Consider any k points. They cannot be shattered (otherwise k would not be a break point)
- $B(N, k)$ is largest such list

How can we bound $B(N, k)$?

$B(N, k)$ for boundary cases

- $B(N, 1) = ?$
- $B(1, k) = ?$ for $k > 1$
- $B(N, N) = ?$

$B(N, k)$ for boundary cases

- $B(N, 1) = ?$
- $B(1, k) = ?$ for $k > 1$
- $B(N, N) = ?$

- $B(N, 1) = 1$
- $B(1, k) = 2$ for $k > 1$
- $B(N, N) = 2^N - 1$

We then assume $N \geq 2$ and $k \geq 2$ and try to develop a **recursion**.

Let us try to bound $B(4, 3)$

How many dichotomies can you list on 4 points with break point 3 (i.e. no subset of 3 is shattered).

x_1	x_2	x_3	x_4
○	○	○	○
○	○	○	●
○	○	●	○
○	●	○	○
●	○	○	○
○	○	●	●
○	●	○	●
●	○	○	●
○	●	●	○
●	○	●	○
●	●	○	○

We need to find a recursion, i.e. bound $B(4, 3)$ using $B(3, \cdot)$.

Let us try to bound $B(4, 3)$

x_1	x_2	x_3	x_4
○	○	○	○
○	○	○	●
○	○	●	○
○	●	○	○
●	○	○	○
○	○	●	●
○	●	○	●
●	○	○	●
○	●	●	○
●	○	●	○
●	●	○	○

- Some dichotomies (out of the 2^4) are missing because of the break point constraint ($k = 3$). The remaining dichotomies are unique.
- What can we say about the prefix (x_1, x_2, x_3) of these remaining dichotomies?

Two types of dichotomies

Prefix (x_1, x_2, x_3) appears **once** or **twice**.

x_1	x_2	x_3	x_4
○	○	○	○
○	○	○	●
○	○	●	○
○	●	○	○
●	○	○	○
○	○	●	●
○	●	○	●
●	○	○	●
○	●	●	○
●	○	●	○
●	●	○	○

→ Let us reorder the dichotomies to simplify the counting.

Reordering the dichotomies

- α : Prefix $(\mathbf{x}_1, \mathbf{x}_2, \mathbf{x}_3)$ appears **once**
- β : Prefix $(\mathbf{x}_1, \mathbf{x}_2, \mathbf{x}_3)$ appears **twice**

	\mathbf{x}_1	\mathbf{x}_2	\mathbf{x}_3	\mathbf{x}_4
α	○	●	●	○
	●	○	●	○
	●	●	○	○
β	○	○	○	○
	○	○	●	○
	○	●	○	○
	●	○	○	○
β	○	○	○	●
	○	○	●	●
	○	●	○	●
	●	○	○	●

Reordering the dichotomies

	x_1	x_2	x_3	x_4
α	○	●	●	○
	●	○	●	○
	●	●	○	○
β	○	○	○	○
	○	○	●	○
	○	●	○	○
	●	○	○	○
β	○	○	○	●
	○	○	●	●
	○	●	○	●
	●	○	○	●

- α : Prefix (x_1, x_2, x_3) appears **once**
- β : Prefix (x_1, x_2, x_3) appears **twice**
- $B(4, 3) = \alpha + 2\beta$

Reordering the dichotomies

	x_1	x_2	x_3	x_4
α	○	●	●	○
	●	○	●	○
	●	●	○	○
β	○	○	○	○
	○	○	●	○
	○	●	○	○
	●	○	○	○
β	○	○	○	●
	○	○	●	●
	○	●	○	●
	●	○	○	●

- α : Prefix (x_1, x_2, x_3) appears **once**
- β : Prefix (x_1, x_2, x_3) appears **twice**
- $B(4, 3) = \alpha + 2\beta$
- Strategy for bounding $B(4, 3) = \alpha + 2\beta$:

$$\alpha + \beta \leq Q$$

$$\beta \leq R$$

$$\implies \alpha + 2\beta \leq Q + R$$

Reordering the dichotomies

	x_1	x_2	x_3	x_4
α	○	●	●	○
	●	○	●	○
	●	●	○	○
β	○	○	○	○
	○	○	●	○
	○	●	○	○
	●	○	○	○
β	○	○	○	●
	○	○	●	●
	○	●	○	●
	●	○	○	●

- α : Prefix (x_1, x_2, x_3) appears **once**
- β : Prefix (x_1, x_2, x_3) appears **twice**
- $B(4, 3) = \alpha + 2\beta$
- Strategy for bounding $B(4, 3) = \alpha + 2\beta$:

$$\alpha + \beta \leq Q$$

$$\beta \leq R$$

$$\implies \alpha + 2\beta \leq Q + R$$

- $Q = ?, R = ? \implies$ exploit the fact that $k = 3$ is a break point.

