

## Uninformed Search

Chapter 3.1 – 3.5

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## Building Goal-Based Agents

What are the key questions needing to be addressed?

- What **goal** does the agent need to achieve?
- What **knowledge** does the agent need?
- What **actions** does the agent need to do?

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## Building Goal-Based Agents

What goal does the agent need to achieve?

- How do you describe the goal?
  - as a task to be accomplished
  - as a situation to be reached
  - as a set of properties to be acquired
- How do you know when the goal is reached?
  - with a **goal test** that defines what it means to have achieved/satisfied the goal
- \* **Determining the goal is difficult and is usually left to the system designer or user to specify.**

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## Building Goal-Based Agents

What knowledge does the agent need?

- The info. needs to be
  - sufficient to describe all relevant aspects to reaching the goal
  - adequate to describe the world state/situation
- **We'll use a *closed world assumption*:**  
*All necessary information about a problem domain is accessible in each percept so that each state is a complete description of the world.*  
*There is no incomplete information at any point in time.*

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## Building Goal-Based Agents

- How should the agent's knowledge be represented?
- knowledge representation problem:
  - What information from the raw percept is relevant?
  - How to represent domain knowledge later...
- \* *Determining what to represent is difficult and is usually left to the system designer to specify.*

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## Building Goal-Based Agents

### What actions does the agent need to do?

- The set of actions/events needs to be
  - decomposed into primitive steps that are discrete, i.e. treated as instantaneous
  - quantified to fully describe initial and final states with no uncertainty
  - sufficient to describe all necessary changes
- \* *The number of actions needed depends on how the world states are represented.*

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## Building Goal-Based Agents

### What actions does the agent need to do?

- Given:
  - an action (operator/move)
  - a description of the current state of the world
- Action completely specifies:
  - if that action can be applied (applicable? legal?)
  - what the exact state of the world will be after the action is performed in the current world (no "history" information needed to compute the new world state).

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## Formalizing Search in a State Space

- A state space is a graph:  $(V, E)$ 
  - $V$  is a set of nodes (vertexes)
  - $E$  is a set of arcs (edges)
    - each arc is directed from one node to another node
- Each node is a data structure that contains:
  - a state description
  - other information such as:
    - link to parent node
    - name of operator that generated this node (from its parent)
    - other bookkeeping data

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## Formalizing Search in a State Space

- Each **arc** corresponds to one of the operators:
  - when the operator is applied to the state associated with the arc's source node
  - then the resulting state is the state associated with the arc's destination node
- Each arc has a fixed, positive **cost**:
  - corresponds to the cost of the operator

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## Formalizing Search in a State Space

- Each node has a set of **successor** nodes:
  - corresponds to all of the legal operators that can be applied at the source node's state
- **Expanding a node** means:
  - generate all of the successor nodes
  - add them and their associated arcs to the state-space search tree

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## Formalizing Search in a State Space

- One or more nodes are designated as **start** nodes
- A **goal test** is applied to a node's state to determine if it is a goal node
- A **solution** is a sequence of operators associated with a path in the state space from a start to a goal node:
  - just the goal state (e.g. cryptarithmic)
  - a path from start to goal state (e.g. 8-puzzle)
- The **cost** of a solution is the sum of the arc costs on the solution path

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## Formalizing Search in a State Space

- **State-space search** is the process of searching through a state space for a solution by making explicit a sufficient portion of an implicit state-space graph to include a goal node
  - initially  $V = \{S\}$ , where  $S$  is the start node
  - then  $S$  is *expanded*, its successors are generated and those nodes are added to  $V$  and the associated arcs are added to  $E$
  - this process continues until a goal node is found

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## Formalizing Search in a State Space

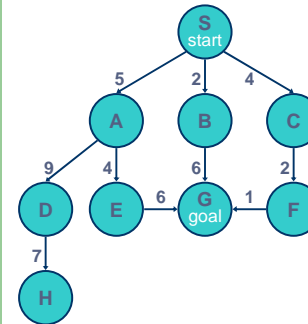
- Each node implicitly or explicitly represents
  - a **partial solution path** from the start node to the given node
  - cost of the partial solution path
- From this node there are:
  - many possible paths that have this partial path as a prefix
  - many possible solutions

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## State Space Diagram



The size of a problem is usually described in terms of the number of possible states:

Tic-Tac-Toe:  $3^9$  states  
 Checkers:  $10^{40}$  states  
 Rubik's Cube:  $10^{19}$  states  
 Chess:  $10^{120}$  states

