# Cache Misses

### **Description**

Modern computers use a **cache** to store a **small amount of data** in a high-speed memory, significantly improving overall system performance. Even though programs frequently access **large amounts of data**, by keeping only a **subset of the main memory** in the cache, access time can be greatly reduced.

When a program runs, it makes a sequence of **memory requests**:

$$\{r1, r2, \dots, rn\}$$

Each request corresponds to a **specific data element** that the program needs to access. If the requested element is already in the cache, it is considered a **cache hit**, meaning it is accessed instantly. However, if the requested element is not in the cache, it results in a **cache miss**, requiring the system to retrieve it from the main memory, which is significantly slower.

Since the cache can only hold **k** elements at a time, when it is full and a new request arrives, the system must decide which element to evict to make space for the new one. The goal of the cache management strategy is to minimize the total number of cache misses over the entire request sequence. An example of a program that accesses four distinct elements {A, B, C, D} might generate the following request sequence:

If the cache size is k = 3, each time a new element is requested, the system must determine which elements to **keep in the cache** and which one to **evict**, ensuring that future requests have the highest chance of being cache hits.

Typically, caching is considered an **on-line problem**, meaning the system must make real-time decisions about which elements to retain in the cache **without knowledge of future requests**. However, in this problem, we consider an **off-line caching version**, where the **entire sequence of requests** is **known in advance**. This allows us to use an optimal **cache replacement strategy**, where elements are evicted **based on their future use**, ensuring the **fewest possible cache misses**.

A crucial rule in this problem is that **every request must be cached if it is not already present**. It is **not allowed** to ignore or skip caching a request, even if it will never be used again. If the cache is full, an eviction **must** take place before adding a new request. The **eviction strategy should** minimize the cache misses.

Your task is to design an **efficient algorithm** to solve this offline caching problem to **minimize** the cache misses?

#### Given:

- An integer cacheSize (K), representing the maximum number of elements the cache can store.
- An array of **N** requests, representing the **sequence of memory requests**.

### Required:

• The minimum number of cache misses that occurs while executing the requests.

# **Complexity**

complexity of your algorithm should be less than  $O(K \times N^2)$ ;

### **Function:** Implement it!

```
static public int RequiredFunction(int cacheSize, string[] requests)
```

PROBLEM\_CLASS.cs includes this method.

## **Examples**

### Test Case 1:

- Cache Size: 2
- Requests: ["R0", "R1", "R2", "R1", "R2", "R2", "R1"]

### **Step-by-Step Execution**

- **1.** R0 is not in the cache  $\rightarrow$  Miss  $\rightarrow$  Cache = [R0]
- 2. R1 is not in the cache  $\rightarrow$  Miss  $\rightarrow$  Cache = [R0, R1]
- 3. R2 is not in the cache  $\rightarrow$  Miss  $\rightarrow$  Evict R0  $\rightarrow$  Cache = [R1, R2]
- **4.** R1 is already in the cache → Hit
- **5.** R2 is already in the cache  $\rightarrow$  Hit
- **6.** R2 is already in the cache  $\rightarrow$  Hit
- 7. R1 is already in the cache  $\rightarrow$  Hit

# ✓ Min Cache Misses = 3

### Test Case 2:

- Cache Size: 3
- Requests: ["R10", "R11", "R12", "R10", "R11", "R12"]

#### **Step-by-Step Execution**

- **1.** R10 is not in the cache  $\rightarrow$  Miss  $\rightarrow$  Cache = [R10]
- 2. R11 is not in the cache  $\rightarrow$  Miss  $\rightarrow$  Cache = [R10, R11]
- 3. R12 is not in the cache  $\rightarrow$  Miss  $\rightarrow$  Cache = [R10, R11, R12]
- **4.** R10 is in the cache  $\rightarrow$  Hit
- **5.** R11 is in the cache  $\rightarrow$  Hit
- **6.** R12 is in the cache  $\rightarrow$  Hit

# ✓ Min Cache Misses = 3

#### Test Case 3:

- Cache Size: 3
- Requests: ["R20", "R21", "R22", "R23", "R24", "R25", "R26"]

