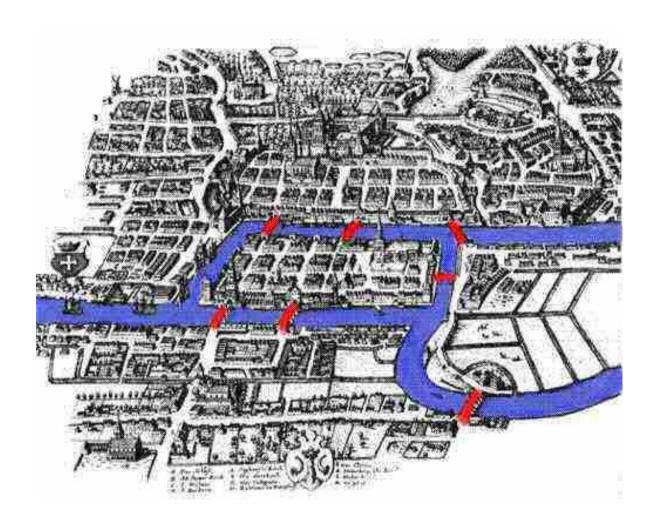
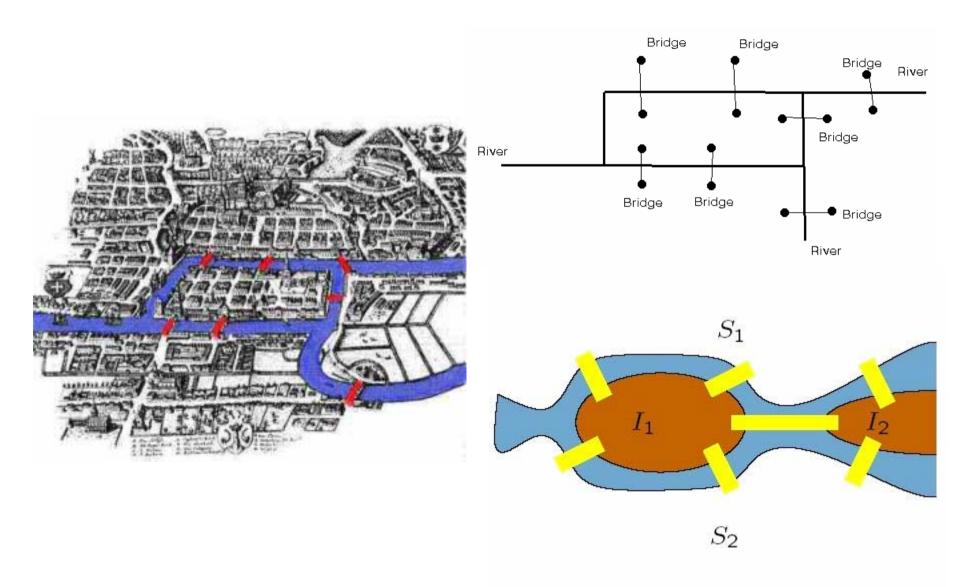
# Elementary Graph Algorithms (1)

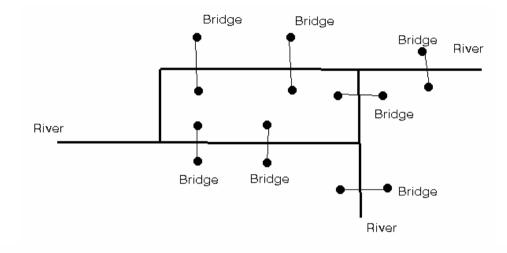
### Konigsberg City

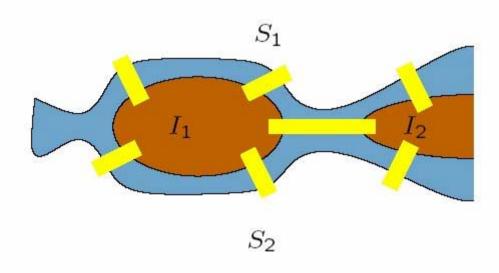


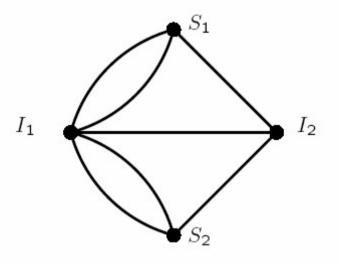


# **Definitions**

- A graph G = (V, E)
  - $V \Rightarrow$  set of vertices (singular: vertex) {node}
  - $E \Rightarrow set of edges$







#### • **Directed Graph** (digraph)

- Edge (u,v) incident from (leaves) vertex u and is incident to (enter) vertex v
- Self-loops—edges from a vertex to itself— are possible
- If (u,v) is an edge in G, v is adjacent to vertex u  $(u\rightarrow v)$   $\checkmark$  Adjacency relation is not necessarily symmetric
- A vertex *out-degree* → number of edges leaving it
- A vertex *in-degree* → number of edges entering it
- A vertex *degree* ← in-degree + out-degree

#### Undirected Graph

- Edge (u,v) and (v,u) are considered to be the same edge and is called incident on vertices u and v
- No self-loops
- If (u,v) is an edge in G, v is adjacent to vertex u
  ✓ Adjacency relation is symmetric
- A vertex *degree* → number of edges incident on it

- A *path* of *length* k from vertex u to a vertex u' in G(V,E) is a sequence  $\langle v_0, v_1, v_2, ..., v_k \rangle$  such that  $u = v_0$ ,  $u' = v_k$  and  $(v_{i-1}, v_i) \in E$  for i=0,1,2,...,k.
- A path *length* is number of edges in a path.
- u' is **reachable** from u if there is a path from u to u'
- A path is *simple* if all vertices are distinct

- A *path*  $\langle v_0, v_1, v_2, ..., v_k \rangle$  form a *cycle* if  $v_0 = v_k$  and the path contains at least one edge.
- A *cycle* is simple if  $\langle v_0, v_1, v_2, \dots, v_k \rangle$  are distinct.
- A self-loop is a cycle of length 1
- A digraph is *simple* if it doesn't contains any self-loop

## **Trees**

- A tree is a connected, acyclic, undirected graph.
- A *forest* is a disconnected, acyclic, undirected graph

# Connectivity

- Undirected Graph
  - Every pairs of vertices is connected by a path.
  - Connected components: classes of vertices under the "is reachable from" relation
- Directed graph
  - Strongly connected if every two vertices are reachable from each other.
  - Strongly connected components: classes of vertices under the "are mutually reachable" relation

## Isomorphism

• For any two graphs G=(V,E) and G'=(V',E'), if there exists a bijection  $f: V \rightarrow V'$  such that  $(u,v) \in E$  iff  $(f(u), f(v)) \in E'$ 

# Special Names Graphs

- A *Complete Graph*: is an undirected graph in which every pairs of vertices is adjacent.
- A *Bipartite Graph*: is an undirected graph G=(V,E) in which V can be partitioned into two sets V1 and V2 such that  $(u,v) \in E$  implies either  $u \in V_1$  and  $v \in V_2$  or  $v \in V_1$  and  $u \in V_2$
- A Weighted Graph: associates weights with either the edges or the vertices

# Representation of Graphs