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## State-Space Search Algorithm

```

//Note: this algorithm doesn't detect loops in the state space
Node generalSearch (Problem problem, List fringe) {
    fringe.add(new Node(problem.getStartState()));
    while (true) {
        if (fringe.isEmpty()) return new Node("failure");
        Node node = fringe.remove();
        if (problem.isGoal(node.getState())) return node;
        fringe.add(expand node given problem operators);

        //expand: generates all of this node's children nodes
        //Note: the goal test is NOT done when nodes are generated
    }
}
    
```

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## Key Issues of State-Space Search Algorithm

- Search process constructs a "search tree"
  - root is the start state
  - leaf nodes are:
    - unexpanded nodes (in the nodes list)
    - "dead ends" (nodes that aren't goals and have no successors because no operators were applicable)
    - goal node is last leaf node found
- Loops in graph may cause "search tree" to be infinite even if state space is small
- Changing the nodes "list" to different data structures leads to different search strategies!

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## Evaluating Search Strategies

- **Completeness**

If a solution exists, will it be found?

- a complete algorithm will find a solution

- **Optimality/Admissibility**

If a solution is found, is it guaranteed to be optimal?

- an admissible algorithm will find a solution with minimum cost

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## Evaluating Search Strategies

- **Time Complexity**

How long does it take to find a solution?

- measured for worst or average case
- measured in number of nodes expanded/tested

- **Space Complexity**

How much space is used by the algorithm?

- measured in terms of the maximum size of the *nodes list* during the search

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## Uninformed Search Strategies

**Uninformed Search:** strategies that order nodes without using any domain specific information

- **BFS: breadth-first search**

- \* *queue (FIFO)* used for the nodes "list"
- remove from front, add to **back**

- **DFS: depth-first search**

- \* *stack (LIFO)* used for the nodes "list"
- remove from front, add to **front**

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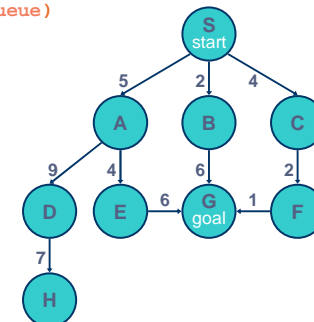
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## Breadth-First Search (BFS)

`generalSearch(problem, queue)`

# of nodes tested: 0, expanded: 0

expnd. node	nodes list
	{S}



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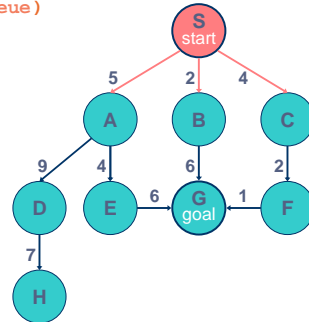
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## Breadth-First Search (BFS)

`generalSearch(problem, queue)`

# of nodes tested: 1, expanded: 1

expnd. node	nodes list
S	{S}
S not goal	{A,B,C}



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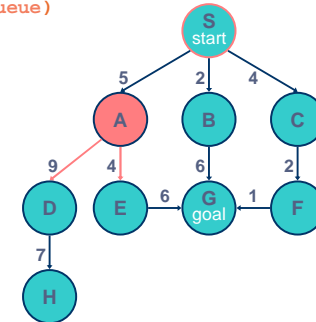
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## Breadth-First Search (BFS)

`generalSearch(problem, queue)`

# of nodes tested: 2, expanded: 2

expnd. node	nodes list
S	{S}
S	{A,B,C}
A not goal	{B,C,D,E}



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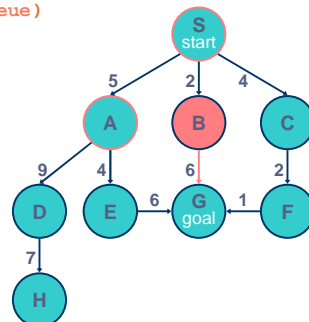
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## Breadth-First Search (BFS)

`generalSearch(problem, queue)`

# of nodes tested: 3, expanded: 3

expnd. node	nodes list
S	{S}
S	{A,B,C}
A	{B,C,D,E}
B not goal	{C,D,E,G}



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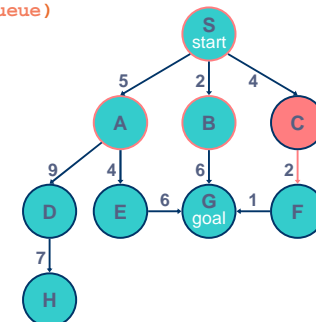
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## Breadth-First Search (BFS)