#### **Step-by-Step Execution**

- **1.** R20 is not in the cache  $\rightarrow$  Miss  $\rightarrow$  Cache = [R20]
- 2. R21 is not in the cache  $\rightarrow$  Miss  $\rightarrow$  Cache = [R20, R21]
- 3. R22 is not in the cache  $\rightarrow$  Miss  $\rightarrow$  Cache = [R20, R21, R22]
- **4.** R23 is not in the cache. Cache is full → Evict R20 → Cache = [R21, R22, R23] → Miss
- **5.** R24 is not in the cache. Cache is full  $\rightarrow$  Evict R21  $\rightarrow$  Cache = [R22, R23, R24]  $\rightarrow$  Miss
- **6.** R25 is not in the cache. Cache is full  $\rightarrow$  Evict R22  $\rightarrow$  Cache = [R23, R24, R25]  $\rightarrow$  Miss
- 7. R26 is not in the cache. Cache is full  $\rightarrow$  Evict R23  $\rightarrow$  Cache = [R24, R25, R26]  $\rightarrow$  Miss
- ✓ Min Cache Misses = 7

#### Test Case 4:

- Cache Size: 3
- Requests: ["R30", "R31", "R32", "R30", "R33", "R34", "R30", "R31", "R32"]

### **Step-by-Step Execution**

- 1. R30 is not in the cache  $\rightarrow$  Miss  $\rightarrow$  Cache = [R30]
- 2. R31 is not in the cache  $\rightarrow$  Miss  $\rightarrow$  Cache = [R30, R31]
- 3. R32 is not in the cache  $\rightarrow$  Miss  $\rightarrow$  Cache = [R30, R31, R32]
- **4.** R30 is in the cache  $\rightarrow$  Hit
- **5.** R33 is not in the cache. Cache is full  $\rightarrow$  Evict R32  $\rightarrow$  Cache = [R30, R31, R33]  $\rightarrow$  Miss
- **6.** R34 is not in the cache. Cache is full  $\rightarrow$  Evict R33  $\rightarrow$  Cache = [R30, R31, R34]  $\rightarrow$  Miss
- 7. R30 is in the cache  $\rightarrow$  Hit
- **8.** R31 is in the cache  $\rightarrow$  Hit
- **9.** R32 is not in the cache. Cache is full  $\rightarrow$  Evict R34  $\rightarrow$  Cache = [R30, R31, R32]  $\rightarrow$  Miss
- ✓ Min Cache Misses = 6

#### Test Case 5:

- Cache Size: 2
- Requests: ["R40", "R41", "R42", "R41", "R41", "R40", "R40", "R40", "R41", "R42"]

#### **Step-by-Step Execution**

- **1.** R40 is not in the cache  $\rightarrow$  Miss  $\rightarrow$  Cache = [R40]
- 2. R41 is not in the cache  $\rightarrow$  Miss  $\rightarrow$  Cache = [R40, R41]
- **3.** R42 is not in the cache. Cache is full  $\rightarrow$  Evict R40  $\rightarrow$  Cache = [R41, R42]  $\rightarrow$  Miss
- **4.** R41 is in the cache  $\rightarrow$  Hit
- **5.** R41 is in the cache  $\rightarrow$  Hit

- **6.** R40 is not in the cache. Cache is full  $\rightarrow$  Evict R42  $\rightarrow$  Cache = [R41, R40]  $\rightarrow$  Miss
- 7. R40 is in the cache  $\rightarrow$  Hit
- **8.** R41 is in the cache  $\rightarrow$  Hit
- **9.** R41 is in the cache  $\rightarrow$  Hit
- **10.** R42 is not in the cache. Cache is full  $\rightarrow$  Evict R40  $\rightarrow$  Cache = [R41, R42]  $\rightarrow$  Miss
- Min Cache Misses = 5

#### Test Case 6:

- Cache Size: 4
- Requests: ["R50", "R51", "R52", "R53", "R50", "R54", "R55", "R56", "R50", "R57", "R56", "R54", "R51", "R57", "R50"]

### **Step-by-Step Execution**

✓ Min Cache Misses = 9

# C# Help

### Getting the size of 1D array

```
int size = array1D.GetLength(0);
```

### **Getting the size of 2D array**

```
int size1 = array2D.GetLength(0);
int size2 = array2D.GetLength(1);
```

### **Creating 1D array**

```
int [] array1D = new int [size]
```

## **Creating 2D array**

```
int [,] array2D = new int [size1, size2]
```

# **Sorting single array**

Sort the given array "items" in ascending order

```
Array.Sort(items);
```

### **Sorting parallel arrays**

Sort the first array "master" and re-order the 2<sup>nd</sup> array "slave" according to this sorting

```
Array.Sort(master, slave);
```