First, bound $\alpha + \beta$

	x_1	x_2	x_3	x_4
α	○	●	●	○
	●	○	●	○
	●	●	○	○
β	○	○	○	○
	○	○	●	○
	○	●	○	○
	●	○	○	○
β	○	○	○	●
	○	○	●	●
	○	●	○	●
	●	○	○	●

● $\alpha + \beta \leq ?$

First, bound $\alpha + \beta$

	x_1	x_2	x_3	x_4
α	○	●	●	○
	●	○	●	○
	●	●	○	○
β	○	○	○	○
	○	○	●	○
	○	●	○	○
	●	○	○	○
β	○	○	○	●
	○	○	●	●
	○	●	○	●
	●	○	○	●

- $\alpha + \beta \leq ?$
- $\alpha + \beta$: the total number of different dichotomies on the first 3 points
- No subset of $k = 3$ of these first 3 points can be shattered (since no 3-subset of all 4 points can be shattered).

First, bound $\alpha + \beta$

	x_1	x_2	x_3	x_4
α	○	●	●	○
	●	○	●	○
	●	●	○	○
β	○	○	○	○
	○	○	●	○
	○	●	○	○
	●	○	○	○
β	○	○	○	●
	○	○	●	●
	○	●	○	●
	●	○	○	●

- $\alpha + \beta \leq ?$
 - $\alpha + \beta$: the total number of different dichotomies on the first 3 points
 - No subset of $k = 3$ of these first 3 points can be shattered (since no 3-subset of all 4 points can be shattered).
- $\implies \alpha + \beta \leq B(3, 3)$

Second, bound β

	\mathbf{x}_1	\mathbf{x}_2	\mathbf{x}_3	\mathbf{x}_4
α	○	●	●	○
	●	○	●	○
	●	●	○	○
β	○	○	○	○
	○	○	●	○
	○	●	○	○
	●	○	○	○
β	○	○	○	●
	○	○	●	●
	○	●	○	●
	●	○	○	●

● $\beta \leq ?$

Second, bound β

	x_1	x_2	x_3	x_4
α	○	●	●	○
	●	○	●	○
	●	●	○	○
β	○	○	○	○
	○	○	●	○
	○	●	○	○
	●	○	○	○
β	○	○	○	●
	○	○	●	●
	○	●	○	●
	●	○	○	●

- $\beta \leq ?$
- If 2 points are shattered, then using the mirror dichotomies you shatter 3 points

Second, bound β

	x_1	x_2	x_3	x_4
α	○	●	●	○
	●	○	●	○
	●	●	○	○
β	○	○	○	○
	○	○	●	○
	○	●	○	○
	●	○	○	○
β	○	○	○	●
	○	○	●	●
	○	●	○	●
	●	○	○	●

- $\beta \leq ?$
- If 2 points are shattered, then using the mirror dichotomies you shatter 3 points
- $\beta \leq B(3, 2)$

Combining, to bound $\alpha + 2\beta$

	x_1	x_2	x_3	x_4
α	○	●	●	○
	●	○	●	○
	●	●	○	○
β	○	○	○	○
	○	○	●	○
	○	●	○	○
	●	○	○	○
β	○	○	○	●
	○	○	●	●
	○	●	○	●
	●	○	○	●

- $B(4, 3) = \alpha + \beta + \beta \leq B(3, 3) + B(3, 2)$
- The argument generalizes to (N, k) :

$$B(N, k) \leq B(N - 1, k) + B(N - 1, k - 1)$$

Numerical computation of $B(N, k)$ bound

Numerical computation of $B(N, k)$ bound

$$B(N, k) \leq B(N-1, k) + B(N-1, k-1)$$



		k						
		1	2	3	4	5	6	..
N	1	1	2	2	2	2	2	..
	2	1	3	4	4	4	4	..
	3	1	4	7	8	8	8	..
	4	1	5	11
	5	1	6	:	.			
	6	1	7	:		.		
	:	:	:	:			.	

Analytic solution for $B(N, k)$ bound

Sauer's Lemma:

$$B(N, k) \leq \sum_{i=0}^{k-1} \binom{N}{i}$$

Proof by induction on N .

Proof of Sauer's Lemma

1. Verify it for the boundary conditions
 - The statement is true whenever $k = 1$ or $N = 1$, by inspection.
 - Since the statement is already true when $k = 1$ (for all values of N) by the initial condition, we only need to worry about $k \geq 2$.
2. Suppose $B(N, k) \leq \sum_{i=0}^{k-1} \binom{N}{i}$ for $k \geq 2$
3. Assuming (2), show that $B(N + 1, k) \leq \sum_{i=0}^{k-1} \binom{N+1}{i}$ for $k \geq 2$.
 - You can use the following Lemma: $\binom{N+1}{i} = \binom{N}{i} + \binom{N}{i-1}$. In other words, when choosing i objects from $N + 1$ objects, either the first object is not included, in $\binom{N}{i}$ ways, or the first object is included in $\binom{N}{i-1}$ ways.