`generalSearch(problem, queue)`

# of nodes tested: 4, expanded: 4

expnd. node	nodes list
S	{S}
S	{A,B,C}
A	{B,C,D,E}
B	{C,D,E,G}
C not goal	{D,E,G,F}



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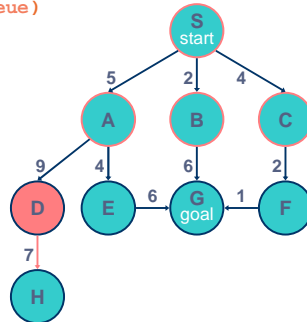
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## Breadth-First Search (BFS)

`generalSearch(problem, queue)`

# of nodes tested: 5, expanded: 5

expnd. node	nodes list
S	{S}
A	{A,B,C}
B	{B,C,D,E}
C	{C,D,E,G}
D not goal	{D,E,G,F}



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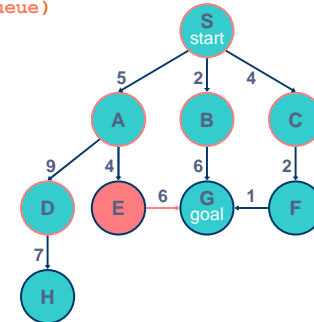
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## Breadth-First Search (BFS)

`generalSearch(problem, queue)`

# of nodes tested: 6, expanded: 6

expnd. node	nodes list
S	{S}
A	{A,B,C}
B	{B,C,D,E}
C	{C,D,E,G}
D	{D,E,G,F}
E not goal	{E,G,F,H}



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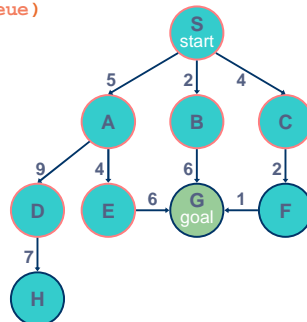
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## Breadth-First Search (BFS)

`generalSearch(problem, queue)`

# of nodes tested: 7, expanded: 6

expnd. node	nodes list
S	{S}
A	{A,B,C}
B	{B,C,D,E}
C	{C,D,E,G}
D	{D,E,G,F}
E	{E,G,F,H}
G goal	{F,H,G} no expand



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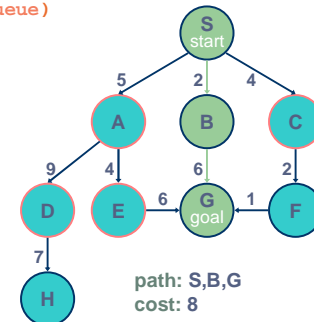
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## Breadth-First Search (BFS)

`generalSearch(problem, queue)`

# of nodes tested: 7, expanded: 6

expnd. node	nodes list
S	{S}
A	{A,B,C}
B	{B,C,D,E}
C	{C,D,E,G}
D	{D,E,G,F}
E	{E,G,F,H}
G	{F,H,G}



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## Breadth-First Search (BFS)

- **Complete**
- **Optimal/Admissible**
  - if all operators (i.e. arcs) have the same cost
  - otherwise, not optimal but does guarantee finding solution of shortest length (i.e. fewest arcs)

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## Breadth-First Search (BFS)

- **Time and space complexity:  $O(b^d)$  (i.e., exponential)**
  - $d$  is the depth of the solution
  - $b$ : the branching factor at each non-leaf node
- \* **Will take a long time to find solutions with a large number of steps because must look at all shorter length possibilities first**

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## Breadth-First Search (BFS)

- **A complete search tree has a total of nodes:**  
 $1 + b + b^2 + \dots + b^d = (b^{(d+1)} - 1) / (b - 1)$ 
  - $d$ : the tree's depth
  - $b$ : the branching factor at each non-leaf node
- **For example:  $d = 12, b = 10$**   
 $1 + 10 + 100 + \dots + 10^{12} = (10^{13} - 1) / 9 = O(10^{12})$ 
  - If BFS expands 1000 nodes/sec and each node uses 100 bytes of storage, then BFS will take 35 years to run in the worst case, and it will use 111 terabytes of memory!

