Advanced Programming Languages. Practice 1.

EPIC Institute of Technology

September 22, 2023

1 Formal grammars

- Non-terminal symbols (A, B, ..., Y, Z)
- Terminal symbols (a, b, ..., y, z)
- Start non-terminal symbol (S)
- Set of rules:
  - $AcB \to \varepsilon$
  - $S \to CCCd$
  - $A \to a|bcD \quad \equiv \quad A \to a, A \to bcD$

1.1 Integers

$$S \to AB$$
$$A \to \varepsilon|+|-$$
$$B \to D|DB$$
$$D \to 0|1|2|3|4|5|6|7|8|9$$

Example: $S \to AB \to -B \to -DB \to -DD \to -4D \to -42$.

1.2 Brackets

$$S \to \varepsilon|(S)|SS$$

$$S \to \varepsilon|(S)S$$

Example: $S \to (S)S \to ()S \to ()(S)S \to ()((S)S)S \to ()(()S)S \to ()(()(S)S)S \to ()(()()S)S \to ()(()())S \to ()(()())$.

2 Types of grammars

Recursively enumerable (Type 0). no restrictions on grammars

Context-sensitive (Type 1). $\alpha A\beta \to \alpha\gamma\beta$.

Context-free (Type 2). $A \to \alpha$

Regular (Type 3). $A \to aB,\ A \to a$

2.1 Determine type of grammar

1.

$$S \to aBC$$
$$aB \to CD$$
$$C \to a$$
$$aC \to BD$$
$$D \to b$$

2.

$$S \to \varepsilon|(S)S$$

3.

$$S \to \varepsilon|aSb$$

4.

$$S \to aSAB|aAB$$
$$BA \to BD$$
$$BD \to ED$$
$$ED \to AD$$
$$AD \to AB$$
$$aA \to ab$$
$$bA \to bb$$
$$bB \to bc$$
$$cB \to cc$$

5.

$$S \to AB$$
$$A \to \varepsilon|+|-$$
$$B \to D|DB$$
$$D \to 0|1|2|3|4|5|6|7|8|9$$

Can we turn it into a regular one?

$$S \to \varepsilon B|+B|-B$$
$$B \to 0|1|2|3|4|5|6|7|8|9|0B|1B|2B|3B|4B|5B|6B|7B|8B|9B$$

3 Regular expressions

- $[abcd]$ — one of the characters from the set $\{a, b, c, d\}$
- a — single character
- (R) — group
- $*$, $+$, $?$ — repetition
- $R|R$ — boolean or

Examples:

- `first|second`
- `[+-]?[0-9]+`
- `(Mr. |Mrs. )?[A-Z][a-z]+ [A-Z][a-z]+`
- `[+-]?([1-9][0-9]+|0x[0-9a-fA-F]+|0[0-7]*)`

4 Backus–Naur form

```
<expr> ::= <sum>
<sum> ::= <mul> | <sum> "+" <mul> | <sum> "-" <mul>
<mul> ::= <last> | <mul> "*" <last> | <mul> "/" <last>
<last> ::= <number> | "+" <last> | "-" <last> | "(" <expr> ")"
<number> ::= <digit> | <digit> <number>
<digit> ::= "0" | "1" | "2" | "3" | "4" | "5" | "6" | "7" | "8" | "9"
```

What kind of grammar can we specify?
Let's try to parse:
$2 * 4 + (2 + -3 * 5)/7$

Syntax for regular expressions above:

```
<regex> ::= <concat> | <regex> "|" <concat>
<concat> ::= <item> | <concat> <item>
<item> ::= <item> <repetition> | "(" <item> ")" | <character> | <group>
<group> ::= "[" <groupItems> "]"
<groupItems> ::= "" | <character> <groupItems>
<character> ::= "a" | "b" | ... | "z" | "A" | ... "Z"
<repetition> ::= "?" | "+" | "*"
```