Proof of Sauer's Lemma

$$\begin{aligned} B(N+1, k) &\leq B(N, k) + B(N, k-1) \\ &\leq \sum_{i=0}^{k-1} \binom{N}{i} + \sum_{i=0}^{k-2} \binom{N}{i} \\ &= \sum_{i=0}^{k-1} \binom{N}{i} + \sum_{i=1}^{k-1} \binom{N}{i-1} \\ &= 1 + \sum_{i=1}^{k-1} \left[\binom{N}{i} + \binom{N}{i-1} \right] \\ &= 1 + \sum_{i=1}^{k-1} \binom{N+1}{i} \\ &= \sum_{i=0}^{k-1} \binom{N+1}{i} \end{aligned}$$

It is polynomial!

Theorem. If k is any break point for \mathcal{H} , i.e. $m_{\mathcal{H}}(k) < 2^k$, then

$$m_{\mathcal{H}}(k) \leq B(N, k) \leq \sum_{i=0}^{k-1} \binom{N}{i}.$$

But, we also have that

$$\sum_{i=0}^{k-1} \binom{N}{i} \leq \begin{cases} N^{k-1} + 1 \\ \left(\frac{eN}{k-1}\right)^{k-1} \end{cases} \quad (\text{polynomial in } N)$$

→ If we can replace M by $m_{\mathcal{H}}(N)$ in the bound, then learning is feasible.

Three examples

$$\sum_{i=0}^{k-1} \binom{N}{i}$$

- \mathcal{H} is **positive rays**: (break point $k = 2$)

$$m_{\mathcal{H}}(N) = N + 1 \leq N + 1$$

- \mathcal{H} is **positive intervals**: (break point $k = 3$)

$$m_{\mathcal{H}}(N) = \frac{1}{2}N^2 + \frac{1}{2}N + 1 \leq \frac{1}{2}N^2 + \frac{1}{2}N + 1$$

- \mathcal{H} is **2D perceptrons**: (break point $k = 4$)

$$m_{\mathcal{H}}(N) = ? \leq \frac{1}{6}N^3 + \frac{5}{6}N + 1$$

Proof that $m_{\mathcal{H}}(N)$ can replace M

Proof that $m_{\mathcal{H}}(N)$ can replace M

What we want

Instead of:

$$\mathbb{P}[|E_{\text{in}}(g) - E_{\text{out}}(g)| > \epsilon] \leq 2 \quad \textcolor{red}{M} \quad e^{-2\epsilon^2 N}$$

We want:

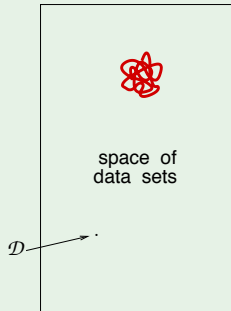
$$\mathbb{P}[|E_{\text{in}}(g) - E_{\text{out}}(g)| > \epsilon] \leq 2 \quad \textcolor{red}{m_{\mathcal{H}}(N)} \quad e^{-2\epsilon^2 N}$$

Pictorial proof ☺

- How does $m_{\mathcal{H}}(N)$ relate to overlaps?
- What to do about E_{out} ?
- Putting it together

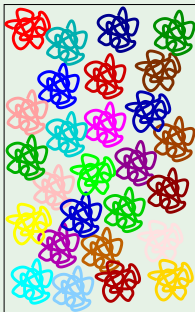
Overlaps between events

Hoeffding Inequality



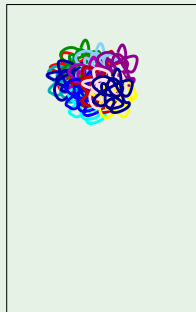
(a)

Union Bound



(b)

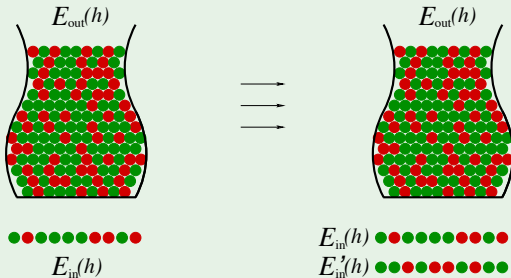
VC Bound



(c)

What to do about E_{out} ?

What to do about E_{out}



The Vapnik-Chervonenkis (VC) Inequality

Putting it together

Not quite:

$$\mathbb{P}[|E_{\text{in}}(g) - E_{\text{out}}(g)| > \epsilon] \leq 2 m_{\mathcal{H}}(N) e^{-2\epsilon^2 N}$$

but rather:

$$\mathbb{P}[|E_{\text{in}}(g) - E_{\text{out}}(g)| > \epsilon] \leq 4 m_{\mathcal{H}}(2N) e^{-\frac{1}{8}\epsilon^2 N}$$

The Vapnik-Chervonenkis Inequality