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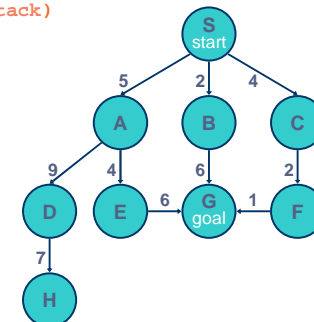
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## Depth-First Search (DFS)

`generalSearch(problem, stack)`

# of nodes tested: 0, expanded: 0

expnd. node	nodes list
	{S}



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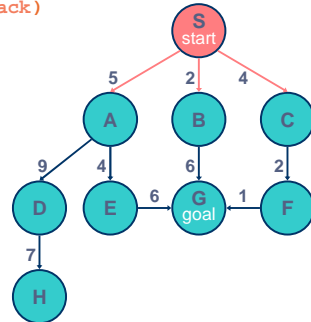


## Depth-First Search (DFS)

`generalSearch(problem, stack)`

# of nodes tested: 1, expanded: 1

expnd. node	nodes list
S	{S}
S not goal	{A,B,C}



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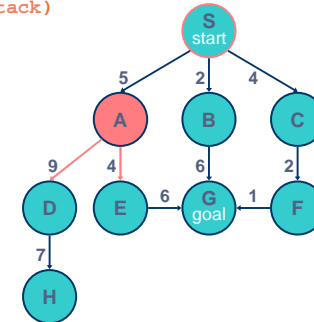
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## Depth-First Search (DFS)

`generalSearch(problem, stack)`

# of nodes tested: 2, expanded: 2

expnd. node	nodes list
S	{S}
S	{A,B,C}
A not goal	{D,E,B,C}



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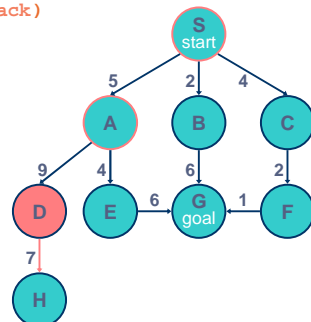
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## Depth-First Search (DFS)

`generalSearch(problem, stack)`

# of nodes tested: 3, expanded: 3

expnd. node	nodes list
S	{S}
S	{A,B,C}
A	{D,E,B,C}
A	{D,E,B,C}
D not goal	{H,E,B,C}



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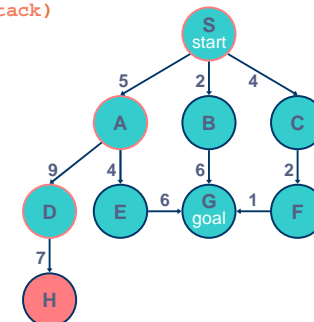
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## Depth-First Search (DFS)

`generalSearch(problem, stack)`

# of nodes tested: 4, expanded: 4

expnd. node	nodes list
S	{S}
S	{A,B,C}
A	{D,E,B,C}
A	{D,E,B,C}
D	{H,E,B,C}
D	{H,E,B,C}
H not goal	{E,B,C}



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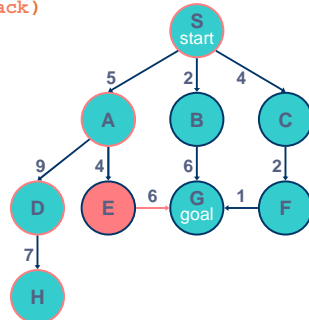
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## Depth-First Search (DFS)

`generalSearch(problem, stack)`

# of nodes tested: 5, expanded: 5

expnd. node	nodes list
S	{S}
A	{A,B,C}
D	{D,E,B,C}
H	{H,E,B,C}
E not goal	{G,B,C}



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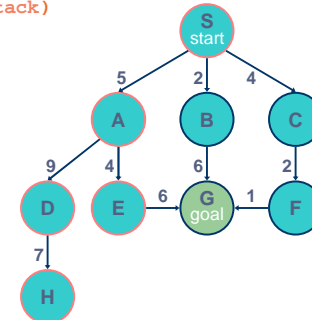
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## Depth-First Search (DFS)

`generalSearch(problem, stack)`

# of nodes tested: 6, expanded: 5

expnd. node	nodes list
S	{S}
A	{A,B,C}
D	{D,E,B,C}
H	{H,E,B,C}
E	{E,B,C}
G goal	{B,C} no expand



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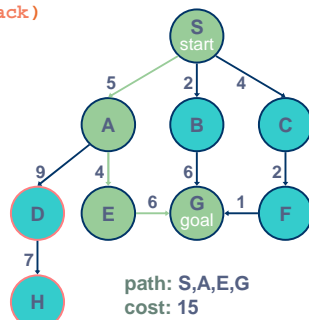
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## Depth-First Search (DFS)

`generalSearch(problem, stack)`

# of nodes tested: 6, expanded: 5

expnd. node	nodes list
S	{S}
A	{A,B,C}
D	{D,E,B,C}
H	{H,E,B,C}
E	{G,B,C}
G	{B,C}



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## Depth-First Search (DFS)

- May not terminate without a **depth bound** i.e., cutting off search below a fixed depth,  $D$

- Not complete
  - with or without cycle detection
  - and, with or without a depth cutoff

- Not optimal/admissible

\* *Can find long solutions quickly if lucky*

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## Depth-First Search (DFS)

- **Time complexity:**  $O(b^d)$  exponential  
**Space complexity:**  $O(bd)$  linear
  - $d$  is the depth of the solution
  - $b$ : the branching factor at each non-leaf node
- **Does chronological backtracking**
  - when search hits a dead end, backs up one level at a time
  - problematic if the mistake occurs because of a bad operator choice near the top of search tree

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## Uninformed Search Strategies

- **UCS: uniform-cost search**
  - \* *priority queue used to order nodes, sort by path cost*
    - let  $g(n)$  = cost of path from start node  $s$  to current node  $n$
    - sort nodes by increasing value of  $g$
    - only uninformed search that worries about costs

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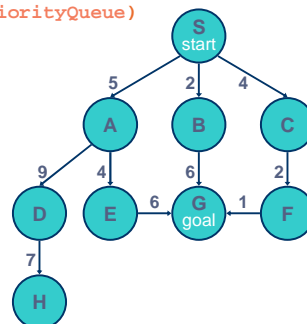
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## Uniform-Cost Search (UCS)

`generalSearch(problem, priorityQueue)`

# of nodes tested: 0, expanded: 0

expnd. node	nodes list
	{S}



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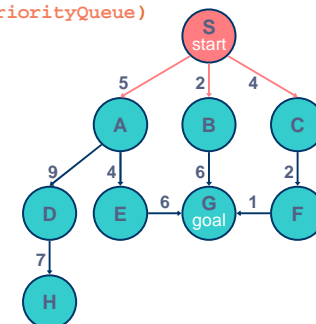
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## Uniform-Cost Search (UCS)

`generalSearch(problem, priorityQueue)`

# of nodes tested: 1, expanded: 1

expnd. node	nodes list
	{S:0}
S not goal	{B:2,C:4,A:5}



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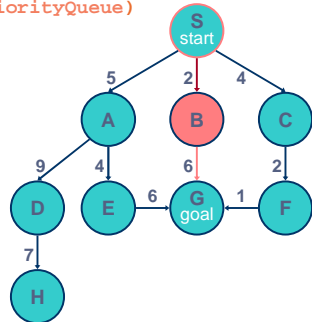
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## Uniform-Cost Search (UCS)

`generalSearch(problem, priorityQueue)`

# of nodes tested: 2, expanded: 2

expnd. node	nodes list
S	{S}
B not goal	{C:4,A:5,G:2+6}



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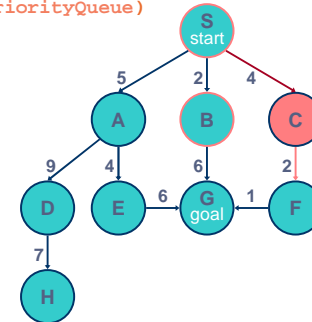
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## Uniform-Cost Search (UCS)

`generalSearch(problem, priorityQueue)`

# of nodes tested: 3, expanded: 3

expnd. node	nodes list
S	{S}
B	{C:4,A:5,G:8}
C not goal	{A:5,F:4+2,G:8}



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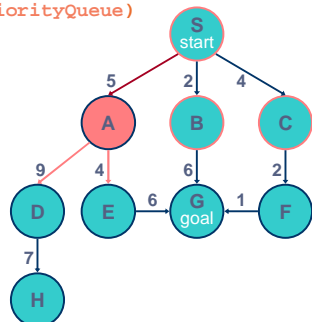
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## Uniform-Cost Search (UCS)

`generalSearch(problem, priorityQueue)`

# of nodes tested: 4, expanded: 4

expnd. node	nodes list
S	{S}
B	{C:4,A:5,G:8}
C	{A:5,F:6,G:8}
A not goal	{F:6,G:8,E:5+4,D:5+9}



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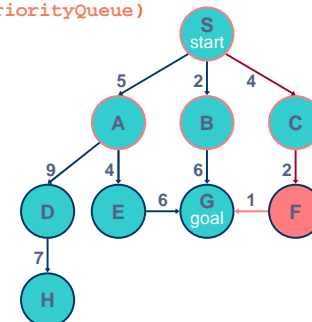
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## Uniform-Cost Search (UCS)

`generalSearch(problem, priorityQueue)`

# of nodes tested: 5, expanded: 5

expnd. node	nodes list
S	{S}
B	{C:4,A:5,G:8}
C	{A:5,F:6,G:8}
A	{F:6,G:8,E:9,D:14}
F not goal	{G:4+2+1,G:8,E:9,D:14}



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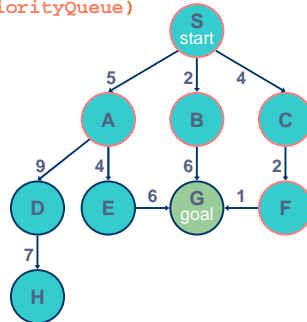
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## Uniform-Cost Search (UCS)

`generalSearch(problem, priorityQueue)`

# of nodes tested: 6, expanded: 5

expnd. node	nodes list
S	{S}
B	{B:2,C:4,A:5}
C	{A:5,F:6,G:8}
A	{F:6,G:8,E:9,D:14}
F	{G:7,G:8,E:9,D:14}
G goal	{G:8,E:9,D:14}
	no expand



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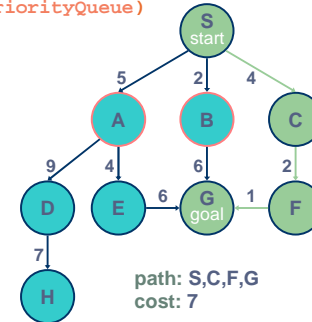
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## Uniform-Cost Search (UCS)

`generalSearch(problem, priorityQueue)`

# of nodes tested: 6, expanded: 5

expnd. node	nodes list
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B	{B:2,C:4,A:5}
C	{A:5,F:6,G:8}
A	{F:6,G:8,E:9,D:14}
F	{G:7,G:8,E:9,D:14}
G	{G:8,E:9,D:14}



path: S,C,F,G  
cost: 7

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## Uniform-Cost Search (UCS)

- Called *Dijkstra's Algorithm* in the algorithms lit.
- Similar to *Branch and Bound Algorithm* in operations research literature
- Complete
- Optimal/Admissible
  - requires that the goal test being applied when a node is removed from the nodes list
  - rather than when the node is first generated while its parent node is expanded

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## Uniform-Cost Search (UCS)

- Time and space complexity:  $O(b^d)$  (i.e., exponential)
  - $d$  is the depth of the solution
  - $b$ : the branching factor at each non-leaf node

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## Uninformed Search Strategies

- **IDS: depth-first, iterative-deepening search**  
requires modification to search algorithm:
  - do DFS to depth 1
  - treat all children of the start node as leafs
  - if no solution found, do DFS to depth 2
  - repeat by increasing depth until a solution found

\* Start node is at depth 0

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## Depth-First, Iterative-Deepening Search (IDS)

```
Node deepeningSearch (Problem problem) {
    int depth = 1;
    Stack DSnodes = new Stack();
    while (true) { // while not solved
        Node node = DFS_depthBound(problem, DSnodes, depth);
        // DFS_depthBound limits DFS search to level <= depth
        if (node isn't "failure") return node; // solved
        depth++; // look deeper
    }
}
```

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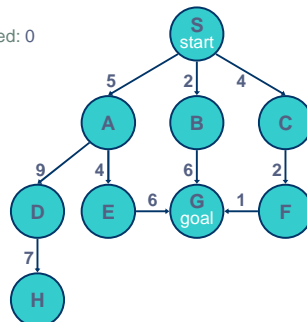
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## Depth-First, Iterative-Deepening Search (IDS)

deepeningSearch(problem)

depth: 1, # of nodes expanded: 0, tested: 0

expnd. node	nodes list
	{S}



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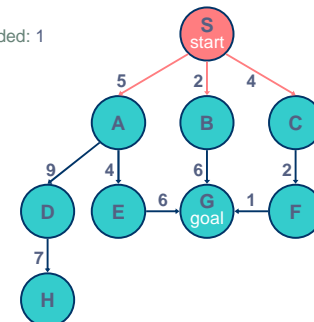
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## Depth-First, Iterative-Deepening Search (IDS)

deepeningSearch(problem)

depth: 1, # of nodes tested: 1, expanded: 1

expnd. node	nodes list
	{S}
S not goal	{A,B,C}



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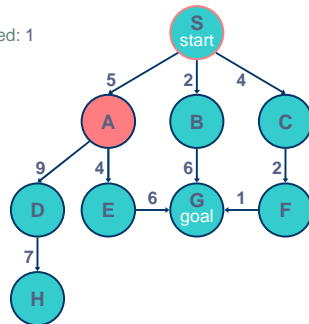
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## Depth-First, Iterative-Deepening Search (IDS)

deepeningSearch(problem)

depth: 1, # of nodes tested: 2, expanded: 1

expnd. node	nodes list
	{S}
S	{A,B,C}
A not goal	{B,C} no expand



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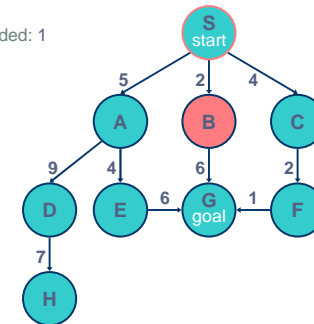
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## Depth-First, Iterative-Deepening Search (IDS)

deepeningSearch(problem)

depth: 1, # of nodes tested: 3, expanded: 1

expnd. node	nodes list
	{S}
S	{A,B,C}
A	{B,C}
B not goal	{C} no expand



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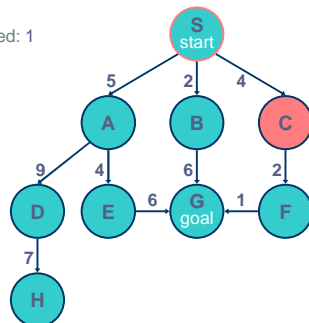
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## Depth-First, Iterative-Deepening Search (IDS)

deepeningSearch(problem)

depth: 1, # of nodes tested: 4, expanded: 1

expnd. node	nodes list
	{S}
S	{A,B,C}
A	{B,C}
B	{C}
C not goal	{ } no expand-FAIL



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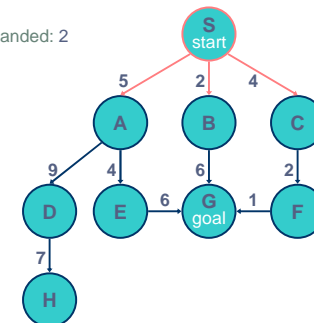
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## Depth-First, Iterative-Deepening Search (IDS)

deepeningSearch(problem)

depth: 2, # of nodes tested: 4(1), expanded: 2

expnd. node	nodes list
	{S}
S	{A,B,C}
A	{B,C}
B	{C}
C	{ }
S no test	{A,B,C}



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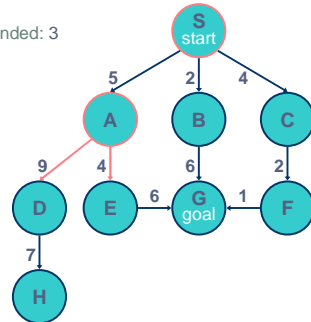
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## Depth-First, Iterative-Deepening Search (IDS)

deepeningSearch(problem)

depth: 2, # of nodes tested: 4(2), expanded: 3

expnd. node	nodes list
S	{S}
A	{A,B,C}
B	{B,C}
C	{C}
S	{A,B,C}
A no test	{D,E,B,C}



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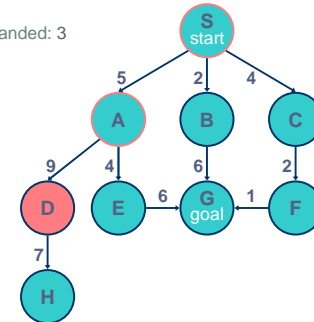
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## Depth-First, Iterative-Deepening Search (IDS)

deepeningSearch(problem)

depth: 2, # of nodes tested: 5(2), expanded: 3

expnd. node	nodes list
S	{S}
A	{A,B,C}
B	{B,C}
C	{C}
S	{A,B,C}
A	{D,E,B,C}
D not goal	{E,B,C} no expand



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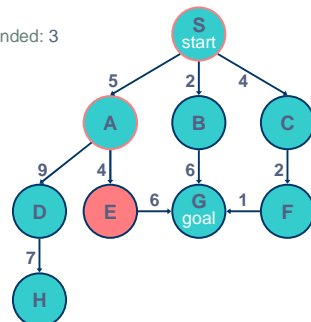
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## Depth-First, Iterative-Deepening Search (IDS)

deepeningSearch(problem)

depth: 2, # of nodes tested: 6(2), expanded: 3

expnd. node	nodes list
S	{S}
A	{A,B,C}
B	{B,C}
C	{C}
S	{A,B,C}
A	{D,E,B,C}
D	{E,B,C}
E not goal	{B,C} no expand



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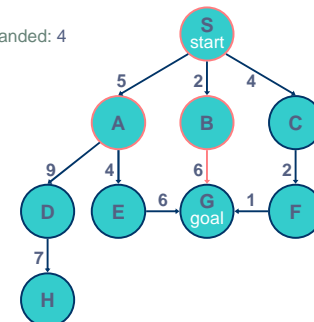
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## Depth-First, Iterative-Deepening Search (IDS)

deepeningSearch(problem)

depth: 2, # of nodes tested: 6(3), expanded: 4

expnd. node	nodes list
S	{S}
A	{A,B,C}
B	{B,C}
C	{C}
S	{A,B,C}
A	{D,E,B,C}
D	{E,B,C}
E	{B,C}
B no test	{G,C}



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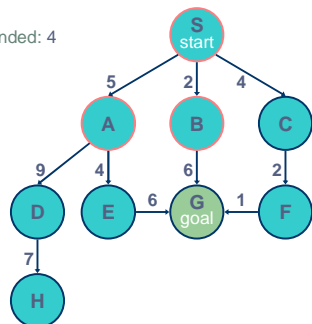


## Depth-First, Iterative-Deepening Search (IDS)

deepeningSearch(problem)

depth: 2, # of nodes tested: 7(3), expanded: 4

expnd. node	nodes list
S	{S}
A	{A,B,C}
B	{B,C}
C	{C}
S	{A,B,C}
A	{D,E,B,C}
D	{E,B,C}
E	{B,C}
B	{G,C}
G	{G,C}
G	goal



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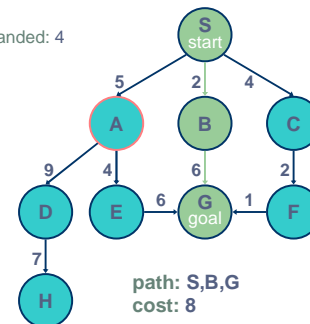
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## Depth-First, Iterative-Deepening Search (IDS)

deepeningSearch(problem)

depth: 2, # of nodes tested: 7(3), expanded: 4

expnd. node	nodes list
S	{S}
A	{A,B,C}
B	{B,C}
C	{C}
S	{A,B,C}
A	{D,E,B,C}
D	{E,B,C}
E	{B,C}
B	{G,C}
G	{G,C}



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## Depth-First, Iterative-Deepening Search (IDS)

- Has advantages of BFS
  - completeness
  - optimality as stated for BFS
- Has advantages of DFS
  - limited space
  - in practice, even with redundant effort it still finds longer paths more quickly than BFS

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## Depth-First, Iterative-Deepening Search (IDS)

- Space complexity:  $O(bd)$  linear like DFS
- Time complexity is a little worse than BFS or DFS
  - because nodes near the top of the search tree are generated multiple times (redundant effort)
- Worst case time complexity:  $O(b^d)$  exponential
  - because most nodes are near the bottom of tree

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## Depth-First, Iterative-Deepening Search (IDS)

How much redundant effort is done?

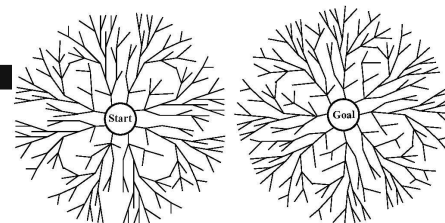
- The number of times the nodes are generated:  
 $1b^d + 2b^{(d-1)} + \dots + db \leq b^d / (1 - 1/b)^2 = O(b^d)$ 
  - $d$ : the solution's depth
  - $b$ : the branching factor at each non-leaf node
- For example:  $b = 4$   
 $4^d / (1 - 1/4)^2 = 4^d / (.75)^2 = 1.78 \times 4^d$ 
  - in the worst case, 78% more nodes are searched (redundant effort) than exist at depth  $d$
  - as  $b$  increases, this % decreases

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## Which Direction Should We Search?



Our choices: Forward, backwards, or bidirectional

The issues: How many start and goal states are there?  
 Branching factors in each direction  
 How much work is it to compare states?

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## General Search with Open and Close

```
//Note: this algorithm does detect loops in the state space
Node generalSearch (Problem problem, List OPEN) {
    OPEN.add(new Node(problem.getStartState()));
    List CLOSE = new List(); //initially empty, just a List DS
    while (true) {
        if (OPEN.isEmpty()) return new Node("failure");
        Node node = OPEN.remove(); //removes front node
        if (problem.isGoal(node.getState())) return node;
        CLOSE.add(node); //remember it was expanded
        OPEN.add(expand node given problem operators);
        //expand: only add successors not already in OPEN or CLOSE
    }
}
```

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